Two Phase Commit



Distributed transactions

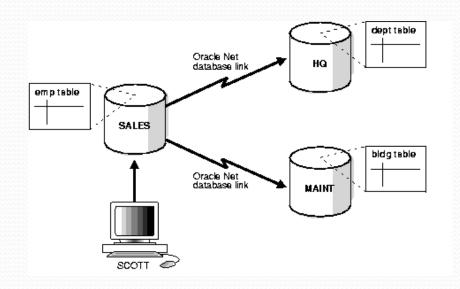
- If transaction modifies objects in multiple databases, a single "two phase commit" can commit changes in all databases, whilst retaining ACID properties.
- Distributed or remote?
 - Note: If all statements of a transaction reference only a single remote node, then the transaction is remote, not distributed.

The databases involved are called nodes; the initiating node is called the 'Global Coordinator'

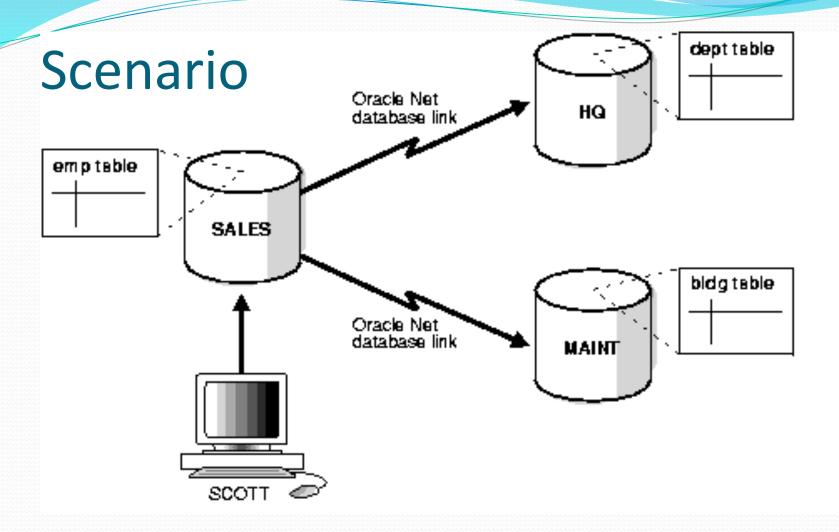


Example

- A distributed transaction updates data on two or more distinct nodes of a distributed database.
- This examples has three databases, one local, and two remote.
- The user is updating tables on all three databases using a distributed transaction.







• User Scott updates the dept table on the remote HQ database, the local emp table and the remote bldg table on the MAINT database, in one transaction.

Example

- This distributed transaction run by scott updates
 - the remote hq database
 - dept table
 - the local sales database
 - emp table
 - the remote maint database
 - bldg table
- Then commits

```
UPDATE
scott.dept@hq.us.acme.com
  SET loc =
  'REDWOOD SHORES'
  WHERE deptno = 10;
UPDATE scott.emp
set deptno = 11
Where deptno = 10;
UPDATE
scott.bldg@maint.us.acme.com
  SET room = 1225
  WHERE room = 1163;
 COMMIT;
```



The two phases

- Prepare phase
 - The global coordinator (initiating node) asks participants to **prepare** (to promise to commit or rollback the transaction, even if there is a failure)
- Commit Phase
 - If all participants respond to the coordinator that they are prepared, the coordinator asks all nodes to **commit** the transaction.
 - If all participants cannot prepare, the coordinator asks all nodes to roll back the transaction.



Prepare Phase

- By preparing, a node:
 - Logs the transaction locally.
 - Places a distributed lock on modified tables, which prevents reads
 - Responds with:
 - Prepared
 - Read-only or
 - Abort





What happens next?



- The prepared nodes then wait until a COMMIT or ROLLBACK request is received from the global coordinator.
- After the nodes are prepared, the distributed transaction is said to be in-doubt until all changes are either committed or rolled back.



Steps in the Commit Phase

- The commit phase consists of the following steps:
 - The global coordinator issues an order to commit.
 - 2. At each node, the local portion of the distributed transaction is committed and locks are released.
 - The participating nodes notify the global coordinator that they have committed.
- When the commit phase is complete, the data on all nodes of the distributed system is consistent.



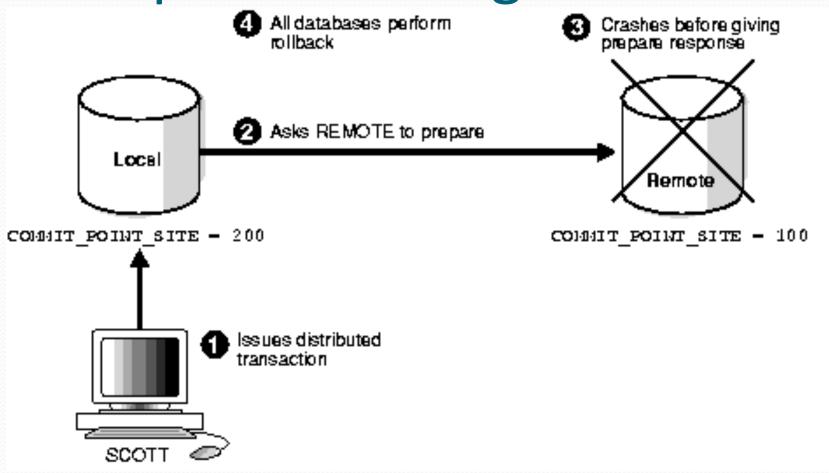
In-Doubt Transactions

- The two-phase commit mechanism ensures that all nodes either commit or perform a rollback together.
- What happens if any of the phases fails because of a system or network error?
- The transaction becomes in-doubt.



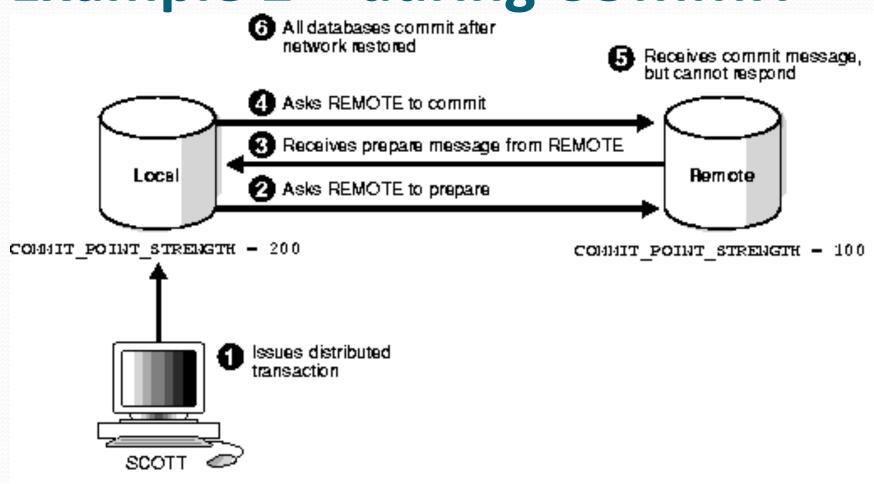


Example 1 – during PREPARE





Example 2 – during COMMIT





Trade-offs



CAP theorem

- There are three core systemic requirements that exist in a special relationship when designing and deploying distributed applications.
 - Consistency,
 - Availability and
 - Partition Tolerance



Partition Tolerance

- Partition tolerance in CAP means tolerance to a network partition.
 - i.e. when two nodes can't talk to each other.
 - A CA system guarantees strong consistency, at the cost of not being able to process requests unless all nodes are able to talk to each other.
 - An AP system is able to function during the network split, while being able to provide various forms of eventual consistency.



Eventual consistency

- Weak eventual consistency
 - This is where a write might not be consistent across the network partition, meaning that it is not possible to read them.
 - BASE (Basically Available, Soft state, Eventual consistency)
- 'Read your writes'
 - Rather than being sure that all reads are consistent, this method reads from N replicated copies and if W writes (< N) agree, then that is considered to be the correct answer.

