

UNIVERSITY OF COPENHAGEN  
Applied High Performance Computing

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## 1 Definitions

**Definition 1.1** (Union-Find). The union-find data structure  $U$  represents a partition of a set  $S$  if:

1.  $a \neq \emptyset$  for all  $a \in \mathcal{C}(U)$
2.  $a \cap b = \emptyset$  for all  $a, b \in \mathcal{C}(U)$  where  $a \neq b$
3.  $\bigcup_{a \in \mathcal{C}(U)} a = S$

where  $\mathcal{C} : \mathbb{P}(U) \rightarrow \mathbb{P}(\mathbb{P}(S))$  converts the data structure  $U$  into a partition of  $S$ .

**Definition 1.2** (Union-Find Set). The set  $\mathbb{U}_n := \mathbb{P}(\mathbb{N}_{\leq n} \times (\mathbb{N}_{\leq n} \cup \{0\}))$  is the set of all union-find data structures of size  $n$ . Each element  $U \in \mathbb{U}_n$  is a set containing tuples  $(i, p)$  where:

1.  $i \in \mathbb{N}_{\leq n}$  is a unique index identifying the element i.e.  $|\{\pi_1(u) : u \in U\}| = n$ .
2.  $p \in \mathbb{N}_{\leq n} \cup \{0\}$  is the index of the parent element. If  $p = 0$  then the element is a root.

**Definition 1.3** (Cycle). A cycle exists in a set  $U \in \mathbb{U}_n$  if:

1.  $i = p$  for some  $(i, p) \in U$  or
2.  $p_1 = i_2, p_2 = i_3, \dots, p_m = i_1$  for some  $(i_1, p_1), (i_2, p_2), \dots, (i_m, p_m) \in U$  where  $m > 1$ .

**Definition 1.4** (Find Root). The root of an element  $i$  in a cycleless union-find data structure  $U \in \mathbb{U}_n$  is defined as:

$$\mathcal{P}_U(i) = \begin{cases} i & \text{if } p = 0 \\ \mathcal{P}_U(p) & \text{if } p \neq 0 \end{cases} \text{ where } (j, p) \in U \wedge i = j.$$

**Proposition 1.1** (Termination of Find Root). The function  $\mathcal{P}_U(i)$  defined in definition 1.4 terminates for all  $i$  where  $1 \leq i \leq n$  if  $U \in \mathbb{U}_n$  has no cycles.

*Proof.* The proof is trivially true. □

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**Proposition 1.2** (Representation of Union-Find). A cycleless union-find structure  $U \in \mathbb{U}_n$  fulfills definition 1.1 for the following conversion function:

$$\mathcal{C}(U) := \{\{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = p\} : p \in \mathbb{N}_{\leq n}\} \setminus \{\emptyset\}$$

where  $f : \mathbb{N}_{\leq n} \rightarrow S$  is a bijective function mapping indices to elements in the set  $S$ .

*Proof.* Let  $U \in \mathbb{U}_n$  be a union-find data structure fulfilling the three criteria in proposition 1.2. We will show that  $U$  fulfills the three criteria in definition 1.1.

1. By definition of  $\mathcal{C}$  it is clear that  $a \neq \emptyset$  for all  $a \in \mathcal{C}(D)$  since  $\mathcal{C}$  only includes non-empty sets.
2. Let  $a, b \in \mathcal{C}(D)$  where  $a \neq b$  then by definition  $a = \{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = p_a\}$  and  $b = \{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = p_b\}$  for some  $p_a, p_b \in \mathbb{N}_{\leq n}$  where  $p_a \neq p_b$ . Since if  $p_a = p_b$  then  $a = b$  it follows that  $a \cap b = \emptyset$ .
3. Let  $s \in S$  then since  $U$  is a union-find data structure of size  $n$  there exists  $(i, p, s) \in U$  for some  $i$  and  $p$ . Let  $r = \mathcal{P}_U(i)$  then by definition of  $\mathcal{C}$  it follows that  $s \in \{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = r\} \in \mathcal{C}(U)$ . Since  $s$  was arbitrary it follows that  $\bigcup_{a \in \mathcal{C}(U)} a = S$ .

□

**Definition 1.5** (Equivalence Set). The set of equivalent indices of an element  $i$  in a cycleless union-find data structure  $U \in \mathbb{U}_n$  is defined as:

$$\mathcal{E}_U(i) := \{j : j \in \mathbb{N}_{\leq n} \wedge \mathcal{P}_U(j) = \mathcal{P}_U(i)\}$$

**Definition 1.6** (Equivalent Union-Find Structures). Two cycleless union-find data structures  $U, U' \in \mathbb{U}_n$  are equivalent if:

$$\mathcal{E}_{U'}(i) = \mathcal{E}_U(i) \text{ for all } i \in \mathbb{N}_{\leq n}$$

**Definition 1.7** (Well-formed Union). For a cycleless union-find data structure  $U \in \mathbb{U}_n$  the union of two elements  $i \sim j$  where  $i, j \in \mathbb{N}_{\leq n}$  is well-formed if it results in a cycleless union-find data structure  $U' \in \mathbb{U}_n$  such that:

1.  $\mathcal{E}_{U'}(i) = \mathcal{E}_U(i) \cup \mathcal{E}_U(j)$  and
2.  $\mathcal{E}_{U'}(k) = \mathcal{E}_U(k)$  for all  $k \in \mathbb{N}_{\leq n} \setminus \{i, j\}$ .

**Definition 1.8** (Sequential Union).