

# UNIVERSITY OF COPENHAGEN

## Union-Find

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## 1 Definitions

**Definition 1.1** (Union-Find). The union-find data structure  $U$  represents a partition of a set  $S$  if:

1.  $a \neq \emptyset$  for all  $a \in \mathcal{C}(U)$
2.  $a \cap b = \emptyset$  for all  $a, b \in \mathcal{C}(U)$  where  $a \neq b$
3.  $\bigcup_{a \in \mathcal{C}(U)} a = S$

where  $\mathcal{C} : \mathbb{P}(U) \rightarrow \mathbb{P}(\mathbb{P}(S))$  converts the data structure  $U$  into a partition of  $S$ .

**Definition 1.2** (Union-Find Set). The set  $\mathbb{U}_n := \mathbb{P}(\mathbb{N}_{\leq n} \times (\mathbb{N}_{\leq n} \cup \{0\}))$  is the set of all union-find data structures of size  $n$ . Each element  $U \in \mathbb{U}_n$  is a set containing tuples  $(i, p)$  where:

1.  $i \in \mathbb{N}_{\leq n}$  is a unique index identifying the element i.e.  $|\{\pi_1(u) : u \in U\}| = n$ .
2.  $p \in \mathbb{N}_{\leq n} \cup \{0\}$  is the index of the parent element. If  $p = 0$  then the element is a root.

**Definition 1.3** (Cycle). A cycle exists in a set  $U \in \mathbb{U}_n$  if:

1.  $i = p$  for some  $(i, p) \in U$  or
2.  $p_1 = i_2, p_2 = i_3, \dots, p_m = i_1$  for some  $(i_1, p_1), (i_2, p_2), \dots, (i_m, p_m) \in U$  where  $m > 1$ .

**Definition 1.4** (Find Root). The root of an element  $i$  in a cycleless union-find data structure  $U \in \mathbb{U}_n$  is denoted  $\mathcal{P}_U : \mathbb{N}_{\leq n} \rightarrow \mathbb{N}_{\leq n}$  defined as:

$$\mathcal{P}_U(i) := \begin{cases} i & \text{if } p = 0 \\ \mathcal{P}_U(p) & \text{if } p \neq 0 \end{cases} \text{ where } (j, p) \in U \wedge i = j.$$

The root represents the set that the element  $i$  belongs to, i.e. all elements with the same root belong to the same set.

**Proposition 1.1** (Termination of Find Root). The function  $\mathcal{P}_U(i)$  defined in definition 1.3 terminates for all  $i$  where  $1 \leq i \leq n$  if  $U \in \mathbb{U}_n$  has no cycles.

*Proof.* The proof is trivially true.  $\square$

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**Proposition 1.2** (Representation of Union-Find). A cycleless union-find structure  $U \in \mathbb{U}_n$  fulfills definition 1.1 for the following conversion function:

$$\mathcal{C}(U) := \{\{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = p\} : p \in \mathbb{N}_{\leq n}\} \setminus \{\emptyset\}$$

where  $f : \mathbb{N}_{\leq n} \rightarrow S$  is a bijective function mapping indices to elements in the set  $S$ .

*Proof.* Let  $U \in \mathbb{U}_n$  be a union-find data structure fulfilling the three criteria in proposition 1.2. We will show that  $U$  fulfills the three criteria in definition 1.1.

1. By definition of  $\mathcal{C}$  it is clear that  $a \neq \emptyset$  for all  $a \in \mathcal{C}(D)$  since  $\mathcal{C}$  only includes non-empty sets.
2. Let  $a, b \in \mathcal{C}(D)$  where  $a \neq b$  then by definition  $a = \{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = p_a\}$  and  $b = \{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = p_b\}$  for some  $p_a, p_b \in \mathbb{N}_{\leq n}$  where  $p_a \neq p_b$ . Since if  $p_a = p_b$  then  $a = b$  it follows that  $a \cap b = \emptyset$ .
3. Let  $s \in S$  then since  $U$  is a union-find data structure of size  $n$  there exists  $(i, p) \in U$  for some  $i$  and  $p$ . Let  $r = \mathcal{P}_U(i)$  then by definition of  $\mathcal{C}$  it follows that  $s \in \{f(\pi_1(u)) : u \in U \wedge \mathcal{P}_U(\pi_1(u)) = r\} \in \mathcal{C}(U)$ . Since  $s$  was arbitrary it follows that  $\bigcup_{a \in \mathcal{C}(U)} a = S$ .

□

**Definition 1.5** (Substitute Parent). The function for substitution a parent of an element  $i$  to  $p$  in a cycleless union-find data structure  $U \in \mathbb{U}_n$  is denoted  $\mathcal{S}_U : \mathbb{N}_{\leq n} \times \mathbb{N}_{\leq n} \rightarrow \mathbb{P}(\mathbb{U}_n)$  and defined as:

$$\mathcal{S}_U(i, p) := (U \setminus \{(i, p')\}) \cup \{(i, p)\} \text{ where } (i, p') \in U$$

**Definition 1.6** (Equivalence Set). The set of equivalent indices of an element  $i$  in a cycleless union-find data structure  $U \in \mathbb{U}_n$  is defined as:

$$\mathcal{E}_U(i) := \{j : j \in \mathbb{N}_{\leq n} \wedge \mathcal{P}_U(j) = \mathcal{P}_U(i)\}$$

**Definition 1.7** (Parent). A element  $i \in \mathbb{N}_{\leq n}$  has a parent  $p \in \mathbb{N}_{\leq n} \cup \{0\}$  if

**Definition 1.8** (Equivalent Union-Find Structures). Two cycleless union-find data structures  $U, U' \in \mathbb{U}_n$  are equivalent if:

$$\mathcal{E}_{U'}(i) = \mathcal{E}_U(i) \text{ for all } i \in \mathbb{N}_{\leq n}$$

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**Lemma 1.1** (Cycleless substitution). A parent substitution  $\mathcal{S}_U(i, p)$  of a cycleless union-find data structure  $U \in \mathbb{U}_n$  will result in a cycleless union-find data structure if:

$$\mathcal{P}_U(i) \neq \mathcal{P}_U(p)$$

*Proof.* □

**Definition 1.9** (Well-formed Union). For a cycleless union-find data structure  $U \in \mathbb{U}_n$  the union of two elements  $i \sim j$  where  $i, j \in \mathbb{N}_{\leq n}$  is well-formed if it results in a cycleless union-find data structure  $U' \in \mathbb{U}_n$  such that:

1.  $\mathcal{E}_{U'}(i) = \mathcal{E}_U(i) \cup \mathcal{E}_U(j)$  and
2.  $\mathcal{E}_{U'}(k) = \mathcal{E}_U(k)$  for all  $k \in \mathbb{N}_{\leq n} \setminus (\mathcal{E}_U(i) \cup \mathcal{E}_U(j))$ .

**Definition 1.10** (Sequential Unions). Given a sequence of union relations  $i_1 \sim j_1, i_2 \sim j_2, \dots, i_m \sim j_m$  where  $i_k, j_k \in \mathbb{N}_{\leq n}$  for all  $1 \leq k \leq m$  and a initial  $U \in \mathbb{U}_n$ . The final union-find data structure  $U' \in \mathbb{U}_n$  is defined as:

$$\begin{aligned} U_0 &:= \{(i, 0) : i \in \mathbb{N}_{\leq n}\} \\ (q_k, p_k) &:= (\mathcal{P}_{U_{k-1}}(i_k), \mathcal{P}_{U_{k-1}}(j_k)) \text{ for } 1 \leq k \leq m \\ U_k &:= \begin{cases} U_{k-1} & \text{if } p_k = q_k \\ \mathcal{S}_{U_{k-1}}(q_k, p_k) & \text{if } p_k \neq q_k \end{cases} \text{ for } 1 \leq k \leq m \\ U' &:= U_m \end{aligned}$$

**Proposition 1.3** (Sequential Unions is Well-formed). Given a sequence of union relations  $i_1 \sim j_1, i_2 \sim j_2, \dots, i_m \sim j_m$  where  $i_k, j_k \in \mathbb{N}_{\leq n}$  for all  $1 \leq k \leq m$  and a cycleless union-find data structure  $U \in \mathbb{U}_n$ . The final union-find data structure  $U' \in \mathbb{U}_n$  defined in definition 1.10 is well-formed.

*Proof.* To prove that  $U'$  is well-formed we need to show that each union in the sequence is well-formed. This can be shown by induction on the index  $k$  of the union in the sequence.

- For  $U_0$  it is clear that  $\mathcal{E}_{U_0}(i) = \{i\}$  for all  $i \in \mathbb{N}_{\leq n}$  since all elements are their own root.
- Assume that  $U_{k-1}$  is well-formed for some  $1 \leq k \leq m$ . We will show that  $U_k$  is well-formed. Let  $q_k = \mathcal{P}_{U_{k-1}}(i_k)$  and  $p_k = \mathcal{P}_{U_{k-1}}(j_k)$ .
  1. If  $p_k = q_k$  then by definition  $U_k = U_{k-1}$  and thus  $U_k$  is well-formed by the inductive hypothesis.

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2. If  $p_k \neq q_k$  then since  $q_k$  is a root  $(q_k, 0) \in U_{k-1}$  and by definition of  $\mathcal{S}$ .

$$U_k := (U_{k-1} \setminus \{(q_k, 0)\}) \cup \{(q_k, p_k)\} \text{ where } (q_k, 0) \in U_{k-1}$$

Now  $p_k = \mathcal{P}_{U_k}(q_k)$  hence  $\mathcal{P}_{U_k}(q) = p_k$  for all  $q \in \mathcal{E}_{U_{k-1}}(q_k)$ . And by definition of  $p_k$  being a root it follows that  $\mathcal{P}_{U_k}(p) = p_k$  for all  $p \in \mathcal{E}_{U_{k-1}}(p_k)$ . Thus  $\mathcal{E}_{U_k}(i_k) = \mathcal{E}_{U_{k-1}}(i_k) \cup \mathcal{E}_{U_{k-1}}(j_k)$  and since all other elements are unaffected by the union it follows that  $\mathcal{E}_{U_k}(l) = \mathcal{E}_{U_{k-1}}(l)$  for all  $l \in \mathbb{N}_{\leq n} \setminus (\mathcal{E}_U(i) \cup \mathcal{E}_U(j))$ .

Finally since  $U_{k-1}$  is cycleless and  $q_k$  is a root it follows that  $U_k$  is cycleless and thus  $U_k$  is well-formed.

□

**Definition 1.11** (Parallel Unions).