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Lambada, Haskell as a better Java

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Abstract

The Lambada framework provides facilities for fluid interoperation between Haskell (currently both Hugs and GHC using non-Haskell98 extensions) and Java. Using Lambada, we can call Java methods from Haskell, and have Java methods invoke Haskell functions. The framework rests on the Java Native Interface (JNI). The Lambada release includes a tool for generating IDL from Java `.class` *files (using reflection), which is fed into our existing HDirect to generate Haskell-callable stubs.*

1 Introduction

It goes without saying that the ability to interact with other languages is of vital importance for the long-term survival of any programming language, in particular for niche languages such as Haskell or ML. Java is an interesting partner for interoperation for a number of reasons:

- Java comes with an large set of stable, and usually well-documented and well-designed libraries and APIs.
- The interaction with Java via the Java Native Interface (JNI) [8,4] effectively allows us to script the underlying JVM. This makes bidirectional interoperability with Java very flexible and vendor independent.
- Because of Java's platform independence, we hope that Lambada gains wider acceptance than our previous work on COM.

Notwithstanding all these advantages, using JNI in its most primitive form is tedious and error prone. Using raw JNI is like assembly programming on the JVM. *Our mission is to make interoperation between Java and Haskell via JNI as convenient as programming in Java directly.*

The architecture of Lambada borrows heavily from our previous work on interfacing Haskell and COM [6,2,3], in particular the binding for Automation [7]. We refer the reader to those papers for a discussion on the rationale behind the object encoding used for Lambada.

2 The Java Native Interface

Simply said, JNI provides a number of invocation functions to initialize and instantiate a Java Virtual Machine (JVM), and defines a COM interface to the JVM (the `JavaVM` interface), and to individual threads that run inside a particular JVM (the `JNIEnv` interface). The `JNIEnv` interface contains roughly 230 methods to interact with objects that live in its associated thread.

2.1 Calling C From Java and vice versa

To explain the way C and Java interact via JNI we use a slightly perverse version of `Hello World!` where a Java method calls a C procedure that calls back into Java to print `Hello World!`.

To use foreign functions in Java we declare a method as `native`. Class `Hello` imports a native method `cayHello` from the `HelloImpl` DLL [1], and defines a normal static Java method `sayHello` that will print the greeting on the standard output:

```
class Hello {
  native cayHello ();
  static void sayHello () {
    System.out.println("Hello World!");
  }
  static {
    System.loadLibrary ("HelloImpl");
  }
}
```

In Java we cannot just call entries of an arbitrary DLL; Java dictates how native methods should look like. For this example `HelloImpl.dll` should have entry `Java_Hello_cayHello` with signature[2]:

```
JNIEXPORT void JNICALL
  Java_Hello_cayHello ([in]JNIEnv* env, [in]jobject this)
```

[1] A dynamic-link library (DLL) is an executable file that acts as a shared library of functions on the Win32 platform.

[2] Throughout this paper we specify the C signatures for the different JNI procedures using IDL attributes such as `[in]` and `[out]`.

In general, each native method implementation gets an additional `JNIEnv` argument, and a `jobject` object reference (or a `jclass` class object for static methods). The `JNIEnv` argument is an interface pointer through which the native method can communicate with Java, for instance to access fields of the object, or to call back into the JVM. The `jobject` (or `jclass`) reference corresponds to the `this` pointer (or class object) in Java.

Procedure `Java_Hello_cayHello` first gets the class object for the `Hello` class, then gets the methodID for the static `void sayHello()` method, and finally calls it:

```
jclass cls = (*env) -> GetObjectClass (env, this);
jmethodID mid = (*env) ->
  GetStaticMethodID (env, cls, "sayHello" "()V");
(*env) -> callStaticVoidMethod (env, cls, mid);
```

In the above example the initiative came from the Java side. It is also possible to start on the C side.

The `JNI_CreateJavaVM` function takes an initialization structure and loads and initializes a JVM instance. It returns both a `JavaVM` interface pointer to the JVM and a `JNIEnv` interface pointer to the main thread running in that JVM instance:

```
jint JNI_CreateJavaVM
  ([out]JavaVM** v, [out]JNIEnv** e, [in]JavaVMInitArgs* a);
```

To call `sayHello` in this manner, we first load and initialize a JVM using `JNI_CreateJavaVM` to obtain a `JNIEnv` pointer, then call the `sayHello` method exactly as we did before except that we get the class object via the `FindClass` function and finish by destroying the JVM:

```
JNIEnv* env;
JavaVM* jvm;
JavaVMInitArgs args;

new_InitArgs (&args);
JNI_CreateJavaVM (&jvm, &env, &args);
jclass cls = (*env) -> FindClass (env, "Hello");
jmethodID mid = (*env) ->
  GetStaticMethodID (env, cls, "sayHello" "()V");
(*env) -> callStaticVoidMethod (env, cls, mid);
(*jvm) -> DestroyJavaVM(jvm);
```

Even though we have left out many details (including dealing with errors) it is clear that interacting with Java at this level of abstraction is quite mind numbing. In particular, the fact that we have to thread around the `JNIEnv` pointer in every call is something we would like to get rid of. In addition, encoding of the result-type in method calls (such as `callStatic`*`Void`*`Method`) and the distinction between static and instance methods is something we want to hide to the user.

3 Essential FFI in Four Easy Steps

We will introduce some basic JNI concepts that underpin our JNI binding for Haskell while reviewing the basics of the Hugs/GHC foreign function interface. For a more thorough explanation of the FFI we refer to our ICFP99 paper [2].

3.1 Importing Function Pointers

In order to call a foreign function from Haskell, we import externally defined functions into Haskell, either by static linking or by dynamic linking.

For example, we can statically import the `LoadLibraryA` entry of the Microsoft Windows kernel32 DLL

```
HMODULE LoadLibraryA([in,string]char*);
```

by giving the following foreign import declaration:

```
type CString = Addr
type HMODULE = Int

foreign import stdcall "kernel32"
  "LoadLibraryA" primLoadLibrary :: CString -> IO HMODULE
```

This foreign declaration defines a Haskell function `primLoadLibrary`, that takes the address of a null terminated string which contains the filename of an executable module, dynamically loads the module into memory and returns a handle to the requested module.

As its name suggests, function `primLoadLibrary` is rather primitive. Using the functions `marshallString` and `freeCString` we can easily write a friendlier version that takes a normal Haskell string as its argument:

```
marshallString :: String -> CString
freeCString :: CString -> IO ()

loadLibrary :: FilePath -> IO HMODULE
loadLibrary = \p -> do{
  s <- marshallString p;
  h <- primLoadLibrary s;
  freeCstring s;
  return h;
  }
```

The kernel32 DLL also exports a function `GetProcAddress` that given a module handle and a pointer to a null-terminated string containing the desired function name returns the address of that function.

```
  FARPROC GetProcAddressA ([in]HMODULE, [in,string]char*);
```

To use function `GetProcAddr` from Haskell we declare the following foreign import:

```
 type FARPROC = Addr

  foreign import "kernel32" "GetProcAddressA"
    primGetProcAddress :: HMODULE -> CString -> IO FARPROC
```

Again we can easily construct a Haskell-friendly version of `primGetProcAddress` that takes a normal Haskell string as its argument:

```
  getProcAddress :: HMODULE -> String -> IO FARPROC
  getProcAddress = \h -> \n -> do{
    pn <- marshallString n;
    pf <- primGetProcAddress h pn;
    freeCString pn;
    return pf;
  }
```

The explicit freeing of `CString`s is already getting tedious but don't despair; in section 3.2 we will instruct the Haskell garbage collector to do this automatically.

When interfacing with Java using JNI, we have to deal with different JVM implementations and different versions of the same JVM. Therefore, we don't want to commit ourselves to a particular VM or version too early. The `loadLibrary` and `getProcAddress` functions together with *dynamic* foreign function import allow us to delay the choice of a particular JVM to the very last moment.

The dynamic variation of the foreign import declaration defines a function that may be used to unmarshall an external function pointer into a Haskell function.

We can use a function pointer that represents the `JNI_CreateJavaVM` DLL entry by declaring the dynamic foreign import `mkPrimCreateJavaVM`. This function will coerce the low-level function pointer into the corresponding Haskell function:

```
  type JavaVM = Addr
  type JNIEnv = Addr
```

```
foreign import dynamic mkPrimCreateJavaVM
  :: FARPROC -> IO (JavaVM -> JNIEnv -> Addr -> IO Int)
```

To write a flexible version of `createJavaVM` we need several helper functions. Function `newJavaVMInitArgs` creates a default JVM initialization structure; it comes with a corresponding free routine `freeJavaVMInitArgs`:

```
newJavaVMInitArgs :: IO Addr
freeJavaVMInitArgs :: Addr -> IO ()
```

Function `newResultAddr` allocates space for a return value; it also comes with a corresponding free routine `freeResultAddr`:

```
newResultAddr :: IO Addr
freeResultAddr :: Addr -> IO ()
```

And finally, a function `deref` to shorten a pointer indirection:

```
deref :: Addr -> IO Addr
```

Using these helper functions, we can complete the tedious code for `createJavaVM`:

```
type JavaVM = Addr
type JNIEnv = Addr

createJavaVM :: FilePath
  -> IO (JavaVM, JNIEnv)
createJavaVM = \p -> do{
  h <- loadLibrary p;
  a <- getprocAddress h "JNI_CreateJavaVM";
  pArgs <- newJavaVMInitArgs;
  pJavaVM <- newResultAddr;
  pJNIEnv <- newResultAddr;
  mkPrimCreateJavaVM a pJavaVM pJNIEnv pArgs;
  freeJavaVMInitArgs pArgs;
  jvm <- deref pJavaVM;
  freeResultAddr pJavaVM;
  jnienv <- deref pJNIEnv;
  freeResultAddr pJNIEnv;
  return (jvm,jnienv);
}
```

There are numerous other situations where we want to call dynamic function pointers. For instance both the `JavaVM` and the `JNIEnv` pointers are a double indirection to a table of function pointers (a *vtable*). The function `deref` we used earlier and the function `getProcAt t i` that returns the `i`-the function pointer in table `t` make it quite easy to call vtable entries. We will give an example of this in section 3.3.

3.2 Interfacing With the GC Using Foreign Objects

One of the nice things about languages such as Haskell and Java is that the garbage collector silently takes care of freeing unused resources. It seems like a great loss therefore that when we import foreign functions, we import the hassle of explicit resource management as well.

What we would like to do instead is to have a hook into the Haskell garbage collector so that we can tell it how to free external resources once they become garbage. As usual, an extra level of indirection does the job. A *foreign object* [5] is a smart pointer to the outside world:

```
data ForeignObject
```

When constructing a foreign object, we pass it a finalizer action that is called by the Haskell garbage collector once the object becomes unreachable. This gives the object a chance to do its last belch.

```
new_ForeignObject :: Addr -> (Addr -> IO ()) -> IO ForeignObject
```

For example, we can define a smart `String` marshaller that frees the string automatically once it has become garbage (to avoid confusion, the 'dumb' version of `marshallString` is qualified with prefix `Prim`):

```
type CString = ForeignObject

marshallString :: String -> IO CString
marshallString = \s -> do{
  ps <- Prim.marshallString s;
  newForeignObject Prim.freeString ps;
}
```

Note that imported functions can directly take a `ForeignObject` as argument. Extracting the underlying `Addr` first, would be rather dangerous as the garbage collector could then free the object before the call is actually made.

3.3 Exporting Closures

Life would be rather boring if we could only *import* foreign functions. Things start to become interesting when we can masquerade Haskell functions as plain function pointers and pass those the outside world.

A *static* foreign export declaration tells the Haskell compiler to export a declared Haskell function behind a C-callable function interface. For example (assuming we are on the Win32 platform) we can create a DLL with an entry `Reverse` that reverses a C string by unmarshalling it, reversing the resulting Haskell string and then marshalling the reversed string back to a C string as follows:

```
type CString = Addr

foreign export stdcall "Reverse"
 primReverse :: CString -> IO CString

primReverse =
  (  unmarshallString ## return.reverse ## marshallString )
```

The double hash operator `##` is just reverse monadic composition: `f ## g = \a -> do{ b <- f a; g b}`. Later, we will also use the single hash operator `#`, which is just reverse function application `a # f = f a`.

Haskell *interpreters*, such as for example Hugs, do not naturally support statically exported functions. Fortunately, we can also *dynamically* export Haskell functions as if they are C function pointers[3]. Dynamically exported Haskell functions are extremely powerful and make programming with callbacks very easy, as the following example demonstrates.

The `JNIEnv` method `RegisterNatives` dynamically registers new native methods with a running JVM. The `RegisterNatives` entry takes a reference to a class object in which the native methods will be registered, and an array of `JNINativeMethod` structures that contain the names, types and implementations of the native methods:

```
typedef struct {
    [string]char* name;
    [string]char* signature;
    [ptr]FARPROC fnPtr;
} JNINativeMethod;

jint RegisterNatives
  ( [in]JNIEnv* env
  , [in]jclass clazz
  , [in,size_is(nMethods)] JNINativeMethod* meths
  , [in]jint nMethods
  );
```

To be able to call `RegisterNatives` in Haskell through a `JNIEnv` pointer, we must first dereference the `JNIEnv` pointer, fetch the 215-th entry in the function table and coerce that into a Haskell callable function. Then we pass

[3] In section 4, we will show how static exports can be simulated using dynamic exports and a statically exported interpreter component.

the resulting function the required arguments to compute the result:

```
getProcAt :: Int -> Addr -> IO FARPROC

foreign import dynamic mkRegisterNatives
  :: FARPROC -> IO (JNIEnv -> Jclass -> Addr -> Int -> IO Int)

registerNatives :: JNIEnv -> Jclass ->
  Addr -> Int -> IO ()
registerNatives = \env -> \clazz -> \meths -> \n ->  do{
    f <- (deref
          ## getProcAt 215
          ## mkRegisterNatives
         ) env;
    f env clazz meths n;
  }
```

Suppose we want to register a native function with signature `void SayHello (JNIEnv*, jobject)` via the `registerNatives` function. We can easily construct the required function pointer from a Haskell function with type `JNIEnv -> Jobject -> IO ()` by using function `mkSayHello`. The latter function is defined using a `foreign export dynamic` declaration:

```
foreign export dynamic
  mkSayHello :: (JNIEnv -> Jobject -> IO ()) -> FARPROC
```

Thus `mkSayHello` is a Haskell function that takes *any* Haskell function of type `JNIEnv -> Jobject -> IO ()` and at run-time creates a C function pointer, that when called will call the exported Haskell function:

```
sayHello_ :: FARPROC
sayHello_ = mkSayHello sayHello

sayHello = \env -> \this -> do{ putStr "Hello From Haskell!"; }
```

Assuming a function `marshallJNINativeMethods` that builds an array of `JNINativeMethod` structures from a Haskell list of name/signature/function pointer-triples

```
marshallJNINativeMethods :: [(String,String,FARPROC)] -> IO Addr
```

it is not hard to define a function `registerHello` that registers the Haskell function `sayHello` as the native implementation of a Java method `void sayHello ()`:

```
registerHello = \env -> \clazz -> do{
 p <- marshallJNINativeMethods [("SayHello","()V", sayHello_)];
  env # registerNatives p 1 clazz;
  }
```

3.4 Interfacing With the GC Using Stable Pointers

When we pass Haskell values (such as dynamically exported functions), to the outside world, we must prevent the Haskell garbage collector from moving them around inside the Haskell heap during a collection. Otherwise, an external function that holds on to the value will suddenly point to a completely different value. In addition, since we have no control over how many copies are made of exported objects, the garbage collector cannot automatically free them anymore. Haskell values are therefore hidden behind an extra level of indirection called *stable pointers* [5,9]. For the remainder of this paper we don't need to understand stable pointers in detail except that we know that they exist behind the scenes.

Now that we have covered the basics of interacting with the outside world and of JNI, we can start looking at the main subject of this paper, the bridge between Java and Haskell via the Java Native Interface (JNI).

4 Calling Java From Haskell and Haskell from Java

Calling Java from Haskell and Haskell from Java is in principle no different from calling Java from C and calling C from Java.

Calling Java from Haskell poses no problems as long as we have static and dynamic foreign import. In order to let Java call Haskell, the only provision is that Java can either load a DLL that contains Haskell implementations of the required native methods (to implement `native` methods directly), or some other DLL that can dynamically register the native methods using the `JRegisterNatives` JNI entry we have seen in section 3.3.

The *DietHEP* component [11] leverages on the foreign function interface we defined for Haskell in previous papers [3] and allows us to view *any* Haskell module as an ordinary DLL. Clients of DietHEP see no difference (except for the extra flexibility of specifying the calling convention at runtime when using `GetProcAddressEx`) between using an ordinary DLL via the `kernel32.dll` or using the `DietHEP` primitives. Going through the extra level of indirection of DietHEP allows us to abstract from the underlying Haskell implementation (provided of course that it supports the DietHEP interface).

The `LoadLibrary` function takes the name of a Haskell module (or GHC compiled binary), loads it and returns a *handle* to the module. Function `GetProcAddress` takes that handle and a function name and returns a 'foreign export'-ed version of the requested Haskell function.

```
  [dllname("DietHEP.dll")]
```

```
module DietHEP { typedef enum {stdcall, ccall} CALLCONV;

   HMODULE LoadLibrary ([in,string]char* n);
   FARPROC  GetProcAddress
     ([in]HMODULE m, [in,string]char* n);
   FARPROC  GetProcAddressEx
     ([in]CALLCONV c, [in]HMODULE m, [in,string]char* n);
};
```

To register native Haskell methods via DietHEP, we wrap the DietHEP component into a Java class `DietHEP` in the expected way:

```
class DietHEP {
  static native int LoadLibrary (String n);
  static native int GetProcAddress (int m, String n);
  ...
}
```

We also assume that we have a class `JNI` that reflects (amongst others) the JNI entry `registernatives` back into Java as the native method `RegisterNatives`:

```
class NativeMethod {
  String name; String signature; int fnPtr;
  ...
}

class JNI {
 static native void RegisterNatives (Class c, NativeMethod[] ms);
  ...
}
```

Using these two classes, the `Hello` example of section 2.1 can be written to use the Haskell function `Hello.sayHello` by first loading the Haskell module `Hello` that contains the definition of the `sayHello` function, and then registering the dynamically exported function pointer for `sayHello` with the Java `Hello` class:

```
class Hello {
 public static native void SayHello ();
 static {
   int h = DietHEP.LoadLibrary ("Hello");
   int sayHello = DietHEP.GetProcAddress (h,"sayHello");
   JNI.RegisterNatives
     ( Hello.class
     , new NativeMethod[]{
        new NativeMethod ("SayHello","()V", sayHello)
        }
      );
    }
  }
```

5 A User Friendly Library

At this stage, we have shown the bare bones of the integration between Haskell and Java. We will make the binding more user friendly in a number of steps. First we will abstract away commonly occurring patterns. Second, we hide the explicit threading of the `JNIEnv` argument through all JNI related code. Next, we will use overloading to hide the encoding of result-types in JNI methods.

5.1 Combining JNI calls

Calling an instance method `void Foo()` on a Java object `obj` and a `JNIEnv` pointer `env` is a three step process:

(i) Get the class object for `obj` using the JNI method `GetObjectClass`:

```
cls <- env # getObjectClass obj
```

(ii) Look up the methodID for the method via its name and type description using the JNI method `GetMethodID`:

```
mid <- env # getMethodID cls "Foo" "()V"
```

(iii) Call the method using the JNI method `Call`*`<t>`*`Method` where *`<t>`* is the result type of the method, passing the object, the methodID and the actual arguments:

```
env # callVoidMethod obj mid nullAddr
```

To call a static or instance method, or to get or set a static or instance field, we have to go through the same steps; fetch a class object, look up a method or field ID based on the name and signature, and eventually perform the desired operation. For the sake of presentation however, we will ignore invoking static methods until section 5.8. Space limitations preclude us from explaining (static) field access, but they don't require new concepts to be dealt with.

Our first step in simplifying interaction with JNI is to combine the basic 3-step JNI call sequence into a family of functions `call`*`<t>`*`Method`. These functions take a `Jobject` pointer, the name and signature of the method or field, a `ForeignObject` pointer to an array of arguments, and return a value of *`<t>`* (where as usual, we qualify the primitive `call`*`<t>Method`* with `Prim`):

```
call<t>Method
  :: Jobject -> String -> String -> ForeignObject
```

```
    -> JNIEnv -> IO <t>
  call<t>Method =
    \this -> \name -> \sig -> \a -> \env -> do{
      cls <- env # getObjectClass this;
      mid <- env # getMethodID cls name mid;
      env # Prim.call<t>Method this mid a;
    }
```

Our next step will be to eliminate the JNIEnv parameter.

5.2 Hiding JNIEnv

In previous examples we have seen that most interaction with Java goes via the JNIEnv interface pointer, and that all methods of the JNIEnv interface also take a 'this' pointer (or a class object pointer) as an additional argument. Having to thread env around is rather tedious, so we are willing to move a few mountains at this moment to save a lot of work later.

The JNIEnv pointer is only valid in its associated thread, so we cannot put it in a global variable just like that. The JavaVM pointer however *does* remain valid across different threads and we can safely put that in a global variable. To do that, we need to get hold of the JavaVM pointer in which the current thread is running, and we postulate that (current and future) JNI implementations provide some way of doing this. We will follow one of the suggestions[4] of Liang ([8], section 8.1.4) that relies on the fact that the current JVM releases do not support the creation of more than one JVM instance inside a single process ([8], page 254).

The JNI library function getCreatedJavaVMs returns a list of all currently running JVMs and as we argued above, we may safely assume that getCreatedJavaVMs always returns a singleton list containing the single running JavaVM. Hence getCreatedJavaVMs is a pure function and thus we can use unsafePerformIO to create a global 'variable' javaVM that contains the pointer to the running JVM:

```
javaVM :: JavaVM
javaVM = unsafePerformIO $ do{
  [javaVM] <- getCreatedJavaVMs;
  return javaVM;
}
```

The JavaVM interface entry attachCurrentThread returns the JNIEnv pointer for the current thread. Using global variable javaVM we can now easily and safely fetch the currently valid JNIEnv pointer in any context:

[4] The book gives several other suggestions, but this one is the simplest to implement. Another possibility would be to capture the current JVM in the JNI_OnLoad event handler.

```
attachCurrentThread :: JavaVM -> IO JNIEnv

getJNIEnv :: IO JNIEnv
getJNIEnv = javaVM # attachCurrentThread
```

Now that we can obtain the right `JNIEnv` in any context, we don't have to supply the `JNIEnv` argument to `call<t>Method` we defined previously any more, since we can fetch it when we need it (again we qualify the more primitive version of a function with `Prim`):

```
call<t>Method
  :: String -> String -> ForeignObject -> Jobject -> IO <t>
call<t>Method = \name -> \sig -> \args -> \this -> do{
    env <- getJNIEnv;
    env # Prim.call<t>Method this name sig args;
  }
```

5.3 Overloading argument types

Before we start using overloading to simplify JNI, we have to make a short digression on marshalling. In section 3.1, we assumed that we had a function `marshallString :: String -> IO CString` that takes a Haskell string and returns a pointer to a null terminated array of characters.

The primitive string marshall function `marshallString` itself is defined in terms two more primitive functions. Function `marshallByRefChar :: Char -> [Addr -> IO ()]` returns (in this case a singleton) list of functions that writes a character at the indicated addresses. Function `writeAt` takes a list of `Addr -> IO ()` functions and a start address and calls each function in the list to write the value at subsequent addresses:

```
writeAt :: [Addr -> IO ()] -> Addr -> IO ()
```

To marshall a string, we allocate enough memory to hold the string, generate a list of marshall functions for each character in the string and then write each one into the allocated memory:

```
marshallString :: String -> IO Addr
marshallString = \s -> do{
  let cs = s++[chr 0];
  p <- malloc $ length cs;
  writeAt (concat (map marshallByRefChar) cs) p;
  return p;
}
```

The free function for strings simply calls the system function `free :: Addr -> IO ()`, which deallocates a memory block that was previously allocated by a call to `malloc`:

```
freeString :: Addr -> IO ()
freeString = \s -> do{ free s; }
```

In general, we have a family of functions `marshall`*<t>*, `marshallByRef`*<t>*, `free`*<t>* for each *<t>* that we can marshall from Haskell to Java:

```
marshall<t> :: <t> -> IO Addr
marshallByRef<t> :: <t> -> [Addr -> IO ()]
free<t> :: Addr -> IO ()
```

At this moment bells should start to ring and alarms should go off: "ad-hoc overloading", this is *exactly* what Haskell type classes were invented for!

So we define a new type class `GoesToJava` with three methods `marshall`, `marshallByRef` and `free`. To be able to overload `free` as well, we add an extra phantom argument to the `free` function in the class:

```
class GoesToJava a where {
  marshall :: a -> IO Addr;
  marshallByRef :: a -> [Addr -> IO ()];
  free :: a -> Addr -> IO ();
}
```

and make all types that can be marshalled from Haskell to Java (`Char`, `Int`, `Integer`, `Double`, `Float`, `Jobject`, `String`, ...) an instance of the `GoesToJava` class:

```
instance GoesToJava <t> where {
  marshall = marshall<t>;
  marshallByRef = marshallByRef<t>;
  free = \a -> free<t>;
 }
```

The `marshallByRef` instances for `Double` and `Integer` return a two element list with the second entry being `\a -> do{ return () }` to ensure that the `marshall` function will allocate enough memory to hold a `double` or a `long`.

We continue to overload `GoesToJava` for tuples to encode Java methods with zero, or more than one argument. This is impossible if we use curried functions to represent multiple argument Java methods in Haskell.

```
instance GoesToJava () where {
  marshall = \() -> do{ malloc 0; };
  marshallByRef = \() -> [];
  free = \() -> \p -> do{ free p; };
}
```

```
instance
 (GoesToJava a, GoesToJava b) =>
  GoesToJava (a,b) where {
    marshall = \(a,b) -> do{
      let ab = marshallByRef a ++ marshallByRef b;
      p <- malloc $ length ab;
      writeAt ab p;
      return p;
    };
    marshallByRef = \(a,b) ->
      marshallByRef a ++ marshallByRef b;
    free = \(a,b) -> \p -> do{
      (deref ## free) p;
      (deref ## free) (incrAddr p);
      free p;
    }
```

By using the `GoesToJava` class we don't have to marshall arguments to a JNI call into an array anymore, but we can simply write (again qualifying the previous version of `call<t>Method` with `Prim`):

```
call<t>Method :: GoesToJava a =>
  String -> String -> a -> Jobject -> IO <t>
call<t>Method =
  \name -> \sig -> \a -> \this -> do{
    p <- marshall a;
    r <- newForeignObject p (free a);
    env <- getJNIEnv;
    env # Prim.call<t>Method this name sig r;
  }
```

5.4 Overloading result types

The signature of JNI calls above obviously cries out to be overloaded as well by means of a type class `ComesFromJava` with an instance for each `<t>` (`()`, `Jobject`, `Bool`, `Byte`, `Char`, `Short`, `Int`, `Integer`, `Float`, `Double`) that can be returned by a JNI method invocation:

```
class ComesFromJava b where {
  callMethod :: GoesToJava a =>
    String -> String -> a -> Jobject -> IO b;
}

instance ComesFromJava <t> where {
  callMethod = call<t>Method;
}
```

At this point we have reduced the complexity for JNI calls quite significantly when compared with the most primitive form, but there are still ample possibilities to improve. The next thing that we are going to attack is the explicit passing of type descriptor strings, which at the moment is rather unsafe.

5.5 Generating type descriptors

The problem with the current version of function `callMethod` is that there is no connection between the *value* of the type descriptor string, which denotes the Java type of a field or method, and the corresponding Haskell *types* of the argument and result of the `callMethod` function. For example, the following definition is accepted by the Haskell type-checker:

```
foo :: () -> Jobject -> IO Int
foo = \() -> \obj -> do{ obj # callMethod "foo" "(I)V" (); }
```

Unfortunately, it leads to a runtime error because the `foo` method is called as if it has Java type `void foo(int)` instead of at the correct type `int foo()` (which corresponds to the type descriptor string should be `"()I"`).

We will generate type descriptors by a family of functions `descriptor`*`<t>`* `::` *`<t>`* `-> String`[5]. As we did in the previous sections, we define a Haskell type class, in this case called `Descriptor`, that witnesses for which *`<t>`* we can obtain a descriptor:

```
class Descriptor a where { descriptor :: a -> String; }

instance Descriptor <t> where { descriptor = descriptor<t>; }
```

The type descriptor for the primitive JNI type `Bool` is `"Z"`:

```
descriptorBool :: Bool -> String
descriptorBool = \z -> "Z"
```

The other basic types are mapped as follows: `Byte` ↦ `"B"`, `Char` ↦ `"C"`, `Short` ↦ `"S"`, `Int` ↦ `"I"`, `Integer` ↦ `"L"`, `Float` ↦ `"F"`, `Double` ↦ `"D"`.

Type descriptors for methods arguments are formed by concatenating the type descriptors of the individual arguments:

```
instance Descriptor () where { descriptor = ""; }

instance (Descriptor a, Descriptor b)
```

[5] This trick is used in the Haskell module `Dynamic` as well, and is attributed to the legendary hacker Lennart Augustsson.

```
    => Descriptor (a,b) where {
         descriptor = \(a,b) -> descriptor a ++ descriptor b;
  }
```

Type descriptors for methods are formed by enclosing the type descriptor for the arguments inside parenthesis followed by the type descriptor of the result type, with the exception that `V` is used for the void return type:

```
methodDescriptor
  :: (Descriptor a, Descriptor b) => a -> b -> String
methodDescriptor = \a -> \b ->
  concat
    [ "(", descriptor a ,")"
    , if null (descriptor b)
      then "V"
      else descriptor b
    ]
```

To deploy function `methodDescriptor` we need to jump one more hurdle. In order to build the method descriptor string, we need to know the type of the result of a call. Since the `methodDescriptor` function does not need the value of its arguments to compute a type descriptor, we can recursively build the right method descriptor from a call which will take that descriptor as an argument:

```
callMethod
  :: ( Descriptor a, GoesToJava a
     , Descriptor b, ComesFromJava b
     ) => String -> -> a -> Jobject -> IO b
callMethod = \name -> \a -> \this
  -> let{
    sig = methodDescriptor a (unsafePerformIO mb);
    mb = this # Prim.callMethod name sig a
} in mb
```

There are several things to note about the above code. First of all, it intimately relies on lazy evaluation, function `methodDescriptor` only depends on the *types* of its arguments and not on their values. Secondly, the call to `unsafePerformIO` is there solely to strip the `IO` from the result type of the `callMethod` function, which is completely safe.

5.6 Encoding Inheritance using phantom types

The `callMethod` function as we have reached by now is not yet as safe and flexible as it could be. In particular, we do not reflect the fact that in Java classes satisfy an inheritance relation. We can however encode inheritance using the same trick we used for COM components [6].

For the Java root class `Object` we define a phantom type `Object_` that contains the low level untyped `Jobject` pointer, and immediately hide `Object_` using type synonym `Object` (we will see in just a minute why):

```
data Object_ a = Object Jobject
type Object a = Object_ a
```

Type `Object_` is called a *phantom* type because its type parameter `a` is not used in its right hand side.

For every class *<class>* that extends its superclass *<super>*, we define a new (empty) phantom type *<class>*_ and a type synonym *<class>* that expresses that *<class>* extends *<super>*:

```
data <class>_ a
type <class> a = <super> (<class>_ a)
```

The definition for *<class>* expands to `Object_ (... (<super>_ (<class>_ a)) ...)`, which precisely expresses the inheritance relationship between `Object` and *<super>* and *<class>*.

An instance of the Java class *<class>* is represented by a value of type *<class>* `()`, which is still the same old `Jobject` pointer, but now guarded by a little type harness. The nice thing about this is that the Haskell type checker prevents us from using a method of the *<class>* class at an instance for which *<class>* is not a superclass.

Of course, when actually calling a method we use the `Jobject` pointer that is hidden inside the `Object` wrapper:

```
callMethod ::
  (Descriptor a, GoesToJava a , Descriptor b, ComesFromJava b)
  => String -> -> a -> Object o -> IO b
callMethod = \name -> \a -> \(Object this) ->
  Prim.callMethod name a this
```

5.6.1 Class descriptors

Using the above encoding of Java classes into Haskell types, we can generate type descriptor strings for Java classes by defining a function `descriptor`*<class>* for each Java class *<class>*:

```
descriptor<class> :: <class> () -> String
descriptor<class> = \_ -> "L<classdescriptor>;"

instance Descriptor (<class> ()) where {
  descriptor = descriptor<class>;
}
```

A *class descriptor* is the fully qualified Java class with all .-s replaced by /-s; for instance the class descriptor for the Java class `Object` is `java/lang/Object`.

For methods that take objects as arguments, JNI expects the class descriptor for the formal parameter and not of the actual argument. This means for example that if we call the method `Component add(Component comp)` of the Java class `Container` with an argument whose run-time class is `Frame` (a subclass of `Container`), we still must pass the class descriptor `Ljava/awt/Component;` instead of `Ljava/awt/Frame;`. The correct version of `add` silently casts its argument to a `Component ()` just before making the actual call to function `add`:

```
add :: Component a -> Container b -> IO (Component a)
add = \comp -> \this -> do{
  this # callMethod "add" (cast comp :: Component ());
}
```

Inside function `add` we use the function `cast :: a -> b` to change the first argument of type `Component a` (which could be any subclass of `Component`) to exactly `Component ()` so that we get the correct method descriptor. Function `add` itself remains completely type safe.

The above situation is the only case where we have to explicitly up an upcast in Haskell; the question remains how to deal with the occasional need for downcasting. In Java, it is possible to unsafely cast between arbitrary classes, which might raise a dynamic `ClassCastException`. Using the Haskell `cast` to do an illegal arbitrary cast will either result in a `ClassCastException` or other exception being thrown in Java (which is checked using the JNI entry `ExceptionOccured`), or a call to one of underlying the JNI entry calls will signal an error (usually by returning `NULL`). The Haskell wrapper will catch these and propagate the errors to the Haskell level by raising a `userError` when returning. In other words, using `cast` to simulate arbitrary Java casting in Haskell is no more safe or unsafe than arbitrary casting in Java.

5.7 Dealing with interfaces

Besides classes, Java also supports *interfaces*. The role of interfaces in Java is quite similar to that of type classes in Haskell. Saying that a Java class implements an interface is like making an instance declaration in Haskell.

To deal with interfaces in Lambada we introduce an empty Haskell type class for each Java interface, possibly by extending some one or more base interfaces, for example:

```
class ImageObserver a
class MenuContainer a
class Serializable a
class EventListener a
class EventListener a => WindowListener a
```

For each method in an interface, we generate a function that requires its argument to satisfy the class constraint, eg for the `MenuContainer` interface we define:

```
getFont :: MenuContainer a => () -> a -> IO (Font ())
getFont = \() -> \this -> do{
  this # callMethod "getFont" ();
}
```

For each Java class *<class>* that implements an interface *<interface>*, we make the corresponding Haskell type *<class>* an instance of the Haskell type class *<interface>* that corresponds to the Java interface *<interface>*. For instance, the Java class `Component` implements the interfaces `ImageObserver`, `MenuContainer`, and `Serializable`, so in Haskell we write:

```
instance ImageObserver (Component a)
instance MenuContainer (Component a)
instance Serializable (Component a)
```

Note that we have to be careful *not* to write `Component ()`, since every subclass of `Component` implements the interfaces that `Component` implements.

Since `Component` is an instance of `MenuContainer`, a call `c # getFont ()` is well-typed whenever `c` has type `Component ()`, `Container ()`, `Window ()`, `Frame ()`, or any subclass of these.

Some methods expect interface 'instances' as arguments:

```
void addWindowListener(WindowListener)
```

To deal with this situation, we have to be able to generate descriptor strings for interfaces. For each interface *<interface>* we define a new type *<interface>* that uniquely identifies that interface and make *<interface>* an instance of `Descriptor`. Of course an interface satisfies its own interface constraint (in reality, we have to do name mangling because in Haskell data types and type classes are in the same name spaces but here no confusion can arise):

```
data <interface>

descriptor<interface> :: <interface> -> String
descriptor<interface> = \_ -> "L<interfacedescriptor>;"
```

```
instance Descriptor <interface> where {
  descriptor = descriptor<interface>;
}

instance interface <interface>
```

Now we can (safely) coerce any object reference that implements a certain interface to that interface and obtain the correct descriptor.

5.8 *Static methods and class objects*

Up to now we have ignored static methods. A static or class method is a method that does not act on a *class* instance, but on the *class object* instance.

Calling a static method, say `void Foo()`, on a class object `cls` and given a `JNIEnv` pointer `env` is quite similar to calling an instance method, except that the method is not invoked on the object instance but on its class object:

(i) Look up the methodID for the static method via its name and type description using the JNI method `GetStaticMethodID`:

```
mid <- env # getStaticMethodID cls "Foo" "()V"
```

(ii) Call the static method using the JNI method `CallStatic<t>Method` where `<t>` is the result type of the method, passing the *class object*, the methodID and the actual arguments:

```
env # callStaticVoidMethod cls mid nullAddr
```

If we continue and mirror the overloading have done for instance methods, we will end up with a function `callStaticMethod` for invoking static methods on a class object:

```
callStaticMethod ::
  (Descriptor a, GoesToJava a , Descriptor b, ComesFromJava b)
  => String -> a -> Jclass -> IO b
```

Although in Java we can call a static method on an object (which the compiler translates to the right class), the recommended style is to call it on the class directly. By *requiring* the latter style, we can distinguish the types of `callStaticMethod` and `callMethod`. This enables us to introduce overloading to hide the distinction between calling static and instance methods:

```
class JNI j where {
  call
    :: ( Descriptor a, GoesToJava a
       , Descriptor b, ComesFromJava b
       ) => String -> a -> j -> IO b;
}
```

For instance method calls we can immediately call `callMethod`:

```
instance JNI (Object a) where { call = callMethod; }
```

For static methods we define a new phantom type `ClassObject` such that `ClassObject` *<class>* uniquely encodes the class object of a Java class *<class>*. Type `ClassObject` just wraps a type harness around an untyped `Jclass` pointer:

```
data ClassObject a = ClassObject Jclass
```

The `JNI` instance declaration for static methods peels off the type harness and calls `callStaticMethod`:

```
instance JNI (ClassObject a) where {
  call = \n -> \a -> (ClassObject clazz) ->
    clazz # callStaticMethod n a;
}
```

So now we can call static methods on class objects, but how do we obtain class objects in the first place? The `JNIEnv` entry `findClass` takes a descriptor for the Java class *<class>* and returns the class object for *<class>*:

```
getClassObject
  :: Descriptor (Object o) => (Object o) -> ClassObject (Object o)
getClassObject = \this ->
  unsafePerformIO $ do{
    env <- getEnv;
    clazz <- env # findClass (descriptor this);
    return (ClassObject clazz);
}
```

Since function `getClassObject` does not use its argument to compute the class object, we can define an overload class object 'constant' with the help of function `stripObject :: Object o -> o`:

```
classObject :: Descriptor (Object o) => ClassObject (Object o)
classObject = let{ c = getClassObject (stripObject c) } in c
```

5.9 Constructing object instances

The distinction between instances (subclasses of `java.lang.Object`) and class objects (subclasses of `java.lang.Class`) can be confusing at first, especially because we can also call non-static methods on class references. One such example are constructor functions to create class instances.

Without using overloading to generate method signatures, we can combine the steps to create a new object into a single function as follows:

```
new :: (GoesToJava a) => String -> Jclass -> a -> IO (Object o)
new = \sig -> \clazz -> \a -> do{
  cid <- env # getMethodId clazz "<init>" sig;
  p <- marshall a;
  r <- newForeignObject (free a) p;
  o <- env # newObject clazz cid r;
  return (Object o);
}
```

As we did before, we want to use function `methodDescriptor` to generate the method descriptor, and furthermore we want to avoid passing the class object as an argument to `new` since we can get that using function `getClassObject` (again we qualify the previous version of `new` with `Prim`):

```
new :: (GoesToJava a , Descriptor (Object o) )
 => a -> IO (Object o)
new = \a -> let{
    o = unsafePerformIO mo;
    clazz = getClassObject o;
    sig = methodDescriptor a o;
    mo = Prim.new sig clazz a;
  } in mo
```

In practice overloading `new` on its return type often requires manual resolution using explicit type signatures. For convenience we will therefor define an additional un-overloaded constructor function `new_`*`<class>`* for each `<class>` that is an instance of `Descriptor` by restricting the type of `new`:

```
new_<class> :: (GoesToJava a) => a -> IO <class>
new_<class> = new
```

6 Simulating anonymous inner classes

Haskell functions/actions are represented in Java as objects of class `Function`, which store the actual exported Haskell function pointer as an `int` in a private field `f` and implement a native method `call` that takes an `Object` array as argument, and an `Object` representing the `this` pointer (or class object) and returns an `Object` as result:

```
class Function {
  private int f;
  Function(int f){ this.f = f; }
  public native Object call
    (Object[] args, Object this);
  ...
}
```

A typical situation is that we have to create an event-handler that implements the `WindowListener` interface. To facilitate the construction of such a handler in Haskell, we define a Java class `WindowHandler` that has a private `Function` object field for each of its methods which are initialized by the constructor function. Each method in the `WindowHandler` class simply delegates its work to the appropriate `Function` object:

```
class WindowHandler
  implements WindowListener {
  private Function activated;
  ...;
  WindowHandler
    (Function activated, ...) {
      this.activated = activated;
      ...
  }
  void windowActivated(WindowEvent e){
    this.activated.call (new Object [] {e}, this);
  }
  ...
}
```

The JNI framework dictates that the native implementation of the `call` method is a function that takes the `JNIEnv` and the `Jobject` this-pointer as all native methods do, a `Jobject` argument that represents the argument array, the explicit `this` argument of the instance in which the call was made, and that returns a `Jobject`:

```
call :: JNIEnv -> Jobject -> (Jobject -> Jobject -> IO Jobject)
```

Function `call` uses the `this` pointer of the `Function` instance to retrieve the function pointer from the private field `f` and then calls that as a function of Haskell signature `(Jobject -> Jobject -> IO Jobject)`:

```
foreign import dynamic importFunction
  :: FARPROC
  -> IO (Jobject -> Jobject -> IO Jobject)

call = \env -> \this -> \args -> \t -> do{
  f <- env # getField this "f";
  importFunction f args t;
}
```

We create `Function` instances in Haskell by creating a C function pointer for the Haskell function we want to export using `exportFunction` and then use the Haskell JNI library to call the `new_Function` constructor:

```
  foreign export dynamic exportFunction
    :: (Jobject -> Jobject -> IO Jobject) -> FARPROC
  new_Function :: Int -> IO (Function ())
```

All this quite low-level and error prone, so let's use a type class to tidy up and do the grungy work for us.

First of all, we extend the `ComesFromJava` class with a new member function `unmarshallArray` that unmarshalls zero or more values from a `Jobject` pointer representing an `Object` array to a tuple of corresponding Haskell values, and a new member function `unmarshall` that marshalls a single `Jobject` pointer into its corresponding Haskell value:

```
  class ComesFromJava b where {
    callMethod :: GoesToJava a =>
      String -> String -> a -> Object o -> IO b;
    unmarshall :: Jobject -> IO b;
    unmarshallArray :: Jarray -> IO b
  }
```

For basic types *<t>* we simply use the unmarshaller for that type, or first get the single element from the array and unmarshall that:

```
 instance ComesFromJava <t> where {
  unmarshall = \o -> do{
   unmarshall<t> o;
   unmarshallArray = \args
    args # (getArrayElement 0 ## unmarshall);
  }
```

For pairs, triples, etc, we recursively unmarshall the elements from the array one at a time.

```
  instance
    (ComesFromJava a, ComesFromJava b)
    => ComesFromJava (a,b) where {
      unmarshall = unmarshallArray;
      unmarshallArray = \args -> do{
        a <- args # (getArrayElement 0 ## unmarshall);
        b <- args # (getArrayElement 1 ## unmarshall);
        return (a,b);
      };
  }
```

Function `marshalFun` takes a high-level Haskell function of type `(ComesFromJava a, GoesToJava b) => (a -> Object o -> IO b)`, turns it into a low-level function of type `Jobject -> Jobject -> IO Jobject` and then dynamically exports it using `exportFunction`:

```
marshallFun
 :: (ComesFromJava a, GoesToJava b)
 => (a -> Object o -> IO b) -> IO FARPROC
marshallFun = \f -> exportFunction $ \args -> \this -> do{
    a <- unmarshall args;
    t <- unmarshallObject this
    b <- f a t; marshall b;
  }
```

With the help of `marshallFun` we can finally make functions an instance of `GoesToJava`. To marshall a function from Haskell to Java, we first use `marshallFun` to get its function pointer representation, then we construct a new Java `Function` object for the pointer and marshall the resulting Java object to obtain a pointer we can throw over to Java:

```
instance
(ComesFromJava a, GoesToJava b)
 => GoesToJava (a -> Object o -> IO b)
  where {
    marshall = \f -> do{
      a <- marshallFun f;
      func <- new_Function (toInt a);
      marshall func
  }
}
```

Since functions are an instance of `GoesToJava`, writing event handlers in Java has now become completely painless. We can call the constructor for our `WindowHandler` example, passing it seven Haskell functions of type `Event e -> WindowHandler w -> IO ()` that handle each of the seven possible `Window` events:

```
new_WindowHandler
  ( \event -> \this -> do{
      this # activated event;
    } , ... )
```

This nearly exactly matches the corresponding code we would write in Java

```
new WindowHandler () { void activated (Event e) { ...; } ... }
```

7 Tool support

To reiterate the introduction of this paper, Java is an interesting partner for a foreign function interface because there are a lot of Java classes out there. However, even with the layers of abstractions we have introduced on top of

JNI, manually coding the Haskell callable stubs for the Java APIs we are interested in is not always practical.

Lambada provides tool support in concert with HDirect[3]. HDirect is IDL based, so provided we can derive an IDL specification from a Java class or interface, HDirect will take care of generating valid Haskell stubs. We do this by providing a tool that uses the Java Reflection API to generate IDL interfaces (annotated with a couple of custom attributes) corresponding to Java classes or interfaces. Provided with that information, HDirect is then able to automatically generate the Haskell callable stubs we have described in this paper. HDirect is also capable of generating Java callable stubs to Haskell implemented classes/methods.

At the Lambada implementation level, HDirect was also used to generate the low-level JNI stubs outlined in Section 3.

8 Related work

Claus Reinke was a remarkably early adaptor of JNI, as early as 1997 he proposed to use JNI to interoperate between Haskell and Java [10]. Reinke's system provides only one-way interop between Haskell and Java using the basic low-level JNI methods, ie he does not support calling back from Java to Haskell. The most important contribution of our work over his is our sophisticated use of overloading and other advanced language features to hide the idiosyncrasies of JNI. The integration that Lambada offers between Haskell and Java is very close to the tight integration offered by MLj [1], with the important exception that Lambada is implemented completely *within* Haskell, whereas MLj needs many language changes.

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