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Co-SRL: A Convex Optimization Algorithm for Anchor Localization in Wireless Sensor Networks

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Abstract

This paper proposed a Convex Optimization method which is called Co-SRL and is used to localize sensor location in Wireless Sensor Networks. Co-SRL can be used to help the node to localize a friend node or mobile node using anchors. In Co-SRL, convex optimization algorithm is used for the estimation of malicious node position. Simulation result shows that Co-SRL is both secure and robust, in an environment without colluding, Co-SRL can identify more than half of the malicious nodes; and in an environment with colluding, no more than 15% of malicious nodes can escape from the identification of our methods.

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1. Introduction and Related Works

As the rapid development of Internet and network technologies especially the wireless network technologies, now there are more and more people working, studying and entertainment from Internet via wireless network. But as the openness of wireless network, there are many security events occurred which lead great loss of money for users using wireless network. So for security consideration and being served as the evidence of computer crimes and further to find the location of the crimes, the localization technique is very important in wireless sensor networks.

Currently, there are mainly the range based localization methods and the range free localization methods in wireless networks. The rangebased localization methods as proposed in [10][8][4], localize nodes with property measurement. On the other hand, the range free localization methods proposed in [5][2][3][9] localize anchors without property measurement. We can tolerate that the website with bad content such as sex and violence are banned for the common good of society [5], but it may be a big pity that some websites with academic disputations are banned so not reachable only for political reason. [6] Moreover, one path may be broken down resulting in some websites not available to certain people, but actually they can still visit the webs after many attempts of trying other way only if the websites have other paths. In addition, sometimes though the websites are still available to us, we can always find a best way, i.e. with the highest speed to use. In one word, we need a put up way which can make the accessibility of the wireless sensor network more robust with better performance.

In fact, there are many papers that study the localization problems in sensor networks using optimization method. For example, [3] described localization information distributing propagation methods. [10] gave a time a localization discovery method in wireless sensor networks. The Co-SRL system is answer we find. In Co-SRL, there is not any center server, and only a register server generate identifications for the nodes who want to join the system, but not one is responsible for monitoring the behaviors single nodes. When strange nodes come to certain nodes and request for service, the requested nodes make their own opinions on response positively or negatively, according to their knowledge of the coming nodes.

This paper is arranged as follows: In section 2, we first introduce some basic principles of localization in wireless sensor networks and the optimization method, and then talk about the Co-SRL. In section 3, we present the simulation results of Co-SRL. And finally, we give the conclusion of this paper.

2. The Algorithm of Co-SRL

For each of the specific sub-problems, the problem specific assumptions will be presented when the problem itself is presented. The sensor network is $A = \{S_i, i = 1, \dots, n\}$ randomly deployed. S_i is an anchor, and $s_i (s_i = (s_{ix}, s_{iy}))$ is the position set of S_i . Both the mobile target and all communications between different anchors are bidirectional.

This paper gave a convex optimization based method to help node to localize anchors, given that at most M of them are malicious in a network. Our enhanced protocol, named Advanced Distance Bounding (AAD) protocol uses the high-speed DB technique. By Using of AAD, it can prevent from many network attacks aiming at wireless sensor networks such as wormhole attacks [1], Sybil attack [1] and Distance reduction attacks etc., because in this case, malicious anchors with faked position information will be detected immediately by a Malicious Node. Why? The reason can be seen as follows.

Basically, we use the convex optimization technique for the Malicious Node position estimation. As shown above, if there exist measurement negative error the intersection region Reg_t of the disks Dsk_{ti} may be empty. The disk Dsk_{ti} corresponds to the circle drawn with the position of anchor s_i being served the center and the distance between s_i and the target t as the radius. In this case, it will result in the increased distance estimation $DE'_{ti} = \frac{DE_{ti}}{1-ERR_{max}}$, if we increase a factor of $\frac{1}{1-ERR_{max}}$ to the distance estimates DE_{ti} for the malicious node MN_t . Let Dsk'_{ti} denotes the increased bound disks, and let BC'_{ti} denotes the corresponding increased bound circles. There must exist some malicious nodes in the non-empty region Reg'_t which is the intersection of Dsk'_{ti} .

All points in Reg'_t may likely be the position of Malicious Node. G_c of Reg'_t is used as the Malicious Node Localization, it would minimize the worst case error in estimation. Where G_c is the point satisfying that $\max_{x \in Reg'_t} \|x - G_c\| \leq \max_{x \in Reg'_t} \|x - y\|$, for any $y \neq G_c$. But it is very difficult to calculate the geometric center. As a result, to obtain G_c , we will find the solution of the following equation (4.1).

$$\begin{aligned} & \max_{x, \alpha} \alpha \\ & \text{subject to } \|x - S_i\|^2 \leq [r'_{ti}\beta]^2, i = 1, \dots, N, \quad (4.1) \\ & x \in E^2, \alpha \geq 0 \end{aligned}$$

Where, $\alpha + \beta = 1$.

In fact, it means that we use a common factor β to simultaneously shrink Dsk'_{ti} (all disks), as much as possible, it only need to ensure that there exist at least one member in the intersection among the disks. The first N constraints in equation (4.1) can ensure the condition of non-empty intersection among the disks. The objective function as shown in equation (4.1) is to maximize α , and finally to minimize β which is the shrinking factor. Clearly, there exists a unique optimal solution (x^*, α^*) for this convex optimization problem. Here x^* is named as the algebraic centre AC of the non-empty region Reg'_t . To simplify the solution of the in (4.1), we transform it into the following equation (4.2).

$$\min_{x, \alpha} -LM \cdot \alpha - \sum_{i=1}^N \log[(DE'_{ti}\beta)^2 - \|x - s_i\|^2] - \log(\alpha) \quad (4.2)$$

Where LM is the Lagrangian multiplier [7]. By using the algorithm which is presented in Algorithm 1, the above problem can be efficiently solved.

As shown in Algorithm 1, in Line 3, it uses the newton method [7] with tolerance $ERR = 1 \times 10^{-6}$ for the minimization.

The α value keeps increasing in algorithm 1, which will result in the reduction of Dsk'_{ti} . Finally, it is stopped and get the solution x^* (the center of Reg'_t) after sometimes of repeat when Reg'_t has been reduced greatly. x^* will be served as the approximation of the localization of the Malicious Node.

Algorithm 1. The optimization method to obtain an approximation to the geometric center

- 1: Initialize x , $LM = LM^{(0)} = 1.00 > 0$, $\mu = 10.00 > 0$, $ERR = 1 \times 10^{-6} > 0$.
- 2: repeat
 - 3: Beginning at x , calculate $x^*(LM)$ by minimizing the objective in (4.2).
 - 4: Update $x := x^*(LM)$; $LM := \mu \cdot LM$.

Algorithm 1: The optimization algorithm

In this research, it is assumed that the malicious node MN_t has adequate ability to perform the localization operations in Algorithm 1. Otherwise, if it cannot perform the localization operations, then the anchors can perform the localization operations.

3. Simulation Results

The proposed method has been implemented using the Matlab software. We apply the algorithm into the design the Co-SRL system which is an application overlay of the physical networks. It is constructed by the nodes and their accessible parts of the Internet. every peer in the system is a proxy which voluntarily relay the traffic of pass through it. Each node of Co-SRL is composed of 4 subsystems: The Functional-Module, the Calculation-Engine, the Friends-List and the Evaluation-Engine. Where, the Functional-Module is used for transferring traffic, registering and other functions. Our policy model for admission control is implemented in the calculation engine. When a new node comes, the calculate engine is responsible to calculate its trust value

and reflect to the friend list to record, and when the node interact with some new node, as the functional module offers or demands the service, the evaluate engine estimate the Recommendation and record in order for further calculation for the calculate engine. We construct a simulation field with $100m \times 100m$. The communication range of the anchors and the Malicious Node was restricted to 50mm. We randomly deployed anchors in the range of the Malicious Node t . We chose $ERR_{max}=0.1$ as the maximum value of measurement of the error proportion. $LP_{max}=1.0$ as the maximum value for the lie proportion. Therefore, a malicious anchor S_i with DE_{ti} as its distance estimate from t , set its distance $[DE_{ti}, DE_{ti} \cdot (1 + LP_{max})]$.

We also studied the effectiveness of our method with different values of LP. When there is no measurement errors, the results were averaged over 100 iterations for a given total anchors number in the range of the Malicious Node. For all the runs, the number of malicious anchors is no more than M . Our method localized the Malicious Node correctly in 100% of the cases and also caught the malicious anchors fully success. We do not present the result here because it is guaranteed by theory as well.

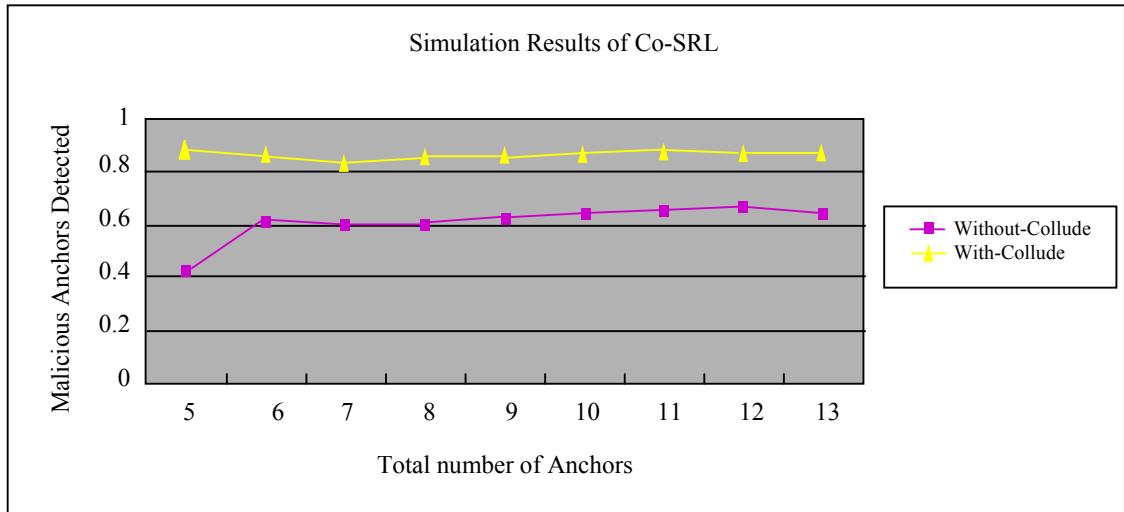


Figure 2. Simulation result of our method

When there exist error prone measurements, we compare the error in localization in our method with that in [8] and [3] to demonstrate the effectiveness of our method. Everyone knows that the method in [3] is prone to large errors when the anchors are lying [6]. Figure 2 shows the simulation results of our method, which are averaged over 50 runs. If the number of the anchors in the range of t is given by N , the number of malicious anchors M belongs to $\{1, \dots, \lfloor \frac{N}{2} \rfloor - 2\}$. In the simulation, the anchors number N in the range of an Malicious Node belonged to $\{5, \dots, 13\}$. Figure 2 shows the average, fraction of malicious anchors that are gotten by the method proposed by us over 50 runs.

4. Conclusion

This paper proposed a Convex Optimization method which is called Co-SRL and is used to localize sensor location in Wireless Sensor Networks. Co-SRL can be used to help the node to localize a friendly target or

mobile node using anchors. Our localization method Co-SRL estimate the malicious nodes position using convex optimization methods.

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