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# Graphs with Girth at Least 8 are b-continuous<sup>1</sup>

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#### Abstract

A b-coloring of a graph is a proper coloring such that each color class has at least one vertex which is adjacent to each other color class. The b-spectrum of G is the set  $S_b(G)$  of integers k such that G has a b-coloring with k colors and  $b(G) = \max S_b(G)$  is the b-chromatic number of G. A graph is b-continous if  $S_b(G) = [\chi(G), b(G)] \cap \mathbb{Z}$ . An infinite number of graphs that are not b-continuous is known. It is also known that graphs with girth at least 10 are b-continuous. In this work, we prove that graphs with girth at least 8 are b-continuous, and that the b-spectrum of a graph G with girth at least 7 contains the integers between  $2\chi(G)$  and b(G). This generalizes a previous result by Linhares-Sales and Silva (2017), and tells that graphs with girth at least 7 are, in a way, almost b-continuous.

Keywords: b-chromatic number; b-continuity; girth; bipartite graphs.

# 1 Introduction

Let G be a simple graph (for basic terminology on graph theory, we refer the reader to [4]). A function  $\psi: V(G) \to \mathbb{N}$  is a proper k-coloring of G if  $|\psi(V(G))| = k$  and  $\psi(u) \neq \psi(v)$  whenever  $uv \in E(G)$ . Because we only deal with proper colorings in this text, from now on we refer to them as simply a coloring. We call the elements of  $\psi(V(G))$  colors. Given a color  $i \in \psi(V(G))$ , the set  $\psi^{-1}(i)$  is called color class i. We say that  $u \in V(G)$  is a b-vertex in  $\psi$ (of color  $\psi(u)$ ) if  $\psi(N[u]) = \psi(V(G))$ . If for some color  $c \in \psi(V(G))$ , the color class c does not contain b-vertices, we can obtain a (k-1)-coloring by changing the color of each vertex  $v \in \psi^{-1}(c)$  to another color in  $\psi(V(G)) \setminus \psi(N[v])$ . We say that this new coloring is obtained from the first

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one by cleaning color c. In a coloring such that we cannot apply this procedure, all color classes have at least one b-vertex. Such a coloring is called a b-coloring of G. Observe that an optimal coloring cannot have the number of colors decreased by the described algorithm; therefore every optimal coloring is also a b-coloring. In [8], the authors define the b-chromatic number of G, denoted by b(G), as the largest natural K for which G has a b-coloring with K colors. In the same article, the authors demonstrated that the problem of finding b(G) is NP-complete in general.

Another interesting aspect about b-colorings concerns its existence for every possible value between  $\chi(G)$  and b(G). In [8], the authors observe that the cube has a b-coloring using 2 colors and 4 colors, but has no b-coloring using 3 colors. Inspired by this result, in [9] it is shown that for any integer  $n \geq 4$  the graph obtained from the complete bipartite graph  $K_{n,n}$  by deleting the edges from a perfect matching has a b-coloring using 2 and n colors, but has no b-coloring using a number of colors between 2 and n. This motivates the definition of the b-spectrum of G, that is the set  $S_b(G)$  containing every integer k such that G has a b-coloring with k colors. A graph G is b-continous if  $S_b(G) = [\chi(G), b(G)] \cap \mathbb{Z}$ . In [2], they prove that for each finite subset  $S \subset \mathbb{N} - \{1\}$ , there exists a graph G such that  $S_b(G) = S$ , and also that deciding if a graph is b-continuous is NP-complete even if colorings with  $\chi(G)$  and h(G) colors are given.

Now, given a b-coloring with k colors, since each b-vertex has at least k-1 neighbors, there exists k vertices with degree at least k-1 (this would be a subset of k b-vertices of the k colors). So if we define m(G) as the largest positive integer k such that there exist at least k vertices with degree at least m(G)-1 in G, we have that  $b(G) \leq m(G)$ . This upper bound was introduced in [8], where the authors show that one can find m(G) in polynomial time using the degree list of the graph. Also, they prove that if G is a tree, then  $b(G) \geq m(G)-1$ , and that one can decide if b(G) = m(G) in polynomial time. Their result was later generalized for graphs with girth at least 7 [6] (the girth of G is the minimum length of a cycle in G). We also mention that there are many results that say that regular graphs with large girth have high b-chromatic number [3,5,13,3]. Indeed, the following conjecture is still open.

**Conjecture 1.1** If G is a d-regular graph with girth at least 5 and G is not the Petersen graph, then b(G) = d + 1.

Because of these results, it makes sense to investigate the b-continuity of graphs with large girth. Indeed, in [1] the authors prove that regular graphs with girth at least 6 and without cycles of length 7 are b-continuous, and in [11], they prove that every graph with girth at least 10 are b-continuous. Here, we improve their result to graphs with girth at least 8.

**Theorem 1.2** If G is a graph with girth at least 8, then G is b-continuous.

In addition, we prove that graphs with girth at least 7 are, in way, almost b-continuous.

**Theorem 1.3** If G is graph with girth at least 7, then  $[2\chi(G), b(G)] \cap \mathbb{Z} \subseteq S_b(G)$ .

Given a graph G and a b-coloring of G with k colors,  $k \geq \chi(G) + 1$ , the proof of Theorems 1.2 and 1.3 consists in trying to obtain a b-coloring with k-1 colors using simple recoloring procedures; when this is not possible, we get that the graph has a special structure and apply non-constructive arguments to obtain the desired b-coloring. We mention that the coloring problem is NP-complete for graphs with girth at least k, for every fixed  $k \geq 3$  [12]. This is why any proof of a result like Theorem 1.2 is expected to have a non-constructive part. In the next section, we present the basic definitions and results, in Section 3 we present our proofs, and in Section 4, we make some further comments on the proof and state some open questions.

## 2 Preliminaries

In [1], a vertex  $u \in V(G)$  is called a k-iris if there exists  $S \subset N(u)$  such that  $|S| \ge k - 1$  and  $d(v) \ge k - 1$  for every  $v \in S$  (observe Figure 1).

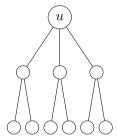


Fig. 1. In the figure, we presente a 4-iris.

This definition is important because of the following important lemma. Observe that the lemma also implies that if G and k satisfies the conditions, then  $b(G) \ge k$ .

**Lemma 2.1** ([1]) Let G be a graph with girth at least 6 and without cycles of lengh 7. If G has a k-iris with  $k \geq \chi(G)$ , then G has a b-coloring with k colors.

As we said before, given a b-coloring of G with k colors,  $k > \chi(G) + 1$ , we try to obtain a b-coloring of G with k-1 colors. However, this is not always possible, and when this happens, it is because we have a k-iris. Our theorem then follows from the lemma above. We mention that the constraint about not having cycles of length 7 appears only in the above lemma, but not on our proof. We now introduce the further needed definitions.

From now on, let G be a simple graph and  $\psi$  be a b-coloring of G with  $k > \chi(G) + 1$  colors. We say that u realizes color i if  $\psi(u) = i$  and u is a b-vertex. We also say that color i is realized by u. For  $x \in V(G)$  and  $i \in \{1, \ldots, k\}$ , let  $N^{\psi,i}(x)$  be the set of vertices of color i in the neighborhood of x, i.e.,  $N^{\psi,i}(x) = N(x) \cap \psi^{-1}(i)$ ; in fact, we omit  $\psi$  in the superscript since it is always clear from the context. This is also done in the next definitions. For a subset  $X \subseteq V(G)$ , let  $N^i(X) = (\bigcup_{x \in X} N^i(x)) \setminus X$ . Let  $B(\psi)$  denote the set of b-vertices in  $\psi$  and, for each  $i \in \{1, \ldots, k\}$ , let  $B_i = B(\psi) \cap \psi^{-1}(i)$  be the set of b-vertices in color class i.

Given a set K such that  $K \subseteq \psi^{-1}(i)$  for some  $i \in \{1, \ldots, k\}$ , we say that a color  $j \in \{1, \ldots, k\} \setminus \{i\}$  is dependent on K if  $N^i(B_j) \subseteq K$ ; denote by U(K) the set of colors depending on K. If  $K = \{x\}$ , we write simply U(x). Given  $x \in V(G) \setminus B(\psi)$ , if  $|U(x)| \geq 2$  we call x a useful vertex; otherwise, we say that x is useless. For  $j \in \{1, \ldots, k\}$ , we say that  $x \in V(G)$  is j-mutable if x is useless and there exists a color c such that we can change the color of x to c without creating any b-vertex of color j; we also say that color c is safe for (x, j). If there is no safe color for (x, j), we say that x is j-imutable.

# 3 Proofs

The next lemma is the main ingredient in our proof. Combined with Lemma 2.1, it immediatly implies Theorem 1.2.

**Lemma 3.1** Let G = (V, E) be a graph with girth at least 7. If G has b-coloring with k colors where  $k \ge \chi(G) + 1$ , then either G has a b-coloring with k - 1 colors, or G contains a (k - 1)-iris.

**Proof.** Our proof is similar to that made in [11], but we concentrate in one color that we want to eliminate.

Suppose that G does not have a b-coloring with k-1 colors; we prove that G has a (k-1)-iris. For this, let  $\psi$  be a b-coloring with k colors that minimizes  $|B_1|$  and then minimizes  $|\psi^{-1}(1)|$  (i.e., it firstly minimizes the number of b-vertices of color 1, then it minimizes the number of vertices of color 1). First, we prove that every  $x \in \psi^{-1}(1) \setminus B_1$  is useful. Suppose otherwise and let x be a useless vertex in color class 1, i.e.,  $|U(x)| \leq 1$ . If  $U(x) = \emptyset$ , then we can recolor x without losing any b-vertex, a contradiction since  $\psi$  minimizes  $|\psi^{-1}(1)|$ . And if  $U(x) = \{d\}$ , then we can obtain a b-coloring with k-1 by recoloring x and cleaning d, again a contradiction. Therefore, the following holds:

(i) Every 
$$x \in \psi^{-1}(1) \setminus B_1$$
 is useful.

Now, we choose any  $u \in B_1$  and analyse its vicinity in order to obtain the desired (k-1)-iris. For this, the following two claims are essential.

**Claim 3.2** Let  $j \in \{2, ..., k\}$ . If every  $x \in N^j(u) \setminus B_j$  is 1-mutable, then one of the following holds:

- (ii)  $N(u) \cap B_j \neq \emptyset$ ; or
- (iii) There exists a color  $d \in \{2, ..., k\} \setminus \{j\}$  such that d depends on  $N^j(u)$ , i.e.,  $N^j(B_d) \subseteq N^j(u)$ .

Proof of claim: Suppose that neither (ii) nor (iii) holds, and let  $\psi'$  be obtained from  $\psi$  by changing the color of each  $x \in N^j(u)$  to a color c safe for (x, 1). Because (iii) does not hold, we get that  $U(N^j(u)) \subseteq \{1\}$ . Therefore, at most one color loses all of its b-vertices, namely color 1, and since every  $x \in N^j(u)$  is 1-mutable, no b-vertices of color 1 is created. But because u is not a b-vertex in  $\psi'$  (it is not adjacent to color j anymore) and  $\psi$  minimizes  $|B_1|$ , we get that  $\psi'$  cannot be a b-coloring, which

means that we can obtain a b-coloring with k-1 colors by cleaning color  $1.\diamondsuit$ 

The following claim tells us that (ii) or (iii) actually always hold.

**Claim 3.3** (iv) Every 
$$x \in N^{j}(u) \setminus B_{j}$$
 is 1-mutable, for every  $j \in \{2, ..., k\}$ .

Proof of claim: Suppose, without loss of generality, that  $d \in \{2, ..., k\}$  is such that the colors in  $\{d+1, ..., k\}$  are exactly the colors that contains some 1-imutable vertex. We count the number of colors with b-vertices in the vicinity of u to get that in fact  $d \geq k$ . So, for each  $i \in \{d+1, ..., k\}$ , let  $w_i \in N^i(u)$  be a 1-imutable vertex. By definition, this means that, for each  $i \in \{d+1, ..., k\}$ , there exists some neighbor of  $w_i$  that would be turned into a b-vertex of color 1 in case we change the color of  $w_i$ ; let  $v_i$  be such a vertex. We then know that  $v_i \in \psi^{-1}(1) \setminus B_1$ , which by (i) gets us that  $|U(v_i)| \geq 2$ . By the definition of U(x) and the fact that every  $x \in \{v_{d+1}, ..., v_k\}$  is colored with color 1, we get:

(1) 
$$U(v_i) \cap U(v_\ell) = \emptyset$$
, for every  $i, \ell \in \{d+1, \ldots, k\}, i \neq \ell$ .

Now, we investigate the b-vertices around colors  $\{2,\ldots,d\}$ . By Claim 3.2, suppose, without loss of generality, that  $p \in \{2,\ldots,d\}$  is such that (ii) holds for colors in  $\{2,\ldots,p\}$ , while (iii) holds for colors in  $\{p+1,\ldots,d\}$ . For each  $i \in \{p+1,\ldots,d\}$ , let  $c_i \in \{2,\ldots,k\} \setminus \{i\}$  be a color depending on  $N^i(u)$ , which means that  $B_{c_i} \subseteq N(N^i(u))$ . Observe that, since G has no cycles of length 3, we get:

$$(2) {2,\ldots,p} \cap {c_{p+1},\ldots,c_d} = \emptyset$$

Also, because G has no cycles of length 4, we get  $c_i \neq c_\ell$  for every  $i \neq \ell$ , i.e.:

$$|\{c_{p+1}, \dots, c_d\}| = d - p$$

Finally, because G has no cycles of length smaller than 6, we get that:

(4) 
$$\{2,\ldots,p,c_{p+1},\ldots,c_d\} \cap \bigcup_{i=d+1}^k U(v_i) = \emptyset.$$

Now, recall that  $\psi(v_i) = 1$  for every  $i \in \{d+1,\ldots,k\}$ , and that  $c_i \neq 1$  for every  $i \in \{p+1,\ldots,d\}$ . This means that  $1 \notin \{c_{p+1},\ldots,c_d\} \cup \bigcup_{i=d+1}^p U(v_i)$ . By combining Equations (1) through (4), we get the following, which implies  $d \geq k$  as desired:

$$k-1 \ge |\{2, \dots, p\} \cup \{c_{p+1}, \dots, c_d\} \cup U(v_{d+1}) \cup \dots \cup U(v_k)|$$

$$= d-1 + \sum_{i=d+1}^k |U(v_i)|$$

$$\ge d-1 + 2(k-d).$$

Now, let  $N = (N(u) \cup N(N(u))) \setminus \{u\}$ . Observe that because (ii) or (iii) holds for every color  $\ell \in \{2, \ldots, k\}$ , we get that  $B(\psi) \subseteq N$ . Suppose that N[u] does not contain a (k-1)-iris, otherwise the proof is done. This means that at least one color in  $\{2, \ldots, k\}$ , say k, is such that (ii) does not hold for k, which by Claim 3.2 implies that (iii) holds, i.e., that there exists a color in  $\{2, \ldots, k-1\}$ , say 2, such that  $N^2(B_k) \subseteq N^2(u)$  (Observe Figure 2). Now, let  $w \in N^1(B_k)$ ; it exists since the vertices in  $B_k$  are b-vertices. By (i), there exists at least two colors in  $\{2, \ldots, k\}$ 

that depend on w. But because  $B(\psi) \subseteq N$ , we get a cycle of length at most 6, a contradiction.

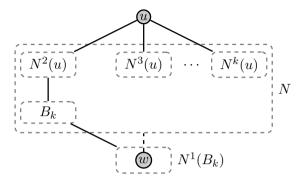


Fig. 2. Structure around u when color k does not satisfy Claim 3.2.(ii).

Now, to prove Theorem 1.3, we apply Lemma 3.1 and the next lemma. A *star* is a tree that has at most one vertex with degree bigger than 1, and the *diameter* of a graph G is the maximum number of edges in a shortest path of G. Here, as happens in Lemma 2.1, we get that the existence of a k-iris in G implies  $b(G) \ge k$ .

**Lemma 3.4** If a graph G has girth at least 7 and a k-iris where  $k \geq 2\chi(G)$ , then G has a b-coloring with k colors.

**Proof.** Let  $u \in V(G)$  be a k-iris with  $k \geq 2\chi(G)$ . Let  $u_2, \ldots, u_k$  be neighbors of u such that  $d(u_i) \geq k-1$  for every  $i \in \{2, \ldots, k\}$ ; let  $N_i$  be a subset of k-2 neighbors of  $u_i$  different from u. Start by coloring u with 1 and, for each  $i \in \{2, \ldots, k\}$ , give color i to  $u_i$  and colors  $\{2, \ldots, k\} \setminus \{i\}$  to  $N_i$ . Denote by T the set  $\{u, u_2, \ldots, u_k\} \cup \bigcup_{i=2}^k N_i$ , i.e., T denotes the set of colored vertices. Observe Figure 3.

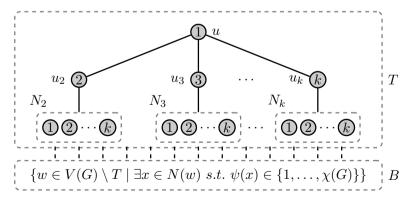


Fig. 3. Subset of vertices around u. The label inside a vertex denotes its color.

Observe that the coloring can be easily done since G[T] - u is a forest formed by k-1 stars. Also, note that we already have k b-vertices of distinct colors, and thus it only remains to extend the partial coloring to the rest of the graph. For this, let B be the set of vertices adjacent to some color class in  $\{1, \ldots, \chi(G)\}$ . We claim that

 $|N(w) \cap T| \leq 1$  for every  $w \in B$ ; indeed, because T induces a tree of diameter 4, if this was not true, then we would get a cycle of length at most 6. By the definition of B, we then get that every  $w \in B$  has no neighbors in color classes  $\chi(G) + 1, \ldots, k$ . Thus, since  $k \geq 2\chi(G)$ , we can color G[B] with colors  $\chi(G) + 1, \ldots, 2\chi(G)$ . Finally, by the definition of B, we know that every  $w \in V(G) \setminus (B \cup T)$  has no neighbors of color 1 through  $\chi(G)$ , which means that we can color G - T - B with these colors.  $\Box$ 

## 4 Conclusion

We have proved that every graph with girth at least 8 is b-continuous, and that graphs with girth at least 7 are in way almost b-continuous. This improves the result presented in [11], where they prove that graphs with girth at least 10 are b-continuous. There, the authors also pose the following questions:

**Question 1** What is the minimum  $\hat{g}$  such that G is b-continuous whenever G is a graph with girth at least  $\hat{g}$ ?

Question 2 Are bipartite graphs with girth at least 6 b-continuous?

Recall that the graph obtained from the complete bipartite graph  $K_{n,n}$  by removing a perfect matching is not b-continuous, for every  $n \geq 4$  [9]. Hence, by our result we get:

$$5 \le \hat{g} \le 8$$
.

We believe that the same techniques might improve this bound to 7, but not further. In particular, we mention that Lemma 3.1 works for graphs with girth 7 and that the bound is 8 because of Lemma 2.1. Therefore, if the following question is answered "yes", then we get  $\hat{g} \leq 7$ .

**Question 3** Let G be a graph with girth at least 7 such that G has a k-iris, with  $k \ge \chi(G) + 1$ . Does G admit a b-coloring with k colors?

As for the case of bipartite graphs, we think it is worth mentioning a known conjecture about their b-chromatic number. Recall the upper bound m(G) for the b-chromatic number b(G), which is the maximum value k for which there exist k vertices with degree at least k-1. The set of all vertices with degree at least m(G)-1 is denoted by D(G), and a graph is said to be tight if |D(G)|=m(G); this means that there is only one candidate set for the b-vertices of a b-coloring of G with m(G) colors. Deciding if b(G)=m(G) is NP-complete even for bipartite tight graphs [9]. In [7], the authors define the class  $\mathcal{B}_m$  that contains every bipartite graph  $G=(A\cup B,E)$  such that m(G)=m, D(G)=A and G has girth at least 6. They conjecture the following:

Conjecture 4.1 [7] For every  $m \geq 3$ , and every  $G \in \mathcal{B}_m$ , we have that:

$$b(G) \ge m(G) - 1.$$

We mention that, if G is a bipartite graph with girth at least 6 and a b-coloring of G with k colors is given,  $k \geq \chi(G)+1$ , then, with a little more work, one can get from the proof of Lemma 3.1 that G contains an induced subgraph H that has a structure

similar to the structure of a graph in  $\mathcal{B}_k$ . Trying to use this structure to obtain a b-coloring of H with k-1 colors could translate into proving Conjecture 4.1. And on the other way around, we believe that a strategy to prove Conjecture 4.1 could help coloring these graphs, which would imply that the answer to Question 2 is "yes". This means that answering Question 2 seems as hard as proving Conjecture 4.1. We also mention that in [10], it is proved that Conjecture 4.1 is a consequence of the famous Erdős-Faber-Lovász Conjecture, which remains open since 1972 and which is largely believed to hold. This is strong evidence that Conjecture 4.1 holds.

Finally, because of the difficulties in obtaining b-continuity already for bipartite graphs with girth at least 6, maybe a good bet would be also to see if the lower bound for  $\hat{g}$  is tight. So, we propose one additional question:

Question 4 Does there exist a graph with girth 5 that is not b-continuous?

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