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Specifying JAVA CARD API in OCL

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Abstract

We discuss the development of an OCL specification for the JAVA CARD API. The main purpose of this specification is to support and aid the verification of JAVA CARD programs in the KeY system. The main goal of the KeY system is to integrate object oriented design and formal methods. The already existing specification written in JML (JAVA Modelling Language) has been used as a starting point for the development of the OCL specification. In this paper we report on the problems that we encountered when writing the specification and their solutions, we present the most interesting parts of the specification, we report on successful verification attempts and finally we evaluate OCL and compare it to JML in the context of JAVA CARD program specification and verification.

*Keywords:* OCL, JML, JAVA CARD, Formal Specification, Formal Verification, Object-Oriented Design

# Introduction

This paper reports on the development of an OCL specification for the JAVA CARD API [[19](#_bookmark26)]. JAVA CARD [[9](#_bookmark16)] is a subset of the JAVA programming language and is used to program smart cards. The JAVA CARD API (Application Pro- gramming Interface) is a set of library classes used in JAVA CARD programs. JAVA CARD API is a much smaller version of the standard JAVA API and is specifically designed for smart card programming. The OCL specification is necessary to perform formal verification of such programs when the implemen- tation of the API classes is not available. Even if the API implementation is

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available, having the OCL specification helps to avoid repetitive work of prov- ing the API implementation each time API method is used in a JAVA CARD program. The secondary purpose of writing the specification is to document the behaviour of the JAVA CARD API in a formal way. We discuss the problems we encountered when writing the specification in OCL and their solutions. We present some of the most interesting parts of the specification and report on successful verification attempts of the reference implementation of JAVA CARD API w.r.t. our specification. Finally, we evaluate OCL and compare it to JML in the context of this work. This paper summarises results from [[11](#_bookmark18)].

In the following section we give more details about the background and motivation of this work. In Section [3](#_bookmark3) we give a detailed report on the devel- opment of the specification, in Section [4](#_bookmark5) we present some interesting parts of our specification, in Section [5](#_bookmark6) we evaluate OCL in the context of the presented work and finally we conclude in Section [6](#_bookmark7).

# Background

* 1. *The KeY Project*

The work presented in this paper has been carried out as part of the KeY project [[1](#_bookmark8),[2](#_bookmark9),[10](#_bookmark17)]. The main goal of the KeY project is to enhance a commer- cial CASE (Computer Aided Software Engineering) tool with functionality for formal specification and deductive verification and, thus, to integrate formal methods into real-world software development processes. Accordingly, the de- sign principles for the software verification component of the KeY system are:

* + 1. The specification language should be usable by people who do not have years of training in formal methods. The Object Constraint Language (OCL) [[14](#_bookmark21)], which is incorporated into the current version of the Unified Modelling Language (UML), is the specification language of our choice.
    2. The programs that are verified should be written in a “real” object- oriented programming language. We decided to use JAVA CARD. This choice is motivated by the following reasons. First of all, many JAVA CARD applications are subject to formal verification, because they are usually security critical (e.g. authentication) and difficult to update in case a fault is discovered. At the same time the JAVA CARD language is easier to handle than full JAVA (for example, there is no concurrency and no GUI—see Section [2.2](#_bookmark2)). Also, JAVA CARD programs are smaller than normal JAVA programs and thus easier to verify.

The architecture of the KeY system is shown in Figure [1](#_bookmark1). It is built on top of a commercial CASE tool (Borland Together Control Center [[7](#_bookmark11)]) and extends



CASE Tool

extension

for formal specification

UML

OCL

Java

Verification Middleware

Dynamic Logic Deduction Component

automated

interactive

Fig. 1. The architecture of the KeY system

it with facilities for formal specification and verification of JAVA programs in the following ways:

* It supports creation and manipulation of OCL constraints, e.g. the KeY system can automatically create a partial OCL specification by instantiat- ing an OCL template (commonly used OCL specification schema) or use a syntax based editor to create OCL expressions.
* The deduction component is used to actually construct proofs for JAVA Dy- namic Logic proof obligations generated from the UML model, OCL con- straints and JAVA implementation. The deduction component is an interac- tive verification system based on JAVA Dynamic Logic, a logic specifically designed for formal verification of JAVA programs [[3](#_bookmark10)].
  1. *JAVA CARD and JAVA CARD API*

JAVA CARD is a technology that provides means to program smart cards with (a subset of) the JAVA programming language. Due to limited resources of smart cards, the JAVA CARD language is limited in a number of ways as compared to full JAVA. The following is the list of features that are not supported in JAVA CARD: large primitive data types (int, long, double, float), characters and strings, multidimensional arrays, dynamic class loading, threads and garbage collection. Most of the remaining JAVA features, in particular object-oriented ones like interfaces, inheritance, virtual methods, overloading, dynamic object

creation are supported by the JAVA CARD language.

The JAVA CARD API is a library that handles smart card specific features, like Application Protocol Data Units (APDUs—used for communication be- tween the card and the rest of the world), Application IDentifiers (AIDs), JAVA CARD specific system routines, PIN codes, etc. [[19](#_bookmark26)]. Some of the packages in- cluded in the JAVA CARD API 2.2 are the following:

* java.lang—provides classes that are fundamental to the design of the JAVA CARD technology subset of the JAVA programming language. The classes in this package are derived from java.lang in the standard JAVA programming language and represent the core functionality required by the JAVA CARD Virtual Machine.
* javacard.framework—provides a framework of classes and interfaces for building, communicating and working with JAVA CARD applets. These classes and interfaces provide the minimum required functionality for a JAVA CARD environment. The key classes and interfaces in this package are:
  + AID—encapsulates the Application IDentifier (AID) associated with an applet.
  + APDU—provides methods for controlling card input and output.
  + Applet—the base class for all JAVA CARD applets on the card. It provides methods for working with applets to be loaded onto, installed into and executed on a JAVA CARD compliant smart card.
  + JCSystem—provides methods for controlling system functions such as transaction management, transient objects, object deletion mechanism, resource management, and inter-applet object sharing.
  + Util—provides convenience methods for working with arrays and array data.

The whole JAVA CARD API consists of 57 classes and interfaces, many of which are very simple (e.g. exception classes).

* 1. *Use Cases for OCL Speciﬁcation of the JAVA CARD API*

One of the purposes of the KeY system is the possibility to formally verify JAVA CARD applications. To successfully verify a program that uses the JAVA CARD API one has to have access to either the implementation of the API or its formal specification. Since the implementation of the API is usually not available (especially when the methods are native), the latter is the solution we are aiming for. Let us look at an example to illustrate how the JAVA CARD API specification is used in the verification process. Suppose we have implemented a method aMethod in our JAVA CARD program. We now want to verify that the implementation satisfies the formal specification (the pair of

pre- and postconditions) of method aMethod:

/\*\*

* @preconditions <pre>
* @postconditions <post>

\*/

public void aMethod(...) {

...

APIClass.apiMethod(...);

...

}

Our method invokes a method from the JAVA CARD API, which we assume has been already specified. The specification of aMethod and its implementation is translated into a proof obligation, which in turn is passed to the KeY de- duction component (prover). When trying to construct a proof for this proof obligation, we sooner or later have to apply a rule that takes care of the in- vocation of the API method apiMethod. If we had no specification of this method we would have to replace the method call with the actual method body. In case the specification of apiMethod is available it is enough to verify that the precondition of apiMethod is satisfied in the state before apiMethod is executed and then we can simply “replace” the method call to apiMethod with its postcondition. This however is not as straightforward as it sounds, there is ongoing work in the KeY project which investigates when and under what conditions such a replacement can be safely done [[6](#_bookmark12)].

In addition to this, having an OCL specification of the API saves a lot of work during verification of JAVA CARD programs in the long run. When there is no specification available, the same API method call has to be replaced by the method’s implementation and proved each time the method in question is used. In practice it can happen that the same piece of API implementation is going to be placed in the proof more than one time in one program.

The secondary purpose of writing the OCL specification for JAVA CARD API is for documentation purposes—an OCL specification can serve as formal documentation of the JAVA CARD API. This is very useful, because the informal specification does not always contain all the necessary information about the behaviour of the API.

* 1. *Related Work*

As already mentioned the starting point for this work was the formal spec- ification of the JAVA CARD API written in JAVA Modelling Language (JML) [[12](#_bookmark19),[13](#_bookmark20),[16](#_bookmark23)]. That work has been done for similar reasons as stated above, the

main difference is the specification language used. The LOOP tool presented in [[20](#_bookmark27)] uses JML and PVS as the means to formally verify JAVA CARD pro- grams, thus the necessity for the API specification written in JML. As we use the industry standard OCL as a specification language in the KeY project we need to have the JAVA CARD API specification formulated in OCL. We also made an effort to have more complete coverage of the JAVA CARD API in our specification.

# The Development of OCL Specification

As stated above, we based our specification on the JML specification of the JAVA CARD API. We then extended it based on the informal specification (API documentation) and we tried to make use of OCL’s expressiveness wherever possible. Later on we tested parts of our specification by formally verifying (using the KeY system) part of the reference implementation of the JAVA CARD API w.r.t. our specification.

We start by giving an overall description of JML and the JML specification of the JAVA CARD API. Based on that we will describe the main problems to be tackled when writing OCL specification for the API.

* 1. *JML vs. OCL*

As in OCL, the specifications in JML are expressed as class invariants and method pre-/postconditions. Class invariants are assertions that should hold for all instances of the class at any time. Pre- and postconditions are contracts between the provider and the user of the method. The user has to fulfil the precondition when he or she calls the method. The provider guarantees that if the precondition holds at the beginning of the method call, then the corresponding postcondition will hold after the method call. In addition, JML allows one to express when a method throws an exception and which attributes of the class can be modified by the method. All the JML specifications are only valid in the context of their JAVA source code and are presented in the form of JAVA comments. Below is the general syntax of JML used to express the method’s behaviour:

/\*\*

@public behavior

@requires <precondition>; @assignable <list of attributes>; @ensures <postcondition>;

@signals (Exception\_1 e1) <ex1postcondition>;

@signals (Exception\_2 e2) <ex2postcondition>;

\*/

public void aMethod() throws Exception { ... }

The @requires clause defines the method’s precondition, the @assignable clause tells which attributes the method can modify. The meaning of the rest of the specification is the following: if the precondition is satisfied then either the method terminates normally (i.e. does not throw any exception) and the postcondition (@ensures) holds or one of the listed exceptions is thrown and then the corresponding postcondition holds.

JML also allows to use a simpler syntax in case the method is not supposed to throw any exceptions, as the example below shows. The example gives a general impression of what the JAVA CARD API specification in JML looks like. The following is a part of the OwnerPIN class:

public class OwnerPIN implements PIN { private byte[] pin;

private byte maxTries; private byte triesRemaining;

public boolean check(byte[] thePin, short offset, byte length) throws ArrayIndexOutOfBoundsException, NullPointerException {

...

}

...

}

The pin array contains the PIN number, maxTries is the maximal number of attempts allowed to present the correct PIN before the card is locked, and triesRemaining the number of attempts left to present the correct PIN. A JML invariant for this class is the following:

/\*\*

@invariant triesRemaining >= 0 && triesRemaining <= maxTries;

\*/

A JML specification of the method check is given below. The arrayCompare method compares length elements of array this.pin starting at element indexed by 0 with length elements of array thePin starting at element indexed by offset:

/\*\*

@public normal\_behavior

@requires triesRemaining > 0 &&

@ Util.arrayCompare(this.pin, (short)0,

@ thePin, offset, length) == 0; @ensures result == true && triesRemaining == maxTries;

\*/

At this point we are ready to define the main differences between JML

and OCL that caused us some problems when writing the JAVA CARD API specification in OCL. The KeY system provides extensions to OCL to overcome most of those problems.

* 1. *Exceptions*

The current version of OCL in its standard form does not provide a straight- forward way to specify that an exception is thrown by a method. A possible solution is to have an association link thrownExceptions in our class, which represents the set of exceptions thrown by methods of that class. Then it is possible to specify that a method aMethod of class MyClass throws an excep- tion of type MyException this way:

context MyClass::aMethod(): pre: true

post: self.thrownExceptions->exists(e : Exception | e.oclIsKindOf(MyException) and e.oclIsNew())

The KeY system has a unified solution for this—one can use an excThrown(My- Exception) clause in the postcondition, which has a very similar meaning. Later on, when the OCL specification is transformed to a JAVA Dynamic Logic proof obligation for the prover, the excThrown clauses are properly translated to corresponding JAVA Dynamic Logic formulas.

Having that, we can now give the general representation of JML’s @be- havior clause in OCL:

context MyClass::aMethod() pre: <precondition>

post: (not excThrown(java::lang::Exception) and <postcondition>)

or (excThrown(Exception\_1) and <ex1postcondition>) or ...

or (excThrown(Exception\_n) and <exnpostcondition>)

* 1. *The null value*

Another thing that is commonly used in JAVA, but which is not supported in the current version of OCL is the null value. This can be handled in OCL in

two ways:

* When one wants to compare a class attribute to a null value, then it is possible to treat the attribute as an association end, which in OCL can be treated as a set. In that case one can simple say attr->isEmpty() to express the fact that attr has a null value.
* When comparing objects other than class attributes (e.g. method argu- ments) to the null value things are a bit more difficult. If such an object is an array or a collection type, one can use the same technique as described above. Otherwise there is no way to specify that an object should (or should not) have the null value.

Fortunately, the KeY system provides a workaround for this problem as well. One can use the null value directly as if it were defined in OCL, and then during the translation to JAVA Dynamic Logic the null values are handled appropriately.

* 1. *Integer Arithmetics*

The main data types that JAVA CARD programs deal with are JAVA shorts, bytes and arrays. Arrays don’t cause much of a problem, in OCL they can be represented as the Sequence type. The JAVA arithmetic types short and byte however don’t have a corresponding type in OCL. The only integer type in OCL is Integer. The most important aspect of JAVA shorts and bytes is that they can overflow (i.e. they are finite types), while the OCL Integer is an infinite type and never overflows. Since the overflow behaviour is a very important aspect of JAVA programs, we have to be able to distinguish between different integer types in OCL. For this purpose we used dummy “wrapper” classes JByte and JShort to represent corresponding JAVA types. They can be used like this:

context PIN::check(pin: Sequence(JByte), offset: JShort,

length: JByte): Boolean

...

This still does not solve the problem of proper interpretation of overflow be- haviour in OCL. Luckily, the KeY system comes to the rescue again. When the OCL specification is translated to a JAVA Dynamic Logic formula, the user can choose how the integer types are interpreted by the prover: either as finite JAVA types short and byte, or as infinite arithmetic types arithShort and arithByte. In both cases the issue of overflow is treated appropriately. More about handling arithmetics in the KeY system can be found in [[5](#_bookmark13),[17](#_bookmark24)]. Also,

[[8](#_bookmark14)] gives insights into problems associated with integer arithmetics in JML.

* 1. *JML assignable clause*

As mentioned before, JML offers a possibility to express (with the @assign- able clause) that a given method is allowed to change a limited set of at- tributes during its execution. OCL does not offer any mechanism or language construct to specify this in a nice way. One can of course state in the post- condition that the value of a given attribute is not changed by the method by saying:

post: self.attr = self.attr@pre

This is not a good solution, though. Suppose we have a class with 20 attributes and we want to express the fact that only one attribute is assignable. That means we have to write 19 expressions like the one above for all the remaining attributes. There is ongoing work that aims at solving this problem in the KeY system [[6](#_bookmark12)]. The work is about how to properly specify attribute modification behaviour and how such specification can be used in proofs. In the current version of our work we left out the parts of the specification corresponding to the @assignable clause in JML.

# The Specification

The present work resulted in an OCL specification for all classes and interfaces of the JAVA CARD API 2.2. This specification expresses, with a few exceptions (some of the signals clauses and the assignable clauses were not possible to be fully expressed in OCL), as much as the JML specification for JAVA CARD API 2.1.1. In some cases the OCL specification expresses more than the JML specification. In the following we illustrate by example how our OCL specification was created and how it was improved (compared to JML).

First, let us look at the PIN interface (which OwnerPIN implements). The informal specification of method check in the PIN interface is the following:

*public boolean check(byte[] pin, short offset, byte length)*

*Compares pin against the PIN value. If they match and the PIN is not blocked, it sets the validated flag and resets the try counter to its maximum. If it does not match, it decrements the try counter and, if the counter has reached zero, blocks the PIN. Even if a transaction is in progress, the internal state such as the try counter, the validated flag and the blocking state must not be conditionally updated.*

*Notes:*

* *If NullPointerException or ArrayIndexOutOfBoundsException is thrown, the validated flag must be set to false, the try counter must be decremented, and the PIN blocked if the counter reaches zero.*
* *If offset or length parameter is negative an ArrayIndexOutOfBoundsException is thrown.*
* *If offset+length is greater than pin.length, the length of the pin array, an ArrayIndexOut- OfBoundsException is thrown.*
* *If pin parameter is null a NullPointerException is thrown.*

*Parameters:*

*pin the byte array containing the PIN value being checked*

*offset the starting offset in the pin array*

*length the length of pin*

*Returns:*

*true if the PIN value matches; false otherwise*

*Throws:*

*ArrayIndexOutOfBoundsException if the check operation would cause access of data outside array bounds.*

*NullPointerException if pin is null.*

The JML specification for this method found in [[15](#_bookmark22)] is the following (the \old

construct corresponds to OCL’s @pre):

/\*\*

@ public normal\_behavior

@ requires triesRemaining == 0; @ assignable \nothing;

@ ensures result == false; @ also

@ public normal\_behavior

@ requires triesRemaining > 0 && pin != null && offset >= 0 @ && length>=0 && offset+length == pin.length &&

@ Util.arrayCompare(this.pin, (short)0, pin, @ offset, length) == 0;

@ assignable isValidated, triesRemaining;

@ ensures result == true && isValidated && @ triesRemaining == maxTries;

@ also

@ public behavior

@ requires triesRemaining > 0 && !(pin != null && @ offset >= 0 && length >= 0 &&

@ offset+length == pin.length &&

@ Util.arrayCompare(this.pin, (short)0, pin, @ offset, length) == 0);

@ assignable isValidated, triesRemaining; @ ensures result == false &&

@ !isValidated && triesRemaining == @ \old(triesRemaining) - 1;

@ signals (NullPointerException) @ !isValidated &&

@ triesRemaining == \old(triesRemaining) - 1; @ signals (ArrayIndexOutOfBoundsException)

@ !isValidated &&

@ triesRemaining == \old(triesRemaining) - 1; @

\*/

public boolean check(byte[] pin, short offset, byte length) throws ArrayIndexOutOfBoundsException, NullPointerException;

It seems that the JML specification mainly agrees with the informal specifi- cation. One subject that is not touched upon in the JML specification is the following sentence from the informal specification: *Even if a transaction is in progress, the internal state such as the try counter, the validated flag and the blocking state must not be conditionally updated.* This is not possible to specify in either JML or OCL, as it has to do with the internal transaction mechanism of the JAVA CARD Runtime Environment. The issue of specifying and verifying the programs involving JAVA CARD’s transaction mechanism has been investigated thoroughly in the KeY project [[4](#_bookmark15)]. For now, however, we decided to leave this issue out in our OCL specification. Another thing to notice is that the informal specification and the JML specification disagree on the subject of whether offset+length must be equal to pin.length or if offset+length might be less than or equal to pin.length. It seems that a mistake has been made in the JML specification, since it clearly disagrees with the informal specification and since there seems to be no good reason to demand that there must be no free elements in the pin array following the actual PIN value. Therefore our resulting OCL specification agrees with the informal specification in this case:

context PIN::check(pin: Sequence(JByte), offset: JShort,

length: JByte): Boolean

pre: true

post: if self.triesRemaining = 0 then result = false endif and if(self.triesRemaining > 0 and pin <> null

and offset >= 0 and length >= 0 and offset+length <= pin->size()

and self.pin->subSequence(1, length) =

pin->subSequence(offset+1, offset+length))

then (

result = true and self.isValidated

and self.triesRemaining = self.maxTries) endif

and if(self.triesRemaining > 0 and

not(pin <> null and offset >= 0 and length >= 0 and offset+length <= pin->size() and self.pin->subSequence(1, length) =

pin->subSequence(offset+1, offset+length)))

then (

not self.isValidated and

self.triesRemaining = self.triesRemaining@pre-1 and ( (not excThrown(java::lang::Exception) and

result = false)

or excThrown(NullPointerException)

or excThrown(ArrayIndexOutOfBoundsException)))

endif

In the next example we show how the specification of method setKey in class DESKey has been enriched compared to JML specification. The method setKey copies the data (an array of bytes) that is passed as an argument and constitutes the actual key to the internal attribute data. Under certain circumstances, this data is not passed to the method in plain text but as a cipher and the method must then decrypt the data before it is copied into the internal representation. Here is the JML specification for this method:

/\*\*

@public behavior

@ requires keyData != null && kOff >= 0 && @ kOff < keyData.length;

@ assignable CryptoException.systemInstance.reason; @ ensures isInitialized();

@ signals (CryptoException e)

@ e.getReason() == CryptoException.ILLEGAL\_VALUE;

\*/

void setKey(byte[] keyData, short kOff) throws CryptoException;

This specification does not give much information about what this method actually accomplishes. In the OCL specification though, we try to give an idea about this:

context DESKey::setKey(keyData: Sequence(JByte), kOff: JShort) pre: not (keyData = null) and kOff >= 0 and

kOff < keyData->size()

post: (not excThrown(java::lang::Exception) and self.isInitialized() and (

not self.oclIsKindOf(javacardx::crypto::KeyEncryption) or self.getKeyCipher() = null implies

)

) or (

self.data->subSequence(1, self.getSize()/8) = keyData->subSequence(kOff+1, kOff+self.getSize()/8)

excThrown(CryptoException) and CryptoException.systemInstance.reason

= CryptoException.ILLEGAL\_VALUE and (

not self.oclIsKindOf(javacardx::crypto::KeyEncryption) or self.getKeyCipher() = null implies

kOff+self.getSize()/8 > keyData->size()

)

)

What we added in this specification is the following. If this particular instance of DESKey is not an instance of javacardx.crypto.KeyEncryption or if this instance is not associated with a Cipher object (the circumstances under which the input keyData have to be decrypted), then the input data is to be copied directly into the internal attribute data.

While studying the JML specification we found a small number of mi- nor inconsistencies. In the class OwnerPIN for example, the invariant states that the internal class attribute pin should not be null at any point, which requires the constructor of that class to set pin (which is initially null) to a non null value. In that case the constructor should be able to modify the pin attribute, but a corresponding @assignable clause is missing in the specification of the OwnerPIN constructor. The informal specification of that constructor also says that two exceptions can be thrown—PINException and SystemException. The condition for throwing the PINException is clearly defined, but this information is not included in the constructor’s specification. We tried to fix all those small deficiencies in our OCL specification and express as much as possible, but, as we mentioned before, giving the full

specification of the JAVA CARD API in OCL is not possible at the moment.

* 1. *Formal Veriﬁcation*

To give our specification a test we looked into the source of the implemen- tation of the JAVA CARD API distributed with SUN’s JAVA CARD Development Kit version 2.1.1 [[18](#_bookmark25)]. We tried to verify this implementation w.r.t. the spec- ification we have written. Due to current limitations of the KeY system this was not done to the extent one might wish for. One of the technical reasons for this is the fact that the KeY system does not handle arrays in

the version we used. Since the arrays are present almost everywhere in the JAVA CARD API this was a major obstacle. We can however report that a number of simple getReason/setReason methods in the exception classes of javacard.framework package have been verified. A more complicated suc- cessful proof attempt was the verification of the reset method in the OwnerPIN class. The specification is the following:

context OwnerPIN inv:

self.maxPINSize > 0 and self.maxTries > 0 and self.triesRemaining >= 0 and self.triesRemaining <= self.maxTries

context OwnerPIN::reset() pre: true

post: not excThrown(java::lang::Exception) and

not self.isValidated and

if self.isValidated@pre then self.triesRemaining = self.maxTries

else

self.triesRemaining = self.triesRemaining@pre endif

A proof obligation generated by the KeY system states the following: the execution of the reset method preserves the invariant and if the precondition is satisfied before reset is executed then the postcondition is satisfied after reset is executed. Explaining what this proof obligation looks like would require introducing the JAVA Dynamic Logic used in the KeY system in more detail. This would go beyond the scope of this paper. One thing we should say though, is that the proof to verify this specification is performed automatically by the KeY prover, reducing the user interaction to absolute minimum.

# Short Evaluation of OCL

There are a few things that we found very useful about OCL. First of all, it is practically an industry standard and is (partially) supported by some CASE tools (e.g. Borland Together Control Center that we use in the KeY project). Second, it seems that the OCL language is richer than JML in some respects, e.g. the whole library of collection type operations makes expressing properties about Sequence (array) type much easier than in JML. Also, for the same reason, we find OCL much easier to read and understand.

When it comes to JAVA specific features, OCL turns out to be not as good as JML. Just to recapitulate the most important findings from Section [3.1](#_bookmark4): there is no standard way in OCL to express the fact that a method throws an exception, there is only one (infinite) integer type in OCL as compared to the whole set of JAVA integer types and there is no JML’s @assignable counterpart in OCL. In this respect JML is a much stronger language than OCL. Of course, this is because JML was designed specifically for JAVA, while OCL was mainly designed for UML.

# Conclusions

In this paper we presented our experience from the development of an OCL specification for the JAVA CARD API 2.2. Despite the mentioned problems with OCL we managed to specify the whole JAVA CARD API to a reasonable extent. The specification is available on-line at:

<http://i12www.ira.uka.de/~key/doc/2003/exjob.html>

The two main purposes of this work were to aid and support formal verification of JAVA CARD programs in the KeY system and to document the JAVA CARD API in a formal way. We tested our specification by formally verifying the reference implementation of the JAVA CARD API with the KeY system, however, due to technical limitations, this was not done to the desirable extent. Still, the proofs we attempted were successful and were performed automatically by the KeY system. In the near future the KeY system will cover the full JAVA CARD standard. Then we plan to continue in this direction and also, based on our specification, perform formal verification of real life JAVA CARD case studies.

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