

Linear Logic

Will Troiani

February 2021

Contents

1	Multiplicatives	1
2	The Sequentialisation Theorem	8
3	Proofs as permutations	17
4	The dynamics of MLL	20
5	Geometry of Interaction Zero, modelling the dynamics of MLL	23

1 Multiplicatives

Definition 1.0.1. There is an infinite set of **unoriented atoms** X, Y, Z, \dots and an **oriented atom** (or **atomic proposition**) is a pair $(X, +)$ or $(X, -)$ where X is an unoriented atom. The set of **pre-formulas** is defined as follows.

- Any atomic proposition is a preformula.
- If A, B are pre-formulas then so are $A \otimes B$, $A \wp B$.
- If A is a pre-formula then so is $\neg A$.

The set of **formulas** is the quotient of the set of pre-formulas by the equivalence relation \sim generated by, for arbitrary formulas A, B and unoriented atom X , the following.

$$\neg(A \otimes B) \sim \neg B \wp \neg A, \quad \neg(A \wp B) \sim \neg B \otimes \neg A, \quad \neg(X, +) \sim (X, -), \quad \neg(X, -) \sim (X, +) \quad (1)$$

Lemma 1.0.2. For all formulas A we have $\neg\neg A = A$.

Proof. We proceed by induction on the number n given by the sum of the occurrences of \otimes and \wp in A .

If $n = 0$ then $A = (X, +)$ or $A = (X, -)$. In either situation the fact that $\neg\neg A = A$ follows from (1).

Now say $n > 0$ and the result holds for all $k < n$. We either have $A = A_1 \otimes A_2$ or $A = A_1 \wp A_2$, each case is proved in a similar way, we show the details for when $A = A_1 \otimes A_2$.

$$\neg\neg A = \neg\neg(A_1 \otimes A_2) = \neg(\neg A_2 \wp \neg A_1) = \neg\neg A_1 \otimes \neg\neg A_2 \quad (2)$$

By the inductive hypothesis we have that $\neg\neg A_1 = A_1$ and $\neg\neg A_2 = A_2$. It thus follows that $\neg\neg A = A$. \square

Definition 1.0.3. A finite sequence of formulas is a **sequent** and we write $\vdash A_1, \dots, A_n$ for the sequent (A_1, \dots, A_n) .

Definition 1.0.4. A **multiplicative, linear logic deduction rule** (or simply **deduction rule**) results from one of the schemata below by a substitution of the following kind: replace A, B by arbitrary formulas, and $\Gamma, \Gamma', \Delta, \Delta'$ by arbitrary (possibly empty) sequences of formulas separated by commas:

- the **identity group**:

- **Axiom**

$$\frac{}{\vdash A, \sim A} \text{ (ax)}$$

- **Cut**:

$$\frac{\vdash \Gamma, A, \Gamma' \quad \vdash \Delta, \sim A, \Delta'}{\vdash \Gamma, \Gamma', \Delta, \Delta'} \text{ (cut)}$$

- the **multiplicative rules**

- **Times**:

$$\frac{\vdash \Gamma, A, \Gamma' \quad \vdash \Delta, B, \Delta'}{\vdash \Gamma, \Gamma', A \otimes B, \Delta, \Delta'} \otimes$$

- **Par**

$$\frac{\vdash \Gamma, A, B, \Gamma'}{\vdash \Gamma, A \wp B, \Gamma'} \wp$$

- the **structural rule**:

- **Exchange**

$$\frac{\vdash \Gamma, A, B, \Gamma'}{\vdash \Gamma, B, A, \Gamma'} \text{ (ex)}$$

Definition 1.0.5. A **proof in MLL** is a finite, rooted, planar, tree where each edge is labelled by a sequent and each node except for the root is labelled by a valid deduction rule. If the edge connected to the root is labelled by the sequent $\vdash \Gamma$ then we call the proof a **proof of** $\vdash \Gamma$ and in such a situation, Γ is the **conclusion** of π .

A proof in MLL is a concept inspired by the proofs that mathematicians write in, for example, algebraic geometry. One of the reasons why it is interesting to consider a formal notion of a proof so that we can analyse the true *essence* of a proof. Many arbitrary choices are made when one writes a proof on paper, for instance, what order two arguments are written down when one does not depend on the other. Hence, a natural idea is to define an equivalence relation on the set of a proofs in MLL which identify proofs which differ only in these insignificant ways. However, there is an even better idea. We begin with the question: “what is good notation”?

Let us begin with an example of *bad* notation. Say we considered proofs in MLL up to equivalence, as suggested by the previous paragraph. For clarity, just in this paragraph and the next we will refer to proofs in MLL as *preproofs*, and equivalence classes of such things as *proofs*. The only reasonable way to denote a proof would be to denote it *indirectly* via a preproof. Hence, this notation *encompasses more than the pure essence of a proof*, as the arbitrary choices belonging to that particular preproof are also included in the notation. The result: a notation which obfuscates the mathematical structure of proofs.

Hence, the desiderata for what constitutes “good notation” at least includes the following condition.:

There should be exactly one way of writing the object in question. (3)

This is what proof structures (Definition 1.0.6) do, and indeed, we will see that this notation leads to many insights (the Sequentialisation Theorem 2.0.17, Geometry of Interaction 5.0.11, etc) which capture some structure of proofs, without reference to the structure of a choice of representing preproof.

Definition 1.0.6. A **proof structure** is a directed multigraph with edges labelled with formulas (Definition 1.0.1) and with nodes labelled with an element of the following set: $\{(\text{ax}), (\text{cut}), \otimes, \wp, \text{c}\}$. The incoming edges of a node are called its **premisses**, the outgoing edges are its **conclusions**. Proof structures are required to adhere to the following conditions.

- Each node labelled (ax) has exactly two conclusions and no premisses, the conclusions are labelled A and $\neg A$ for some A .
- Each node labelled (cut) has exactly two premisses and no conclusion, where the premisses are labelled A and $\neg A$ for some A .
- Each node labelled \otimes has exactly two premisses and one conclusion. These two premisses are ordered. The smallest one is called the *left* premise of the node, the biggest one is called the *right* premise. The left premise is labelled A , the right premise is labelled B and the conclusion is labelled $A \otimes B$, for some A, B .
- Each node labelled \wp has exactly two ordered premisses and one conclusion. The left premise is labelled A , the right premise is labelled B and the conclusion is labelled $A \wp B$, for some A, B .
- Each node labelled c has exactly one premise and no conclusion. Such a premise of a node labelled c is called the **conclusion** of the proof structure.

Let π be a proof structure. An **axiom link** of π is a subgraph consisting of a node labelled (ax) along with its conclusions labelled $A, \neg A$ for some formula A . A **tensor link** of π consists of a node labelled \otimes along with its premisses and conclusion labelled respectively $A, B, A \otimes B$ for some pair of formulas (A, B) . A **par link** consists of a node labelled \wp along with its premisses and conclusion labelled respectively $A, B, A \wp B$ for some pair of formulas (A, B) . A **cut link** consists of a node labelled (cut) along with its premisses labelled $A, \neg A$ for some formula A .

Remark 1.0.7. Our reference for the definition of a proof structure is [7], see there for more details.

Definition 1.0.8. An **occurrence of a formula** A in a proof structure π is an edge e labelled by A .

Loosely speaking, logic is about determining correct arguments. That is, from the space of arguments (either correct or incorrect), logic determines whether an argument A lies in the subspace of *correct arguments* or not. In the current context, we take the set of proof structures to be the space of proofs, both correct and incorrect, and we take the subset of so called *proof nets* to be the subspace of correct arguments. Proof nets are the proof structures which lie in the image of a translation map (Definition 1.0.10 below) between sequent style proofs and proof structures, we now explain this sequent style logical system.

Definition 1.0.9. The set of **intuitionistic formulas** is defined as follows.

- Any atomic proposition (Definition 1.0.1) is an intuitionistic formula.
- If A, B are formulas then so are $A \otimes B, A \multimap B$.

Let \mathcal{P}^n be the set of all length n sequences of labelled intuitionistic variables with $\mathcal{P}^0 := \{\emptyset\}$, and $\mathcal{P} := \bigcup_{n=0}^{\infty} \mathcal{P}^n$. A **sequent** is a pair (Γ, A) where $\Gamma \in \mathcal{P}$ and $A \in I\Psi$, written $\Gamma \vdash A$. We call Γ the **antecedent** and A the **succedent** of the sequent. Given Γ and a labelled intuitionistic formula A we write Γ, A for the element of \mathcal{P} given by appending A to the end of Γ . We write $\vdash A$ for $\emptyset \vdash A$.

An **intuitionistic, multiplicative deduction rule** (or simply **deduction rule**) results from one of the schemata below by a substitution of the following kind: replace A, B, C by arbitrary labelled intuitionistic formulas, and Γ, Δ, Θ by arbitrary (possibly empty) sequences of labelled intuitionistic formulas separated by commas:

- The **identity group**:

$$\text{Axiom} \quad \frac{}{A \vdash A} \text{ (ax)} \quad (4)$$

$$\text{Cut} \quad \frac{\Gamma \vdash A \quad \Delta, A, \Theta \vdash B}{\Gamma, \Delta, \Theta \vdash B} \text{ (cut)} \quad (5)$$

- The logical rules:

$$\begin{array}{c} \text{Left/right} \\ \text{times} \end{array} \quad \frac{\Gamma, A, B, \Gamma' \vdash C}{\Gamma, A \otimes B, \Gamma' \vdash C} \text{ (L } \otimes \text{)} \quad \frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \otimes B} \text{ (R } \otimes \text{)} \quad (6)$$

$$\begin{array}{c} \text{Right/left} \\ \text{implication} \end{array} \quad \frac{\Gamma, A, \Gamma' \vdash B}{\Gamma, \Gamma' \vdash A \multimap B} \text{ (R } \multimap \text{)} \quad \frac{\Gamma \vdash A \quad \Delta, B, \Delta' \vdash C}{A \multimap B, \Gamma, \Delta \vdash C} \text{ (L } \multimap \text{)} \quad (7)$$

- The structural rule

$$\text{Exchange} \quad \frac{\Gamma, A, B, \Gamma' \vdash C}{\Gamma, B, A, \Gamma' \vdash C} \text{ (ex)} \quad (8)$$

A **proof in IMLL** is defined the same way as in Definition 1.0.5.

Definition 1.0.10. Let Σ denote the set of multiplicative, linear logic proofs and MPS the set of multiplicative proof structures. We let $T : \Sigma \rightarrow \text{MPS}$ denote the function defined inductively by associating to each deduction rule of Definition 1.0.4 a multiplicative proof structure. More precisely, we simultaneously inductively prove that if π has height n and is constructed from either one proof π' with height less than n or from two proofs π_1, π_2 each with height less than n , then $T(\pi'), T(\pi_1), T(\pi_2)$ have conclusions corresponding to the conclusions of π', π_1, π_2 , and we use this fact to inductively define $T(\pi)$ which in turn has conclusions corresponding to those of π .

Given a proof π , the following notation:

$$\begin{array}{c} T(\pi) \\ | \\ A \end{array} \quad (9)$$

means the translation $T(\pi)$, which admits a conclusion A , with the conclusion node c removed.

$$\begin{array}{ccc} \text{Axiom} & \frac{}{\vdash A, \sim A} \text{ (ax)} & \xrightarrow{T} \begin{array}{c} \text{(ax)} \\ \curvearrowright \\ A \downarrow c \quad \neg A \downarrow c \end{array} \\ \\ \text{Cut} & \frac{\frac{\pi_1}{\vdash \Gamma, A, \Gamma'} \quad \frac{\pi_2}{\vdash \Delta, \sim A, \Delta'}}{\vdash \Gamma, \Gamma', \Delta, \Delta'} \text{ (cut)} & \xrightarrow{T} \begin{array}{c} T(\pi_1) \quad T(\pi_2) \\ | \quad | \\ A \quad \neg A \\ \curvearrowright \quad \curvearrowleft \\ \text{(cut)} \end{array} \end{array}$$

Times	$\frac{\frac{\pi_1}{\vdots} \quad \frac{\pi_2}{\vdots}}{\vdash \Gamma, A, \Gamma' \quad \vdash \Delta, B, \Delta'} \otimes$	\xrightarrow{T}	
Par	$\frac{\frac{\pi}{\vdots}}{\vdash \Gamma, A, B, \Gamma'} \wp$	\xrightarrow{T}	
Exchange	$\frac{\frac{\pi}{\vdots}}{\vdash \Gamma, A, B, \Gamma'} \text{ (ex)}$	\xrightarrow{T}	$T(\pi)$

A **multiplicative proof net** (or simply **proof net**) is a multiplicative proof structure which lies in the image of T .

Definition 1.0.11. Let Π denote the set of intuitionistic, multiplicative, linear logic proofs. Then again, there is a translation

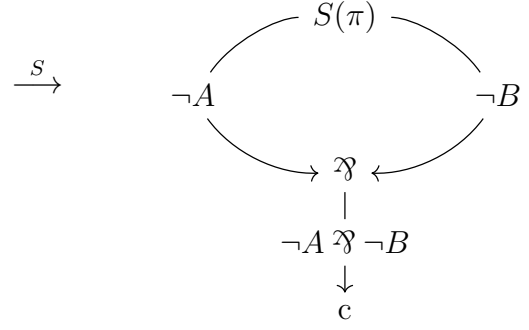
$$S : \Pi \longrightarrow \text{MPS} \quad (10)$$

defined inductively:

Axiom	$\overline{A \vdash A} \text{ (ax)}$	\xrightarrow{S}	
Cut	$\frac{\frac{\pi_1}{\vdots} \quad \frac{\pi_2}{\vdots}}{\Gamma \vdash A \quad \Delta, A, \Theta \vdash B} \text{ (cut)}$	\xrightarrow{S}	

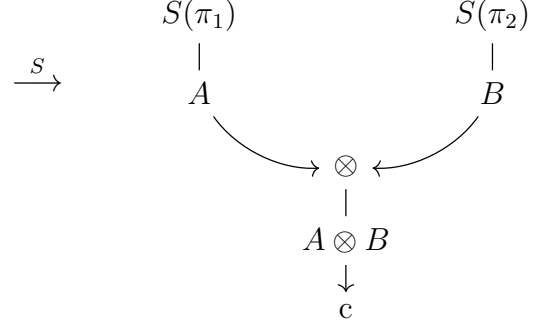
Left times

$$\frac{\frac{\pi}{\vdots}}{\Gamma, A, B, \Gamma' \vdash C} \quad \frac{\Gamma, A, B, \Gamma' \vdash C}{\Gamma, A \otimes B, \Gamma' \vdash C} \text{ (L } \otimes \text{)}$$



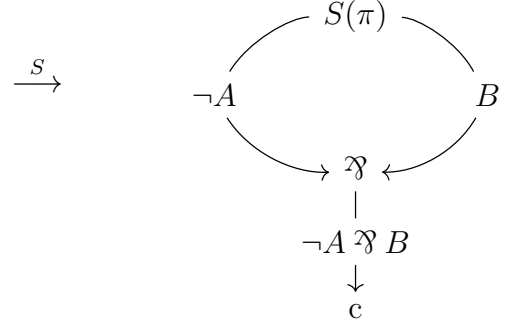
Right times

$$\frac{\frac{\pi_1}{\vdots}}{\Gamma \vdash A} \quad \frac{\frac{\pi_2}{\vdots}}{\Delta \vdash B} \quad \frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \otimes B} \text{ (R } \otimes \text{)}$$



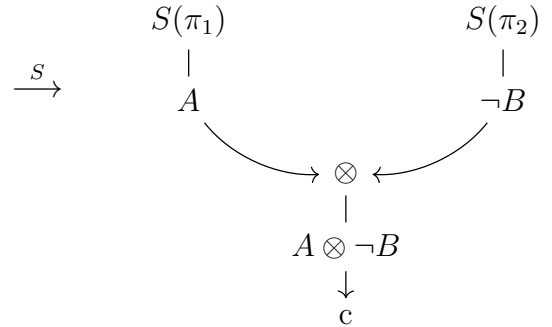
Right implication

$$\frac{\frac{\pi}{\vdots}}{\Gamma, A, \Gamma' \vdash B} \quad \frac{\Gamma, A, \Gamma' \vdash B}{\Gamma, \Gamma' \vdash A \multimap B} \text{ (R } \multimap \text{)}$$



Left implication

$$\frac{\frac{\pi_1}{\vdots}}{\Gamma \vdash A} \quad \frac{\frac{\pi_2}{\vdots}}{\Delta, B, \Delta' \vdash C} \quad \frac{\Gamma \vdash A \quad \Delta, B, \Delta' \vdash C}{A \multimap B, \Gamma, \Delta \vdash C} \text{ (L } \multimap \text{)}$$



Exchange

$$\frac{\frac{\pi}{\vdots}}{\Gamma, A, B, \Gamma' \vdash C} \quad \frac{\Gamma, A, B, \Gamma' \vdash C}{\Gamma, B, A, \Gamma' \vdash C} \text{ (ex)}$$

$$\xrightarrow{S} S(\pi)$$

A **intuitionistic, multiplicative proof net** (or simply **intuitionistic proof net**) is an proof structure which lies in the image of S .

Definition 1.0.12. There is also a translation $R : \Pi \longrightarrow \Sigma$ which we now present. In what follows, if $\Gamma = A_1, \dots, A_n$ is a sequence of formulas, then $\neg\Gamma$ denotes the sequence given by $\neg A_1, \dots, \neg A_n$.

Axiom	$\frac{}{A \vdash A} \text{ (ax)}$	\xrightarrow{R}	$\frac{}{\vdash \neg A, A} \text{ (ax)}$
Cut	$\frac{\frac{\pi_1}{\vdots} \quad \frac{\pi_2}{\vdots}}{\Gamma \vdash A \quad \Delta, A, \Theta \vdash B} \text{ (cut)}$	\xrightarrow{R}	$\frac{\frac{R(\pi_1)}{\vdots} \quad \frac{R(\pi_2)}{\vdots}}{\vdash \neg \Gamma, A \quad \vdash \neg \Delta, \neg A, \neg \Theta, B} \text{ (cut)}$
Left times	$\frac{\frac{\pi}{\vdots}}{\Gamma, A, B, \Gamma' \vdash C} \text{ (L } \otimes)$	\xrightarrow{R}	$\frac{\frac{R(\pi)}{\vdots}}{\vdash \neg \Gamma, \neg A, \neg B, \neg \Gamma', C} \otimes$
Right times	$\frac{\frac{\pi_1}{\vdots} \quad \frac{\pi_2}{\vdots}}{\Gamma \vdash A \quad \Delta \vdash B} \text{ (R } \otimes)$	\xrightarrow{R}	$\frac{\frac{R(\pi_1)}{\vdots} \quad \frac{R(\pi_2)}{\vdots}}{\vdash \neg \Gamma, A \quad \vdash \neg \Delta, B} \otimes$
Right implication	$\frac{\frac{\pi}{\vdots}}{\Gamma, A, \Gamma' \vdash B} \text{ (R } \multimap)$	\xrightarrow{R}	$\frac{\frac{R(\pi)}{\vdots}}{\vdash \neg \Gamma, \neg A, \neg \Gamma', B} \text{ (ex)}$ $\frac{\vdash \neg \Gamma, \neg \Gamma', \neg A, B}{\vdash \neg \Gamma, \neg \Gamma', \neg A \wp B} \wp$
Left implication	$\frac{\frac{\pi_1}{\vdots} \quad \frac{\pi_2}{\vdots}}{\Gamma \vdash A \quad \Delta, B, \Delta' \vdash C} \text{ (L } \multimap)$	\xrightarrow{R}	$\frac{\frac{R(\pi_1)}{\vdots} \quad \frac{R(\pi_2)}{\vdots}}{\vdash \neg \Gamma, A \quad \vdash \neg \Delta, \neg B, \neg \Delta', C} \otimes$
Exchange	$\frac{\frac{\pi}{\vdots}}{\Gamma, A, B, \Gamma' \vdash C} \text{ (ex)}$	\xrightarrow{R}	$\frac{\frac{\pi}{\vdots}}{\vdash \neg \Gamma, \neg A, \neg B, \neg \Gamma', C} \text{ (ex)}$

It is easy to see that the following diagram commutes.

$$\begin{array}{ccc}
 \Pi & \xrightarrow{R} & \Sigma \\
 & \searrow S & \downarrow T \\
 & & \text{MPN}
 \end{array} \tag{11}$$

Lemma 1.0.13. *The map R is injective.*

Proof. By inspection of the defining rules of R . □

Lemma 1.0.14. *The map R is not surjective.*

Proof. There is not proof π in Π such that $R(\pi)$ is the following.

$$\frac{}{\vdash \neg A, A} \text{ (ax)}$$

□

Remark 1.0.15. Neither of the maps T, S are injective. By definition of proof nets, the map T is surjective. The image of S is the set of **multiplicative, intuitionistic proof nets**. These will not be considered again in these notes, but it would be interesting to find a correctness criterion for intuitionistic proof nets similar to the *long trip condition* (Section 2).

In Section 2 we will see what the image of the map T is.

2 The Sequentialisation Theorem

Definition 2.0.1. Let π be a proof structure and denote the set of tensor and par links of π by $\text{Link}_{\otimes, \wp} \pi$ (or simply $\text{Link} \pi$). A **switching** of π is a function

$$S : \text{Link} \pi \longrightarrow \{L, R\} \quad (12)$$

A **switching** of a particular link l is a choice of L, R associated to l .

Definition 2.0.2. Let π be a proof structure. Let $\mathcal{O}(\pi)$ denote the set of occurrences of formulas in π (Definition 1.0.8). We consider two disjoint copies of this set

$$\mathcal{U}(\pi) := \mathcal{O}(\pi) \amalg \mathcal{O}(\pi) \quad (13)$$

where elements from the first copy are the **up elements**, and elements from the second copy are the **down elements**. We write $\uparrow A$ for the up element corresponding to an occurrence of a formula A in π , and similarly for $A \downarrow$. Given a switching S of π , a **pretrip of π with respect to S** is a finite sequence (x_1, \dots, x_n) of elements of $\mathcal{U}(\pi)$ satisfying the following.

1. The sequence is a loop, that is, $x_1 = x_n$, and all elements (except the first and the last) are distinct.
2. If $x_j = \uparrow A$ and A is part of an axiom link then $x_{j+1} = \sim A \downarrow$.
3. If $x_j = A \downarrow$ and A is part of a cut link then $x_{j+1} = \uparrow \sim A$.
4. For any tensor link l with premises A, B such that l has switching L , we have:
 - if $x_j = A \downarrow$ then $x_{j+1} = (A \otimes B) \downarrow$,
 - if $x_j = \uparrow (A \otimes B)$ then $x_{j+1} = \uparrow B_j$,
 - if $x_j = B \downarrow$ then $x_{j+1} = \uparrow A$.

and if l has switching R , we have:

- if $x_j = A \downarrow$ then $x_{j+1} = \uparrow B$,
- if $x_j = \uparrow (A \otimes B)$ then $x_{j+1} = \uparrow A$,
- if $x_j = B \downarrow$ then $x_{j+1} = (A \otimes B) \downarrow$.

(see Figure 1)

5. for any par link l with premises A, B such that l has switching L , we have:
 - if $x_j = \uparrow (A \wp B)$ then $x_{j+1} = \uparrow A$,

- if $x_j = A \downarrow$ then $x_{j+1} = (A \wp B) \downarrow$,
- if $x_j = B \downarrow$ then $x_{j+1} = \uparrow B$.

and if l evaluates under S to R , we have:

- if $x_j = A \downarrow$ then $x_{j+1} = \uparrow A$,
- if $x_j = \uparrow (A \wp B)$ then $x_{j+1} = \uparrow B$,
- if $x_j = B \downarrow$ then $x_{j+1} = (A \wp B) \downarrow$.

(see Figure 2)

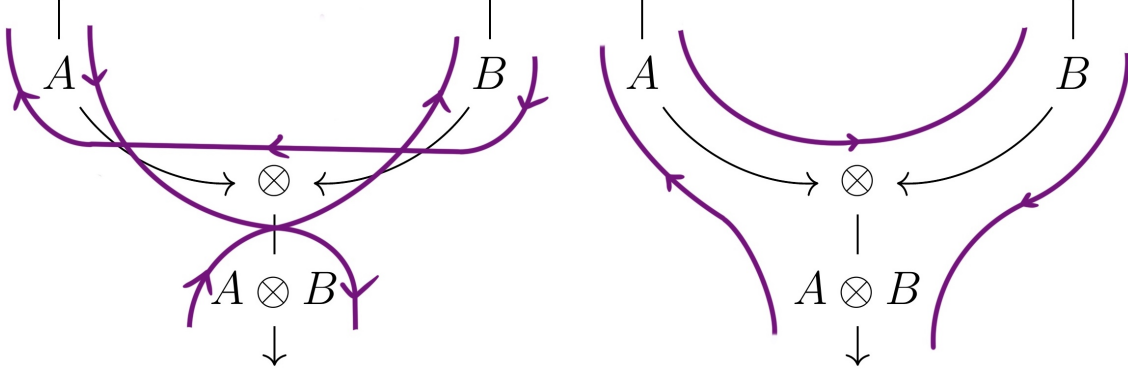


Figure 1: Tensor link, L switching, R switching

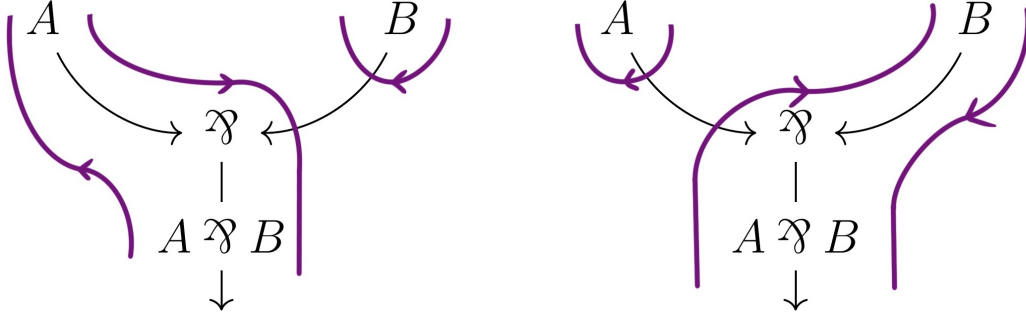


Figure 2: Par link, L switching, R switching.

Definition 2.0.3. Let $\text{Pre}\mathcal{T}(\pi, S)$ denote the set of all pretrips of π with respect to S . We define an equivalence relation on this set \sim where two pretrips (x_1, \dots, x_n) and (y_1, \dots, y_m) are equivalent if $n = m$, and there exists an integer k such that $x_{i+k} = y_i$ (where $i + k$ means mod n) for all $i = 1, \dots, n$.

A **trip** of π with respect to S is an equivalence class of pretrips. We denote the set of all trips by $\mathcal{T}(\pi, S)$. If the set $\mathcal{T}(\pi, S)$ admits more than one element, these elements are called **short trips**, and if it admits only one element, this element is the **long trip**. We refer to the statement “for all switchings S , the set $\mathcal{T}(\pi, S)$ contains exactly one element” as the **long trip condition**.

A **short pretrip** is a choice of representative for a pretrip, and a **long pretrip** is a choice of representative of a long trip.

Given a proof structure π satisfying the long trip condition and a tensor link l with premises A, B say, let S be a switching of π and $t := (x_1, \dots, x_n)$ be the long pretrip of π satisfying $x_1 = A \downarrow$. Since π satisfies the long trip condition, it must be the case that $\uparrow (A \otimes B)$ and $B \downarrow$ occur somewhere in t , can we determine

which occurs earlier? Let $m, l > 0$ be such that $x_m = \uparrow (A \otimes B)$, $x_l = B \downarrow$ and assume $l < m$. Say $S(\tau) = L$, then t has the shape

$$(A \downarrow, (A \otimes B) \downarrow, \dots, B \downarrow, \uparrow A, \dots, \uparrow (A \otimes B), \uparrow B, \dots, A \downarrow) \quad (14)$$

Now consider the switching given by

$$\hat{S}(\sigma) = \begin{cases} S(\sigma), & \sigma \neq \tau \\ R, & \sigma = \tau \end{cases}$$

Then (14) becomes:

$$(A \downarrow, \uparrow B, \dots, A \downarrow) \quad (15)$$

which is a short pretrip, contradicting the assumption that π satisfies the long trip condition. Thus $m < l$. We have proven (the first half) of the following.

Lemma 2.0.4. *Let π be a proof structure satisfying the long trip condition, l be a tensor link with premises A, B say, S be a switching of π and (x_1, \dots, x_n) the long pretrip satisfying $x_1 = A \downarrow$. If $m, l > 0$ are such that $x_m = \uparrow (A \otimes B)$, $x_l = B \downarrow$, then:*

- if $S(\tau) = L$ then $m < l$,
- if $S(\tau) = R$ then $l < m$.

The proof of the other half is similar to what has already been written, however since Lemma 2.0.4 contradicts [2, Lemma 2.9.1] we write out the details here:

Proof. Say $m < l$, then t has the shape

$$(A \downarrow, \uparrow B, \dots, \uparrow (A \otimes B), \uparrow A, \dots, B \downarrow, (A \otimes B) \downarrow, \dots, A \downarrow) \quad (16)$$

Now consider the switching given by

$$S'(\sigma) = \begin{cases} S(\sigma), & \sigma \neq \tau \\ L, & \sigma = \tau \end{cases}$$

Then (16) becomes:

$$(A \downarrow, (A \otimes B) \downarrow, \dots, A \downarrow) \quad (17)$$

which is a short pretrip. □

Lemma 2.0.5. *Let π be a proof structure satisfying the long trip condition, l be a par link with premises A, B say, S be a switching of π and (x_1, \dots, x_n) be the long pretrip satisfying $x_1 = A \downarrow$. If $m, l > 0$ are such that $x_m = \uparrow (A \wp B)$, $x_l = B \downarrow$, then*

- if $S(\tau) = L$ then $m < l$,
- if $S(\tau) = R$ then $l < m$

Remark 2.0.6. Lemma 2.0.4 gives a nice interpretation of Lemma 2.0.4 that long trips *return to where they left* at each tensor link.

The situation is a bit different for par links; the relevant slogan is long trips *visit the premises before returning to the conclusion*.

Say π satisfies the long trip condition and moreover π admits a tensor link l (with premises A, B say) such that if l is removed, the resulting proof structure consists of two disjoint proof structures π_1, π_2 each satisfying the long trip condition. It is necessarily the case that any pretrip ρ of π starting at $\uparrow A$ visits the entirety of $\mathcal{U}(\pi_1)$ before returning to the tensor link l , lest π_1 admit a short trip. Moreover, it must be the case that σ admits no occurrence of formulas in π_2 lest the result of removing the tensor link l not result in disjoint proof structures. Thus, if such a link l exists, it is *maximal* in the sense that there is no other tensor link l' where a pretrip starting at a premise of l' contains the entirety of any pretrip starting at A . The remainder of this Section will amount to proving the converse, that any such maximal tensor link “splits” π .

Definition 2.0.7. Let π be a proof structure satisfying the long trip condition, S a switching of π , and A an occurrence of a formula in π . Consider the long pretrip (x_1, \dots, x_n) satisfying $x_1 = \uparrow A$. We denote by

$$\text{PTrip}(\pi, S, A, \uparrow) \quad (18)$$

the subsequence (x_1, \dots, x_m) of (x_1, \dots, x_n) satisfying $x_m = A \downarrow$. We define

$$\text{PTrip}(\pi, S, A, \downarrow) \quad (19)$$

similarly.

Also, for $a \in \{\uparrow, \downarrow\}$ we define the following set

$$\text{Visit}_S(A, a) := \{C \in \mathcal{O}(\pi) \mid \uparrow C, C \downarrow \text{ occur in } \text{PTrip}(\pi, S, A, a)\} \quad (20)$$

The **up empire of A** is the following set:

$$\text{Emp}_\uparrow A := \{C \in \mathcal{O}(\pi) \mid \text{For all switchings } S \text{ we have } \uparrow C, C \downarrow \text{ occur in } \text{PTrip}(\pi, S, A, \uparrow)\} \quad (21)$$

The **down empire of A** is defined symmetrically.

With this new terminology we now have some corollaries of Lemmas 2.0.4 and 2.0.5:

Corollary 2.0.8. Let π be a proof structure satisfying the long trip condition, and let S be a switching of π , for a formula A and $a \in \{\uparrow, \downarrow\}$, denote $\text{PTrip}(\pi, S, A, a)$ by $\text{PTrip}(A, a)$:

1. if A is part of an axiom link then

$$\text{PTrip}(A, \uparrow) = \uparrow A, \text{PTrip}(\sim A, \downarrow), A \downarrow \quad (22)$$

2. if l is a tensor link with conclusion $A \otimes B$:

(a) if $S(l) = L$:

$$\text{PTrip}(A, \downarrow) = A \downarrow, \text{PTrip}(A \otimes B, \downarrow), \text{PTrip}(B, \uparrow), \uparrow A \quad (23)$$

$$\text{PTrip}(B, \downarrow) = B \downarrow, \text{PTrip}(A, \uparrow), \text{PTrip}(A \otimes B, \downarrow), \uparrow B \quad (24)$$

$$\text{PTrip}(A \otimes B, \uparrow) = \uparrow A \otimes B, \text{PTrip}(B, \uparrow), \text{PTrip}(A, \uparrow), A \otimes B \downarrow \quad (25)$$

(b) if $S(l) = R$:

$$\text{PTrip}(A, \downarrow) = A \downarrow, \text{PTrip}(B, \uparrow), \text{PTrip}(A \otimes B, \downarrow), \uparrow A \quad (26)$$

$$\text{PTrip}(B, \downarrow) = B \downarrow, \text{PTrip}(A \otimes B, \downarrow), \text{PTrip}(A, \uparrow), \uparrow B \quad (27)$$

$$\text{PTrip}(A \otimes B, \uparrow) = \uparrow A \otimes B, \text{PTrip}(A, \uparrow), \text{PTrip}(B, \uparrow), A \otimes B \downarrow \quad (28)$$

3. if A is a premise of a par link l with conclusion $A \wp B$:

(a) if $S(l) = L$:

$$\text{PTrip}(A, \downarrow) = A \downarrow, \text{PTrip}(A \wp B, \downarrow), \uparrow A \quad (29)$$

$$\text{PTrip}(B, \downarrow) = B \downarrow, \uparrow B \quad (30)$$

$$\text{PTrip}(A \wp B, \uparrow) = \uparrow A \wp B, \text{PTrip}(A, \uparrow), A \wp B \downarrow \quad (31)$$

(b) if $S(l) = R$:

$$\text{PTrip}(A, \downarrow) = A \downarrow, \uparrow A \quad (32)$$

$$\text{PTrip}(B, \downarrow) = B \downarrow, \text{PTrip}(A \wp B, \downarrow), \uparrow B \quad (33)$$

$$\text{PTrip}(A \wp B, \uparrow) = \uparrow A \wp B, \text{PTrip}(B, \uparrow), A \wp B \downarrow \quad (34)$$

In particular:

Corollary 2.0.9. *For any formula A which is a premise to either a tensor or par link, and any $a \in \{\uparrow, \downarrow\}$, we have:*

$$\uparrow C \text{ occurs in } \text{PTrip}(\pi, S, A, \uparrow) \quad \text{if and only if} \quad C \downarrow \text{ occurs in } \text{PTrip}(\pi, S, A, \downarrow)$$

and similarly for $\text{PTrip}(\pi, S, A, \downarrow)$.

Proof. By induction on the length of the sequence $\text{PTrip}(\pi, S, A, a)$ and appealing to Corollary 2.0.8. \square

Corollary 2.0.10. *Let π be a proof structure satisfying the long trip condition, we have the following.*

1. *For any axiom link with conclusions $A, \neg A$:*

$$\text{Emp}_{\uparrow} A = \text{Emp}_{\downarrow}(\neg A) \cup \{A\} \quad (35)$$

2. *For any cut link with premises $A, \neg A$:*

$$\text{Emp}_{\downarrow} A = \text{Emp}_{\uparrow}(\neg A) \cup \{A\} \quad (36)$$

3. *For any tensor link with premises A, B :*

$$\text{Emp}_{\uparrow} A \cap \text{Emp}_{\uparrow} B = \emptyset \quad (37)$$

4. *For any tensor or par link with premises A, B and conclusion C :*

$$\text{Emp}_{\uparrow} C = \text{Emp}_{\uparrow} A \cup \text{Emp}_{\uparrow} B \cup \{C\} \quad (38)$$

5. *For any tensor link with premises A, B :*

$$\text{Emp}_{\downarrow} B = \text{Emp}_{\uparrow} A \cup \text{Emp}_{\downarrow}(A \otimes B) \cup \{B\} \quad (39)$$

and similarly:

$$\text{Emp}_{\downarrow} A = \text{Emp}_{\uparrow} B \cup \text{Emp}_{\downarrow}(A \otimes B) \cup \{A\} \quad (40)$$

Definition 2.0.11. Given any link l we write $B \in l$ if B occurs as either a premise or a conclusion of l .

Let π be a proof structure satisfying the long trip condition, and $a \in \{\uparrow, \downarrow\}$. The set of **links of A with respect to S** is the set

$$\text{Link}_a A := \{l \in \text{Link } \pi \mid \forall B \in l, B \in \text{Emp}_a A\} \quad (41)$$

Definition 2.0.12. Let π be a proof structure satisfying the long trip condition and let $a \in \{\uparrow, \downarrow\}$. Define the set

$$\text{Link}_{\mathcal{R},a}^0 A := \{l \in \text{Link } \pi \mid \text{Exactly one premise of } l \text{ is in } \text{Emp}_a A\} \quad (42)$$

Lemma 2.0.13 (Realisation Lemma). *Let π be a cut-free proof structure satisfying the long trip condition, let $a \in \{\uparrow, \downarrow\}$ and A an occurrence of a formula in π . Define the following function:*

$$S : \text{Link}_{\mathcal{R},a}^0 A \longrightarrow \{L, R\}$$

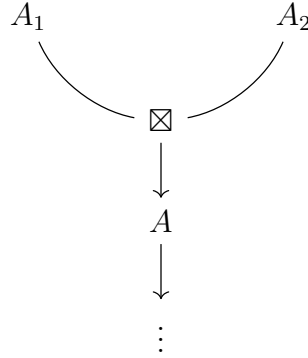
$$l \longmapsto \begin{cases} L, & \text{if the right premise of } l \text{ is in } \text{Emp}_a A \\ R, & \text{if the left premise of } l \text{ is in } \text{Emp}_a A \end{cases}$$

and extend this to a switching $\hat{S} : \text{Link } \pi \longrightarrow \{L, R\}$ arbitrarily. Then

$$\text{Emp}_a A = \text{Visit}_{\hat{S}}(A, a) \quad (43)$$

Proof. We proceed by induction on the size $|\text{Link}_a(A)|$ of the set $\text{Link}_a(A)$. For the base case, assume $|\text{Link}_a(A)| = 0$. The formula A is part of an axiom link and so $\text{Emp}_\uparrow A = A, \sim A$ and $\text{Emp}_\downarrow A = A$, the result follows easily.

Now assume that $|\text{Link}_a A| = n > 0$ and the result holds for any formula B such that $|\text{Link}_a B| < n$. First say $a = \uparrow$, and A is a conclusion of either a tensor or a par link



where $\boxtimes \in \{\otimes, \wp\}$ and $A = A_1 \otimes A_2$ or $A = A_1 \wp A_2$. By (4) we have

$$\begin{aligned} \text{Emp}_\uparrow A &= \text{Emp}_\uparrow A_1 \cup \text{Emp}_\uparrow A_2 \cup \{A\} \\ &= \text{Visit}_{\hat{S}}(A_1, \uparrow) \cup \text{Visit}_S(A_2, \uparrow) \cup \{A\} \\ &= \text{Visit}_{\hat{S}}(A, \uparrow) \end{aligned}$$

where the second equality follows from the inductive hypothesis.

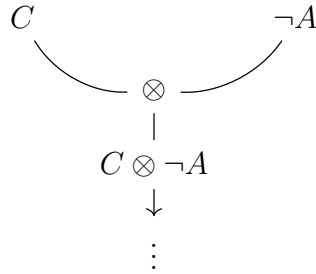
Assume A is part of an axiom link. By (1)

$$\text{Emp}_\uparrow A = \text{Emp}_\downarrow(\neg A) \cup \{A\} \quad (44)$$

with

$$|\text{Link}_\uparrow A| = |\text{Link}_\downarrow(\neg A)| \quad (45)$$

Since $|\text{Link}_\downarrow(\sim A)| > 0$ we necessarily have that $\sim A$ is not a conclusion. Thus, since π is cut-free, A is connected to an occurrence $\sim A$ which is a premise to either a tensor link or a par link. In the case of the former, we have:

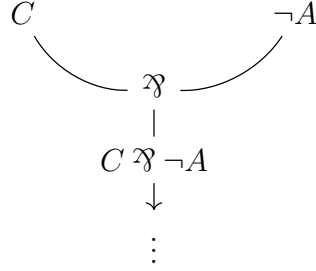


then by (5):

$$\begin{aligned} \text{Emp}_\downarrow(\neg A) &= \text{Emp}_\uparrow C \cup \text{Emp}_\downarrow(C \otimes \neg A) \cup \{\neg A\} \\ &= \text{Visit}_{\hat{S}}(C, \uparrow) \cup \text{Visit}_{\hat{S}}(C \otimes \neg A, \downarrow) \cup \{\neg A\} \\ &= \text{Visit}_{\hat{S}}(\neg A, \downarrow) \end{aligned}$$

where the second equality follows from the inductive hypothesis.

If $\sim A$ is a premise of a par link



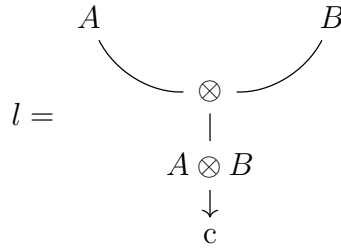
then by construction of \hat{S} , where we use the specific definition of S for the first time,

$$\begin{aligned}
\text{Emp}_{\downarrow}(\neg A) &= \{\neg A\} \\
&= \text{Visit}_{\hat{S}}(\neg A, \downarrow)
\end{aligned}$$

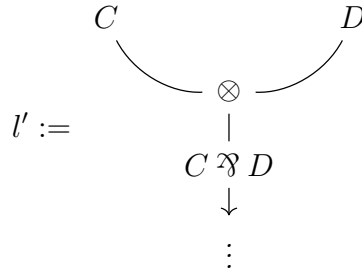
The case when $a = \downarrow$ is exactly similar and so we omit the proof. □

Definition 2.0.14. A tensor or par link is **terminal** if it is a conclusion.

Corollary 2.0.15. Let π be a cut-free proof structure satisfying the long trip condition. Let



be a terminal tensor link of π . Then π admits a par link



such that either $C \in \text{Emp}_{\uparrow} A$ and $D \in \text{Emp}_{\uparrow} B$ or $C \in \text{Emp}_{\uparrow} B$ and $D \in \text{Emp}_{\uparrow} A$ if and only if for any switching S of π we have that either

$$\text{Emp}_{\uparrow} A \subsetneq \text{Visit}_S(A, \uparrow) \quad \text{or} \quad \text{Emp}_{\uparrow} B \subsetneq \text{Visit}_S(B, \uparrow)$$

Proof. Say π admitted l' and $C \in \text{Emp}_{\uparrow} A$ and $D \in \text{Emp}_{\uparrow} B$. If the switching S is such that $S(l) = L$ then $C \wp D \in \text{Visit}_S(B) \setminus \text{Emp}_{\uparrow} B$ and if $S(\tau) = R$ then $C \wp D \in \text{Visit}_S(A) \setminus \text{Emp}_{\uparrow} A$. The other case is similar.

Conversely, say π admits no such par link l' , that is, assume

$$\text{Link}_{\wp, \uparrow}^0(A) \cap \text{Link}_{\wp, \uparrow}^0(B) = \emptyset \tag{46}$$

Then there is by Lemma 2.0.13 a well defined function

$$S : \text{Link}_{\wp, \uparrow}^0(A) \cup \text{Link}_{\wp, \uparrow}^0(B) \longrightarrow \{L, R\}$$

which extends to a switching \hat{S} such that

$$\text{Emp}_{\uparrow} A = \text{Visit}_{\hat{S}}(A, \uparrow) \quad \text{and} \quad \text{Emp}_{\uparrow} B = \text{Visit}_{\hat{S}}(B, \uparrow) \tag{47}$$

□

Lemma 2.0.16 (Separation Lemma). *A cut-free proof structure π satisfying the long trip condition, with only tensor links amongst its conclusions admits a tensor link*

$$l := \begin{array}{c} A \quad \quad B \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \otimes \\ \quad \quad | \\ A \otimes B \\ \quad \quad \downarrow \\ c \end{array}$$

satisfying

$$\mathcal{O}(\pi) = \text{Emp}_\uparrow A \cup \text{Emp}_\uparrow B \cup \{A \otimes B\} \quad (48)$$

Moreover, removing $A \otimes B$ results in a disconnected graph with each component a proof structure satisfying the long trip condition.

Proof. Consider the set of tensor links $\text{Link}_\otimes(\pi)$ of π . We endow this with the following partial order \leq : a pair of links:

$$l' := \begin{array}{c} A \quad \quad B \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \otimes \\ \quad \quad | \\ A \otimes B \\ \quad \quad \downarrow \\ \vdots \end{array} \quad \quad l'' := \begin{array}{c} C \quad \quad D \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \otimes \\ \quad \quad | \\ C \otimes D \\ \quad \quad \downarrow \\ \vdots \end{array}$$

are such that $l' \leq l''$ if $\text{Emp}_\uparrow A \cup \text{Emp}_\uparrow B \subseteq \text{Emp}_\uparrow C \cup \text{Emp}_\uparrow D$. Let l (with conclusion $A \otimes B$ say) be a tensor link maximal with respect to \leq . We show that l satisfies the required property.

Say $\mathcal{O}(\pi) \neq \text{Emp}_\uparrow A \cup \text{Emp}_\uparrow B \cup \{A \otimes B\}$. Then by Lemma 2.0.15 there exists a par link

$$l' = \begin{array}{c} C \quad \quad D \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \wp \\ \quad \quad | \\ C \wp D \\ \quad \quad \downarrow \\ \vdots \end{array}$$

such that either $C \in \text{Emp}_\uparrow A$ and $D \in \text{Emp}_\uparrow B$ or $C \in \text{Emp}_\uparrow B$ and $D \in \text{Emp}_\uparrow A$. We show the proof in the case of the former. Since π admits no terminal par links, this link is above a tensor link

$$l'' := \begin{array}{c} E \quad \quad F \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \otimes \\ \quad \quad | \\ E \otimes F \\ \quad \quad \downarrow \\ \vdots \end{array}$$

Notice that if $l'' = l$, then either $C \wp D \in \text{Emp}_\uparrow A$ or $C \wp D \in \text{Emp}_\uparrow B$ which in either case implies $\text{Emp}_\uparrow A \cap \text{Emp}_\uparrow B \neq \emptyset$, contradicting Corollary 2.0.10, 3, and so $l'' \neq l$. Without any loss of generality, assume that σ sits above F . Let S be a switching of π so that $\text{Emp}_\uparrow F = \text{Visit}_S(F, \uparrow)$ and so that $S(l') = L$,

which exists by Lemma 2.0.15. Let $t = (x_1, \dots, x_n)$ be the long pretrip of π with respect to S satisfying $x_1 = F \uparrow$. We have by Lemma 2.0.5 that t takes the following shape:

$$\uparrow F, \dots, \uparrow (C \wp D), \uparrow C, \dots, D \downarrow, \uparrow D, \dots, C \downarrow, (C \wp D) \downarrow, \dots, F \downarrow, \dots \quad (49)$$

We have that $D \in \text{Emp}_\uparrow B$ so for simplicity, rewrite (49) as $t' = (x_{1+k}, \dots, x_{n+k})$ for some $k > 0$ (where $i + k$ means $i + k \bmod n$) so that t takes the shape

$$\dots, \uparrow F, \dots, \uparrow (C \wp D), \uparrow C, \dots, D \downarrow, \uparrow D, \dots, C \downarrow, (C \wp D) \downarrow, \dots, F \downarrow, \dots \quad (50)$$

with $\uparrow B$ occurring to the left of $D \downarrow$ and $B \downarrow$ occurring to the right of $\uparrow D$. We have that $C \notin \text{Emp}_\uparrow B$ and so by Corollary 2.0.9:

$$\uparrow B \text{ occurs in } \uparrow C, \dots, D \downarrow \text{ and } B \downarrow \text{ occurs in } \uparrow D, \dots, C \downarrow \quad (51)$$

However, this implies that $B \in \text{Visit}_S(F, \uparrow)$ which by Lemma 2.0.13 implies $B \in \text{Emp}_\uparrow F$.

By reversing the switching of l' and interchanging the rolls of C, D in the above argument, we also have that $A \in \text{Emp}_\uparrow F$, contradicting the maximality of l . This proves the first claim.

For the second claim, since $\mathcal{O}(\pi) = \text{Emp}_\uparrow A \cup \text{Emp}_\uparrow B \cup \{A \otimes B\}$ we have by Lemma 2.0.15 that

$$\text{Link}_{\wp, \uparrow}^0(A \otimes B) = \emptyset \quad (52)$$

and we saw in the proof of Lemma 2.0.13 that a switching S which realises $\text{Emp}_\uparrow A$ is given by setting all switchings arbitrarily except for those in $\text{Link}_{\wp, \uparrow}^0(A \otimes B)$. This means that for any switching S of π :

$$\text{Visit}_S(A, \uparrow) = \text{Emp}_\uparrow A \quad \text{and} \quad \text{Visit}_S(B, \uparrow) = \text{Emp}_\uparrow B \quad (53)$$

which is to say the two subproof structures given by removing $A \otimes B$ never admit a short trip, that is, they each satisfy the long trip condition. \square

Theorem 2.0.17 (The Sequentialisation Theorem). *A proof structure π (possibly with cuts) satisfies the long trip condition if and only if π is a proof net.*

Proof. First assume that π is cut-free.

We proceed by induction on the size $|\text{Link } \pi|$ of the set $\text{Link } \pi$. If there this is zero then π consists of a single axiom link and so the result is clear.

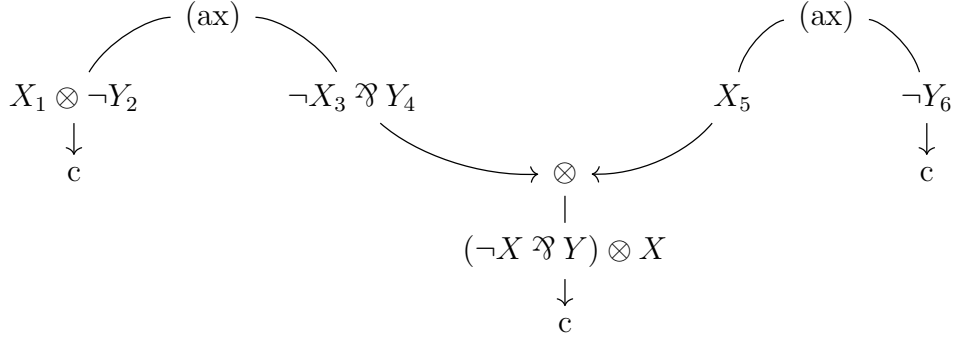
For the inductive step, we consider two cases, first say π admits a par link for a conclusion. Then removing this par link clearly results in two cut-free subproof structures satisfying the long trip condition and so the result follows from the inductive hypothesis. If no such terminal par link exists, then by the Separation Lemma there exists some tensor link in the conclusion for which we can remove and apply the inductive hypothesis.

Now say that π contained cuts. We replace each cut with a tensor link to create a new proof ζ . That there exists a proof Ξ which maps to ζ follows from the part of the result proved already as ζ is cut-free. We adapt Ξ appropriately by replacing \otimes -rules by cut-rules and we are done. \square

3 Proofs as permutations

Definition 3.0.1. Let π be a proof net. Let $\mathcal{P}(\pi)$ denote the disjoint union of all the unoriented axioms of all formulas which are conclusions to axiom links in π .

Example 3.0.2. Let π denote the following proof net, where X, Y are atomic. For simplicity, we usually attach artificial labels to the conclusions of axiom links. So, in the following example, for any integer i and any $Z = X, Y$, the notation Z_i simply means Z , also Z means $(Z, +)$ and $\neg Z$ means $(Z, -)$.



Then:

$$\mathcal{P}(\pi) = \{X\} \coprod \{Y\} \coprod \{X\} \coprod \{Y\} \coprod \{X\} \coprod \{Y\} = \{X_1, Y_2, X_3, Y_4, X_5, Y_6\} \quad (54)$$

Definition 3.0.3. Let A be a formula with sequence of oriented atoms $((X_1, x_1), \dots, (X_n, x_n))$. The **sequence of unoriented atoms** of A is (X_1, \dots, X_n) and the **set of unoriented atoms** of A is the disjoint union $\{X_1\} \coprod \dots \coprod \{X_n\}$.

Definition 3.0.4. Let π be a proof net with axiom links l_1, \dots, l_n say. For each $i = 1, \dots, n$ the link l_i defines a permutation τ_{l_i} on the set $\mathcal{P}(\pi)$ in the following way: if l_i has conclusions $A, \neg A$ then the j^{th} element of the sequence of unoriented atoms of A is mapped via τ_{l_i} to the j^{th} element of the sequence of unoriented atoms of $\neg A$. We define α_π to be the product of all these permutations.

$$\alpha_\pi := \tau_{l_1} \dots \tau_{l_n} \quad (55)$$

We call this permutation the **axiom link permutation associated to π** .

There is another permutation on $\mathcal{P}(\pi)$ defined symmetrically but where the cut links are considered rather than the axiom links. We denote this permutation γ_π .

There is yet another permutation on $\mathcal{P}(\pi)$. Let S be a switching of π . For each unoriented axiom $X \in \mathcal{P}(\pi)$, corresponding to the formula A say, let $\beta_\pi^S(X)$ denote the unoriented axiom corresponding to the first occurrence in $\text{PTrip}(\pi, S, A, \downarrow)$ (Definition 2.0.7) of the form $\uparrow B$ where B is a formula labelling a conclusion of an axiom link in π .

The set of all permutations of the second form is denoted:

$$\Sigma(\pi) := \{\beta_\pi^S \mid S \text{ is a switching of } \pi\} \quad (56)$$

We will often denote elements of $\beta_\pi^S \in \Sigma(\pi)$ simply by β .

Example 3.0.5. Continuing with Example 3.0.2, and using the artificial labels we attributed to the atomic formulas of conclusions, we have $\alpha_\pi = (13)(24)(56)$ and since π admits no cut links we have $\gamma_\pi = \text{id}_{\mathcal{P}(\pi)}$.

The set $\Sigma(\pi)$ is uninteresting so we consider a more complicated example. Before doing this though, we introduce a simple Lemma for the sake of easing notation.

Lemma 3.0.6. Let π be a proof net with conclusions A_1, \dots, A_n and let ζ be a proof net obtained by beginning with π and in any order forming par links which connect all the conclusions A_1, \dots, A_n so that ζ has conclusions B_1, \dots, B_m where $m \leq n$ and each B_i is constructed only by \wp and a subset of the formulas A_1, \dots, A_n . Then $\Sigma(\pi) = \Sigma(\zeta)$.

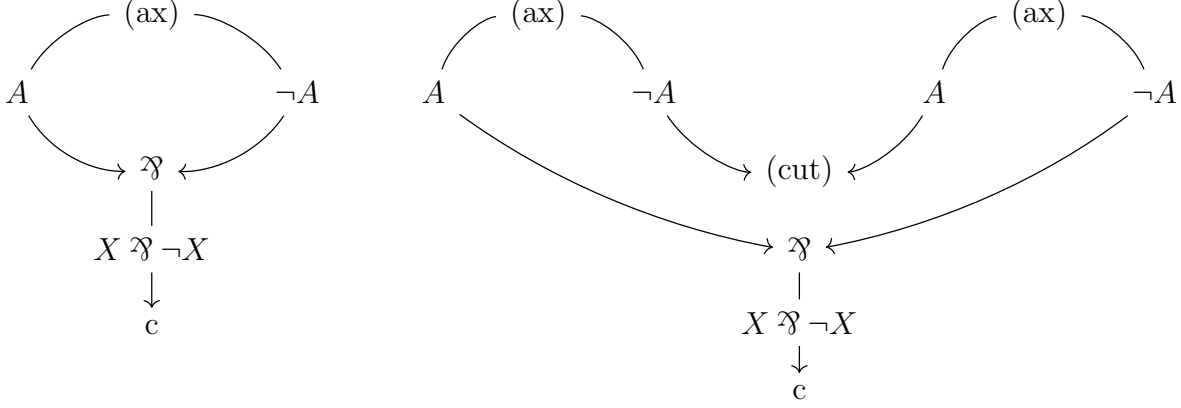
Proof. Easy proof by induction on the integer given by the number of par links in ζ minus the number of par links in π . \square

[illegible]
$$\Sigma(\pi_1) = \{(1375), (1357), (1753), (1573)\} \quad (57)$$

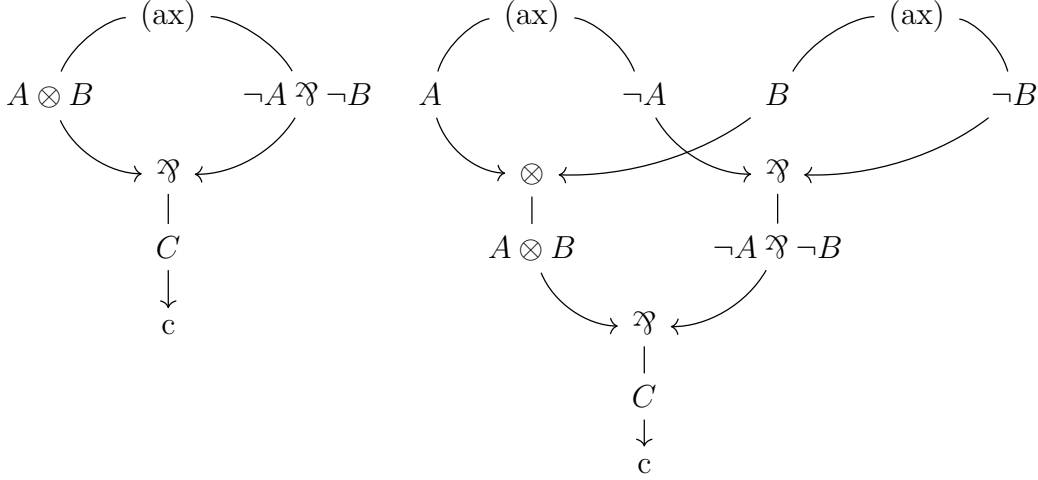
We can now rephrase the longtrip condition of Section 2 in terms of permutations.

18

If π is a proof net admitting a single conclusion, then it is *not* the case that α_π uniquely determines this proof net, for instance, the following two pairs of proof nets admit the same axiom link permutations.



In what follows, C denotes $(A \otimes B) \wp (\neg A \wp \neg B)$.



The following Proposition says that these are the only cases when this happens.

Proposition 3.0.10. *Let π be a cut-free proof net with only one conclusion, assume also that all conclusions of all axiom links are atomic. Then π is determined uniquely by α_π and the conclusion A .*

Proof. Let A be the soul conclusion of π . Consider the typing tree of A , this is a binary tree with all leaves labelled by atomic propositions. The proof net π is constructed from this by choosing axiom links, this choice of axiom links is determined by α_π . \square

Corollary 3.0.11. *Let π be a cut-free proof net with conclusions A_1, \dots, A_n , assume also that all conclusions of all axiom links are atomic. Then π is determined uniquely by α_π and the conclusions A_1, \dots, A_n .*

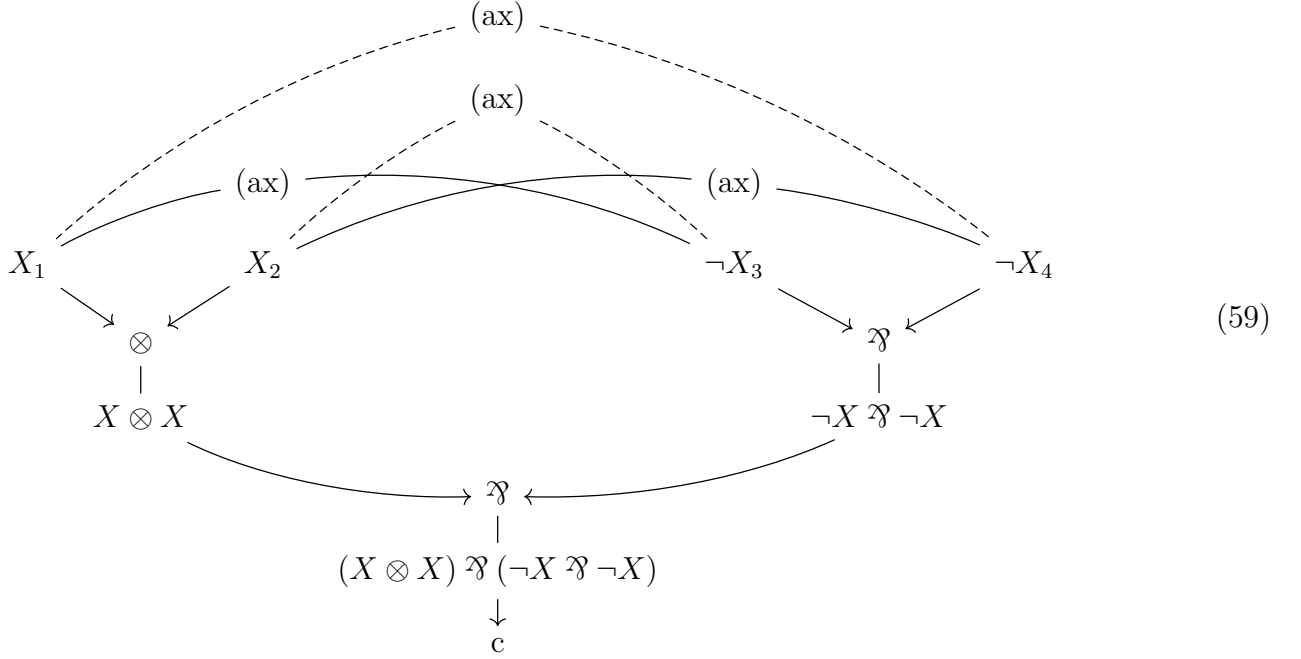
Proof. Let π_1, π_2 be two such proof nets. Construct a new proof net π'_1 from π_1 by creating par links (in any order) so that π'_1 has a unique conclusion $A_1 \wp \dots \wp A_n$. Construct a new proof net π'_2 in a similar way and make the same choices of par links made when π'_1 was constructed. No axiom links were altered in this process, and so $\alpha_{\pi_1} = \alpha_{\pi_2} = \alpha_{\pi'_1} = \alpha_{\pi'_2}$. It follows from Proposition 3.0.10 that $\pi'_1 = \pi'_2$. For $i = 1, 2$, remove the par links which appear in π'_i but not in π_i and conclude that $\pi_1 = \pi_2$. \square

Example 3.0.12. Consider the following formula, where X is atomic.

$$((X, +) \otimes (X, +)) \wp ((X, -) \wp (X, -)) \quad (58)$$

corresponds to the sub-proof structure given by ignoring the dashed lines and the axiom links of (59), for clarity, we label the conclusions of axiom links explicitly with an integer $i \in \{1, 2, 3, 4\}$, but each X_i just

means the formula X , for simplicity, we write X for $(X, +)$ and $\neg X$ for $(X, -)$, the important point is that X is atomic.



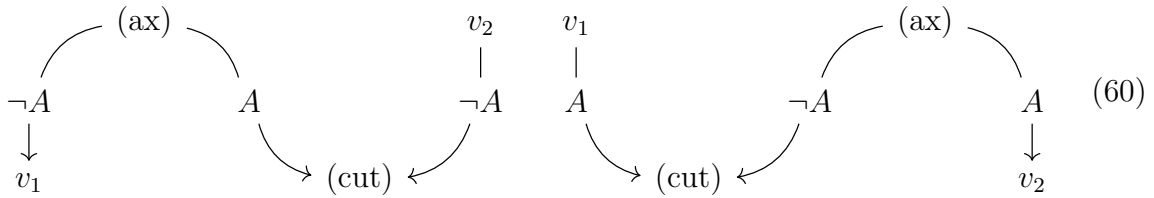
The proof net given by ignoring the dashed lines in (59) corresponds to the permutation (12)(34) (ie, the permutation $X_1 \leftrightarrow X_2, X_3 \leftrightarrow X_4$), and that given by ignoring the axiom links and including the dashed lines is (14)(23).

Definition 3.0.13. A **super normal form** is a proof net π which is cut-free and is such that all conclusions of all axiom links are atomic.

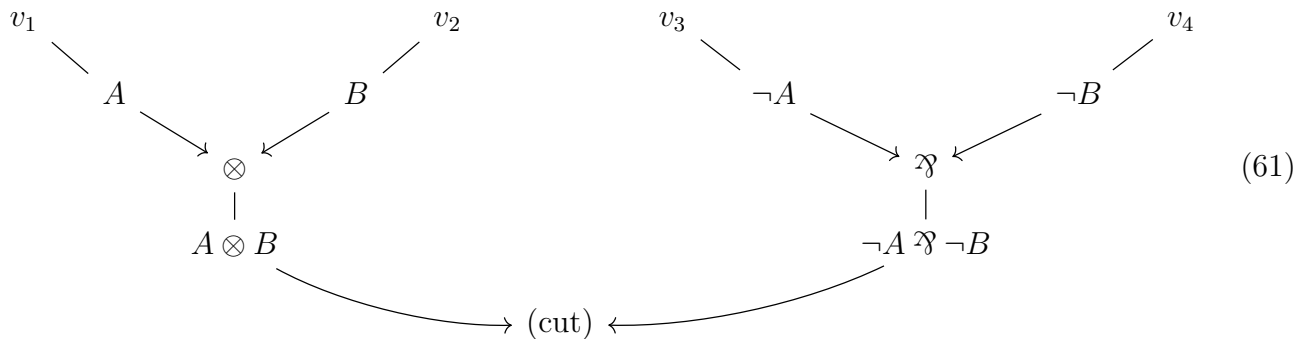
In light of Corollary 3 we wish to canonically associate an arbitrary proof net π to a super normal form.

4 The dynamics of MLL

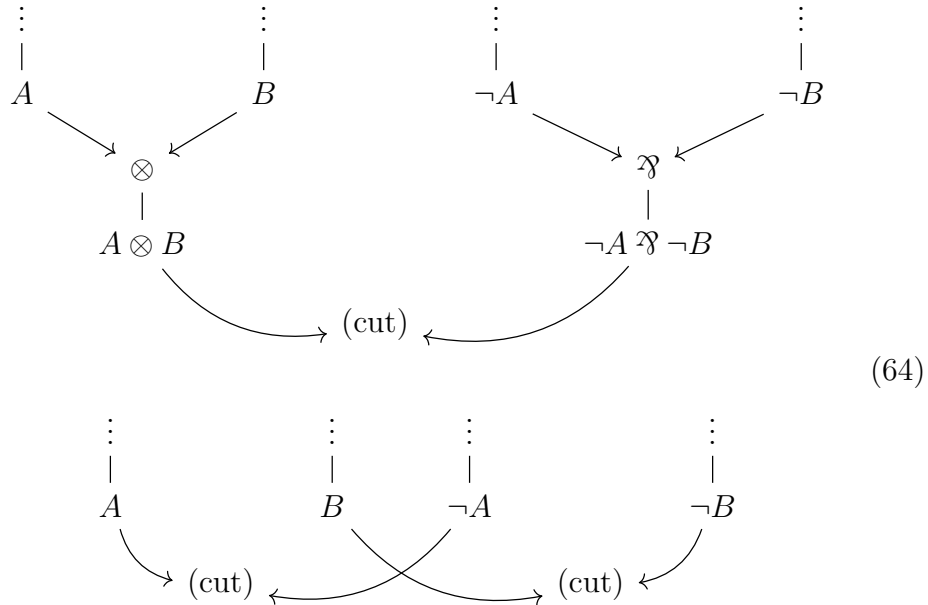
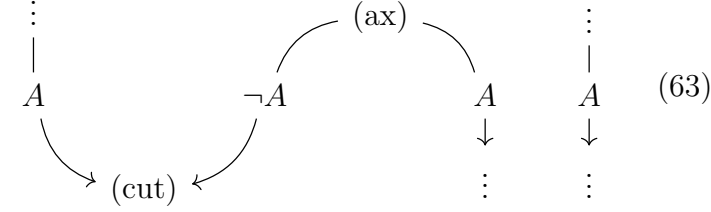
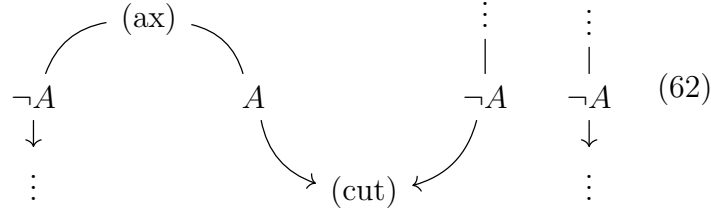
Definition 4.0.1. A subgraph of a proof structure π of one of the following forms, where $v_1 \neq v_2$ are distinct vertices in π , is an *a-redex*.



A subgraph of a proof structure of the following form where v_1, v_2, v_3, v_4 are vertices in π is an *m-redex*.



Definition 4.0.2. Multiplicative linear logic proof structures come equipt with three reduction rules, these reduction rules apply to proofs which admit either an a -redex or an m -redex. More precisely, given a multiplicative linear logic proof structure π admitting an a -redex ζ of the form given on the left in (62), the reduction of π is the proof π' given by replacing the subgraph ζ in π by what is displayed on the right in (62). Similarly for if π admits an a -redex of the form given on the left of (63) or if π admits an m -redex.



A **reduction** is a pair of proof structures (π, π') where π' is the result of applying one of the reduction rules just described to π . We write $\pi \longrightarrow_{(\text{cut})} \pi'$ when (π, π') is a reduction.

Proposition 4.0.3 (Church-Rosser). *If π_1 is a proof structure and $\pi_1 \longrightarrow_{\text{cut}} \pi_2, \pi_1 \longrightarrow_{\text{cut}} \pi_3$ then there exists a proof structure π_4 such that $\pi_2 \longrightarrow_{\text{cut}} \pi_4, \pi_3 \longrightarrow_{\text{cut}} \pi_4$.*

Proof sketch. The key observation is that reducing any redex in a proof does not eliminate any other redex. \square

Definition 4.0.4. Let π be a proof net possibly containing cut links. A **reduction sequence** is a sequence

$$\pi = \pi_0 \longrightarrow_{\text{cut}} \pi_1 \longrightarrow_{\text{cut}} \dots \longrightarrow_{\text{cut}} \pi_n \quad (65)$$

with π_n cut-free.

Lemma 4.0.5. *Every proof net π admits a reduction sequence.*

Proof. Given a cut link l with premises $A, \neg A$ say, the **complexity of l** , $c(l)$ is the sum of the number of occurrences of \otimes and the number of occurrences of \wp in A . We proceed by induction on the maximum of the complexities of all cut links in π .

Say this maximum is 0. Then all cut-links have the shape of either (62) or (63) (using the fact that π is a *proof net*, the hypothesis that $v_1 \neg v_2$ in Definition 4.0.2 is satisfied). We can reduce these redexes (in any order) to deduce the result.

Now say the maximum is $n > 0$. We then apply (64) to all cut links of complexity n (in any order) to obtain a new proof structure ζ . We wish to use the inductive hypothesis on ζ but we must make sure that ζ satisfies the longtrip condition. This follows easily by considering the contrapositive. Any pretrip (long or short) of ζ appears as a subsequence of some pretrip of π , so if ζ admits a short trip so does π . We can now apply the inductive hypothesis and we are done. \square

Definition 4.0.6. Let $\text{Red } \pi$ denote the set of all reduction sequences of π . The **length** $l(\underline{x})$ of a reduction sequence $\underline{x} \in \text{Red } \pi$ is the length of the sequence \underline{x} .

Corollary 4.0.7. *The length of a reduction path is independent of the choice of reduction path.*

Proof. The proof is purely geometric. Let

$$\underline{x} := (\pi = \pi_1 \longrightarrow_{\text{cut}} \dots \longrightarrow_{\text{cut}} \pi_n) \quad (66)$$

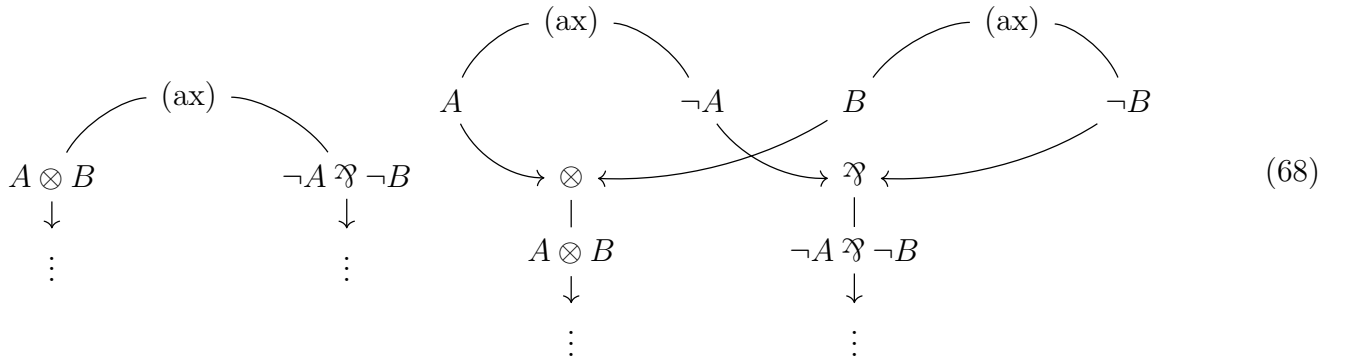
be the reduction path described by Lemma 4.0.5 and let

$$\underline{y} := (\pi = \zeta_0 \longrightarrow_{\text{cut}} \dots \longrightarrow_{\text{cut}} \zeta_n) \quad (67)$$

be any other reduction sequence. By Lemma 4.0.3 we have $\pi_n = \zeta_n$. Also using 4.0.3, the pair of reduction paths can be completed to some grid defined by a subset of $\mathbb{N} \times \mathbb{N}$. All paths p consisting of only upwards steps or right steps such that p is bound to this grid have the same length and so $l(\underline{x}) = l(\underline{y})$. \square

Definition 4.0.8. The proof of Corollary 4.0.7 shows that every reduction path of a proof net π leads to the same cut-free proof ζ . We call ζ the **normal form** of π .

Definition 4.0.9. A pair of proof nets (π, π') where π' is obtained from π via replacing some subgraph of π of the form on the left of (68) is an η -**expansion**. We write $\pi \longrightarrow_{\eta} \pi'$. An η -**increx** is a subgraph of a proof net of the form given on the left of (68).



Lemma 4.0.10. *If $\pi \longrightarrow_{\eta} \pi'$ and $\pi \longrightarrow_{\eta} \pi''$ then there exists π''' such that $\pi' \longrightarrow_{\eta} \pi'''$ and $\pi'' \longrightarrow_{\eta} \pi'''$.*

Proof. Any η -expansion does not alter any other η -index. \square

Due to the finiteness of proof nets and their formulas, the process of choosing any η -index and reducing it eventually terminates and results in a proof net with no η -indexes.

Corollary 4.0.11. *The process of reducing all redexes in a proof net π and then expanding all η -indexes terminates and results in a cut-free proof net where all conclusions of all axiom links are atomic.*

Definition 4.0.12. The result of applying process 4.0.11 to a proof net π is the **super normal form** of π .

5 Geometry of Interaction Zero, modelling the dynamics of MLL

Geometry of interaction zero requires that all formulas occurring in axiom links are atomic, every proof structure not of this form can be related to one which is by η -equivalence: We begin with the following crucial observation:

Lemma 5.0.1. *Let π be a proof structure and assume there is a cut in π of A against $\neg A$. Write*

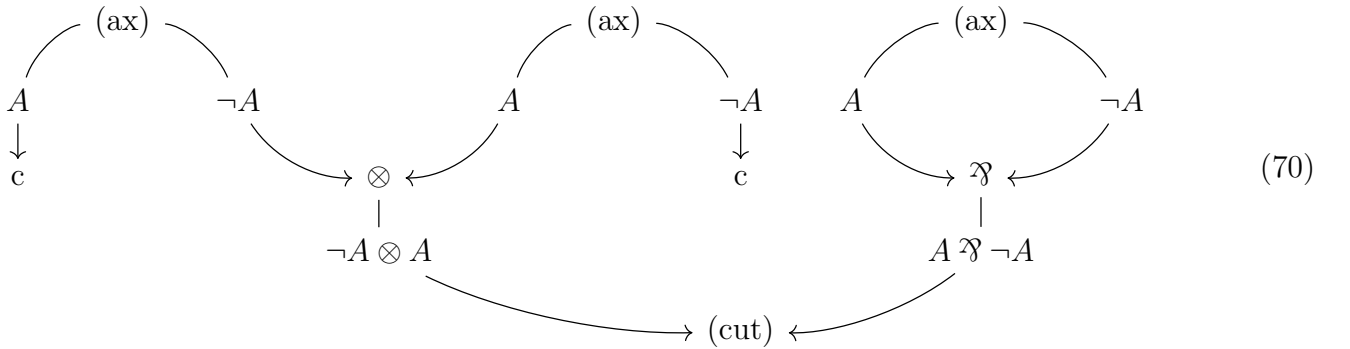
$$A := A_1 \boxtimes_1 \dots \boxtimes_{n-1} A_n \quad (69)$$

where for each i we have $\boxtimes_i \in \{\otimes, \wp\}$. Let ζ be a proof structure equivalent to π under cut-reduction which is obtained by reducing all m -redexes (Definition 4.0.2). Then for each i there exists a unique cut link l_i with premises $A_i, \neg A_i$.

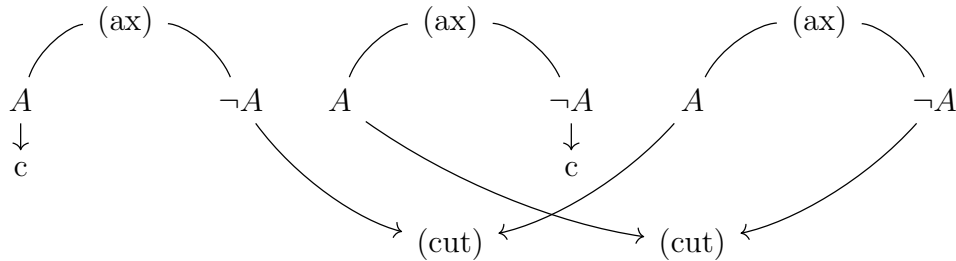
Proof. By induction on n , where the base case follows trivially and the inductive step by inspection of (64). \square

Definition 5.0.2. We define a permutation γ_π on $\mathcal{P}(\pi)$. Let l be a cut link in π with premises $A, \neg A$, say. For each atomic formula A_i of A there exists by Lemma 5.0.1 a cut link l_i in the normal form ζ of π so that the premises of l_i are $A_i, \neg A_i$. Let $A_i, \neg A_i$ have corresponding unoriented atoms X_i, Y_i . Let γ be the permutation which swaps X_i and Y_i . For all other elements of $X \in \mathcal{P}(\pi)$ we define $\gamma_\pi(X) = X$.

Example 5.0.3. We denote by π the following proof net.



Reducing the cut-redex we obtain:



The cut link in π consists of $\neg A \otimes A$ and $A \wp \neg A$, the order of these formulas was respected by the cut-reduction step, in the sense made precise by Lemma 5.0.1, as seen in this example as the resulting cuts are between “the two first elements” and “the two last elements”, ie, between $\neg A$ and A , and between A and $\neg A$.

Lemma 5.0.1 will be used to prove both Geometry of Interaction zero and Geometry of Interaction One.

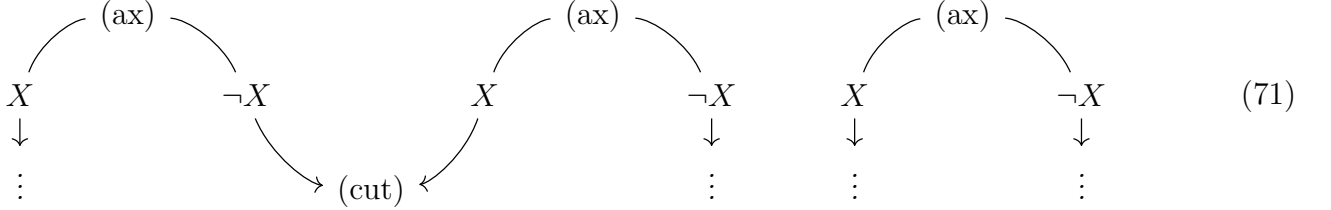
Lemma 5.0.4. *The set $\mathcal{P}(\pi)$ is invariant under reduction of m -redexes and η -expansion (defined in Remark 3.0.8). More precisely, we have the following two statements.*

- Say π' produced by reducing an m -redex (64) in π , then $\mathcal{P}(\pi) = \mathcal{P}(\pi')$.
- Say $\pi \rightarrow_\eta \pi'$ (see Definition 4.0.9), then $\mathcal{P}(\pi) = \mathcal{P}(\pi')$.

Proof. For the first claim we simply notice that rule (64) has no effect on the axiom links of π . For the second we see that the order of the sequence of unoriented atoms of A, B is explicated by the axiom links produced by an η -expansion. \square

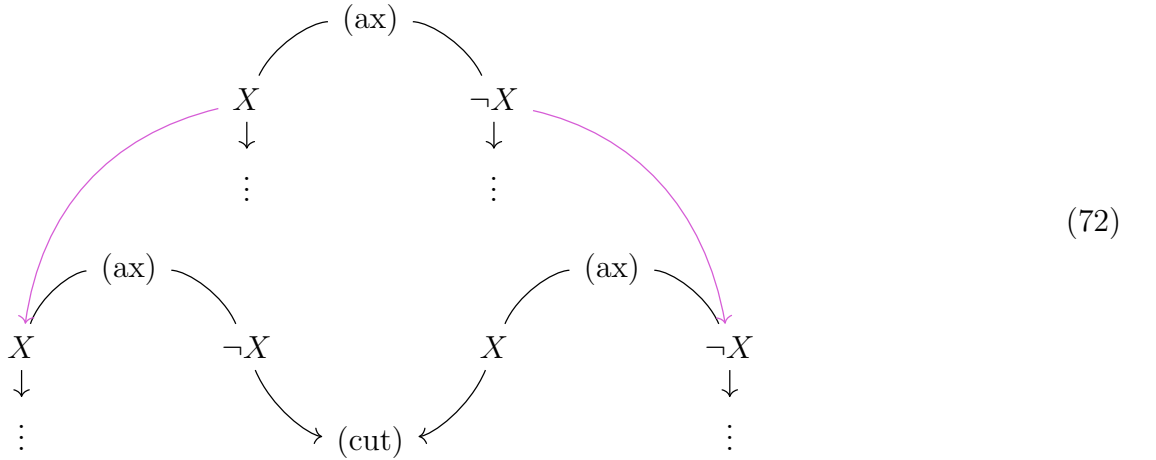
Hence, when considering $\mathcal{P}(\pi)$, we can always assume without loss of generality that π contains no m -redexes and that all conclusions of all axiom links of π are atomic.

Lemma 5.0.5. *Let π be a proof net admitting no m -redexes and assume that all conclusions of all axiom links of π are atomic. All redexes $\pi \rightarrow_{(\text{cut})} \pi'$ are necessarily of the following form.*



Proof. All redexes of π are a -redexes, so all redexes of π are of the form (62) or (63), but since all the axiom links have atoms as conclusions, it must be the case that the cut link in (62), (63) have premises which are also atoms. These atoms can only possibly exist if they are conclusions to an axiom link, and so we obtain the form given in the statement. \square

Definition 5.0.6. Say π is a proof net with no m -redexes and all conclusions of all axiom links are atomic. Moreover, say there is a reduction $\pi \rightarrow \pi'$ which by Lemma 5.0.4 is of the form (71). We define a function $\iota : \mathcal{P}(\pi') \rightarrow \mathcal{P}(\pi)$ given by the following schema:



The following proposition states that super normal forms are completely determined by their axiom links, and conclusion.

Definition 5.0.7. Let π be a proof net and consider $\mathcal{P}(\pi)$, in light of Lemma 5.0.4 we can assume without loss of generality that π admits no m -redexes and that all conclusions of all axiom links in π are atomic. Let ζ be the corresponding super normal form established by Corollary 3. Let $(\pi = \pi_1, \dots, \pi_n = \zeta)$ be a sequence of cut reductions. These induce a family of functions:

$$\mathcal{P}(\zeta) = \mathcal{P}(\pi_n) \rightarrow \mathcal{P}(\pi_{n-1}) \rightarrow \dots \rightarrow \mathcal{P}(\pi_2) \rightarrow \mathcal{P}(\pi_1) = \mathcal{P}(\pi) \quad (73)$$

Composing these determines a function $\iota_\pi : \mathcal{P}(\zeta) \rightarrow \mathcal{P}(\pi)$.

Remark 5.0.8. It follows from Proposition 4.0.3 that the function ι_π is independent of the choice of reduction path used to define it.

We give an alternate characterisation of the image of ι_π .

Lemma 5.0.9. *Let π be a proof net admitting no m -redexes and assume all conclusions of all axiom links are atomic. A formula A in π is in $\text{im } \iota_\pi$ if and only if it is not premise to a cut link.*

Proof. Say A is premise to a cut link. Since $A \in \mathcal{P}(\pi)$ it is also the case that A is conclusion to a cut link. Hence there exists a cut reduction which removes A , and so A is not in the image of ι_π .

Now say A is *not* premise to a cut link and so A is necessarily *not* part of an a -redex. There are no m -redexes in π and so all cut reductions reduce a -redexes. Hence A survives the cut reduction process. In other words, $A \in \text{im } \iota_\pi$. \square

Definition 5.0.10. Let π be a proof net. We describe a final permutation δ_π on $\mathcal{P}(\pi)$. Recall the injective function ι_π of Definition 5.0.7. For each $X \in \mathcal{P}(\pi)$ let d_i denote the least integer such that

$$(\alpha_\pi \circ \gamma_\pi)^{d_i}(X) \in \text{Ess}(\pi) \quad (74)$$

Notice that such an integer d_i always exists as π is a proof net (as π satisfies the longtrip condition, see Section 2).

We then define the following permutation on $\mathcal{P}(\pi)$, the permutations α_π, γ_π are defined in Definition 3.0.4:

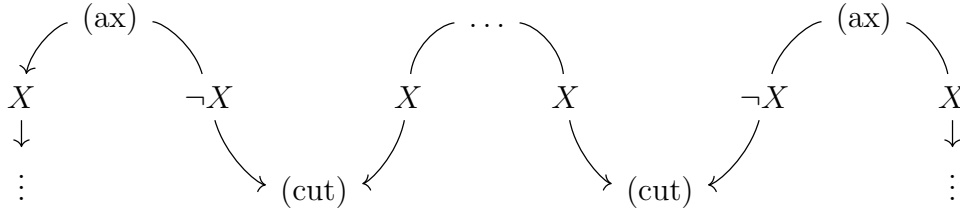
$$\delta_\pi(X) = (\alpha_\pi \circ \gamma_\pi)^{d_i}(X) \quad (75)$$

Theorem 5.0.11. *[Geometry of Interaction zero] Let π be a proof net possibly with cuts and let ζ be the normal form of π (Definition 4.0.8). Then*

$$\delta_\pi = \iota_\pi \alpha_\zeta \quad (76)$$

Proof. By inspection of (64) we have that γ_π is invariant under reduction of m -redexes. Also, α_π is clearly invariant under reduction of m -redexes, thus we can assume that π admits no m -redex. Furthermore, by inspection of (68) we see that α_π is invariant under η -expansion, it is also clear that γ_π is invariant under η -expansion. Thus we can also assume that all conclusions of all axiom links of π are atomic.

All cut links appear inside “chains” of axiom and cut links, such as in the following Diagram.



By Lemma 5.0.9, all formulas in the “interior” of these chains are not in $\text{im } \iota_\pi$. Hence, δ_π is a product of transposes where the formulas on the two extreme ends of these chains are swapped. By considering the cut elimination rules (62), (63) we see that this is exactly the behaviour of α_ζ , and the that these two formulas are the images of the corresponding formulas in ζ . \square

References

- [1] Intuitionistic, Linear Sequent Calculus, W. Troiani.
- [2] *Linear Logic*, J.Y. Girard
- [3] *Multiplicatives*, J.Y. Girard.
- [4] *Geometry of Interaction I*, J.Y. Girard.
- [5] *Intuitionistic, linear sequent calculus* W. Troiani.
- [6] *proof nets*, W Troiani
- [7] *An introduction to proof nets*