

15-381 Artificial Intelligence  
Homework 4  
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### 1.1: Nash Equilibrium

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Let  $p_1 = (U, 1 - U)$  and  $p_2 = (L, 1 - L)$  be the probabilities of the choices for player 1 and player 2, and let  $P_1$  and  $P_2$  be matrices for their payoffs. We calculate the expected values for each of the players, as follows:

$$P_1 p_2 = \begin{bmatrix} 2 & -2 \\ 1 & 2 \end{bmatrix} \begin{bmatrix} L \\ 1 - L \end{bmatrix} = \begin{bmatrix} 2L + (-2)(1 - L) \\ L + 2(1 - L) \end{bmatrix} = \begin{bmatrix} 4L - 2 \\ 2 - L \end{bmatrix}$$
$$p_1 \cdot P_1 p_2 = \begin{bmatrix} U \\ 1 - U \end{bmatrix} \cdot \begin{bmatrix} 4L - 2 \\ 2 - L \end{bmatrix} = 4UL - 2U + 2 - L - 2U + UL = 5UL - 4U - L + 2$$

Take the derivative of this with respect to U and set equal to 0, then obtain the following:

$$5L - 4 = 0 \implies L = \frac{4}{5}$$

Similarly, find the expected value for player 2, as follows:

$$P_2 p_2 = \begin{bmatrix} 0 & 2 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} L \\ 1 - L \end{bmatrix} = \begin{bmatrix} 2(1 - L) \\ L + (-1)(1 - L) \end{bmatrix} = \begin{bmatrix} 2 - 2L \\ 2L - 1 \end{bmatrix}$$
$$p_1 \cdot P_2 p_2 = \begin{bmatrix} U \\ 1 - U \end{bmatrix} \cdot \begin{bmatrix} 2 - 2L \\ 2L - 1 \end{bmatrix} = 2U - 2UL + 2L - 1 - 2UL + U = -4UL + 3U + 2L - 1$$

Again, take the first derivative with respect to L and set equal to 0, then obtain the following:

$$-4U + 2 = 0 \implies U = \frac{1}{2}$$

Therefore, our Nash Equilibrium is:  $p_1 = \left(\frac{1}{2}, \frac{1}{2}\right)$  and  $p_2 = \left(\frac{4}{5}, \frac{1}{5}\right)$ .

To verify Nash equilibrium, we ensure first that for any mixed strategy of player 1, say  $p'_1 = (U', 1 - U')$ , we have  $p'_1 \cdot P_1 p_2 \leq p_1 \cdot P_1 p_2$ :

$$\begin{aligned} p'_1 \cdot P_1 p_2 &= 5UL - 4U - L + 2 \\ &= 5U \left(\frac{4}{5}\right) - 4U - \frac{4}{5} + 2 \\ &= \frac{3}{5} \end{aligned}$$

Also,

$$\begin{aligned} p_1 \cdot P_1 p_2 &= 5UL - 4U - L + 2 \\ &= 5 \left(\frac{1}{2}\right) \left(\frac{4}{5}\right) - 4 \left(\frac{1}{2}\right) - \frac{4}{5} + 2 \\ &= 2 - 2 - \frac{4}{5} + 2 \\ &= \frac{3}{5} \end{aligned}$$

Now, we ensure that for any mixed strategy of player 2, say  $p'_2 = (L', 1 - L')$ , we have that  $p_1 \cdot P_2 p'_2 \leq p_1 \cdot P_2 p_2$ :

$$\begin{aligned} p_1 \cdot P_1 p'_2 &= -4UL + 3U + 2L - 1 \\ &= -4 \left( \frac{1}{2} \right) L + 3 \left( \frac{1}{2} \right) + 2L - 1 \\ &= -2L + \frac{3}{2} + 2L - 1 \\ &= \frac{1}{2} \end{aligned}$$

Also,

$$\begin{aligned} p_1 \cdot P_2 p_2 &= -4UL + 3U + 2L - 1 \\ &= -4 \left( \frac{1}{2} \right) \left( \frac{4}{5} \right) + 3 \left( \frac{1}{2} \right) + 2 \left( \frac{4}{5} \right) - 1 \\ &= \frac{1}{2} \end{aligned}$$

Thus, the Nash Equilibrium holds.

## 1.2: Extensive Form to Normal Form

1. From left to right, denote the  $P_2$  nodes as 1, 2, and 3.  $P_2$  will look to maximize their score, so in node 1, we take path  $b$ , in node 2, we take path  $e$ , and in node 3, we take path  $g$ . Now, we have node 1 labelled as  $(0, 2)$ , node 2 labelled as  $(2, 2)$ , and node 3 labelled as  $(-1, 1)$ . Now that it is  $P_1$ 's turn, we find the maximum value for player 1 given these values, and we find that path  $Y$  gives us this. Therefore, the solution to the game is the path  $\boxed{Ye}$ .

We have the following conversion from extensive form game to normal form game:

	a/d/g	a/d/h	a/e/g	a/e/h	a/f/g	a/f/h	b/d/g	b/d/h	b/e/g
$X$	(2,1)	(2,1)	(2,1)	(2,1)	(2,1)	(2,1)	(0,2)	(0,2)	(0,2)
$Y$	(1,-1)	(1,-1)	(2,2)	(2,2)	(0,1)	(0,1)	(1,-1)	(1,-1)	(2,2)
$Z$	(-1,1)	(0,0)	(-1,1)	(0,0)	(-1,1)	(0,0)	(-1,1)	(0,0)	(-1,1)
	b/e/h	b/f/g	b/f/h	c/d/g	c/d/h	c/e/g	c/e/h	c/f/g	c/f/h
$X$	(0,2)	(0,2)	(0,2)	(0,0)	(0,0)	(0,0)	(0,0)	(0,0)	(0,0)
$Y$	(2,2)	(0,1)	(0,1)	(1,-1)	(1,-1)	(2,2)	(2,2)	(0,1)	(0,1)
$Z$	(0,0)	(-1,1)	(0,0)	(-1,1)	(0,0)	(-1,1)	(0,0)	(-1,1)	(0,0)

2. We have the following conversion from extensive form game to normal form game:

	a/d	a/e	a/f	b/d	b/e	b/f	c/d	c/e	c/f
$X$	(2,-2)	(2,-2)	(2,-2)	(1,-1)	(1,-1)	(1,-1)	(-1,1)	(-1,1)	(-1,1)
$Y$	(3,-3)	(2,-2)	(0,0)	(3,-3)	(2,-2)	(0,0)	(3,-3)	(2,-2)	(0,0)
$Z$	(1,-1)	(2,-2)	(0,0)	(1,-1)	(2,-2)	(0,0)	(1,-1)	(2,-2)	(0,0)

To solve the game, notice that the  $a/e$ , the  $b/d$  column, the  $a/d$  column, and the  $b/e$  columns are strictly dominated by the  $c/f$  column for player 2, so we can simply remove those columns.

By backward induction now, from the node connecting edges  $a$ ,  $b$ , and  $c$ , we have that player 2 will take edge  $c$  because it provides the most utility for player 2, and from the next two nodes connecting edges  $d, e$ , and  $f$ , player 2 will select edge  $f$ . Player 1 will now select the node containing  $(0, 0)$ , so the solution is either  $Yf$  or  $Zf$ .

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### 1.3: Electing the Median

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We can create the following cases. Suppose student  $i$  selects the temperature  $S$  such that  $S < T_i < T^*$  or  $T_i < S < T^*$ . In either case, the median will remain unchanged regardless of if he chose  $S$  or  $T_i$ , and because the utility is even less for selecting  $S$  as it is even further away from  $T^*$  than  $T_i$ , there is no gain in selecting  $S < T_i$ . Suppose we have that  $T_i < T^* < S$  or  $T^* < T_i < S$ . We have that the symmetric result will also occur, such that there is no gain.