GPU hardware

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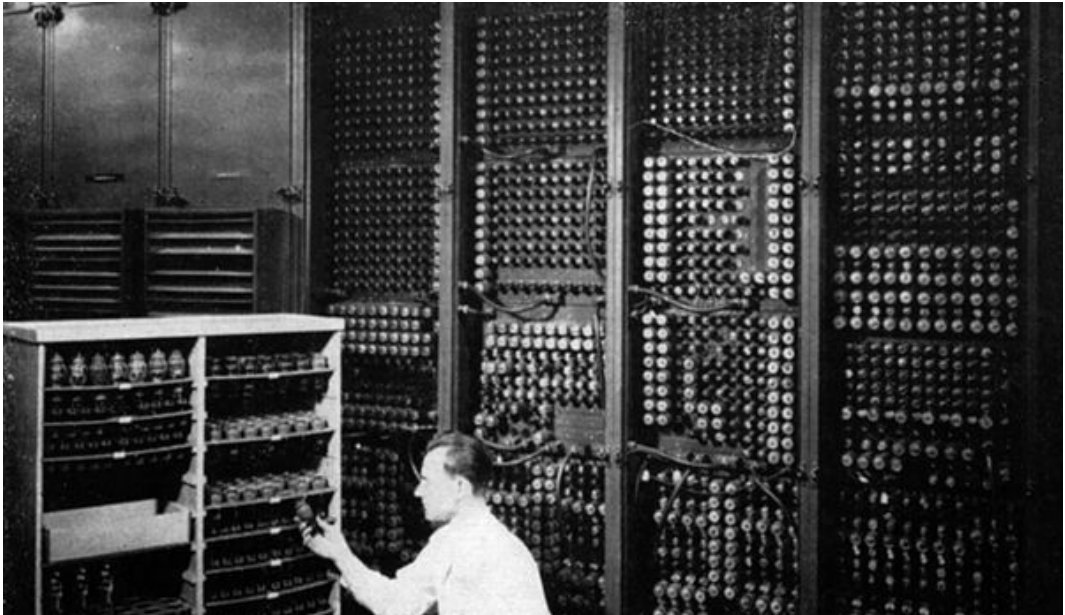
Spetember 2025

Hardware Trends

The GPU Architecture

The OpenCL Programming Model

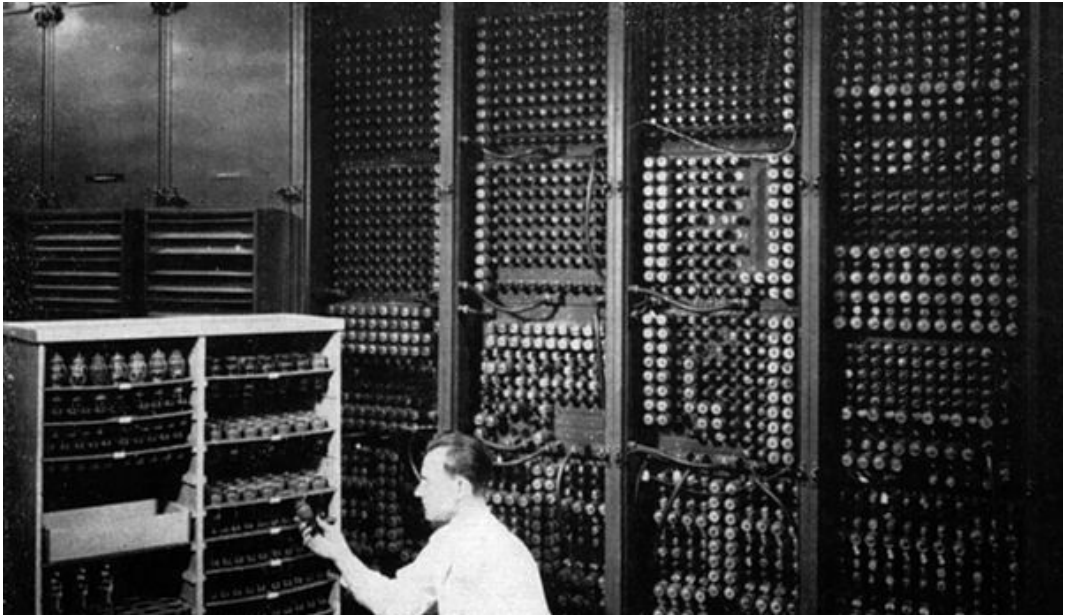
The first computers were not this



But this

And if you had a larger problem

But then they started looking like this



Then this

Then this

Then this

Then this

Then this

Then, from around 2005

Then, from around 2005

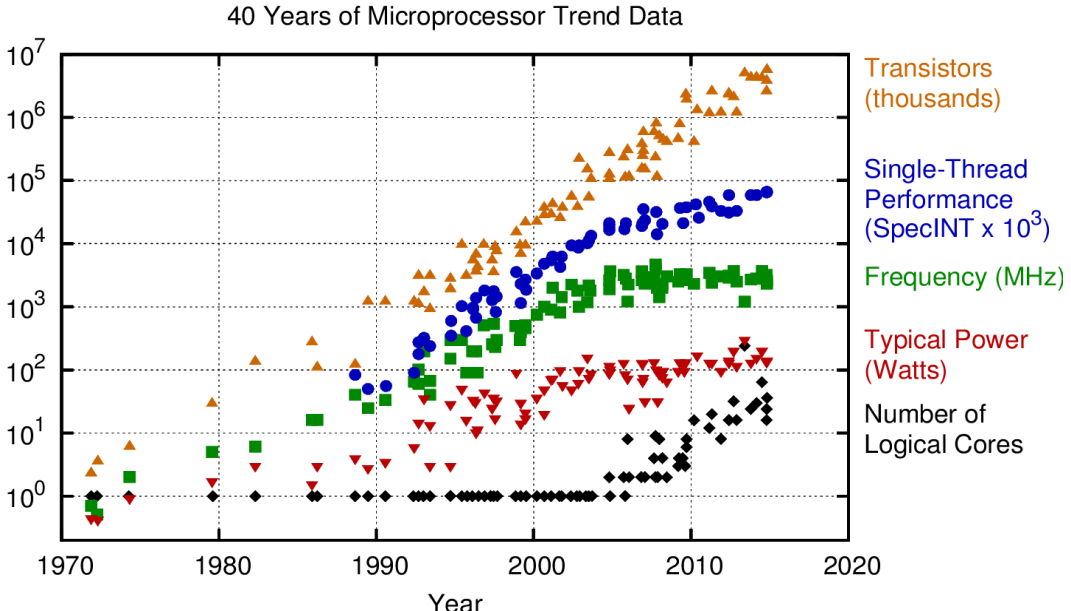
Then, from around 2005

Improvements in *sequential performance* stalled, although computers still got smaller and faster.

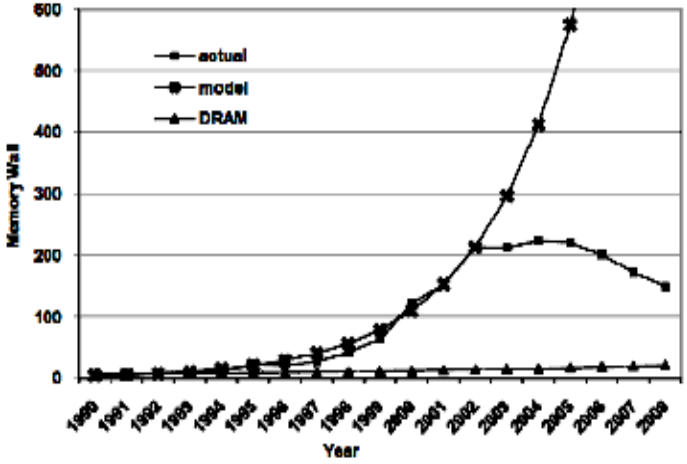
What Changed?

- *Power complexity* $P_{dynamic} \sim Freq^3$, preventing us from increasing processor frequency.
- *Memory wall*, ever-increasing performance gap between processor and memory (which means that *memory* becomes bottleneck, not processor speed).

CPU progress



The Memory Wall



Memory Wall = processor cycles/memory cycles

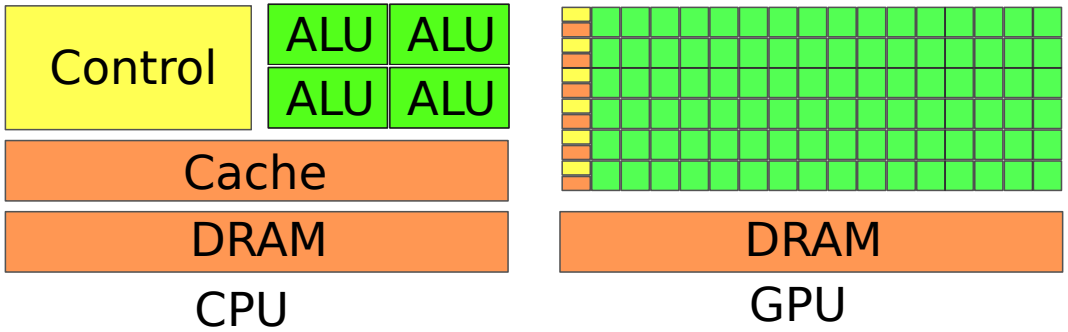
This is why GPUs are useful

The design of GPUs directly attacks these two problems.

- **Frequency scaling** becomes less of an issue because we can instead use thousands of (slower) cores.
- The **memory wall** is partially circumvented by using faster and smaller memory, but mostly by *latency hiding*. With tens of thousands of threads, we can probably find something else to do while some threads are waiting for memory!

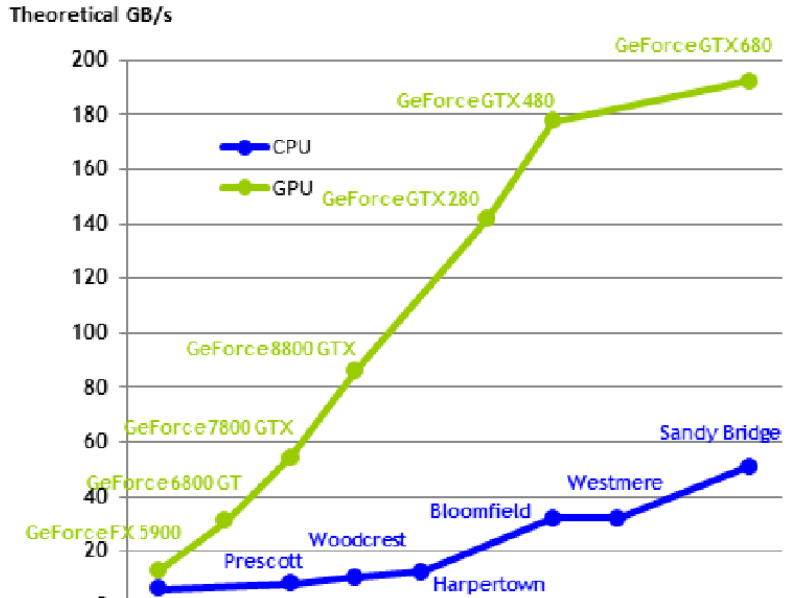
Ultimately, GPUs do *throughput processing*, and operations have (relatively) high latency.

CPUs compared to CPUs

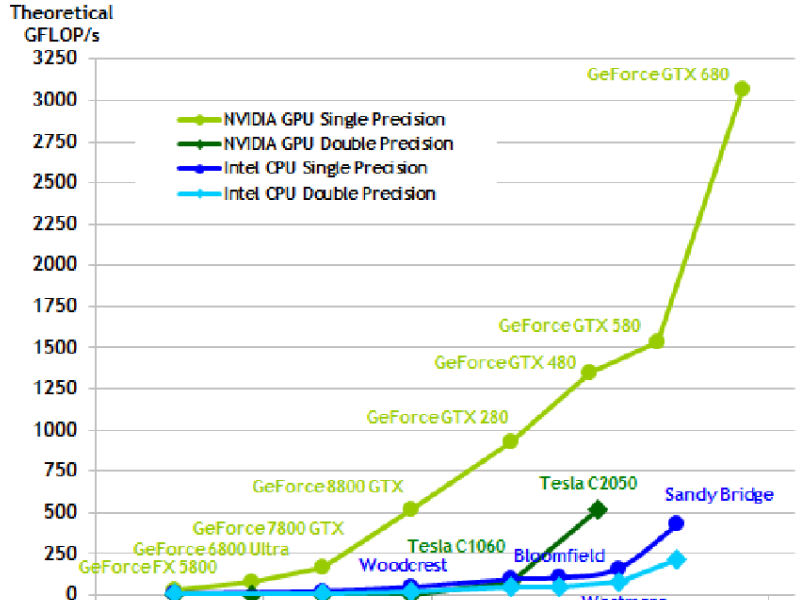


- GPUs have *thousands* of simple cores and taking full advantage of their compute power requires *tens of thousands* of threads.
- GPU threads are very *restricted* in what they can do: no stack, no allocation, limited control flow, etc.
- Potential *very high performance* and *lower power usage* compared to CPUs, but programming them is *hard*.

GPUs and Memory



GPUs and GFLOPS



Hardware Trends

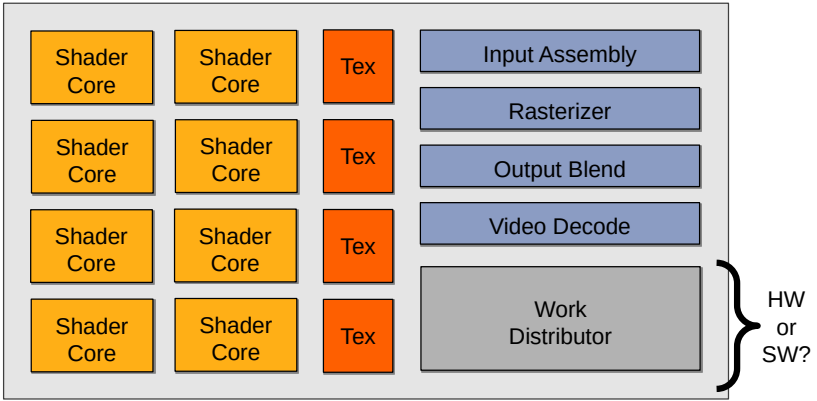
The GPU Architecture

The OpenCL Programming Model

The following slides are taken from the presentation *Introduction to GPU Architecture* by Ofer Rosenberg of AMD.

What's in a GPU?

A GPU is a heterogeneous chip multi-processor (highly tuned for graphics)



A diffuse reflectance shader

```
sampler mySamp;
Texture2D<float3> myTex;
float3 lightDir;

float4 diffuseShader(float3 norm, float2 uv)
{
  float3 kd;
  kd = myTex.Sample(mySamp, uv);
  kd *= clamp( dot(lightDir, norm), 0.0, 1.0);
  return float4(kd, 1.0);
}
```

Shader programming model:

Fragments are processed independently,
but there is no explicit parallel programming

Compile shader

```
sampler mySamp;
Texture2D<float3> myTex;
float3 lightDir;

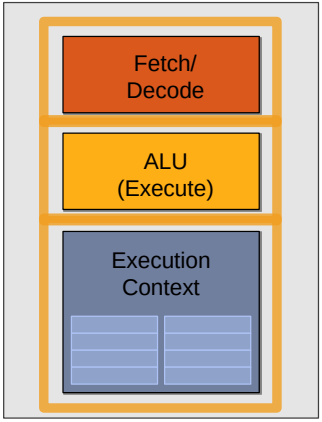
float4 diffuseShader(float3 norm, float2 uv)
{
  float3 kd;
  kd = myTex.Sample(mySamp, uv);
  kd *= clamp( dot(lightDir, norm), 0.0, 1.0);
  return float4(kd, 1.0);
}
```

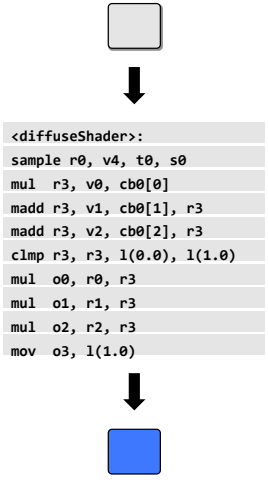
1 unshaded fragment input record

```
<diffuseShader>:
sample r0, v4, t0, s0
mul  r3, v0, cb0[0]
madd r3, v1, cb0[1], r3
madd r3, v2, cb0[2], r3
clmp r3, r3, l(0.0), l(1.0)
mul  o0, r0, r3
mul  o1, r1, r3
mul  o2, r2, r3
mov  o3, l(1.0)
```

1 shaded fragment output record

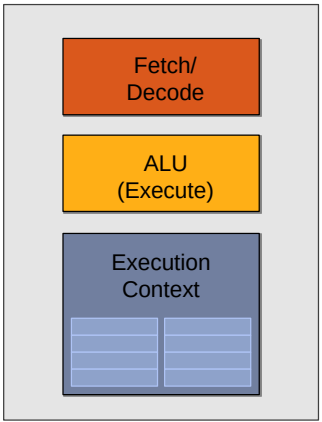
Execute shader

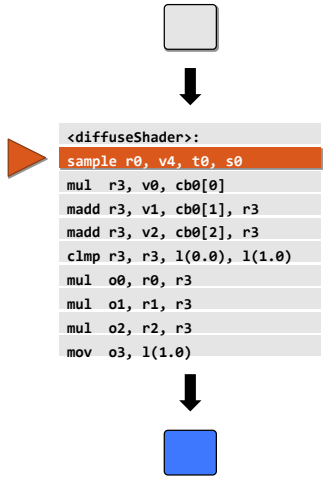




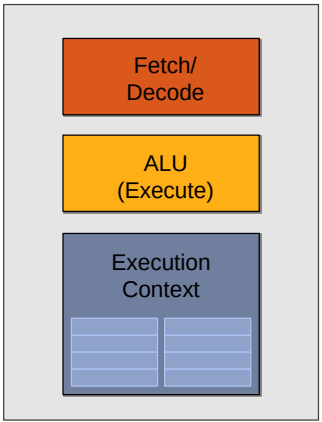
```
<diffuseShader>:
sample r0, v4, t0, s0
mul  r3, v0, cb0[0]
madd r3, v1, cb0[1], r3
madd r3, v2, cb0[2], r3
clmp r3, r3, l(0.0), l(1.0)
mul  o0, r0, r3
mul  o1, r1, r3
mul  o2, r2, r3
mov  o3, l(1.0)
```

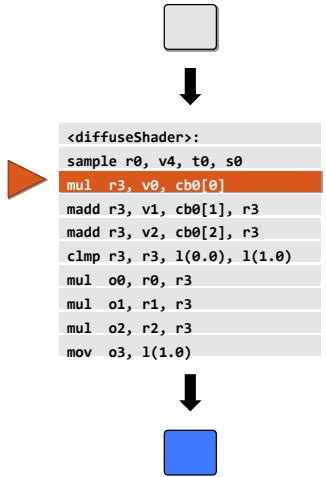
Execute shader



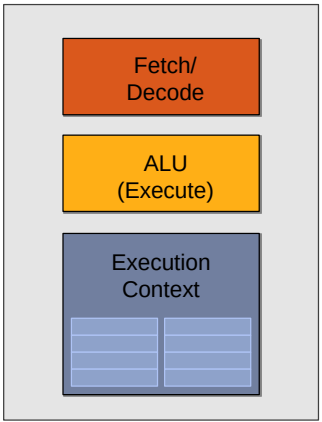


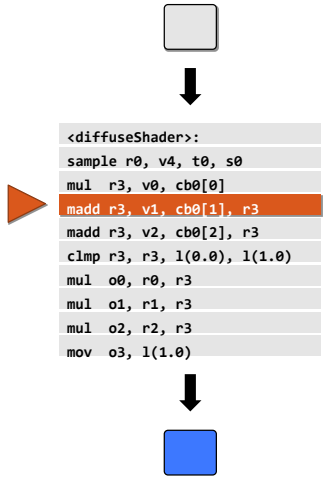
Execute shader



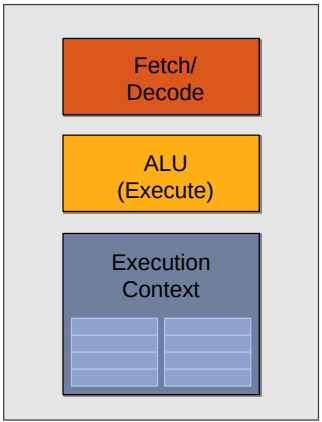


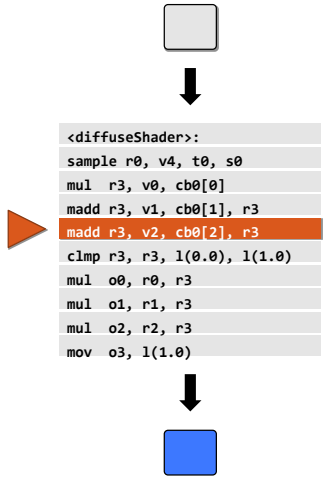
Execute shader



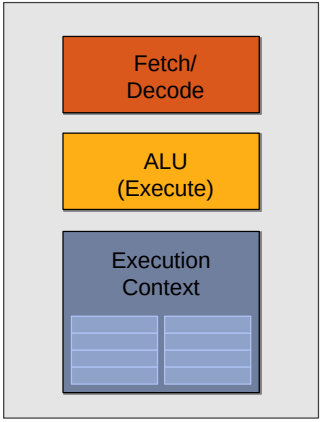


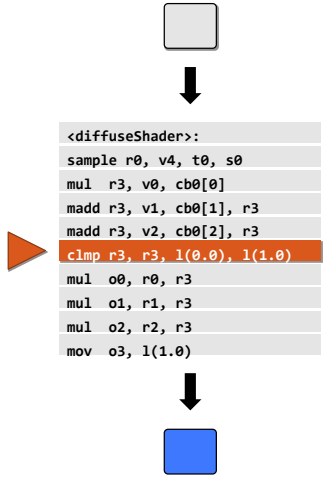
Execute shader



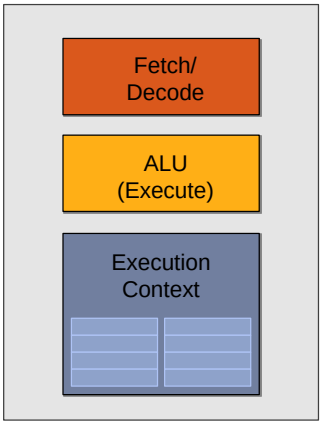


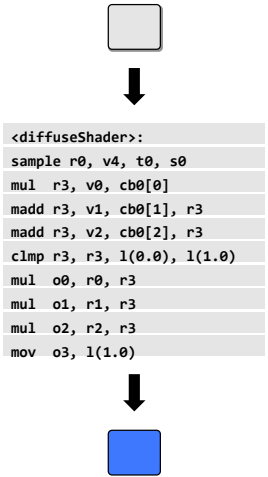
Execute shader



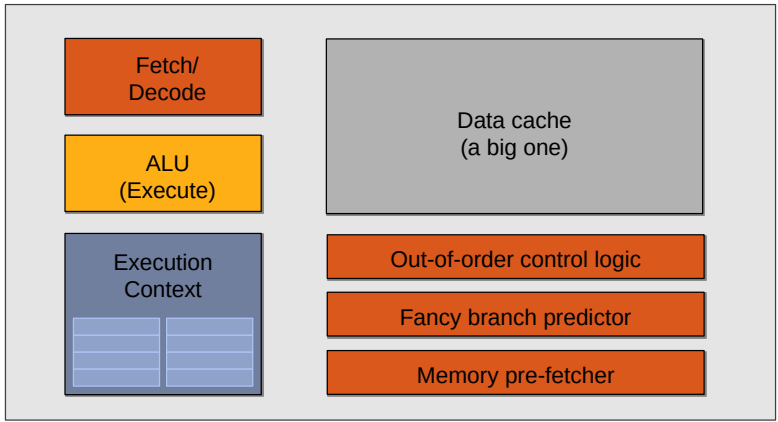


Execute shader

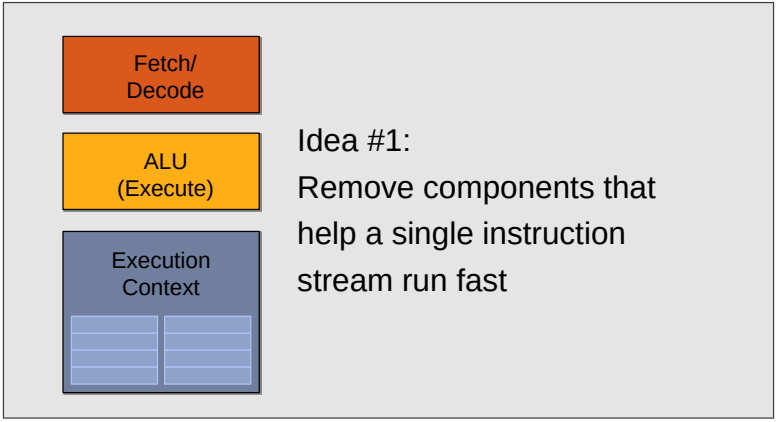




“CPU-style” cores



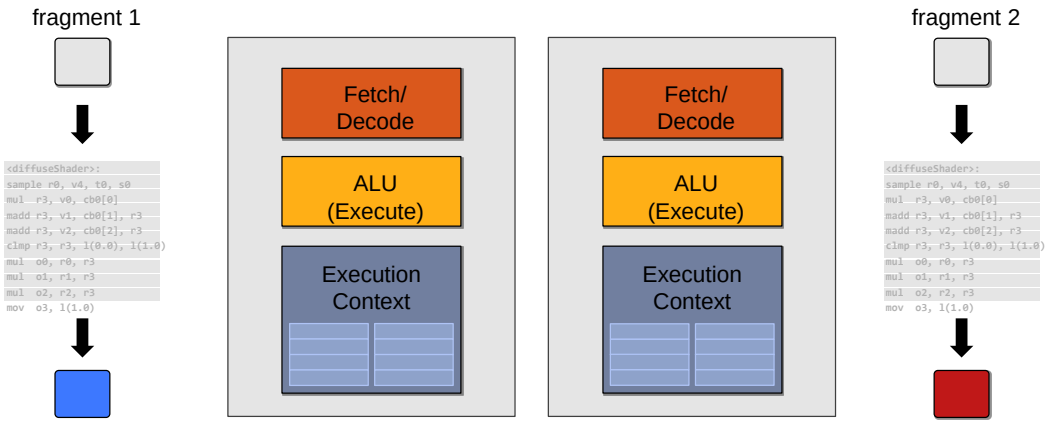
Slimming down



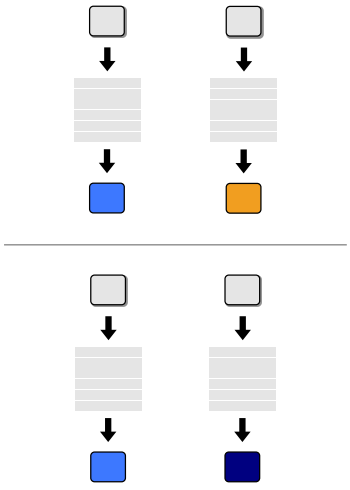
Idea #1:

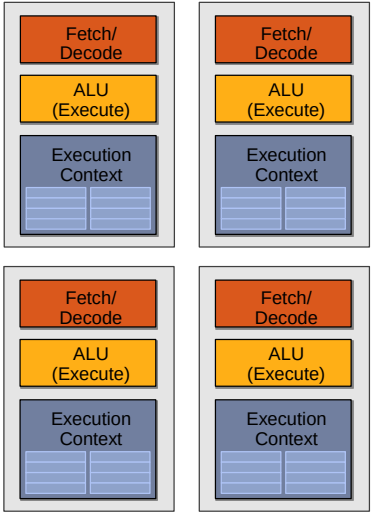
Remove components that help a single instruction stream run fast

Two cores (two fragments in parallel)

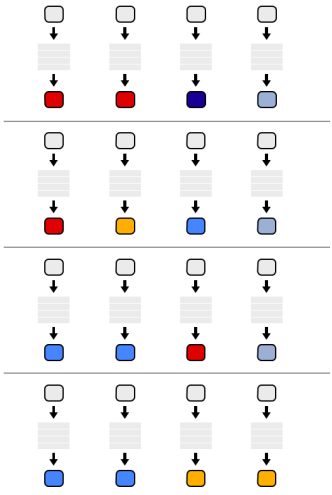


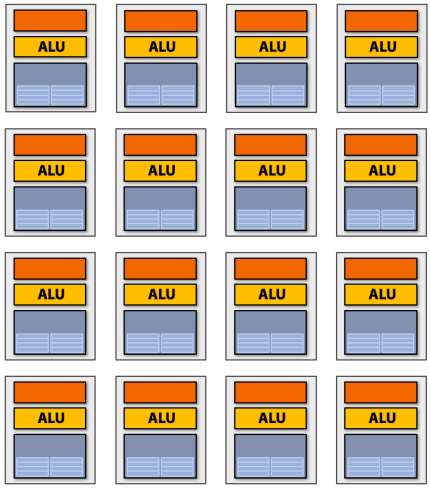
Four cores (four fragments in parallel)





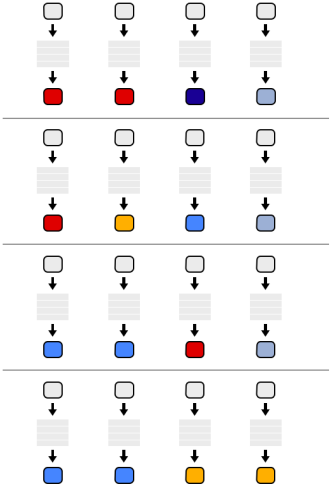
Sixteen cores (sixteen fragments in parallel)





16 cores = 16 simultaneous instruction streams

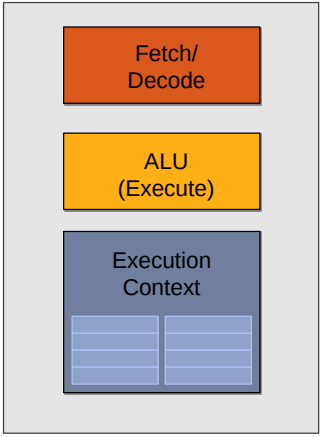
Instruction stream sharing



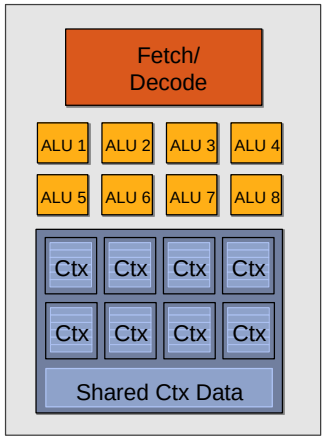
But ... many fragments should be able to share an instruction stream!

```
<diffuseShader>:
sample r0, v4, t0, s0
mul  r3, v0, cb0[0]
madd r3, v1, cb0[1], r3
madd r3, v2, cb0[2], r3
clmp r3, r3, l(0.0), l(1.0)
mul  o0, r0, r3
mul  o1, r1, r3
mul  o2, r2, r3
mov  o3, l(1.0)
```

Recall: simple processing core



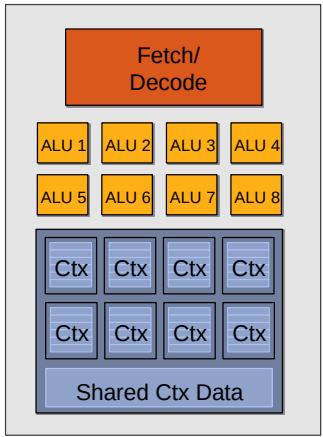
Add ALUs



Idea #2:
Amortize cost/complexity of managing an instruction stream across many ALUs

SIMD processing

Modifying the shader

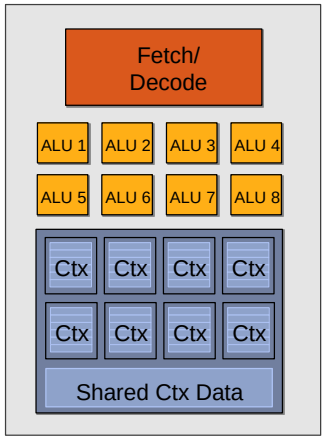



```
<diffuseShader>:
sample r0, v4, t0, s0
mul  r3, v0, cb0[0]
madd r3, v1, cb0[1], r3
madd r3, v2, cb0[2], r3
clmp r3, r3, l(0.0), l(1.0)
mul  o0, r0, r3
mul  o1, r1, r3
mul  o2, r2, r3
mov  o3, l(1.0)
```

Original compiled shader:

Processes one fragment using scalar ops on scalar registers

Modifying the shader

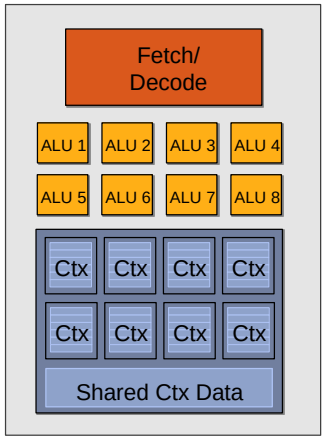


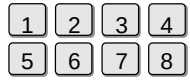

```
<VEC8_diffuseShader>:
VEC8_sample vec_r0, vec_v4, t0, vec_s0
VEC8_mul    vec_r3, vec_v0, cb0[0]
VEC8_madd   vec_r3, vec_v1, cb0[1], vec_r3
VEC8_madd   vec_r3, vec_v2, cb0[2], vec_r3
VEC8_clmp   vec_r3, vec_r3, l(0.0), l(1.0)
VEC8_mul    vec_o0, vec_r0, vec_r3
VEC8_mul    vec_o1, vec_r1, vec_r3
VEC8_mul    vec_o2, vec_r2, vec_r3
VEC8_mov    o3, l(1.0)
```

New compiled shader:

Processes eight fragments using vector ops on vector registers

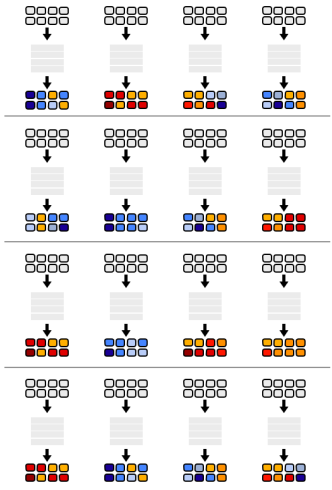
Modifying the shader



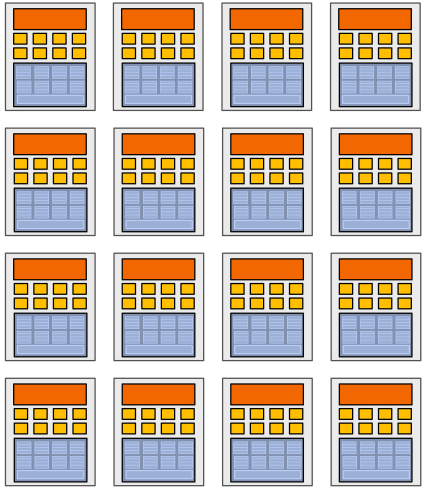



```
<VEC8_diffuseShader>:
VEC8_sample vec_r0, vec_v4, t0, vec_s0
VEC8_mul  vec_r3, vec_v0, cb0[0]
VEC8_madd vec_r3, vec_v1, cb0[1], vec_r3
VEC8_madd vec_r3, vec_v2, cb0[2], vec_r3
VEC8_clmp vec_r3, vec_r3, l(0.0), l(1.0)
VEC8_mul  vec_o0, vec_r0, vec_r3
VEC8_mul  vec_o1, vec_r1, vec_r3
VEC8_mul  vec_o2, vec_r2, vec_r3
VEC8_mov  o3, l(1.0)
```

128 fragments in parallel

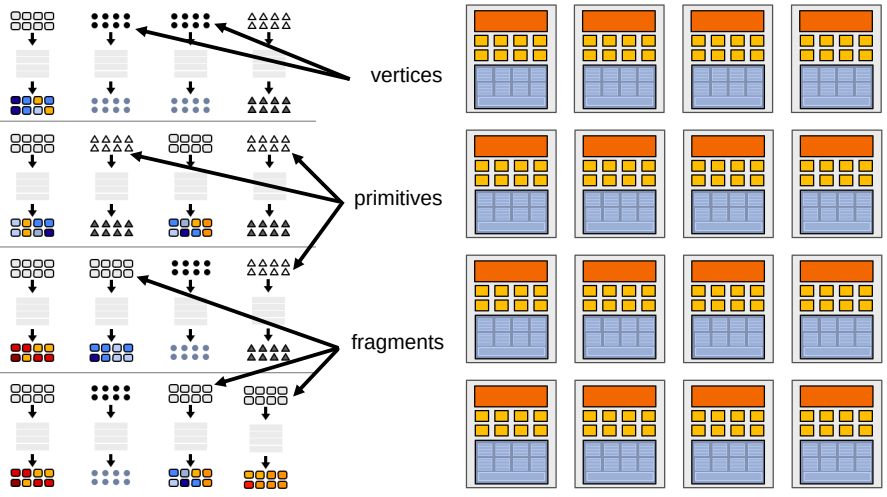


16 cores = 128 ALUs

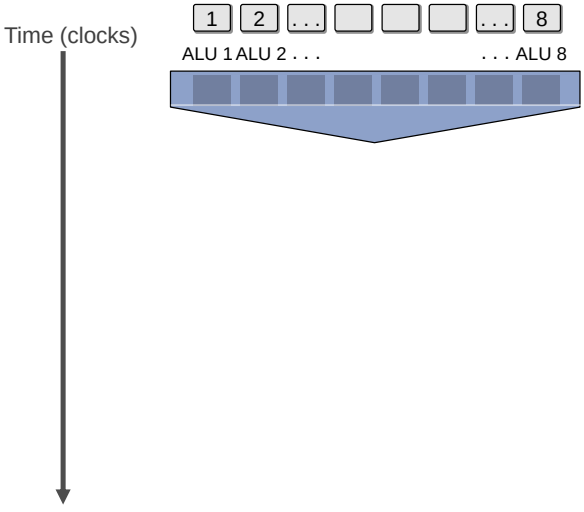


, 16 simultaneous instruction streams

128 [vertices/fragments primitives OpenCL work items] in parallel



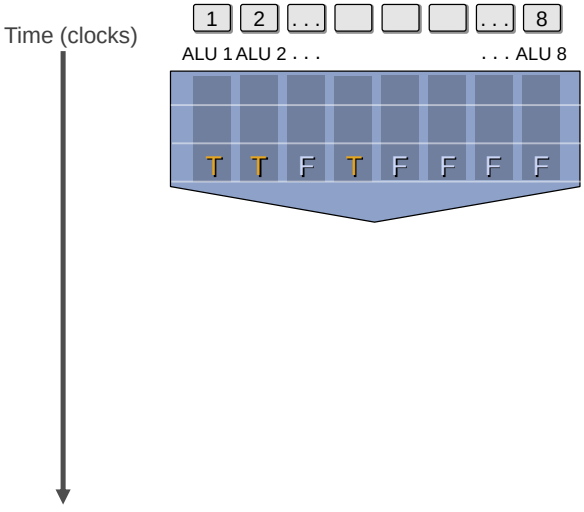
But what about branches?




```
<unconditional
 shader code>

if (x > 0) {
    y = pow(x, exp);
    y *= Ks;
    refl = y + Ka;
} else {
    x = 0;
    refl = Ka;
}
<resume unconditional
 shader code>
```

But what about branches?

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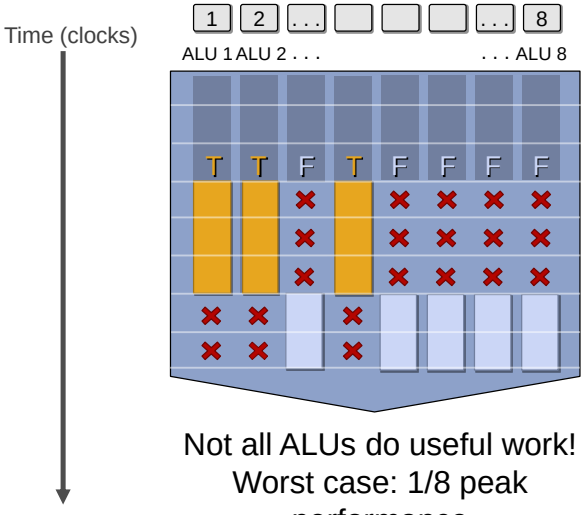
But what about branches?

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if (x > 0) {
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} else {
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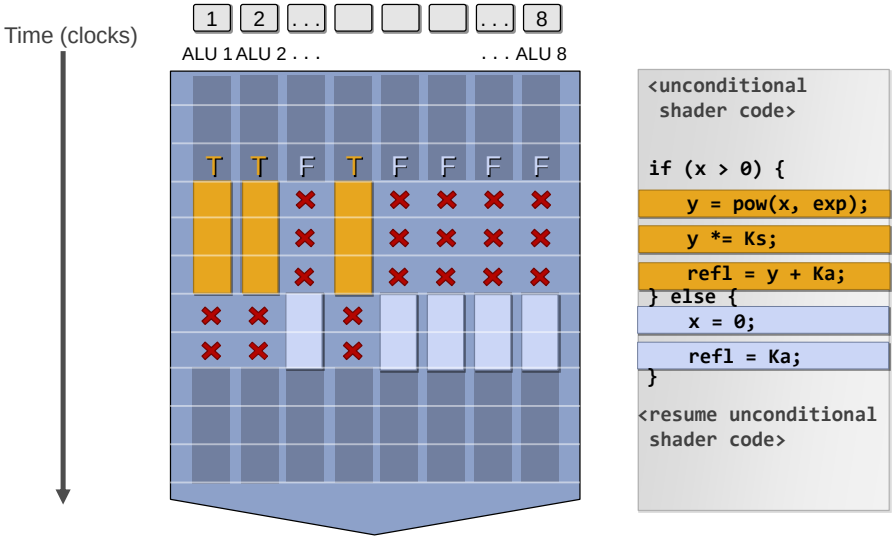
    refl = Ka;
}

<resume unconditional
 shader code>
```

Not all ALUs do useful work!
Worst case: 1/8 peak performance

But what about branches?



Clarification

SIMD processing does not imply SIMD instructions

- Option 1: explicit vector instructions
  - x86 SSE, AVX, Intel Larrabee
- Option 2: scalar instructions, implicit HW vectorization
  - HW determines instruction stream sharing across ALUs (amount of sharing hidden from software)
  - NVIDIA GeForce (“SIMT” warps), ATI Radeon architectures (“wavefronts”)

In practice: 16 to 64 fragments share an instruction stream.

Stalls!

Stalls occur when a core cannot run the next instruction because of a dependency on a previous operation.

Texture access latency = 100’s to 1000’s of cycles

We’ve removed the fancy caches and logic that helps avoid stalls.

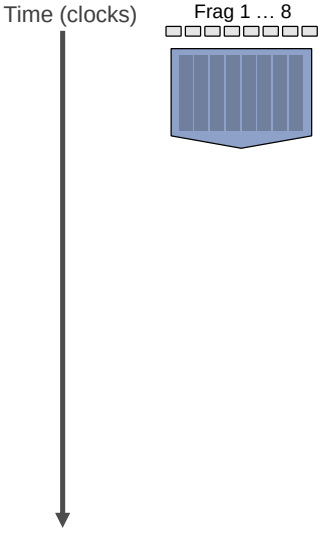
But we have LOTS of independent fragments.

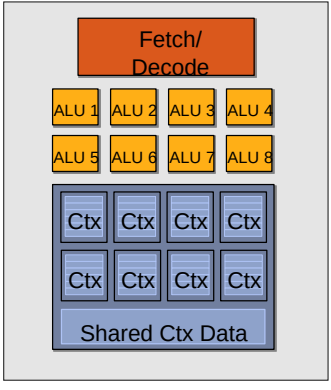
Idea #3:

Interleave processing of many fragments on a single core to avoid stalls caused by high latency operations.

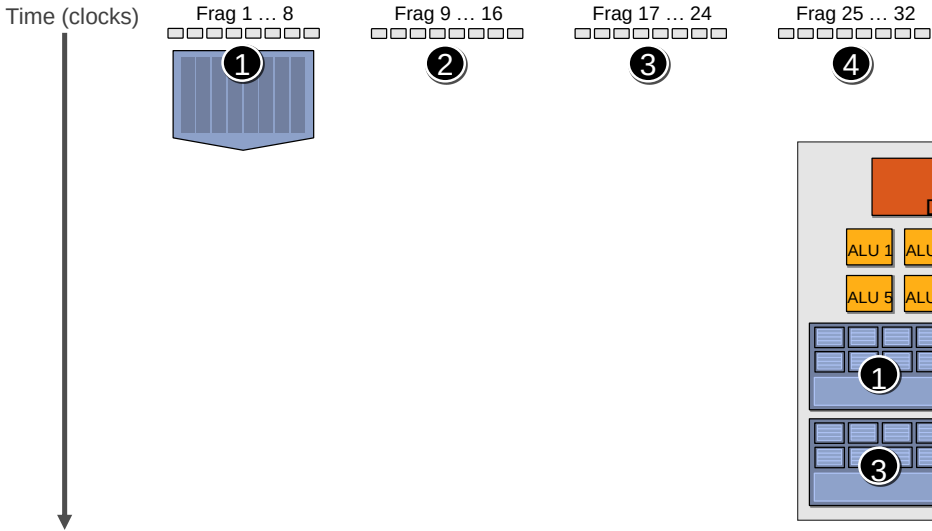
Hiding shader stalls

Time (clocks)

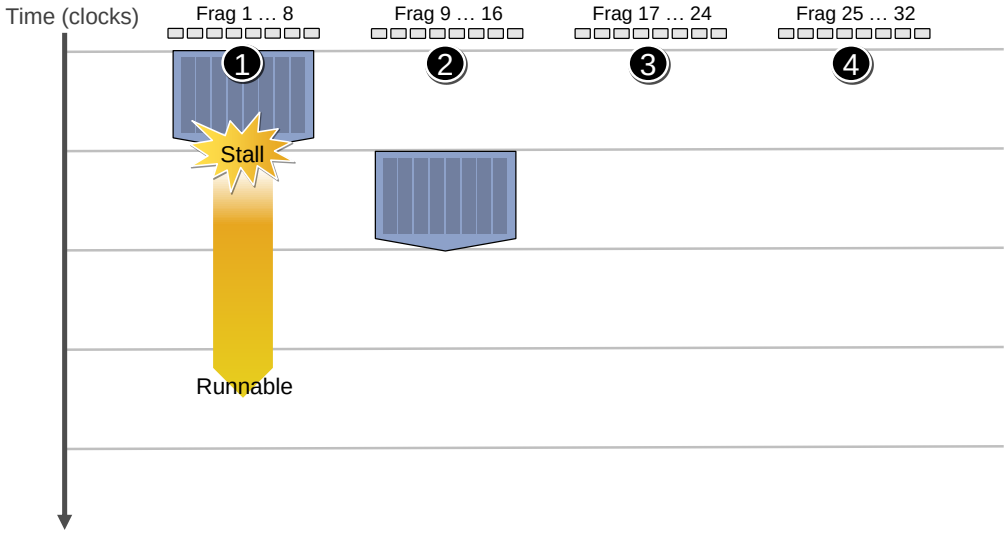




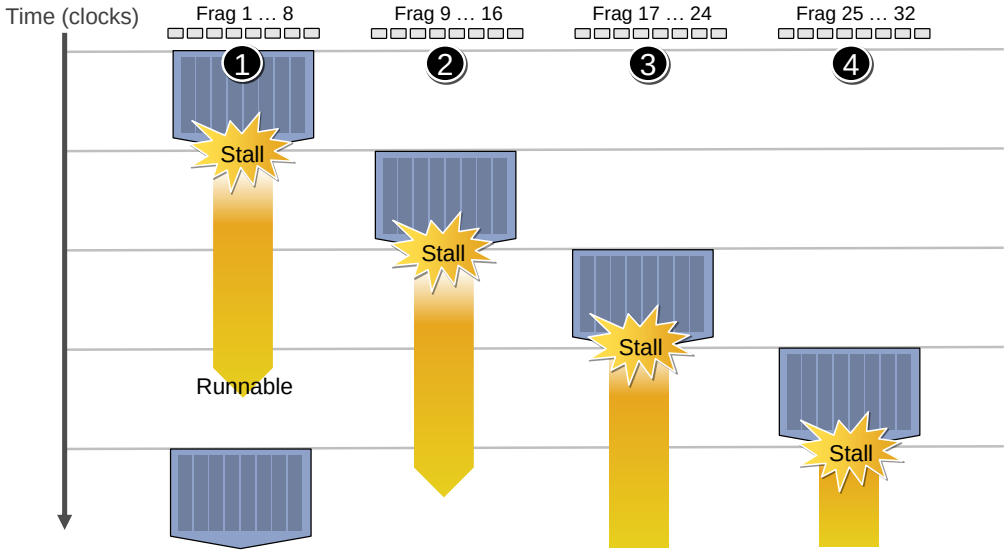
Hiding shader stalls



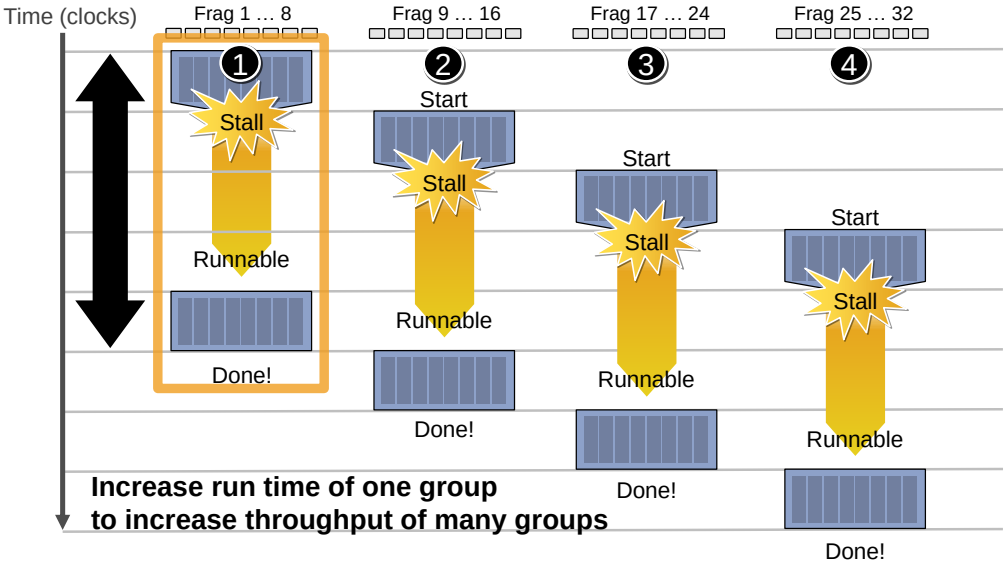
Hiding shader stalls



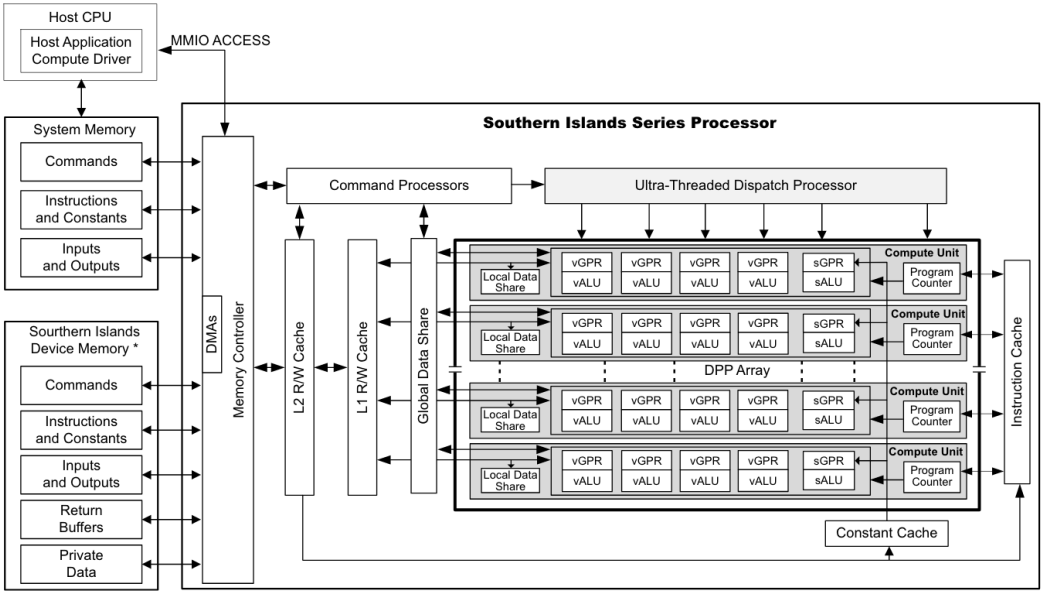
Hiding shader stalls



Throughput!

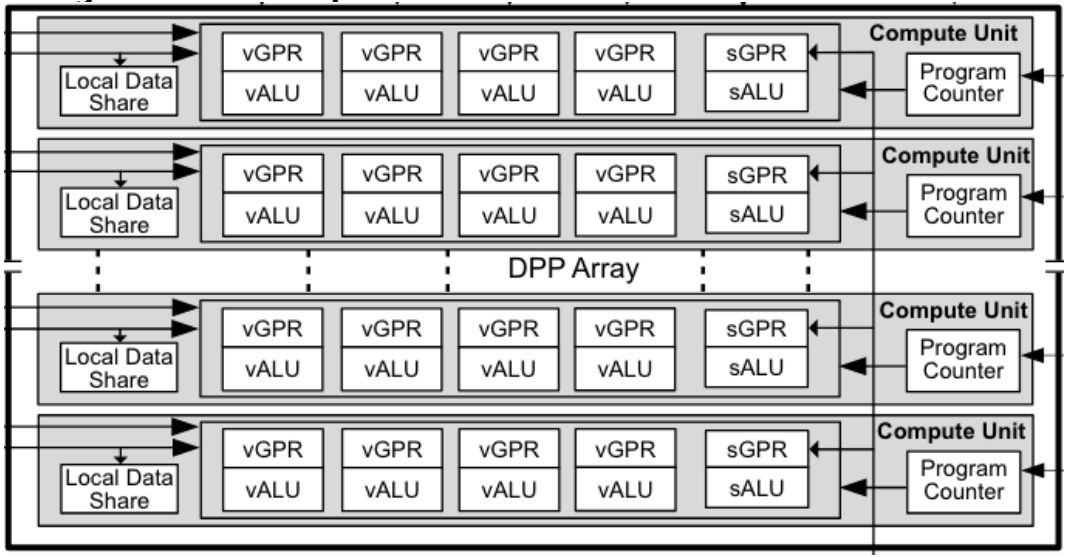

Time (clocks)
Frag 1 … 8
Frag 9 … 16
Frag 17 … 24
Frag 25 … 32
1
2
3
4
Stall
Start
Stall
Start
Stall
Start
Stall
Runnable
Runnable
Runnable
Runnable
Done!
Done!
Done!
Done!
Increase run time of one group
to increase throughput of many groups

The GPU we will be using: Radeon HD 7800[1]



[1] https://developer.amd.com/wordpress/media/2012/12/AMD_Southern

Zooming in on the Compute Units



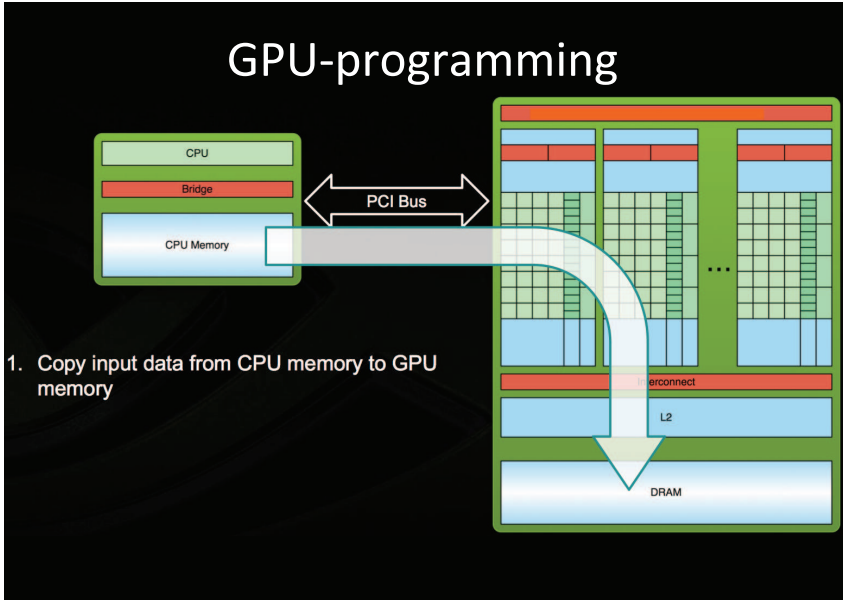
- Each vector-ALU executes a *wavefront* of 64 work-items over four clock cycles.
- Many wavefronts in flight at once to hide latency

Hardware Trends

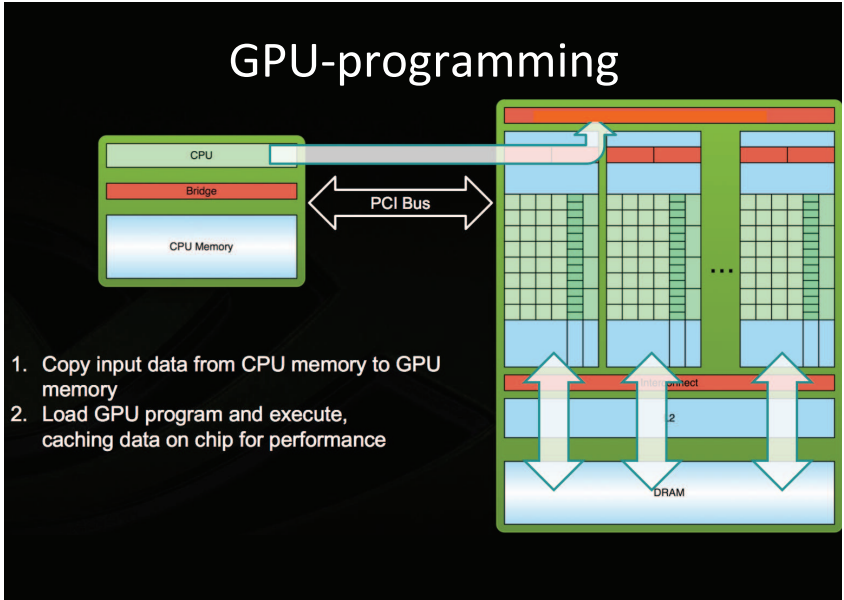
The GPU Architecture

The OpenCL Programming Model

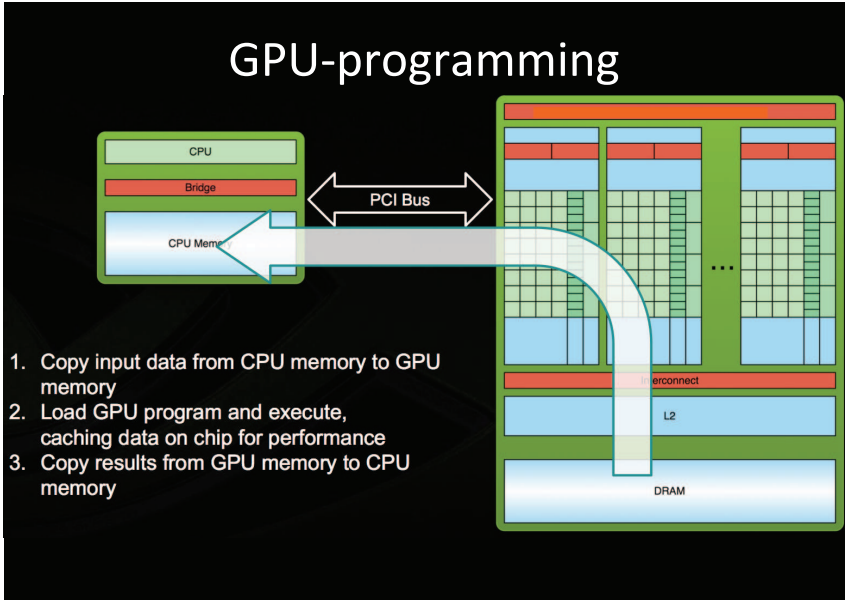
GPU programming



GPU programming



GPU programming



OpenCL (CUDA) is an SIMT model

Single Instruction Multiple Threads means we provide a *sequential function* that is executed in parallel by multiple threads ("work items" in OpenCL).