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Notes on the

Dynkin Classification

31st March to 4th April 2025

BOCHUM SPRING SCHOOL

Dynkin Classification

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Abstract

These are my notes from the spring school “The Dynkin Classification”. This spring school was run from 31st March to 4th April 2025, taking place at Ruhr Universität Bochum in Germany as part of the “Combinatorial Synergies” program. I am grateful for funding recieved from “DFG priority program 2458 Combinatorial Synergies” and “EPSRC AGQ CDT (EP/Y035232/1)” which made this trip possible.

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Chapter 0

Dynkin Classification

0.1 On the Ubiquity of Dynkin Classifications

The Dynkin classification arises in many areas of mathematics, from representation theory to algebraic geometry, and from combinatorics to string theory and spin structures, and many more. Of these we will only touch on the first three.

The structure of these notes follows that of the original spring school. Each day was focused on a different topic, and each day refers to a chapter here. Some of these days focused on the classification more than others, but all were about objects for which the Dynkin classification plays an important role.

Exactly which Dynkin diagrams are allowed in a given classification (or which of the related Coxeter diagrams appears for the first topic) depends on the situation at hand, but in all cases we get at least the A, D, and E types, leading to the so-called ADE-classification.

The precise reason why Dynkin diagrams, particularly the ADE-types, appear so frequently is somewhat of a mystery. Often, but not always, at least not in an obvious way, it's because there is some root system involved, which is a set of vectors in \mathbb{R}^n subject to some conditions (\cdot), the structure of which may be used to generate more complicated objects. Each root system has a corresponding bilinear form, and it is the properties of this inner product which usually restrict which root systems appear in a given classification. The Dynkin diagrams then are just a neat way of encapsulating the minimum amount of information required to define a root system. This is done by first picking a special basis for the root system (a set of so-called simple roots) and then encoding the relative positions of these basis vectors in a graph. The vertices of the graph correspond to elements of this basis, and the number of edges is controlled by the value of the bilinear form when evaluated at the corresponding basis vectors. The direction of any directed edges is chosen such that the arrow points towards the shorter vector. One may then argue that the omnipresence of the Dynkin classification is simply because many problems come down to, or can be phrased as, arranging some points in space subject to some relation between the position vectors relative angles and lengths.

0.2 Coxeter and Dynkin Diagrams

Throughout the first day the following Dynkin diagrams were on the board for reference. We start with our main focus, the (finite type) Dynkin diagrams. We'll then list the (finite type) Coxeter diagrams, which are similar but distinct. Finally, we'll list the affine Coxeter diagrams.

Dynkin Diagrams

Here are the **(finite type) Dynkin diagrams**. The subscript in each case is the number of nodes, which is also known as the **rank**. Dotted edges just mean a chain of nodes connected by single edges. When there are multiple edges there is an arrow, which indicates the direction of all of those edges. These finite-type Dynkin diagrams correspond to root systems which generate a positive definite bilinear form, giving a genuine inner product.

$$A_n = \bullet \text{---} \bullet \text{---} \bullet \text{---} \cdots \bullet \text{---} \bullet \text{---} \bullet \quad (0.2.1)$$

$$B_n = \bullet \text{---} \bullet \text{---} \bullet \text{---} \cdots \bullet \xrightarrow{\quad} \bullet \quad (0.2.2)$$

$$C_n = \bullet \text{---} \bullet \text{---} \bullet \text{---} \cdots \bullet \xleftarrow{\quad} \bullet \quad (0.2.3)$$

$$D_n = \bullet \text{---} \bullet \text{---} \bullet \text{---} \cdots \bullet \begin{array}{l} \nearrow \bullet \\ \searrow \bullet \end{array} \quad (0.2.4)$$

$$E_6 = \begin{array}{c} \bullet \\ | \\ \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \end{array} \quad (0.2.5)$$

$$E_7 = \begin{array}{c} \bullet \\ | \\ \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \end{array} \quad (0.2.6)$$

$$E_8 = \begin{array}{c} \bullet \\ | \\ \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \text{---} \bullet \end{array} \quad (0.2.7)$$

$$G_2 = \bullet \xrightarrow{\quad} \bullet \quad (0.2.8)$$

$$F_4 = \bullet \xrightarrow{\quad} \bullet \text{---} \bullet \text{---} \bullet \quad (0.2.9)$$

Note that some of these are the same, we have the following graph isomorphisms:

- $A_1 \cong B_1 \cong C_1 \cong D_1$, all of which are just a single node;
- $B_2 \cong C_2$, both of which are two nodes connected by a double directed edge;
- $A_3 \cong D_3$, both of which are three nodes connected in a row by single edges;

Note that the sensible definition of D_2 is two disconnected nodes, just the ones appearing in the far right of the picture above. This gives us the graph isomorphism $D_2 \cong A_1 \sqcup A_1$.

There are two ways around these isomorphisms, one is to restrict the valid indices as follows:

- for B_n we require $n \geq 2$;
- for C_n we require $n \geq 3$;
- for D_n we require $n \geq 4$.

The other is to just allow these, but to keep track of whether we think of one of these graphs as being part of, say the A or D series. I prefer this second approach, since it will give us interesting isomorphisms between the objects we classify. These so-called exceptional isomorphisms are best known on the level of Lie algebras, where they correspond to

- $\mathfrak{sl}_2 \cong \mathfrak{so}_3 \cong \mathfrak{sp}_2$;
- $\mathfrak{so}_5 \cong \mathfrak{sp}_4$;
- $\mathfrak{sl}_4 \cong \mathfrak{so}_6$;
- $\mathfrak{sl}_2 \oplus \mathfrak{sl}_2 \cong \mathfrak{so}_2$.

Coxeter Diagrams

The **Coxeter diagrams** are similar to Dynkin diagrams, and the two are often conflated. Coxeter diagrams are undirected graphs, with edge labellings. The rule is that any unlabelled edges are labelled 3, and any two nodes not connected by an edge actually have an undrawn edge labelled 2. Every Dynkin diagram has, ignoring ordering, a corresponding Coxeter diagram, in which the directed edges are replaced by labelled edges with labels greater than 3. There are also some Coxeter diagrams which don't have corresponding Dynkin diagrams. We will see why these distinctions are made later, for now we just list the Coxeter diagrams. The Coxeter diagrams are as below, they have the same labelling as the corresponding Dynkin diagrams, further confusing the two.

$$A_n = \quad \blacksquare - \blacksquare - \blacksquare - \cdots - \blacksquare - \blacksquare - \blacksquare \quad (0.2.10)$$

$$B_n = C_n = \quad \blacksquare - \blacksquare - \blacksquare - \cdots - \blacksquare - \overset{4}{\blacksquare} - \blacksquare \quad (0.2.11)$$

$$D_n = \quad \begin{array}{c} \blacksquare \\ | \\ \blacksquare - \blacksquare - \blacksquare - \cdots - \blacksquare - \blacksquare \\ \diagup \quad \diagdown \\ \blacksquare \quad \blacksquare \end{array} \quad (0.2.12)$$

$$E_6 = \quad \begin{array}{c} \blacksquare \\ | \\ \blacksquare - \blacksquare - \blacksquare - \blacksquare - \blacksquare \\ | \\ \blacksquare \end{array} \quad (0.2.13)$$

$$E_7 = \quad \begin{array}{c} \blacksquare \\ | \\ \blacksquare - \blacksquare - \blacksquare - \blacksquare - \blacksquare - \blacksquare \\ | \\ \blacksquare \end{array} \quad (0.2.14)$$

$$E_8 = \begin{array}{c} \blacksquare \\ | \\ \blacksquare - \blacksquare - \blacksquare - \blacksquare - \blacksquare - \blacksquare - \blacksquare - \blacksquare \end{array} \quad (0.2.15)$$

$$G_2 = \begin{array}{c} 6 \\ | \\ \blacksquare - \blacksquare \end{array} \quad (0.2.16)$$

$$F_4 = \begin{array}{c} 4 \\ | \\ \blacksquare - \blacksquare - \blacksquare - \blacksquare \end{array} \quad (0.2.17)$$

$$H_2 = \begin{array}{c} 5 \\ | \\ \blacksquare - \blacksquare \end{array} \quad (0.2.18)$$

$$H_3 = \begin{array}{c} 5 \\ | \\ \blacksquare - \blacksquare - \blacksquare \end{array} \quad (0.2.19)$$

$$H_4 = \begin{array}{c} 5 \\ | \\ \blacksquare - \blacksquare - \blacksquare - \blacksquare \end{array} \quad (0.2.20)$$

$$I_2(m) = \begin{array}{c} m \\ | \\ \blacksquare - \blacksquare \end{array} \quad (0.2.21)$$

As with the Dynkin diagrams there are certain graph isomorphisms here for specific values of indices. In addition to the isomorphisms of the Dynkin diagrams we also have the following:

- $G_2 \cong I_2(6)$;
- $H_2 \cong I_2(5)$.

Affine Coxeter Diagrams

The final type of diagram we'll need are the affine Coxeter diagrams. We won't be particularly interested in the structures that these classify, they tend to be infinite versions of the finite versions the finite types generate, but they are needed to prove the classification of Coxeter groups. These affine types are formed from the finite types by adding another node. There are generally many ways to do this, but most end up giving isomorphic graphs. In fact, for every outer automorphism of the original graph there is a unique (up to isomorphism) way to add an extra node. We will give only the simplest here, which are the ones corresponding to the identity automorphism. The other types are called "twisted" affine Coxeter graphs. There are two common notations for the affine Coxeter graph, they are \tilde{A}_n and $A_n^{(1)}$. The advantage of the second is that the superscript can index the corresponding outer automorphism, since we're not interested in the twisted case the tilde will be sufficient for our purposes.

Here are all the untwisted affine Coxeter graphs. The new node is the one in colour. Note that the index is now one fewer than the number of nodes.

$$\tilde{A}_1 = \begin{array}{c} \blacksquare - \textcolor{red}{\blacksquare} \end{array} \quad (0.2.22)$$

$$\tilde{A}_n = \begin{array}{c} \textcolor{red}{\blacksquare} \\ / \quad \backslash \\ \blacksquare - \blacksquare - \blacksquare - \cdots - \blacksquare - \blacksquare \end{array} \quad (0.2.23)$$

(0.2.24)

(0.2.25)

(0.2.26)

(0.2.27)

(0.2.28)

(0.2.29)

$$(0.2.30)$$

(0.2.31)

First note that \tilde{A}_1 is “weird”, it’s sort of a degenerate case in which adding the extra node as we would for \tilde{A}_n gives us two nodes with edges between them and this 2-cycle somehow explodes into an infinite loop. Note also that \tilde{B}_n and \tilde{C}_n are different Coxeter diagrams, even though B_n and C_n are the same Coxeter diagram.

Part I

Coxeter Groups

Chapter 1

Reflection Groups

The recommended reference for this topic is [1]. Throughout this part of the notes V will always be a real Euclidean space. That is, V is equipped with some positive definite symmetric bilinear form, $\langle -, - \rangle: V \times V \rightarrow \mathbb{R}$.

1.1 Definition

Definition 1.1.1 — Reflection Fix some $\alpha \in V$. We say that $s \in GL(V)$ is a **reflection** along $\alpha \in V \setminus 0$ if $s(\alpha) = -\alpha$ and the hyperplane orthogonal to α ,

$$H_\alpha := (\mathbb{R}\alpha)^\perp = \{\lambda \in V \mid \langle \alpha, \lambda \rangle = 0\}, \quad (1.1.2)$$

is fixed pointwise, so $s(\lambda) = \lambda$ for $\lambda \in H_\alpha$, in other words s acts as the identity on H_α , so $s|_{H_\alpha} = \text{id}_{H_\alpha}$.

Lemma 1.1.3 The map defined by

$$s_\alpha(\lambda) = \lambda - 2 \frac{\langle \alpha, \lambda \rangle}{\langle \alpha, \alpha \rangle} \alpha \quad (1.1.4)$$

is a reflection in H_α .

Proof. First note that s_α is clearly linear. Then we have

$$s_\alpha(\alpha) = \alpha - 2 \frac{\langle \alpha, \alpha \rangle}{\langle \alpha, \alpha \rangle} \alpha = -\alpha \quad (1.1.5)$$

and if $\lambda \in H_\alpha$, so $\langle \alpha, \lambda \rangle = 0$, we simply have

$$s_\alpha(\lambda) = \lambda - 2 \frac{\langle \alpha, \lambda \rangle}{\langle \alpha, \lambda \rangle} \alpha = \lambda. \quad (1.1.6)$$

Finally, s_α is its own inverse:

$$s_\alpha(s_\alpha(\lambda)) = s_\alpha\left(\lambda - 2\frac{\langle\alpha, \lambda\rangle}{\langle\alpha, \alpha\rangle}\alpha\right) \quad (1.1.7)$$

$$= s_\alpha(\lambda) - 2\frac{\langle\alpha, \lambda\rangle}{\langle\alpha, \alpha\rangle}s_\alpha(\alpha) \quad (1.1.8)$$

$$= \lambda - 2\frac{\langle\alpha, \lambda\rangle}{\langle\alpha, \alpha\rangle}\alpha + 2\frac{\langle\alpha, \lambda\rangle}{\langle\alpha, \alpha\rangle}\alpha \quad (1.1.9)$$

$$= \lambda. \quad \square$$

Notice that by the definition of H_α we can always decompose our space as $V = H_\alpha \oplus \mathbb{R}\alpha$, so the information about what happens to α and anything orthogonal to it is enough to fully determine s_α . Further, note that any nonzero scalar multiple of α also determines the same reflection. All reflections may be specified in this way, and we'll simply write s_α for the reflection in the hyperplane orthogonal to α .

Lemma 1.1.10 Reflections are orthogonal transformations.

Proof. To show that $s_\alpha \in O(V)$ we need to show that $\langle s_\alpha(\lambda), s_\alpha(\mu) \rangle = \langle \lambda, \mu \rangle$ for all $\lambda, \mu \in V$. We can do this with the explicit formula above:

$$\langle s_\alpha(\lambda), s_\alpha(\mu) \rangle = \left\langle \lambda - \frac{\langle\alpha, \lambda\rangle}{\langle\alpha, \alpha\rangle}\alpha, \mu - \frac{\langle\alpha, \mu\rangle}{\langle\alpha, \alpha\rangle}\alpha \right\rangle \quad (1.1.11)$$

$$= \langle \lambda, \mu \rangle - 2\frac{\langle\alpha, \lambda\rangle}{\langle\alpha, \alpha\rangle}\langle\alpha, \mu\rangle + 2\frac{\langle\alpha, \mu\rangle}{\langle\alpha, \alpha\rangle}\langle\lambda, \alpha\rangle \quad (1.1.12)$$

$$+ 4\frac{\langle\alpha, \lambda\rangle\langle\alpha, \mu\rangle}{\langle\alpha, \alpha\rangle^2}\langle\alpha, \alpha\rangle$$

$$= \langle \lambda, \mu \rangle. \quad \square$$

Definition 1.1.13 — Reflection Group A **reflection group** is a group, W , is a subgroup of $O(V)$ generated by reflections.

We will almost entirely restrict our study to finite reflection groups, where W is a finite group. In this part of the notes W is a finite reflection group unless stated otherwise.

We will see that finite reflection groups are classified by finite type Dynkin diagrams. We will write $W(R)$ for the finite reflection group corresponding to the Dynkin type R under this correspondence.

1.2 Examples

Dihedral Groups: $I_2(m)$

Consider the dihedral group, $W(I_2(m)) = \text{Dih}_m$, of symmetries of the regular m -gon. Note that this group has order $2m$. We will illustrate things with Dih_5 ,

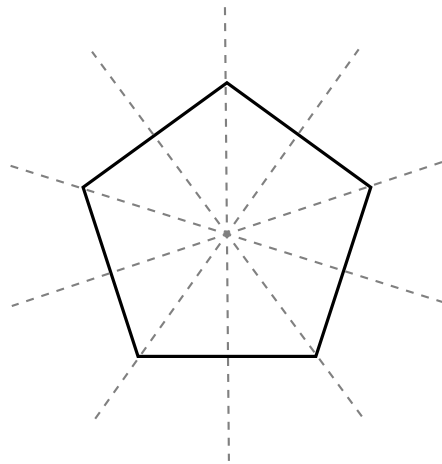


Figure 1.1: The regular pentagon. A reflection in any of the five dashed lines is a symmetry.

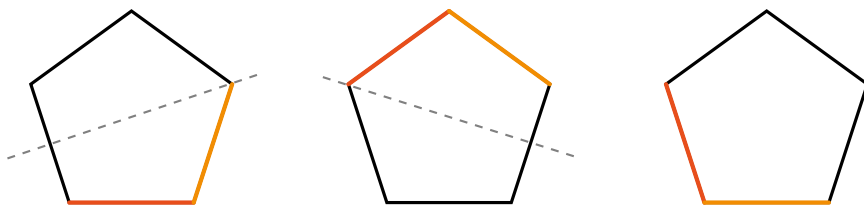


Figure 1.2: Successive reflections generate a rotation by $2\pi/5$. The coloured edges are just there to help visualise the reflections occurring.

the symmetries of the regular pentagon, as seen in [Figure 1.1](#). With $V = \mathbb{R}^2$ we can naturally extend the symmetries of the regular m -gon (centred on the origin) to the entire plane. A rotation by $2\pi/m$ can be generated by two reflections. For example, [Figure 1.2](#) shows two reflections generating a rotation of $2\pi/5$ by applying two reflections, and from this we can generate all rotations.

Platonic Solids: A_3 , B_3 , H_3

The platonic solids, the simplex, cube, octahedron, dodecahedron, and icosahedron, are the strictest generalisation of regular shapes to three dimensions. Like the regular m -gons the symmetry groups of the platonic solids are all reflection groups upon extending the action from the polyhedron to all of \mathbb{R}^3 .

The 3-simplex, also known as the tetrahedron, has symmetry group $W(A_3) = \text{Sym}_4$, which simply acts by permuting the faces.

The cube and its dual, the octahedron, have $\text{Sym}_4 \times \mathbb{Z}_2$. The \mathbb{Z}_2 corresponds to a reflection in some chosen plane through three non-adjacent vertices, inverting the cube. The action of Sym_4 is to permute the diagonals, which are the lines that are invariant under the different choices of inversion above.

However, it turns out that if we want to generalise in the future it is actually better to think of the symmetries of the cube as being $W(B_3) = \text{Sym}_3 \ltimes \mathbb{Z}_2^3$. Here

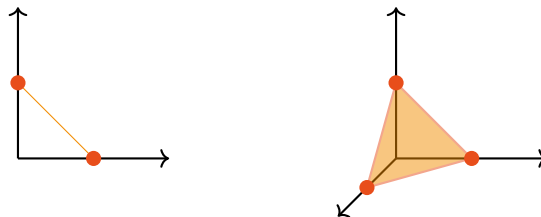


Figure 1.3: Simplices: The 1-simplex is just an interval and the 2-simplex is a triangle.

Sym_3 acts on \mathbb{R}^3 by permuting the standard basis, and each copy of \mathbb{Z}_2 acts as a reflection in one of the coordinate planes. Note that in this description the cube is thought of as $[-1, 1]^3$, centred on the origin.

The dodecahedron has symmetry group $W(H_3)$, of order 120. It's a bit complicated to specify how it acts here, and it also doesn't generalise so we won't go into details.

Simplices and Permutations: A_n

The n -simplex is the convex hull of the n -basis vectors of \mathbb{R}^{n+1} . This is perhaps a confusing definition, but hopefully the pictures of the 1-, and 2-simplices in Figure 1.3 help. The 1-simplex is simply a line segment, the 2-simplex is an equilateral triangle, and the 3-simplex is the regular tetrahedron. The symmetries of the n -simplex are simply given by permuting n -cells (faces of the tetrahedron, edges of the triangle, the two end points of the line), and there are always $n + 1$ n -cells. The symmetry group of the n -simplex is thus $W(A_n) = \text{Sym}_{n+1}$.

Recall that Sym_{n+1} is generated by transpositions of neighbouring elements, $(i \ i + 1)$ for $1 \leq i \leq n$. When realised as a reflection group the transposition $(i \ i + 1)$ corresponds to the reflection $s_{e_i - e_{i+1}}$.

Signed Permutations: B_n

The n -cube is $[-1, 1]^n \subset \mathbb{R}^n$. The symmetry group of the n -cube is the group of *signed* permutations of n -letters.