

The written document focuses on explaining how my program is executed to solve  $A\mathbf{x} = \mathbf{b}$  using Gaussian Elimination on page 67 of the textbook.

Since we are given a fact that the matrix  $A$  can be factorized to a lower-triangular and an upper-triangular matrix  $L$  and  $U$ , which is  $A = LU$ . The first step is to decompose  $A$ .

Step 1: Decompose  $A$  to  $L$  and  $U$

(1) Using gaussian Elimination by obtaining elementary matrices(elimMat.m)

```
function [M_k,L_k] = elimMat(A, k)
```

```
[n,~] = size(A);
```

```
M_k = eye(n);
```

```
% Initialize M_k as an identity matrix
```

```
M_k_entry = -A(k+1:n, k)/A(k,k);
```

```
%Vectorization
```

```
M_k(k+1:n, k) = M_k_entry;
```

```
%Obtain the entries in the lower triangular of M_k
```

```
L_k = inv(M_k);
```

```
%Obtain L_k according to the formula that  $L_k = \text{inv}(M_k)$ 
```

```
end
```

Every time when computing an elementary matrix, multiply it to matrix  $A$  on the left hand side until we derive an upper-triangular matrix by doing so. We know that

$$A = LU = LM_{n-1}M_{n-2} \cdots M_1A$$

Thus, we can derive that  $L = (M_{n-1}M_{n-2} \cdots M_1)^{-1} = M_1^{-1}M_2^{-1} \cdots M_n^{-1}$

(2) Therefore, we can simultaneously update the temporary matrix  $L$  by inverting the elementary matrix we have obtained.

```

function [L, U] = myLU(A)

[n,~] = size(A);
L = eye(n);
U = A;

for k = 1:n-1
    %The reason for the termination of for loop is (n-1) is
    that
    %we start every time from (k+1)th column to get M_k entry,
    thus k+1 <=n
    %
    [M_k, L_k] = elimMat(U, k);
    U = M_k * U;
    L = L * L_k;
end
end

```

After LU decomposition, we get a new logically equivalent equation  $LU\mathbf{x} = \mathbf{b}$ .

Step 2: Let  $\mathbf{y} = U\mathbf{x} \Rightarrow L\mathbf{y} = \mathbf{b}$

Forward Substitution to solve  $\mathbf{y}$

Since zeros occupy the top right portion of L, we can solve  $y_1$  first from the top, which is forward substitution.

```

function y = fwdSubst(L, b, k)
%Forward substitution
[m,n]=size(L);
if ~exist('k') % If first call no k param given, but k=1
    k=1;
end

y=b(k)/L(k,k);
if k < n % Recursion step
    1 = [zeros(k,1);L(k+1:m,k)];
    y = [y;fwdSubst(L,b-y*1,k+1)];
end

```

Step 3:  $U\mathbf{x} = \mathbf{y}$  Backward Substitution to solve  $\mathbf{x}$

We have  $U\mathbf{x} = \mathbf{y}$  and  $\mathbf{y}$  has been derived in the previous step. Likewise, since zeros occupy the bottom left portion of  $U$ , we can solve  $x_n$  first from the top, which is forward substitution. We can obtain the equation system below:

$$\begin{aligned} u_{11}x_1 + u_{12}x_2 + \cdots + u_{1n}x_n &= y_1 \\ u_{22}x_2 + \cdots + u_{2n}x_n &= y_2 \\ &\vdots \\ u_{nn}x_n &= y_n \end{aligned}$$

We can derive that  $x_n = \frac{y_n}{u_{nn}}$ , iterating  $x_n$  to equations above, we can obtain that

$$x_{n-1} = \frac{y_{n-1} - u_{(n-1)n}x_n}{u_{(n-1)(n-1)}}, \quad x_{n-2} = \frac{y_{n-2} - (u_{(n-2)(n-1)}x_{n-1} + u_{(n-2)n}x_n)}{u_{(n-2)(n-2)}}$$

Thus, we can induct the general formula for  $x_i$ , which is

$$x_i = \frac{y_i - \sum_{j=i+1}^n u_{ij}x_j}{u_{ii}}$$

**function**  $\mathbf{x} = \text{backSubst}(U, \mathbf{y}, k)$

$\mathbf{x} = \text{zeros}(k, 1);$

**for**  $i=k:-1:1$

$\mathbf{x}(i) = (\mathbf{y}(i) - U(i, i+1:k) * \mathbf{x}(i+1:k)) / U(i, i);$

**end**

**end**

Finally, we can solve  $\mathbf{x}$  after executing the program.