Mark each statement true or false (2 points each, 20 cents)	—、Mark each statement <i>tru</i> e or <i>fals</i> e (2 points each)	
1) A LR(1) parser cannot parse any left-recursive CFG without ambiguity. (F) 2) LL(1) grammar cannot be left-recursive. (T)	1. A grammar is ambiguous if it has two different derivations or two different	1. Strings over the alphabet {a,b,c} where the length of the string is even can
3) Given a legal string of tokens for a CFG, there must be a unique parsing tree to derivate the string. (F)		be generated by regular expression:←
4) The language L={a^b^ n>=1} can't be generated by any regular expression. (T)	parse trees for a <u>sentence.</u>	[A]. ((a b c)(a b c))*↩
5) There is only one parse tree for the string of an unambiguous grammar. (T)	2. If a grammar is LR(I), but not LALR(I). There are not shift-reduce conflicts	[B]. (a b c)(aa bb cc)*(a b c)↩
6) If a grammar is LR(1), but not LALR(1). There may be shift-reduce conflicts in its parsing table of	in its parsing table of LALR(1).	[C]. (aa bb cc)*←
LALR(1). (F) 7) Finding the next handle is the man task of a LR parser. (T)	In its parsing table of <u>LALK(</u> 1).	
8) Both DFA and NFA can recognize regular set. (T)	二、Single Choice (2 points each)⊲	[D]. (a b c)*↩
9) The parse tree will completely reflect the derivation steps for a string. (F)	A 1. In the Top-Down Parsing, the action () will never be used.	4
10) Left recursion is commonly used to make operations left associative. (T)	[A] Chiff [D] Matala [O] Carranta [D] Asserts	2. Which of the following is not only found in a stack frame (activation record)?
 Multiple Choice (3 points each, 30 cents) Given the CFG: E'→E, E→E + n n, is(are) the viable prefix(es) of the right sentential form 'n+n'. 	[A] Shift[B] Match [C] Generate [D] Accept	[A] static variables↩
(ACD)	A 2、In the Top-Down Parsing, the action () will never be used. ←	[B] return address∈
[A] n [B] n+ [C] E [D] E+	[A] Shift[B] Match [C] Generate [D] Accept⊲	[C] saved registers↩
2) If a LL(1) grammar contains the rules: $A \rightarrow \alpha_1 \alpha_2$; $B \rightarrow \beta_1 \beta_2$, then the following condition must be satisfied.		[D] parameters⊬
(AD) [A] First(β_1) \cap First(β_2) is empty [B] First (B) \cap Follow (A) is empty	B 3、IF one CFG grammar contains two non-terminals 'A','B' and two terminal	₽ parameters.
[C] First (A) \cap Follow (A) is empty [D] First (α_1) \cap First (α_2) is empty	'a'.'b' where 'A' is the start symbol, then the Follow set of 'A' may be	3. The elements in a LR Parsing table are labeled with the following kinds of
3) Given the production A \rightarrow B α C, we have (B)		5
$ [A] \ Follow \ (C) \subset Follow \ (A), \ First(B) \subset First(A) \\ [B] \ Follow \ (A) \subset Follow \ (C), \ First(B) \subset First(A) $	()	actions except? ←
[C] Follow (C) \subset Follow (A), First(B) [D] Follow (A) \subset Follow (C), First(A) \subset First(B) 4) Given the grammar rules: A' \rightarrow A \rightarrow A (A) A ϵ and the rightmost derivation A'=>A=> (A)(A) => (A)(A) =>	[A] {a, b} [B] {a, b, \$} [C] {a, b, ε } [D] {a, b, B} \leftarrow	[A]. Shift[B]. Goto[C]. Accept[D]. Push
(A)()=>()() is the viable prefix the sentential form (A)(A). (B)		4
[A] A) [B] (A [C] ((A)A)A [D] ()		4. In the production $B\rightarrow \alpha Ay$, Which would not be in the FOLLOW(A) Set? \leftarrow
5) Give the LR(1) item $[A \rightarrow \alpha Bc, a/b]$, we have (C)	- Single Choice (12 cents)	[A]. ε[B]. FIRST(<u>γ)</u> [C]. FOLLOW(<u>B)</u> [D]. \$ <i>←</i>
(A) $\{a,b\}\subset First(A)$ (B) $\{a,b\}\subset First(B)$	1, which of the following language is generated by the context free grammar G:	4
(C) {a,b}⊂Follow(A) (D) {a,b}⊂Follow(B)	$S \rightarrow xSx \mid y$	5. Which grammer person helpsy in the least powerful?
6) Here is a grammar: (C) X→a X→ ε Y→b Y→X	[A] xyx [B] $(xyx)^*$ [C] $x^nyx^n(n>=0)$ [D] x^*yx^*	5. Which grammar parser below is the least powerful?
	2. If the context-free grammar G is not ambiguous, for any sentence generated by	[A]. <u>LL(</u> 0) [B]. <u>LR(</u> 0) [C]. SLR[D]. <u>LR(</u> 1)←
	G, which of the following description is correct?	\leftarrow
$Z \rightarrow c$ $Z \rightarrow XYZ$ $W \rightarrow d$ $W \rightarrow XY$	[A] the parse tree corresponding to the left-most derivation must be the same to that	6. Here is a grammar: ←
Which symbol is not nullable?		
[A]. X [B].Y [C].Z [D].W	of the right-most derivation.	X→a X→ Y→b Y→X ←
7) The parsing method of YACC is (A)	[B] the parse tree corresponding to the left-most derivation may not be the same to	$Z \rightarrow c$ $Z \rightarrow XYZ$ $W \rightarrow d$ $W \rightarrow XY \leftarrow$
[A] LALR(1) [B] LR(1) [C] SLR(1) [D] LL(1)	that of the right-most derivation.	Which symbol is not nullable?↩
8) Which action is not in a LL(1) parsing table? (AD)	[C] the left-most derivation must be the same to the right-most derivation.	[A]. X[B].Y [C].Z [D].W←
[A] Shift [B] Generate [C] Accept [D] Reduce	[D] there will be more than one left-most derivation, but the parse tree is the same.	4
9) The output of the parser is (B)	3. If one CFG grammar is LL(1) and contains the rules: $A\rightarrow\alpha_1 \alpha_2$; $B\rightarrow\beta_1 \beta_2$, then the	7. Which element would not appear in the stack of LR parser?
[A] token [B] syntax tree [C] target code [D]intermediate code	following condition () must be satisfied.	[A <u>].terminal</u> [B].nonterminal [C].\$ [D].state⊖
	[A] First (A) \cap First (B) is empty [B] First(α_1) \cap First(α_2) is empty	4
· · · · · · · · · · · · · · · · · · ·	[C] First (A) ∩Follow (A) is empty [D] First (B) ∩Follow (A) is empty	8. Normally, we do error recovery with methods except? ←
[A] Left factoring [B] First set and follow set	4. LR(1) item $[A \rightarrow \alpha . B \gamma , a]$, follow(B)= {	[A]. add error production[B]. modify the parsing tables ✓
[C] Remove left recursion [D] Remove right recursion		[C]. eliminate the conflict [D]. modify the parsing engine
	[A] a [B] γ [C] $\{\gamma, a\}$ [D] $\{\gamma a\}$	4
	5, the parsing method of YACC is ()	9. Which of the following is commonly found in a stack frame (activation
— Mark each statement true or false (2 points each, 6 cents)	$ [A] LALR(1) \qquad [B] LR(1) \qquad [C] SLR(1) \qquad [D] LL(1) $	record)? ←
The same language token may be generated by many different regular expressions.	6. In the Bottom-Up Parsing, the action () will never be used.	[A] fp pointer←
(T)	[A] Reduce [B] Match [C] Shift [D] Accept	[B] sp pointer⊲
To any regular expression, we can find a context-free grammar defining the same		[C] static variable
language. (T)		[D] global variable
The LL(1) parsing algorithm parses an input string of tokens by tracing out the steps		Many compilers are divided into two largely independent passes: a front
in a rightmost derivation. (F)	—、 Mark each statement true or false (20 cents) ←	end, responsible for analyzing source code, and a back end, responsible for
二、Single Choice (1 points each ,5 cents)	 Scopes of the variables are intercrossed sometimes. 	generating target code. (They may also include a third, intermediate phase,
1. The concept () is not related to the LL(1) parsing method.	2. <u>Yacc</u> can not use ambiguous grammars. ✓	responsible for code improvement.) What is the most important motivation for
[A] Left-factoring [B]. First set and follow set	3. Both DFA and NFA can recognize regular set.	this division into passes?
[C.] Left recursion removal [D]. Shift and reduce	4. A grammar is ambiguous if it has two different derivations or two different	[A] to provide portability of compiler ←
2、Which one below is not a part of a compiler? (这个题目不是太好,就不扣分了)	parse trees for a sentence.	[B] to facilitate division of labor when a compiler is built by a large team of
[A] Symbol table [B] Assembler	■ 5. Grammars with left recursion eliminated are LL(1).	programmers←
[C] Code optimizer [D] Parser	6. All LL(1) grammars can be parsed by LALR parser. ←	[C] to minimize memory requirements on modern machines⊲
3. In the production A \rightarrow B α C, we have		[D] to facilitate debugging the compiler←
[A] Follow (C) \subset Follow (A), First(B) \subset First(A)	7.Code generation depends on detailed information about the target	4
[B] Follow (C) \subset Follow (A), First(A) \subset First(B)	architecture, and doesn't care the characteristics of the source language. S. The best obside of data structure of the sourch late is IASU table. S. The best obside of data structure of the sourch late is IASU table. S. The best obside of data structure of the sourch late is IASU table. S. The best obside of data structure of the sourch late is IASU table. S. The best obside of data structure of the sourch late is IASU table. S. The best obside of data structure of the source language. S. The best obside of data structure of the source language. S. The best obside of data structure of the source language. S. The best obside of data structure of the source language. S. The best obside of data structure of the source language. S. The best obside of data structure of the source language. S. The best obside of the source language. S. The bes	el .
[C] Follow (A) \subset Follow (C), First(B) \subset First(A)	8. The best choice of data structure of the symbol table is HASH table.	Mark each statement true or false (20 cents) ←
[D] Follow (A) \subset Follow (C), First(A) \subset First(B)	9. An L-value can not occur on the right of an assignment statement. ←	1. (F) 2. (F) 3. (T) 4. (F) 5. (F) \leftarrow
4、IF one CFG grammar contains two non-terminals 'A', 'B' and two terminal 'a', 'b', where	10.An intermediate representation is a kind of machine language with	7. (F) 7. (F) 8. (T) 9. (F) 10. (F) □
'A' is the start symbol, then the Follow set of 'A' may be (committing to too much machine-specific detail. ←	
[A] $\{a, b\}$ [B] $\{a, b, \$\}$ [C] $\{a, b, \epsilon\}$ [D] $\{a, b, B\}$	·	四、 Single <u>Choice</u> (15 cents) (4
5、In the Top-Down Parsing, the action () will never be used.		1. <u>A</u> 2. A 3. D 4. A 5. A ← 6. C 7. C 8. C 9_ A 10. A←
[A] Shift [B] Match [C] Generate [D] Accept		0.01.00.0 <u>0 n</u> 10.n

—、Mark each statement true or false ←	1. An LR(1) parser can detect errors earlier than an LR(0) parser.	 Given the following declarations: t1 = array [10] of int;
1、Both DFA and NFA can recognize regular set. ←	判断题 (4.0 分) (难易度:中) A. True	t2 = array [10] of int; if t1 and t2 isequivalent, it follows equivalence.
2. Context free grammar can generate language L = {anbncm n >=0 , m>=0}	B. False	单选题 (4.0 分) (难易度:中)
3. Intermediate Code generation depends on detailed information about the	正确答案: A	A. structural
target architecture, and doesn't care the characteristics of the source	答案解释: 哲无	B. name C. declaration
language. ←	$2. \text{If a grammar is LR} (1), then the LALR (1) parsing table cannot have any reduce-reduce conflicts.} \\$	D. none of the above
4、A display is a data structure that may be used as a alternative to static links	判断題 (4.0 分) (难易度:中)	正确答案: A
for maintaining access to the variables. ←	A. True B. False	省案解释 : 到无
A grammar is ambiguous if it has two different derivations for a sentence.	D. False ∠ 正确答案: B	8. When inheriting a previously computed synthesized attribute during LR parsing, it is suitable to treat the computed synth
6. In grammar, nullable symbol can only derive the empty string.	答案解释 : 智无	attribute as 单选题 (4.0 分) (难易度:中)
7. All LR(0) grammars can be parsed by SLR parser. ←	3. The topological sort of attribute computation of is unique.	A. return value
	判断題 (4.0 分) (难易度:中)	B. passing as parameter
8. There might be shift-reduce, shift-shift and reduce-reduce conflicts during	A. True	C. external data structure
the LR parsing. ←	B. False	D. just leaving it on the value stack 正确答案: C
	正确答案: B 答案解释: 图无	答案解释: 哲无
— Single Choice		A Yacc-generated parser employ method to parse the input token stream.
1. Which of the following string can be defined by the regular expression	 An inherited attribute can be calculated by a postorder traversal of the parse tree. 判断題 (4.0 分) (难易度:中) 	単选題 (4.0 分) (难易度:中)
((<u>blc</u>)*a(<u>blc</u>)*a) (<u>blc</u>)*.↩	A. True	A. top-down
[A] abbcab[B] aaaa [C] abbbbbc [D] bbacc	B. False	B.LL(1)
2、 the output of the scanner is: ←	正确答案: B 答案解释: 哲无	C. SLR(1) D. LALR(1)
[A] token[B] syntax tree	東京中で・ 日ル	正确答案:D
3、Lex is a tool that ()←	All inherited attributes can be changed into synthesized attributes by suitable modification of the grammar, without ch language of the grammar.	hanging the 答案解释: 哲无
[A] break the source file into individual words, or tokens	判断题 (4.0 分) (唯易度·中)	10. The symbol tables will not carry theinformation about the variables.
[B] analyze the phrase structure of the program ←	A. True	单选题 (4.0 分) (难易度:中)
[C] is a lexical analyzer generator ←	B. False	A. data type
[D] is a parser generator ✓	正确答案: A 答案解释: 哲无	B. scope
4. In the following which is(are) not commonly found in a stack frame		C. liveness D. location in memory
(activation record)?←	6. The input of semantic analysis phase is 单选题 (4.0 分) (难易度:中)	正确答案: A
[A] return address[B] static variables <-	A. tokens	
[C] saved registers[D] arguments⊬	B. CFG	>•\$,\$
5. which does the operation pushing the return address in the stack frame?	C. an abstract syntax tree S→	>•a,\$ T→T,S•,)
[A] the caller procedure⊲	D. a parse tree S ·· 正确答案: C	>•(T), \$
[B] the called procedure		T→T, • S,)
[C] the operation system⊲	5. (15 cents) Consider the following augmented grammar.	$(S \rightarrow (T) \ S \rightarrow (T), S \rightarrow (T),$
[D] CPU 4	S' \rightarrow S	S → (•T), \$ T→•T, S,) S → (•T), \$ T→T•, S,)
6. Which action is not in a LR Parsing table? ←	$S \rightarrow a (T)$	T→S,) (a
o. William action is not in a Err Faising table:	$T \rightarrow T, S S$	$S \rightarrow a, 1$ $S \rightarrow (T), 1$ $S \rightarrow a, 1$
[A] Shift[B] Push [C] Accept [D] Reduced 7. In the production B→αAy, <u>Which</u> would not be in the FOLLOW(A) Set? ✓	(1) Construct the LR(1) DFA. (10 cents) (12 cents)	$S \rightarrow (T) \bullet,)$
[A] ϵ [B] FIRST(γ) [C] FOLLOW(B) [D] \$ \leftarrow	(2) Is the given left-recursive CFG LR(1)? (3 cents)	
8. Which element would not appear in the stack of LR parser? ← [A] nonterminal[B] terminal [C] state [D] \$←	(2) is the given lett-recursive of a Er(1): (a certa)	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$
9. Which grammar parser below is the most powerful?		T→ s,) T→ t•, s,)
[A] LL(1) [B] LR(0) [C] SLR [D] LR(1)↩ 10. the parsing method_of_YACC_is()↩		S → (T),) a
$ \begin{array}{llllllllllllllllllllllllllllllllllll$		
一、Mark each statement <i>true</i> or <i>false</i> (16 cents) (每题 2 分) ↩ (エ)(F) (F) (T) (F) (T)(F)↩		
二、Single Choice (10 cents) (每题 1 分) ↔ A&CBBBADDA↔		