

* Queue Reconstruction by Height *

$[7,0], [4,4], [7,1], [5,0], [6,1], [5,2]$

store s.t. # of $\geq \text{curr}$ in left equals to $\text{curr}[1]$

→ sort in descending

$[max, 0]$

add lesser value to left makes no difference

→ for same $\text{curr}[0]$ sort by $\text{curr}[1]$ asc.

$O(n \log n)$
 $O(n)$