## **Review For Final**

**Structure of the Exam.** You will have 2 ½ hours for the exam.

- 1. True/False and Multiple Choice
- 2. Very Short Answer
- 3. Short Answer
- 4. SCI (3 points)

## **Review Points**

- 1. Defintions Be familiar with definitions of the following concepts
  - heap
  - red-black tree
  - vertex cover
  - Hamiltonian graph / Hamiltonian cycle
  - connected graph
  - simple cycle in a graph
  - sparse graph / dense graph
  - tree (as a graph)
  - $\bullet$  P = NP
- 2. Be familiar with the formulas related to red-black tree (discussed in class): If *T* is a red-black tree with *n* nodes, height *h*, and root *r*, then the following hold true for *T*:

$$bh(r) \geq \frac{h}{2}$$

$$n \geq 2^{bh(r)} - 1$$

$$h \leq 2\log(n+1)$$

- 3. Be able to carry out the steps of the insertion algorithm for red-black trees.
- 4. Be able to do HeapSort by hand, clearly indicating steps for Phase I and Phase II.
- 5. Know the brute force algorithm for determining the smallest size of a vertex cover for a graph. Know also the VertexCoverApprox algorithm and be able to explain why its running time is polynomial.
- 6. Know in detail the BFS algorithm and how it can be used for various tasks (spanning tree/forest, shortest path, num components, connected graph, existence of path between two given vertices)
- 7. Be able to do simple reasoning about graphs. Know the following results (you do not need to know the proofs).

- i. A graph *G* is **connected** if for any two vertices *u*, *v* in *G*, there is a path from *u* to *v*.
- ii. **Theorem.** For any graph G, if

$$\epsilon > \binom{\nu-1}{2}$$

then G is connected.

iii. **Theorem.** If G is a tree, then  $\varepsilon = v - 1$ .

Exercise. Suppose G is a graph with  $\nu-1$  edges. Show that the following are equivalent:

- a. G is connected
- b. G is acyclic
- c. G is a tree
- iv. **Theorem**. If  $\varepsilon < v 1$ , then G is not connected.
- v. **Theorem**. If  $\varepsilon \ge v$ , G contains a cycle.
- vi. **Ore's Theorem**. If in a connected graph G with n vertices (n > 2), the sum of the degrees of nonadjacent vertices is always no smaller than n, then G is Hamiltonian.
- 8. Know how to prove that the running time of insertion, deletion and search in a heap is  $O(\log n)$
- 9. Be able to carry out the steps of Kruskal's algorithm
- 10. Know how to do the steps of slow version of Dijkstra.
- 11. Be able to verify that a decision problem belongs to *NP*.
- 12. Be able to establish that a decision problem is NP-complete, given that some other (related) problem is already known to be NP-complete (as in the slides).
- 13. Be familiar with the results about the SubsetSum problem.
- 14. Be able to illustrate with an example the fact that the HamiltonianCycle problem is polynomial reducible to TSP (as in the slides and Lab 12)
- 15. Be familiar with the graph package that was developed in the class. You may be asked to implement a method in the Graph class by devising your own subclass of BreadthFirstSearch.
- 16. Be able to explain why the usual IsPrime algorithm is in reality exponential.
- 17. NOTE: You will not be required to reproduce any proofs