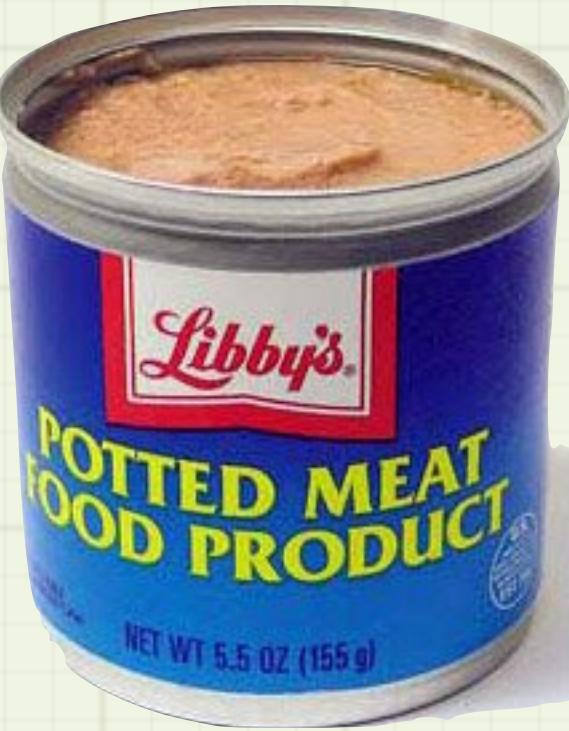


rquery and **rqdatatable**: **R** tools for data manipulation

John Mount
Win-Vector LLC

Outline

- Who I am.
- What is the unmet need for **R** users working with big data?
- Solution context: a potted history of data manipulation concepts.
- The solution: **rquery** and **rqdatatable**.
- Performance comparisons.
- Conclusion.



John Mount

Win-Vector LLC

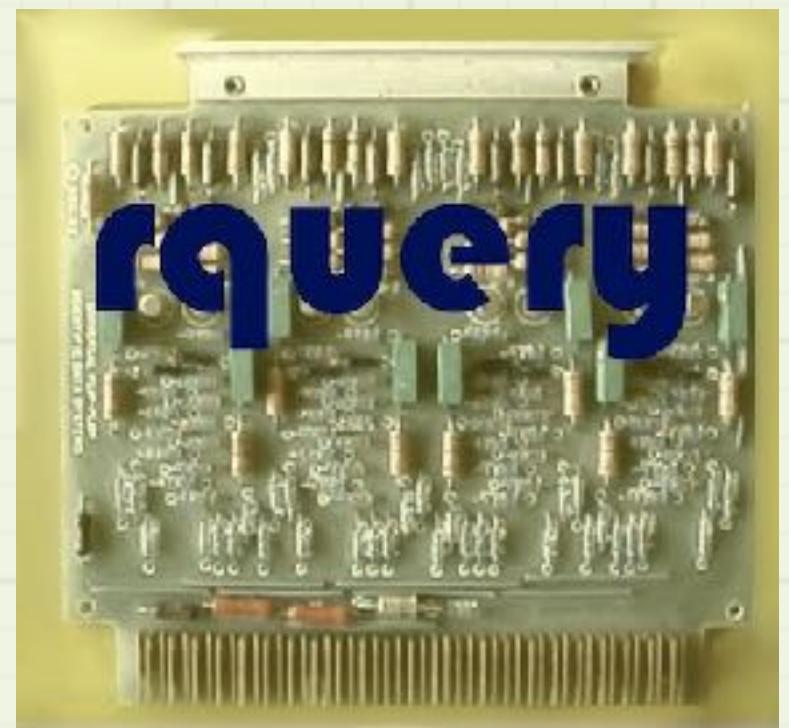
<http://www.win-vector.com/>

- One of the authors of the book *Practical Data Science with R*, Zumel Mount, (Manning 2014).
- We (at Win-Vector LLC) provide **R**, statistics, and data science training (both live and pre-recorded) and consulting.
- One of the authors of **vtreat**: statistically sound preparation of data for predictive modeling.
- Frequent contributor to [the Win-Vector blog](#) and conference speaker.



What I Want to Share

- How to wrangle data using the **rquery** SQL query generator and the **rqdatatable** implementation.



Not going to be able to demonstrate/explain *everything*.

If you want to build a ship, don't drum up the [men people] to gather wood, divide the work and give orders. Instead, teach them to yearn for the vast and endless sea.
Antoine de Saint-Exupery, "The Wisdom of the Sands"

R is about ~~statistics~~ programming data



Adapted from “Excursionist Drama 2”, p. 72 of Ben Katchor’s “Julius Knipl Real Estate Photographer”, Little Brown and Company, 1996.

The trouble with data

- You go to a lot of trouble to acquire it, and it gets large and unwieldy.
- One needs tools to work with it at scale.



Michael Pangrazio matte painting of the final warehouse in
“Raiders of the Lost Arc” (Paramount Pictures, 1981).

The problem

- Need an **R** grammar of data transforms that works well and works the same the same both in-memory and on large data systems (e.g., **Apache Spark**).
- Candidates
 - **SQL** (run in-memory with **sqldf** or by other round-tripping through the data store).
 - Too verbose and hard to maintain.
 - **dplyr / dbplyr**
 - **dbplyr** (despite claims) does not work the same as **dplyr**.
 - User-facing lazy-evaluation semantics break user expectations and make cross-team development difficult.

The solution

- Go back to influential sources for ideas.
- Implement Codd's relational algebra in a piped notation.
- Explicitly manage a separation between specification and execution.

Zeus giving birth to Athena, Rudolph Tegner



Athena may have leaped from Zeus's head, fully grown and armed; but even she didn't appear out of nowhere.

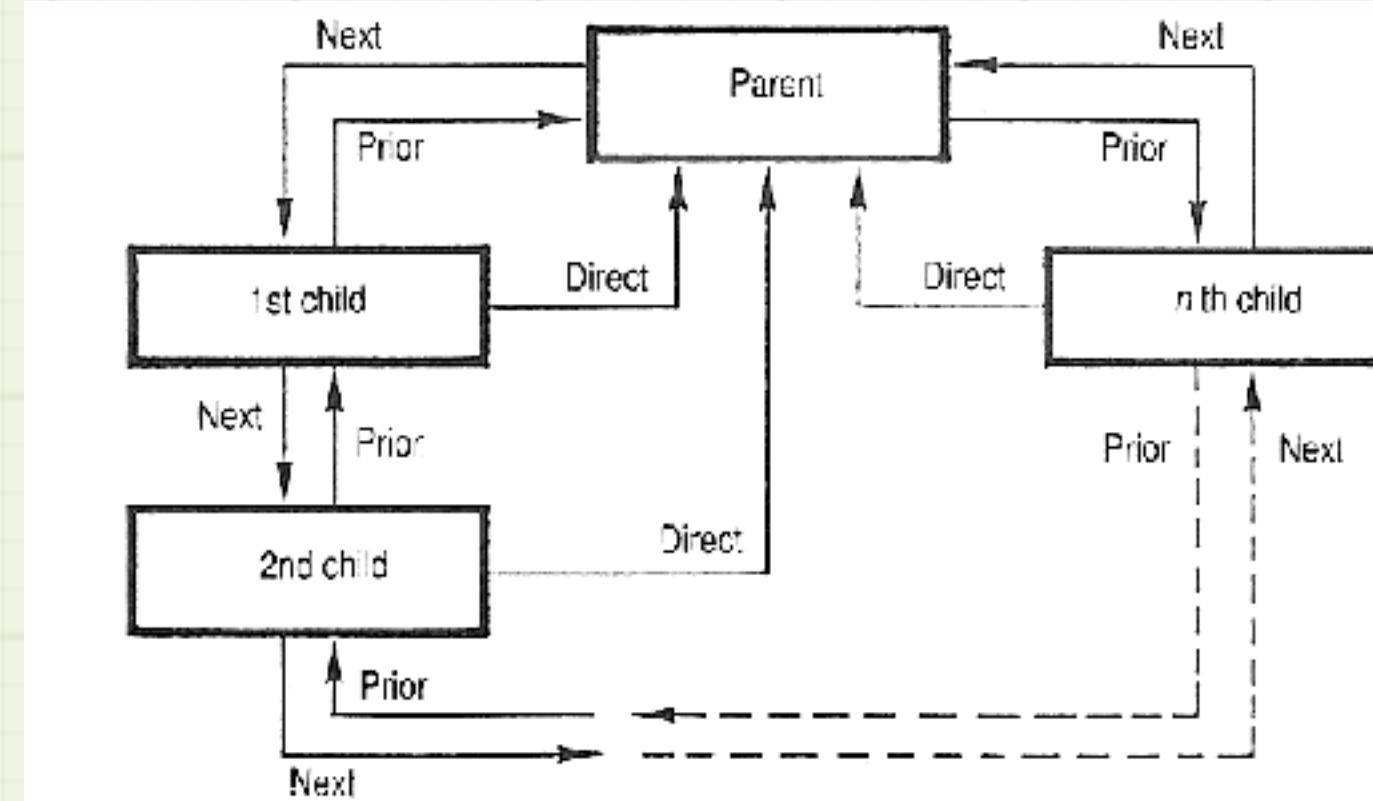
Data Manipulation

- First big idea: pointer chasing
- **CODASYL** (Conference/Committee on Data Systems Languages, 1959).
- **MongoDB** / NoSQL
- **JSON**
- **ORM** (Object Relational Mapping)

CODASYL

- Data is found by chasing around unidirectional references or pointers.
- Hierarchical or network data model.
- Influential *to this day*.
- Data is navigated by a cursor in an imperative style.

```
OBTAINT CALC CUSTOMER.  
PERFORM ORDER-LOOP UNTIL END-OF-SET.  
    ORDER-LOOP.  
        OBTAIN NEXT ORDER WITHIN CUSTOMER-ORDER.  
        MOVE ORDER-NO TO OUT-REC.  
        WRITE OUT-REC.
```



A closed chain of records in a navigational database model (e.g. CODASYL), with **next pointers**, **prior pointers** and **direct pointers** provided by keys in the various records.

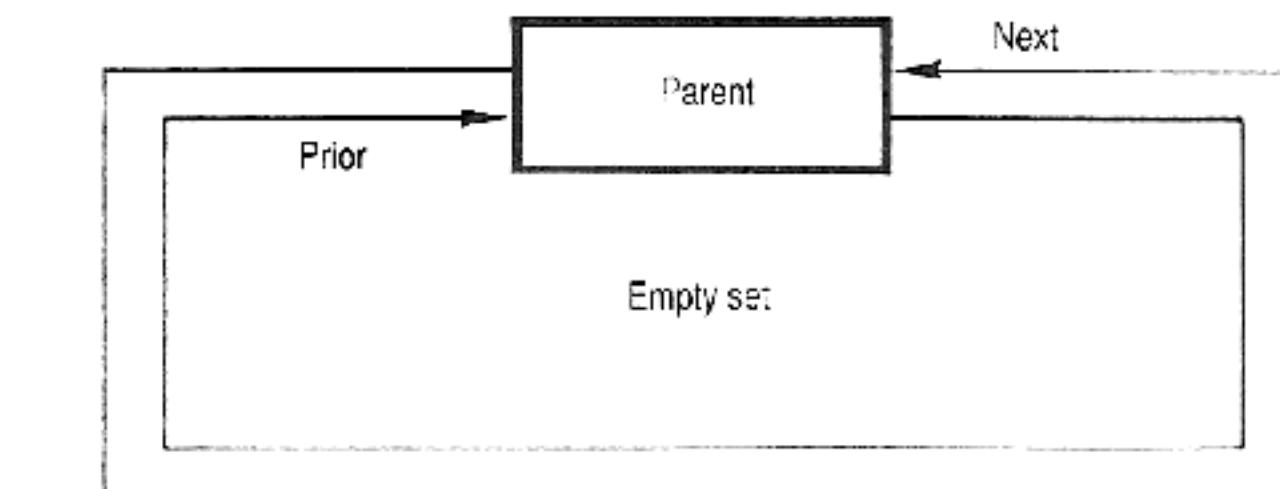


Illustration of an **empty set**

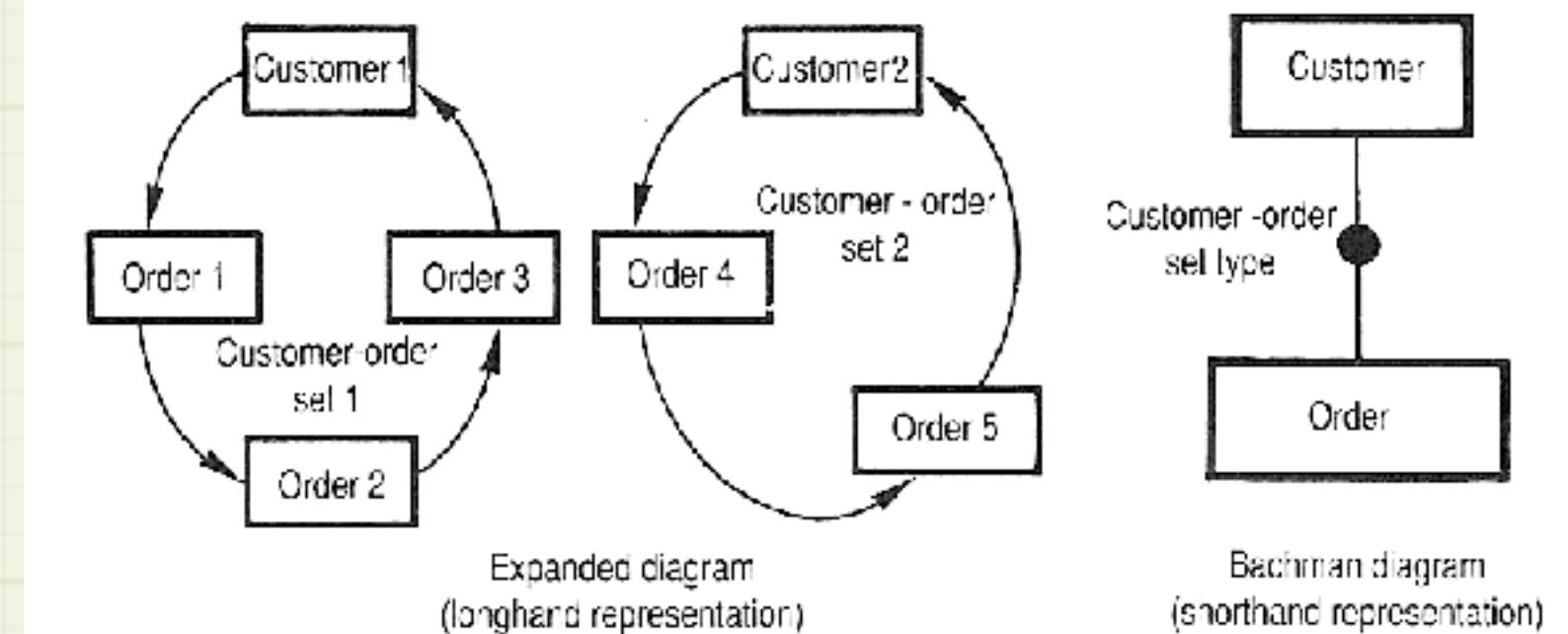
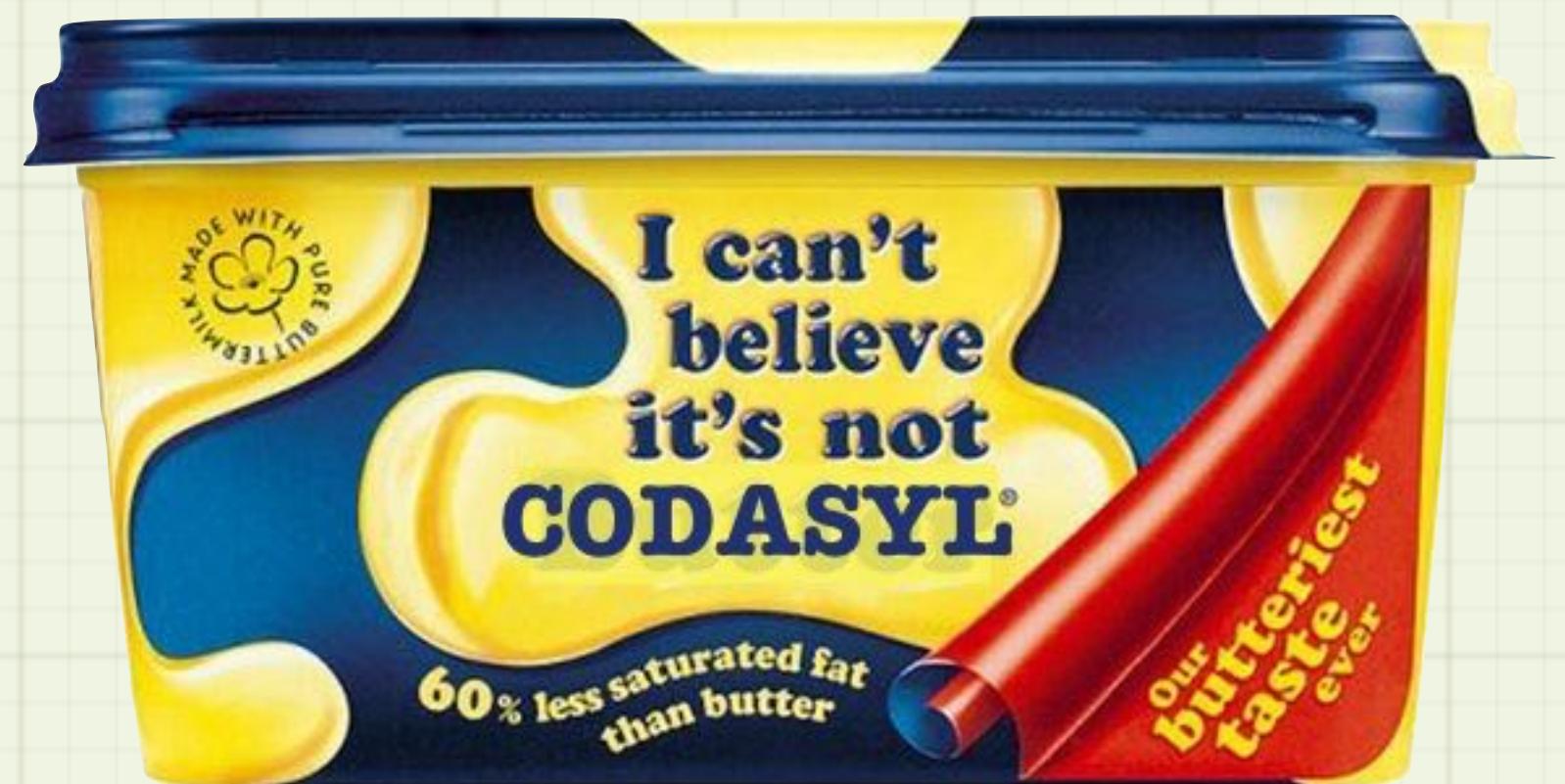
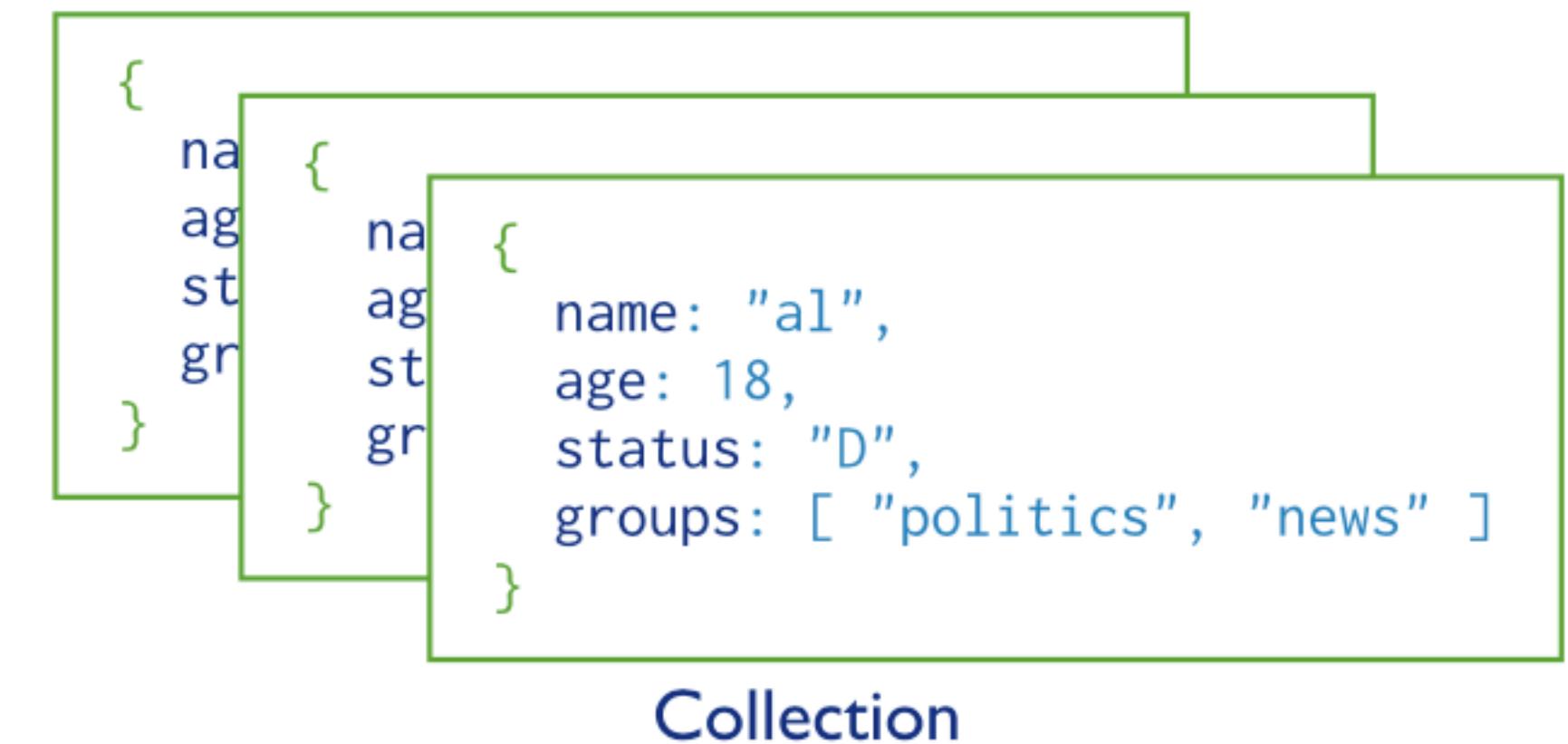


Illustration of a set type using a **Bachman diagram**

The record set, basic structure of navigational (e.g. CODASYL) database model. A set consists of one parent record (also called "the owner"), and n child records (also called members records)

MongoDB

MongoDB stores **BSON documents**, i.e. data records, in **collections**; the collections in databases.



<https://docs.mongodb.com/manual/core/databases-and-collections/>

JSON

- Nested arrays and maps.

```
{  
    "firstName": "John",  
    "lastName": "Smith",  
    "isAlive": true,  
    "age": 27,  
    "address": {  
        "streetAddress": "21 2nd Street",  
        "city": "New York",  
        "state": "NY",  
        "postalCode": "10021-3100"  
    },  
    "phoneNumbers": [  
        {  
            "type": "home",  
            "number": "212 555-1234"  
        },  
        {  
            "type": "office",  
            "number": "646 555-4567"  
        },  
        {  
            "type": "mobile",  
            "number": "123 456-7890"  
        }  
    ],  
    "children": [],  
    "spouse": null  
}
```

<https://en.wikipedia.org/wiki/JSON>

“Do you want *nested for-loops*?! Because that's how you get *nested for-loops*!”

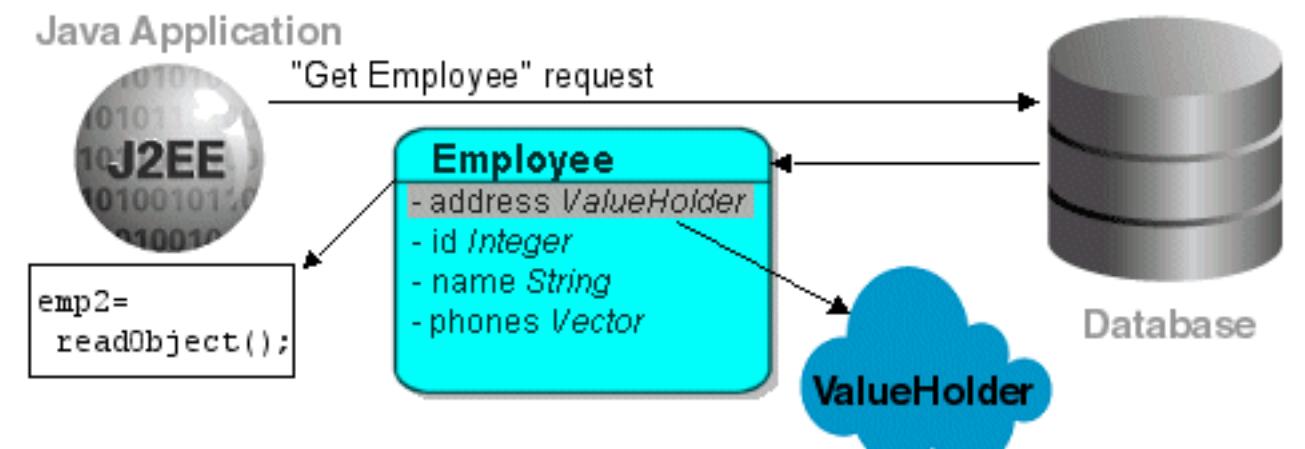


Adapted from Malory Archer, in “Archer” pilot episode.

ORM

[Figure 7-9](#) shows the Employee object being read from the database. The Address object is not read and will not be created unless it is accessed.

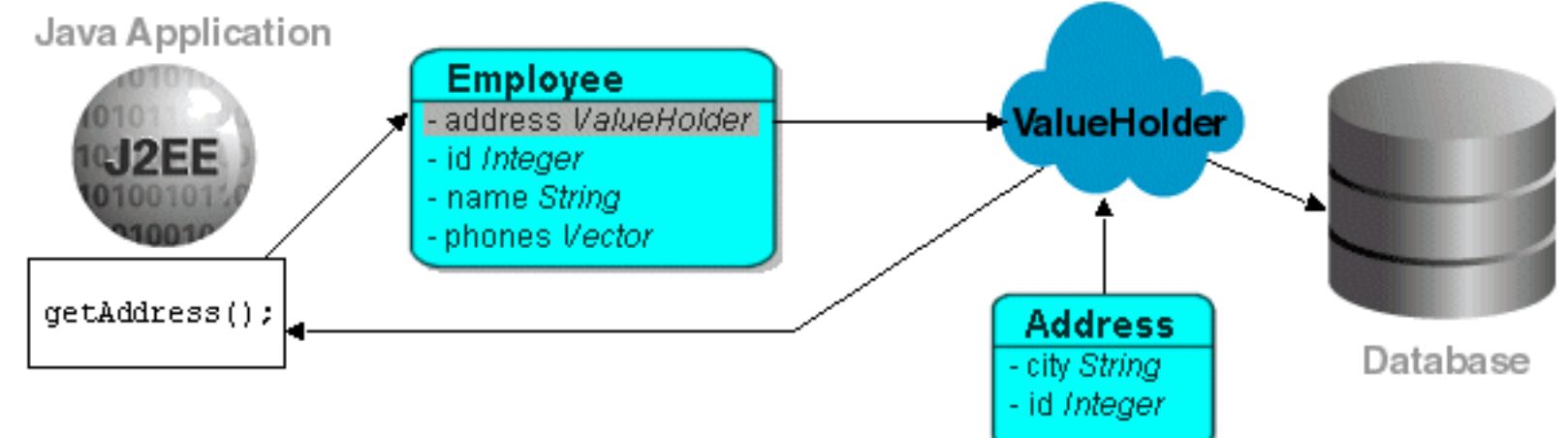
Figure 7-9 Address Object Not Read



[Description of "Figure 7-9 Address Object Not Read"](#)

The first time the address is accessed, as in [Figure 7-10](#), the ValueHolder reads and returns the Address object.

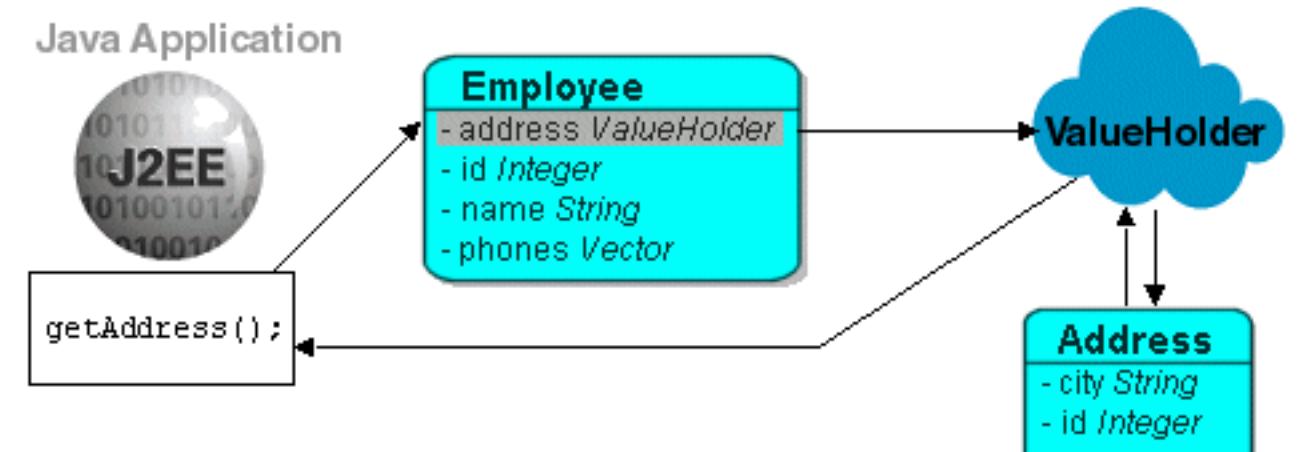
Figure 7-10 Initial Request



[Description of "Figure 7-10 Initial Request"](#)

Subsequent requests for the address do not access the database, as shown in [Figure 7-11](#).

Figure 7-11 Subsequent Requests



[Description of "Figure 7-11 Subsequent Requests"](#)

<http://www.eclipse.org/eclipselink/documentation/2.4/concepts/mappingintro002.htm>

“Well that just sounds like pointer chasing with extra steps.”



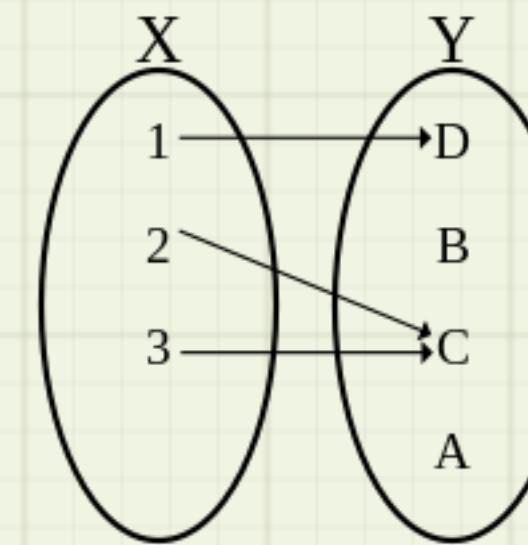
Adapted from: Rick and Morty "The Ricks Must Be Crazy."

A very big idea

- Relational algebra
 - Codd, E.F., “A Relational Model of Data for Large Shared Data Banks”, Communications of the ACM 13, (June 1970).
 - Delegate all of your data wrangling tasks to a small set of powerful operators.
 - Example archetypal powerful operator for data scientists: left join (we will return this after deriving the relational model).

Abstract and Generalize

- Pointers are an example of functions.
- Functions can be written as tables with the condition that the domain column has unique entries.
- Relax the table conditions to have merely every row be unique (or a set of rows) and you have what is called a relation.



[https://en.wikipedia.org/wiki/Function_\(mathematics\)](https://en.wikipedia.org/wiki/Function_(mathematics))

domain	range
1	D
2	C
3	C

X	Y
1	D
2	C
3	C
ω	A
ω	B

Same data structure (table) can be used to represent data or relations between tables!

ω (lower-case omega) represents no-value, similar to **NULL**, **NA**, or, \perp .

An entire theory of data wrangling

- Codd could define complex operators as equivalent to a sequence of simpler operators.
 - Example: in Codd's theory the full outer join (\bowtie) is defined as:
$$R \bowtie S := (R \bowtieleft S) \cup (R \bowtiright S).$$
(where \bowtieleft and \bowtiright are the left and right joins respectively).
- Could also prove different re-arrangements of operators were formally equivalent.
 - Basis for optimizing query planners to this day.
 - Example: pivoting a row selection prior to an expensive operation such as a join.

The realization

- **SQL** (Structured query language, 1974).
- *Far better system than the ugly syntax would suggest.*
- Essentially made Oracle Corporation
 - Oracle released a commercial **SQL** offering in 1979. This is shortly after IBM's **System R** (first customer: Pratt & Whitney in 1977) and *before* IBM's general commercial offering: **DB2** (1983).

SQL

- Further relaxed the mathematics from tables that represent sets of rows to arbitrary tables (collection of rows now a “bag” or “multiset”).
 - Improves the ability to represent data and makes some operations faster.
 - Don’t have to de-duplicate rows.
 - Makes joins more confusing.
 - Joins longer defined in terms of set-operations such as union.
 - Invites tons of ugly questions on how many rows each join situation generates.
 - No longer the case that all tables are relations.
 - Some optimizations no longer possible due to some theorems not carrying over.
 - Example failing theorem: distributive law of intersection over union fails

$$R \cap (S \cup T) \neq (R \cap S) \cup (R \cap T)$$

(mostly due to “`U`” being re-defined as essentially `rbind()`).

- Good description of the extended theory: Hector Garcia-Molina; Jeffrey D. Ullman; Jennifer Widom, *Database systems: the complete book (2nd ed.)*, Pearson Prentice Hall, (2009).

Relational Databases

- Essentially known by the join operation.
 - In a relational database you assume you have an effective join operator and write many other tasks in terms of join (and other steps).
- Declarative.
 - Instead specifying how to walk through the data the user declares the desired transformation. Like working with thick gloves on, awkward but safe and strong.
- For data scientists the natural left-join is *very* useful.
 - Example task: add columns from a new table **S** to matching rows in our left table **R**.

<https://www.jpl.nasa.gov/blog/2012/7/remote-controlled-manipulators>



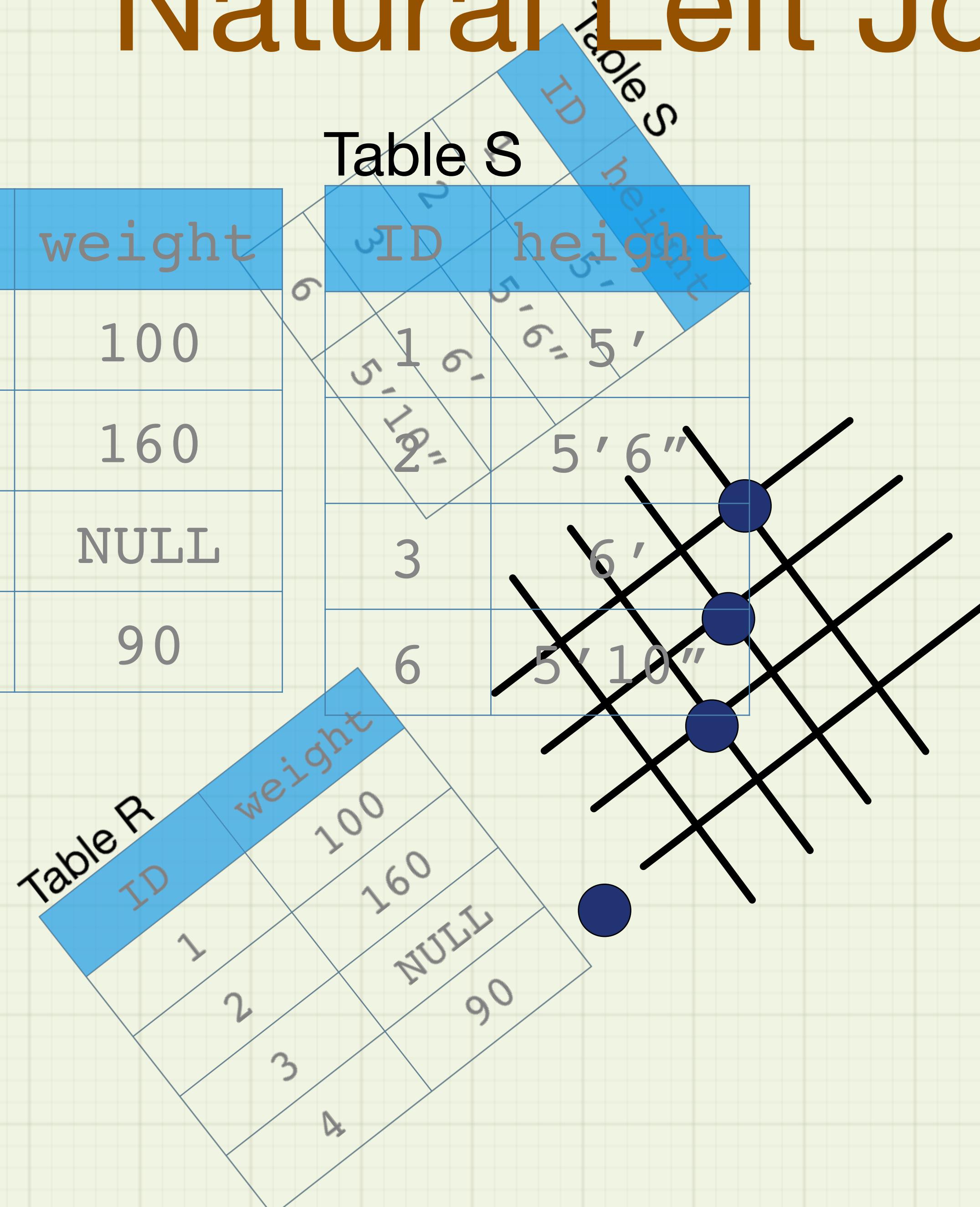
Natural Left Join Example

Table R

ID	weight
1	100
2	160
3	NULL
4	90

Table S

ID	height
1	5'
2	6'
3	5' 6"
4	5' 10"
5	6'
6	5' 6"
7	5' 10"
8	6'
9	5' 6"
10	5' 10"
11	6'
12	5' 6"
13	5' 10"
14	6'
15	5' 6"
16	5' 10"
17	6'
18	5' 6"
19	5' 10"
20	6'



`left_join(R, S)`

ID	weight	height
1	100	5'
2	160	5' 6"
3	NULL	6'
4	90	NULL

joins

- Joins replace for-loops and pointer chasing.
- We now think in bulk: annotate all rows with this new column (instead of for each row find an annotation).

Codd relational operators

- Small set of operators that most data-wrangling tasks can be decomposed into.
 - **join** (we just saw)
 - **project/aggregate** (delete columns and also produces group-summaries).
 - **extend** (add new calculated columns such as $x+y$).
 - **select_rows** (take a subset of rows based on a criterion).
- Complex transforms are expressed as a sequence of simpler transforms.
- Notable gaps:
 - Window-functions (such as ranking or grouped ranking).
 - Standardized in **SQL92**, commonly available.
 - Transitive closure / graph reachability.
 - Part of **SQL3** (1999), mostly still non-standard and not usually offered as part of non-graph databases.
 - Common table expressions and correlated sub-queries.

rquery

- **rquery** is a grammar for data wrangling based on Codd's relational algebra and experience working with **SQL** and **dplyr** at big data scale.
 - **rquery** uses the **wrapr** “dot arrow” to supply legible left to right pipe notation.
 - Consider **x %.>% f(.)** as an approximate synonym for **f(x)**.
 - Doesn't seem like much but it turns out **x %.>% f(.) %.>% g(.)** is easier to build up piece by piece than **g(f(x))** (and doesn't require reading backwards).
 - **rquery** is primarily a **SQL** query generator.
 - **rquery** depends on external systems (such as **SparklyrR**, **SparkR**, **PostgreSQL**) for implementation.
- **rqdatatable** is an in-memory implementation of the **rquery** grammar based on **data.table**.

What do we mean by SQL query generation?

Work through

[**rquery/extras/NarrowExample.Rmd**](#)

Rosetta Stone

Details on pipe-ready R-equivalents at https://github.com/WinVector/wrapr/blob/master/extras/pipe_base.R

Primary **dplyr** verbs can be found at <https://dplyr.tidyverse.org>

relational algebra	SQL	rquery	R / data.table	dplyr
π project (column restriction)	SELECT expressions	select_columns()	.[, cols, drop = FALSE]	select()
aggregation	GROUP BY	project_nse() / project_se()	tapply() / aggregate() "j=" (data.table)	group_by() summarize() ungroup()
extend (extended projection)	SELECT expressions	extend_nse() / extend_se()	transform() / := (data.table)	mutate()
X	ORDER BY	orderby()	.[order(), , drop = FALSE]	arrange()
σ select (row restriction)	WHERE	select_rows_nse() / select_rows_se()	subset()	filter()
\bowtie (left outer join)	LEFT JOIN	natural_join(jointype = "LEFT")	merge(, all.x = TRUE)	left_join()

The **rquery *_nse()** forms are the expected expression capturing interfaces, and the ***_se()** are the value oriented interfaces needed for effective abstraction and programming.

An example of the benefit can be found here <https://github.com/WinVector/rquery/blob/master/extras/CollectExprs.md>.

A substantial example

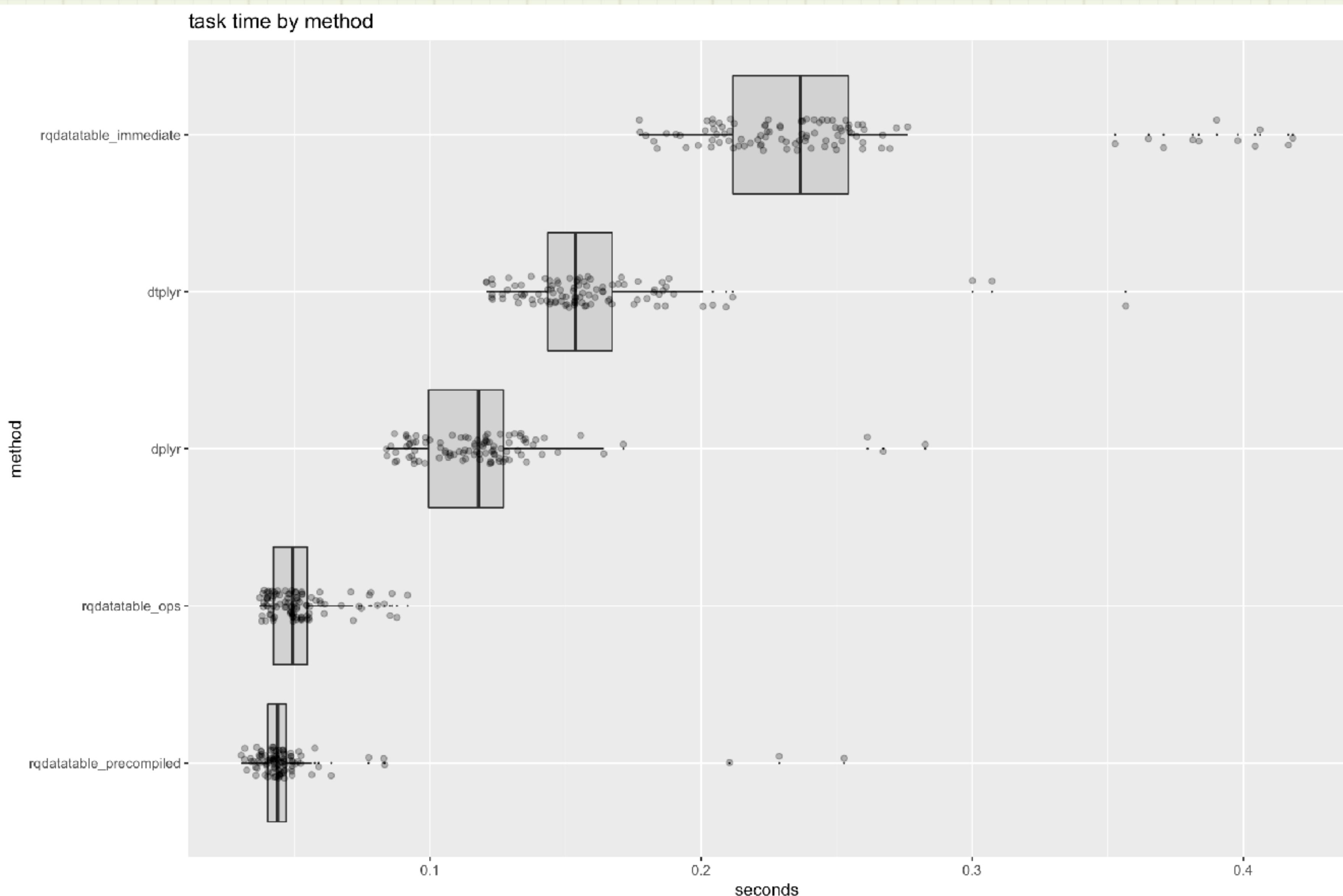
Work through

rquery/extras/IrisExample.Rmd

Immediate Mode

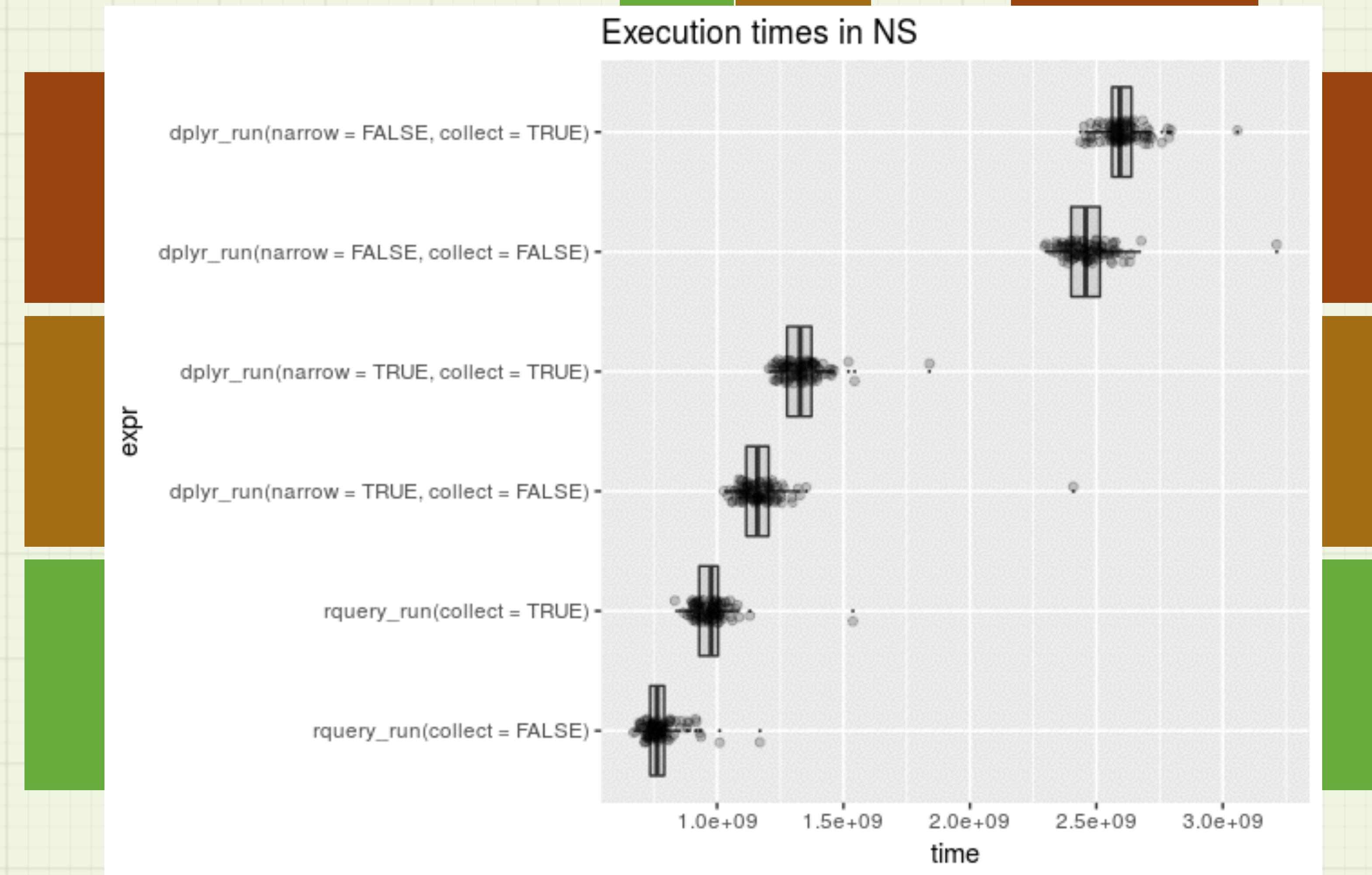
- **rqdatatable** in immediate mode is myopic (only can see one stage at a time) and fighting to bridge the difference between **data.table** reference semantics and expected **R** value semantics.

- Fully avoidable by building an **rquery** operator tree object and then piping data into that object.
- **dplyr** documents having similar issue (though no current way to avoid it).

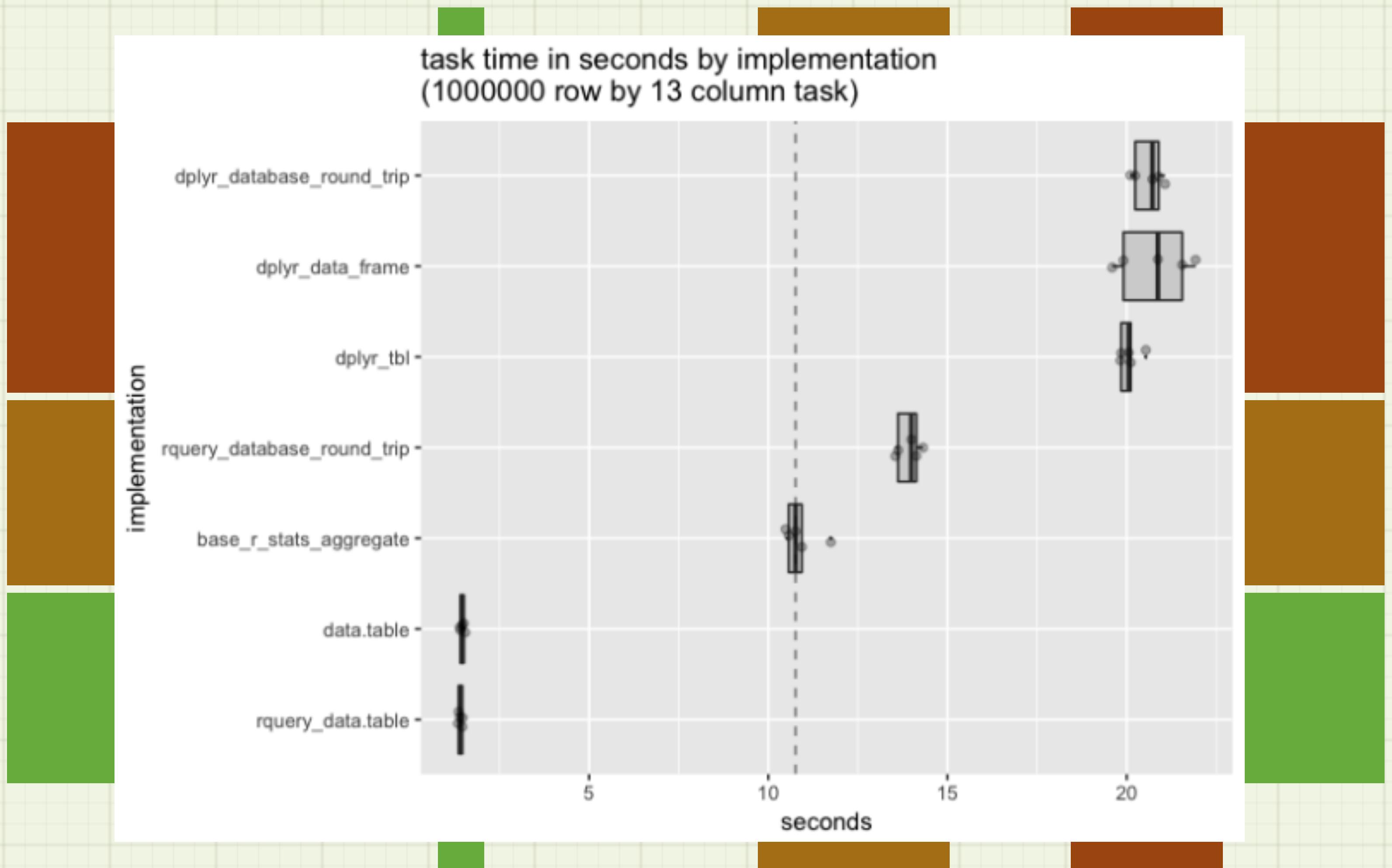


More Results

Performance (**sparklyr**, 40000 rows 1003 columns)



rqdatatable



Notice both **rquery** database round-trip and base R are much faster than **dplyr**. This is common, but contrary to many unfounded claims.

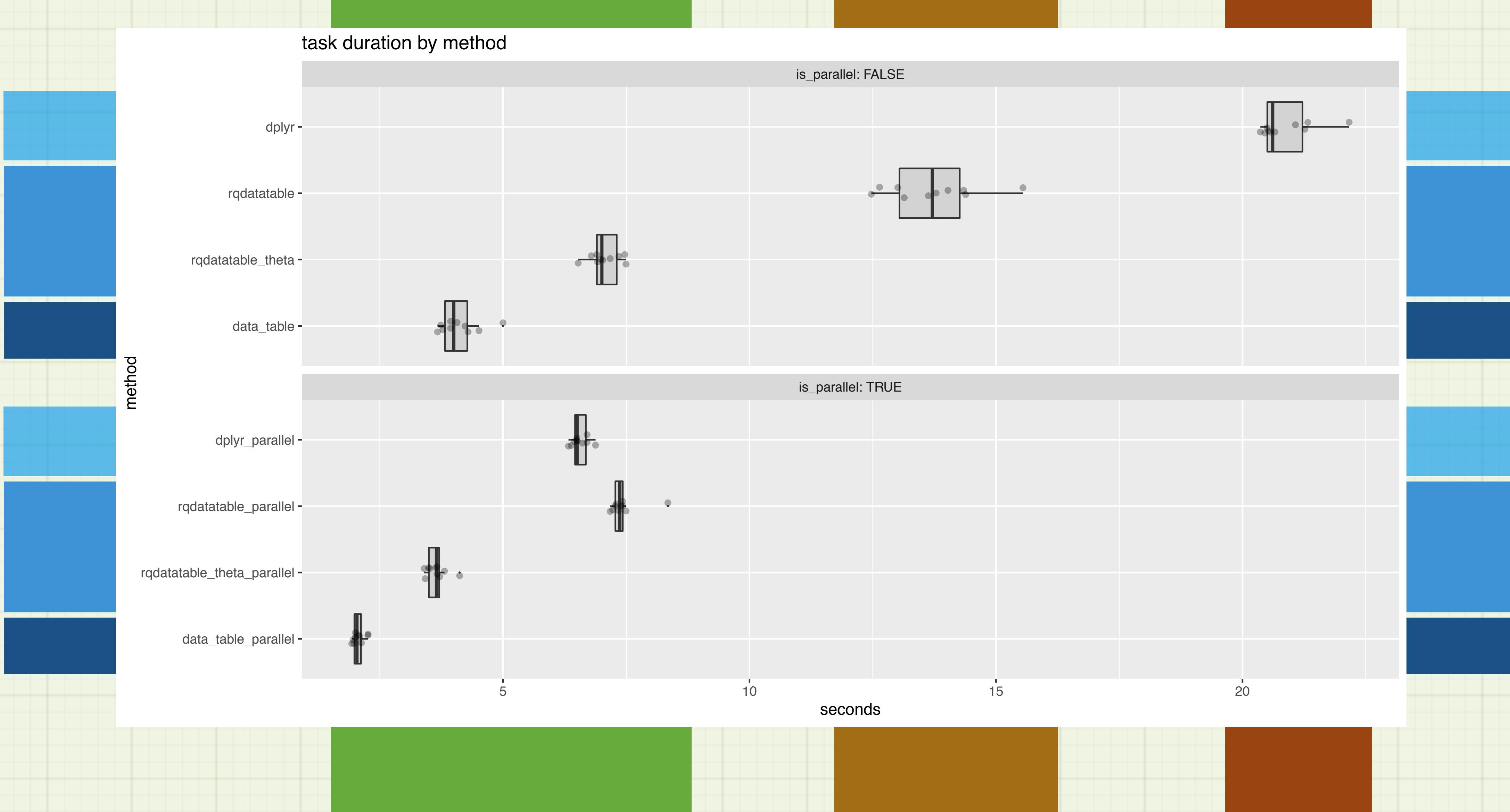


Notice *nothing* prior to these rows is in fact fast. Both these results are due to **data.table**.

`rquery_data.table == rqdatatable`, database is PostgreSQL

<http://www.win-vector.com/blog/2018/06/rqdatatable-rquery-powered-by-data-table/>

wrapr::execute_parallel()!!



rquery on SparkR (with DataBricks)!



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rquery: Practical Big Data Transforms for R-Spark Users

How to use rquery with Apache Spark on Databricks



by Nina Zumel and John Mount

Posted in ENGINEERING BLOG | July 26, 2018

This is a guest community blog from [Nina Zumel](#) and [John Mount](#), data scientists and consultants at [Win-Vector](#). They share how to use rquery with Apache Spark on Databricks

Try this notebook in Databricks

Introduction

In this blog, we will introduce [rquery](#), a powerful query tool that allows R users to implement powerful data transformations using [Apache Spark on Databricks](#). [rquery](#) is based on [Edgar F. Codd's relational algebra](#), informed by our experiences using SQL and R packages such as [dplyr](#) at big data scale.

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<https://databricks.com/blog/2018/07/26/rquery-practical-big-data-transforms-for-r-spark-users.html>

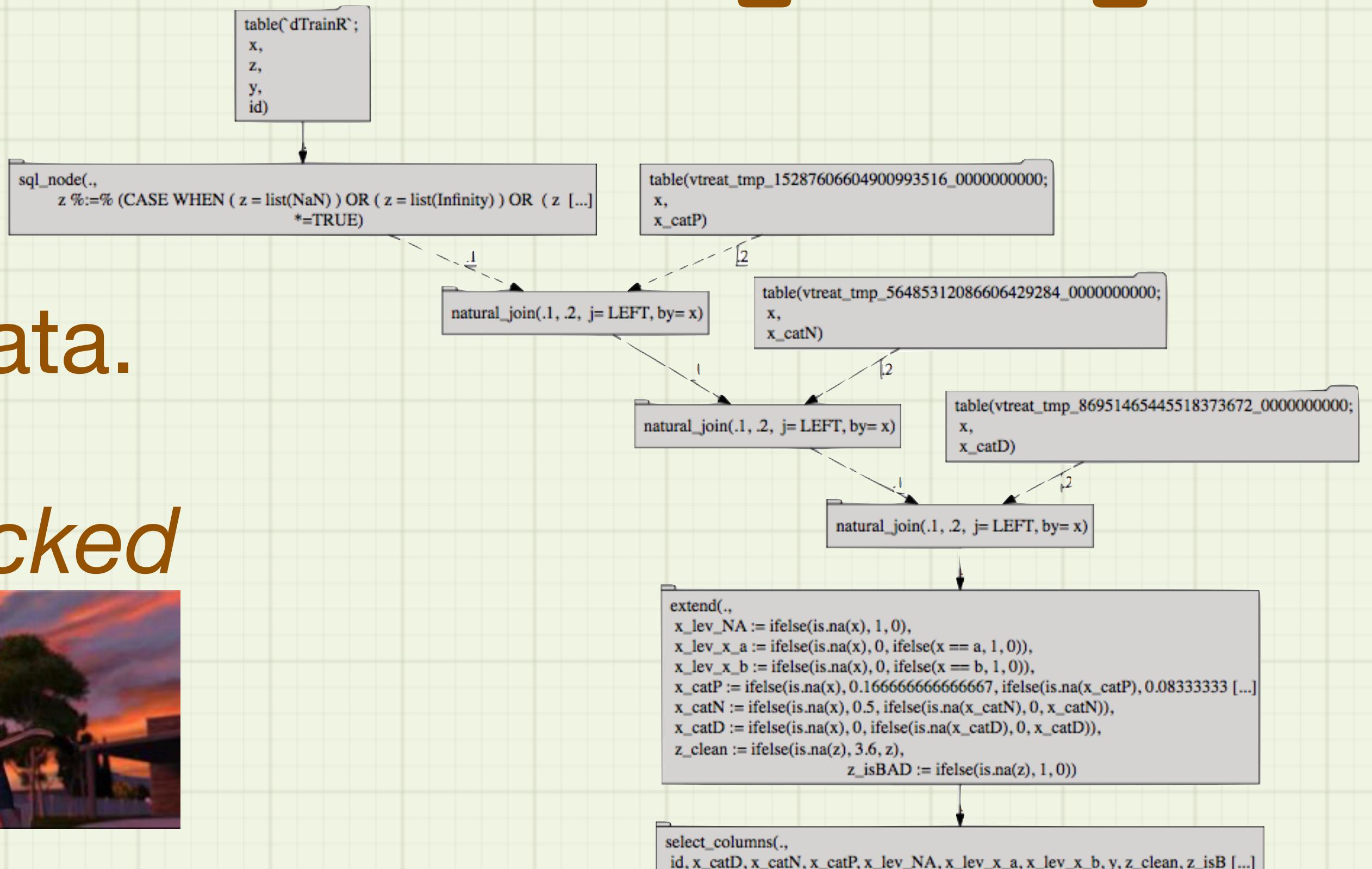


vtreat hosted on rquery

- Enables **vtreat** on big data.



- Makes for some *totally wicked* op diagrams.



- https://github.com/WinVector/vtreat/blob/master/extras/rquery_vtreat.md
- <https://github.com/WinVector/vtreat/blob/master/extras/vtreatOnSpark.md>

Conclusion

- **rquery** is an excellent query generator for **R** in terms of performance and usability. Using it can increase your team's productivity on **R** projects.
 - **rquery** is a best of breed solution in terms of:
 - Error Checking
 - Correctness
 - Usability
 - Performance.
 - Building up experience with it mostly with **PostgreSQL** and **Spark**.
- **rqdatatable** is a fast in-memory realization of **rquery** supplied by **data.table**.
- I would *love* to explore ways to collaborate and get further introductions.
 - **Please** reach out to me at jmount@win-vector.com. I would especially like to meet with groups considering working with **R** and databases or **Spark**.

Thank You