

# Logic

### Problem 1

Let  $F$  be the set of well-formed formulas with propositional variables from PROP. Define a relation,  $R \subseteq F \times F$  by  $(\varphi, \psi) \in R$  if  $\varphi \models \psi$ . Prove or give a counter-example to disprove:  $\neg(\psi, \varphi) \in R$ .

- (a)  $R$  is a partial order.  $\phi \models \psi : \phi \rightarrow \psi$  is tautology.  $(\phi, \psi) \in R \iff \phi \models \psi$ .  
 (b)  $R \cup R^{\leftarrow}$  is an equivalence relation. Suppose  $(a, b) \in R$ . Then  $(b, a) \in R^{\leftarrow}$ .  
 (c)  $R \cap R^{\leftarrow}$  is an equivalence relation. Suppose  $(a, b) \in R \cap R^{\leftarrow}$ . Then  $(a, b) \in R$  and  $(b, a) \in R$ .

### Problem 2

Prove that  $\neg N$  follows logically from  $H \wedge \neg R$  and  $(H \wedge N) \rightarrow R$ .

### Problem 3

Consider the formulae  $\phi_1 = (r \rightarrow p)$  and  $\phi_2 = (p \rightarrow (q \vee \neg r))$ . Transform the formula  $\phi = (\neg q \rightarrow (\phi_1 \wedge \phi_2))$  into

- (a) **DNF**, and  
(b) **CNF**.

Simplify the result as much as possible.

### Problem 4

Let  $(T, \wedge, \vee, ', 0, 1)$  be a Boolean Algebra. Define  $\oplus : T \times T \rightarrow T$  as follows:

$$x \oplus y = (x \wedge y') \vee (x' \wedge y)$$

- Prove using the laws of Boolean Algebra that for all  $x \in T$ ,  $x \oplus 1 = x'$ .
- Prove using the laws of Boolean Algebra that  $x \wedge (y \oplus z) = (x \wedge y) \oplus (x \wedge z)$ .
- Find a Boolean Algebra (and  $x, y, z$ ) which demonstrates that  $x \oplus (y \wedge z) \neq (x \oplus y) \wedge (x \oplus z)$

### Problem 5

- How many well-formed formulas can be constructed from one  $\vee$ ; one  $\wedge$ ; two parenthesis pairs  $(, )$ ; and the three literals  $p$ ,  $\neg p$ , and  $q$ ?
- Under the equivalence relation defined by **logical equivalence**, how many equivalence classes do the formulas in part (a) form?