

# Spatial Social Community

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## 1 Preliminary

- Trusses
- Triangle connected k-truss
- K-truss community
- Finding k-truss communities

## 2 Spatial Constraints

- An example
- Basic solution
- Example

## Definition

A  $k$ -truss is a none-trivial, one-component subgraph such that each edge is reinforced by at least  $k-2$  pairs of edges making a triangle with the edge. (Non-trivial here excludes an isolated vertex as a truss)

# Triangle connected k-truss

- $k \geq 3$
- Triangle adjacency: given two triangles, they are adjacent if and they share a common edge
- Triangle connectivity: given any two triangles  $\triangle_s$  and  $\triangle_t$  in  $G$ , they are connected if there exist a series of triangles  $\triangle_1, \dots, \triangle_n$  in  $G$ , where  $n \geq 2$ , such that,  $\triangle_1 = \triangle_s$ ,  $\triangle_n = \triangle_t$  and for  $1 \leq i < n$ ,  $\triangle_i$  and  $\triangle_{i+1}$  are adjacent

## Definition

K-truss community: 1)k-truss, 2)triangle connected, and 3)maximal subgraph

## Basic

- Edge trussness index: running k-truss decomposition
- Query k-truss communities from a vertex  $v$ : running BFS search from edges containing  $v$

## Better

- TCP index: it is built on top of edge trussness index
- Query k-truss communities from a vertex  $v$ : running BFS search from  $v$

# Finding k-truss communities

- Observation: Given a  $k$ , a edge  $e$  in  $G = (V, E)$  can only be contained by at most one  $k$ -truss community
- Finding  $k$ -truss communities: For each unvisited edge  $e$  with trussness no less than  $k$  in  $G = (V, E)$ , we run BFS from it.



- Observation: A spatial constrained  $k$ -truss community must be contained by a  $k$ -truss community or equal to a  $k$ -truss community
- Based on the observation, in the following part, the discussion focuses on finding spatial constrained  $k$ -truss communities in a  $k$ -truss community.
- A  $k$ -truss community may contain a set of two-vertexes-pairs. Any pair of vertexes in the set, the distance of the two vertexes is larger than  $d$ .
- Observation: for each pair of vertexes, we only need to remove one of the vertex (different removal choices induce different combinations)
- Observation: For the set of pairs of vertexes, the processing order of pairs of vertexes does not affect the result but affects the performance
- Observation: After dealing with all pairs of vertexes whose distances are larger than  $d$ , the remaining of  $k$ -truss community is subgraph of the  $k$ -truss community (could be disconnected).  $K$ -truss communities in the remaining subgraph are spatial constrained  $k$ -truss communities if they exist.





# Table

<b>Treatments</b>	<b>Response 1</b>	<b>Response 2</b>
Treatment 1	0.0003262	0.562
Treatment 2	0.0015681	0.910
Treatment 3	0.0009271	0.296

Table: Table caption