# FIT5037: Network Security Computer system security and malicious code

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## Lecture 9: Computer system security and malicious code

#### Lecture Topics:

- Symmetric key cryptography
- Asymmetric key cryptography
- Pseudorandom Number Generators and hash functions
- Authentication Methods and AAA protocols
- Security at Network layer (IPsec)
- Security at Network layer (firewalls and wireless security)
- Security at Transport layer
- Security at Application layer
- Computer system security and malicious code
- Computer system vulnerabilities and penetration testing
- Intrusion detection
- Denial of Service Attacks and Countermeasures / Revision

#### Outline

- Malicious code definition and types
- Remote code execution
- Malware incident prevention
  - policy
  - awareness
  - Vulnerability mitigation
  - Threat mitigation
  - Defensive architecture
- Malware incident response
  - Preparation
  - Detection and analysis
  - Eradication and recovery
  - Post-incident activities



#### Malicious Code: Definition

• RFC 4949: Internet Security Glossary, Version 2 defines "malicious logic" as:

Hardware, firmware, or software that is intentionally included or inserted in a system for a harmful purpose.

- the RFC categorises the potential impact of malicious logic as:
  - corruption
  - incapacitation
  - masquerade
  - and misuse
- examples of malicious logic in RFC4949:
  - virus
  - worm
  - logic bomb
  - Trojan horse
  - spyware (deprecated term due to broad classification)

### Malicious Code: Types

#### virus

- a self-replicating (and usually hidden) section of computer software (usually malicious logic)
- propagates by infecting i.e., inserting a copy of itself into and becoming part of another program
  - A virus cannot run by itself; it requires that its host program be run to make the virus active.
- NIST SP 800-83r1 divides viruses into two classes
  - compiled: executed by OS such as file infector viruses and boot sector viruses
  - interpreted: executed by an application such as macro viruses infecting application documents

#### worm

- a computer program that can run independently
- propagates a complete working version of itself onto other hosts on a network
  - may consume system resources destructively
- NIST SP 800-83r1 divides worms into two categories
  - Network service worms which exploits a network service vulnerability to propagate itself
  - Mass mailing worms which is a self contained malicious code that uses email to propagate itself
- logic bomb
  - activates when specified conditions are met
  - Usually intended to cause denial of service or otherwise damage system resources.



## Malicious Code: Types (continued)

- Trojan horse
  - a computer program that appears to have a useful function
  - also has a hidden and potentially malicious function that evades security mechanisms
    - sometimes by exploiting legitimate authorizations of a system entity that invokes the program.
- spyware
  - this term is widely used to describe different malicious logic (deprecated according to RFC4949 for this reason)
  - Software that an intruder has installed surreptitiously (secretly) on a networked computer to:
    - gather data from that computer
    - and send it through the network to the intruder or some other interested party.
- ransomware
  - a type of malicious code that blocks or limits access to a system until a ransom is paid
    - lockers: lock down the system
    - encryptors: encrypt specific files e.g. WannaCry
- malicious mobile code (NIST)
  - software with malicious intent that is transmitted from a remote host to a local host (hence mobile)
  - executed locally without the user's explicit instruction
  - popular languages: Java, ActiveX, JavaScript, and VBScript
- blended attacks (NIST)

#### Attacker tools

#### NIST SP 800-83r1 provides a list of attacker tools that may be delivered to an infected host

- backdoors
  - malicious code that listens for commands on a certain TCP or UDP port
- allows an attacker to perform certain tasks e.g. acquire passwords or execute commands
- Zombies or Botnets are a collection of hosts with backdoors under an attacker control
- key loggers
  - monitors and records keyboard use
  - attacker may retrieve data or key logger may transfer data to the command and control host
- rootkits
  - collection of files to change the functionality of the infected host in a malicious way
  - tries to hide its existence
- web browser plug-ins
  - malicious plug-ins can monitor all use of a browser
- email generators
  - can be used as a propagation engine for the attacker to infect other hosts in infected host network

Giffensive security tools: e.g. Wireshark, Kali Linux, nmap, openvas, Nessus, John the Ripper etc.

- attacker tool kit
- different types of utilities such as packet sniffers, port scanners, vulnerability scanners, password crackers etc.

#### Malicious code obfuscation

One of the main approaches of anti-virus applications is signature detection, to evade this security measure malicious code use obfuscation techniques

- encryption: the body of the malicious code is encrypted except for a small part that performs decryption
  - the encryption key changes with each infection hence signature (bit pattern) of the code changes
  - the key must be stored with the virus

#### Polymorphic mutation

- encryption + decryptor mutation
  - in addition to encrypting its code it mutates its decryptor using other obfuscation techniques

Metamorphic mutation techniques (changing appearance/signature/bit pattern)

- insertion of no operation or useless instructions (junk code)
  - malicious code has a mutation engine that changes the signature of the code with each infection
  - the code changes by insertion of no operation instructions between malicious code instructions
    - most instructions can be used here as longs the inserted instruction does not change the logic

## Malicious code obfuscation (continued)

#### Metamorphic mutation techniques ...

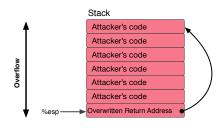
- register renaming/variable substitution
  - most processors provide multiple registers to store variables and intermediate values
    - mutation engine changes the registers used in each instruction without changing the logic
- subroutine permutation
  - the order of defined subroutines are changed which does not affect the logic of the code
  - the overall bit pattern of the code changes
- code reordering through jumps
  - the code is divided into sections and permuted and jump instructions are inserted to preserve logic
- equivalent code substitution/instruction replacement
  - part of the code (one or more instructions) is replaced by an equivalent code with different set of instructions (different bit pattern)

#### Remote Code Execution

- attacker exploits a vulnerability and is able to execute code on the target system
- examples of vulnerabilities that could lead to remote code execution
  - stack-based buffer overflow
  - heap-based buffer overflow
  - command injection
- exploitation techniques
  - shell code injection
  - return-to-libc (and return to library in general)
  - Return Oriented Programming
  - Jump Oriented Programming
- exploitation methods can be built into the malicious code as infection and propagation method

### Remote Code Execution: stack-based BOF shell code injection

- when a function is called current function state is pushed unto the stack
- problem is some control information is also pushed to the stack
  - current function stack frame address.
  - return address (to continue execution after the call)
- buffer over flow vulnerability
  - allows the user to write beyond assigned memory cells to a variable
- shell code injection
  - attacker injects instructions into the stack (overflow buffer)
    - generally to call system to execute a shell
  - attacker overwrites the return address with the address of the injected code



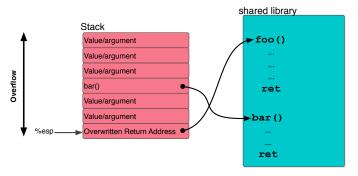
### Remote Code Execution: Mitigating shell code injection

- non-executable stack
  - marks the stack memory page as read/write but not executable
  - jump to the stack addresses results in an exception
- stack canaries or cookies
  - random values inserted between return address and called function parameters and stack variables
    - assumption: to overwrite the return address the canary value is also overwritten
  - check the random value when returning (function return is replaced with jump to verification code)
  - problem: array index (integer) overflow techniques can be used to overwrite return address without overwriting the cookie
    - cookie can be overwritten one byte at a time to find out the value of cookie (works in some implementation for fork)
- Address Space Layout Randomisations (ASLR)
  - attacker needs to know the address of the injected shell code to jump to it (replace return address with)
    - calculated based on the address of the stack
  - ASLR randomises the process address space
    - program load address, shared libraries, heap and stack base addresses are randomised and will be different with each execution
- Shadow stack
  - return address is also written in a shadow stack
  - the return address is verified at return from function call with value stored in shadow stack

### Remote Code Execution: return to library

#### return to library

- attacker replaces the return address with address of a function in a shared library
  - bypasses the non-executable stack
    - injected value is no longer code and is not executed
    - converts return into a function call



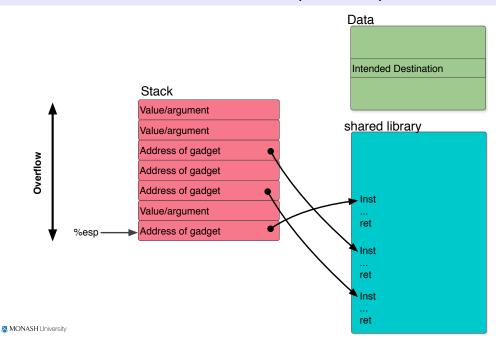
• return address can be an address in the middle of a function in shared library

#### Remote Code Execution: ROP

- instead of using the complete library function small pieces of code in library are used
  - each piece of code is called a gadget
  - a gadget is comprised of few instructions ending with return instruction
    - to get back to the stack that attacker controls
  - each gadget performs a task equivalent to a single machine instruction (in a typical processor)
    - e.g. adding two registers and storing the results
    - however use push and pop (to read from stack and write to stack) to read and write values
- attacker injects addresses and values into the overflown buffer in stack
  - generated by the ROP compiler
- it is shown that a Turing-complete machine (machine capable of executing programs) can be built using these gadgets
  - attacker writes the intended code using high level program
  - a tool compiles the code using gadgets in a specified target library
    - can find the gadgets within the code of a target program



# Remote Code Execution: ROP (diagram)



## Remote Code Execution: Mitigating ret-to-lib and ROP

- Address Space Layout Randomisations (ASLR)
  - attacker needs to know the address of shared libraries or the executable code itself
    - to inject address of library functions in case of ret-to-lib
    - to inject address of gadgets in case of ROP
- stack canaries/cookies
  - are effective against both
  - can be bypassed by other techniques
- Control Flow Integrity
  - techniques used to check whether jumps and returns follow Control Flow Graph of the program
    - statically generated CFG are prone to attacks (more permissive)
    - dynamically generated CFG are more accurate but more expensive (still active topic in research)
- Shadow stack
  - keeping a duplicate stack is expensive



#### Malware Incident Prevention

NIST SP 800-83r1 proposes the following elements to prevent malware incidents

- policy
- awareness
- vulnerability mitigation
- threat mitigation
- defensive architecture

### Malware Incident Prevention: Policy<sup>1</sup>

- organisation security policy should address malware incident prevention
  - policy statement can then be used to drive the incident prevention efforts
    - IT staff awareness, vulnerability mitigation, defensive architecture
  - without policy the efforts may not be effective and consistent
- Malware prevention policies should be general to be flexible in policy implementation
  - reduce policy update
  - however must be clear in intent and scope
  - should include remote workers
    - including those using hosts outside organisation's control
- Examples of Malware prevention policies:
  - require scanning of media from outside organisation
  - require scanning of email attachment
  - prohibit send/receive of certain file types
  - restrict or prohibit use of unnecessary software
  - restrict use of removable media
  - specify the required type of antivirus software for various types of hosts

 $\frac{1}{8}$  MNIST SP 800-83r1: Guide to Malware Incident Prevention and Handling for Desktops and Laptops

### Malware Incident Prevention: Awareness<sup>2</sup>

- explains proper rules of behaviour for use of the organisation IT hosts
- provides guidance to users on malware incident prevention
  - reduces the frequency and severity of incidents
- inform all users of
  - the ways malware enters and infects hosts
  - the risks that malware poses
  - technical controls (security measures) cannot prevent all incidents
- Examples of recommended practices for avoiding malware incident (informed users)
  - not opening suspicious emails from known or unknown entities, clicking on attachments, hyperlinks etc.
  - not clicking on suspicious browser pop-ups
  - not opening files with extensions that are likely be associated with malware
  - not disabling security mechanisms such as antivirus, personal firewall etc.
  - not downloading or executing applications from untrusted sources

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### Malware Incident Prevention: Vulnerability Mitigation<sup>3</sup>

- malicious code often exploits existing vulnerability to infect and propagate
- vulnerability mitigation can be achieved by
  - updating OS and applications
  - OS/software reconfiguration to disable vulnerable services
- automated configuration checklist for OS and applications
  - to check the configuration or to apply settings automatically
    - e.g. Microsoft Active Directory Policies
- automated patch management
- use host hardening principles (e.g. NIST SP 800-123: Guide to General Server Security)
  - least privilege
  - disable/remove unneeded services
  - remove unsecured file share
  - remove/change default username and passwords for OS and applications
  - disable automatic execution of binaries and scripts
  - change default file association for files commonly used by malware (e.g. .pif, .vbs)

### Malware Incident Prevention: Threat Mitigation<sup>4</sup>

- is still important even if all vulnerabilities are mitigated
  - e.g. social engineering: trick users to run malware
- antivirus software
  - recommended capabilities
    - scanning critical host components
    - monitor host for suspicious activities
    - monitor applications behaviour
    - file scan for known viruses (less effective against poly/metamorphic malicious code)
    - disinfecting files, quarantine infected files
  - use both host-based and network-based antivirus scanning tools
  - keep antivirus software as well as its signature database up-to-date
  - centrally managed
    - prevent users from disabling/removing the antivirus software
    - automatically update signature database
  - use multiple antivirus software from different vendors for key hosts

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# Malware Incident Prevention: Threat Mitigation<sup>5</sup> (continued)

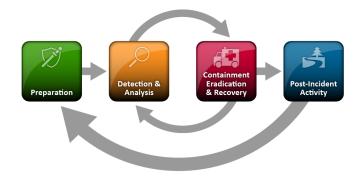
- intrusion prevention system
  - can prevent some network attacks before reaching hosts
  - are customisable allowing to add new attack signatures
  - some IPS can detect anomalous behaviour by monitoring normal network traffic patterns (used as baseline)
  - host-based IPS can monitor a single host using similar principle as network-based IPS
- firewalls
  - deny-by-default policy will block unintentionally open ports (unneeded services)
  - can block a port used by malware exploiting a vulnerable service
  - in case of a major incident firewall can be used to block the traffic at network level
  - or block external sources used by malware to download other malicious code
- Content Filtering/Inspection
  - used to prevent email-based malware propagation
  - spam filtering to prevent spam emails to deliver malware
  - block attachments with types associated with malware and suspicious file extensions
  - used to stop web-based malware threats
  - blacklist or reputation information
  - block browser pop-up windows
- A Practice Militarian
- Application Whitelisting
   MONASHUMVERINY provide a list of authorised applications for use on a host

### Malware Incident Prevention: Defensive Architecture<sup>6</sup>

- used in addition to vulnerability and threat mitigation, to reduce the impact
- BIOS protection
  - protects modification of BIOS
- sandboxing
  - isolate processes or threads within a controlled environment with restricted access to system resources
    - only *safe* operations are allowed within the sandbox
  - can potentially prevent all or some of the malware malicious activity
  - sandbox can be reset to a known-good state at initialisation
- browser separation
  - multiple brands are installed on a single host
    - different browsers are used for different activities (reducing the likelihood of compromised browser be used for sensitive data)
  - allows tighter control e.g. disabling all forms of mobile code for browsers used for corporate tasks
- segregation through virtualization
  - similar to sandboxing in concept however to segregate applications and OS
  - provides a more rigorous segregation
  - e.g. an OS for corporate tasks another for other activities
    - a compromise in one image does not affect the other unless the virtualization software is also compromised
- images can be reset to a known-good state when restarted

# Malware Incident Response<sup>7</sup>

- no matter how many protective measures are put in place malware incidents will happen
  due to unknown types or threats, human error etc.
- must have malware incident response plan to limit the damage and restore data and services
- NIST SP 800-83r1 defines the following life cycle for malware incident response
- Preparation
- Detection and Analysis
- Containment, Eradication, and Recovery
- Post-Incident Activity



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### Malware Incident Response: Preparation<sup>8</sup>

- building and maintaining malware-related skills
  - how each major category of malware infects and propagates
  - be familiar with organisation's malware detection tool configuration
  - be able to analyse supporting data
  - be able to identify characteristics of threats
  - have incident handlers temporarily work as antivirus engineers to learn new skills
- facilitating communication and coordination
  - have a coordination team for malware incident response
  - coordinator to provide guidance and instructions to all staff assisting with
    - containment
    - eradication
    - recovery
  - coordinator to update management
  - establish point of contact to answer questions about malware alerts
- acquire tools and resources
  - malware analysis toolkits

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# Malware Incident Response: Detection and Analysis<sup>9</sup>

- identifying malware incident characteristics
  - validate malware incident
    - false positives: benign activity as malicious
  - when in doubt treat as a malware incident until confirmed
  - since waiting can have significant impact on the organisation
  - helps with assigning priority to response efforts
  - collaborate with security administrators to identify data source aiding in detection
    - antivirus
    - IDS
    - Security Information and Event Management
  - search for identified characteristics
  - malware category
    - services, ports, protocols attacked
    - exploited vulnerabilities
    - malicious filenames, size, content
    - OS, devices, applications
- identify infected hosts
  - to undergo containment, eradication, and recovery

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# Malware Incident Response: Detection and Analysis<sup>10</sup> (continued)

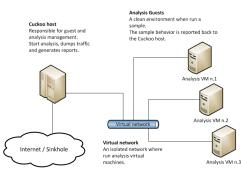
- identify infected hosts ...
  - forensic identification
    - searching for evidence of recent infections
    - security logs, dns server logs, application server logs, network forensic tools
  - active identification
    - security automation: continuous monitoring e.g. network access control technologies
    - custom network-based IDS and IPS: dedicated IDS/IPS with custom signatures for malware detection
  - packet sniffers and protocol analysers
  - configured to look for characteristics of a particular malware
  - manual identification
    - when automated methods are not feasible
- prioritising incident response
  - malware propagation speed e.g. worms
  - value of data when responding to Trojan horse incidents
  - establish a set of criteria to identify the proper level of response for various malicious code

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- type of malware
- point of entry and propagation method
- types of attacker tools placed in infected host by malware
- what hosts and networks are affected
- MONASH University how the impact is going to grow if not contained

# Malware Incident Response: Detection and Analysis<sup>11</sup> (continued)

- malware analysis
  - infect and monitor an isolated and clean test host to study behaviour of malware
  - virtualisation can be used where guest OS is isolated
- Example: cuckoo sandbox: open source automated malware analysis tool
- cuckoo host
  - runs the core components of the sandbox
  - manages the entire analysis process
- cuckoo guests
  - physical or virtual machines
  - run malware in an isolated environment
- malware analysis
  - trace malware calls
  - monitor any file system activity of malware
  - dump memory of the malware process
  - monitor network traffic
  - screen shots taken during malware execution
  - memory dump of the machine



Cuckoo automated malware analysis architecture<sup>a</sup>

<sup>a</sup>What is cuckoo?

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### Malware Incident Response: Containment<sup>12</sup>

- two major components
  - stop malware propagation
  - prevent further damage to hosts

#### Containment through user participation:

- important in large scale incidents and in non-managed environments
- should have alternate mechanisms to inform users in case email is unavailable
- organisation should not rely (entirely) on this means for containing malware
  - not all users may be able to follow provided instructions successfully
  - users may ignore instructions and carry on performing their usual tasks

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# Malware Incident Response: Containment<sup>13</sup> (continued)

#### Containment through **automated detection**:

- antivirus is the preferred method of automated detection and removal
  - however often antivirus software fail to detect novel malware threats
  - signature updates should be tested before deployment
    - to avoid potential problems introduced with update itself
- content filtering:
  - email servers and clients, anti-spam software etc. configured to
    - block emails or email attachments with certain characteristics
    - only when malware has static characteristics
  - web content filtering
- network based IPS software
  - when deployed inline with malware signature can stop detected packets to reach hosts
  - can prevent incoming and outgoing attempts
  - the effectiveness depends on accuracy of the signature
- executable black-listing
  - OS capability to prevent execution of files named in a list

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# Malware Incident Response: Containment<sup>14</sup> (continued)

#### Containment through **disabling services**:

- when malware use a service extensively, containment may require disabling the service
  - shutting down service used by malware at application layer or network layer
  - service can be disabled at the perimeter of the network
  - portions of a service is disabled e.g. large mailing lists
- a service may provide a channel for infection
- need to understand the consequence of disabling a service
  - for instance the impact on other services
  - should have a list of dependencies between services

# Malware Incident Response: Containment<sup>15</sup> (continued)

#### Containment through **disabling connectivity**:

- placing restrictions on network connectivity
  - blocking all access of infected hosts to external network
    - prevent download of attacker rootkits
  - block access to local network
  - prevent propagation of malware to other hosts
  - more drastic measures: isolate subnets/groups of hosts
- can design the network to make containment through loss of connectivity easier
  - separate subnets for servers and workstations
    - in malware incidents targeting workstations the workstation subnet can be isolated
  - separation using VLAN
    - separate VLAN for infected hosts
    - NAS servers can place hosts in different VLAN based on security check lists

## Malware Incident Response: Eradication<sup>16</sup>

- primary goal is to remove malware
  - also require eradication of vulnerability used for infection
    - vulnerable service
    - unsecured file share
- antivirus software are most common tools for eradication
  - remote scanning is more effective
  - isolation of infected hosts may make it difficult to remotely access to update
    - placing infected hosts in separate VLAN can provide remote access to incident handlers
- may require user participation
- for may incidents rebuild of all infected hosts are required
  - rebuild and securing OS
  - restoring data from known-good backups
- rebuild any host if:
  - attacker gained administrative level access
  - administrative level access was available to anyone through a backdoor
  - system files were replaced by malware
  - host is unstable
- doubt about the nature and extent of the infection

## Malware Incident Response: Recovery<sup>17</sup>

- restoring functionality and data of infected hosts
- removing temporary containment measures
  - determining when to remove
  - during major incidents may require restoring connectivity while still going through eradication and recovery
     when vulnerable hosts are patched so restring network connectivity will not lead to new hosts being infected
- should consider worst-case scenarios requiring rebuild of thousands of hosts
  - who will perform the recovery tasks
  - estimating the time required
  - priority of recovery efforts

## Malware Incident Response: Post-incident Activities

- Changes to security policy
  - to prevent similar incidents
- Changes to awareness program
  - training for users
- Software reconfiguration
  - OS
  - application
- Malware detection software deployment
  - an incident provides justification to deploy malware detection if previously was not used
- Malware detection software reconfiguration
  - increasing signature update frequency
  - increase scope of monitoring
  - change default response to malware detection



#### References

The materials in this document are reproduced, at times without modification, from the following sources:

- NIST SP 800-83r1: Guide to Malware Incident Prevention and Handling for Desktops and Laptops
- cuckoo sandbox: open source automated malware analysis tool