

114學年度第1學期 組合語言與嵌入式系統 — Final Project

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班組別：資訊二乙-3

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一、程式說明

1. name.s 和 id.s 都是沿用期中 project 的程式碼

2. extern 的三大功用

1. 跨檔案使用函式或變數

2. 避免重複定義符號

3. 讓 linker 正確進行符號解析 (symbol resolution)

3.

本程式最核心的部分為雙層迴圈結構，用來逐一處理畫面中每一個像素點。外層迴圈 loopwidth 負責控制畫面在水平方向 (x 軸) 的掃描，從左至右依序處理每一個位置，內層迴圈 loopheight 則負責垂直方向 (y 軸) 的掃描，從上到下依序計算每個像素。

在每一組 x、y 座標中，程式會先計算對應的初始複數值，接著進入 while 迴圈進行 Julia Set 的反覆運算。此 while 迴圈會根據運算結果是否超出範圍或是否達到最大迭代次數來決定是否跳出。當跳出迴圈後，會依據剩餘的迭代次數計算顏色，並將結果寫入 frame buffer，完成單一像素的繪製。透過此多層迴圈設計，程式能完整繪製整張 Julia Set 圖形。

二、設計重點說明

1.

frame 宣告是 : int16_t (*frame)[FRAME_WIDTH]

代表它指向的是一個一個 row，每個 row 有 FRAME_WIDTH 個 int16_t

sizeof(int16_t) = 2 bytes

=> 用 LDRH / STRH

=> FRAME_WIDTH = 640, FRAME_HEIGHT = 480, int16_t = 2 bytes

=> frame 總大小 = $640 \times 480 \times 2 = 614,400$ bytes

=> 轉成 16 進位 : 0x96000

假設我 frame 起始位置為 0xBEF69498

$0xBEF69498 + 0x00096000 = 0xBFFF498$ (結束位置)

2. Operand2 格式 :

mov r5, #1

sub r2, r5, r2

add r0, r0, r5, lsl #1

cmp r5, #0

3. 非 Branch 指令的 Conditional Execution:

movge r5, #0

movne r6, #0

moveq r6, #0

4. 指定的必執行指令 :

adds lr, sp, pc (第四道) 直接加上去程式碼不會編譯錯誤

三、程式驗證結果

frame start

```
id.s main.c name.s drawJuliaSet.s
104 ldr r0, -0xffff
105 bic r11, r0, r11
106
107 @ .L1:           base + (y * 640 * 2) + (x * 2)
108 mov r0, r7          @ r0 = frame buffer base
109 mov r1, #1280        @ .L2:           1280 bytes (640 * 2)
110 mul r1, r1, r6        @ r1 = 1280 * y
111 add r0, r0, r1        @ base + offset_y
112 add r0, r0, r5, lsl #1    @ + (x * 2)
113 strh r11, [r0]        @ .L3:           loopheight
114
115 add r6, r6, #1
116 b loopheight
```

frame end

```
id.s main.c name.s drawJuliaSet.s
109 mov r1, #1280        @ .L1:           1280 bytes (640 * 2)
110 mul r1, r1, r6        @ r1 = 1280 * y
111 add r0, r0, r1        @ base + offset_y
112 add r0, r0, r5, lsl #1    @ + (x * 2)
113 strh r11, [r0]        @ .L2:           loopheight
114
115 add r6, r6, #1
116 b loopheight
117
118 endloopheight:
119 add r5, r5, #1
120 b loopwidth
```

name.s function start

The screenshot shows the Code::Blocks IDE interface with the QEMU debugger. The assembly code for 'name.s' is displayed in the main editor window. A red dot marks the instruction at address 0x892c, which is the start of the 'name()' function. The CPU Registers window shows various寄存器 (Registers) with their hex and integer values. The Memory window shows the memory dump starting at address 0x892c.

```
int fd;
printf("Function1: Name\n");
name();
printf("Function2: ID\n");
id();
//Dummy printout. Please refer to the specification of Project 1.
printf("Main Function:\n");
printf("*****Dont Allxxxxxx\n");
```

name function end

The screenshot shows the Code::Blocks IDE interface with the QEMU debugger. The assembly code for 'name.s' is displayed in the main editor window. A red dot marks the instruction at address 0x8930, which is the end of the 'name()' function. The CPU Registers window shows various寄存器 (Registers) with their hex and integer values. The Memory window shows the memory dump starting at address 0x8930.

```
int fd;
printf("Function1: Name\n");
name();
printf("Function2: ID\n");
id();
//Dummy printout. Please refer to the specification of Project 1.
printf("Main Function:\n");
printf("*****Dont Allxxxxxx\n");
```

id.s function start

The screenshot shows the Code::Blocks IDE interface with the QEMU debugger. The assembly code for `id.s` is displayed in the main window:

```
int fd;
printf( "Function1: Name\n" );
name();
printf( "Function2: ID\n" );
id();
//Dummy printout. Please refer to the specification of Project 1.
printf( "Main Function:\n" );
printf( "*****Point All*****\n" );
```

The CPU Registers window shows the following values:

Reg#	Hex	Integer
r10	0x40026000	107389747
r11	0xbefff4b4	320444536
r12	0x0	0
sp	0xbef69490	320383092
lr	0x4020b8b4	107588626
pc	0x8938	3512
cpsr	0x60000010	161061275

The Memory window shows the memory starting at address 0x8938:

Address	Bytes
0x8938	32

Output:

```
0x8938: Be ff ff eb fc 01 9f e5|05 ff ff eb f8 01 9f e5 .xxxxxx.yyyy.aaaa
0x8948: 03 ff ff eb f4 31 9f e5|03 00 a0 e1 f0 11 9f e5 .yyyy01xxxxx.. aa.xxxx
```

id function end

The screenshot shows the Code::Blocks IDE interface with the QEMU debugger. The assembly code for `id.s` has been modified:

```
int fd;
printf( "Function1: Name\n" );
name();
printf( "Function2: ID\n" );
id();
//Dummy printout. Please refer to the specification of Project 1.
printf( "Main Function:\n" );
printf( "*****Point All*****\n" );
```

The CPU Registers window shows the following values:

Reg#	Hex	Integer
r10	0x40026000	107389747
r11	0xbefff4b4	320444536
r12	0x0	0
sp	0xbef69490	320383092
lr	0x893c	3513
pc	0x893c	3513
cpsr	0x60000010	161061275

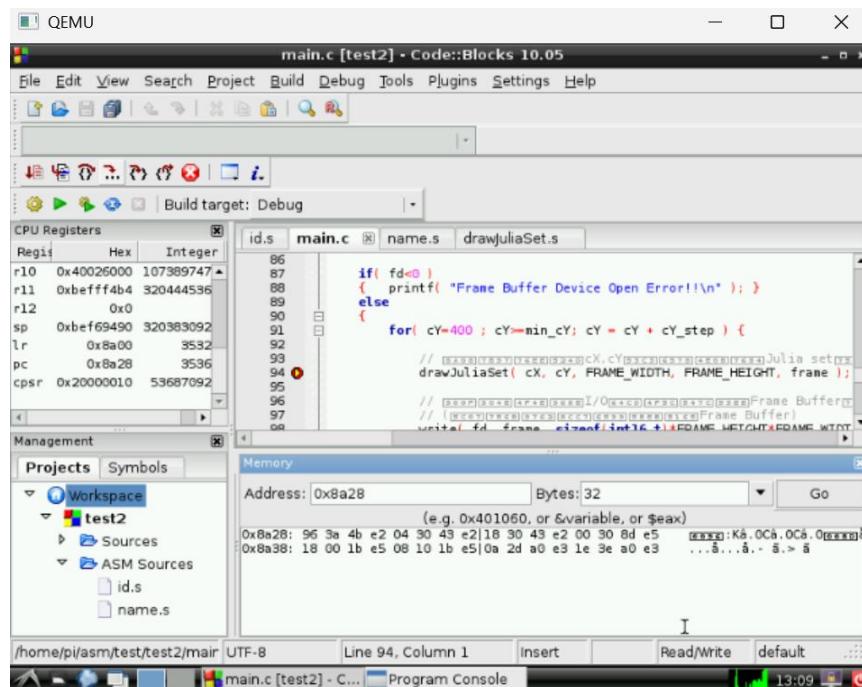
The Memory window shows the memory starting at address 0x893c:

Address	Bytes
0x893c	32

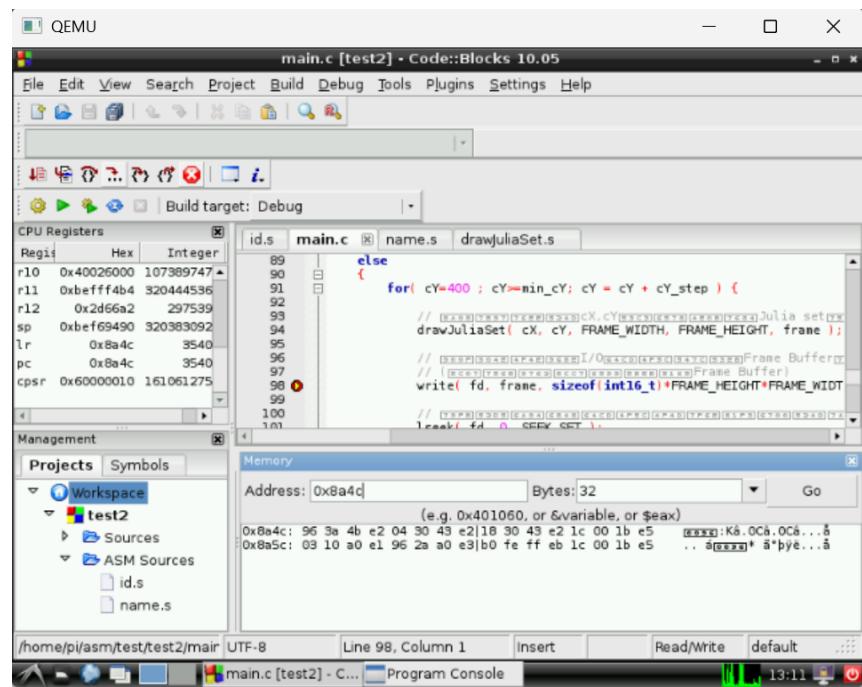
Output:

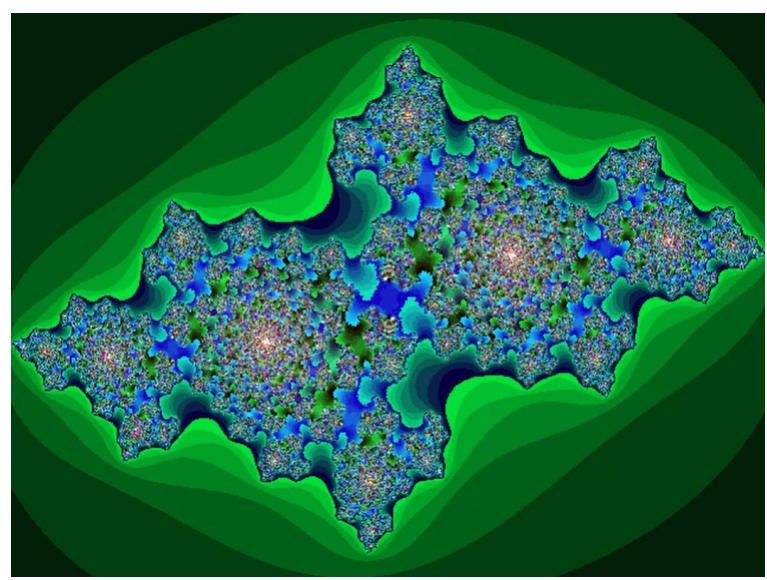
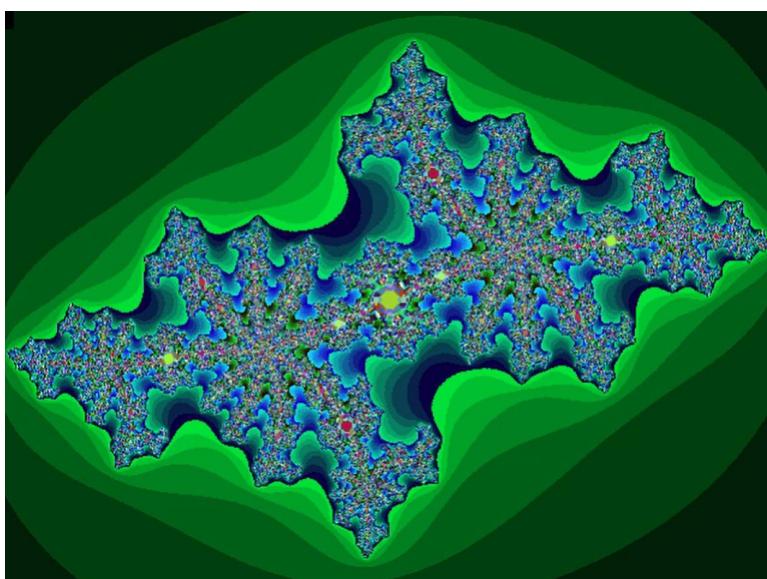
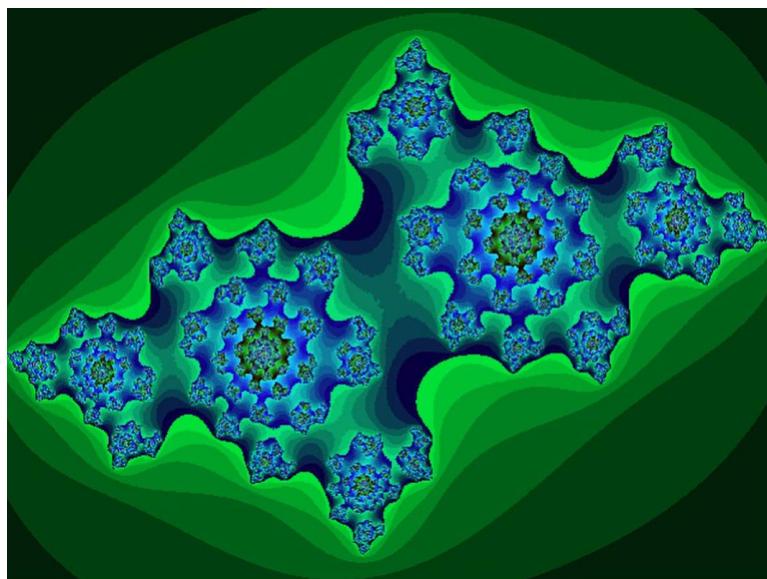
```
0x893c: fc 01 9f e5 05 ff ff eb|f8 01 9f e5 03 ff ff eb 0.xxxx.yyyy.aaaa
0x894c: f4 31 9f e5 03 00 a0 e1 f0 11 9f e5 f9 fe ff eb 01xxxxx.. aa.xxxx
```

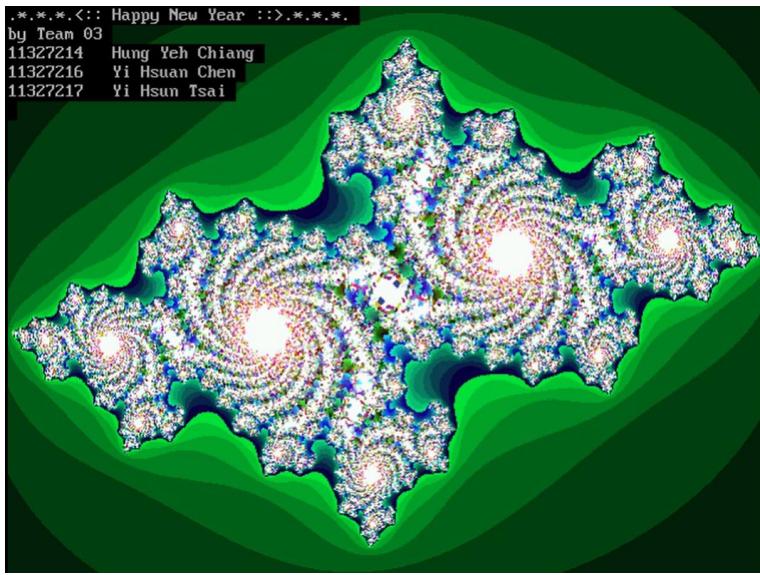
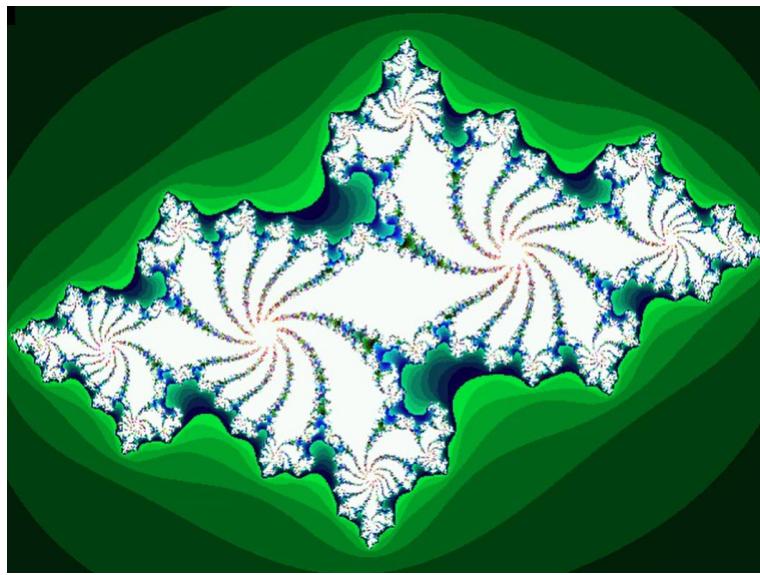
drawJuliaSet function start



drawJuliaSet function end







四、心得感想

這次的 function 實作相較上一次明顯困難許多，不僅程式中包含更多層的迴圈與條件判斷，也花一些時間在完成指令要求。在組合語言中，迴圈與流程控制需要自行處理每一次比較與跳轉，只要有一個細節出錯就會影響整體結果。實作過程中，常需要反覆測試與除錯，確認暫存器的數值是否正確。

五、各組員分工方式與負責項目

共同完成