

# MIT 6.035

## Top-Down Parsing

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# Orientation

- Language specification
  - Lexical structure – regular expressions
  - Syntactic structure – grammar
- This Lecture - recursive descent parsers
  - Code parser as set of mutually recursive procedures
  - Structure of program matches structure of grammar

# Starting Point

- Assume lexical analysis has produced a sequence of tokens
  - Each token has a type and value
  - Types correspond to terminals
  - Values to contents of token read in
- Examples
  - Int 549 – integer token with value 549 read in
  - if - if keyword, no need for a value
  - AddOp + - add operator, value +

# Basic Approach

- Start with Start symbol
- Build a leftmost derivation
  - If leftmost symbol is nonterminal, choose a production and apply it
  - If leftmost symbol is terminal, match against input
  - If all terminals match, have found a parse!
  - Key: find correct productions for nonterminals

# Graphical Illustration of Leftmost Derivation

Sentential Form



# Grammar for Parsing Example

*Start*  $\rightarrow$  *Expr*

*Expr*  $\rightarrow$  *Expr* + *Term*

*Expr*  $\rightarrow$  *Expr* - *Term*

*Expr*  $\rightarrow$  *Term*

*Term*  $\rightarrow$  *Term* \* Int

*Term*  $\rightarrow$  *Term* / Int

*Term*  $\rightarrow$  Int

- Set of tokens is  $\{ +, -, *, /, \text{Int} \}$ , where  $\text{Int} = [0-9][0-9]^*$
- For convenience, may represent each Int n token by n

# Parsing Example

Parse  
Tree

*Start*



*Current Position in Parse Tree*

Remaining Input

2-2\*2

Sentential Form

*Start*

# Parsing Example

Parse  
Tree

*Start*



*Expr*



*Current Position in Parse Tree*

Remaining Input

2-2\*2

Sentential Form

*Expr*

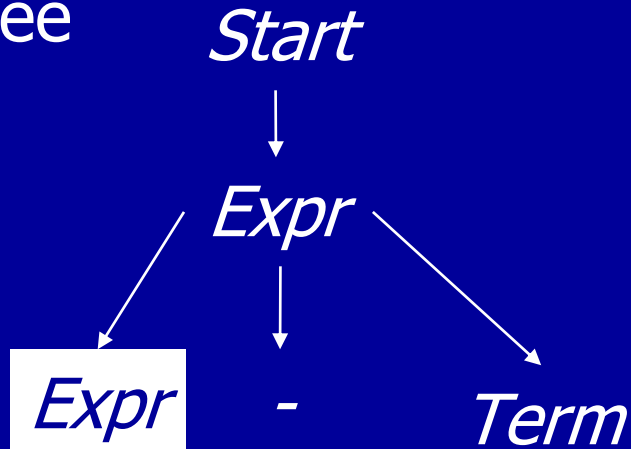
Applied Production

*Start*  $\rightarrow$  *Expr*



# Parsing Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

*Expr - Term*

*Expr*  $\rightarrow$  *Expr* + *Term*

*Expr*  $\rightarrow$  *Expr* - *Term*

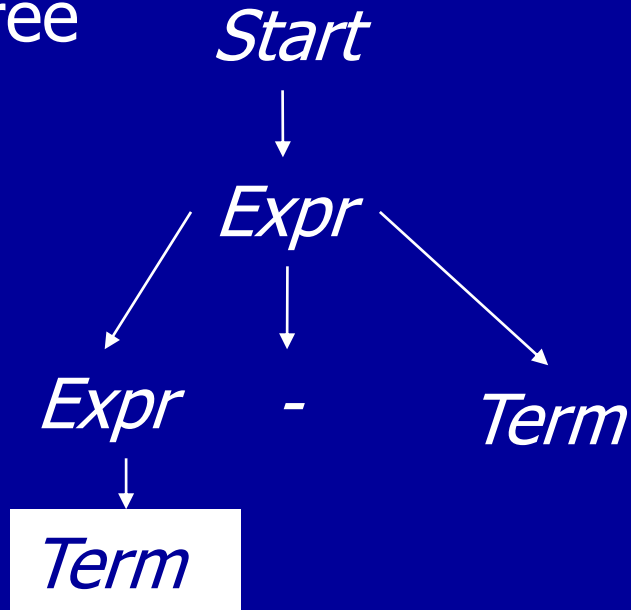
*Expr*  $\rightarrow$  *Term*

Applied Production

*Expr*  $\rightarrow$  *Expr* - *Term*

# Parsing Example

Parse  
Tree



$Expr \rightarrow Expr + Term$

$Expr \rightarrow Expr - Term$

$Expr \rightarrow Term$

Remaining Input

2-2\*2

Sentential Form

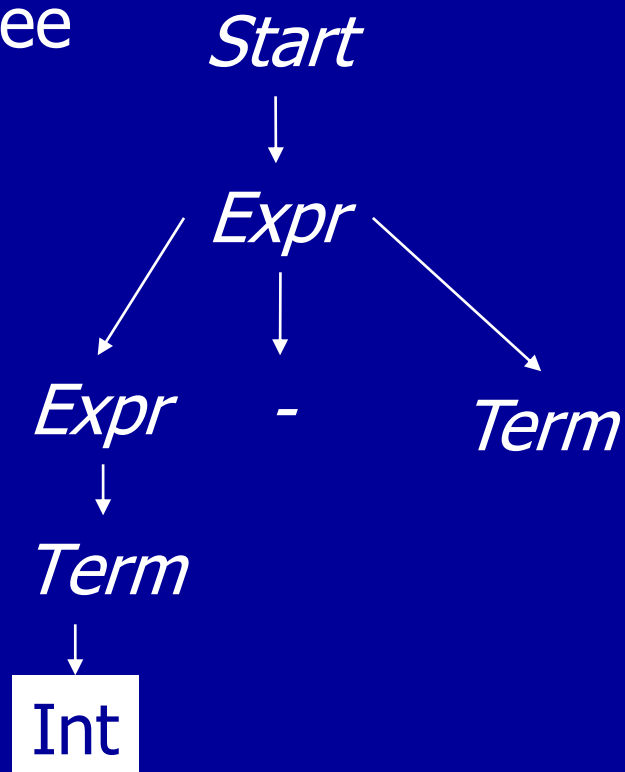
$Term - Term$

Applied Production

$Expr \rightarrow Term$

# Parsing Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

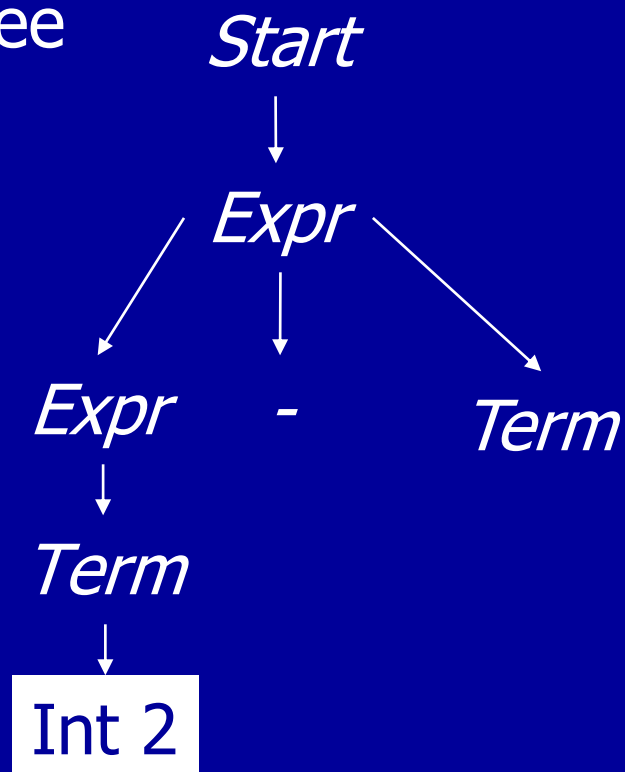
Int - Term

Applied Production

$Term \rightarrow \text{Int}$

# Parsing Example

Parse  
Tree



Match  
Input  
Token!

Remaining Input

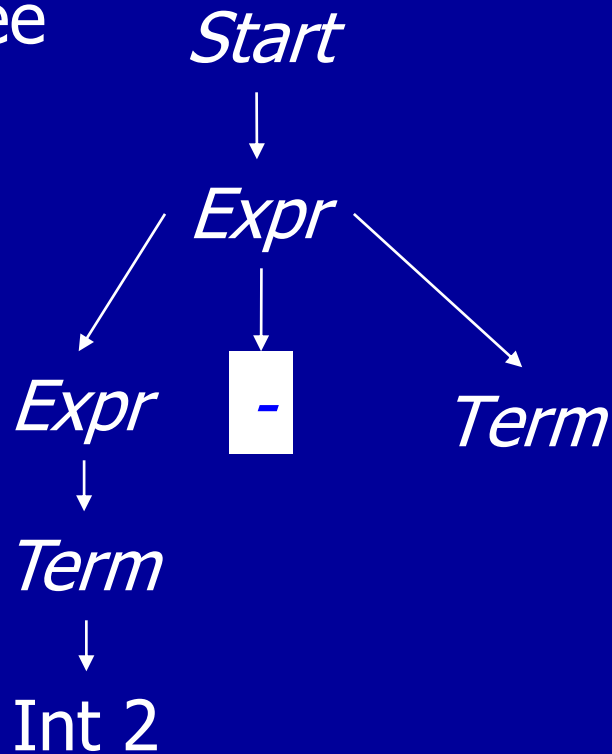
2-2\*2

Sentential Form

2 - *Term*

# Parsing Example

Parse  
Tree



Match  
Input  
Token!

Remaining Input

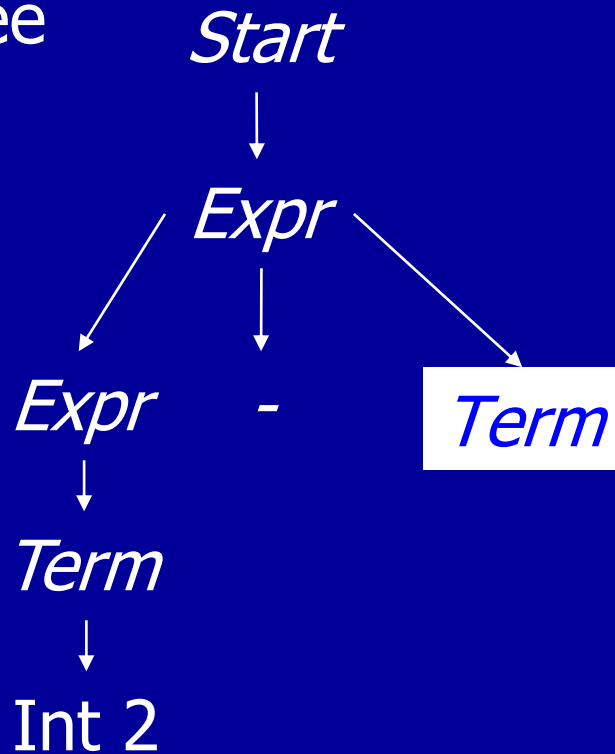
-2\*2

Sentential Form

2 - Term

# Parsing Example

Parse  
Tree



Match  
Input  
Token!

Remaining Input

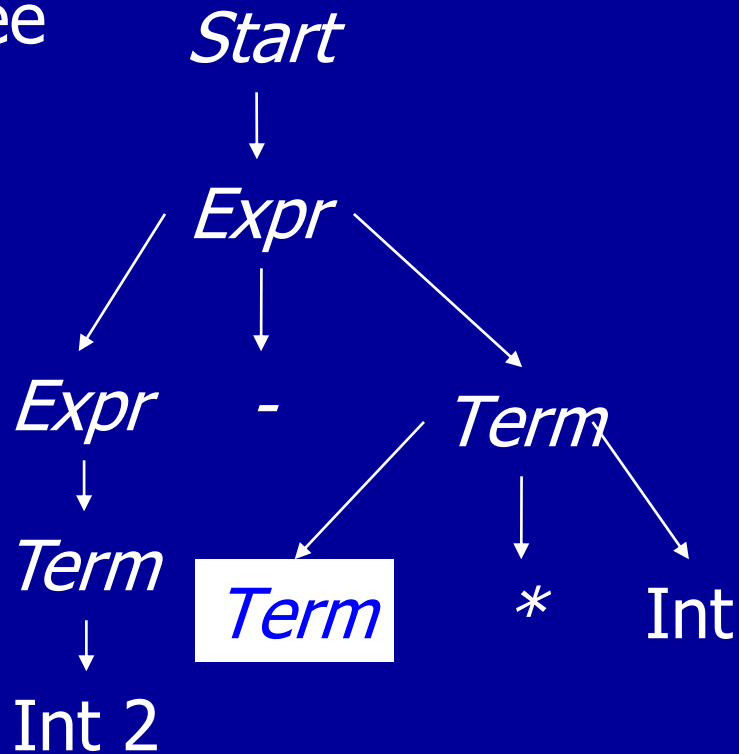
2\*2

Sentential Form

2 - Term

# Parsing Example

Parse  
Tree



Remaining Input

2\*2

Sentential Form

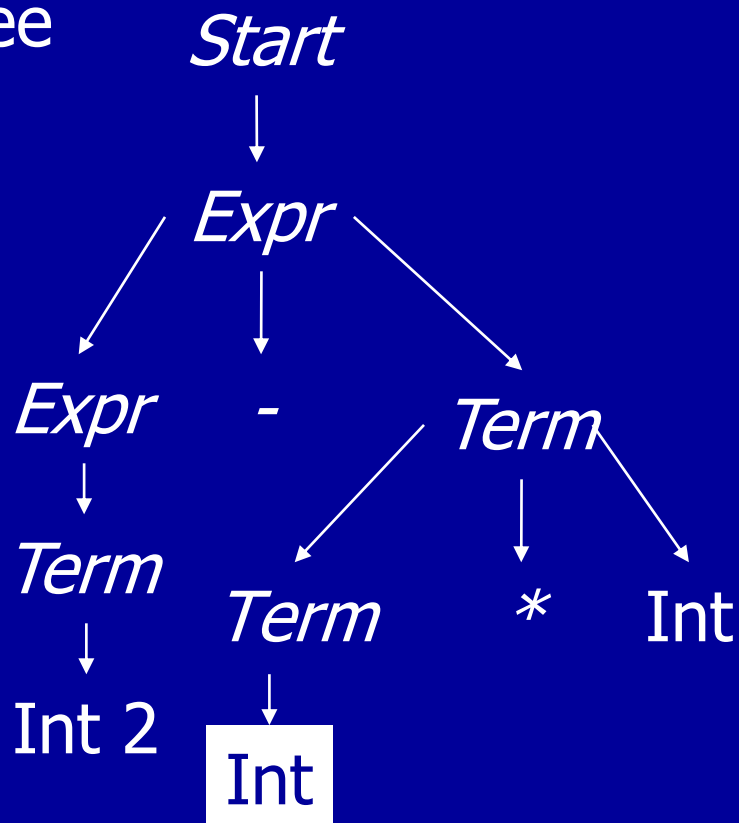
2 - *Term*\*Int

Applied Production

*Term* → *Term* \* Int

# Parsing Example

Parse  
Tree



Remaining Input

2\*2

Sentential Form

2 - Int \* Int

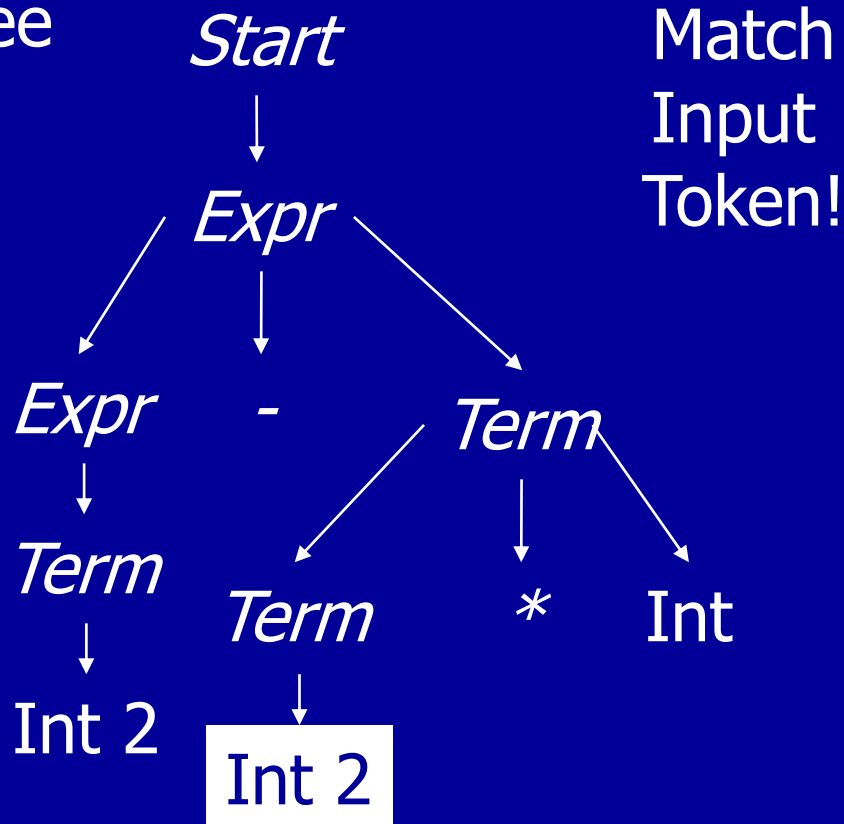
Applied Production

$Term \rightarrow Int$



# Parsing Example

Parse  
Tree



Remaining Input

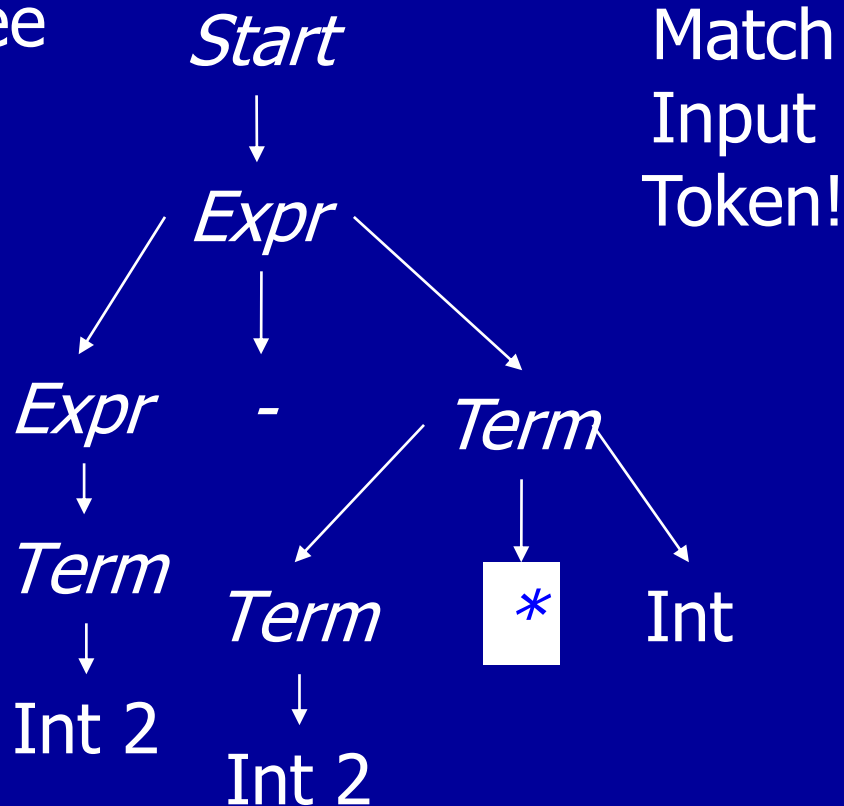
$2 * 2$

Sentential Form

$2 - 2 * \text{Int}$

# Parsing Example

Parse  
Tree



Match  
Input  
Token!

Remaining Input

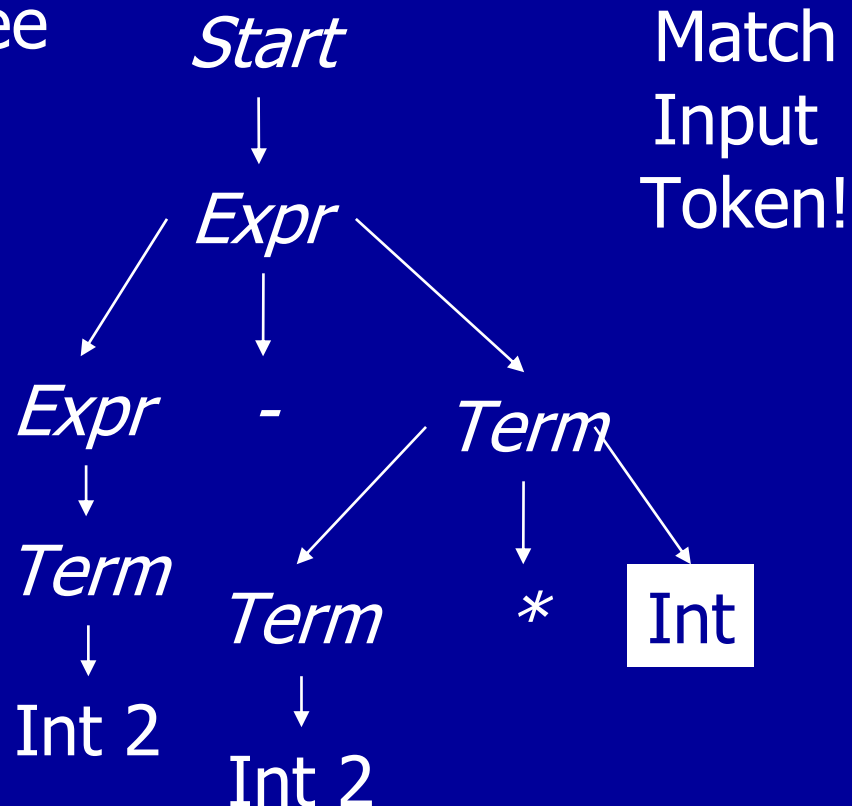
\*2

Sentential Form

2 - 2 \* Int

# Parsing Example

Parse  
Tree



Remaining Input

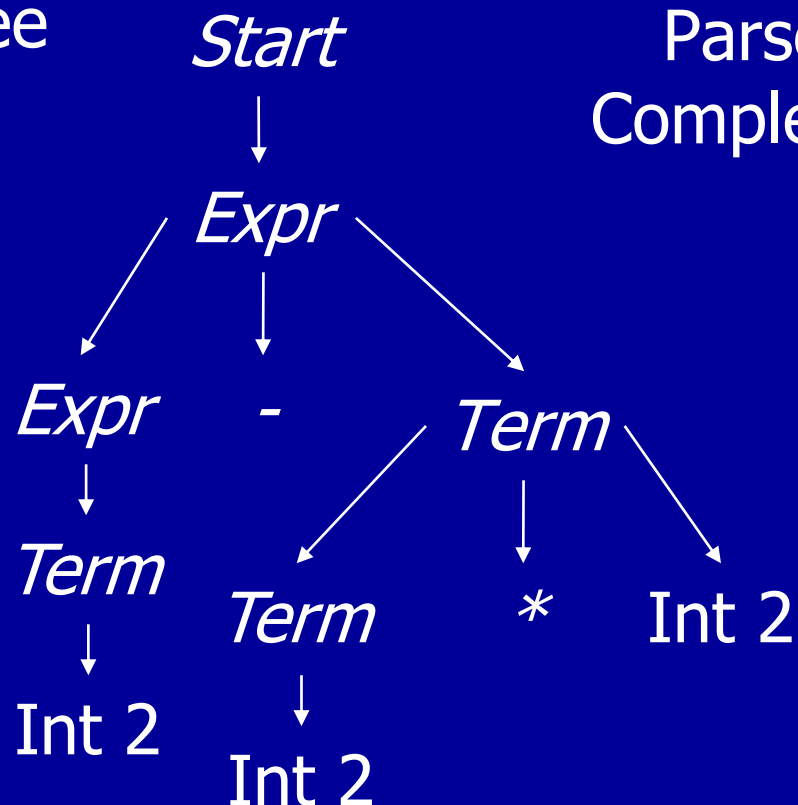
2

Sentential Form

2 - 2 \* Int

# Parsing Example

Parse  
Tree



Parse  
Complete!

Remaining Input

2

Sentential Form

2 - 2 \* 2

# Summary

- Three Actions (Mechanisms)
  - Apply production to expand current nonterminal in parse tree
  - Match current terminal (consuming input)
  - Accept the parse as correct
- Parser generates preorder traversal of parse tree
  - visit parents before children
  - visit siblings from left to right

# Policy Problem

- Which production to use for each nonterminal?
- Classical Separation of Policy and Mechanism
- One Approach: Backtracking
  - Treat it as a search problem
  - At each choice point, try next alternative
  - If it is clear that current try fails, go back to previous choice and try something different
- General technique for searching
- Used a lot in classical AI and natural language processing (parsing, speech recognition)

# Backtracking Example

Parse  
Tree

*Start*

Remaining Input

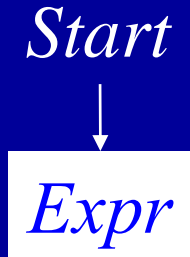
$2-2*2$

Sentential Form

*Start*

# Backtracking Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

*Expr*

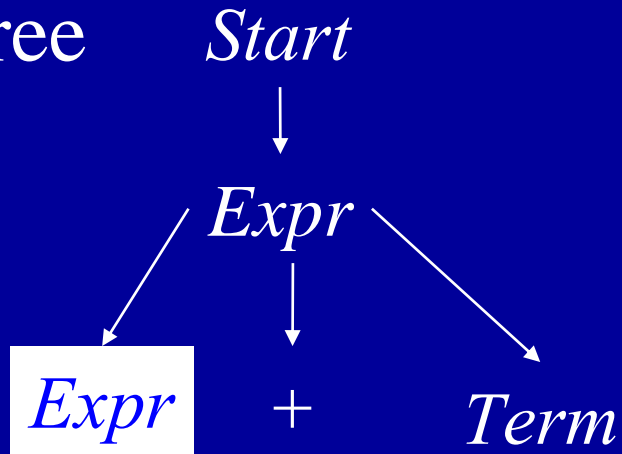
Applied Production

*Start*  $\rightarrow$  *Expr*



# Backtracking Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

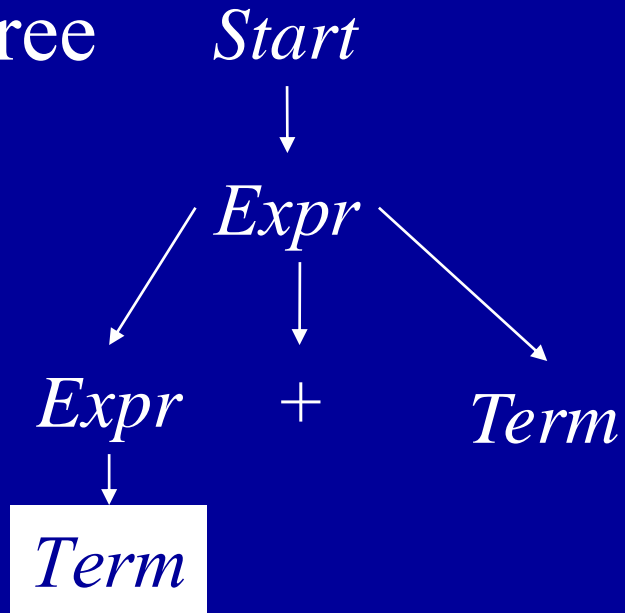
$Expr + Term$

Applied Production

$Expr \rightarrow Expr + Term$

# Backtracking Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

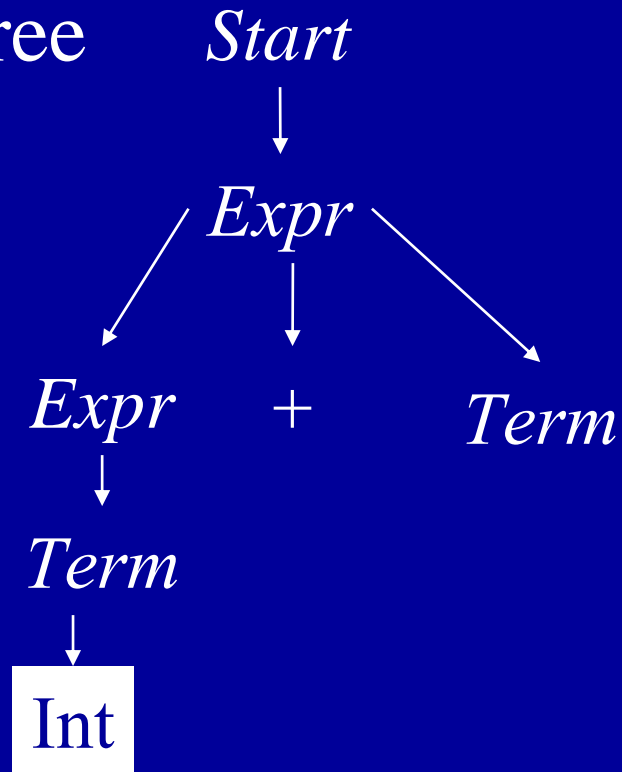
$Term + Term$

Applied Production

$Expr \rightarrow Term$

# Backtracking Example

Parse  
Tree



Match  
Input  
Token!

Remaining Input

2-2\*2

Sentential Form

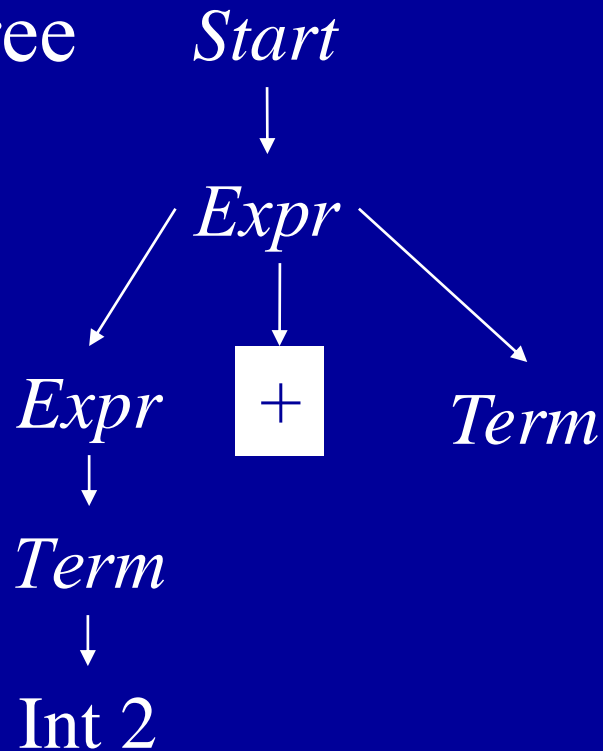
*Int* + *Term*

Applied Production

*Term*  $\rightarrow$  *Int*

# Backtracking Example

Parse  
Tree



Can't  
Match  
Input  
Token!

Remaining Input

-2\*2

Sentential Form

2 - Term

Applied Production

$Term \rightarrow Int$

# Backtracking Example

Parse  
Tree

*Start*



*Expr*

So  
Backtrack!

Remaining Input

2-2\*2

Sentential Form

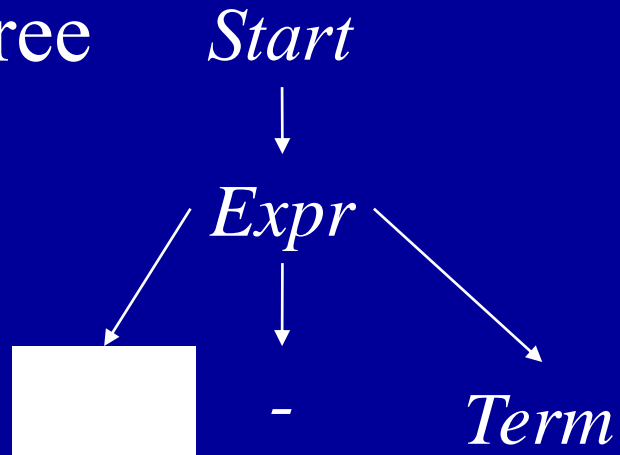
*Expr*

Applied Production

*Start*  $\rightarrow$  *Expr*

# Backtracking Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

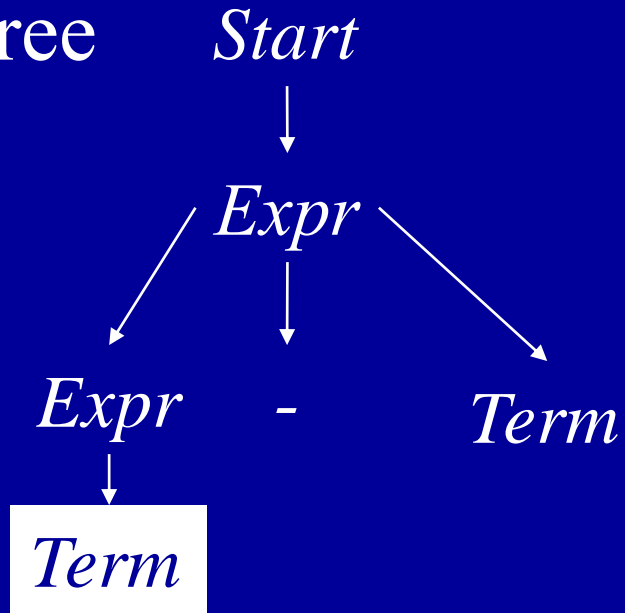
$Expr - Term$

Applied Production

$Expr \rightarrow Expr - Term$

# Backtracking Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

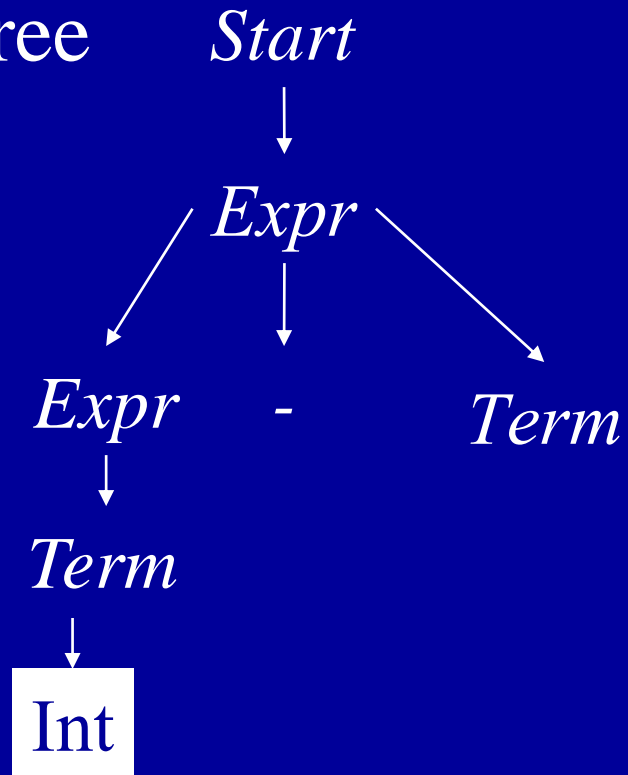
$Term - Term$

Applied Production

$Expr \rightarrow Term$

# Backtracking Example

Parse  
Tree



Remaining Input

2-2\*2

Sentential Form

Int - Term

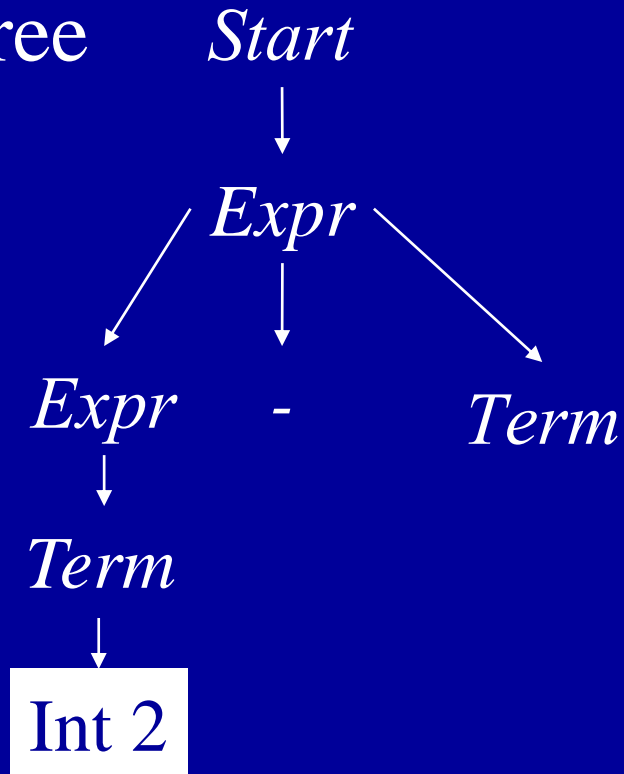
Applied Production

$Term \rightarrow Int$



# Backtracking Example

Parse  
Tree



Match  
Input  
Token!

Remaining Input

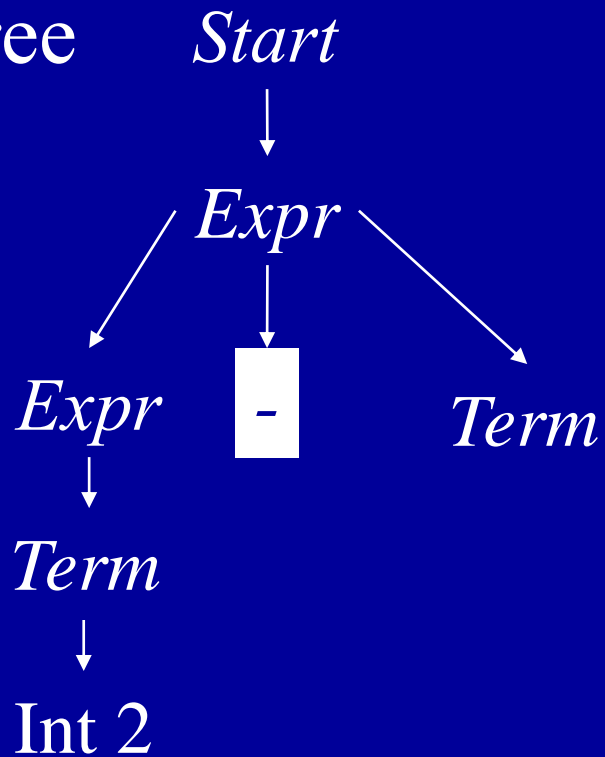
-2\*2

Sentential Form

2 - *Term*

# Backtracking Example

Parse  
Tree



Match  
Input  
Token!

Remaining Input

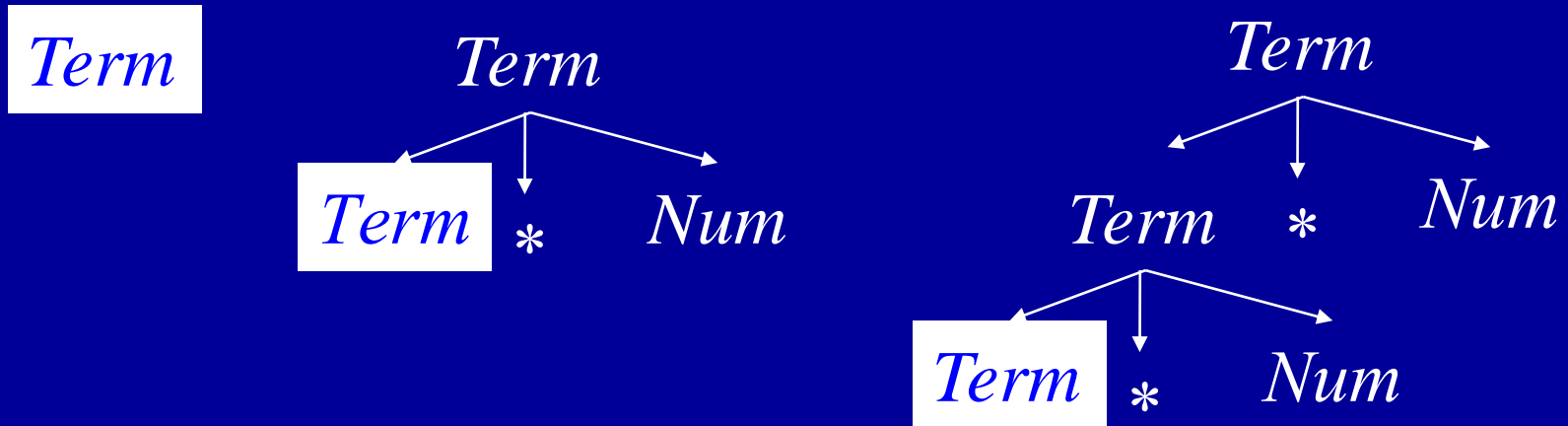
$2 * 2$

Sentential Form

$2 - \text{Term}$

# Left Recursion + Top-Down Parsing = Infinite Loop

- Example Production:  $Term \rightarrow Term * Num$
- Potential parsing steps:

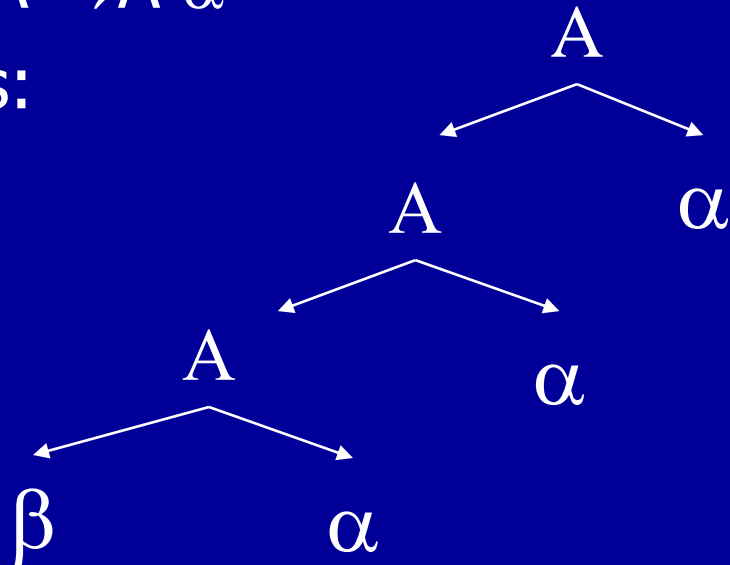


# General Search Issues

- Three components
  - Search space (parse trees)
  - Search algorithm (parsing algorithm)
  - Goal to find (parse tree for input program)
- Would like to (but can't always) ensure that
  - Find goal (hopefully quickly) if it exists
  - Search terminates if it does not
- Handled in various ways in various contexts
  - Finite search space makes it easy
  - Exploration strategies for infinite search space
  - Sometimes one goal more important (model checking)
- For parsing, hack grammar to remove left recursion

# Eliminating Left Recursion

- Start with productions of form
  - $A \rightarrow A \alpha$
  - $A \rightarrow \beta$
  - $\alpha, \beta$  sequences of terminals and nonterminals that do not start with  $A$
- Repeated application of  $A \rightarrow A \alpha$  builds parse tree like this:



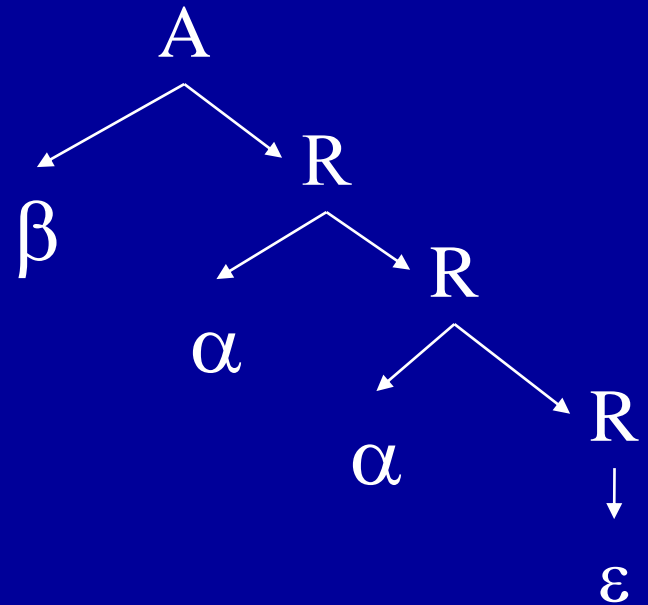
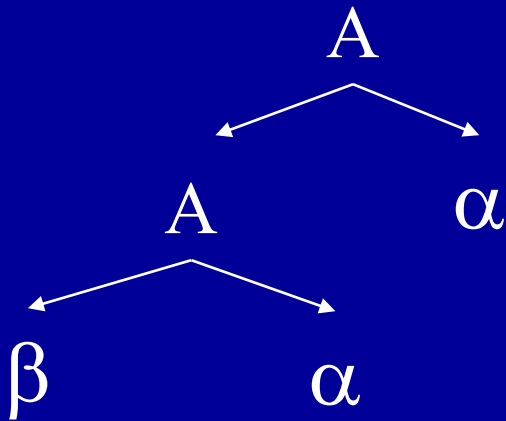
# Eliminating Left Recursion

- Replacement productions

- $A \rightarrow A \alpha$        $A \rightarrow \beta R$        $R$  is a new nonterminal
- $A \rightarrow \beta$        $R \rightarrow \alpha R$
- $\quad\quad\quad R \rightarrow \varepsilon$

New Parse Tree

Old Parse Tree



# Hacked Grammar

Original Grammar  
Fragment

$Term \rightarrow Term * Int$

$Term \rightarrow Term / Int$

$Term \rightarrow Int$

New Grammar Fragment

$Term \rightarrow Int \ Term'$

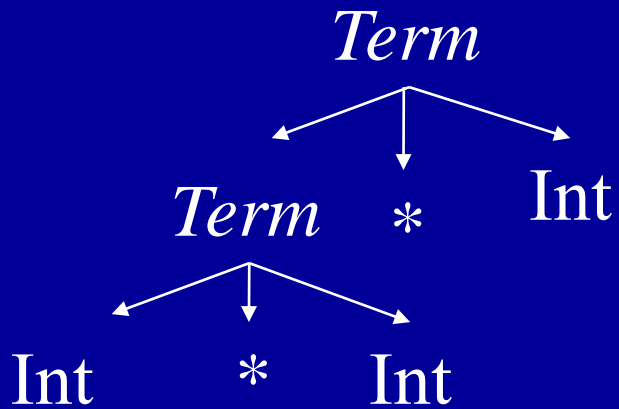
$Term' \rightarrow *Int \ Term'$

$Term' \rightarrow /Int \ Term'$

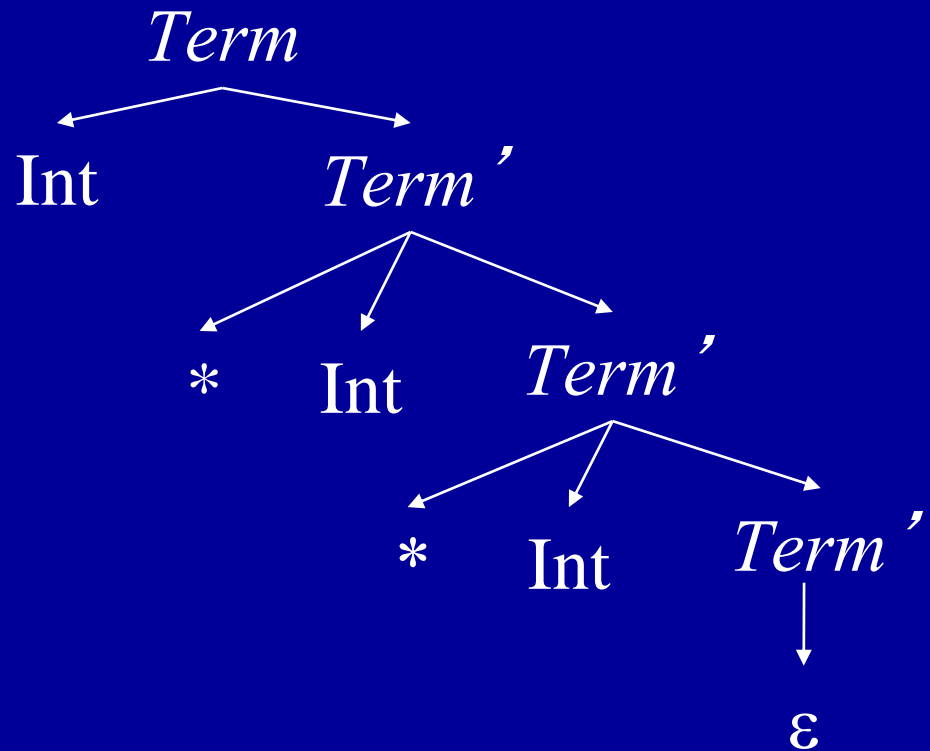
$Term' \rightarrow \varepsilon$

# Parse Tree Comparisons

Original Grammar



New Grammar





# Eliminating Left Recursion

- Changes search space exploration algorithm
  - Eliminates direct infinite recursion
  - But grammar less intuitive
- Sets things up for predictive parsing

# Predictive Parsing

- Alternative to backtracking
- Useful for programming languages, which can be designed to make parsing easier
- Basic idea
  - Look ahead in input stream
  - Decide which production to apply based on next tokens in input stream
  - We will use one token of lookahead

# Predictive Parsing Example Grammar

*Start*  $\rightarrow$  *Expr*

*Expr*  $\rightarrow$  *Term Expr'*

*Expr'*  $\rightarrow$  + *Term Expr'*

*Expr'*  $\rightarrow$  - *Term Expr'*

*Expr'*  $\rightarrow$   $\epsilon$

*Term*  $\rightarrow$  Int *Term'*

*Term'*  $\rightarrow$  \* Int *Term'*

*Term'*  $\rightarrow$  / Int *Term'*

*Term'*  $\rightarrow$   $\epsilon$

# Choice Points

- Assume  $Term'$  is current position in parse tree
- Have three possible productions to apply
$$Term' \rightarrow *Int\ Term'$$
$$Term' \rightarrow /Int\ Term'$$
$$Term' \rightarrow \varepsilon$$
- Use next token to decide
  - If next token is  $*$ , apply  $Term' \rightarrow *Int\ Term'$
  - If next token is  $/$ , apply  $Term' \rightarrow /Int\ Term'$
  - Otherwise, apply  $Term' \rightarrow \varepsilon$

# Predictive Parsing + Hand Coding = Recursive Descent Parser

- One procedure per nonterminal  $NT$ 
  - Productions  $NT \rightarrow \beta_1, \dots, NT \rightarrow \beta_n$
  - Procedure examines the current input symbol  $T$  to determine which production to apply
    - If  $T \in \text{First}(\beta_k)$
    - Apply production  $k$
    - Consume terminals in  $\beta_k$  (check for correct terminal)
    - Recursively call procedures for nonterminals in  $\beta_k$
  - Current input symbol stored in global variable `token`
- Procedures return
  - true if parse succeeds
  - false if parse fails

# Example

Boolean Term()

```
if (token = Int n) token = NextToken(); return(TermPrime())  
else return(false)
```

Boolean TermPrime()

```
if (token = *)  
    token = NextToken();  
    if (token = Int n) token = NextToken(); return(TermPrime())  
    else return(false)
```

```
else if (token = /)  
    token = NextToken();  
    if (token = Int n) token = NextToken(); return(TermPrime())  
    else return(false)
```

```
else return(true)
```

$Term \rightarrow \text{Int } Term'$

$Term' \rightarrow * \text{Int } Term'$

$Term' \rightarrow / \text{Int } Term'$

$Term' \rightarrow \varepsilon$

# Multiple Productions With Same Prefix in RHS

- Example Grammar

$NT \rightarrow \text{if then}$

$NT \rightarrow \text{if then else}$

- Assume  $NT$  is current position in parse tree, and  $\text{if}$  is the next token
- Unclear which production to apply
  - Multiple  $k$  such that  $T \in \text{First}(\beta_k)$
  - $\text{if} \in \text{First}(\text{if then})$
  - $\text{if} \in \text{First}(\text{if then else})$

# Solution: Left Factor the Grammar

- New Grammar Factors Common Prefix Into Single Production

$NT \rightarrow \text{if then } NT'$

$NT' \rightarrow \text{else}$

$NT' \rightarrow \varepsilon$

- No choice when next token is if!
- All choices have been unified in one production.



# Nonterminals

- What about productions with nonterminals?

$$NT \rightarrow NT_1 \alpha_1$$

$$NT \rightarrow NT_2 \alpha_2$$

- Must choose based on possible first terminals that  $NT_1$  and  $NT_2$  can generate
- What if  $NT_1$  or  $NT_2$  can generate  $\varepsilon$ ?
  - Must choose based on  $\alpha_1$  and  $\alpha_2$

# $NT$ derives $\varepsilon$

- Two rules
  - $NT \rightarrow \varepsilon$  implies  $NT$  derives  $\varepsilon$
  - $NT \rightarrow NT_1 \dots NT_n$  and for all  $1 \leq i \leq n$   $NT_i$  derives  $\varepsilon$  implies  $NT$  derives  $\varepsilon$

# Fixed Point Algorithm for Derives $\varepsilon$

for all nonterminals  $NT$

    set  $NT$  derives  $\varepsilon$  to be false

for all productions of the form  $NT \rightarrow \varepsilon$

    set  $NT$  derives  $\varepsilon$  to be true

while (some  $NT$  derives  $\varepsilon$  changed in last iteration)

    for all productions of the form  $NT \rightarrow NT_1 \dots NT_n$

        if (for all  $1 \leq i \leq n$   $NT_i$  derives  $\varepsilon$ )

            set  $NT$  derives  $\varepsilon$  to be true

# First( $\beta$ )

- $T \in \text{First}(\beta)$  if  $T$  can appear as the first symbol in a derivation starting from  $\beta$ 
  - 1)  $T \in \text{First}(T)$
  - 2)  $\text{First}(S) \subseteq \text{First}(S\beta)$
  - 3)  $NT$  derives  $\varepsilon$  implies  $\text{First}(\beta) \subseteq \text{First}(NT\beta)$
  - 4)  $NT \rightarrow S\beta$  implies  $\text{First}(S\beta) \subseteq \text{First}(NT)$
- Notation
  - $T$  is a terminal,  $NT$  is a nonterminal,  $S$  is a terminal or nonterminal, and  $\beta$  is a sequence of terminals or nonterminals

# Rules + Request Generate System of Subset Inclusion Constraints

Grammar

$Term' \rightarrow *Int\ Term'$

$Term' \rightarrow /Int\ Term'$

$Term' \rightarrow \varepsilon$

Rules

- 1)  $T \in \text{First}(T)$
- 2)  $\text{First}(S) \subseteq \text{First}(S\beta)$
- 3)  $NT$  derives  $\varepsilon$  implies  
 $\text{First}(\beta) \subseteq \text{First}(NT\beta)$
- 4)  $NT \rightarrow S\beta$  implies  
 $\text{First}(S\beta) \subseteq \text{First}(NT)$

Request: What is  $\text{First}(Term')$ ?

Constraints

$\text{First}(*Num\ Term') \subseteq \text{First}(Term')$

$\text{First}(/Num\ Term') \subseteq \text{First}(Term')$

$\text{First}(*) \subseteq \text{First}(*Num\ Term')$

$\text{First}(/) \subseteq \text{First}(/Num\ Term')$

$* \in \text{First}(*)$

$/ \in \text{First}(/)$

# Constraint Propagation Algorithm

## Constraints

$$\text{First}( * \text{ Num } Term' ) \subseteq \text{First}( Term' )$$

$$\text{First}( / \text{ Num } Term' ) \subseteq \text{First}( Term' )$$

$$\text{First}( * ) \subseteq \text{First}( * \text{ Num } Term' )$$

$$\text{First}( / ) \subseteq \text{First}( / \text{ Num } Term' )$$

$$* \in \text{First}( * )$$

$$/ \in \text{First}( / )$$

## Solution

$$\text{First}( Term' ) = \{ \}$$

$$\text{First}( * \text{ Num } Term' ) = \{ \}$$

$$\text{First}( / \text{ Num } Term' ) = \{ \}$$

$$\text{First}( * ) = \{ * \}$$

$$\text{First}( / ) = \{ / \}$$

Initialize Sets to  $\{ \}$

Propagate Constraints Until  
Fixed Point

# Constraint Propagation Algorithm

## Constraints

$\text{First}( * \text{ Num } \text{Term}' ) \subseteq \text{First}( \text{Term}' )$

$\text{First}( / \text{ Num } \text{Term}' ) \subseteq \text{First}( \text{Term}' )$

$\text{First}( * ) \subseteq \text{First}( * \text{ Num } \text{Term}' )$

$\text{First}( / ) \subseteq \text{First}( / \text{ Num } \text{Term}' )$

$* \in \text{First}( * )$

$/ \in \text{First}( / )$

## Solution

$\text{First}( \text{Term}' ) = \{ \}$

$\text{First}( * \text{ Num } \text{Term}' ) = \{ \}$

$\text{First}( / \text{ Num } \text{Term}' ) = \{ \}$

$\text{First}( * ) = \{ * \}$

$\text{First}( / ) = \{ / \}$

## Grammar

$\text{Term}' \rightarrow * \text{Int } \text{Term}'$

$\text{Term}' \rightarrow / \text{Int } \text{Term}'$

$\text{Term}' \rightarrow \varepsilon$

# Constraint Propagation Algorithm

## Constraints

$\text{First}( * \text{ Num } Term' ) \subseteq \text{First}( Term' )$

$\text{First}( / \text{ Num } Term' ) \subseteq \text{First}( Term' )$

$\text{First}( * ) \subseteq \text{First}( * \text{ Num } Term' )$

$\text{First}( / ) \subseteq \text{First}( / \text{ Num } Term' )$

$* \in \text{First}( * )$

$/ \in \text{First}( / )$

## Solution

$\text{First}( Term' ) = \{ \}$

$\text{First}( * \text{ Num } Term' ) = \{ * \}$

$\text{First}( / \text{ Num } Term' ) = \{ / \}$

$\text{First}( * ) = \{ * \}$

$\text{First}( / ) = \{ / \}$

## Grammar

$Term' \rightarrow * \text{Int } Term'$

$Term' \rightarrow / \text{Int } Term'$

$Term' \rightarrow \varepsilon$



# Constraint Propagation Algorithm

## Constraints

$\text{First}( * \text{ Num } Term' ) \subseteq \text{First}( Term' )$

$\text{First}( / \text{ Num } Term' ) \subseteq \text{First}( Term' )$

$\text{First}( * ) \subseteq \text{First}( * \text{ Num } Term' )$

$\text{First}( / ) \subseteq \text{First}( / \text{ Num } Term' )$

$* \in \text{First}( * )$

$/ \in \text{First}( / )$

## Solution

$\text{First}( Term' ) = \{ *, / \}$

$\text{First}( * \text{ Num } Term' ) = \{ * \}$

$\text{First}( / \text{ Num } Term' ) = \{ / \}$

$\text{First}( * ) = \{ * \}$

$\text{First}( / ) = \{ / \}$

## Grammar

$Term' \rightarrow * \text{Int } Term'$

$Term' \rightarrow / \text{Int } Term'$

$Term' \rightarrow \varepsilon$

# Constraint Propagation Algorithm

## Constraints

$\text{First}( * \text{ Num } Term' ) \subseteq \text{First}( Term' )$

$\text{First}( / \text{ Num } Term' ) \subseteq \text{First}( Term' )$

$\text{First}( * ) \subseteq \text{First}( * \text{ Num } Term' )$

$\text{First}( / ) \subseteq \text{First}( / \text{ Num } Term' )$

$* \in \text{First}( * )$

$/ \in \text{First}( / )$

## Solution

$\text{First}( Term' ) = \{ *, / \}$

$\text{First}( * \text{ Num } Term' ) = \{ * \}$

$\text{First}( / \text{ Num } Term' ) = \{ / \}$

$\text{First}( * ) = \{ * \}$

$\text{First}( / ) = \{ / \}$

## Grammar

$Term' \rightarrow * \text{Int } Term'$

$Term' \rightarrow / \text{Int } Term'$

$Term' \rightarrow \varepsilon$

# Building A Parse Tree

- Have each procedure return the section of the parse tree for the part of the string it parsed
- Use exceptions to make code structure clean

# Building Parse Tree In Example

Term()

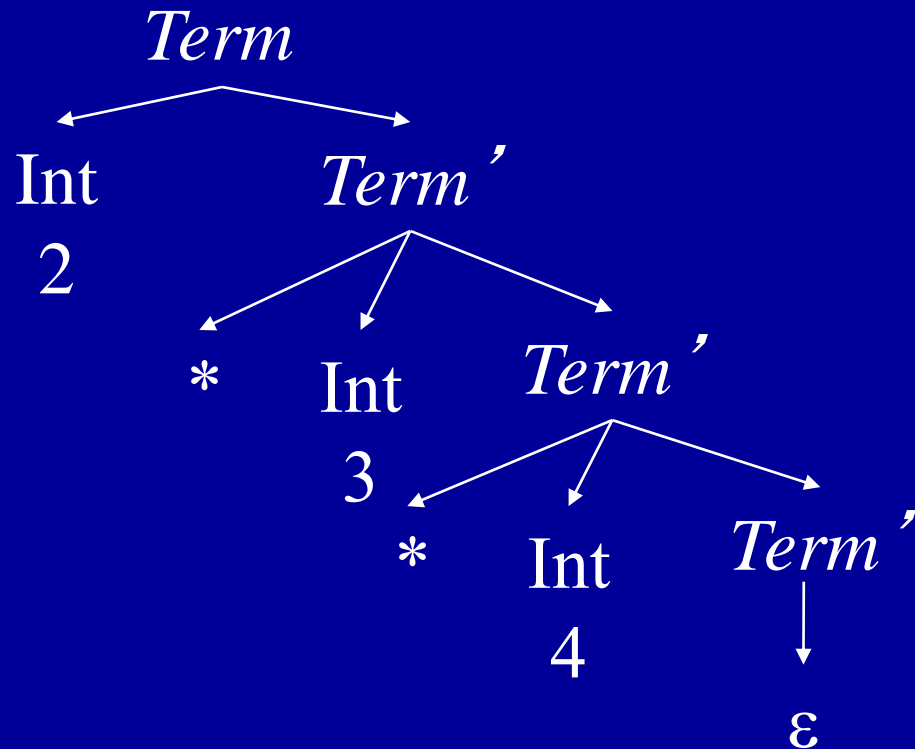
```
if (token = Int n)
    oldToken = token; token = NextToken();
    node = TermPrime();
    if (node == NULL) return oldToken;
    else return(new TermNode(oldToken, node);
else throw SyntaxError
```

TermPrime()

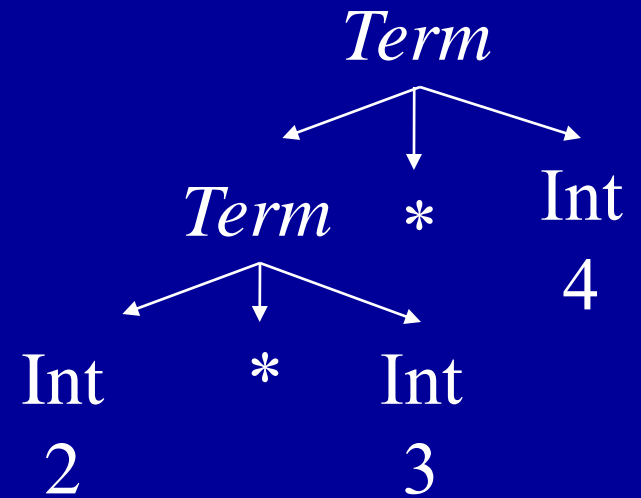
```
if (token = *) || (token = /)
    first = token; next = NextToken();
    if (next = Int n)
        token = NextToken();
        return(new TermPrimeNode(first, next, TermPrime()))
    else throw SyntaxError
else return(NULL)
```

# Parse Tree for $2*3*4$

Concrete  
Parse Tree



Desired  
Abstract  
Parse Tree



# Why Use Hand-Coded Parser?

- Why not use parser generator?
- What do you do if your parser doesn't work?
  - Recursive descent parser – write more code
  - Parser generator
    - Hack grammar
    - But if parser generator doesn't work, nothing you can do
- If you have complicated grammar
  - Increase chance of going outside comfort zone of parser generator
  - Your parser may NEVER work

# Bottom Line

- Recursive descent parser properties
  - Probably more work
  - But less risk of a disaster - you can almost always make a recursive descent parser work
  - May have easier time dealing with resulting code
    - Single language system
    - No need to deal with potentially flaky parser generator
    - No integration issues with automatically generated code
- If your parser development time is small compared to rest of project, or you have a really complicated language, use hand-coded recursive descent parser

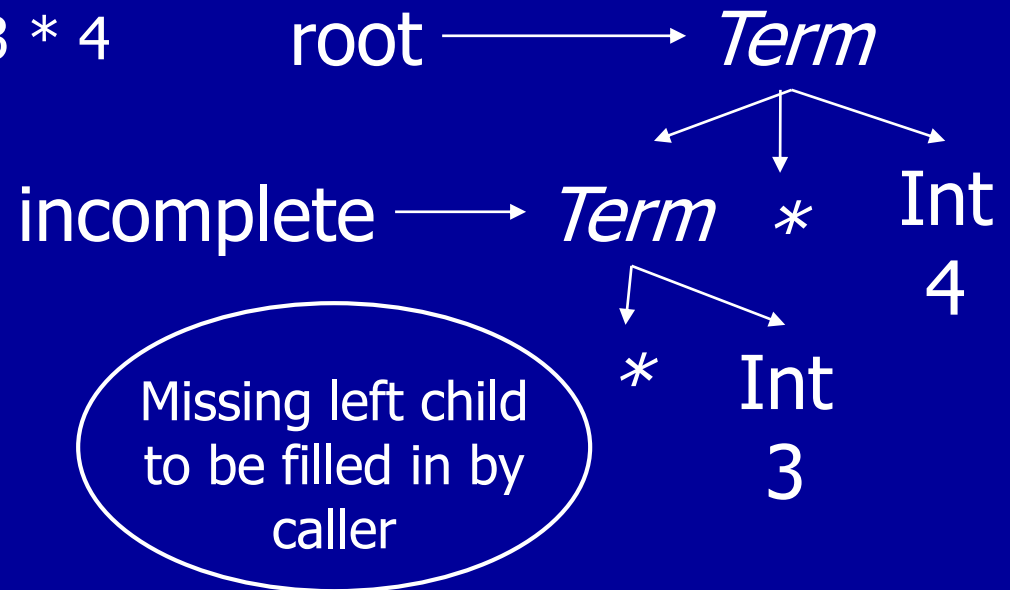
# Summary

- Top-Down Parsing
- Use Lookahead to Avoid Backtracking
- Parser is
  - Hand-Coded
  - Set of Mutually Recursive Procedures




# Direct Generation of Abstract Tree

- TermPrime builds an incomplete tree
  - Missing leftmost child
  - Returns root and incomplete node
- $(\text{root}, \text{incomplete}) = \text{TermPrime}()$ 
  - Called with token = \*
  - Remaining tokens = 3 \* 4




# Code for Term

Term()

```
if (token = Int n)   
    leftmostInt = token; token = NextToken();  
    (root, incomplete) = TermPrime();  
    if (root == NULL) return leftmostInt;  
    incomplete.leftChild = leftmostInt;  
    return root;  
else throw SyntaxError
```

Input to  
parse

2\*3\*4  


token  Int  
2

# Code for Term

Term()

```
if (token = Int n)
```

```
    leftmostInt = token; token = NextToken();
```

```
    (root, incomplete) = TermPrime();
```

```
    if (root == NULL) return leftmostInt;
```

```
    incomplete.leftChild = leftmostInt;
```

```
    return root;
```

```
else throw SyntaxError
```

Input to  
parse

2\*3\*4



token → Int  
2

# Code for Term

Term()

```
if (token = Int n)
    leftmostInt = token; token = NextToken();
    (root, incomplete) = TermPrime(); ←
    if (root == NULL) return leftmostInt;
    incomplete.leftChild = leftmostInt;
    return root;
else throw SyntaxError
```

Input to  
parse

2\*3\*4  
↑

token → Int  
2

# Code for Term

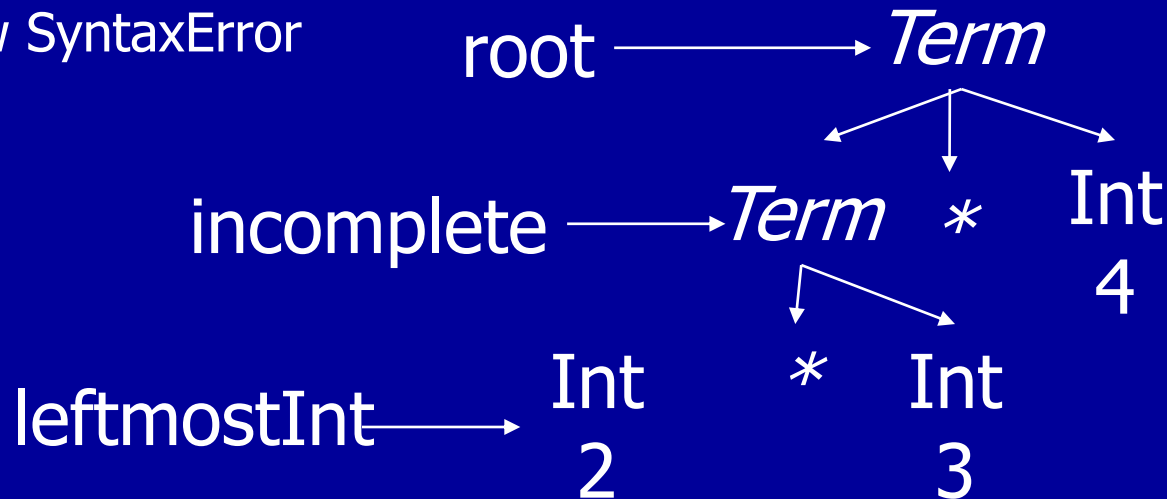
Term()

```
if (token = Int n)
    leftmostInt = token; token = NextToken();
    (root, incomplete) = TermPrime();
    if (root == NULL) return leftmostInt;
    incomplete.leftChild = leftmostInt;
    return root;
```

else throw SyntaxError

Input to  
parse

2\*3\*4  
↑



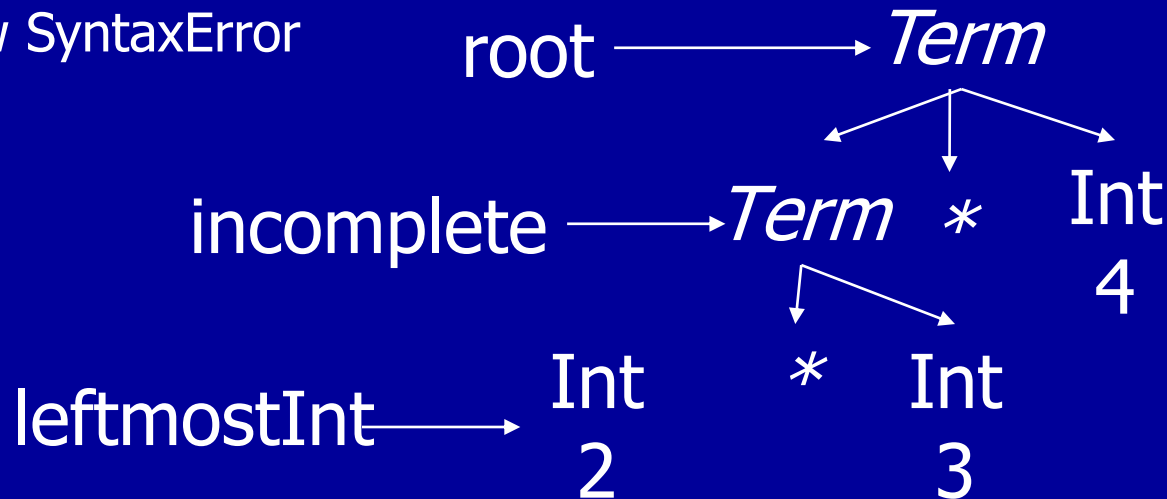
# Code for Term

Term()

```
if (token = Int n)
    leftmostInt = token; token = NextToken();
    (root, incomplete) = TermPrime();
    if (root == NULL) return leftmostInt;
    incomplete.leftChild = leftmostInt;
    return root;
else throw SyntaxError
```

Input to  
parse

2\*3\*4  
↑



# Code for Term

Term()

if (token = Int n)

leftmostInt = token; token = NextToken();

(root, incomplete) = TermPrime();

if (root == NULL) return leftmostInt;

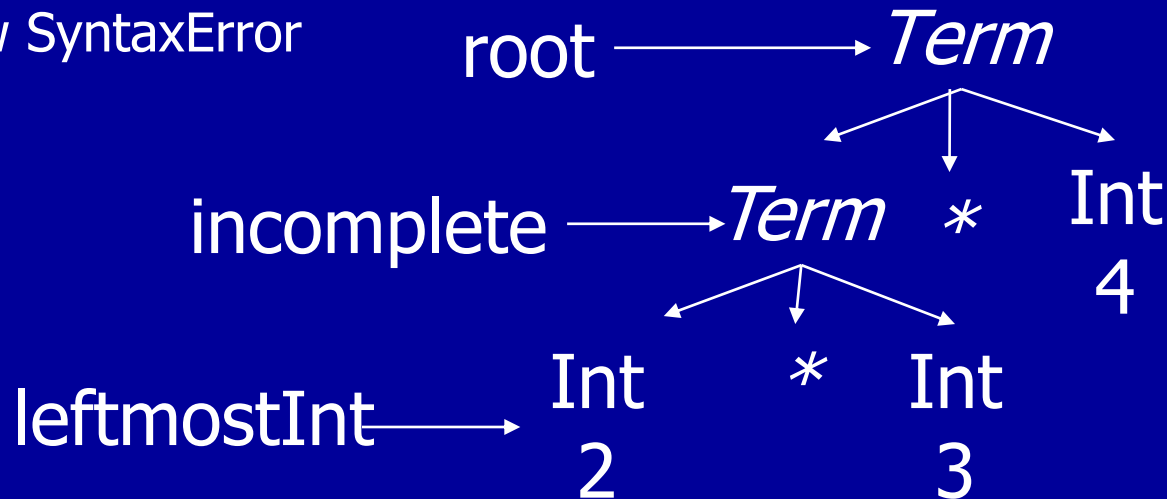
incomplete.leftChild = leftmostInt;

return root; ←

else throw SyntaxError

Input to  
parse

2\*3\*4  
↑



# Code for TermPrime

TermPrime()

```
if (token = *) || (token = /)
    op = token; next = NextToken();
    if (next = Int n)
        token = NextToken();
        (root, incomplete) = TermPrime();
        if (root == NULL)
            root = new ExprNode(NULL, op, next);
            return (root, root);
        else
            newChild = new ExprNode(NULL, op, next);
            incomplete.leftChild = newChild;
            return(root, newChild);
    else throw SyntaxError
else return(NULL, NULL)
```

Missing left child  
to be filled in by  
caller

