

Implementation of Circular Queue

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Lesson Plan

Subject/Course	Competitive Coding
Lesson Title	Implementation of Circular Queue

Lesson Objectives

To explain the concept and purpose of a circular queue and how it differs from a linear queue.

To implement circular queue operations (enqueue, dequeue, front, rear, isFull, isEmpty) using an array-based approach in code.

Evaluate the time and space complexity of common circular queue operations.

To identify suitable scenarios in computer science or real-life where Circular Queue is applicable and efficient.

Problem Statement:

Write a program to design a circular queue(k) which Should implement the below functions

- a. Enqueue
- b. Dequeue
- c. Front
- d. Rear

Introduction to the Problem

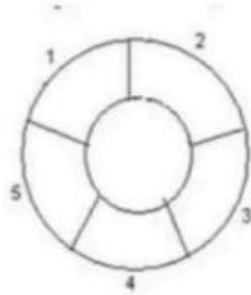
- A Circular Queue is a linear data structure that follows the FIFO (First In First Out) principle.
- Unlike a normal queue, the last position is connected back to the first position to make a circle.
- This structure efficiently utilizes memory by reusing empty spaces created by dequeued elements.

Concept and Background

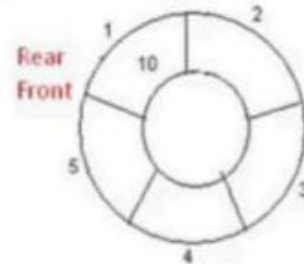
- Circular Queue connects the end of the queue back to the front.
- It prevents wastage of space in a linear queue when elements are dequeued.
- It can be implemented using Arrays, Linked List etc.
- Two pointers are maintained:
 - front \rightarrow points to the first element.
 - rear \rightarrow points to the last element.
- The queue is full when $(\text{rear} + 1) \% \text{size} == \text{front}$.
- The queue is empty when $\text{front} == -1$.

Visualization

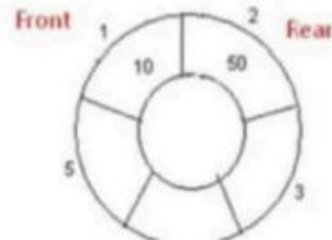
1.



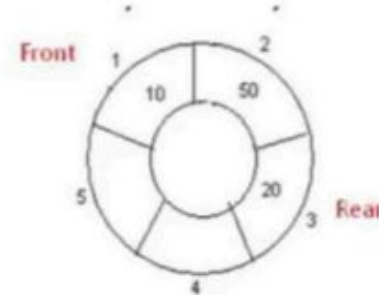
2. Insert 10, Rear = 1, Front = 1.



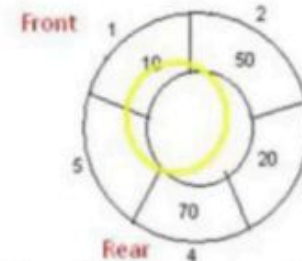
3. Insert 50, Rear = 2, Front = 1.



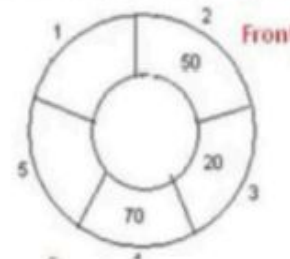
4. Insert 20, Rear=3, Front =1



5. Insert 70, Rear = 4, Front = 1.



6. Delete front, Rear = 4, Front = 2.



Queue Operations

1. Enqueue (Insert): Add an element at the rear.

- Check if the queue is full.
- Update $\text{rear} = (\text{rear} + 1) \% \text{size}$.
- Insert element at rear.

2. Dequeue (Delete): Remove an element from the front.

- Check if the queue is empty.
- Update $\text{front} = (\text{front} + 1) \% \text{size}$.

3. Peek: Retrieve the front element without removing it.

4. Display: Print elements from front to rear circularly.

Conditions for Queue State

- **Queue Full Condition:**

$(\text{rear} + 1) \% \text{size} == \text{front}$

- **Queue Empty Condition:**

$\text{front} == -1$

- **Reset Condition after last Dequeue:**

If $\text{front} == \text{rear}$, then set $\text{front} = \text{rear} = -1$

Algorithm/Logic

❑ **ENQUEUE(x):**

1. If $(\text{rear} + 1) \% \text{size} == \text{front} \rightarrow \text{Overflow}$
2. Else if $\text{front} == -1 \rightarrow \text{front} = \text{rear} = 0$
3. Else $\rightarrow \text{rear} = (\text{rear} + 1) \% \text{size}$
4. $\text{queue}[\text{rear}] = x$

❑ **DEQUEUE():**

1. If $\text{front} == -1 \rightarrow \text{Underflow}$
2. $x = \text{queue}[\text{front}]$
3. If $\text{front} == \text{rear} \rightarrow \text{front} = \text{rear} = -1$
4. Else $\rightarrow \text{front} = (\text{front} + 1) \% \text{size}$

Algorithm/Logic

□ **Front(x):**

1. If $\text{front} \neq -1$, then $\text{queue}[\text{front}]$ is the front item.

□ **Rear():**

1. If $\text{rear} \neq -1$, then $\text{queue}[\text{rear}]$ is the rear item.

Code Implementation

```
1 public class Practical4_CircularQueue {  
2  
3     static class CircularQueue {  
4         int[] queue;  
5         int front, rear, size, capacity;  
6  
7         CircularQueue(int k) {  
8             capacity = k;  
9             queue = new int[k];  
10            front = 0;  
11            rear = -1;  
12            size = 0;  
13        }  
14  
15        boolean enqueue(int value) {  
16            if (isFull()) return false;  
17            rear = (rear + 1) % capacity;  
18            queue[rear] = value;  
19            size++;  
20            return true;  
21        }  
    }  
}
```

```
22  boolean dequeue() {
23      if (isEmpty()) return false;
24      front = (front + 1) % capacity;
25      size--;
26      return true;
27  }
28
29  int Front() { return isEmpty() ? -1 : queue[front]; }
30  int Rear() { return isEmpty() ? -1 : queue[rear]; }
31  boolean isEmpty() { return size == 0; }
32  boolean isFull() { return size == capacity; }
33  }
34
35  public static void main(String[] args) {
36      CircularQueue q = new CircularQueue(3);
37      q.enqueue(10);
38      q.enqueue(20);
39      q.enqueue(30);
40      System.out.println("Front Element: " + q.Front());
41      q.dequeue();
42      q.enqueue(40);
43      System.out.println("Rear Element: " + q.Rear());
44  }
45 }
```

Output :

Front Element: 10

Rear Element: 40

Output

```
Front: 10  
Rear: 70  
Deleted: 10  
Front: 50  
Rear: 70  
Queue Elements: 50 20 70
```

Time & Space Complexity

- Enqueue Operation: $O(1)$
- Dequeue Operation: $O(1)$
- Peek Operation: $O(1)$
- Display Operation: $O(n)$
- Space Complexity: $O(n)$

Circular Queue provides efficient memory use and constant-time operations.

Summary

- Implemented a Circular Queue of capacity k supporting operations: enqueue, dequeue, front, and rear.
- Used modular arithmetic to connect the last index back to the first, avoiding wasted space.
- Handled overflow (queue full) and underflow (queue empty) conditions effectively.
- Demonstrated the difference between linear queue and circular queue, emphasizing better memory utilization.
- Strengthened understanding of queue data structure, front/rear pointers, and FIFO principle in data handling.

Practice Questions:

1.Design Circular Deque -- LeetCode#641

<https://leetcode.com/problems/design-circular-deque/>

2.Find the Winner of the Circular Game -- LeetCode#1823

<https://leetcode.com/problems/find-the-winner-of-the-circular-game/>

Thanks