

1 Design

This is an implementation in pseudocode of all 4 components of the radix sort variations included in this project: Radix sort starting from the most significant and least significant digit and then the integer sorting methods Counting sort and Pigeonhole sort. Each permutation of Radix sort takes the input array A of integers and treats each item $A[i]$ as a series of digits d within the range $0 \leq d < k$ where k is the radix or the number of unique digits.

1.1 Integer Sorting

Integer sorting is the underlying sorting method that is applied to each digit of the input array, these methods are non-comparitive algorithms that distribute items based on the value of each digit.

1.1.1 Pigeonhole Sort

Pigeonhole sort operates by instantiating an array B of length k (1) containing empty arrays (3). The input array A is iterated over and each item $A[i]$ is appended to the array in B at the position that corresponds to the current digit of the current item. B is then iterated over (10) and each item in each array overwrites the original array in order (11).

Algorithm 1 Pigeonhole Sort(A)

Input: $A \Leftarrow$ Unsorted Array of length n with maximum value k

Output: Sorted Array

```
1:  $B \Leftarrow [1..k]$  ▷ Buckets array
2: for  $i \leftarrow 1$  to  $k$  do
3:    $B[i] \Leftarrow []$  ▷ Each bucket contains empty array
4: end for
5: for  $j \leftarrow 1$  to  $k$  do
6:   append  $A[j]$  to  $B[A[j]]$  ▷ Move item to bucket
7: end for
8:  $c \Leftarrow 1$  ▷ Count items in A
9: for  $i \leftarrow 1$  to  $k$  do
10:  for all item in  $B[i]$  do
11:     $A[c] \Leftarrow item$  ▷ Move items from bucket to A
12:     $c \Leftarrow c + 1$ 
13:  end for
14: end for
15: return  $A$ 
```

1.1.2 Counting Sort

Counting sort instantiates 2 arrays; B , the output array equal in length to the input array A (1), and C , a count array to store the frequency of each digit which is of length k and is instantiated to contain only integers of value 0 (2, 4). The input array is iterated over and the integer in C at the position equal to the current digit of the current item is incremented (7), producing a histogram of the frequency of each digit. This count list is then accumulated such that each item is added to the sum of the previous values (10). Finally the input array is iterated over in reverse and each item is moved to the output array in the position that is referred to by the value expressed in the count array at the position of the current digit of the current item (13), to ensure that values do not overwrite one another in the output array each item in the count array is decremented when a position in the output array is taken from it (14).

Algorithm 2 Counting Sort(A)

Input: $A \Leftarrow$ Unsorted Array of length n with maximum value k

Output: Sorted Array

```
1:  $B \Leftarrow [1..n]$                                  $\triangleright$  Temp array
2:  $C \Leftarrow [1..k]$                                  $\triangleright$  Count array
3: for  $i \leftarrow 1$  to  $k$  do
4:    $C \Leftarrow 0$ 
5: end for
6: for  $j \leftarrow 1$  to  $n$  do
7:    $C[A[j]] \Leftarrow C[A[j]] + 1$                  $\triangleright$  Count occurrence of value
8: end for
9: for  $i \leftarrow 2$  to  $k$  do
10:   $C[i] \Leftarrow C[i] + C[i - 1]$                  $\triangleright$  Accumulate Count array
11: end for
12: for  $j \leftarrow n$  to  $1$  by  $-1$  do
13:   $B[C[A[j]]] \Leftarrow A[j]$                      $\triangleright$  Move item from input to temp array
14:   $C[A[j]] \Leftarrow C[A[j]] - 1$                  $\triangleright$  Decrement count array
15: end for
16: return  $B$ 
```

1.2 Radix Sort

Radix sort is the method of applying the integer sorting method to each digit of the input, there are two approaches to this: Least significant digit first (LSD) or Most significant digit first (MSD).

1.2.1 Least Significant Digit

Least Significant Digit Radix sort (LSD) takes an input array A and sorting function $f()$ which can be either of the previous integer sorting methods. The input array is sorted via the chosen method once for each of digits in the chosen radix of the largest value (3) This sorting process is applied first at the least significant or rightmost digit and moves incrementally to the most significant or leftmost digit. Values can have an unequal number of digits provided any digits farther left than their most significant are taken as 0 and the chosen integer sorting method is stable as this means they maintain their relative position with respect to values that have an equal or lesser number of digits.

Algorithm 3 LSD($A, f()$)

Input: $A \Leftarrow$ Unsorted Array of length n with maximum absolute Value k

Input: $f() \Leftarrow$ Sorting function \triangleright Pigeonhole or Counting Sort

Output: Sorted Array

1: $d \Leftarrow$ Number of digits of k \triangleright Varies based on radix chosen

2: $A(i) \Leftarrow$ i -th digit of each item in A

3: **for** $i \leftarrow 1$ to d **do**

4: $A \Leftarrow f(A(i))$ \triangleright Sorts Array using sort function

5: **end for**

6: **return** A

1.2.2 Most Significant Digit

Most Significant Digit Radix sort (MSD) takes an input array A , sorting function $f()$ and current digit i which for the first call of the function should be the number of digits of the largest value in the input. This algorithm operates recursively, and the termination condition occurs when i is equal to 0 meaning that all digits of the input have been sorted (1,2). After ensuring the termination condition has not been met the input array A is then sorted using the chosen function with respect to the current digit (5). The input array is then iterated over and items that have an equal value at the current digit are identified (8) and taken as a slice of the input array. This slice is then sorted recursively using a function call to $MSD()$ with a decremented value for i meaning they are sorted over the next least significant digit (11). Once the termination condition has been reached within this slice and all sub-slices created by recursive calls the algorithm resets the start point of the slice and identifies further slices to sort (12).

Algorithm 4 $MSD(A, i, f())$

Input: $A \Leftarrow$ Unsorted Array of length n with maximum absolute Value k

Input: $i \Leftarrow$ Digit to sort Array on ▷ start with number of digits of k

Input: Sorting function $f()$ ▷ Pigeonhole or Counting Sort

Output: Sorted Array

```
1: if  $i == 0$  then
2:   return  $A$ 
3: end if
4:  $A(i) \Leftarrow$   $i$ -th digit of each item in  $A$ 
5:  $A \Leftarrow f(A(i))$  ▷ Sorts Array using sort function
6:  $s \Leftarrow 1$  ▷ Starting position of recursive call
7: for  $j \leftarrow 1$  to  $s$  do
8:   if  $A(i)[j] == A(i)[s]$  then
9:     Do Nothing
10:  else
11:     $A[s..j - 1] \Leftarrow MSD(A[s..j - 1](i), i - 1)$ 
12:     $s \Leftarrow j$  ▷ Reset start position of array slice
13:  end if
14: end for
15: return  $A$ 
```

2 Time complexity

2.1 Integer Sorting

Counting and Pigeonhole sort both require one iteration over the input list which contains n items and one pass over the intermediary buckets array for pigeonhole sort or the count array for counting sort, both of these arrays contain k items where k is the maximum range of input values. The combination of these two passes gives both integer sorting methods a time complexity of

$$\mathcal{O}(n+k)$$

2.2 Radix Sort

Radix sort repeatedly applies integer sorting to each successive digit of the input array, it differs from counting sort in that it does not sort into k buckets but a constant base b is chosen. It repeats the integer sorting method d times where d is the number of passes necessary to sort each digit of the maximum value in the input with respect to the base b and can be expressed as $d = \log_b k$, giving a time complexity of:

$$\mathcal{O}((n + b) \cdot d)$$

or if b is a constant:

$$\mathcal{O}(n \cdot d)$$

However, in the application of this project, radix sort will only be applied to integers represented in 64-bits, meaning there is a defined upper bound to the max value k . This is important because it means that k is polynomial in b or that k is always able to be represented by b raised to some exponent c where c is some constant value. This can be expressed in the form $k = b^{\mathcal{O}(1)}$ or $\mathcal{O}(1) = \log_b k$

$\mathcal{O}(1)$ can be substituted in for d in the time complexity analysis meaning Radix sort will run in $\mathcal{O}(n)$ or linear time in the context of this project.

3 Radix Selection

3.1 Introduction

Radix sort utilises repeated rounds of integer sorting across each digit of the input list, therefore, it is necessary to make a clear definition of a digit so that an effective method of separating an integer into its component digits can be implemented. When expressing a number of any kind, it is necessary to establish the numerical base that will be used, this is the number of unique values that can be displayed in each digit and when a value exceeds the maximum value within this base, another digit must be added; i.e. in standard decimal or base 10 numbers 0 through 9 can be expressed in a single digit but increasing beyond this range requires adding a second digit to create the number 10.

Though numbers in typical day to day use by humans are typically thought of as arabic numerals in the decimal system, this is a social convention which this algorithm is not beholden to, and as Radix sort fundamentally relies on the processing of the component digits of an integer it is of great importance to select an appropriate base as well as identify an efficient method for processing each digit.

3.2 Explanation

Python stores integers in the format "two's complement", which is a type of binary and therefore uses a base of 2. This means that bases that are of the order 2^n where n is some integer are the most suitable for a typical radix sort algorithm, this is because bases that fit this specification allow each digit to be directly accessed from an integer's binary digits.

Any base of 2^n will utilise all possible combinations of n digits of binary to express a single digit and so the rightmost n digits represent the first digit in 2^n , the next n digits represent the next digit etc.

3.3 Implementation

The method of identifying each digit of an integer from the binary representation first involves manipulating the integer such that the desired digit is in the rightmost position, in base n this means discarding the rightmost $n \cdot d$ binary digits where d is the index (from the right) of the required digit in base 2^n . Following this, isolating the rightmost n binary digits produces the correct digit.

In order to complete this task, bitwise operations have been utilised, and are explained fully in section 5.2.

| | | | | | | | | | | | | | | | |
|---------|---|---|---|---------|---|---|---|---------|---|---|---|---------|---|---|---|
| 1 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |
| Digit 4 | | | | Digit 3 | | | | Digit 2 | | | | Digit 1 | | | |
| 14 | | | | 10 | | | | 6 | | | | 0 | | | |

The 4 digits in base 2^4 of a 16 bit integer: 60000

4 Negative Integers

4.1 Two's Complement

When handling integers in two's complement, positives are represented in binary, but with a leading 0 in the leftmost position to indicate the positive sign of the integer. Negative integers are represented with a leading 1 which represents that the binary value of this digit is to be subtracted i.e. negative and all other binary digits are positive such that the sum of all the digits equals the integer being represented. table reference has been included to illustrate this, as you can see the value +3 has a leading 0 at index 2, representing 2^2 is not subtracted, this is followed by 1 in all remaining digits representing the addition of 2^1 and 2^0 or $2 + 1 = 3$. -3 is represented by a leading 1 to represent -2^2 followed by 0 1 to represent no addition of 2^1 and the addition of 2^0 or $-4 + 1 = -3$.

The effect of performing Radix sorting without any accommodation for negativity is that due to the leading 1, negative values will be incorrectly sorted ahead of positive values as though they were larger. They will also be sorted into the correct order relative to other negative values (increasing from most negative to least negative), as among values that can be represented in the same number of bits the more negative numbers have fewer additional bits in positions to the right, for example -4 is 100 and -3 is 101.

The method of correcting for this problem is to invert the bits representing the sign of the largest absolute value of the input list (see 5.2.5 for further explanation). This operation causes the items to retain their relative position amongst elements of the same sign but causes negative values to be sorted to their correct position as less than positive values.

| Bits | Unsigned Value | Signed Value (Two's Complement) |
|------|----------------|------------------------------------|
| 000 | 0 | 0 |
| 001 | 1 | 1 |
| 010 | 2 | 2 |
| 011 | 3 | 3 |
| 100 | 4 | -4 |
| 101 | 5 | -3 |
| 101 | 6 | -2 |
| 111 | 7 | -1 |

Table 1: 3 bit integers and their respective values when interpreted in unsigned binary and two's complement

4.2 Implementation

The purpose of this project was to incorporate this sorting algorithm into the standard library for Pypy, meaning it had to be written in RPython, the language used for Pypy. RPython is a restricted form of python, specifically it is a form of Python 2, which is relevant to the handling of integers as RPython inherits Python2's method of allocation of a quantity of bits for each integer equal to the word size of the machine it is being run on (typically this is 32 or 64 bits), even if the integer can be expressed in fewer bits. The leftmost bit represents whether the integer is positive or negative and the remaining bits represent the magnitude of the integer, all digits between the sign bit and the most significant bit needed to represent the integer are copies of the sign bit. To put this in the context of two's complement, this representation can be understood as the first bit representing -2^m where m is the machine word length and all following digits are additions of 2^n where n is $[m - 1..0]$.

The true sign bit, meaning the bit that always and exclusively represents the sign of the integer, is the one in the leftmost position, however it is not necessarily the most optimal solution to this problem of negatives being sorted into the incorrect position to flip this bit. Firstly, when sorting integers that take up a small number of bits it would require some additional logical test to identify if the algorithm is performing the first (for MSD) or last (for LSD) pass of the data to then instruct the algorithm to jump to the leftmost bit. Additionally, it is necessary to identify the machine word length of the machine on which this code is running as flipping the 64th bit on a 32 bit machine is not possible, this eliminates the possibility of a hardcoded value and essentially forces the use of the "sys.maxint" function from the Python 2 standard library. I decided that using this function was not in the best interest of this project for posterity as it was removed from Python 3 and replaced with "sys.maxsize" which means if RPython ever moved on to Python 3 this function would stop working; whilst this is a minor issue that would likely never cause problems there is an alternative solution without any issues.

The approach used in this project was to identify the number of digits needed to express the largest absolute value of the input in the chosen base and then flip the leftmost bit of the digit preceding this digit. So for example the integer 246 when expressed in base 4 (equal to 2^2 or 8 bits) requires 4 digits [3312], so in RPython the preceding bits are all 0 to indicate the sign being positive, this method then identifies the bits that would represent the next digit (the bits 9th and 10th from the right) and flips the leftmost digit giving it the binary value of [10] and a digit value of 2 in base 4. This dynamic approach ensures that the relative correct order within the positive and negative integers is preserved and is suitable for use with any integers being sorted in any base.

5 Implementation

5.1 Extra Methods and Class

Explanation of all functions and classes that exist outside of the radixsort algorithms. Each function is documented in isolation in the order they appear in the source files, for a full view of each of the files refer to [reference](#)

5.1.1 Function absolute

The purpose of this helper method was to allow for the calculation of absolute values without overflow issues specifically caused by using the minimum possible value. This value's positive counterpart lies outside of the range of the 64 bits assigned for each integer in RPython, this causes unpredictable errors and even though there is only one specific case that necessitates the use of this method over the regular `abs()` inbuilt function, the runtime impacts are low so it's a worthwhile helper method to include.

```
import sys

def absolute(num):
    if num == (-sys.maxint) - 1:
        return sys.maxint
    return -num if num < 0 else num
```

5.1.2 Function int_bytes

Returns the minimum quantity of bytes in a given radix necessary to express an integer, including a sign bit. This function is used to calculate the number of iterations that the sorting algorithm must make in order to sort all values up to and including the given input i .

The function calculates $\log_{radix}(i)$, i.e. to what power i must be raised to get $radix$, uses `ceil` to get the next highest integer and adds 1 to the value to allow for the sorting of negatives. It is necessary to return an integer in this case because this value is used to determine the number of bits used to identify each digit and the number of iterations over the list that must be made, both of these are discrete values. `ceil()` has been used instead of `int()` or `floor()` as a value lower than needed would result in incorrect sorting of values the larger values.

```
from math import ceil, log

def int_bytes(i, radix):
    return ceil(log(absolute(i) / log(radix))) + 1
```

5.1.3 Function `make_radixsort_class`

The container function that instantiates all other functions, this is modelled on the existing function in `listsort.py` [reference](#) to minimize compatability issues when including this project in the pypy library.

Almost all functions have been removed or altered in a major way, particularly any functions that pertain to comparative sorting, it is the general structure that has been emulated.

```
def make_radixsort_class(
    setitem=None,
):
```

5.1.4 Function `setitem`

This function has been imported from the existing `listsort.py` file within `rlib` [Reference? - This whole needs more research, I'm not certain it is correct.](#) It serves the function of enabling the input list to be edited by functions within the Radixsort class. When the `.sort()` function is called by PyPy, it initialises the Radixsort class with the list as a constructor, this means that the list attribute is not mutable from within the class. Importing `setitem` *setitem* into the class allows the input list to be updated.

This is the only existing function of the original `listsort.py` that was not removed.

```
if setitem is None:
    def setitem(list, item, value):
        list[item] = value
```

5.1.5 Class `Radixsort`

This is the class that contains all other functions of the radixsort algorithm, it creates an object that is used to ease the manipulation of the input list and store attributes that are referenced and edited from multiple different functions.

in particular the list attribute contains the input list, which is used by every other function of the object.

The function *setitem* is imported here as mentioned above so that the list can be edited outside of the scope of the class.

```
class Radixsort(object):
    def __init__(self, list, listlength=None):
        self.list = list
        self.base = 0
        self.listlength = len(self.list)
        self.radix = 0

    def setitem(self, item, value):
        setitem(self.list, item, value)
```

5.1.6 Function `list_abs_max`

This function is an implementation of a combination of `min()` and `max()` functions (neither of which are implemented in RPython [reference](#)) as well as an optional module to identify whether the list is sorted in ascending or descending order. The function also compares the maximum and minimum values and returns the value that is farthest away from 0.

The primary purpose of this function is to return the value that is used to identify the number of passes over the input necessary for sorting, this is why it is necessary to identify the farthest value from 0 as this is the value that will contain the most digits.

The secondary purpose of the function is to update the class attributes *ordered* and *reverseOrdered*, which are used to identify whether the sort function has been called on an already sorted list in which case there is no purpose in running the algorithm.

The reason that these actions are performed at the same time is because iterating over the list is necessary for both purposes, and both must be performed before sorting can take place (or not take place if the list is sorted), therefore it is more efficient to perform both tasks simultaneously.

```
self.ordered = True
self.reverseOrdered = True

def list_abs_max(self, checkorder=False):

    assert len(self.list) != 0
    m = self.list[0]
    n = self.list[0]
    prev = self.list[0]
    for i in range(1, len(self.list)):
        if self.list[i] > m:
            m = self.list[i]
        if self.list[i] < n:
            n = self.list[i]
        if checkorder:
            self.ordered &= self.list[i] >= prev
            self.reverseordered &= self.list[i] <= prev
            prev = self.list[i]
    return m if absolute(m) > absolute(n) else n
```

5.1.7 Function insertion_sort

This is an implementation of insertion sort, it is used in MSD sorts when the length of a sublist falls below the critical value. [reference](#)

```
def insertion_sort(self, start, end):
    for step in xrange(start, end):
        key = self.list[step]
        j = step - 1
        while j >= 0 and key < self.list[j]:
            self.setitem(j + 1, self.list[j])
            j = j - 1
        self.setitem(j + 1, key)
```

5.1.8 function reverse_slice

This reverses the items of the input inbetween the provided start and stop indexes. By default is reverses te entire list, but includes provisions to reverse a smaller slice of the list (this was from an earlier version of the algorithm, this feature is currently unused). [I should probably just remove that at this point](#)

```
def reverseSlice(self, start=0, stop=0):
    if stop == 0:
        stop = self.listlength - 1
    while start < stop:
        i = self.list[start]
        j = self.list[stop]
        self.setitem(start, j)
        self.setitem(stop, i)
        start += 1
        stop -= 1
```

5.2 Bitwise operations

Bitwise operations operate on binary numerals at the level of individual bits, they allow for fast and consistent manipulation that bypasses arithmetic functions. Through the use of bitwise operations, radix sort can be performed using different numeral bases, meaning a different maximum value at each digit and a different total number of digits. By reducing the number of digits the amount of iterations over the data is reduced thereby improving run time, however this comes at the cost of an increase in the amount of time each round of integer sorting takes and an increase in the memory consumption due to the larger count list or quantity of bins.

Bitwise operations are particularly useful in this regard as it makes processing integers in bases of the order 2^n where n is some integer very easy (see section 3.2) whereas changing base from the standard base 10 requires time consuming arithmetic methods such as euclidian division.

In order to properly implement bitwise operations it was necessary to properly explore their implementation in RPython as well as the implementation of the Integer type. In RPython each integer is comprised of exactly 64 bits, where 1 bit is assigned to represent the integer's sign, leaving a maximum of 63 bits to express the magnitude of the integer, this gives integers in RPython a range of $[-2^{63} .. 2^{63}-1]$. The PyPy interpreter considers any integer outside of this range to instead be of the 'Long' type, and as such a list containing even a single value outside of this range does not have the Integer specific sorting method applied to it, thus values outside of this range are not within the scope of this project.

I have already covered in general terms how these are used in this project to identify the value of each digit in section 3 and handle negative integers in 4.2 but in this section I will cover the implementation of the bitshift operations used in RPython.

5.2.1 Left shift

A ‘Left shift’ refers to moving each bit of an integer a number of digits to the left, in RPython this is performed using: $n \ll s$ where n is the integer to be shifted and s is the number of digits the integer should be shifted to the left.

RPython utilises Arithmetic Shifting, this means when a shift is performed and the rightmost bit is no longer in the rightmost position, all digits in the range $[0..s]$ are overwritten with the value 0 and an equal quantity of digits in the leftmost positions are discarded.

In this project, left shifting is used to create integers that occupy a specific number of bits and is used for the accomodations necessary to sort negative integers (further explained in 5.2.5).

| Value | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
|------------|---|---|---|---|---|---|---|---|
| 13 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 1 |
| $13 \ll 2$ | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 |

Table 2: Demonstration of shifting 13 left 2 bits
only rightmost 8 bits are shown

5.2.2 Right shift

A ‘Right shift’ is the inverse of left shift, referring to moving each bit of an integer to the right, using the operator $n \gg s$ where n is the integer to be shifted and s is the number of digits the integer should be shifted to the right.

Like Left shifting, Right shifting utilises Arithmetic shifting, meaning the rightmost s digits are discarded. Right shift differs from leftshift however as the digits in the the leftmost positions that become empty are overwritten with a value equal to the most significant (leftmost) bit, which would typically represent the sign of the integer. This means that a Right shifted integer retains it’s original sign.

Right shifting is used in this project to manipulate integers from the input list in order to move a specific digit into the rightmost position so that the value of this digit can be extracted using Bitwise And (see 5.2.3). The bases selected for this project are of the order 2^n where n is some integer, meaning a single digit of the integer occupies n bits and therefore; Right shifting $n \cdot d$ positions places the digit d from the right in the rightmost position, occupying the rightmost n bits.

| Value | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
|------------|---|---|---|---|---|---|---|---|
| 13 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 1 |
| $13 \gg 2$ | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |

Table 3: Shift 13 right 2 bits. Only
rightmost 8 bits are shown

| Value | 63 | 62 | 61 | 60 | 59 | 58 | 57 | 56 |
|-----------------|----|----|----|----|----|----|----|----|
| -2^{61} | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |
| $-2^{61} \gg 2$ | 1 | 1 | 1 | 1 | 1 | 0 | 0 | 0 |

Table 4: Shift -2^{61} right 2 places. Only leftmost
8 bits are shown.

5.2.3 And

The ‘And’ operator takes two integers and compares each bit across all binary digits, returning 1 if both bits are equal to 1 or returning 0 if either or both are equal to 0.

This is utilised in this project to isolate a specified number of binary digits of an integer, specifically n digits in binary represent one digit of the integer in base 2^n . This is done by ‘Right shifting’ $n \cdot d$ positions (as stated in 5.2.2) and then by performing an ‘And’ operation with $2^n - 1$, which in binary is represented as a 1 in the rightmost n digits, and 0 in all other digits, which will retrieve the rightmost n bits and 0 for all other bits.

In pypy, the And operation is performed using the $x \& y$ operator resulting in an ‘And’ operation between x and y .

| Value | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
|---------|---|---|---|---|---|---|---|---|
| 59 | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 1 |
| 15 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 |
| 59 & 15 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 1 |

Table 5: Demonstration of ‘And’ on 15 and 59

This has isolated the rightmost 4 digits which is one digit in 2^4

5.2.4 Not

The bitwise not operator produces the inverse of all binary digits in an integer, giving a 1 if an integer contains a 0 and vice versa. It is unique in the sense that it is the only operation that is performed on a solitary input. This is used in the process of inverting the sign of integers detailed in 5.2.5

| Value | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
|------------|---|---|---|---|---|---|---|---|
| 102 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 |
| ~ 102 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 |

Table 6: Demonstration of ‘Not’ performed on 102

All bits have been inverted

5.2.5 Exclusive Or

‘Exclusive Or’ compares the bits of two integers and returns a value of 1 for each digit where the two bits are different and 0 if they are the same. It is performed in pypy using $x \hat{y}$ which results in an ‘Exclusive Or’ performed across all 64 bits of the two integers x and y .

The use case for this in this project occurs when inverting the sign bit to ensure negative and positive values are sorted correctly (see 4.1 for details). This is done by identifying the quantity of binary digits needed to express the largest absolute value of the input using the function ‘int.bytes’ (shown in 5.1.2) with a radix of 2, this quantity is given as the integer d . Left shifting an integer of value 1 d times and then subtracting 1 produces the maximum integer expressable in the same number of bits as the largest absolute input value, this is given as integer v . In binary v contains 1 in all digits up to the most significant digit of the largest absolute input value excluding any sign bits and a 0 in all other digits.

This value v is then inverted using bitwise ‘Not’ (see 5.2.4) so that it instead contains 0 in the bits representing the max input and 1 in all other bits. As explained in 4.2, in two’s complement binary all digits to the left of the most significant are equal and represent the sign of an integer, therefore v contains 1 in all the sign bits of the largest max input. Performing the Bitwise ‘Exclusive Or’ operation with any value in the list against the value v inverts all values in these sign bits. As shown in 4.2, negative integers are originally stored in the correct order but placed higher than positive values, performing these operations ensures that the relative order is not altered but the negatives are positioned correctly.

Depending on the base selected and the number of digits of the largest absolute value, multiple sign bits may be contained in the leftmost digit of some integers being sorted, for example 300 contains 3 digits in base 2^4 , but only 9 digits in binary, therefore the leftmost digit includes three sign bits. Therefore all sign bits must be flipped to ensure that sorting an integer across digits containing sign bits does not alter their position in the sorted list.

| Value | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
|---------------------|---|---|---|---|---|---|---|---|
| -13 | 1 | 1 | 1 | 1 | 0 | 0 | 1 | 1 |
| ~ 15 | 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 |
| $-13 \hat{\sim 15}$ | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |

Table 7: Performing ‘Exclusive Or’ on -13 with ~ 15
All bits in range [63..8] now equal 0

5.2.6 Optimizations **Incomplete**

Some of the code outlined in section [Reference code explanation](#) differ in relatively substantial ways from the pseudocode outlined in section 1, many of these differences arise from increasing the performance of the algorithm. In this section I will explain the function and reasoning of these differences.

5.2.6.1 LSD counting sort - creating all count lists in one pass

LSD counting sort is a combination of LSD sort ([1.2.1](#)) which calls counting sort ([1.1.2](#)) on every digit. It iterates over the input list for each digit: once whilst counting the occurrence of the values at the digit and a second time when moving the input list items into the output list based on these count lists.

An optimization I have used is to make a single pass over the input list whilst creating the count lists for every digit simultaneously instead of once for each digit. The reason this is possible is that the occurrence of the values for each digit do not change based on the order of the input list, items are moved from the input list to their new position based on the value stored in the count list representing the value of the current digit of the item, the order of the input list has no bearing on the order or values of the count list.

The values stored in the count list represent the position of the items relative to one another after being sorted across a given digit; due to MSD sort creating sublists in a specific order and then sorting the sublists' items amongst one another, a count list created for a different digit was therefore created for a different list and so does not represent the position of the items relative only to other items in the sublist.

The way that this has been implemented in this project is seen below. It involves creating a 2 dimensional array with an outer list of d sublists each of length r where d is the number of digits and r is the radix. Following this the input list is iterated over and for each item the count list of each digit is incremented i.e. the sublist at position $[0..d]$ represents the count of the digit at the same position. After this, list of count lists is iterated over and each list is incremented, then when moving the input items based on the count lists the relevant count list can be accessed by accessing the outer list at the index equal to the index of the current digit.

```
#Note: this code has been simplified
counts = [[0 for _ in range(radix)] for _ in range(min_bytes + 1)]

for num in self.list:
    disc = 0
    for i in xrange(min_bytes):
        val = get_digit(num, i)
        counts[i][val] += 1

for i in range(min_bytes):
    for j in xrange(1, self.radix):
        counts[i][j] += counts[i][j - 1]
```

5.2.7 Order checking **This needs testing further, it might be incorrect**

Radix sort has no advantage when called on already sorted lists compared to unsorted due to the non-comparative nature of the algorithm. Simple comparative algorithms such as bubble sort will actually out perform radix sort when called on sorted lists and as Timsort is highly efficient at sorting them there is substantial motivation for improvements to be made with regards to handling sorted data.

Identifying sorted lists is not computationally expensive, it involves iterating over and comparing the list items in order and ensuring that each value is larger or equal to the previous value, essentially only involving $n-1$ comparisons where n is the input list length giving a time complexity of $\mathcal{O}(n)$. As this algorithm already involves iterating over the data, the run time cost of adding a comparison function to check for sortedness is extremely low, thus a check to identify whether a list is sorted as well as whether a list is reverse sorted is included in the first iteration over the data, which is the point at which the absolute maximum value of the list is identified (a necessary step for this algorithm).

```
def list_abs_max(self, checkorder=False):

    m = self.list[0]
    n = self.list[0]
    prev = self.list[0]
    (ordered, reverseordered) = (True, True)
    for i in range(1, len(self.list)):
        ordered &= self.list[i] >= prev
        reverseordered &= self.list[i] <= prev
        prev = self.list[i]
    if checkorder:
        self.ordered = ordered
        self.reverseOrdered = reverseordered
```

Checking whether the list is ordered or reverse ordered whilst
iterating over the list to identify the absolute maximum value

An extension of this idea is to perform this order checking every time the input list is rearranged (once for each digit). Unlike the initial order checking which has a minor footprint on the run time of the algorithm, checking order on each digit has a moderately significant impact, therefore it should only be implemented on the conditions that it is reasonably likely that sorted lists will occur and that when the test does identify a sorted list the time saved from early termination of the algorithm is greater than the time added by performing the test.

Testing involving the comparison of the performance of algorithms including and excluding order checking showed that MSD radix sorts can benefit from order checking performed on sublists which trigger positive ordered tests frequently and early termination of the sorting of these sublists improved performance **Obviously you have to prove this**. In contrast, LSD sorts rarely trigger positive results with the frequency of these events dropping to 0 when negative integers are included. The lack of occurrences of sorted lists in this circumstance is likely because the least significant digits

have weaker influence on the size of integers meaning aside from random events and specific circumstances the input list is simply less likely to become sorted earlier than expected; the inclusion of negatives precludes the possibility of a list being sorted because of the leading sign bits causing negatives (and particularly negatives closer to 0) to be sorted as though they are the largest integers i.e. in the wrong order.

```
def is_sorted(self, start=0, end=-1):
    end = end if end > -1 else self.listlength
    assert start >= 0; assert end >= 0
    sortd, rsortd = True, True
    for i, el in enumerate(self.list[start + 1 : end]):
        if el < self.list[start + i]:
            sortd = False
            if rsortd == False: return sortd, rsortd
        if el > self.list[start + i]:
            rsortd = False
            if sortd == False: return sortd, rsortd
    return sortd, rsortd
```

Checking whether a slice of the list is ordered or reverse ordered
after integer sorting has been applied to each digit

5.2.8 skipping iterations

The time complexity of radix sort is $\mathcal{O}(N \cdot k)$ where k is the number of digits that must be sorted, therefore it is possible (and effective) to have a meaningful effect on the runtime of the algorithm by reducing k , this is why the use of different bases has been focussed on as this also reduces k (the number of iterations over the input list that needs to be made).

Skipping iterations over the data can be made by identifying digits on which no sorting is necessary, the specific method of identification differs based on the implementation of radix sort.

For counting sort, if every single list item contain an identical value for a digit there will be no change in the order of the list when sorting across that digit. This circumstance is identified by a single count list item being equal to the list length and can be identified by iterating over the count list, a positive result means the process of moving the list items can be skipped. As the count list is already iterated over to accumulate the values, this test is computationally inexpensive.

```
skip = []
for i in xrange(min_bytes):
    for j in xrange(self.radix):
        if counts[i][j] == self.listlength:
            skip.append(i)
            break
    if j==1: continue
    counts[i][j] += counts[i][j - 1]
```

Identifying if all items have an identical value for each digit in LSD counting sort. The same process is used for MSD counting sort but the sublist length is used instead of total list length

In pigeonhole sort a similar opportunity arises in the form of identifying whether all items have been placed into any of the buckets which would mean overwriting the input list with the bucket wouldn't change the order. Testing for this condition involves iterating over the buckets list and counting the number of non-empty buckets, if the count increases past 1 bucket it signals the condition has not been met and the test iteration can be terminated. This test takes up no space and at worst $\mathcal{O}(k)$ time to iterate over the bucket list but negative resulting tests take far less time typically.

```
count = 0
for b in bucket:
    if b!=[]:
        count+=1
    if count>1:
        break
if count==1:continue
```

Identifying if more than 1 bucket has been used

5.2.9 Insertion sort

Radix sort has a significant overhead cost due to the establishment of the count lists and the list of bins that must be instantiated or cleared on each digit, this cost has little to no relation to the length of the list and so it is often highly inefficient when compared to other algorithms to perform radix sort.

This issue becomes exponentially worse when radix sort is implemented in MSD fashion due to the nature of establishing sub-lists and thus performing sorting on an ever increasing number of lists containing progressively fewer and fewer items (sorting 150000 items in base 2^{16} can create over 65000 lists with an average length of 2 for which the count list must be established separately).

Using insertion sort in situations where the input list is reasonably small can be much more efficient than performing continuous rounds of radix sort, this effect is compounded on small lists of large items. Identifying the exact point at which it becomes more efficient to use insertion sort was done by [Not yet decided](#)

5.3 Algorithms

5.3.1 LSD Counting sort

In lines 2-18 the algorithm is instantiating the variables it needs to perform calculations:

listmax contains the largest absolute value of the list and is used to calculate the passes over the data needed to complete sorting as well as the integer mask **is integer mask the correct term**_j used to invert the sign bit.

min_bytes Uses the listmax variable and the selected radix to calculate the total iterations over the data (rounds of integer sorting) that will be needed, which is equal to the number of digits required to express *list_max* in the current radix, to account for the sign of the integer, *min_bytes* includes an extra digit incase . However if *min_bytes* is equal to the maximum possible number of digits, *min_bytes* is reduced by 1, this is a safeguard put in place to prevent the algorithm from using a rightshift value greater than the machine word length as this results in an overflow error, when *min_bytes* **Can this be done by checking if min bytes is divisible by the radix?**

uint_63 Integer mask used to invert the sign bit. Line 16 is the default value, but in the case of a listmax occupying the same quantity of digits as the largest possible value line 13 is used. $(1 \ll \text{int_bytes}(\text{listmax}, 2))$ is represented in binary by a 1 in the position directly to the left of the leftmost 1 of the list max,.

```

1  def sort(self):
2      if self.listlength < 2:
3          return
4      listmax = self.list_abs_max(checkorder=True)
5      min_bytes = int_bytes(listmax, self.radix)
6      if self.ordered == True:
7          return
8      if self.reverseOrdered == True:
9          self.reverseSlice()
10         return
11
12     if min_bytes == int_bytes((-sys.maxint) - 1, self.radix):
13         uint_63 = ~((1 << int_bytes(listmax, 2) - 1) - 1)
14         min_bytes -= 1
15     else:
16         uint_63 = ~((1 << int_bytes(listmax, 2)) - 1)
17
18     counts = [[0 for _ in xrange(self.radix)] for _ in xrange(min_bytes)]
19
20     for num in self.list:
21         for i in xrange(min_bytes):
22             shift = (self.base) * i
23             sortkey = (num) ^ uint_63
24             val = (sortkey >> shift) & self.radix - 1
25             counts[i][val] += 1
26
27     skip = []
28     for i in xrange(min_bytes):
29         for j in xrange(1, self.radix):
30             if counts[i][j] == self.listlength:
31                 skip.append(i)
32                 counts[i][j] += counts[i][j - 1]
33     temp_list = [0 for _ in xrange(self.listlength)]
34     for i in xrange(min_bytes):
35         if i in skip:
36             continue
37         shift = (self.base) * i
38         for j in xrange(self.listlength - 1, -1, -1):
39             num = self.list[j]
40             sortkey = (num) ^ uint_63
41             val = (sortkey >> shift) & self.radix - 1
42             temp_list[counts[i][val] - 1] = self.list[j]
43             counts[i][val] -= 1
44     self.setslice(temp_list, 0)

```

LSD counting sort