Structured matrix

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ABSTRACT

The documentation for fast MSSA and relationship between multi-dimensional convolution and structured matrix, such as Circulant, Toeplitz and Hankel matrix.

INTRODUCTION

(Korobeynikov, 2009) proposed the computation and memory efficient method for MSSA.

METHOD

We start from the definition of discrete fourier transform.

Discrete Fourier transform (DFT)

For a vector \mathbf{s} length of N, the discrete Fourier transform is defined as

$$\tilde{s}_k = \sum_{l=0}^{N-1} e^{-i2\pi kl/N} s_l \,, \tag{1}$$

where the index $k = 0, 1, 2, \dots, N-1$. Suppose the vector **s** represent a time series length of N with time sampling interval is Δt . So the frequency sampling interval Δf can be computed as

$$\Delta f = 1/\Delta t/N = \frac{1}{N\Delta t} \tag{2}$$

The Fourier transform for this time series in continuous case can be expressed as

$$\tilde{s}(f) = \int e^{-i2\pi f t} s(t) dt \tag{3}$$

In discrete case and assume the frequency f can be expressed as $f = k\Delta f$ and using the expression in equation 2

$$\tilde{s}_{k\Delta f} = \sum_{l=0}^{N-1} e^{-i2\pi k\Delta f l\Delta t} s_{l\Delta t}$$

$$\tilde{s}_{k\Delta f} = \sum_{l=0}^{N-1} e^{-i2\pi k \frac{1}{N\Delta t} l\Delta t} s_{l\Delta t}$$

$$\tilde{s}_{k} = \sum_{l=0}^{N-1} e^{-i2\pi k l/N} s_{l}.$$
(4)

The the last equation is exactly same as the equation 1. Through this derivation, we can understand the physical meaning of the Fourier coefficients obtained via discrete Fourier transform. Note that the first coefficient corresponding to 0 frequency and so as to the time series. The first element of \mathbf{s} is the sample at 0 time.

The above operation can be expressed into the matrix-vector form. To simplify the definition of Fourier transform matrix, we define a variable $z = e^{-i\pi/N}$, so the Fourier transform matrix, which is a N-by-N square matrix and the element $F_{kl} = z^{kl}$. It can be expanded as

$$\mathbf{F}_{N} = \begin{bmatrix} 1 & 1 & \cdots & 1 \\ 1 & z & \cdots & z^{N-1} \\ 1 & z^{2} & \cdots & z^{2(N-1)} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & z^{N-1} & \cdots & z^{(N-1)(N-1)} \end{bmatrix}$$
 (5)

As \mathbf{F}_N is a orthogonal (but not orthonormal) matrix, we can be easily get the its inverse matrix

$$\mathbf{F}_N^{-1} = \frac{1}{N} \mathbf{F}_N^H = \frac{1}{N} \mathbf{F}_N^* \tag{6}$$

where the super-script H represent complex conjugate transpose and * indicate conjugate transpose. As the matrix \mathbf{F}_{N} is a symmetrical matrix, that's why we can get equation 6.

Circulant matrix

An 9×9 circulant matrix is defined as

$$\mathbf{C} = \begin{bmatrix}
c_1 & c_9 & c_8 & c_7 & c_6 & c_5 & c_4 & c_3 & c_2 \\
c_2 & c_1 & c_9 & c_8 & c_7 & c_6 & c_5 & c_4 & c_3 \\
c_3 & c_2 & c_1 & c_9 & c_8 & c_7 & c_6 & c_5 & c_4 \\
c_4 & c_3 & c_2 & c_1 & c_9 & c_8 & c_7 & c_6 & c_5 \\
c_5 & c_4 & c_3 & c_2 & c_1 & c_9 & c_8 & c_7 & c_6 \\
c_6 & c_5 & c_4 & c_3 & c_2 & c_1 & c_9 & c_8 & c_7 \\
c_7 & c_6 & c_5 & c_4 & c_3 & c_2 & c_1 & c_9 & c_8 \\
c_8 & c_7 & c_6 & c_5 & c_4 & c_3 & c_2 & c_1 & c_9 \\
c_9 & c_8 & c_7 & c_6 & c_5 & c_4 & c_3 & c_2 & c_1
\end{bmatrix} \tag{7}$$

A circulant matrix is fully specified by its first column $\mathbf{c} = \begin{bmatrix} c_1 & c_2 & c_3 & c_4 & c_5 & c_6 & c_7 & c_8 & c_9 \end{bmatrix}^T$, we know that the product between circulant matrix and a vector \mathbf{v} length of N is equivalent to the circular convolution between vector $\mathbf{b}f\mathbf{c}$ and \mathbf{v} , and the circular convolution can be computed efficiently via Fast Fourier transform (FFT). So we have

$$\mathbf{C}\mathbf{v} = \mathbf{F}_N^{-1} \left(\mathbf{F}_N \mathbf{c} \circ \mathbf{F}_N \mathbf{v} \right) = \mathbf{F}_N^{-1} \left(diag(\mathbf{F}_N \mathbf{c}) \mathbf{F}_N \mathbf{v} \right) , \tag{8}$$

where the symbol \circ and diag represents Hardmard (element-wise) multiplication and building diagonal matrix from a vector.

Toeplitz matrix

Toeplitz matrix is associated with linear convolution, suppose we have two vectors $\mathbf{s} = \begin{bmatrix} s_1 & s_2 & s_3 & s_4 & s_5 \end{bmatrix}^T$ and $\mathbf{f} = \begin{bmatrix} f_1 & f_2 & f_3 & f_4 & f_5 \end{bmatrix}^T$, the linear convolution of these two vectors is vector \mathbf{r} length of 9. We can represent the convolution into matrix-vector form as

$$\mathbf{r} = \mathbf{h} * \mathbf{f} = \begin{bmatrix} r_1 \\ r_2 \\ r_3 \\ r_4 \\ r_5 \\ r_6 \\ r_7 \\ r_8 \\ r_9 \end{bmatrix} = \begin{bmatrix} s_1 \\ s_2 & s_1 \\ s_3 & s_2 & s_1 \\ s_4 & s_3 & s_2 & s_1 \\ s_5 & s_4 & s_3 & s_2 \\ s_5 & s_4 & s_3 & s_2 \\ s_5 & s_4 & s_3 & s_2 \\ s_5 & s_4 & s_3 \\ s_5 & s_4 & s_3 \end{bmatrix} \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ f_4 \\ f_5 \end{bmatrix}$$
(9)

The efficient linear convolution usually padding 0 to vectors make their length equal to 9, after padding zeros, we get new vector $\hat{\mathbf{s}} = \begin{bmatrix} s_1 & s_2 & s_3 & s_4 & s_5 & 0 & 0 & 0 \end{bmatrix}$ and $\hat{\mathbf{f}} = \begin{bmatrix} f_1 & f_2 & f_3 & f_4 & f_5 & 0 & 0 & 0 \end{bmatrix}$, the circular convolution between $\hat{\mathbf{s}}$ and $\hat{\mathbf{f}}$ can be expressed as

$$\hat{\mathbf{r}} = \hat{\mathbf{s}} \circledast \hat{\mathbf{f}} = \begin{bmatrix} s_1 & 0 & 0 & 0 & s_5 & s_4 & s_3 & s_2 \\ s_2 & s_1 & 0 & 0 & 0 & 0 & s_5 & s_4 & s_3 \\ s_3 & s_2 & s_1 & 0 & 0 & 0 & 0 & s_5 & s_4 \\ s_4 & s_3 & s_2 & s_1 & 0 & 0 & 0 & 0 & s_5 \\ s_5 & s_4 & s_3 & s_2 & s_1 & 0 & 0 & 0 & 0 \\ 0 & s_5 & s_4 & s_3 & s_2 & s_1 & 0 & 0 & 0 \\ 0 & 0 & s_5 & s_4 & s_3 & s_2 & s_1 & 0 & 0 \\ 0 & 0 & 0 & s_5 & s_4 & s_3 & s_2 & s_1 & 0 \\ 0 & 0 & 0 & s_5 & s_4 & s_3 & s_2 & s_1 & 0 \\ 0 & 0 & 0 & s_5 & s_4 & s_3 & s_2 & s_1 \end{bmatrix} \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ f_4 \\ f_5 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

When can see that $\hat{\mathbf{r}} = \mathbf{r}$, so linear convolution can be efficiently computed by padding zeros to vectors then compute circular convolution in frequency domain.

Let's consider another Toeplitz matrix built from a vector $\mathbf{c} = \begin{bmatrix} c_1 & c_2 & c_3 & c_4 & c_5 & c_6 & c_7 & c_8 & c_9 \end{bmatrix}$ with length N = 9, we first determine two integer: L and K make L + K - 1 = N. To make the matrix close to square form, we determine them by

$$L = \left| \frac{N}{2} \right| + 1, \ K = N + 1 - L \tag{11}$$

in this special case, N=9, so L=K=5, the Toeplitz matrix build from this vector can be expressed as

$$\mathbf{T} = \begin{bmatrix} c_5 & c_4 & c_3 & c_2 & c_1 \\ c_6 & c_5 & c_4 & c_3 & c_2 \\ c_7 & c_6 & c_5 & c_4 & c_3 \\ c_8 & c_7 & c_6 & c_5 & c_4 \\ c_9 & c_8 & c_7 & c_6 & c_5 \end{bmatrix}$$
(12)

The vector \mathbf{c} can be uniquely determined from the last column and last row. There are two way to embedding this Teoplitz matrix into a circular matrix, the first to embedding this matrix at the lower-left part of the circular matrix created by vector \mathbf{c} , which shows as below as

The above circulant matrix is called C_1 . We can also create a new vector $\hat{\mathbf{c}} = \begin{bmatrix} c_5 & c_6 & c_7 & c_8 & c_9 & c_1 & c_2 & c_3 \end{bmatrix}$ which is obtained by permutating the elements of the original vector, the Toeplitz matrix can be embedded at the upper-left part of the circulant matrix created from $\hat{\mathbf{c}}$ shows as bellow

and call this new matrix as C_2 .

Suppose we need to compute the product between Toeplitz matrix T and a vector v

$$\mathbf{r} = \mathbf{T}\mathbf{v} \tag{13}$$

where the length of \mathbf{v} is K. There are two way to do it efficiently via FFT, and these two ways corresponding to the two embedding methods. The first step is to padding vector \mathbf{v} with L-1 0s get a new vector $\hat{\mathbf{v}}$. Following above example, K=L=5 and $\mathbf{v}=\begin{bmatrix}v_1 & v_2 & v_3 & v_4 & v_5\end{bmatrix}^T$, and the padded vector $\hat{\mathbf{v}}=\begin{bmatrix}v_1 & v_2 & v_3 & v_4 & v_5 & 0 & 0 & 0\end{bmatrix}^T$, The product between \mathbf{C}_1 and $\hat{\mathbf{v}}$ are

$$\hat{\mathbf{r}}_{1} = \mathbf{C}_{1}\hat{\mathbf{v}} = \begin{bmatrix}
c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} \\
c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} \\
c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} \\
c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} \\
c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} \\
c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} \\
c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} \\
c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} \\
c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} \\
c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1}
\end{bmatrix} \begin{bmatrix} v_{1} \\ v_{2} \\ v_{3} \\ v_{4} \\ v_{5} \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$(14)$$

where the last L elements of $\hat{\mathbf{r}}_1$ equal to \mathbf{r} . On the other hand, we can take advantage of the circulate matrix \mathbf{C}_2 , we have

$$\hat{\mathbf{r}}_{2} = \mathbf{C}_{2}\hat{\mathbf{v}} = \begin{bmatrix} c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} \\ c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} \\ c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} \\ c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} & c_{9} \\ c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} & c_{1} \\ c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} & c_{2} \\ c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} & c_{3} \\ c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} & c_{4} \\ c_{4} & c_{3} & c_{2} & c_{1} & c_{9} & c_{8} & c_{7} & c_{6} & c_{5} \end{bmatrix} \begin{bmatrix} v_{1} \\ v_{2} \\ v_{3} \\ v_{4} \\ v_{5} \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

$$(15)$$

So the first L elements of $\hat{\mathbf{r}}_2$ equal to \mathbf{r} . We summarize the efficient computation of the product between Toeplitz matrix and a vector in Algorithm 1 and Algorithm 2

Algorithm 1: Teoplitz-Times-Vector

```
1 Function \mathbf{v} = ttv(\mathbf{c}, \mathbf{v})

2 | length of \mathbf{c}: N = length(\mathbf{c})

3 | compute L = \lfloor \frac{N}{2} \rfloor, K = N + 1 - L

4 | padding L - 1 0s to the end of \mathbf{v} to get a new vector \hat{\mathbf{v}} with length of N

5 | \mathbf{r}_1 = \mathbf{F}_N^{-1} (\mathbf{F}_N \mathbf{c} \circ \mathbf{F}_N \hat{\mathbf{v}})

6 | take the last L elements from \mathbf{r}_1
```

Another method is summarized in Algorithm 2 as

Algorithm 2: Teoplitz-Times-Vector

```
1 Function \mathbf{v} = ttv(\mathbf{c}, \mathbf{v})

2 | length of \mathbf{c}: N = length(\mathbf{c})

3 | compute L = \lfloor \frac{N}{2} \rfloor, K = N + 1 - L

4 | padding L - 1 0s to the end of \mathbf{v} to get a new vector \hat{\mathbf{v}} with length of N

5 | take the first L - 1 element of \mathbf{c} to the end, get a new vector \hat{\mathbf{c}}

\mathbf{r}_2 = \mathbf{F}_N^{-1} (\mathbf{F}_N \hat{\mathbf{c}} \circ \mathbf{F}_N \hat{\mathbf{v}})

6 | take the first L elements from \mathbf{r}_2
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REFERENCES

Korobeynikov, A., 2009, Computation-and space-efficient implementation of ssa: arXiv preprint arXiv:0911.4498.