

Make flows small again: revisiting the flow framework

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[TACAS'23]

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Frame Rule

$$\frac{\{ P \} \ com \ \{ Q \}}{\{ P * F \} \ com \ \{ Q * F \}}$$

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Usage

$$\frac{\{ x \mapsto 5 \} [x] = 7 \{ x \mapsto 7 \}}{\{ x \mapsto 5 * F \} [x] = 7 \{ x \mapsto 7 * F \}}$$

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(small axioms)

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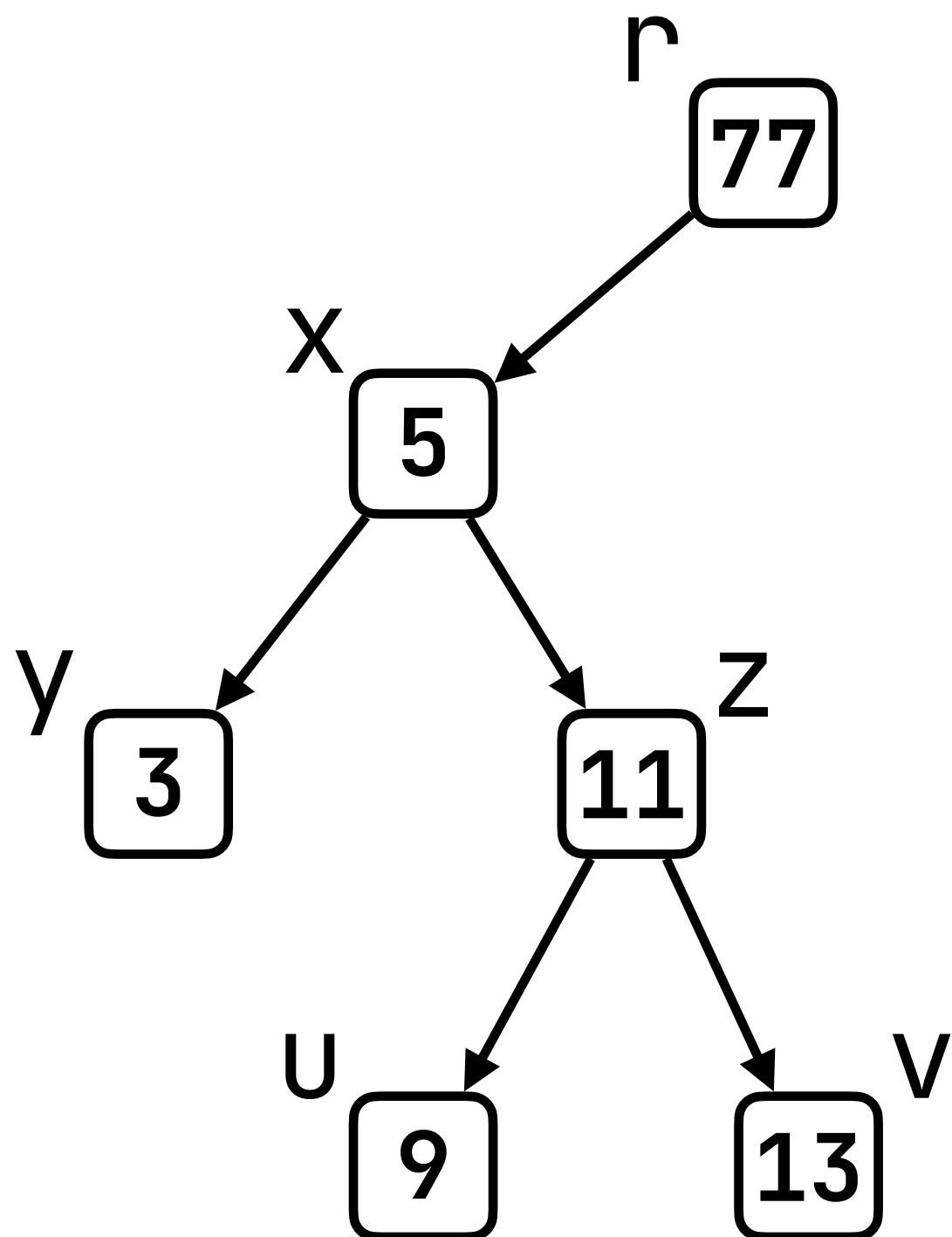
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Usage

Frame inference is a key challenge for proof automation.

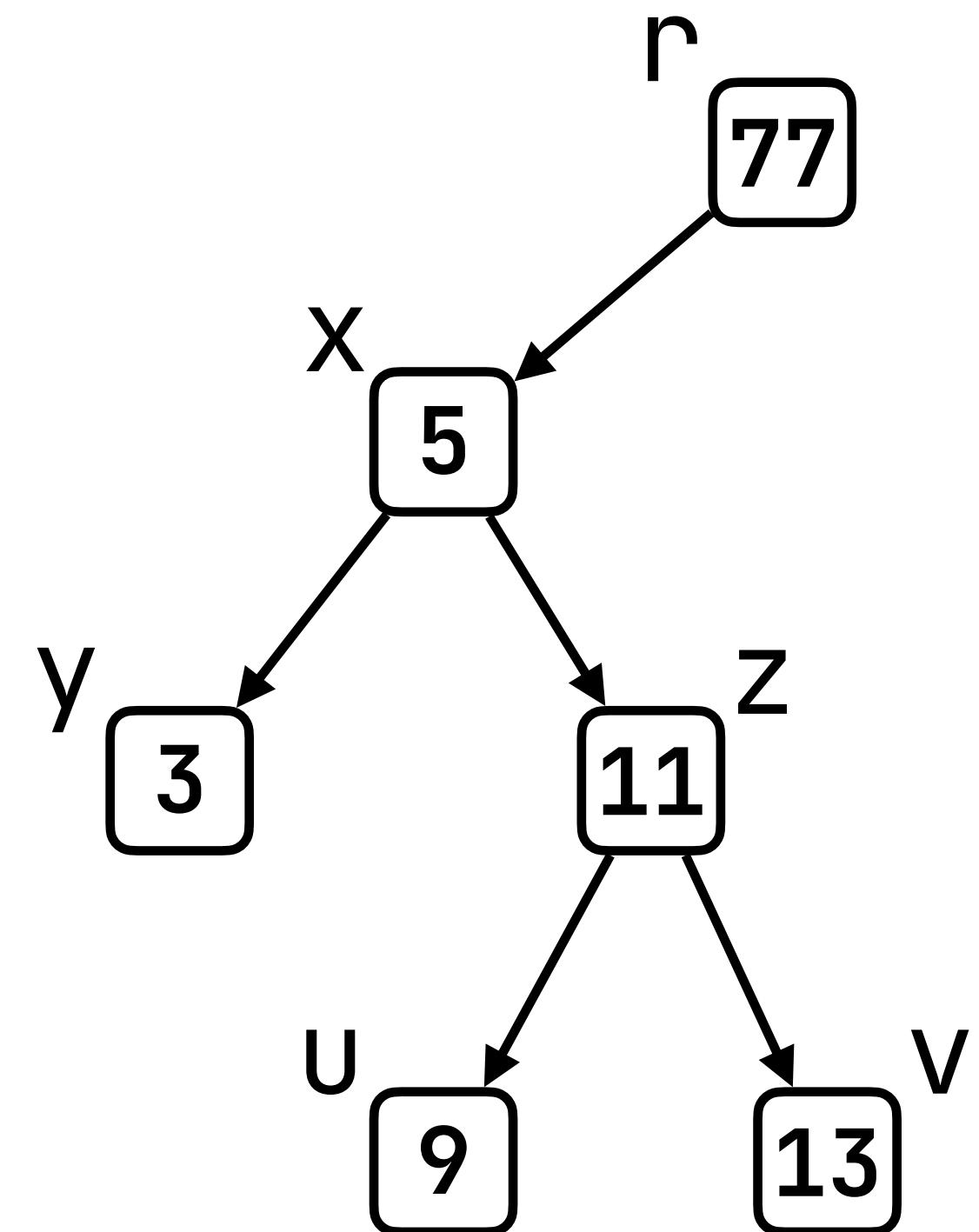
State

- Physical state
 - *heap graph*
 - e.g. $r \mapsto 77, x, \perp * x \mapsto 5, y, z * \dots$



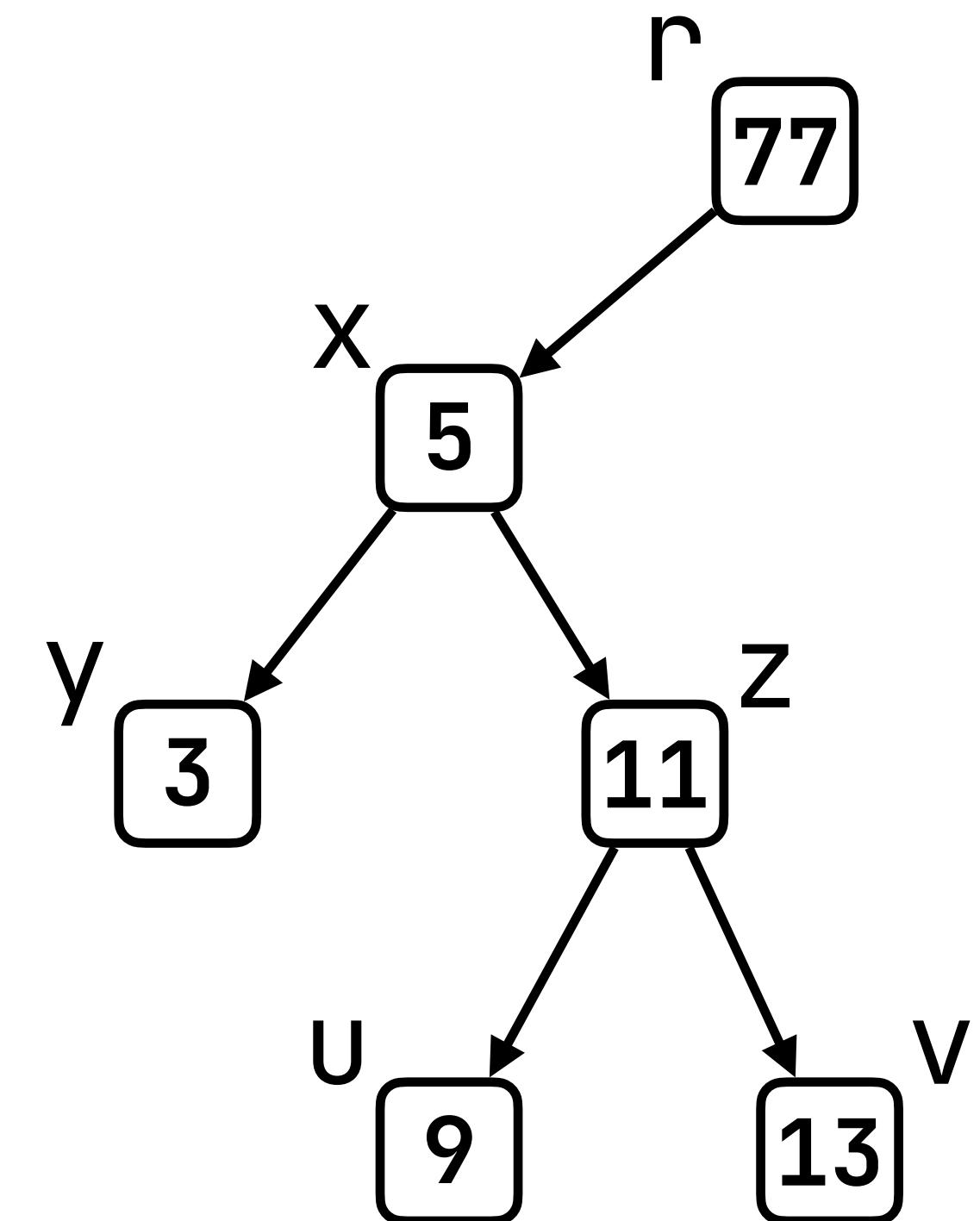
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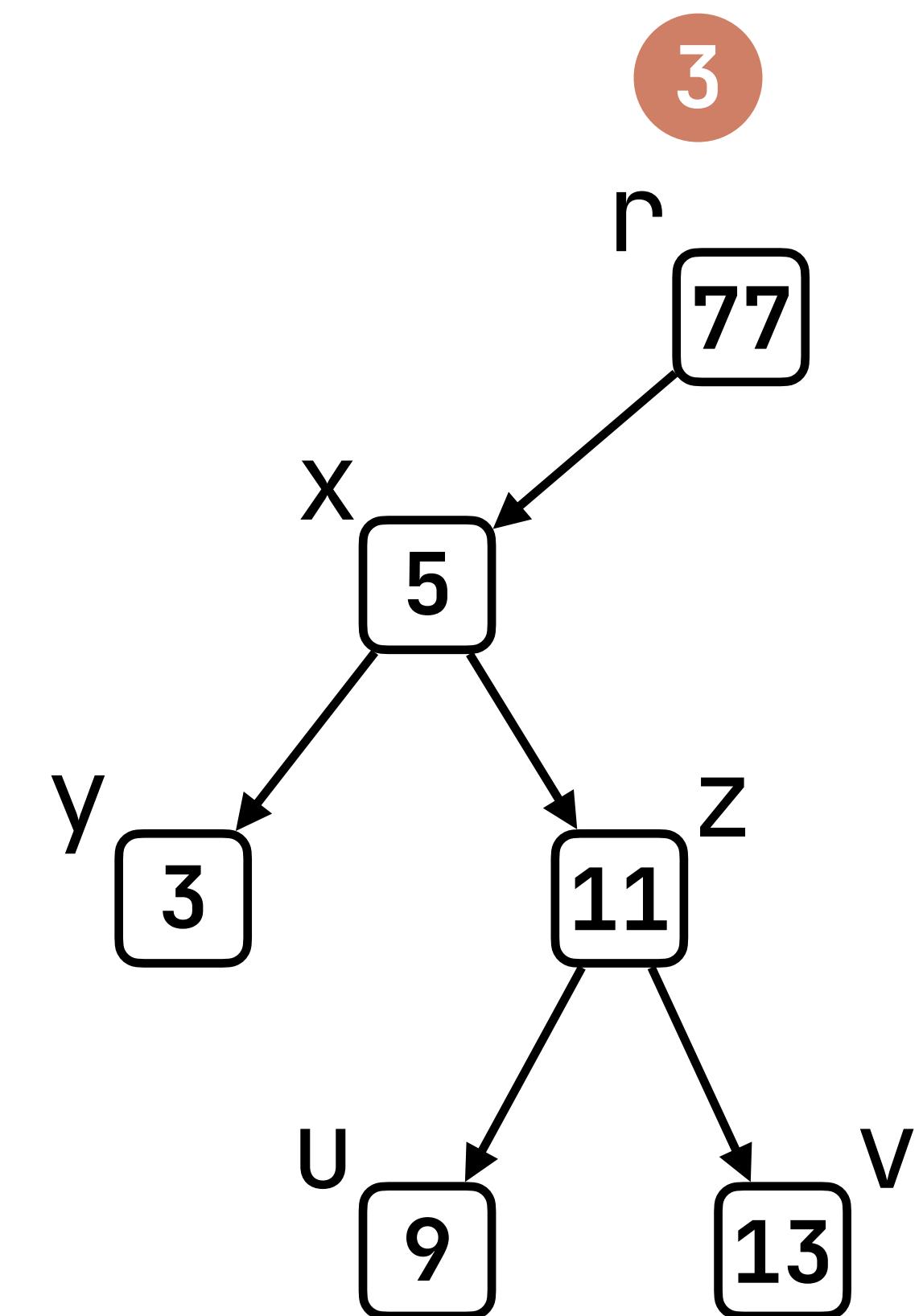
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to show linearizability of data structures



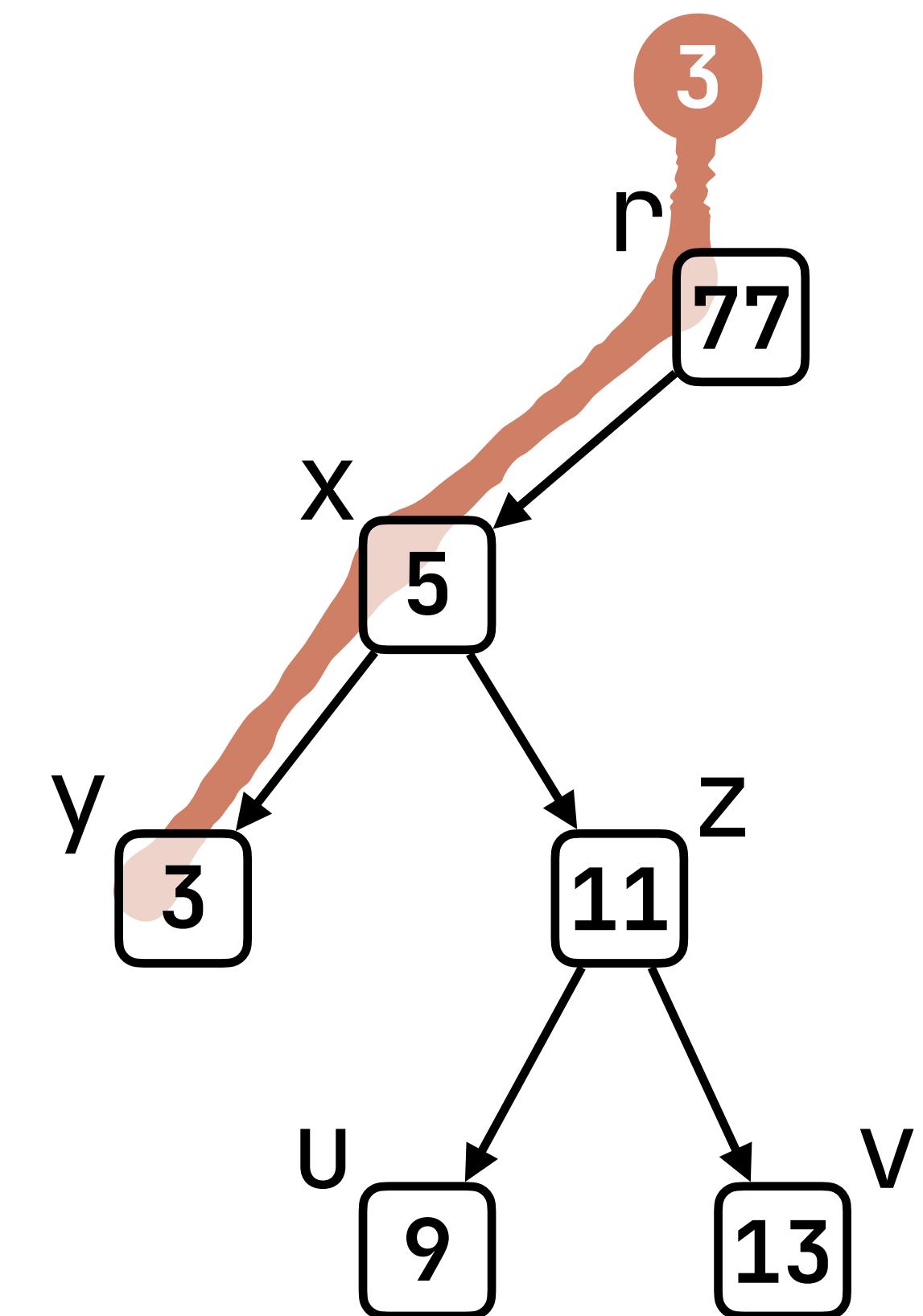
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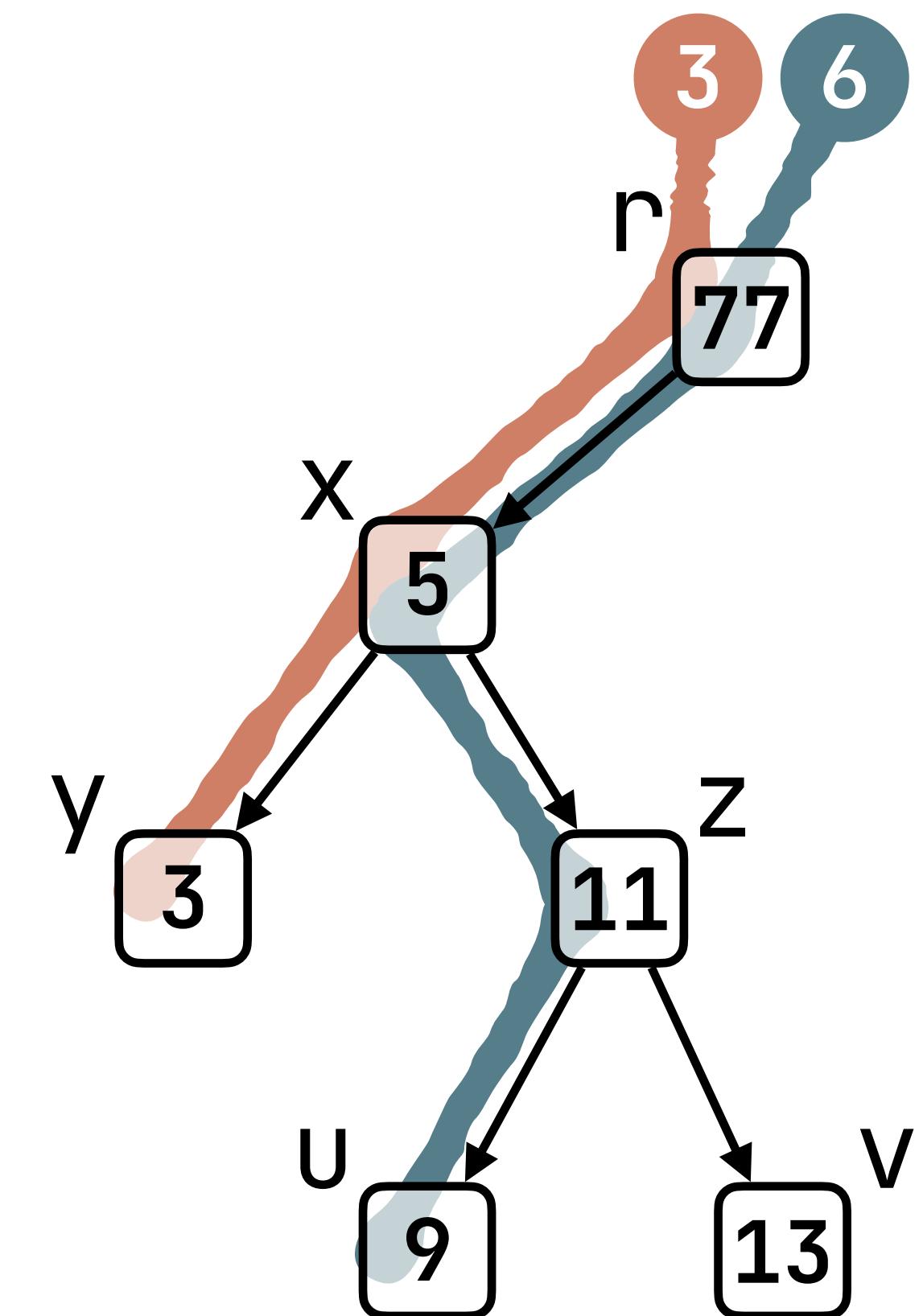
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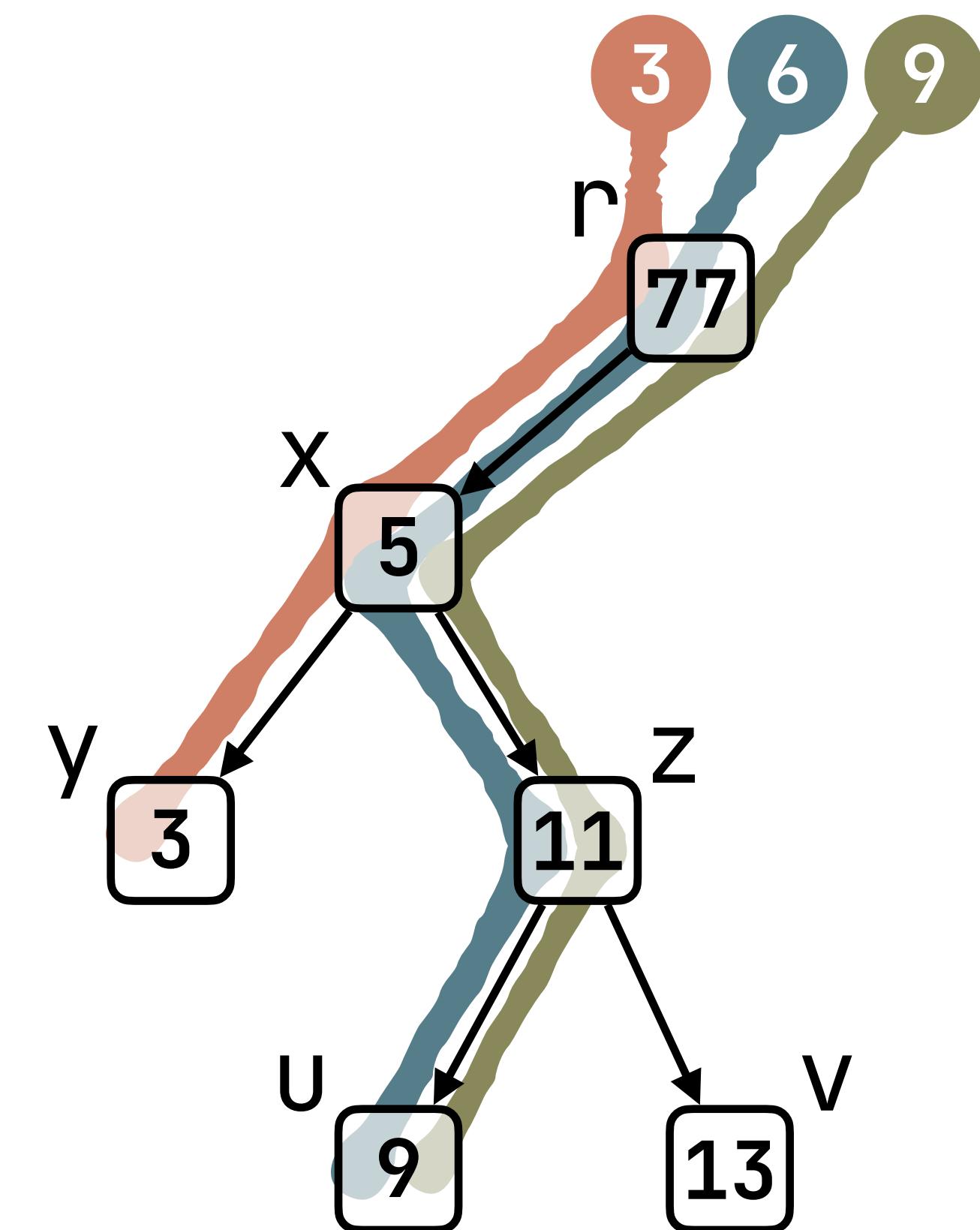
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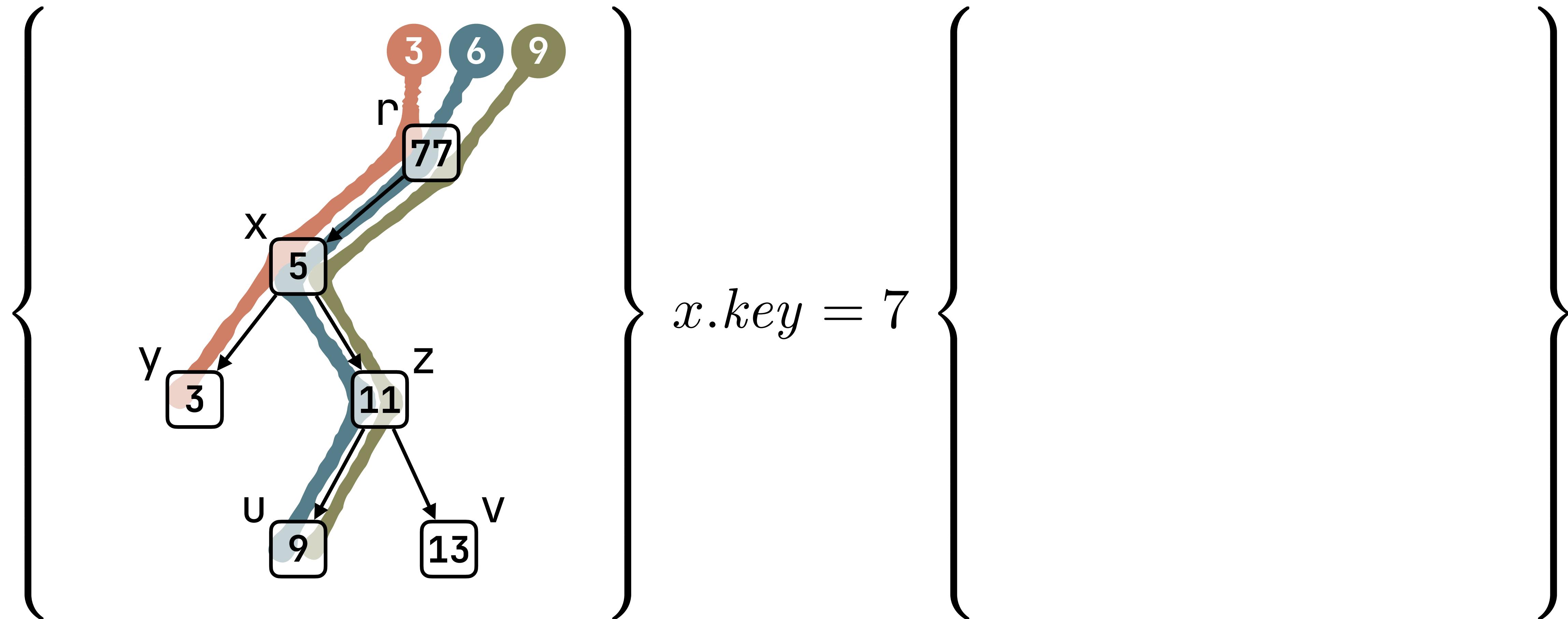


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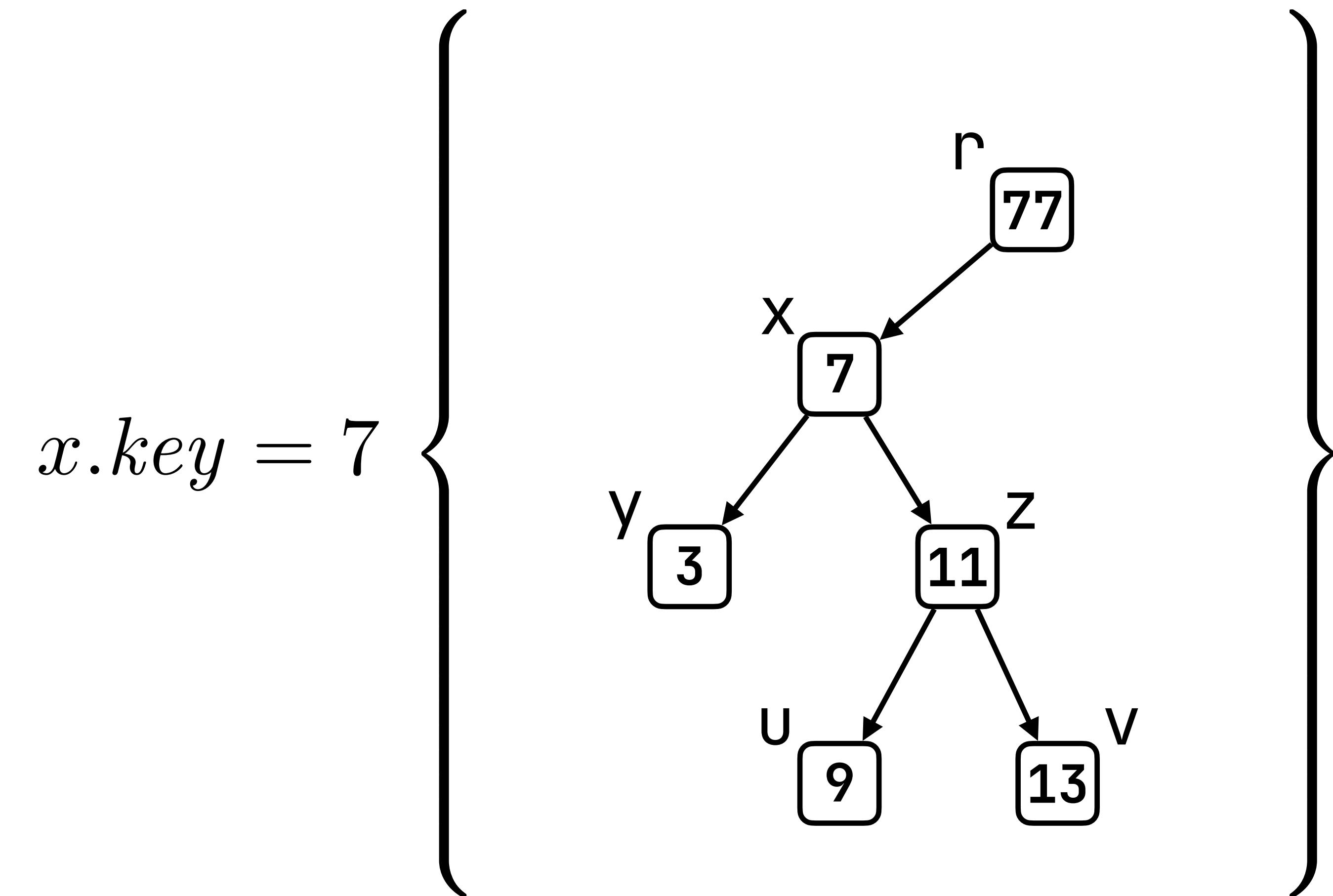
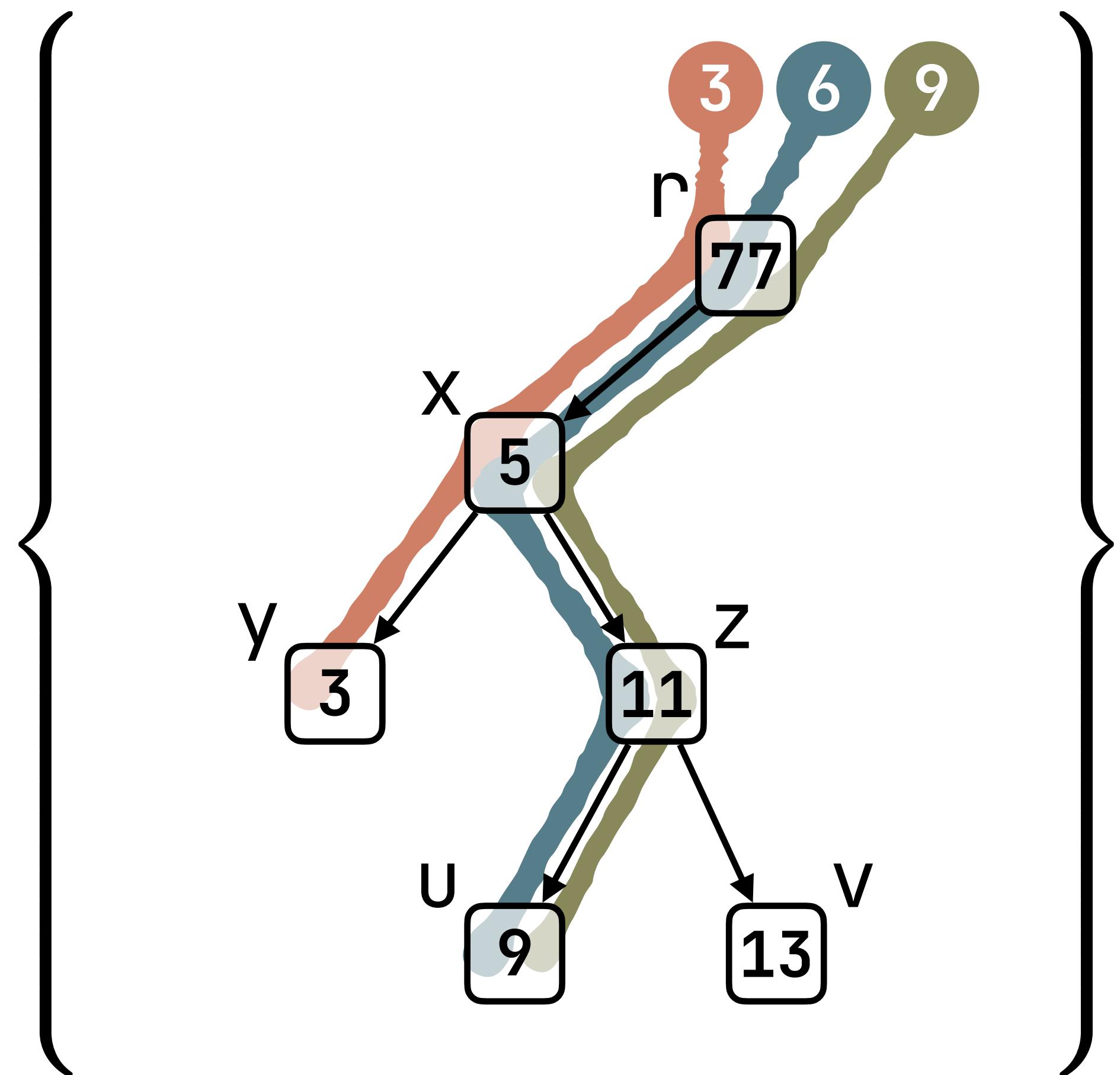
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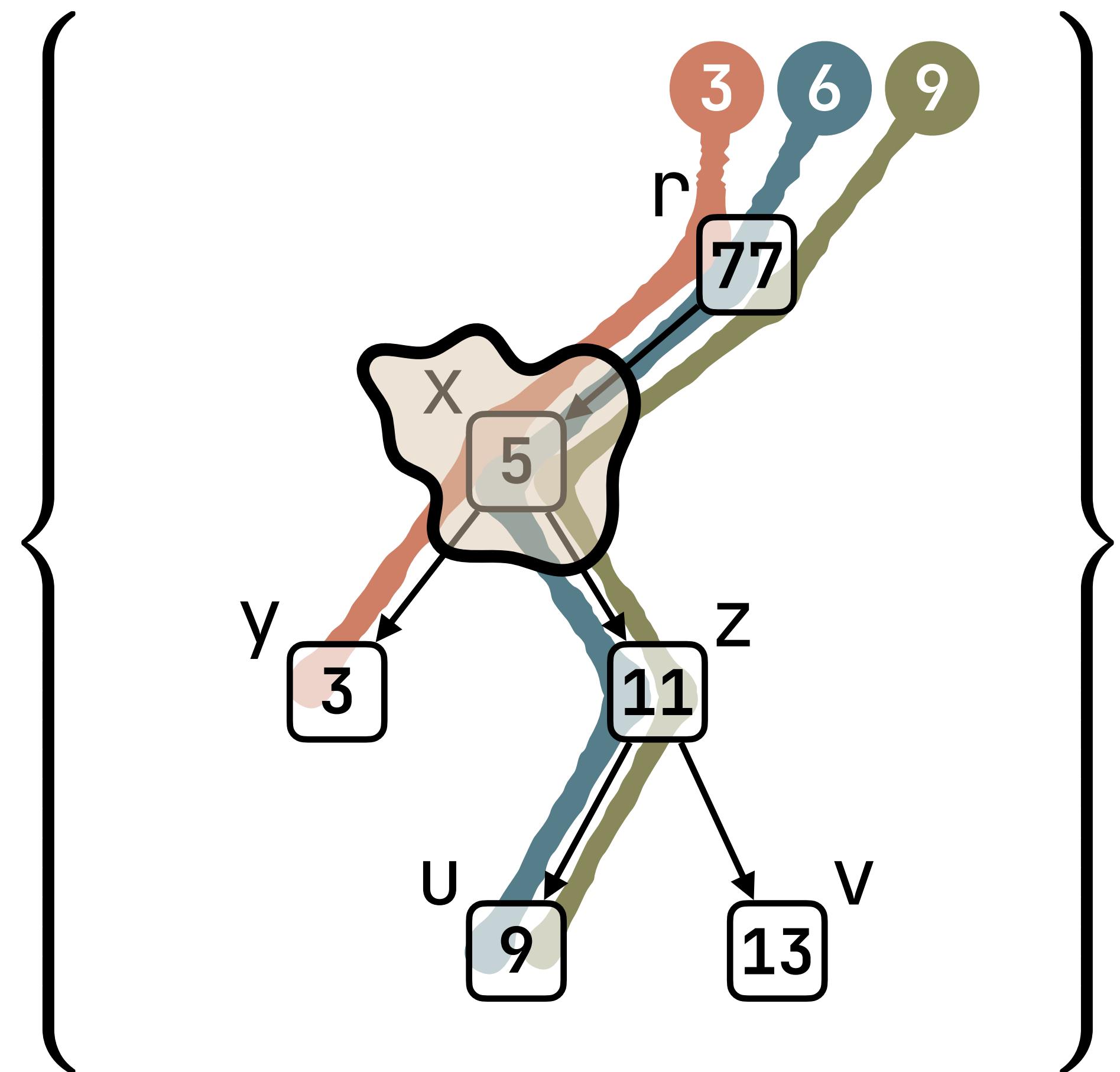
Goal: Frame Inference



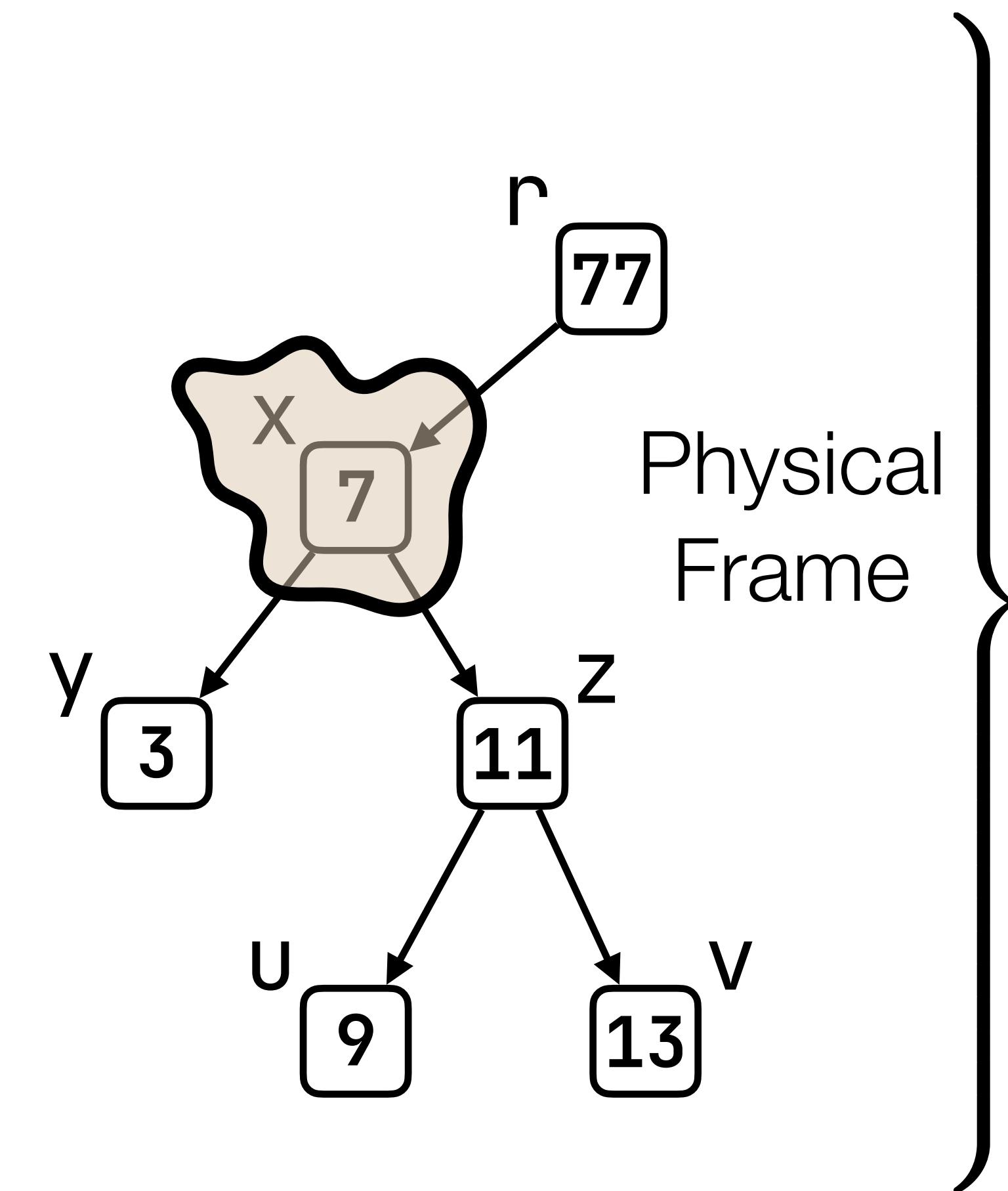
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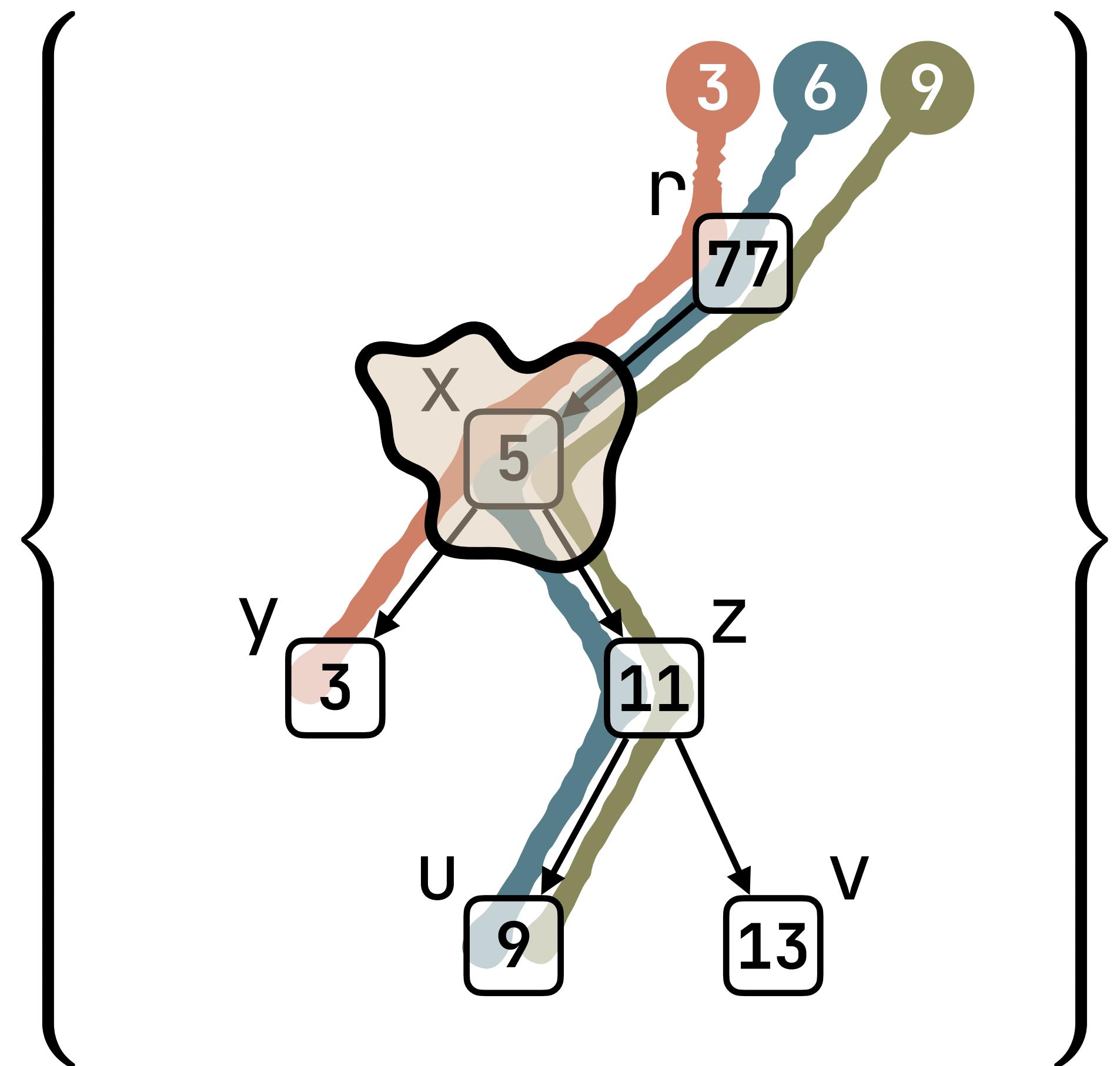


$$x.key = 7$$

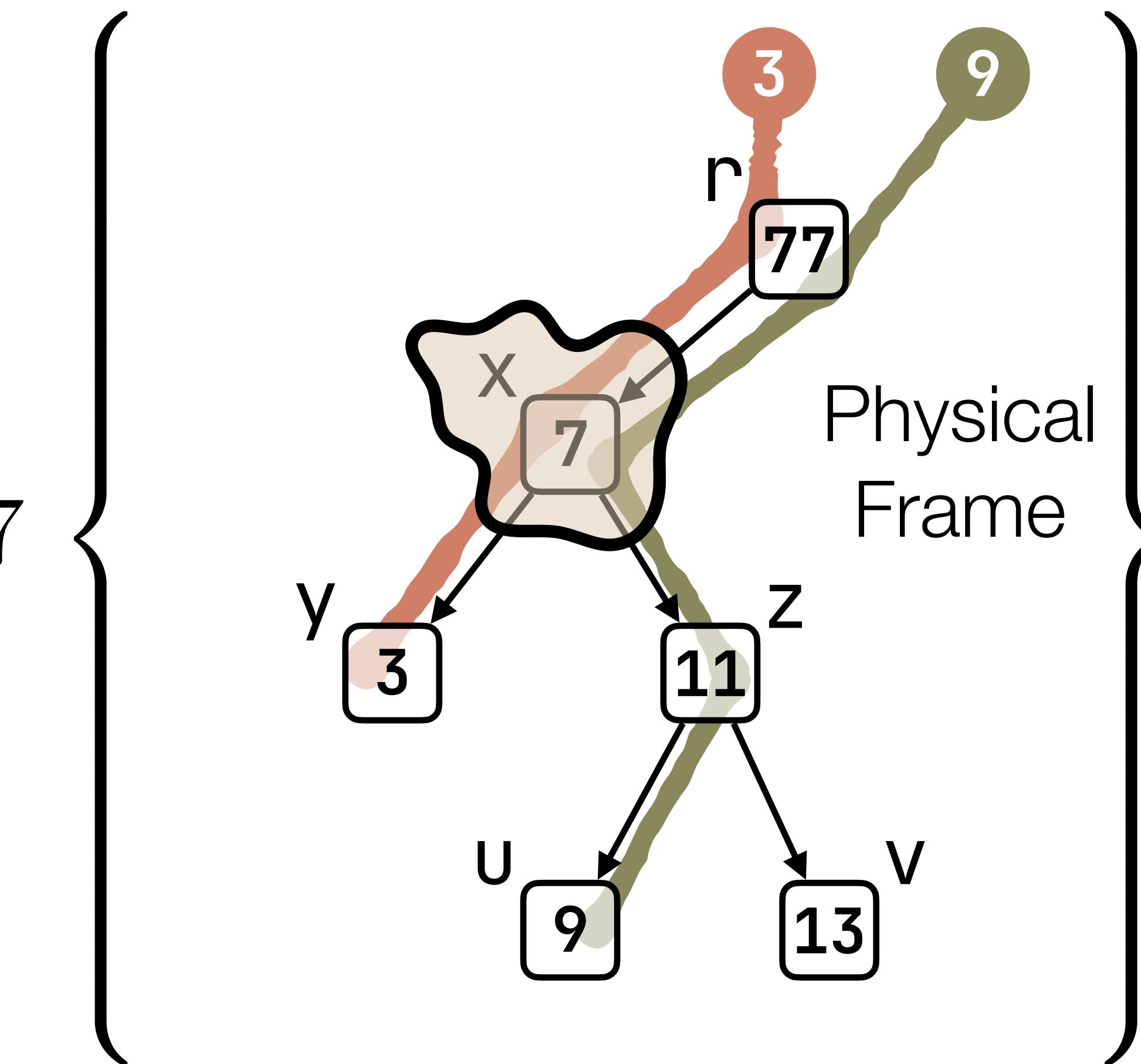


Physical
Frame

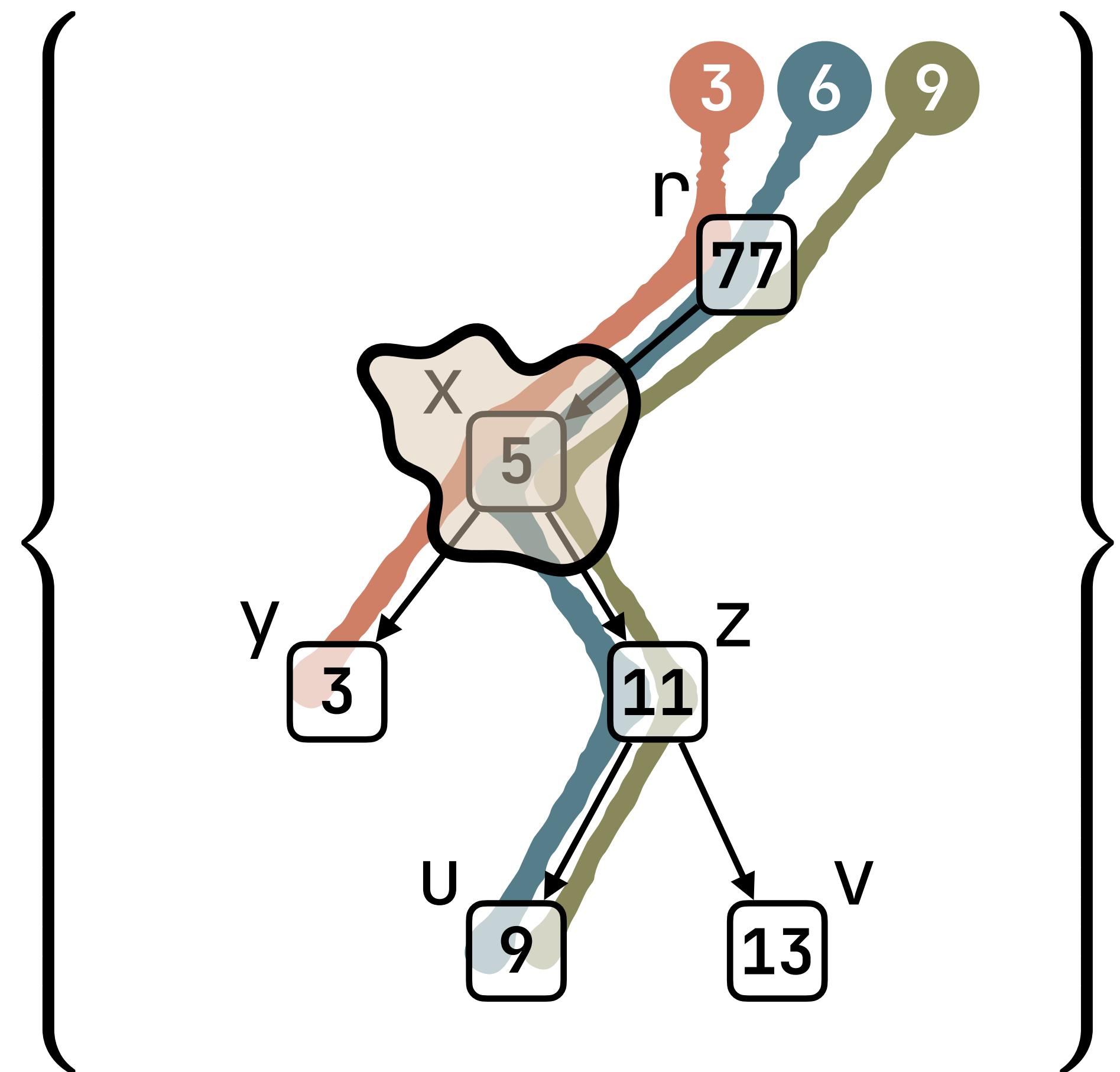
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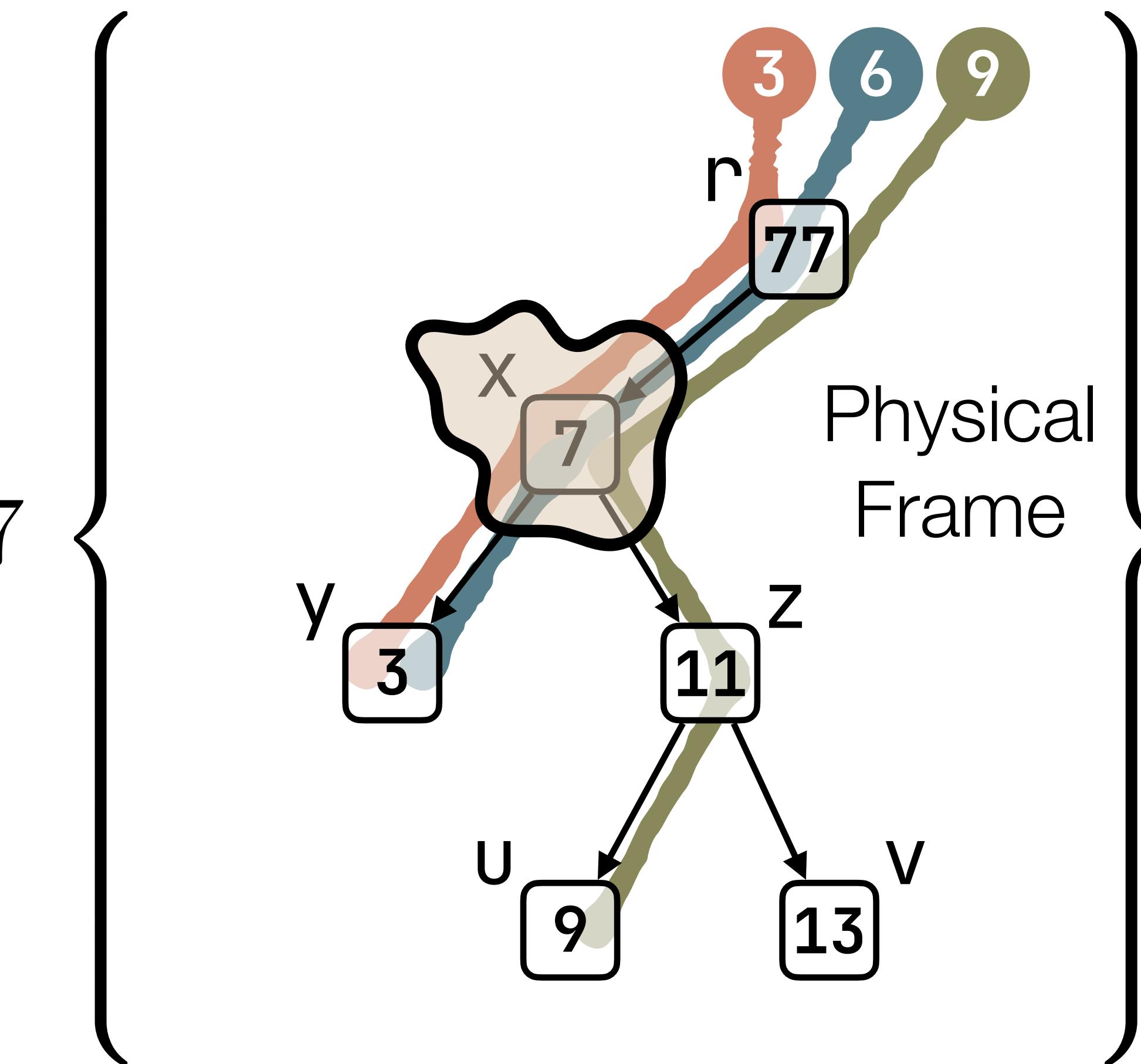
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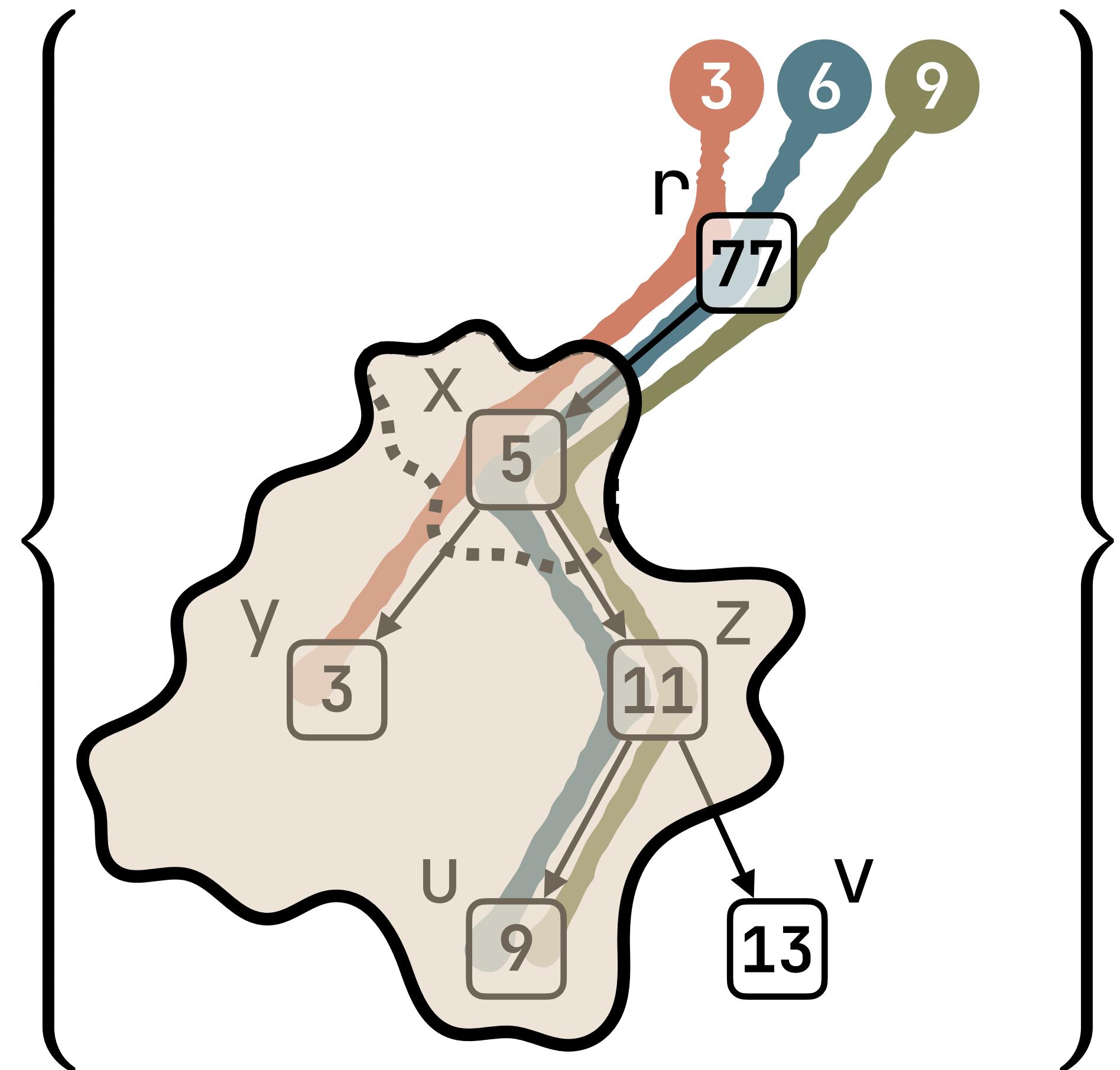
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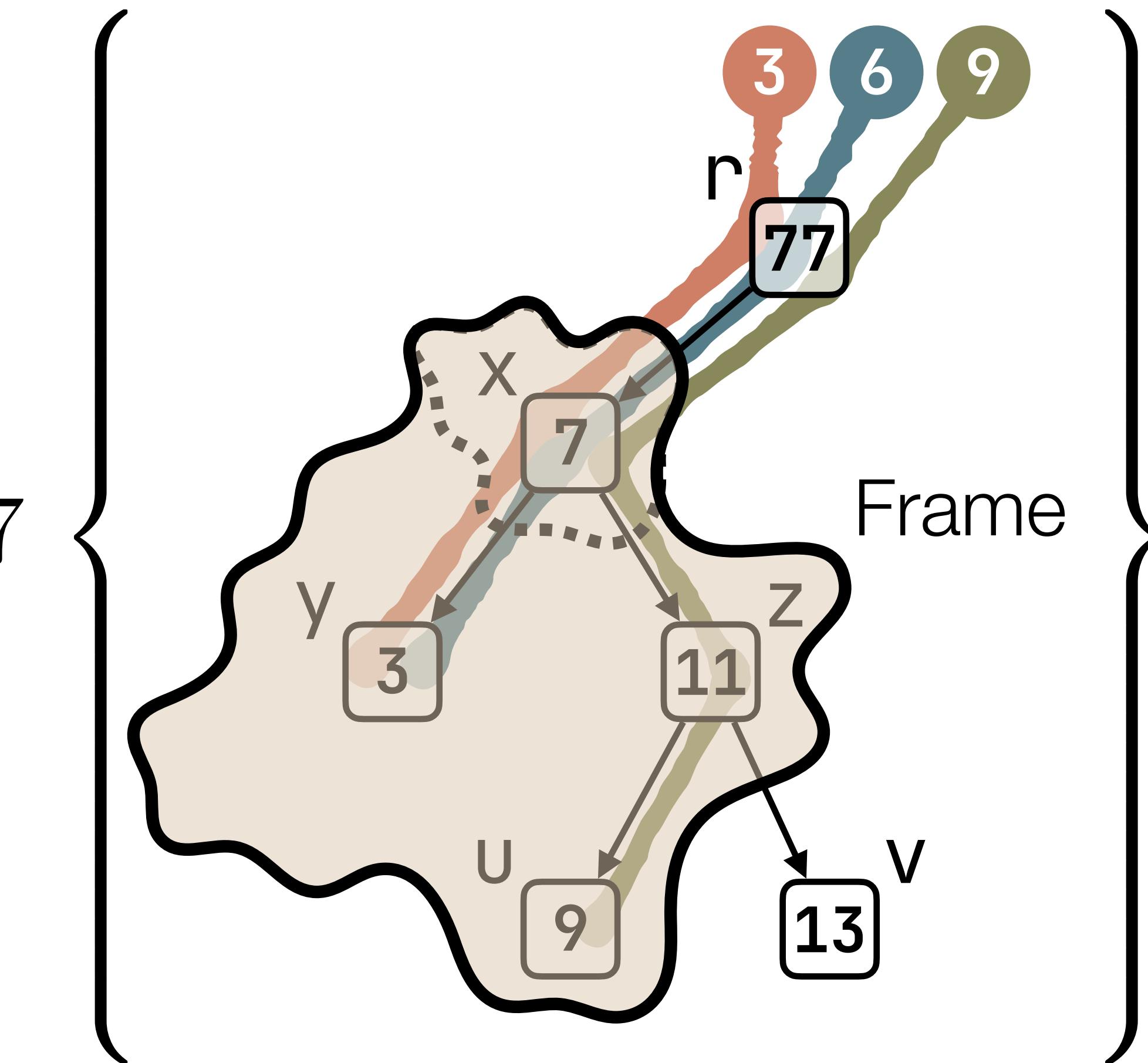
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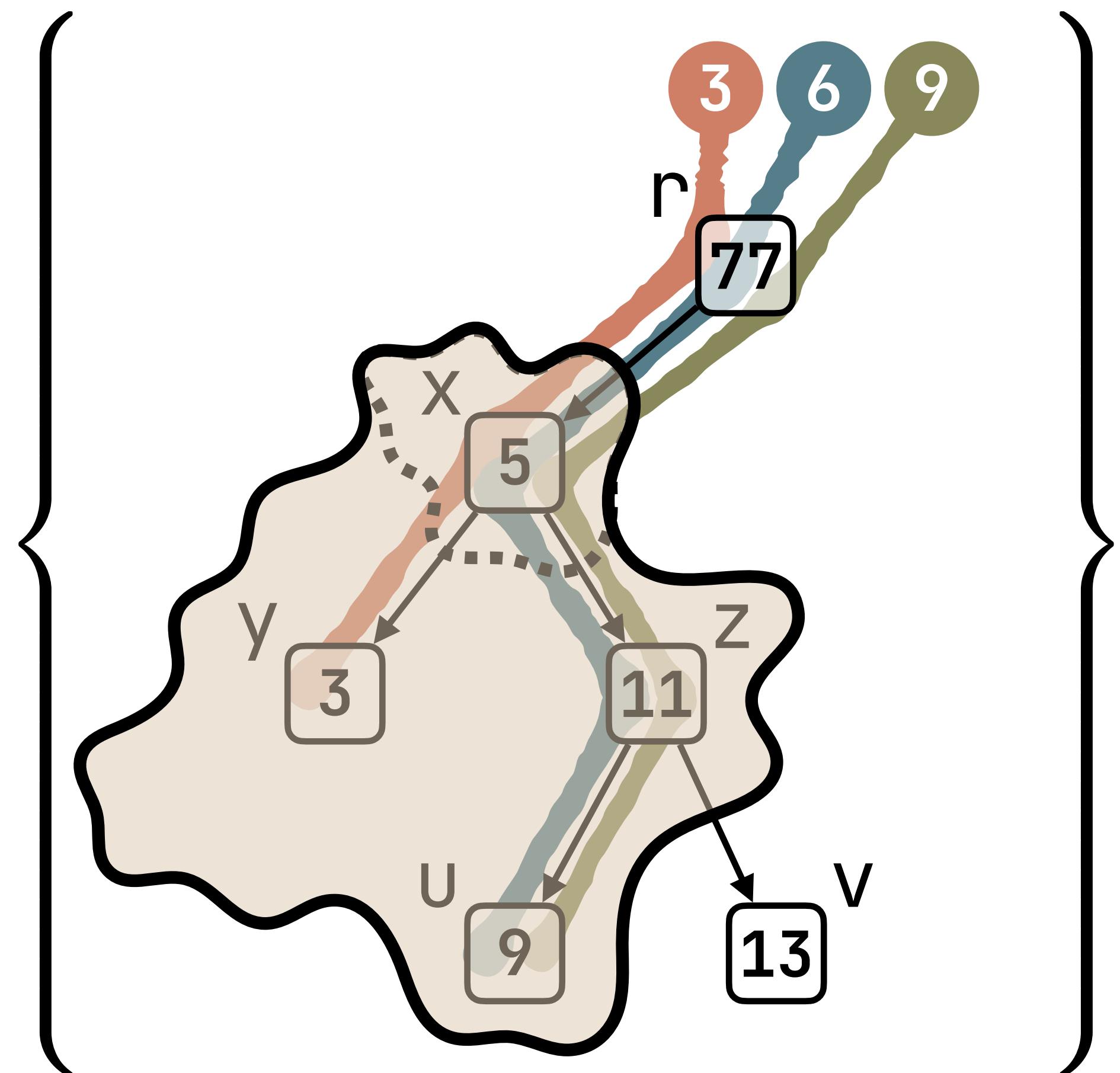
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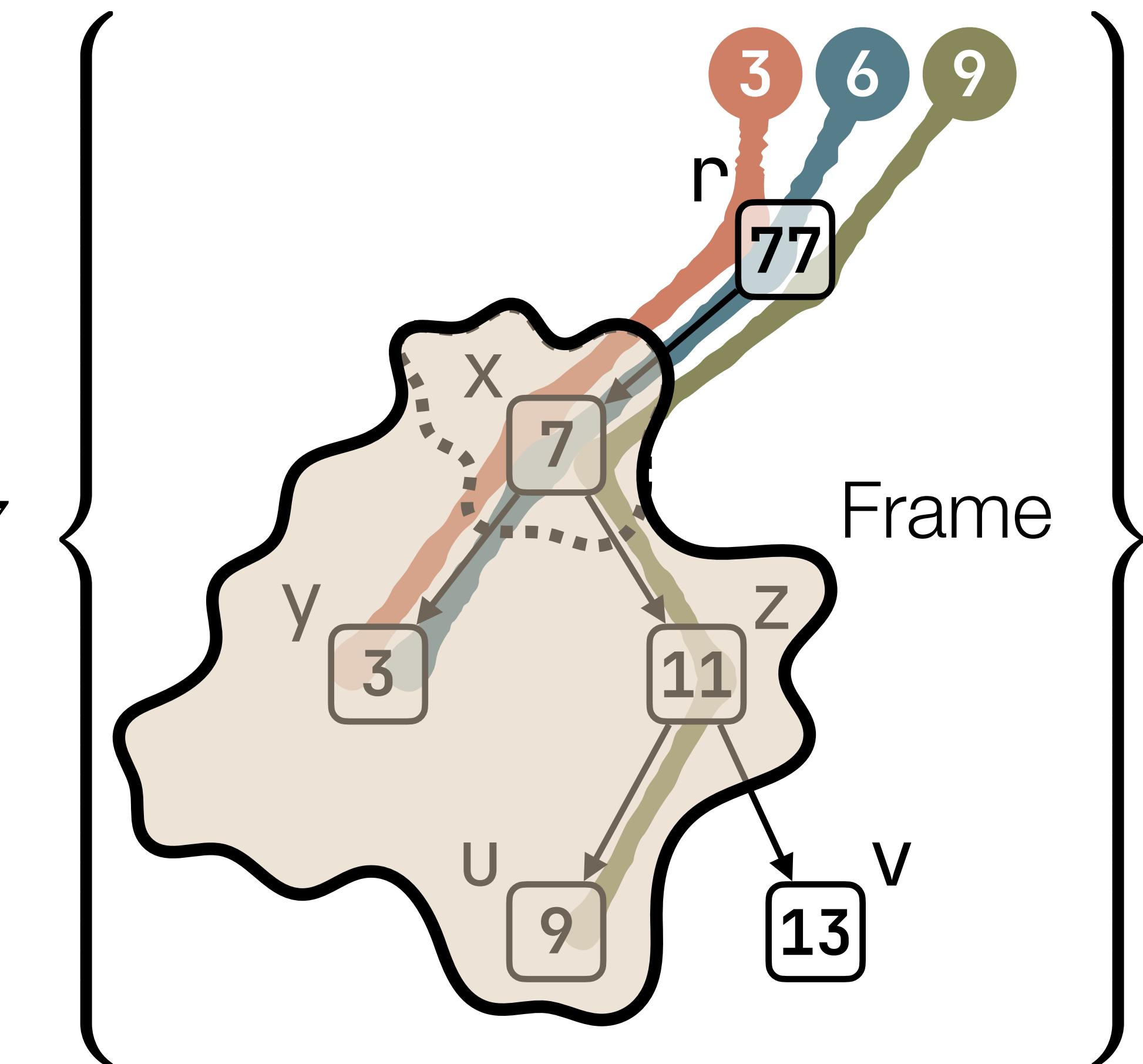
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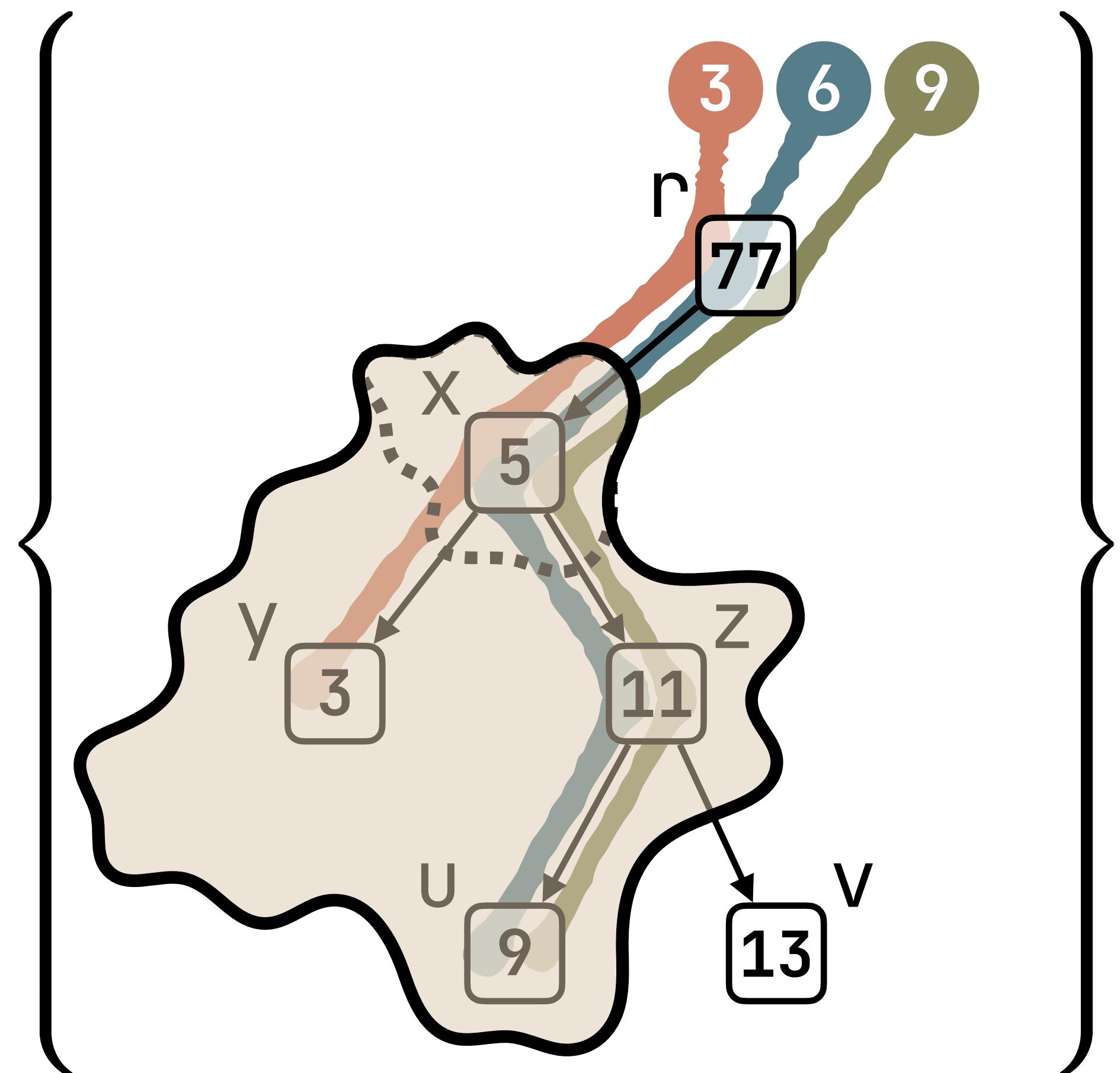


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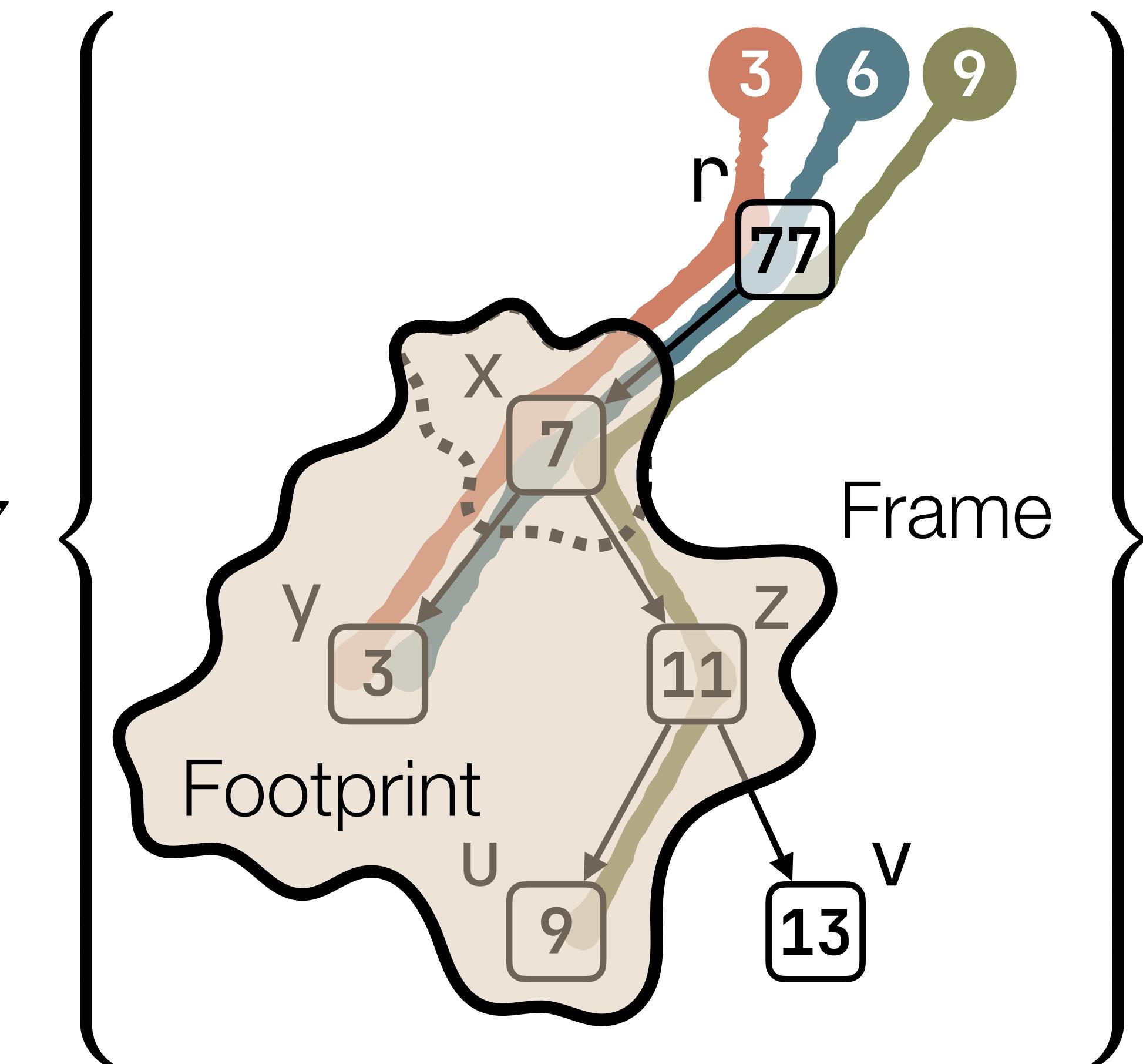


Goal: ***automatically*** find frame.

Goal: Frame Inference



$x.key = 7$



Goal: **automatically** find **footprint**.



Flow Framework

- Ghost state for heap graphs
- Inspired by data-flow analysis
- Formalizes inductive heap invariants



Frame Inference

- Separation & flows
- Frame-preserving updates
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Comparing Footprints

- Check if update is frame-preserving
- Efficient checks for general graphs



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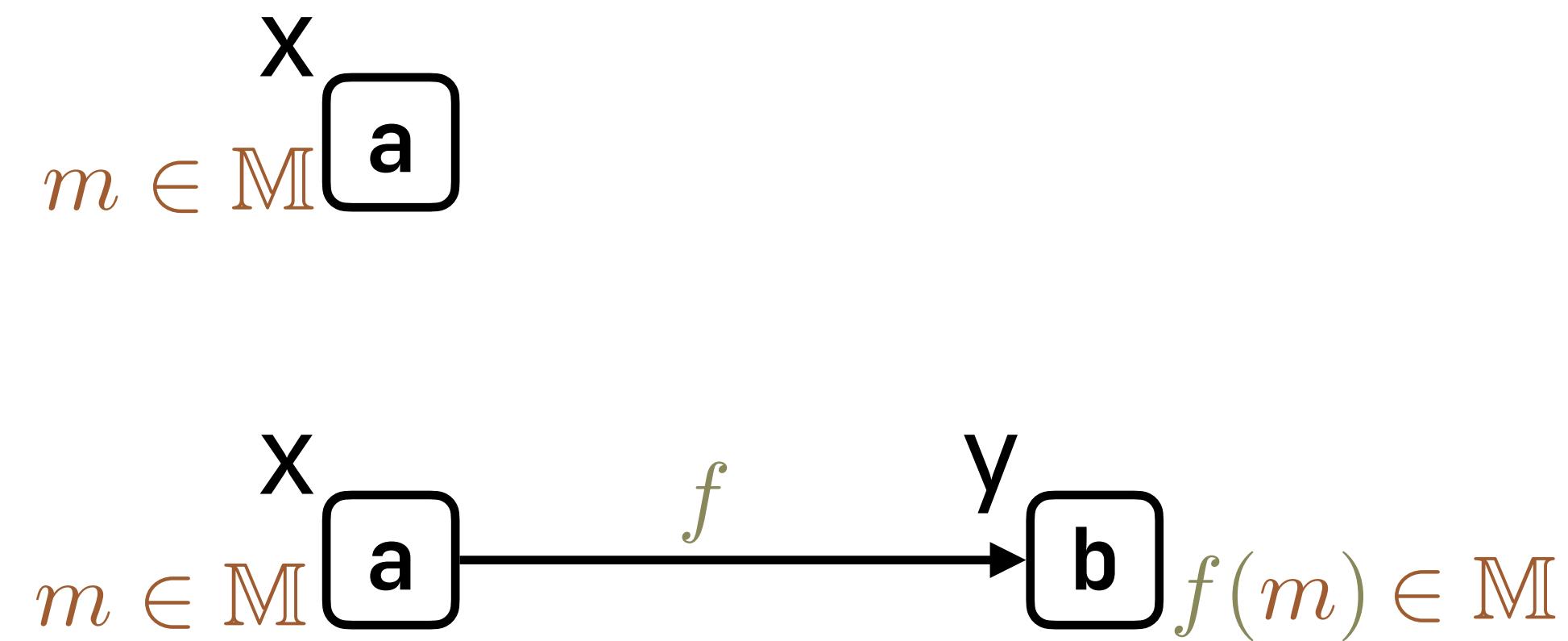
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$$m \in \text{M}^{\boxed{a}}$$

Flow Framework

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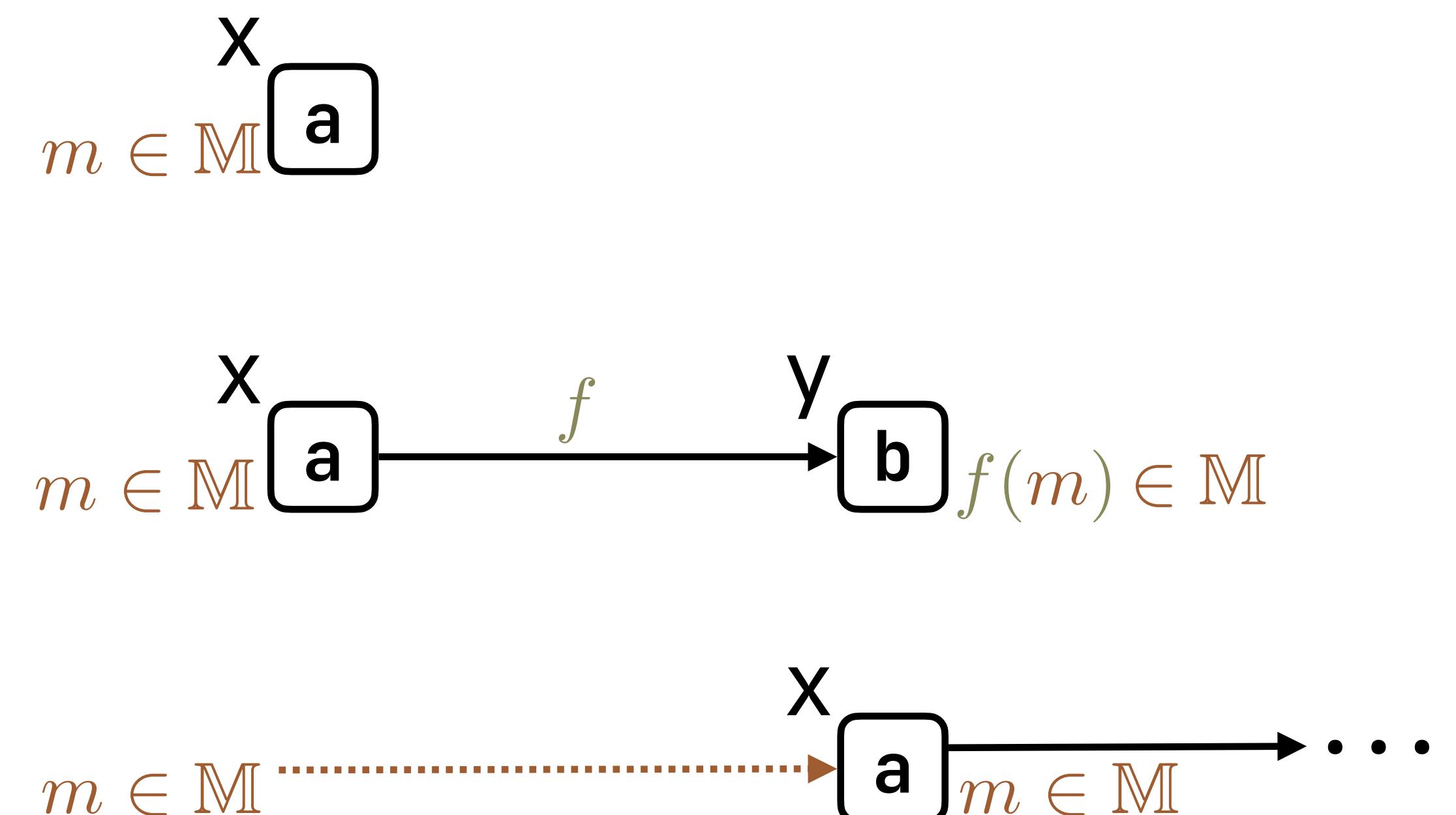
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- **The flow:** least fixed point, wrt. initial value



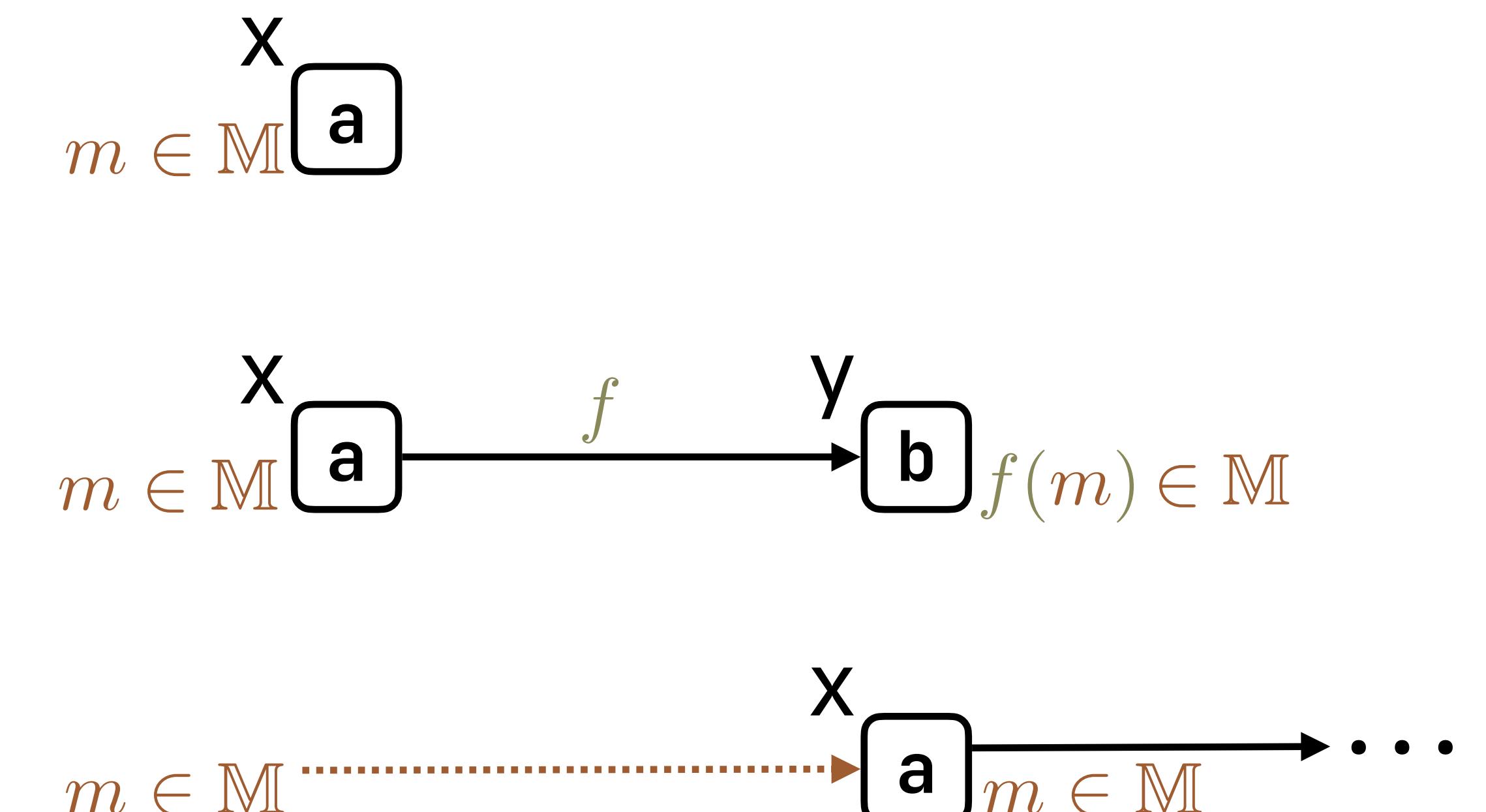
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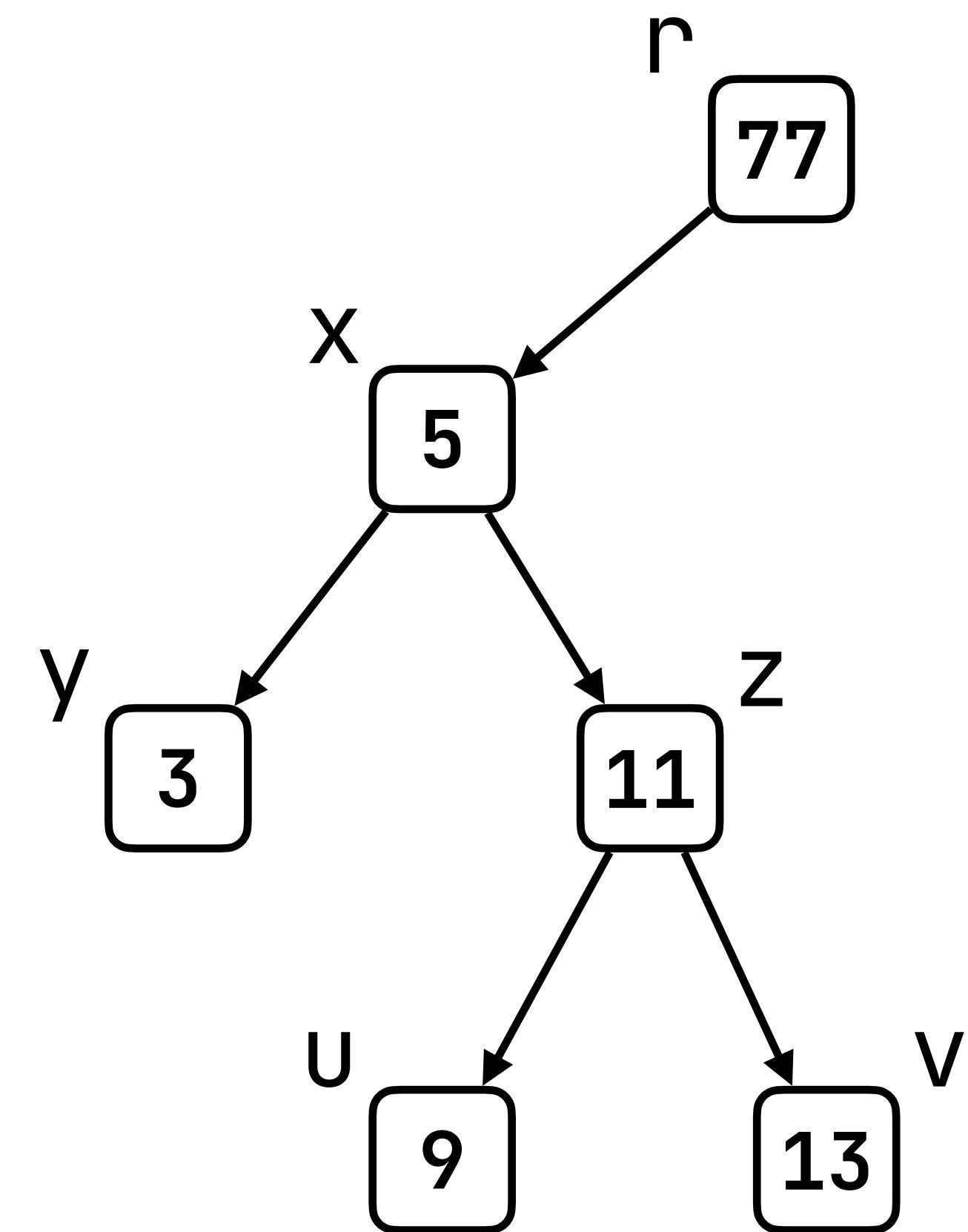
Exists if: \leq is ω -cpo and $+, \sup$ commute



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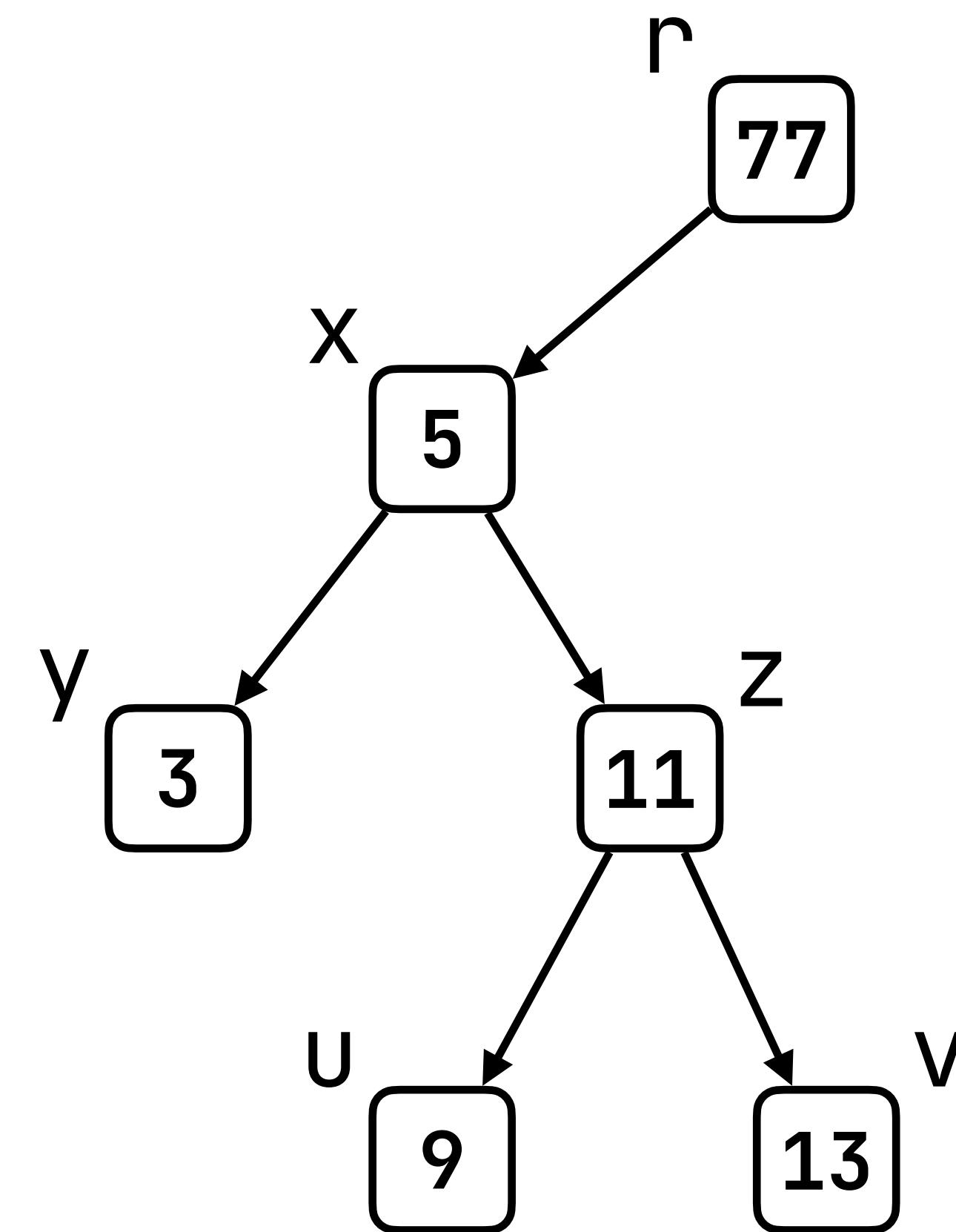
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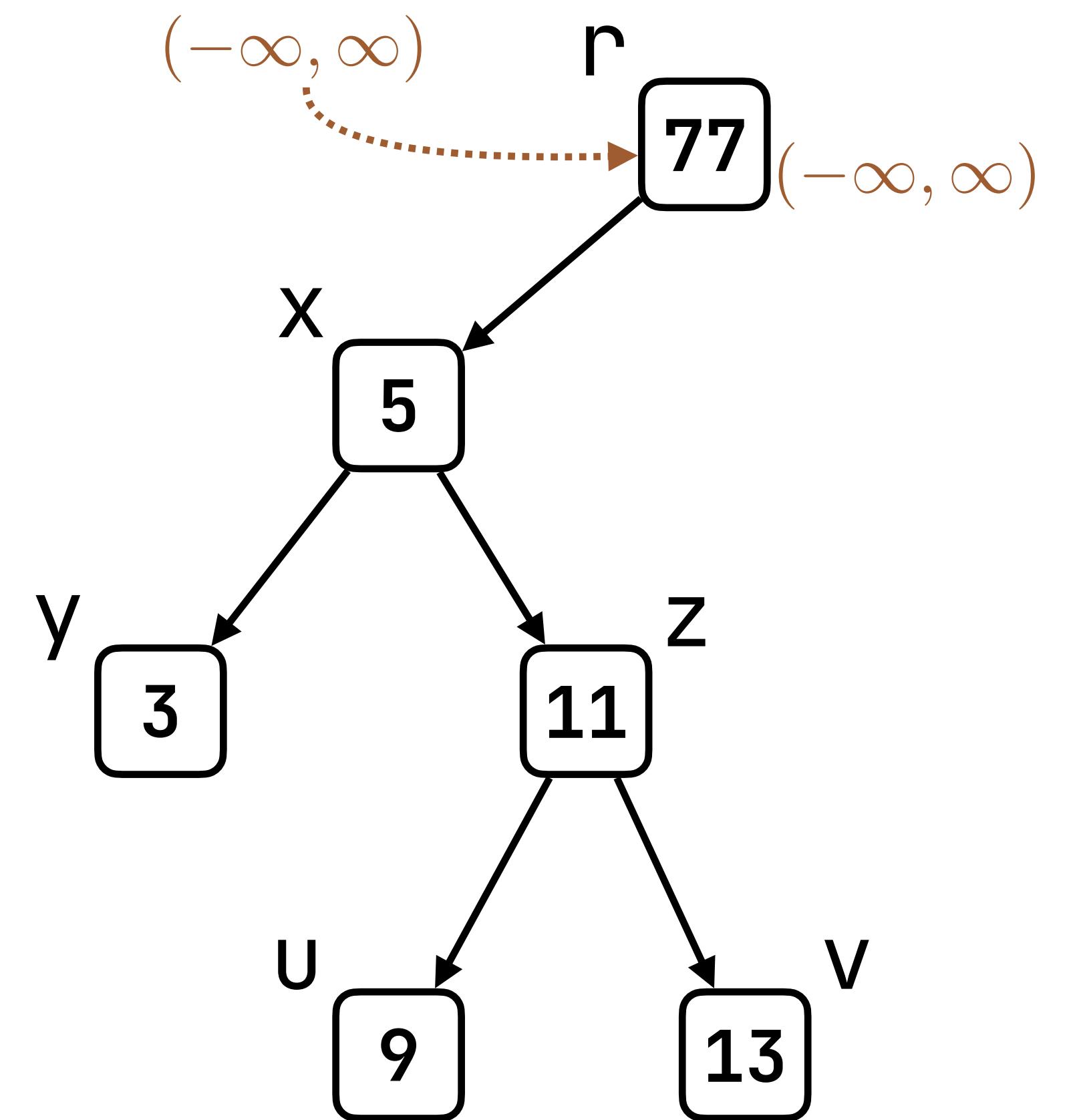
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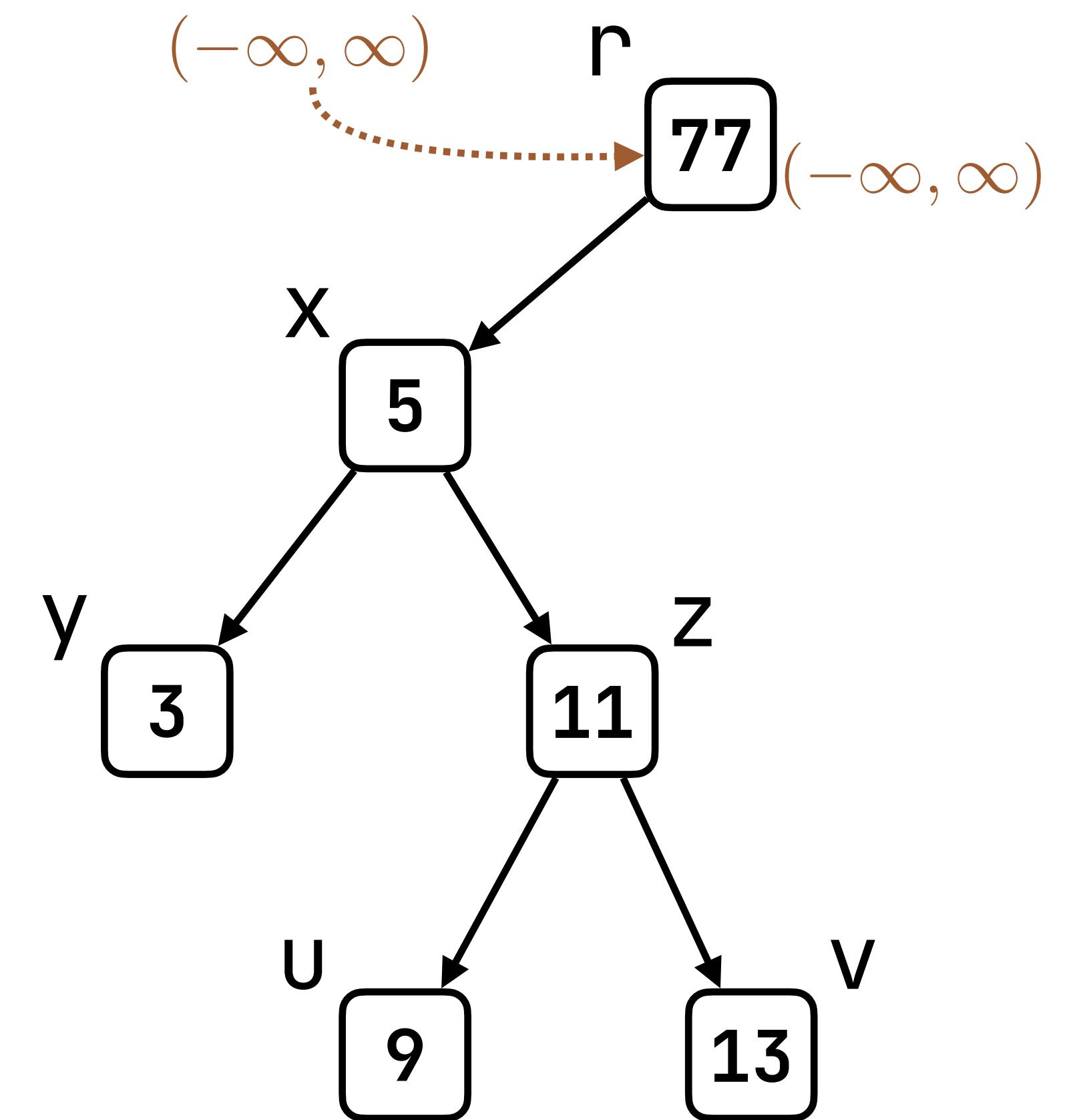
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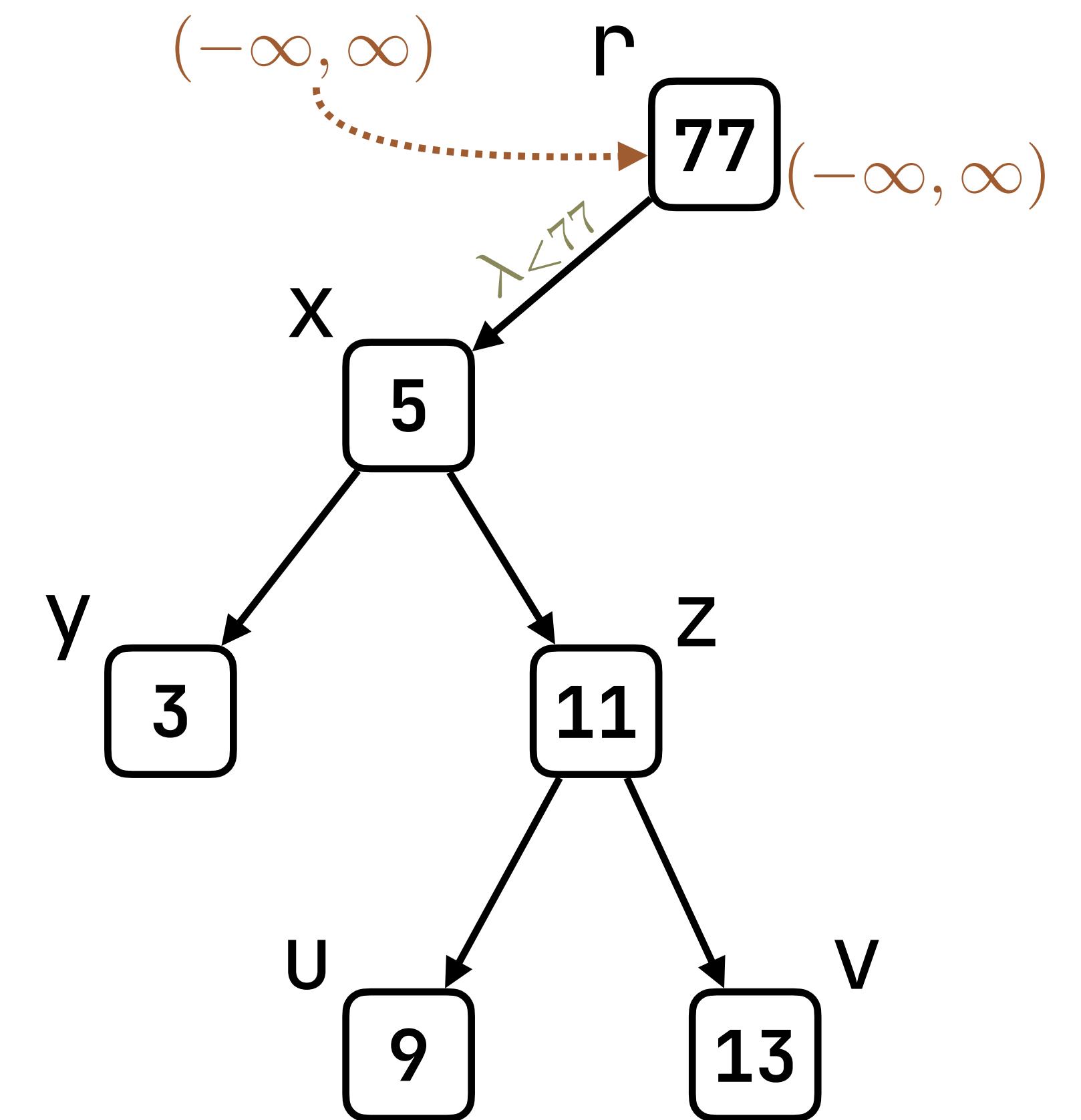
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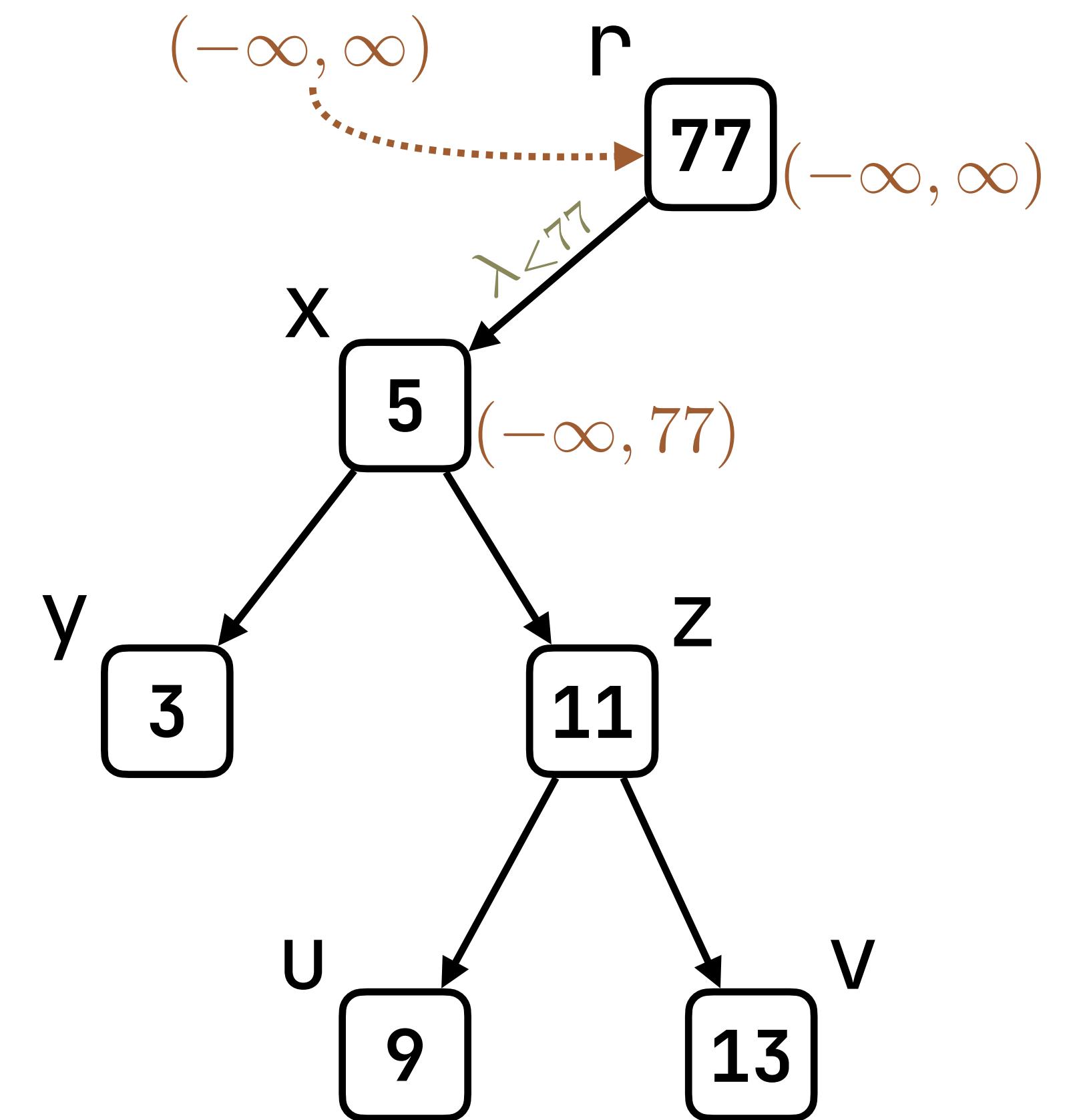
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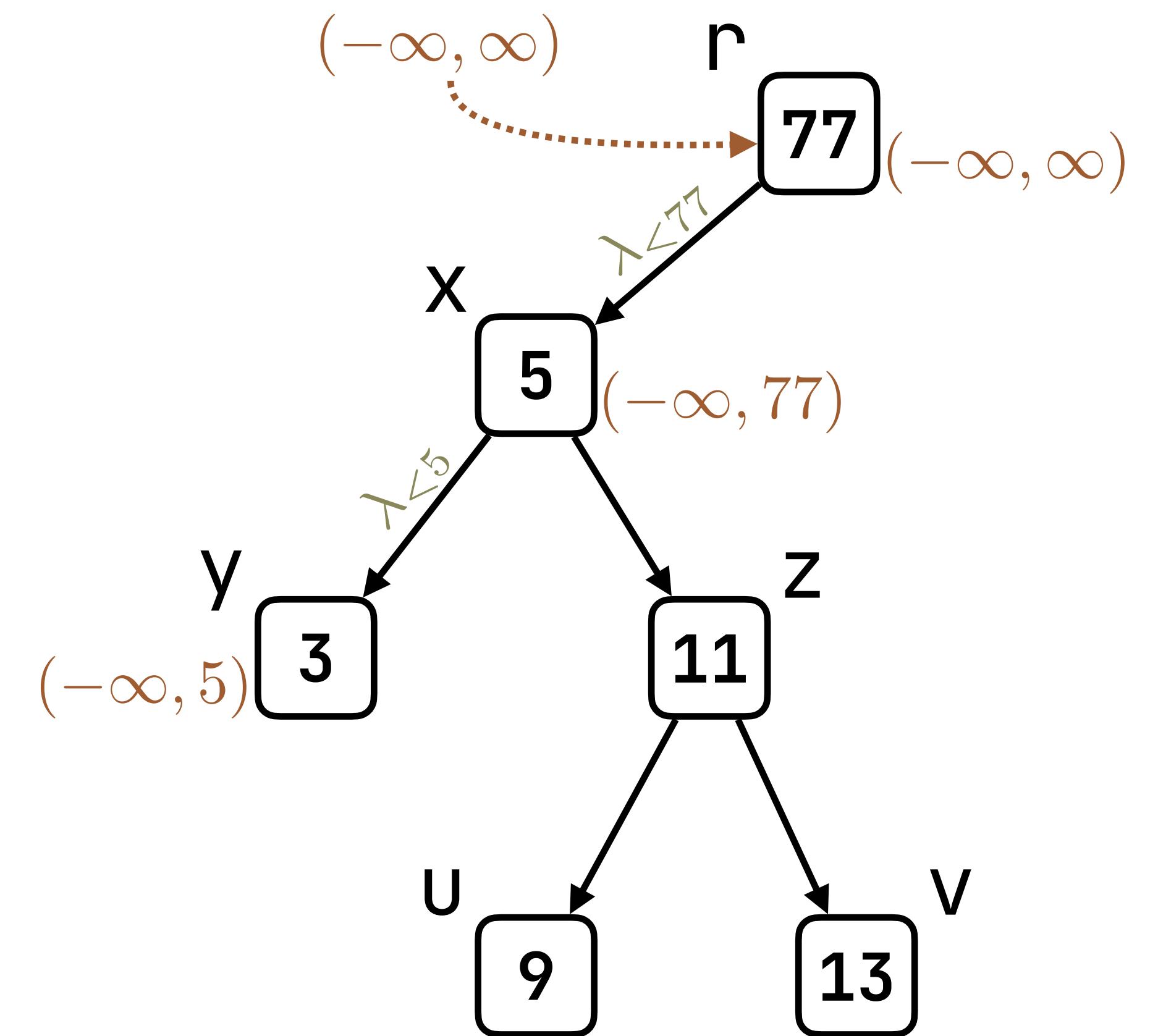
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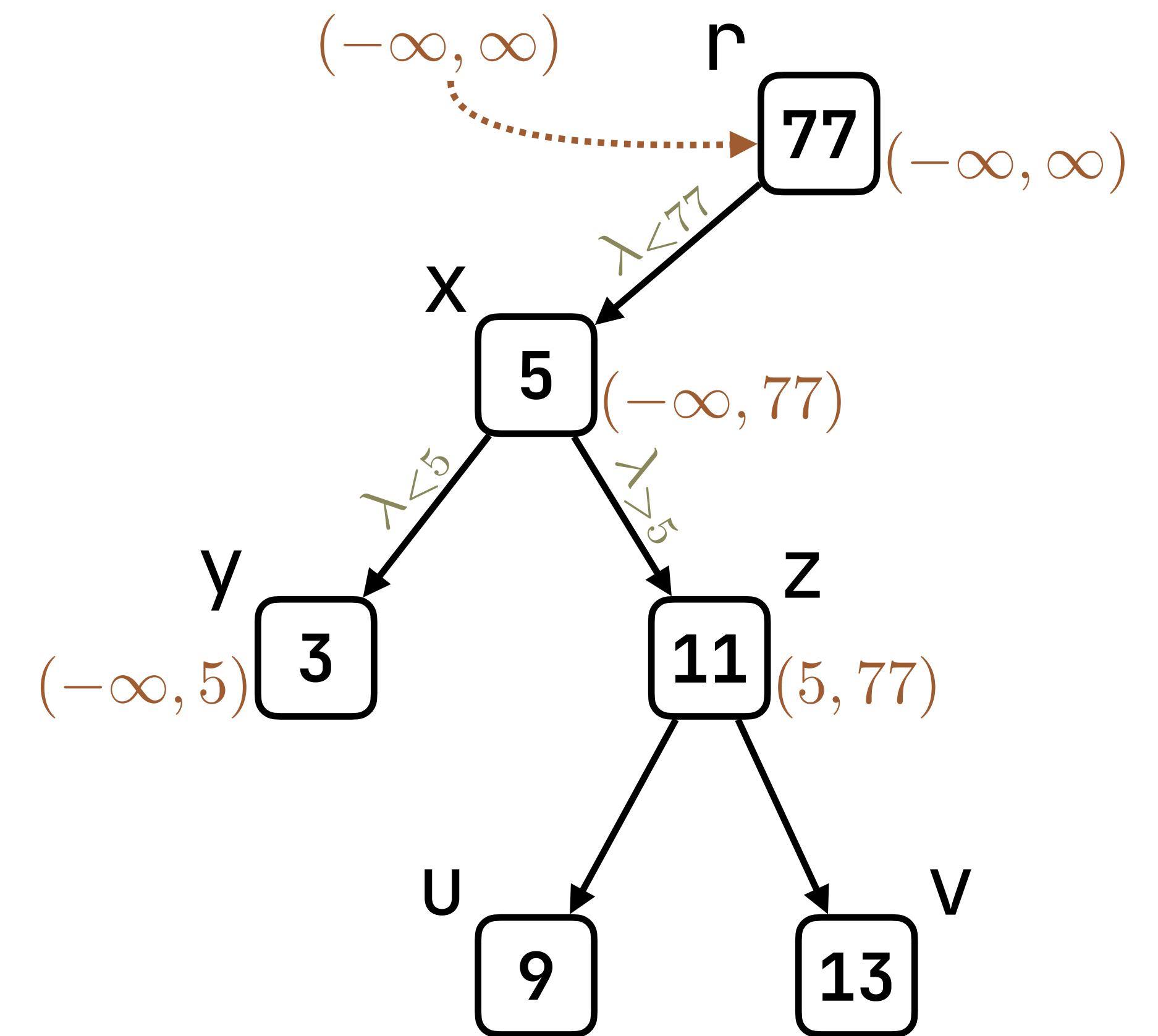
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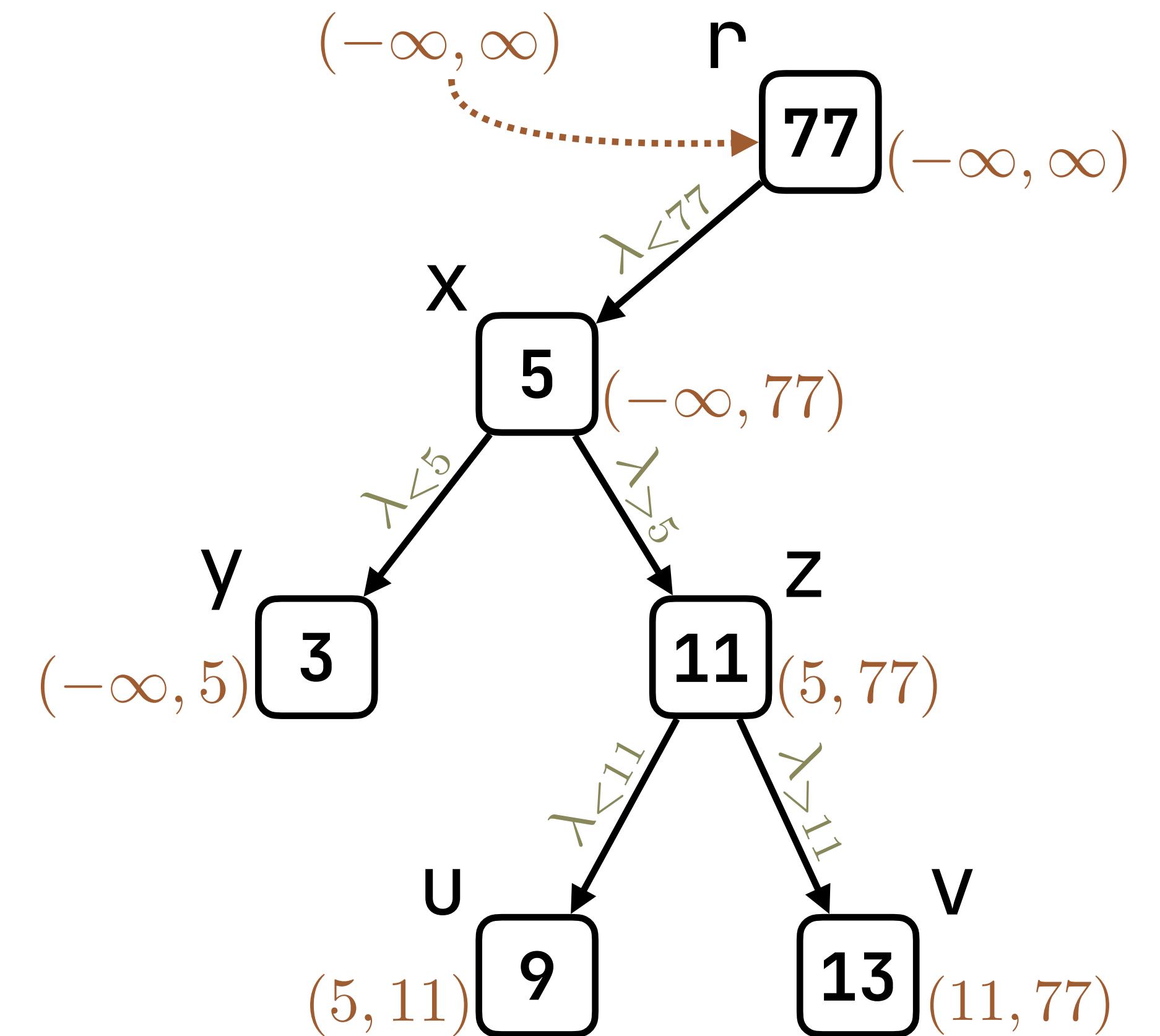
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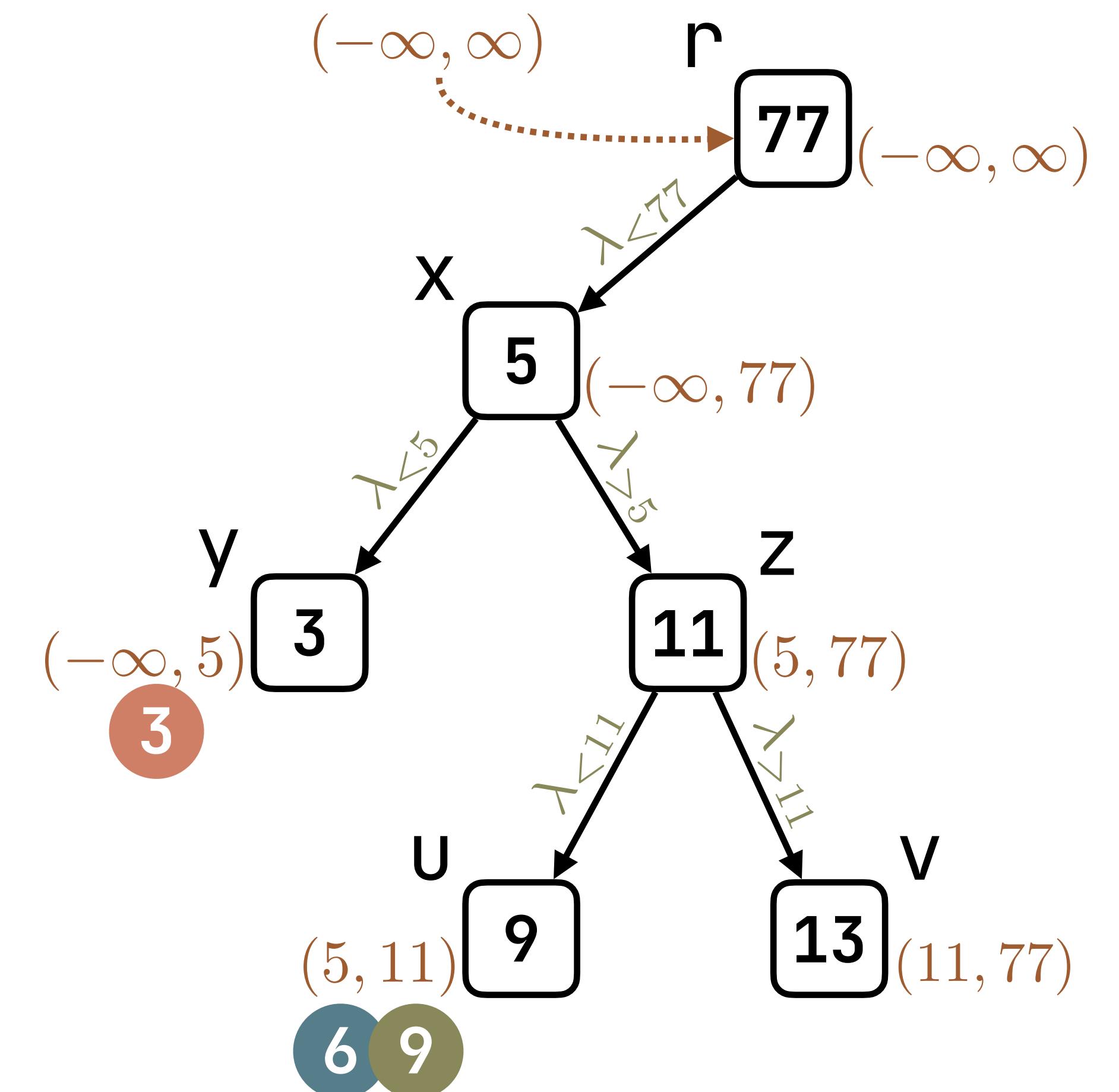
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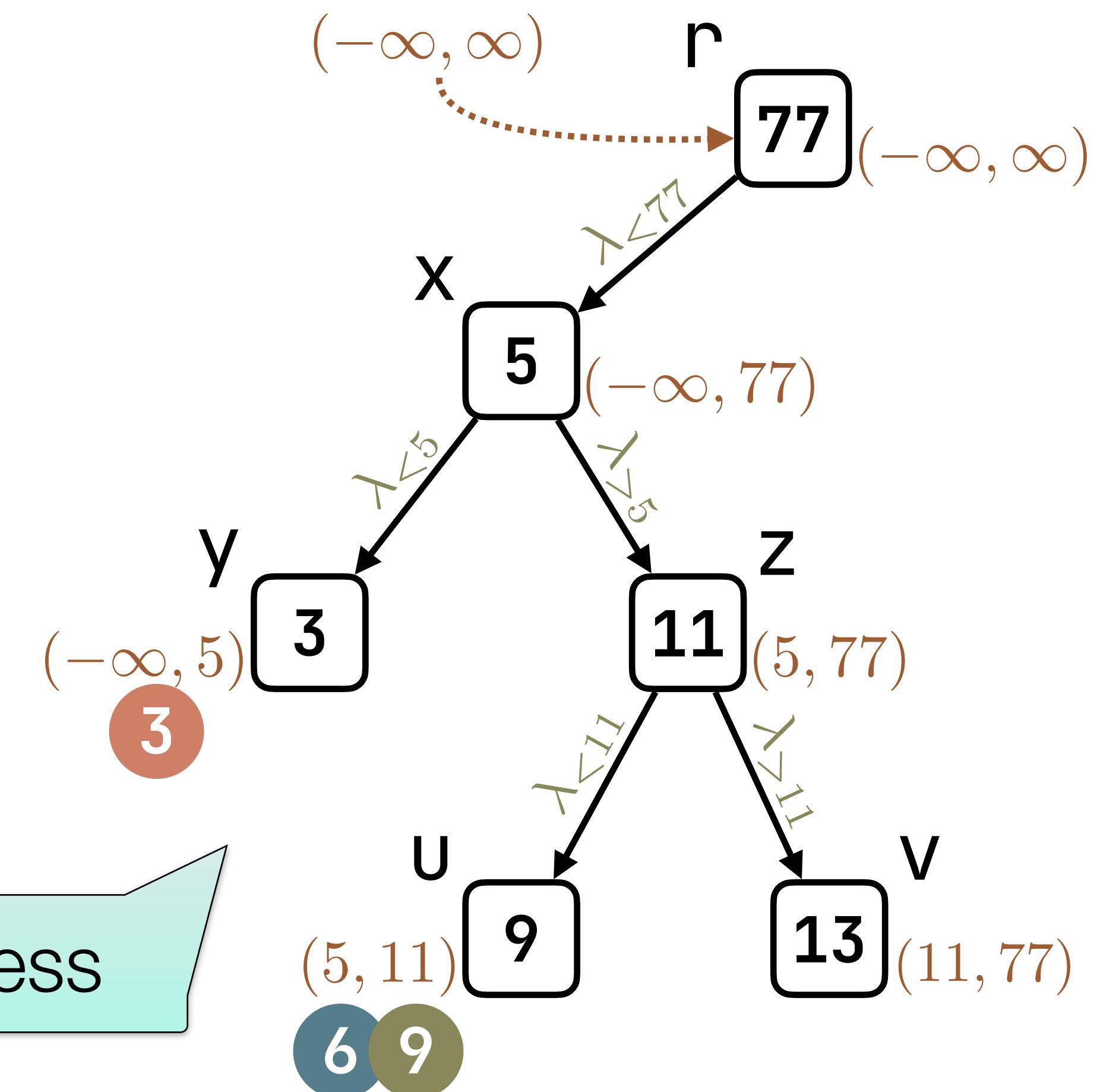


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Sufficient information for functional correctness





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Frame Inference

- Separation & flows
- Frame-preserving updates
- Finding footprints algorithmically



Comparing Footprints

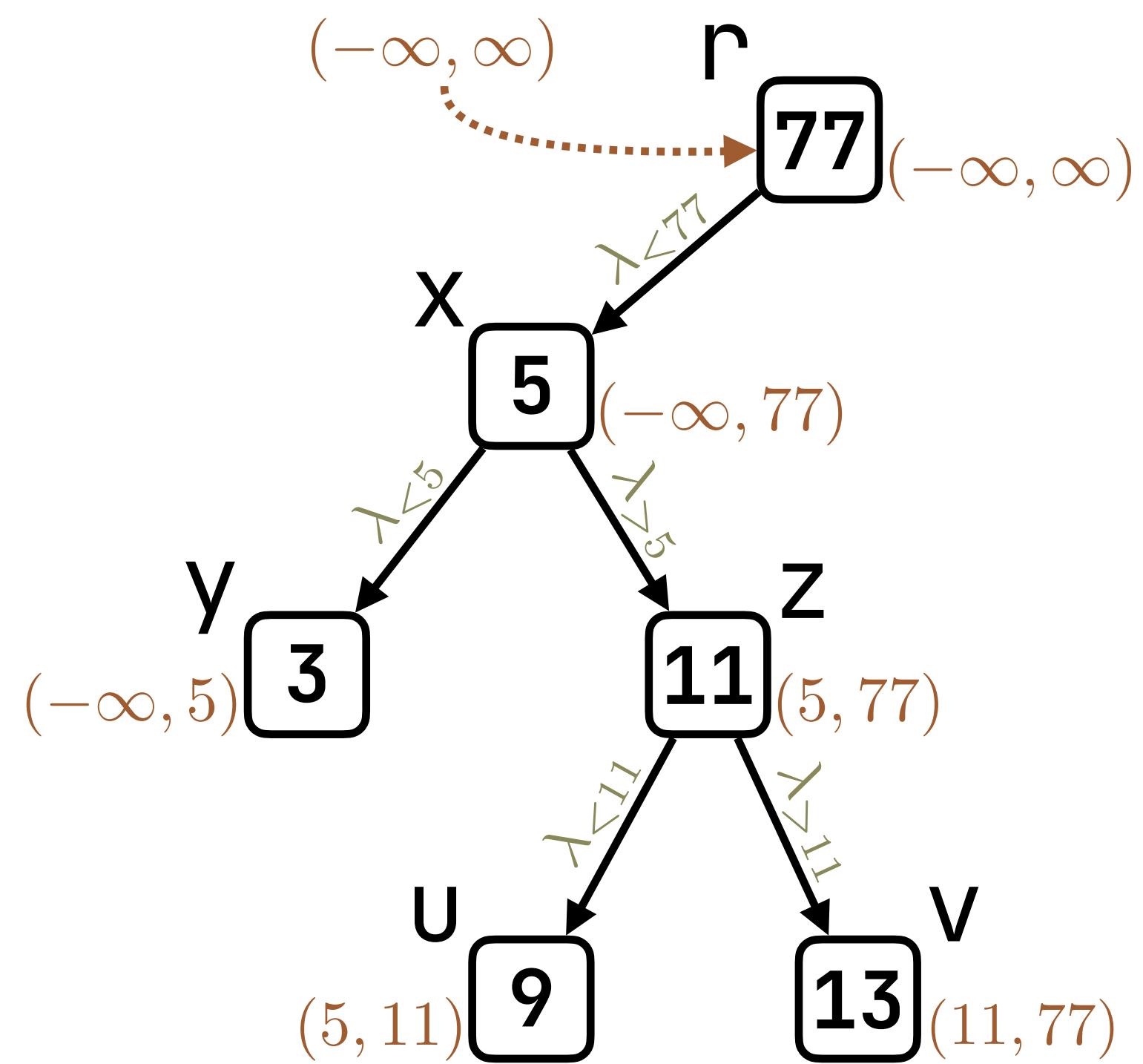
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Separation

- Flows graphs form a separation algebra
 - framing = "cutting the graph & flow"

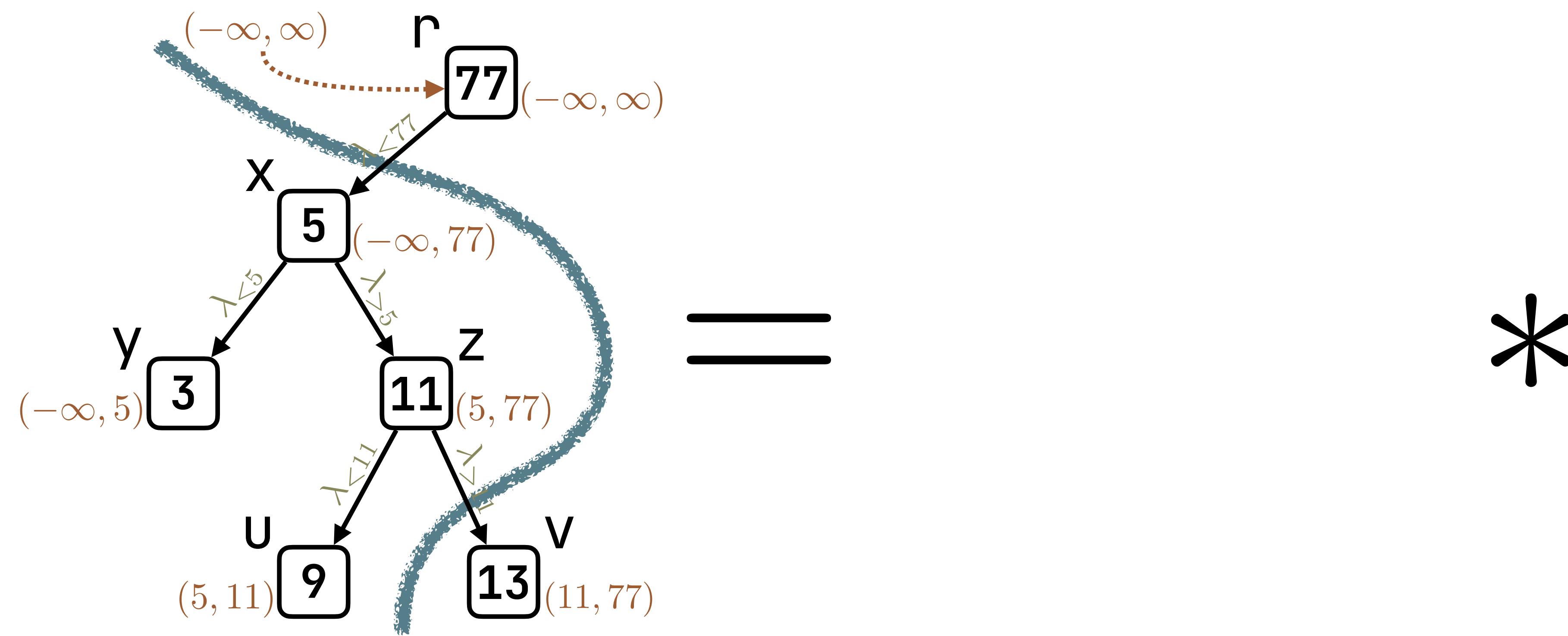
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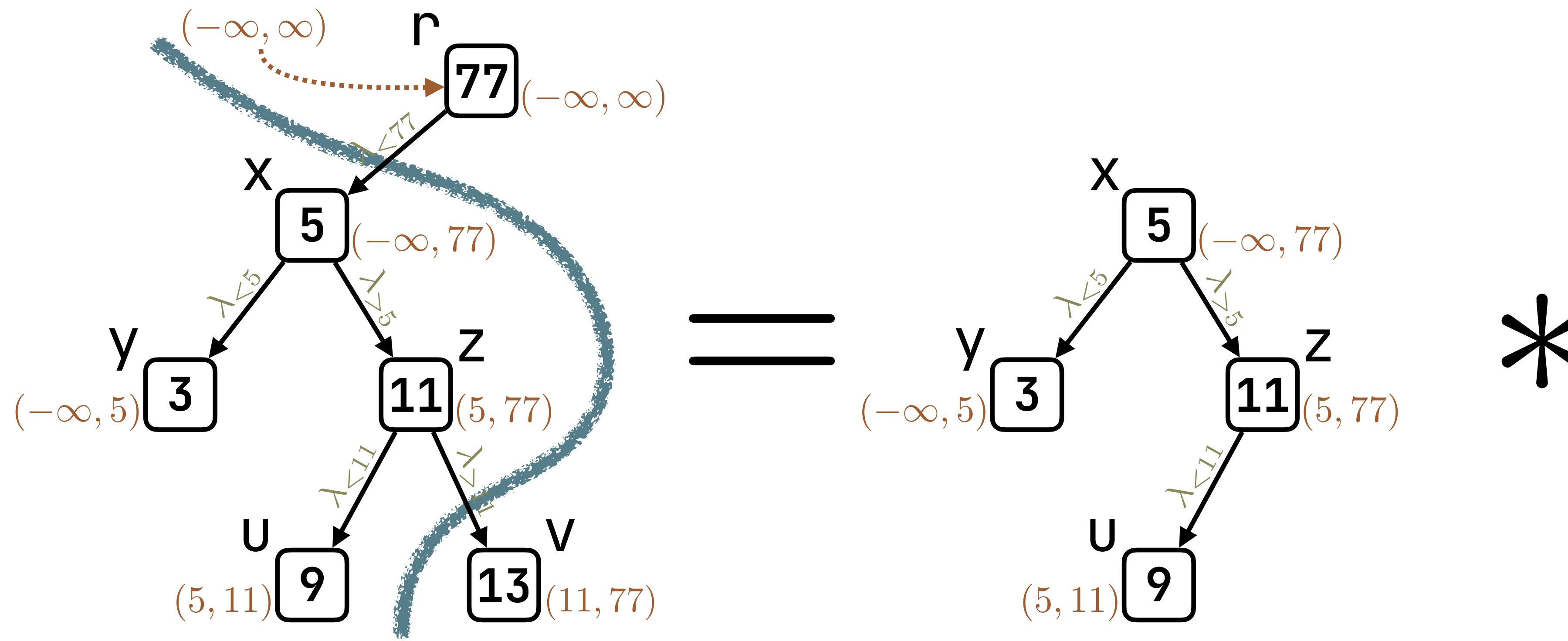
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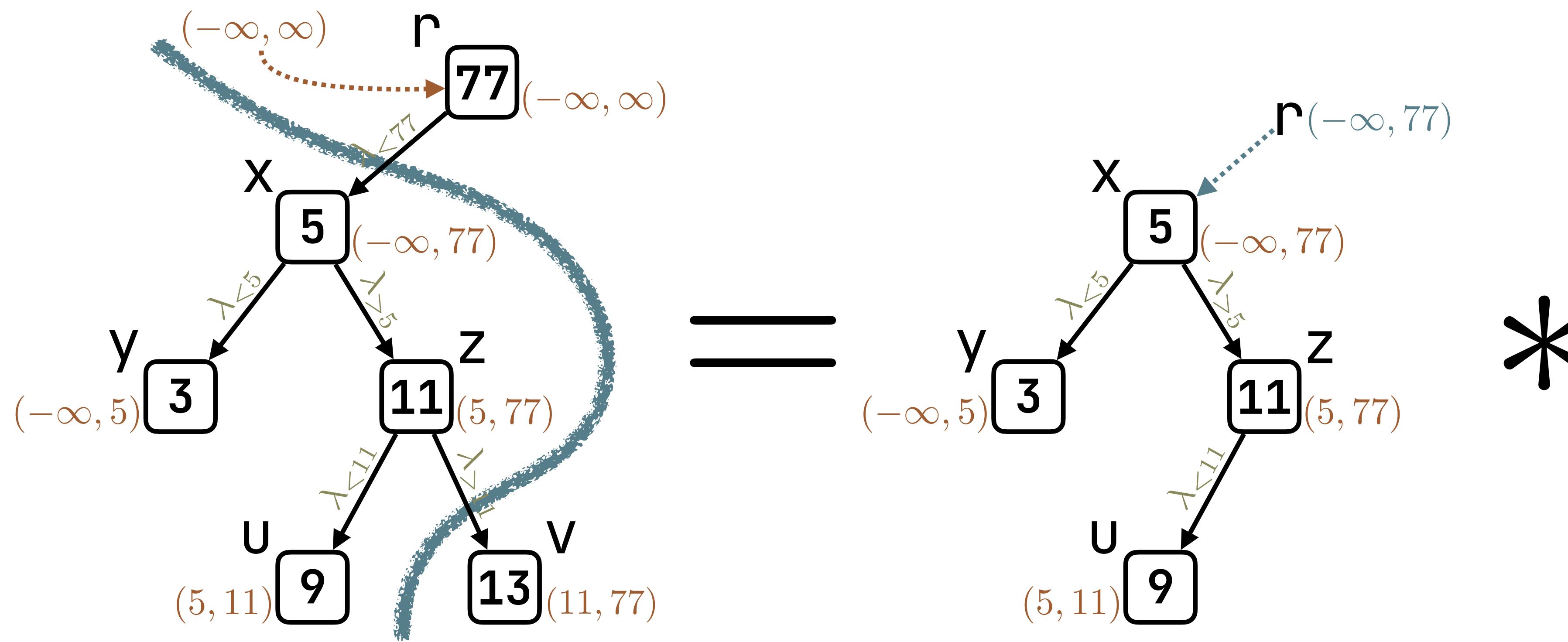
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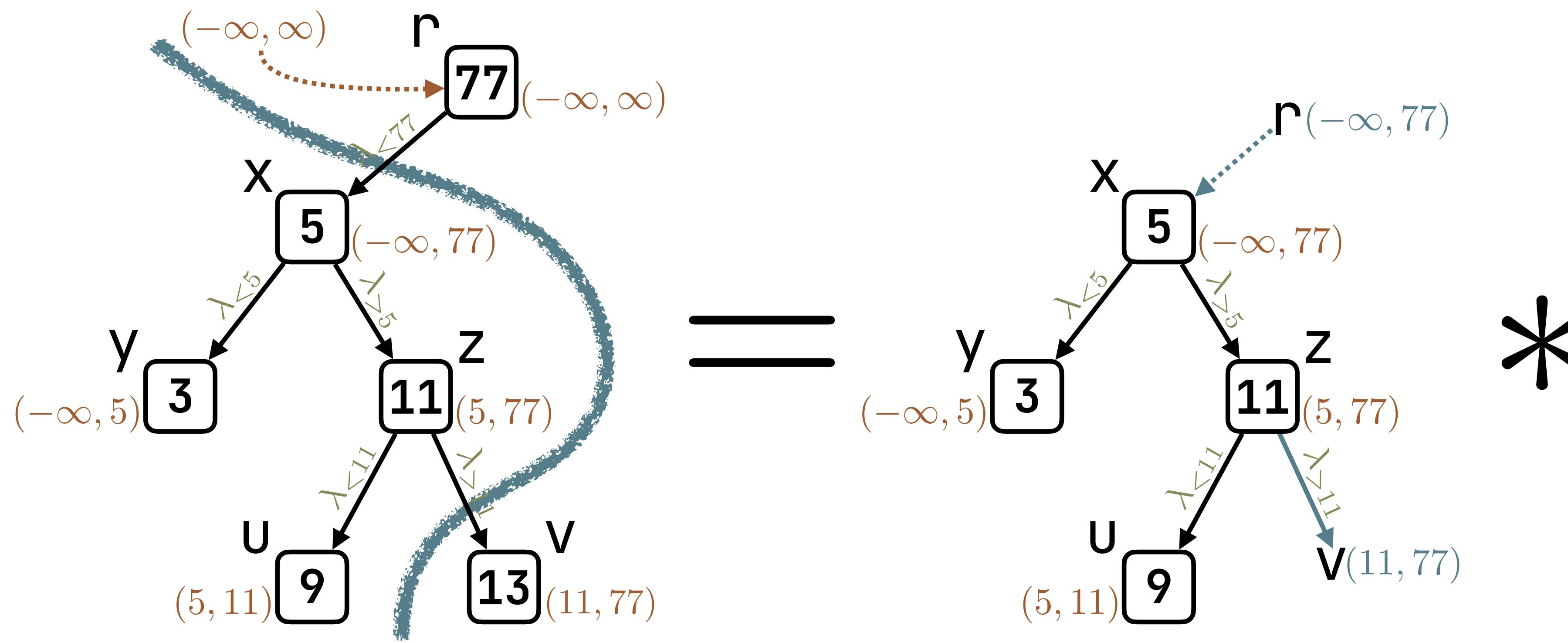
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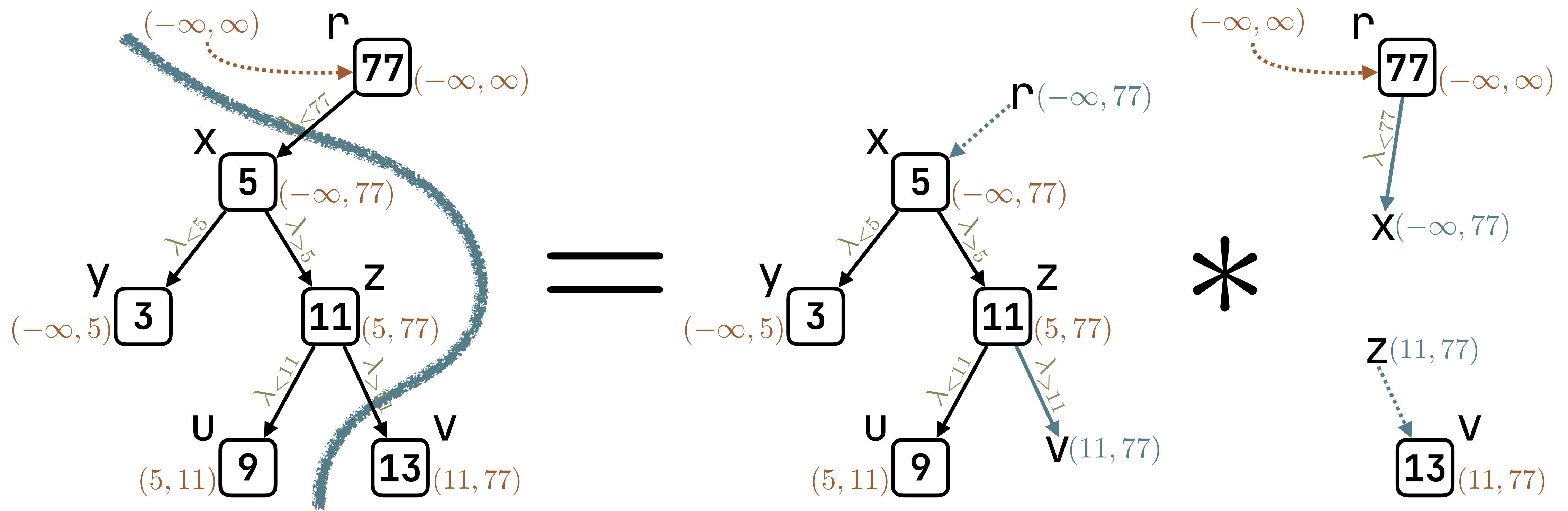
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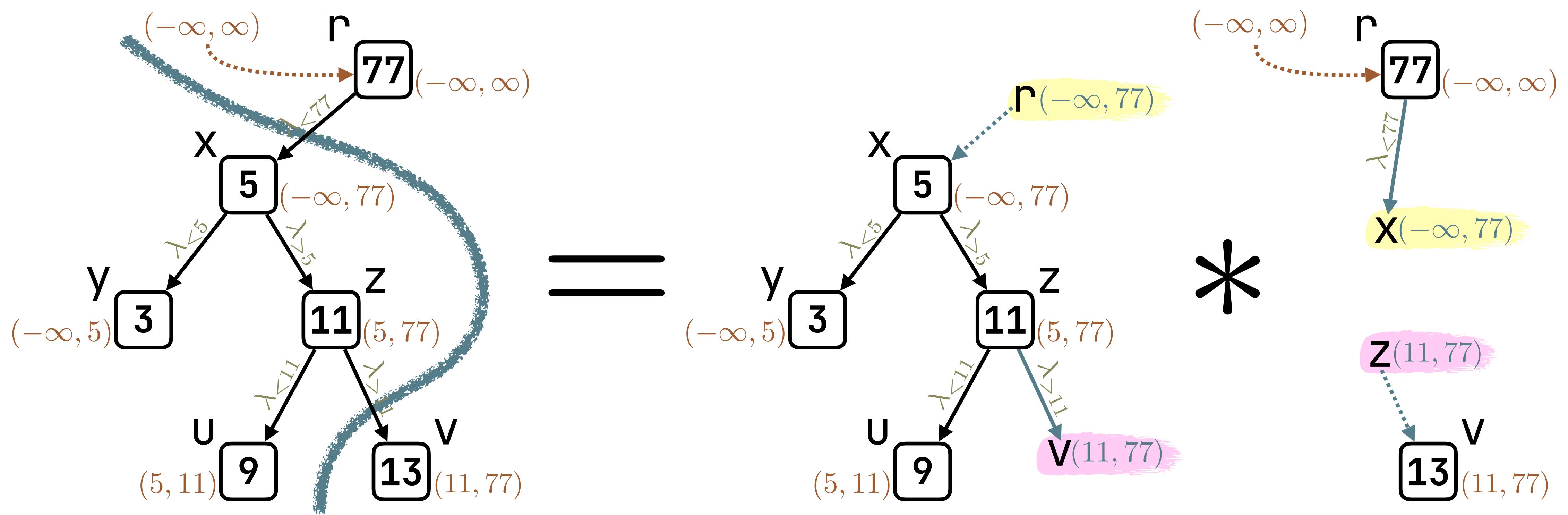
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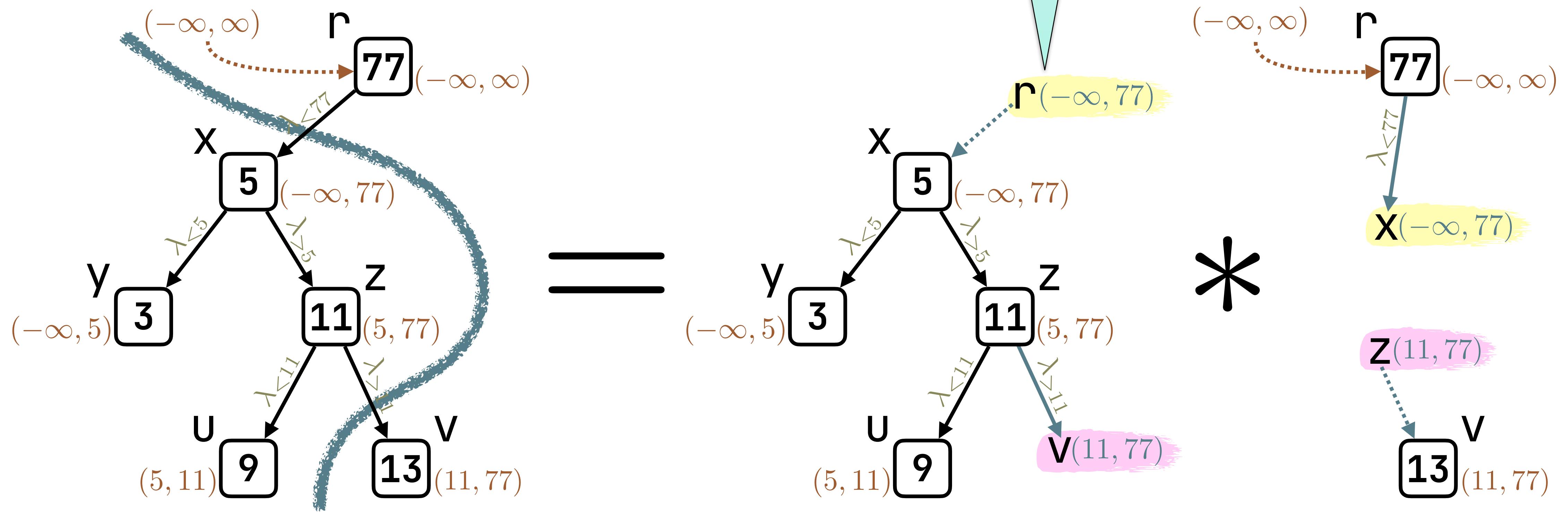
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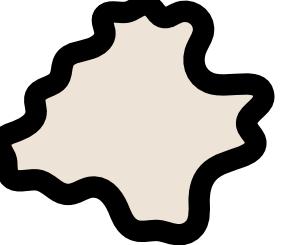
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Composition $*$ defined only if inflow & outflow match.

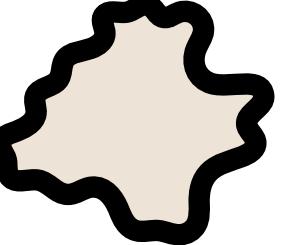


Frame-preserving Updates

- Footprint
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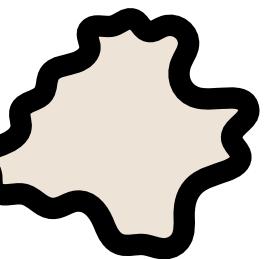
- Footprint

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- must be frame-preserving (not affect the frame):

$$\left\{ \begin{array}{c} P \\ \text{---} \\ \text{---} \end{array} \right\} \underset{\text{com}}{\sim} \left\{ \begin{array}{c} Q \\ \text{---} \\ \text{---} \end{array} \right\} \Rightarrow \left\{ \begin{array}{c} P \\ \text{---} \\ \text{---} \end{array} * F \right\} \underset{\text{com}}{\sim} \left\{ \begin{array}{c} Q \\ \text{---} \\ \text{---} \end{array} * F \right\}$$

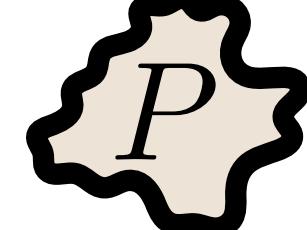
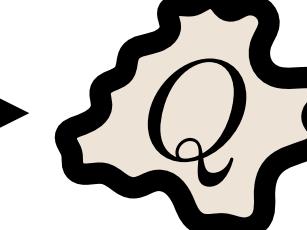
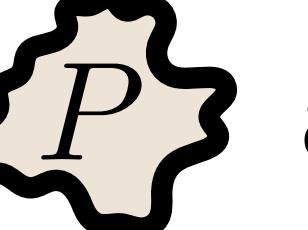
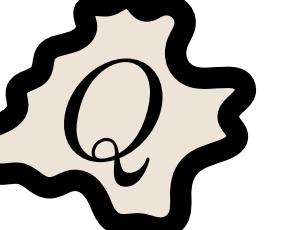
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- Theorem:

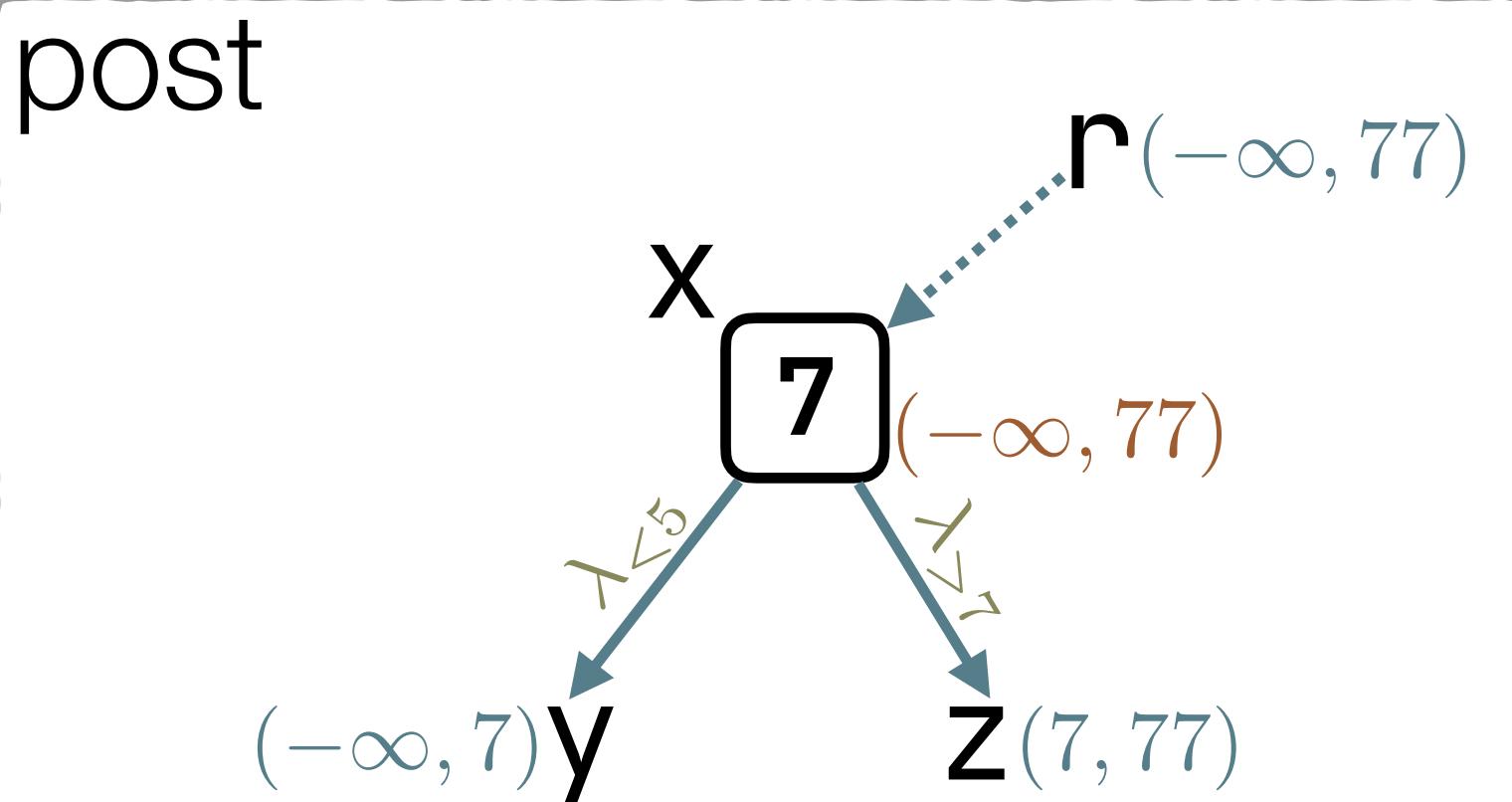
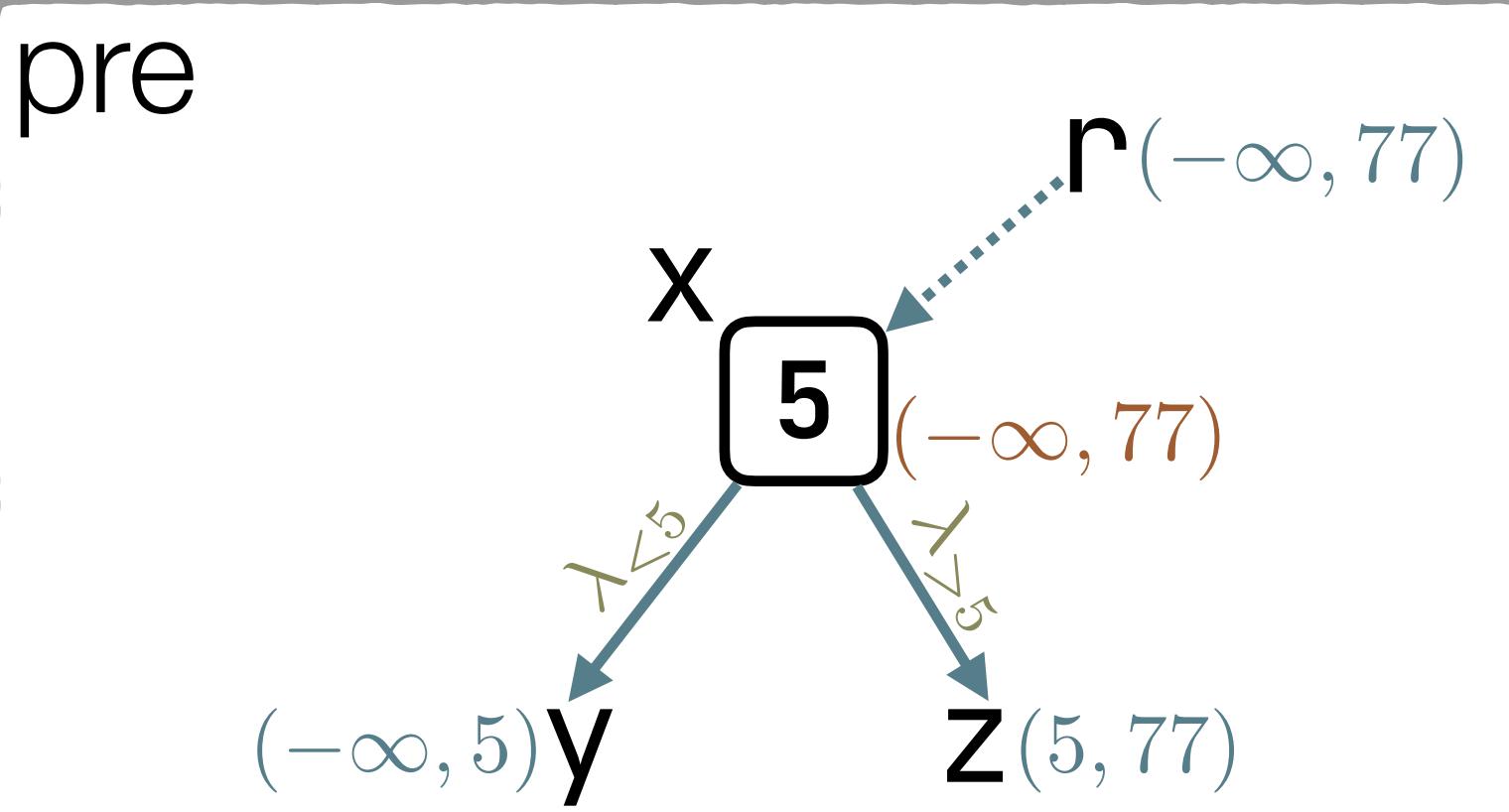
Update  \rightarrow  is frame-preserving if  and  have the same outflow, *for all inflows*.

Finding Footprints

- Algorithm:
 1. add physical footprint
 2. compute outflow (for all inflows)
 3. add nodes if pre/post outflow differs
 4. repeat until fixed point

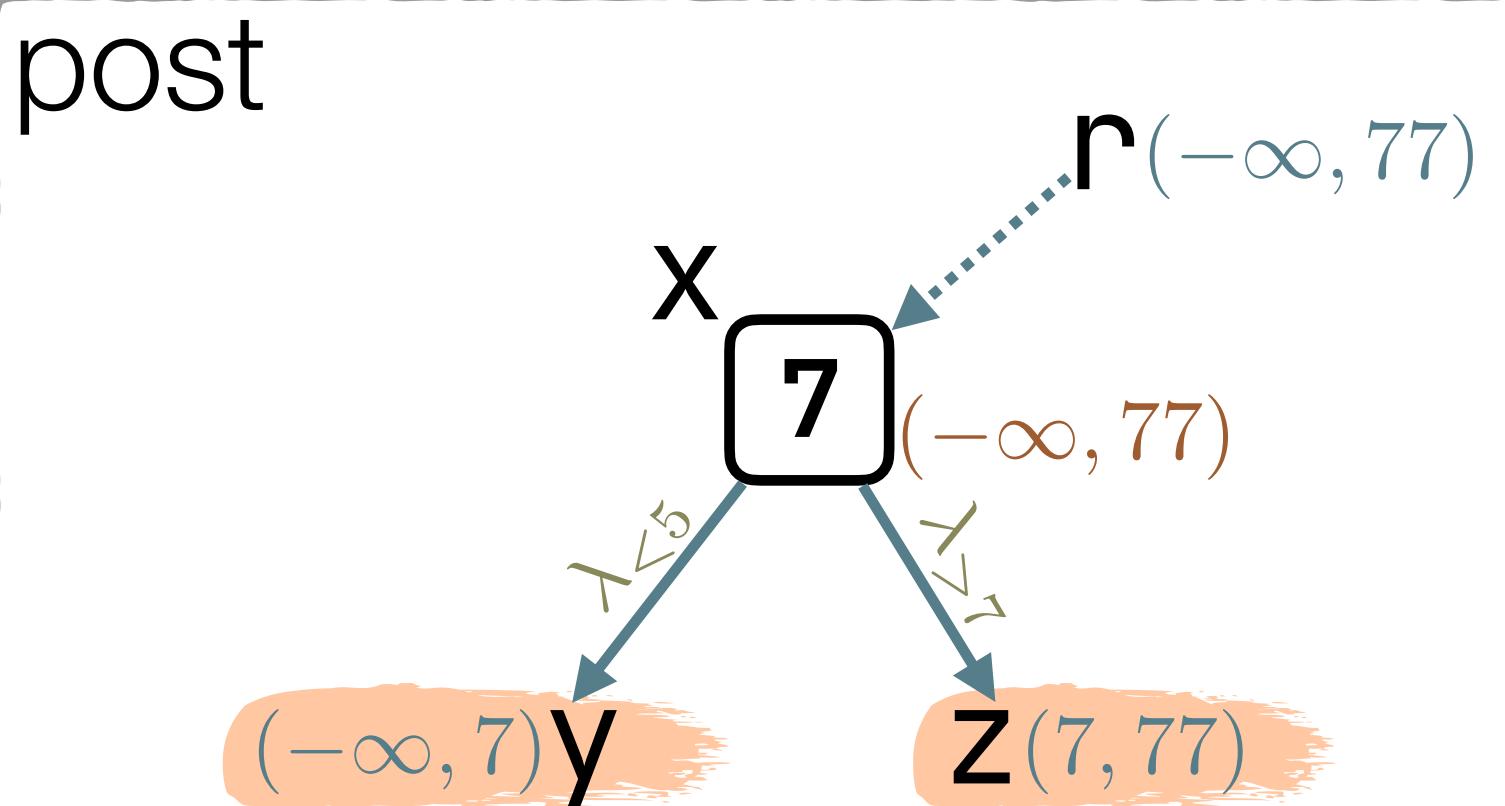
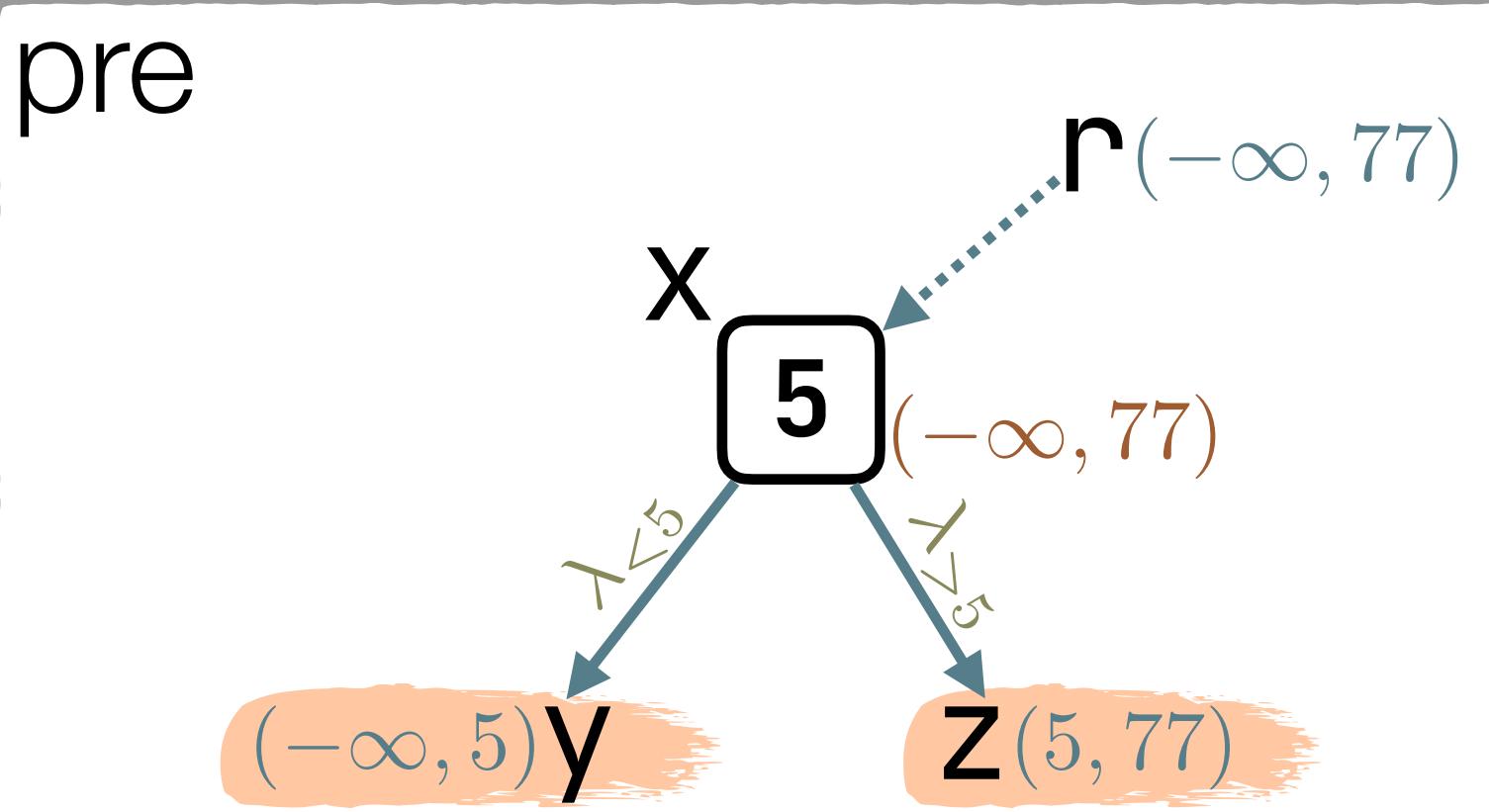
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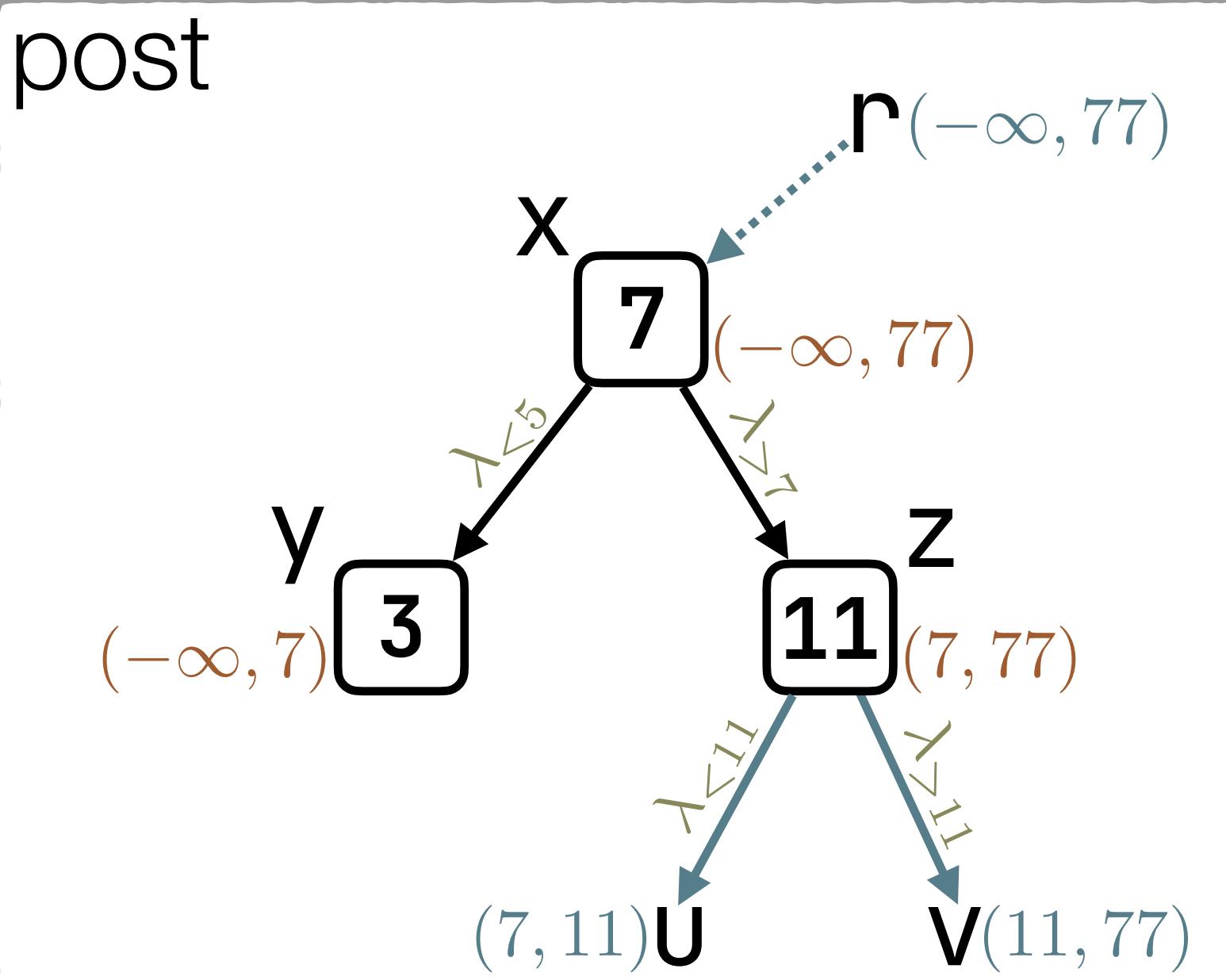
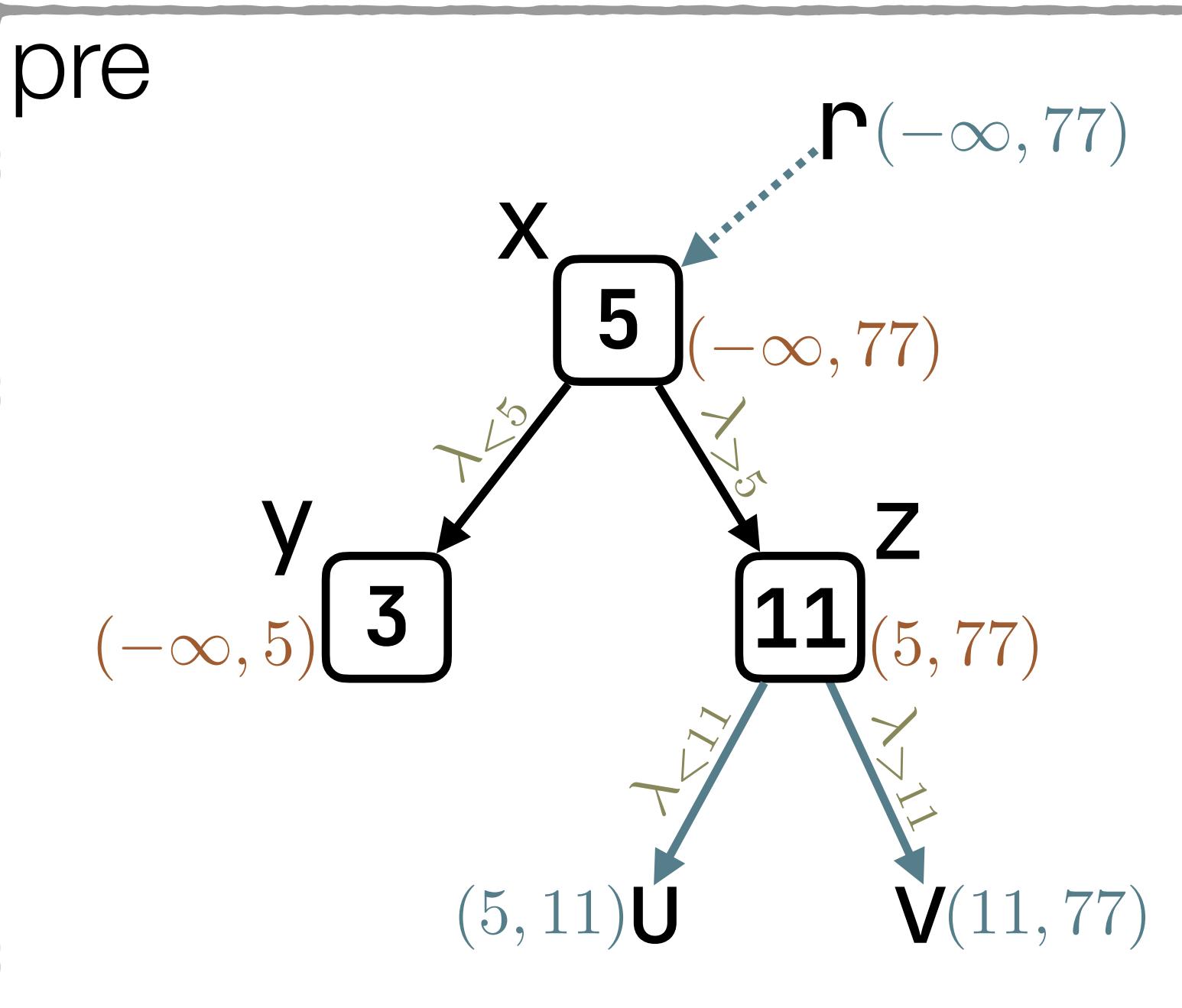
Finding Footprints

- Algorithm:
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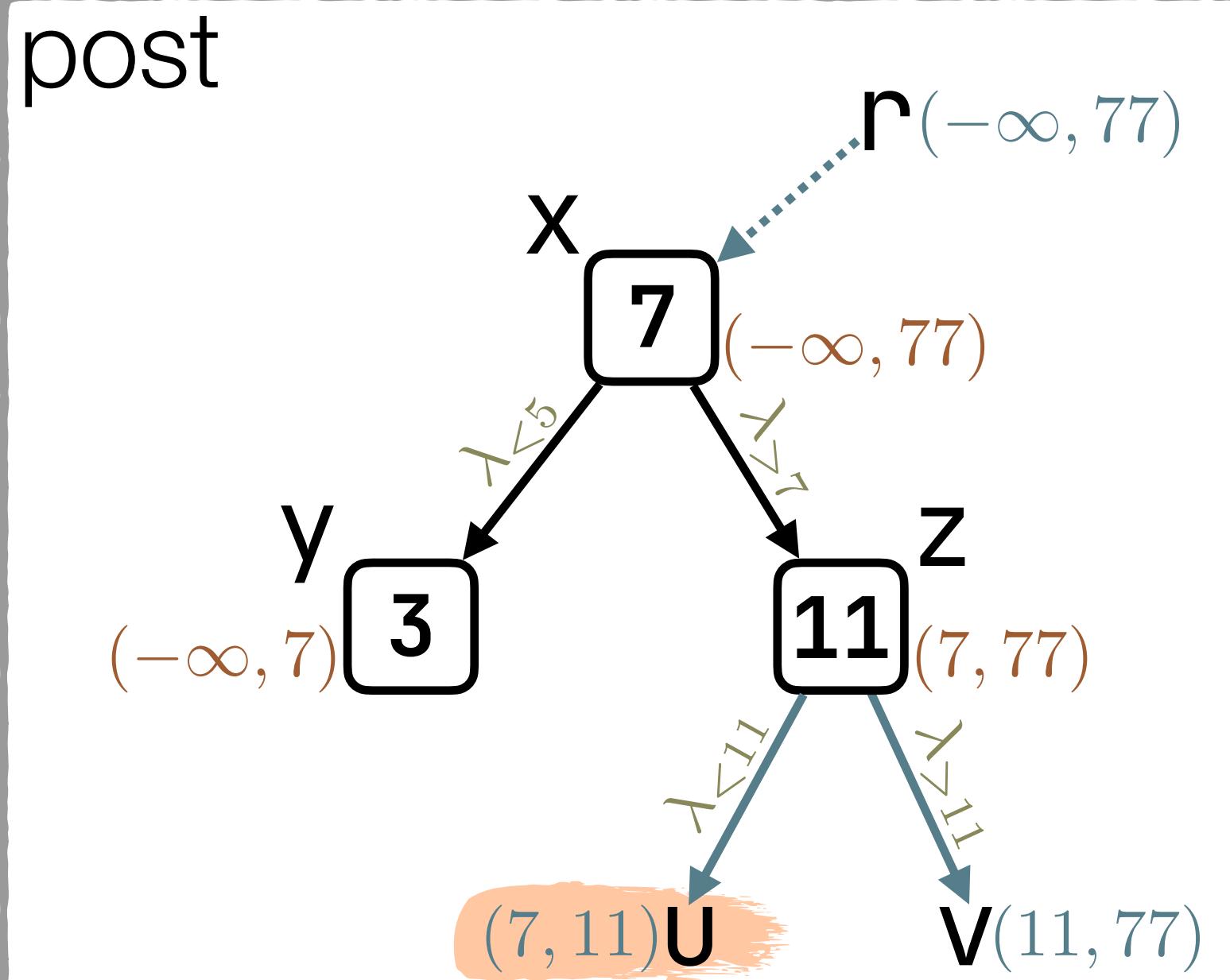
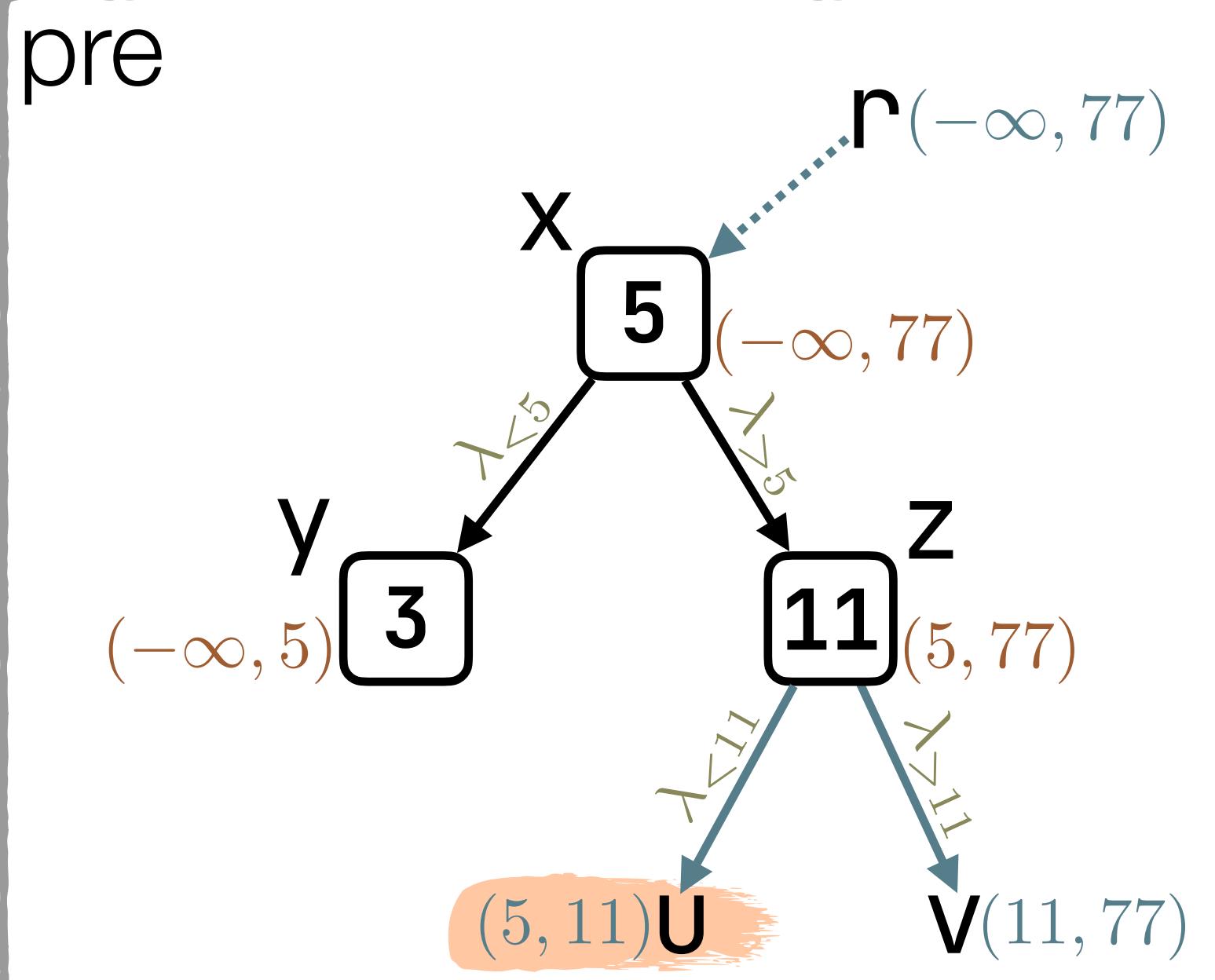
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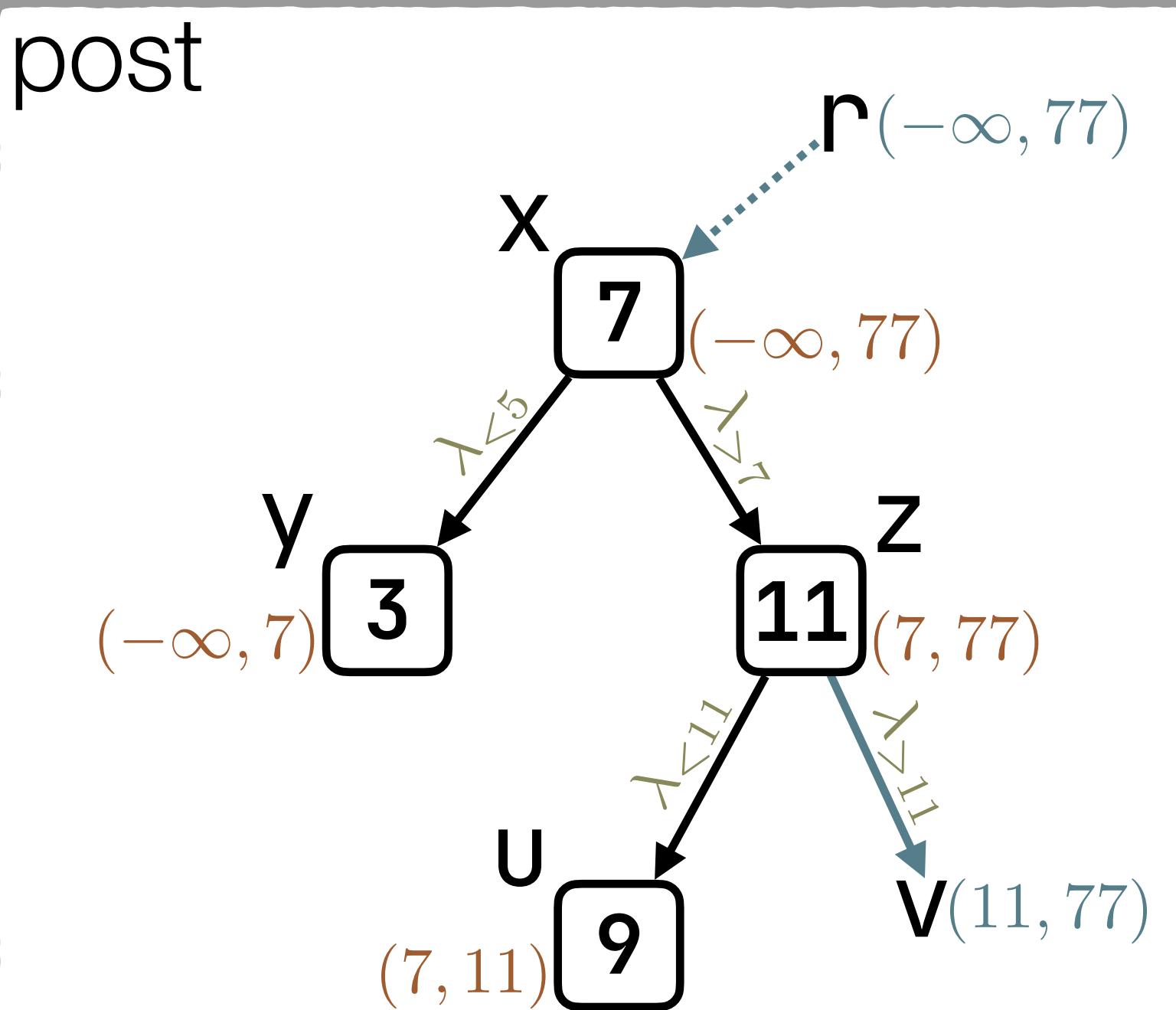
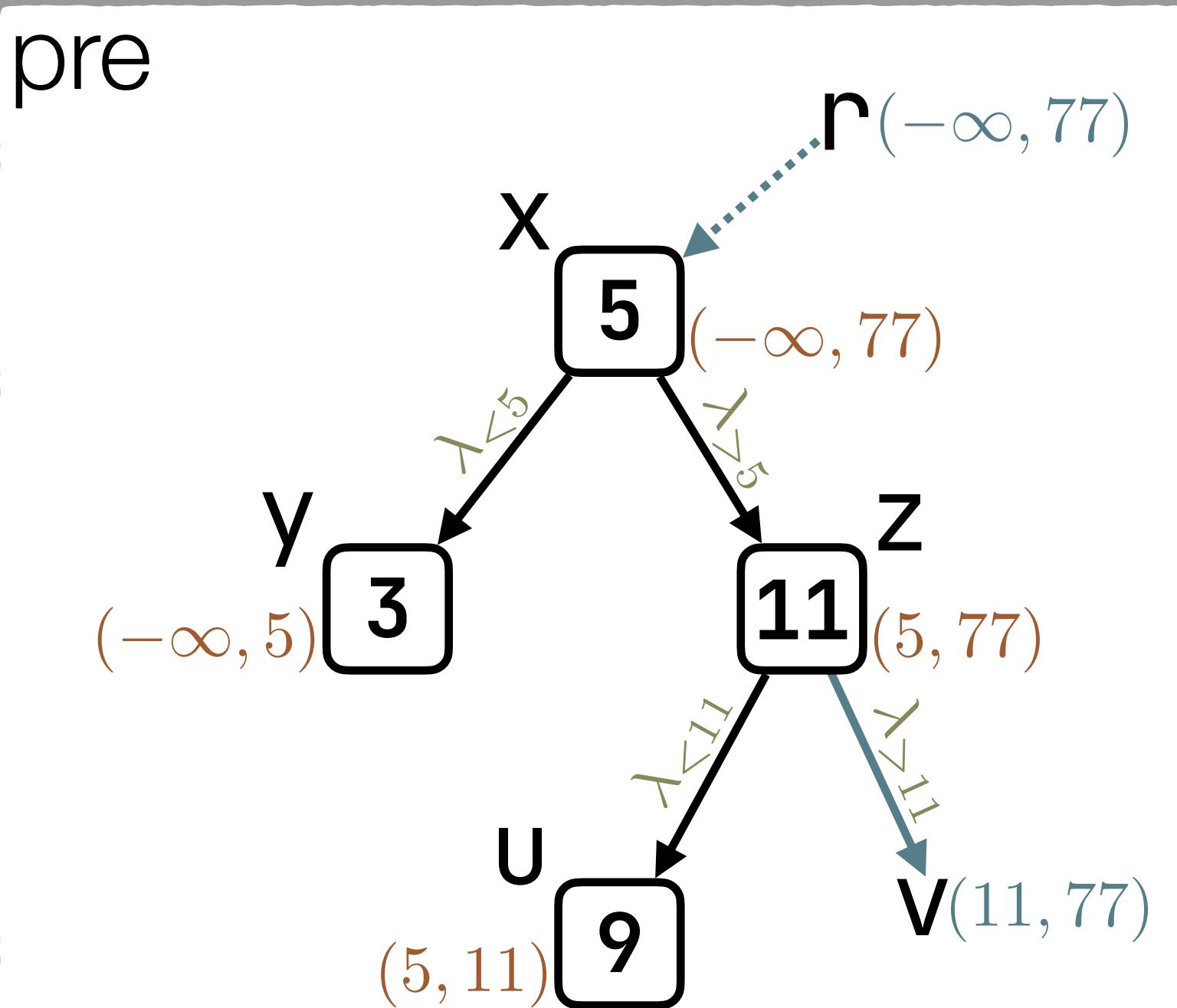
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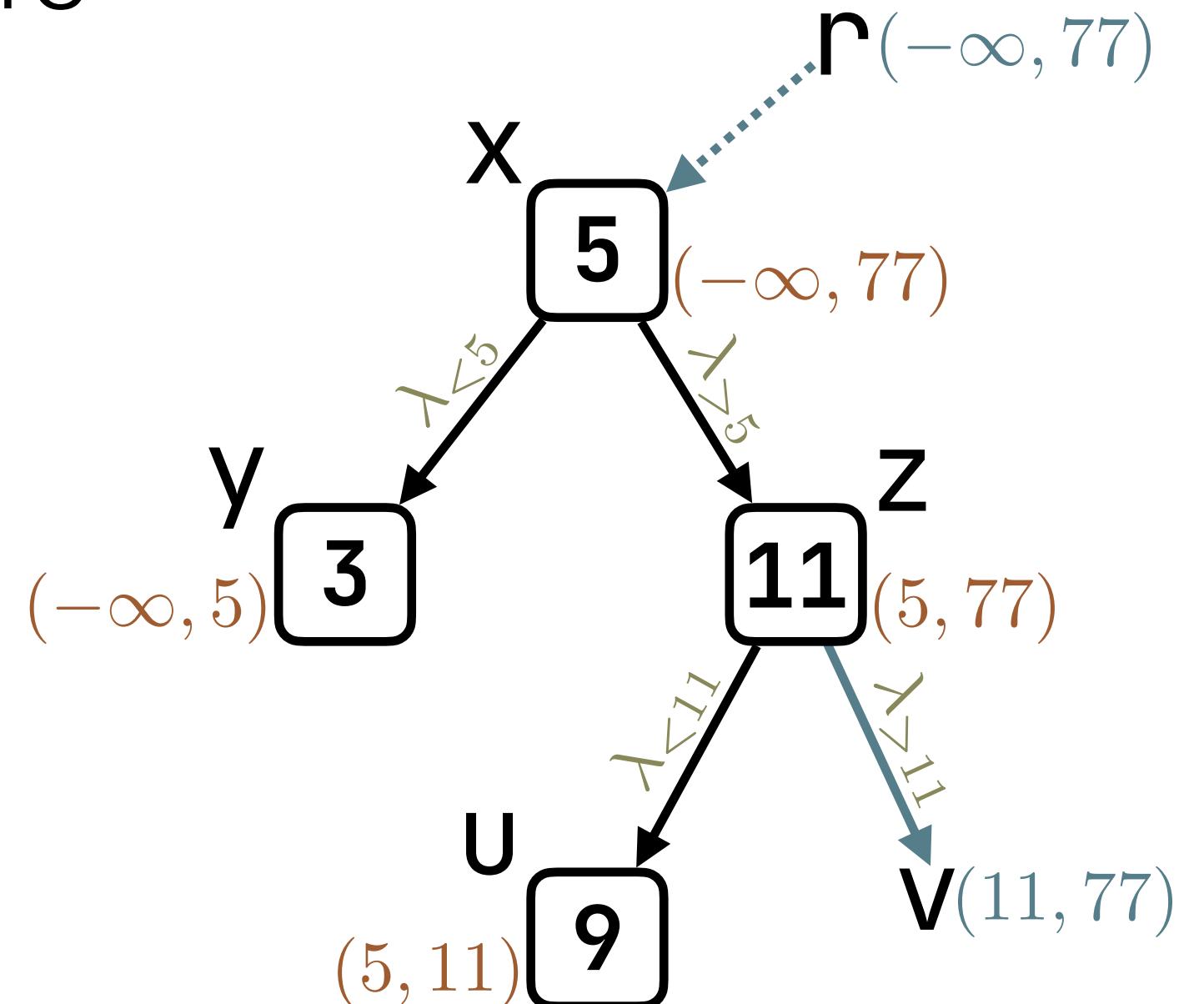
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Not minimal.

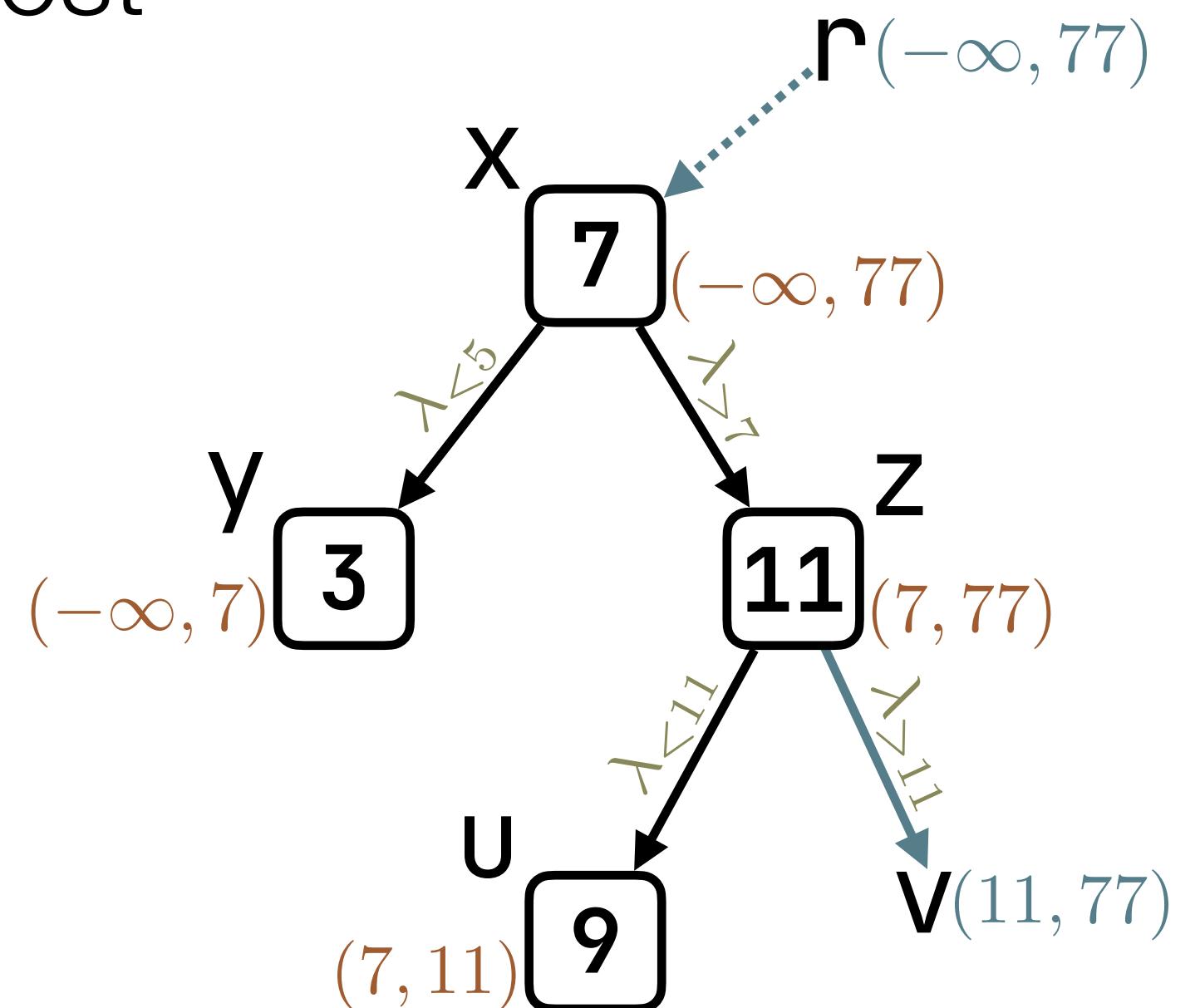
Incomplete.

Works well in practice.

pre



post





Flow Framework

- Ghost state for heap graphs
- Inspired by data-flow analysis
- Formalizes inductive heap invariants



Frame Inference

- Separation & flows
- Frame-preserving updates
- Finding footprints algorithmically

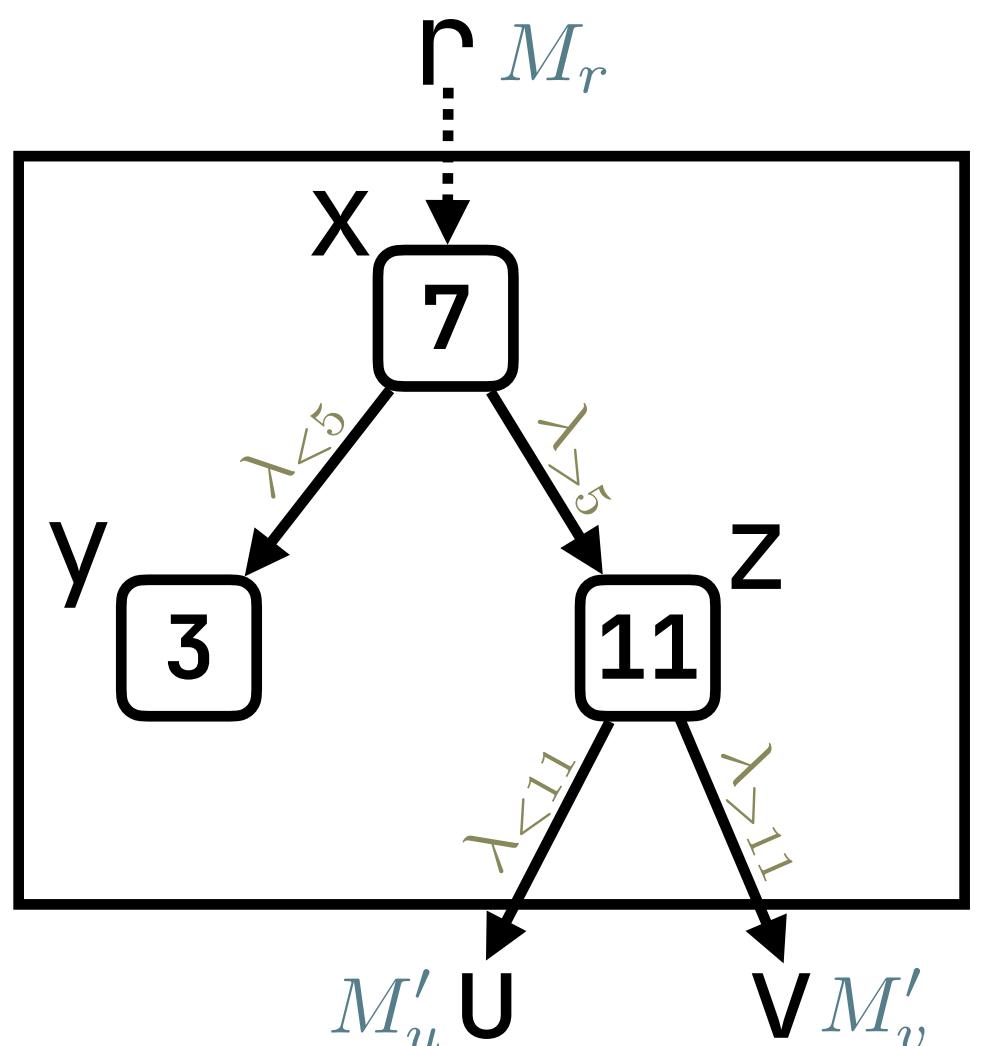
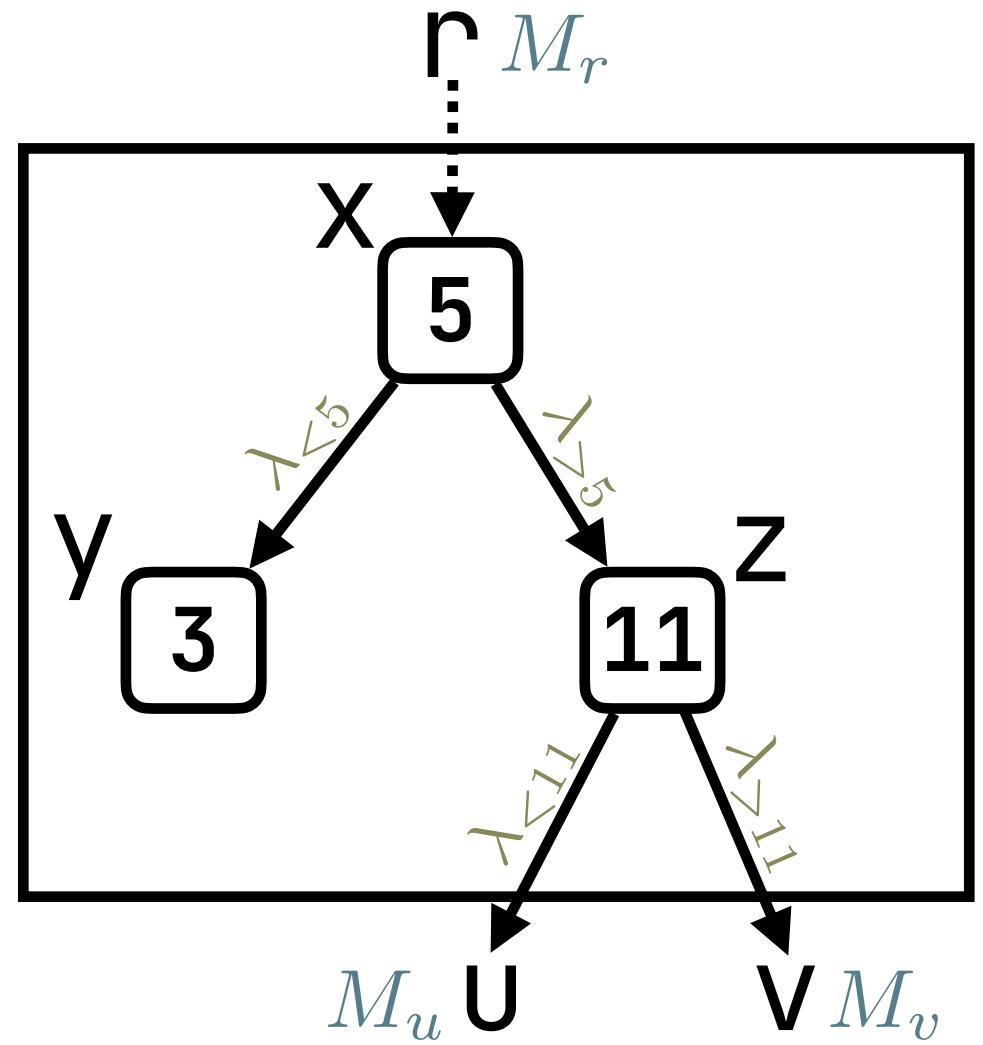


Comparing Footprints

- Check if update is frame-preserving
- Efficient checks for general graphs

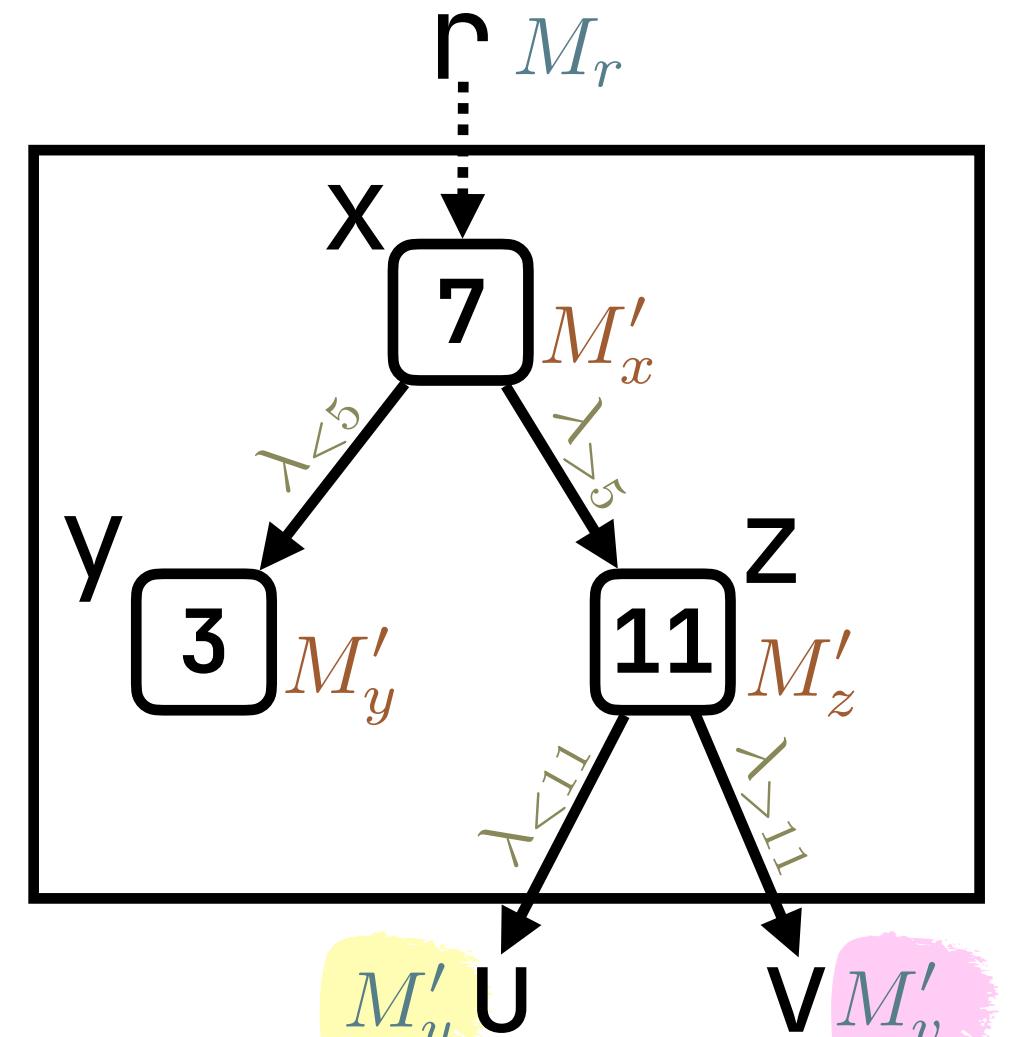
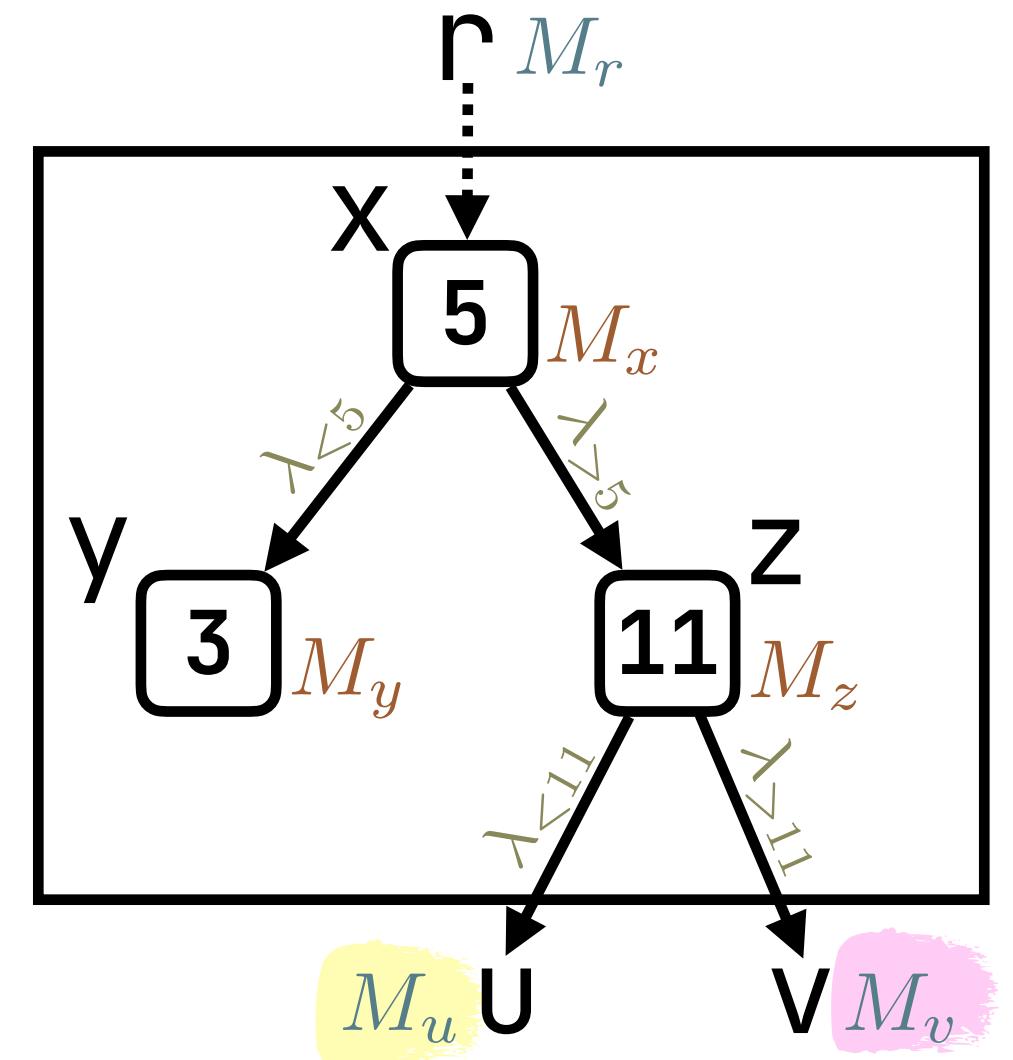
Overview

- Question: does $M_u = M'_u$ and $M_v = M'_v$ hold?



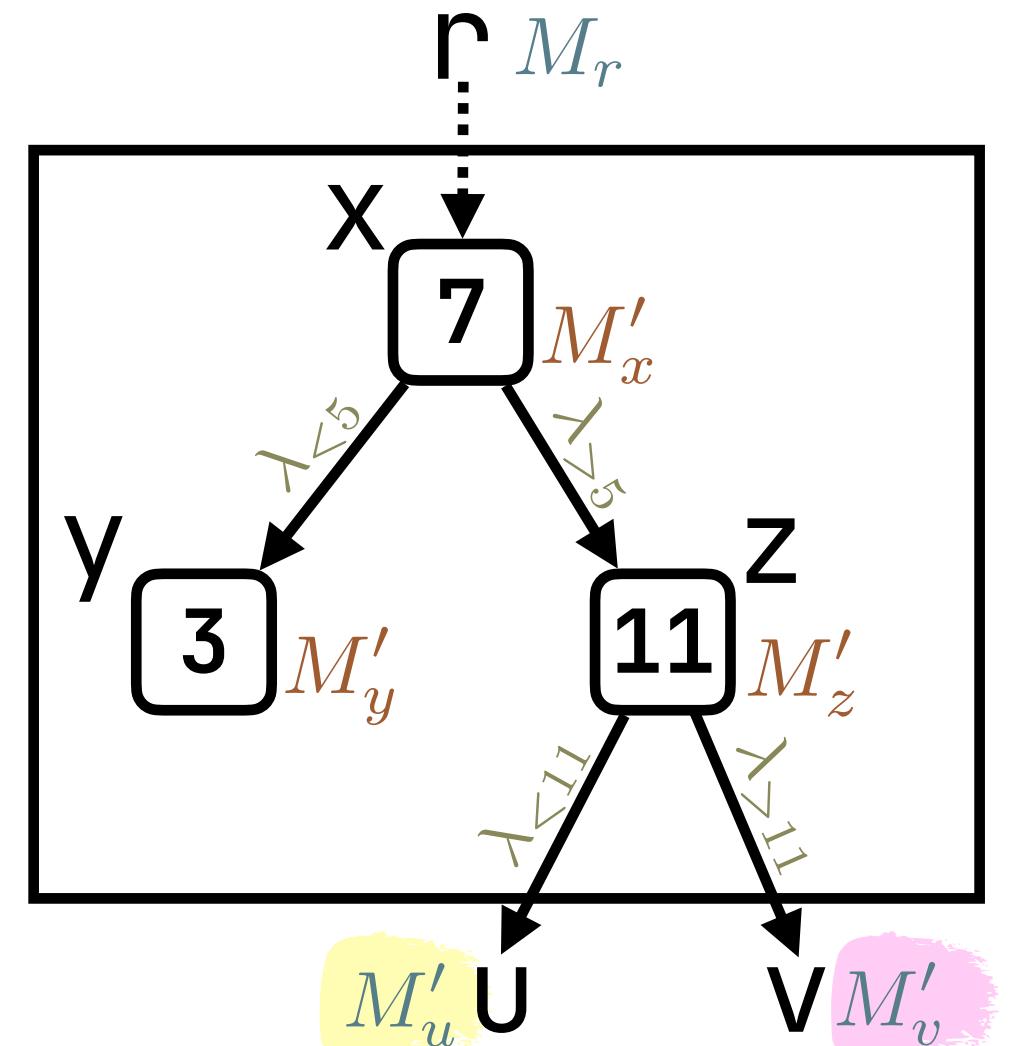
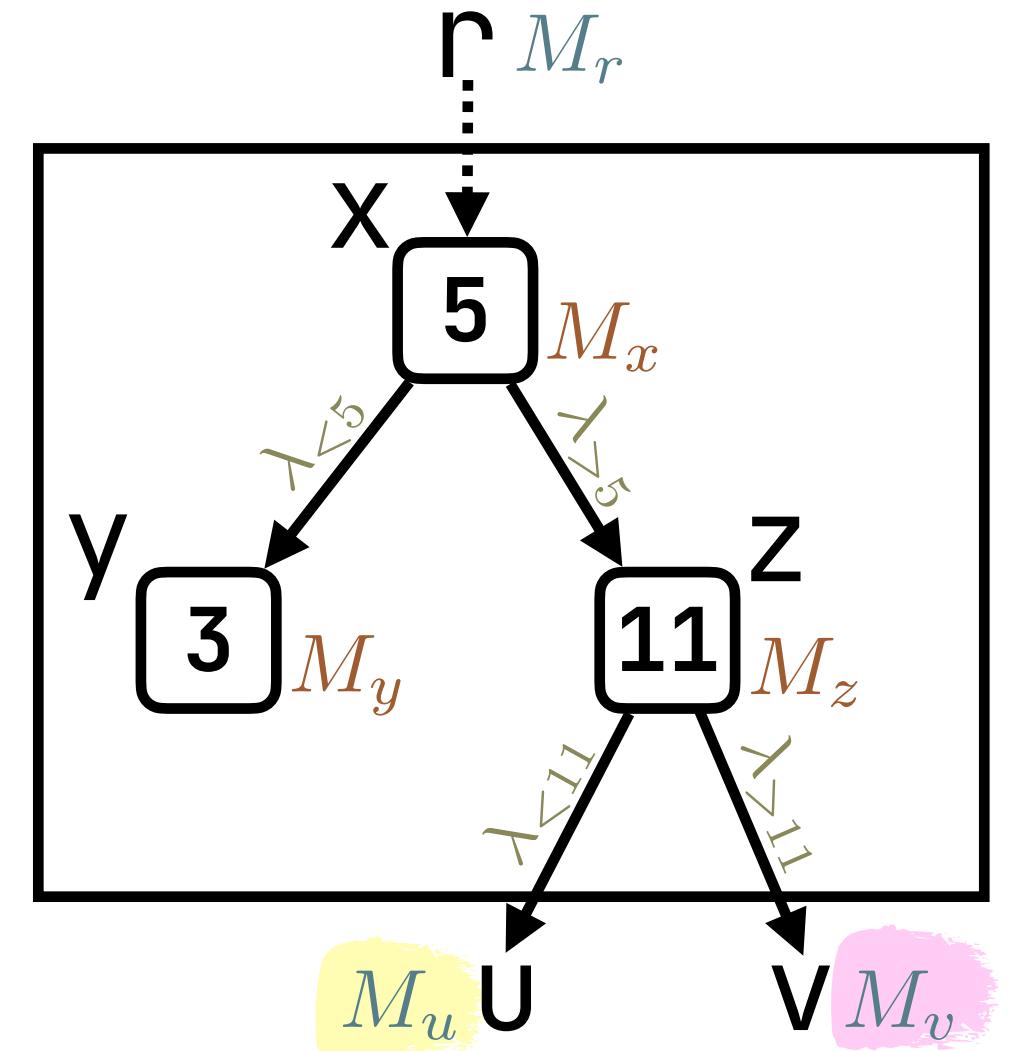
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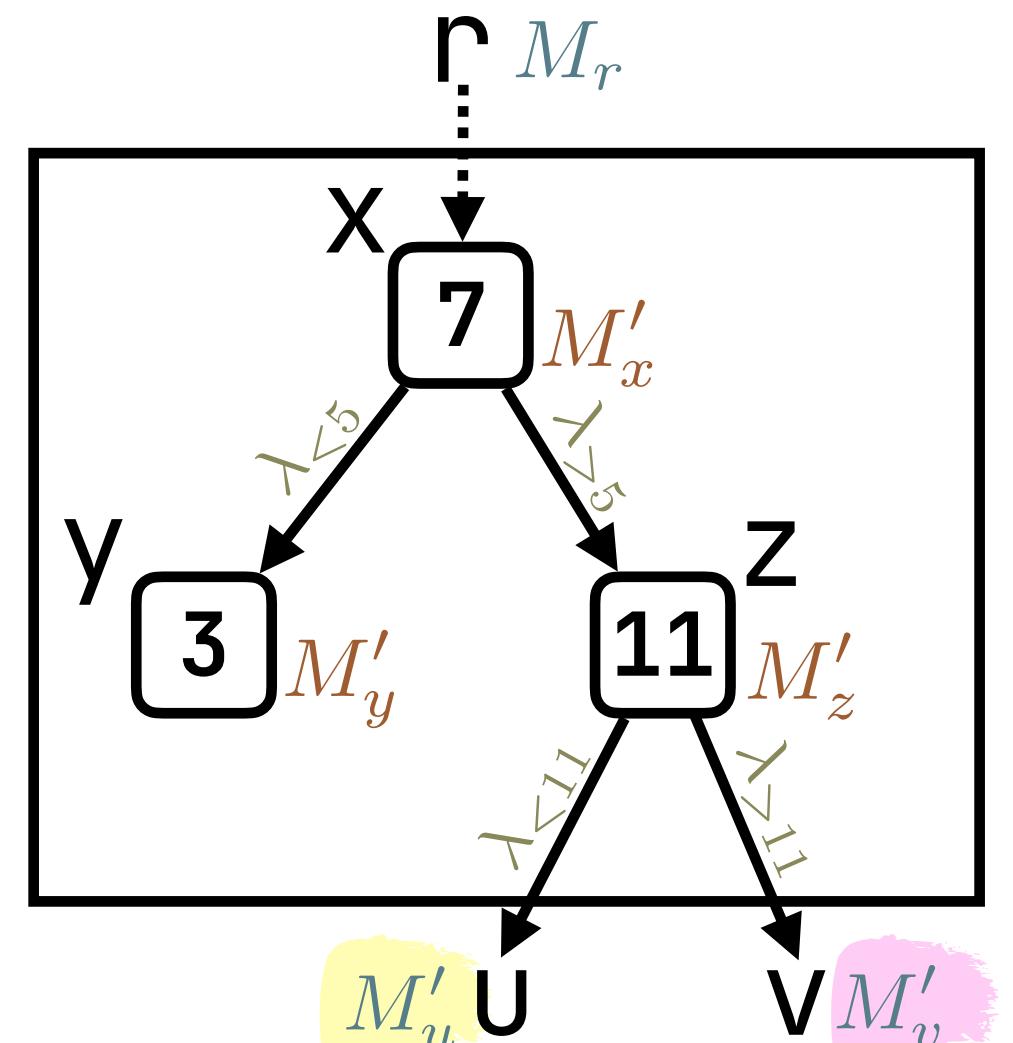
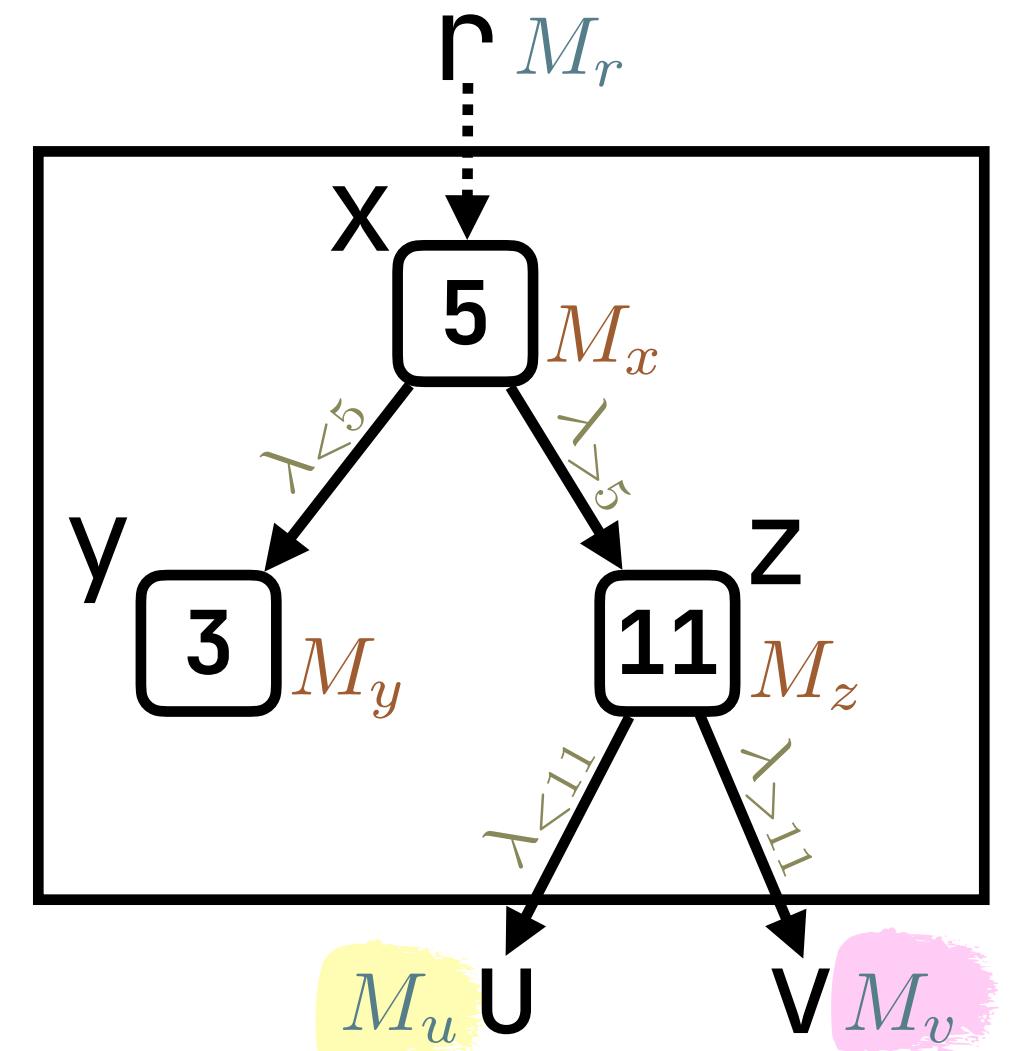
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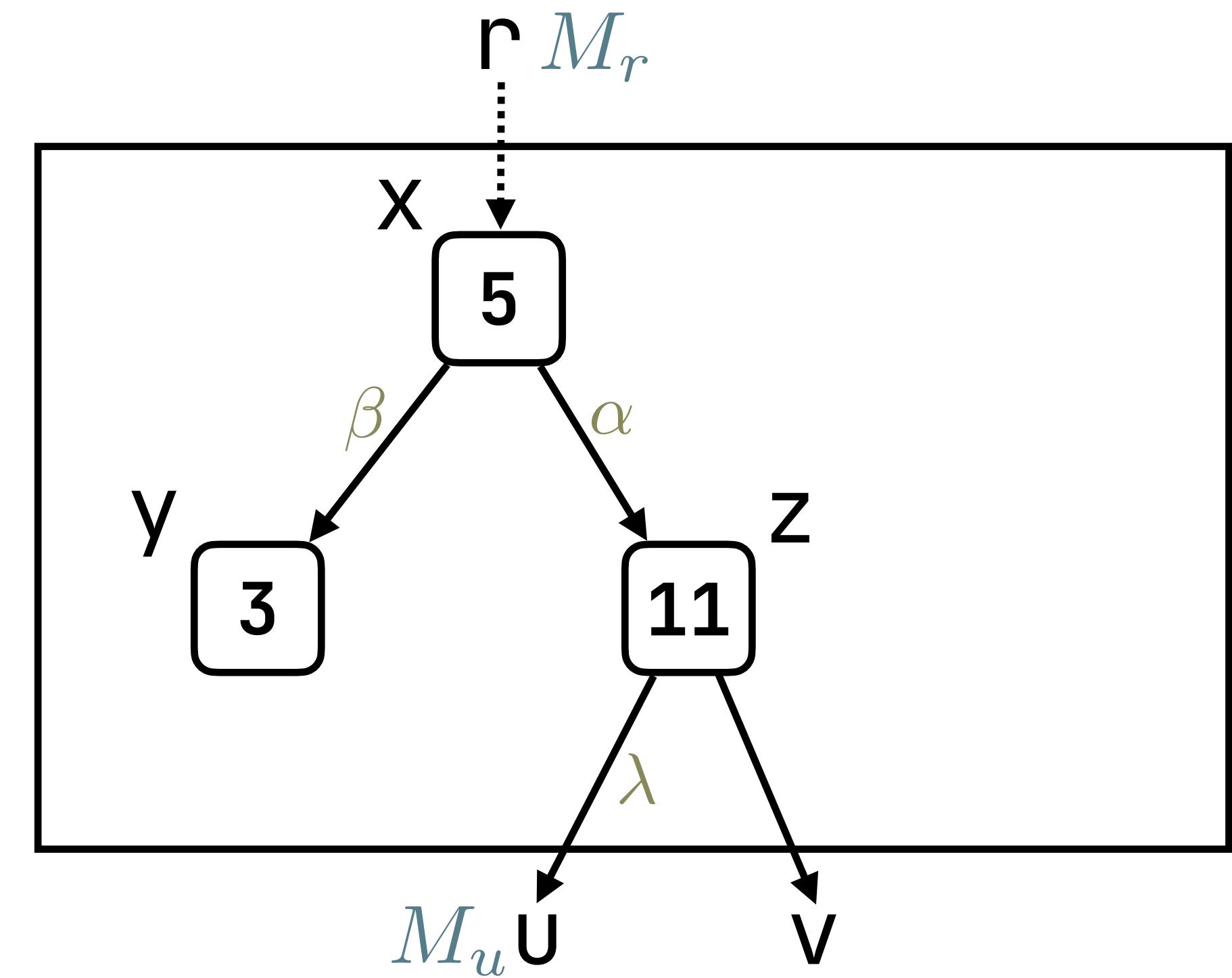
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- Goal: automated & efficient approach



Avoiding Fixed Points

- Observation for trees:

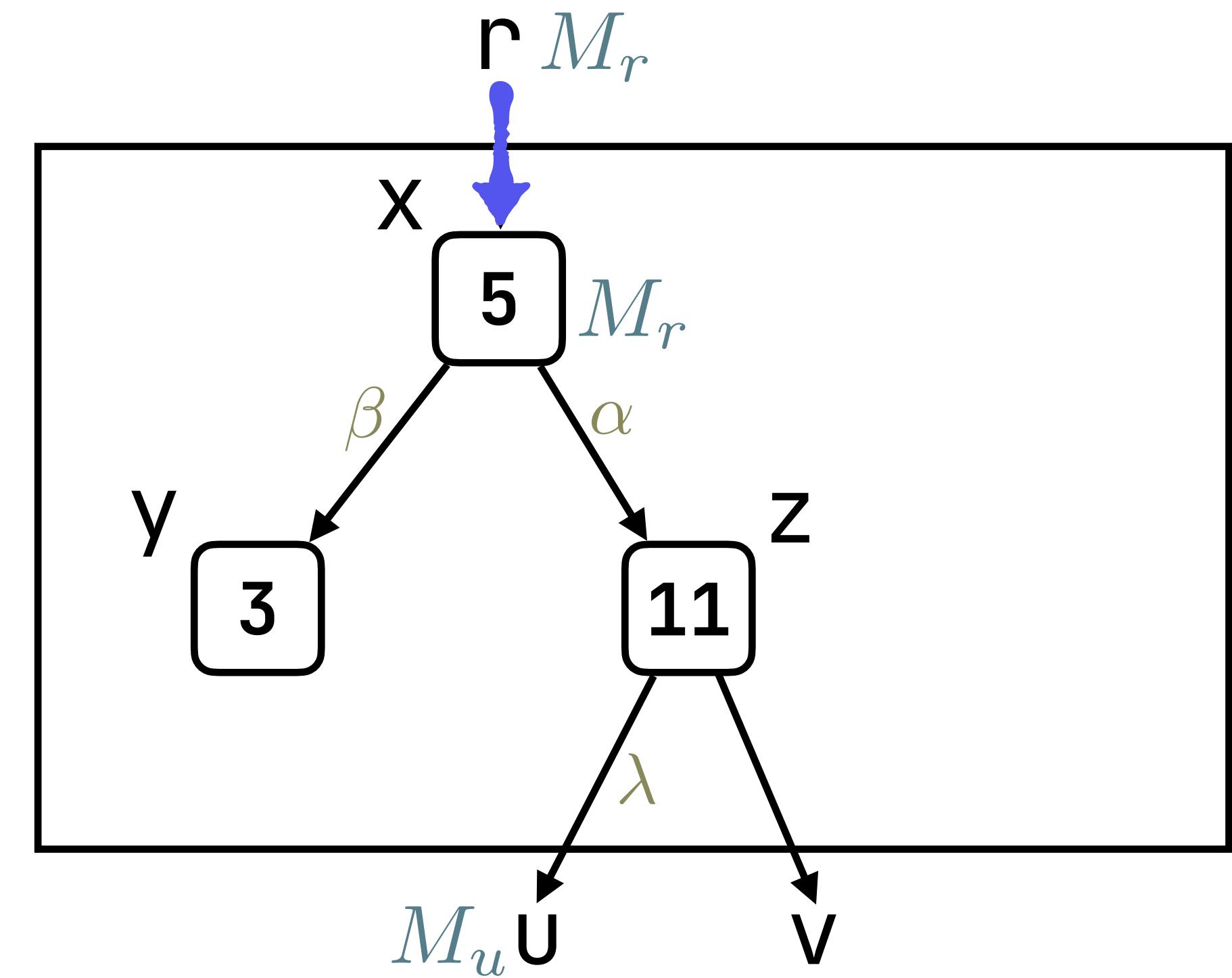
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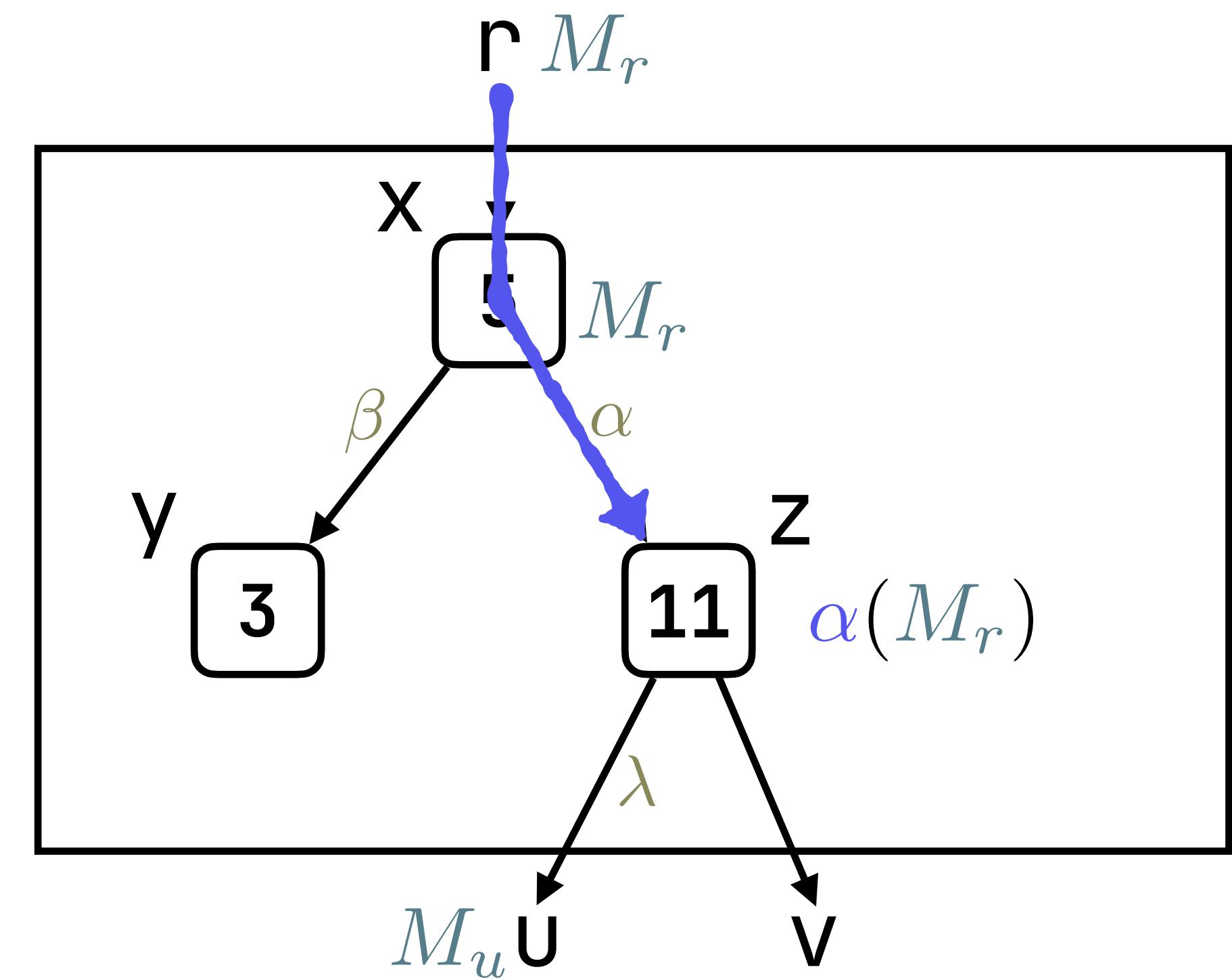
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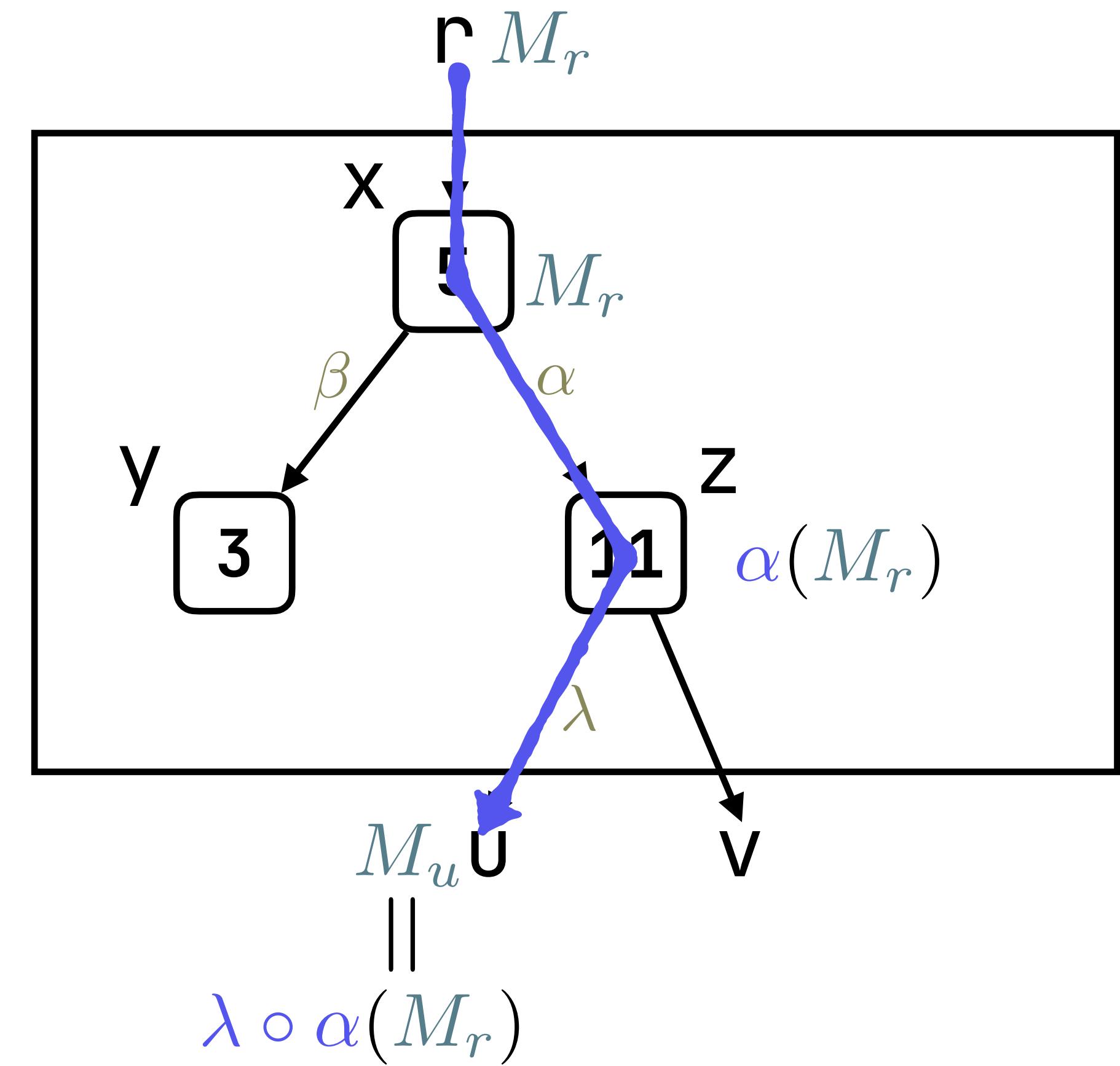
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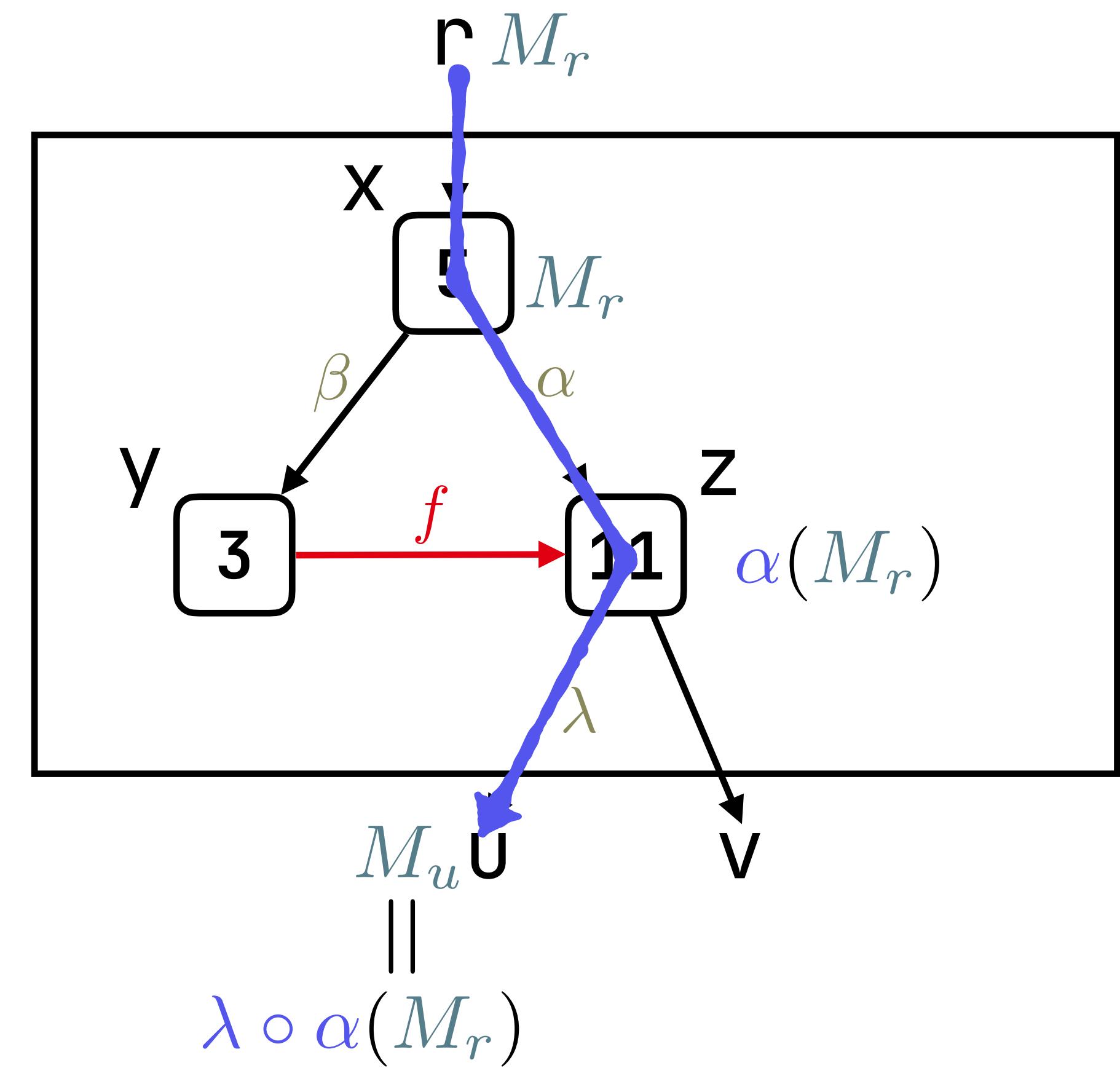


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→ not true in general

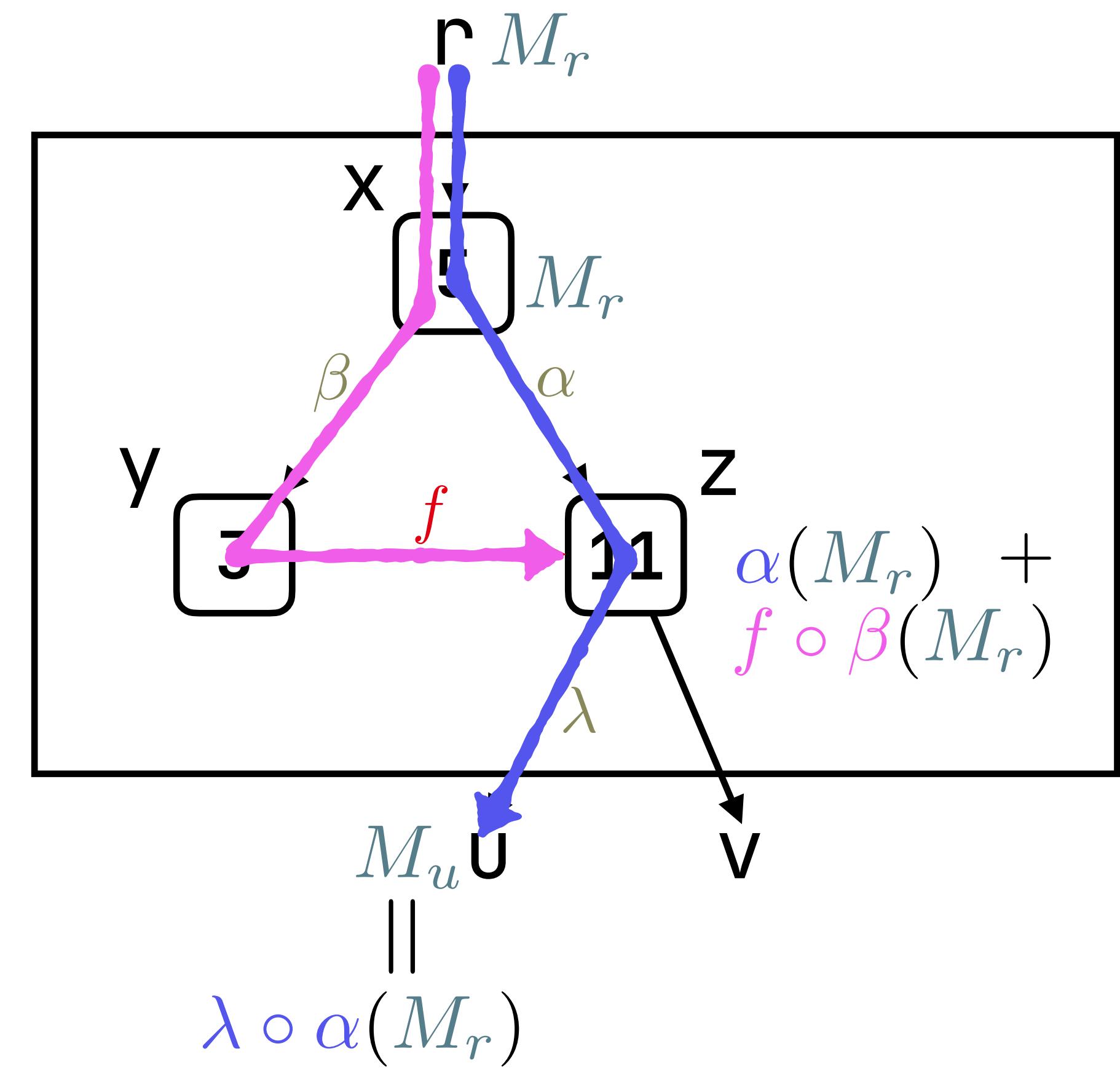


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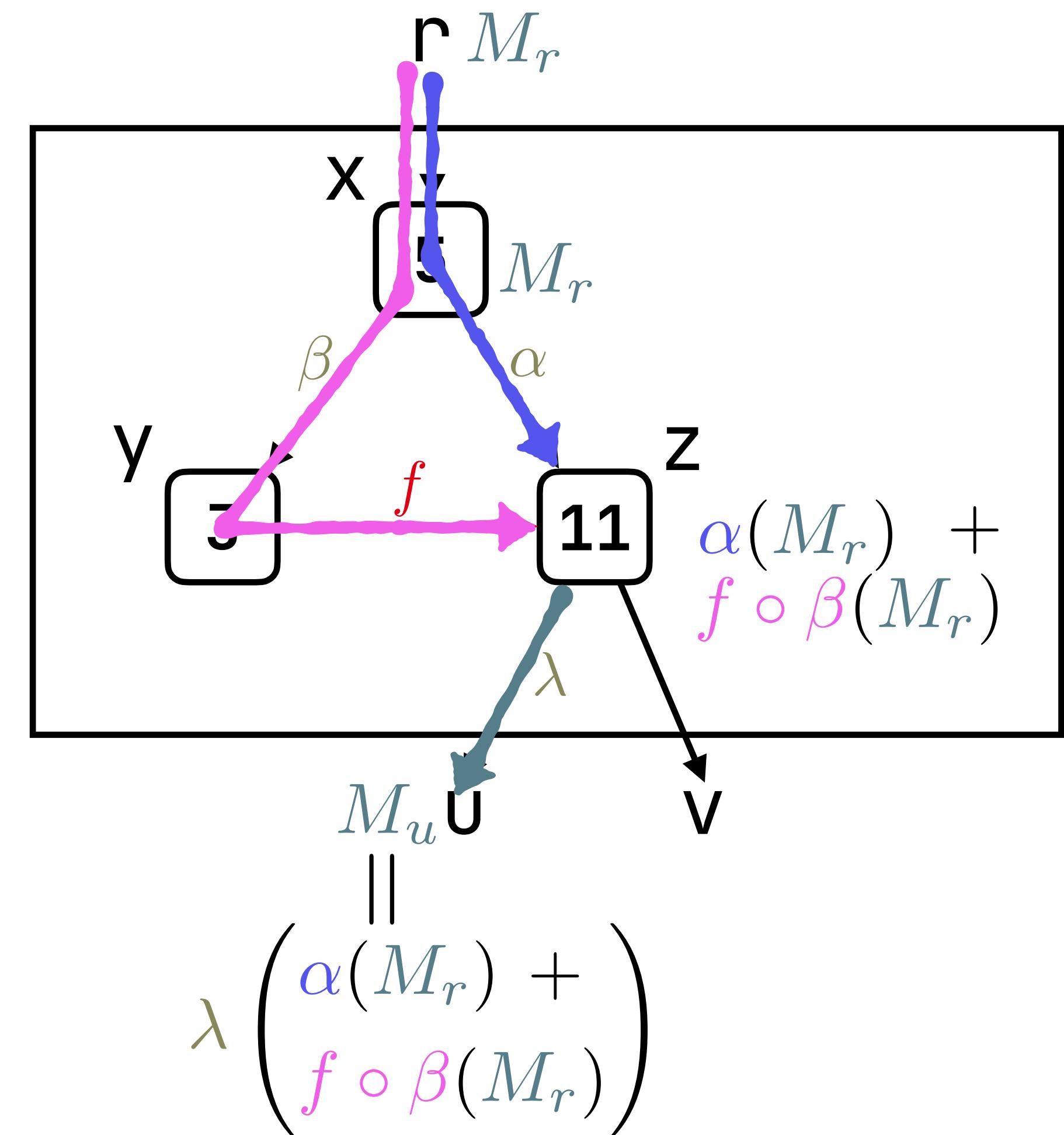


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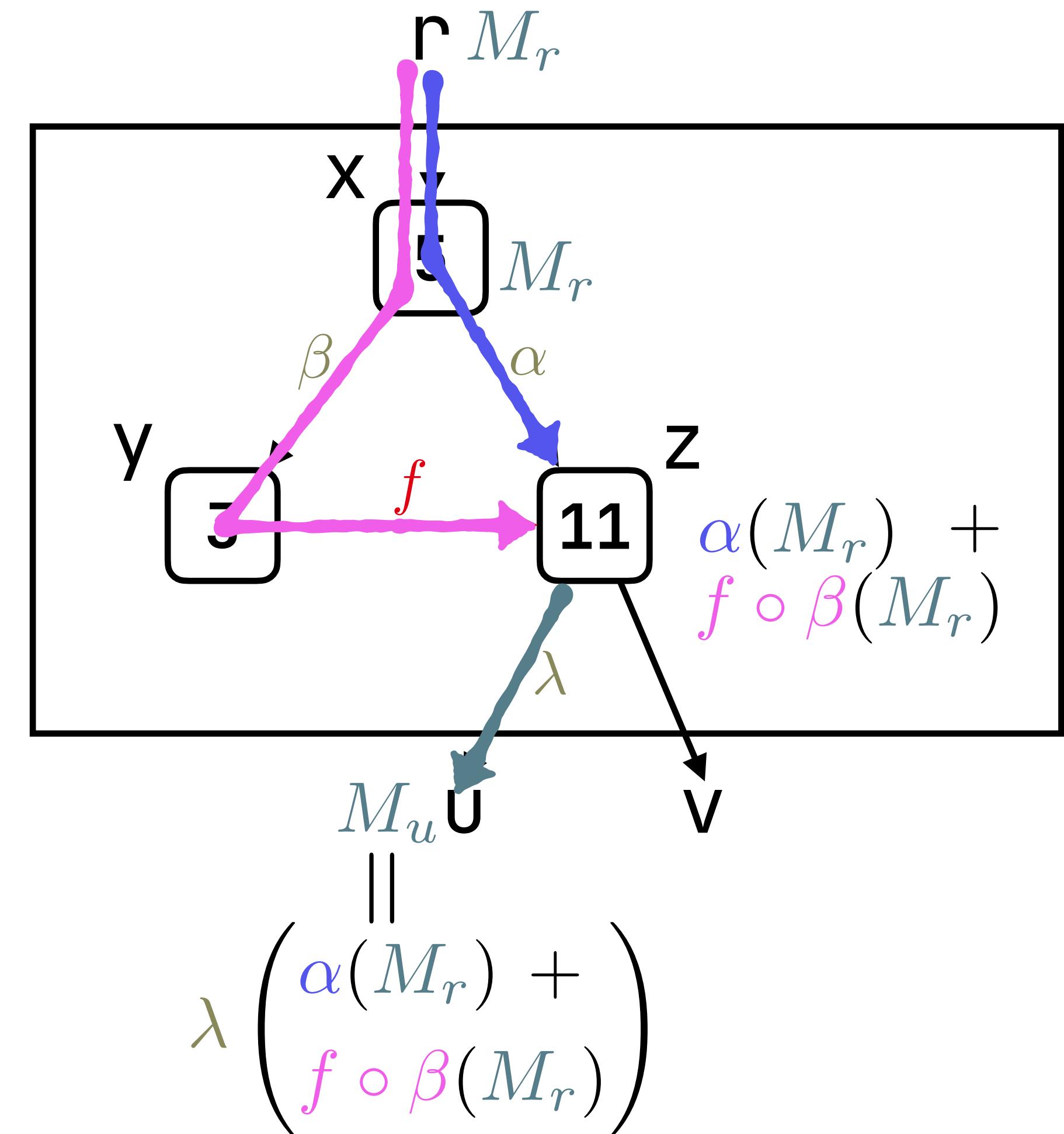
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$$f(m + n) = f(m) + f(n)$$

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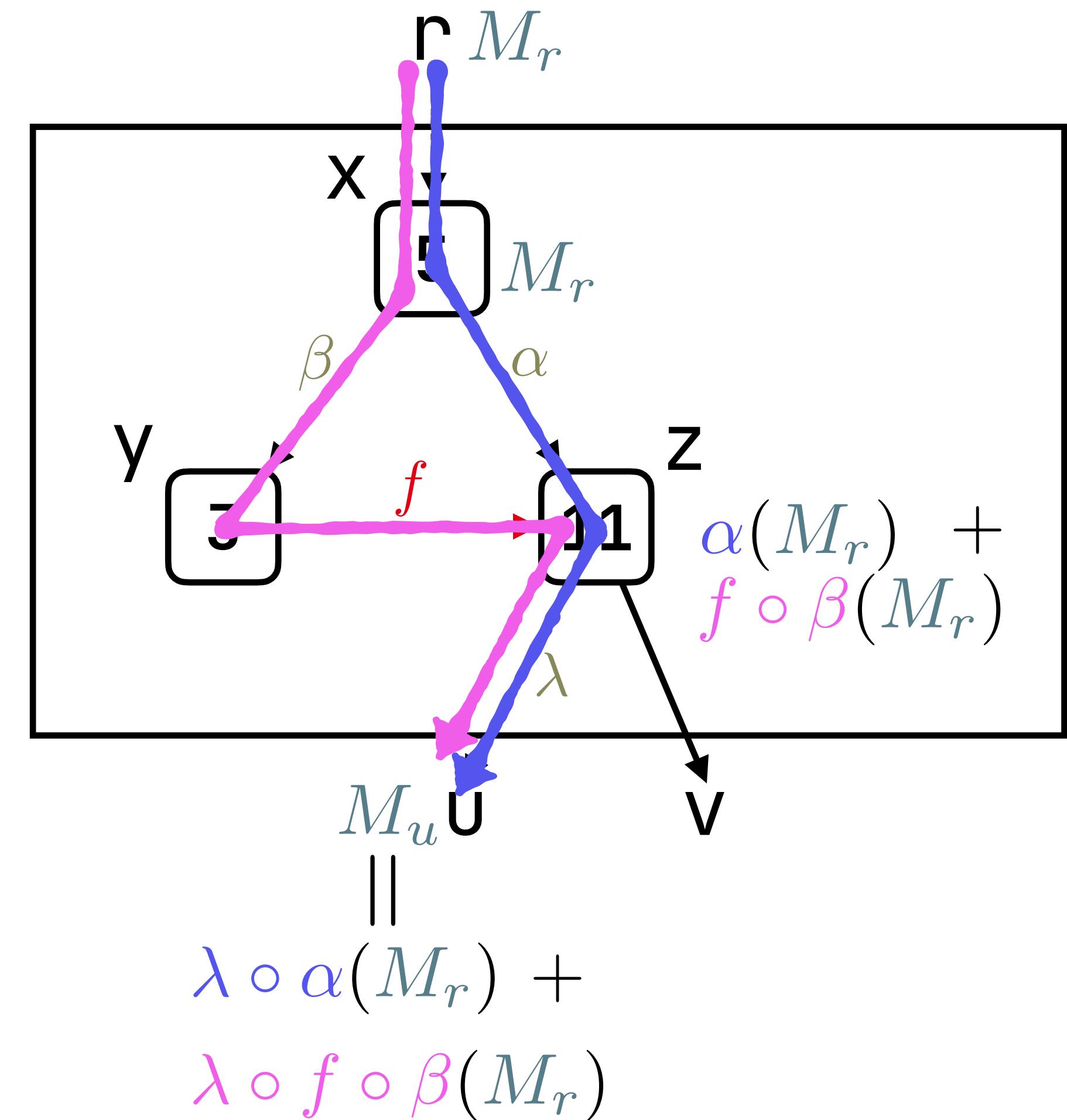
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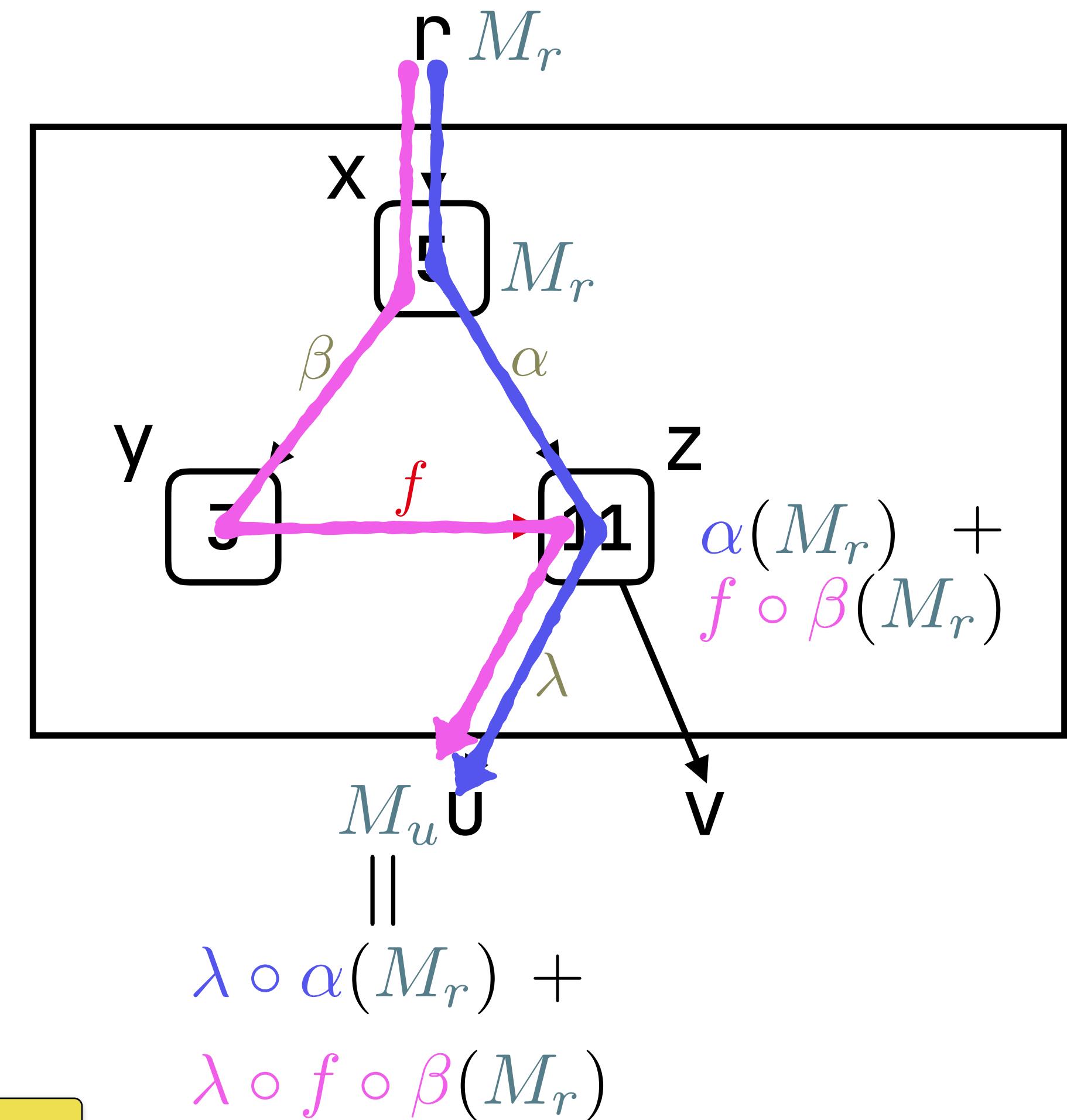
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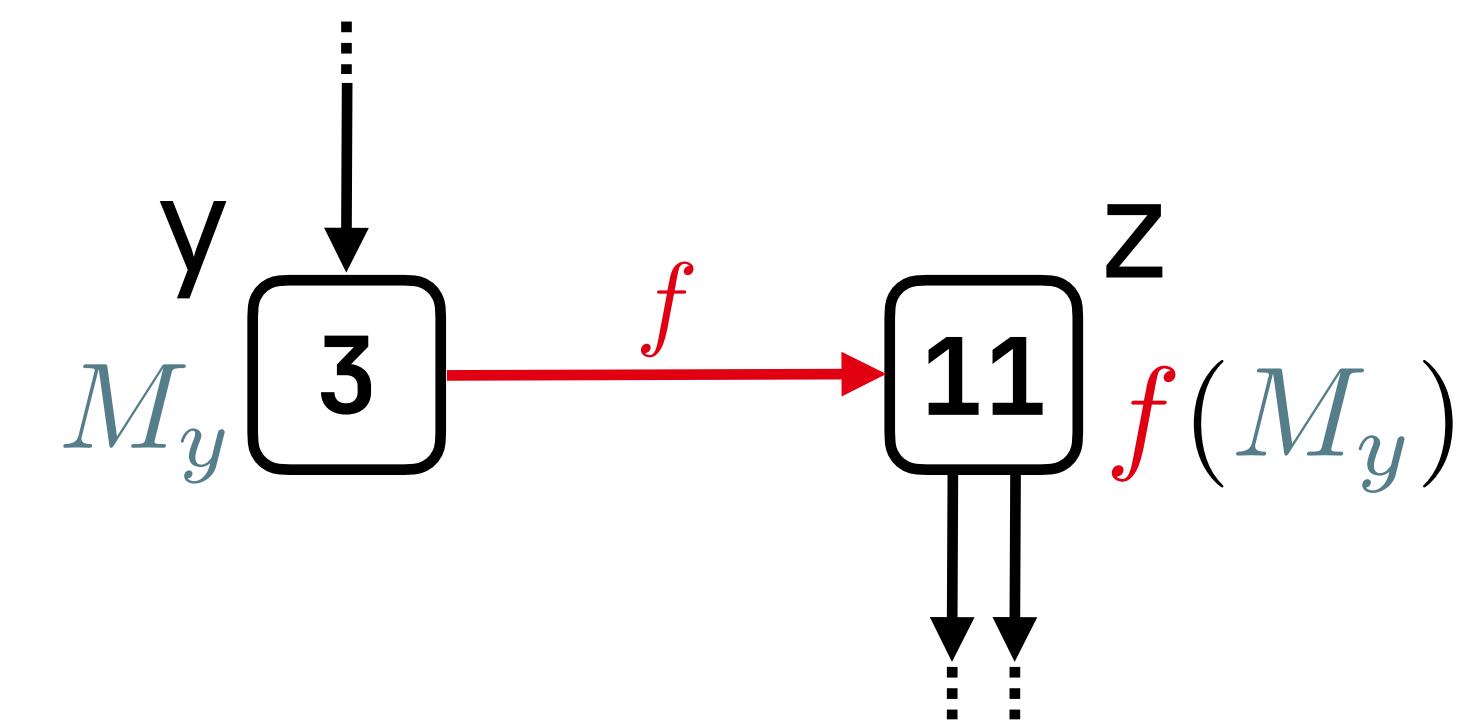
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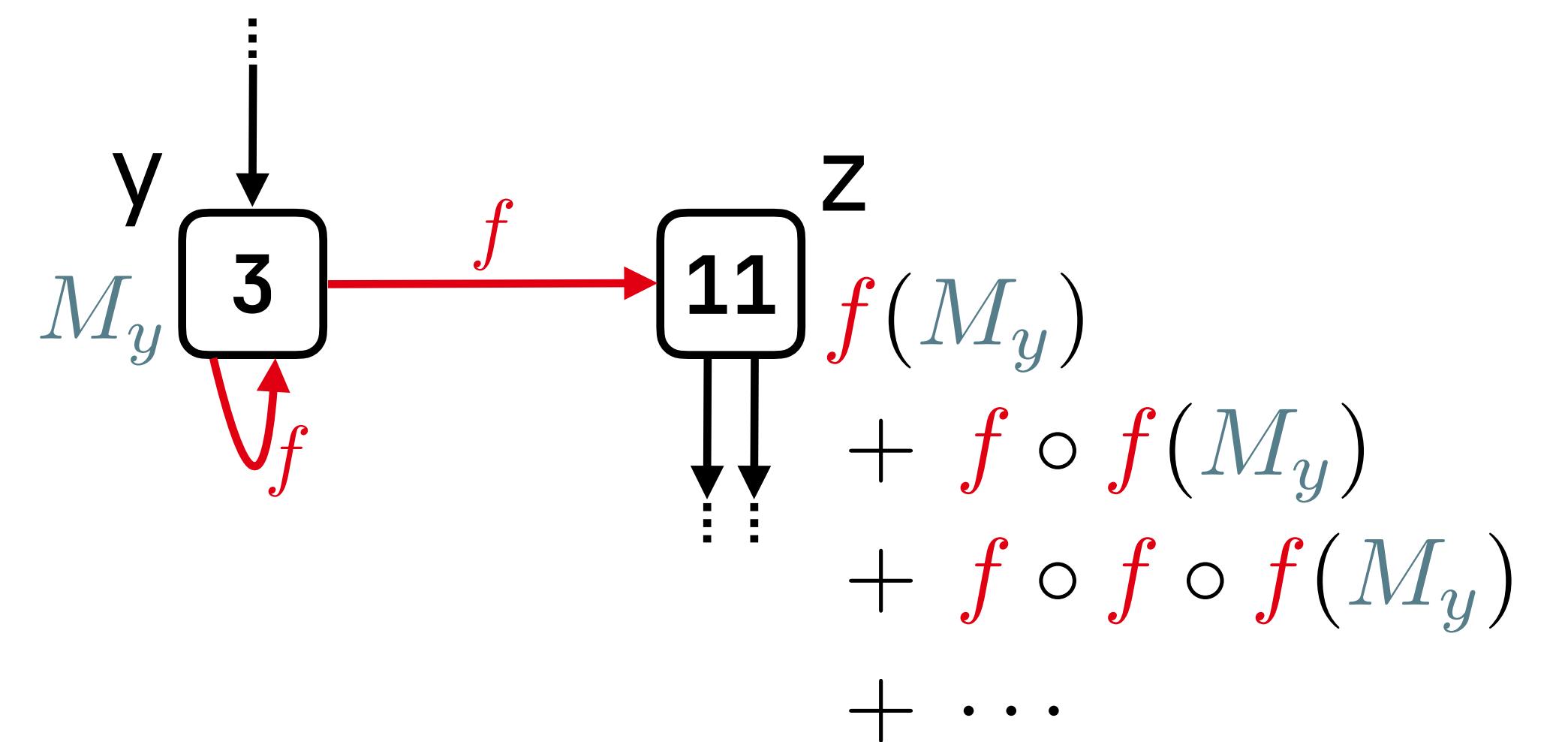
Infinite sum for cyclic graphs.



Handling Cycles



Handling Cycles

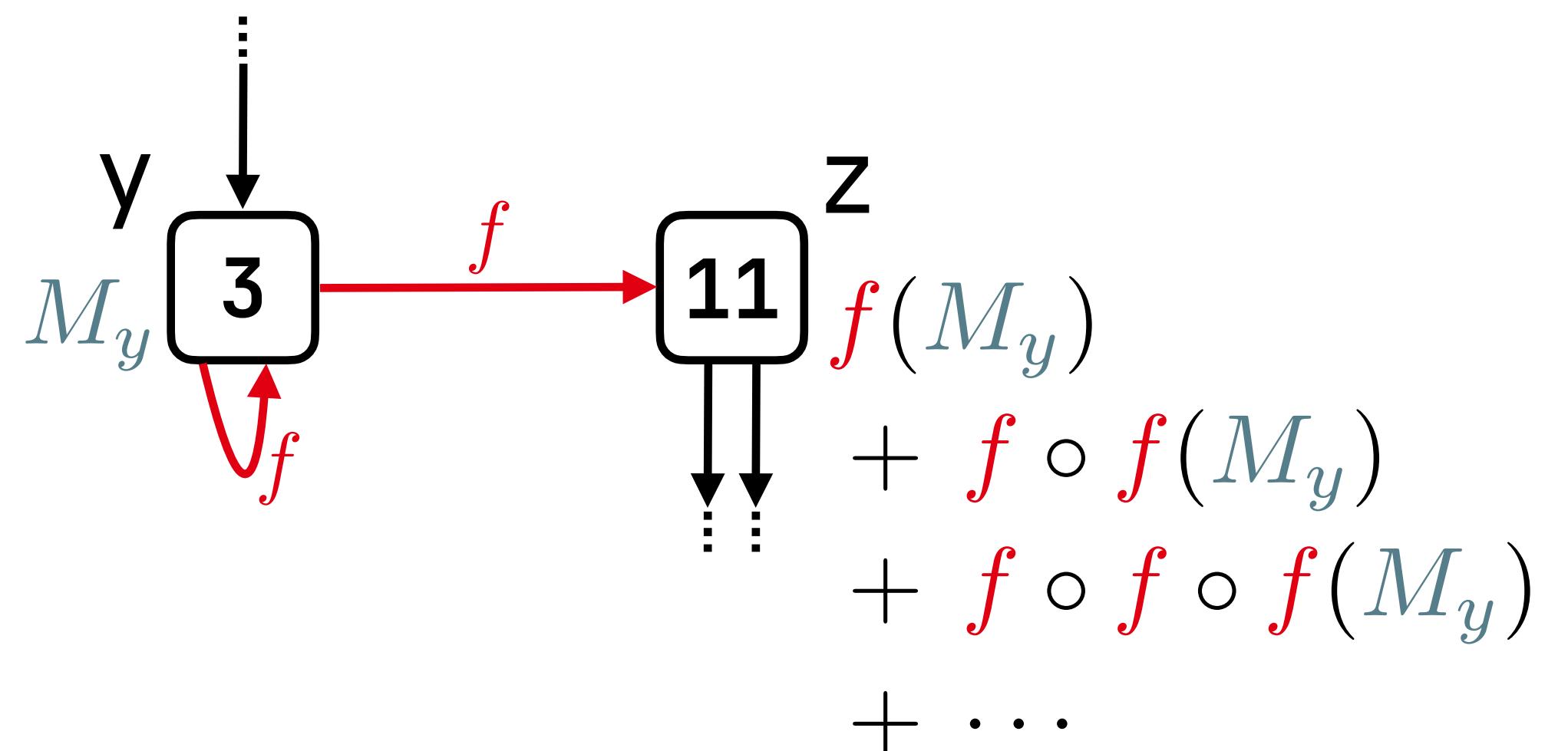


Handling Cycles

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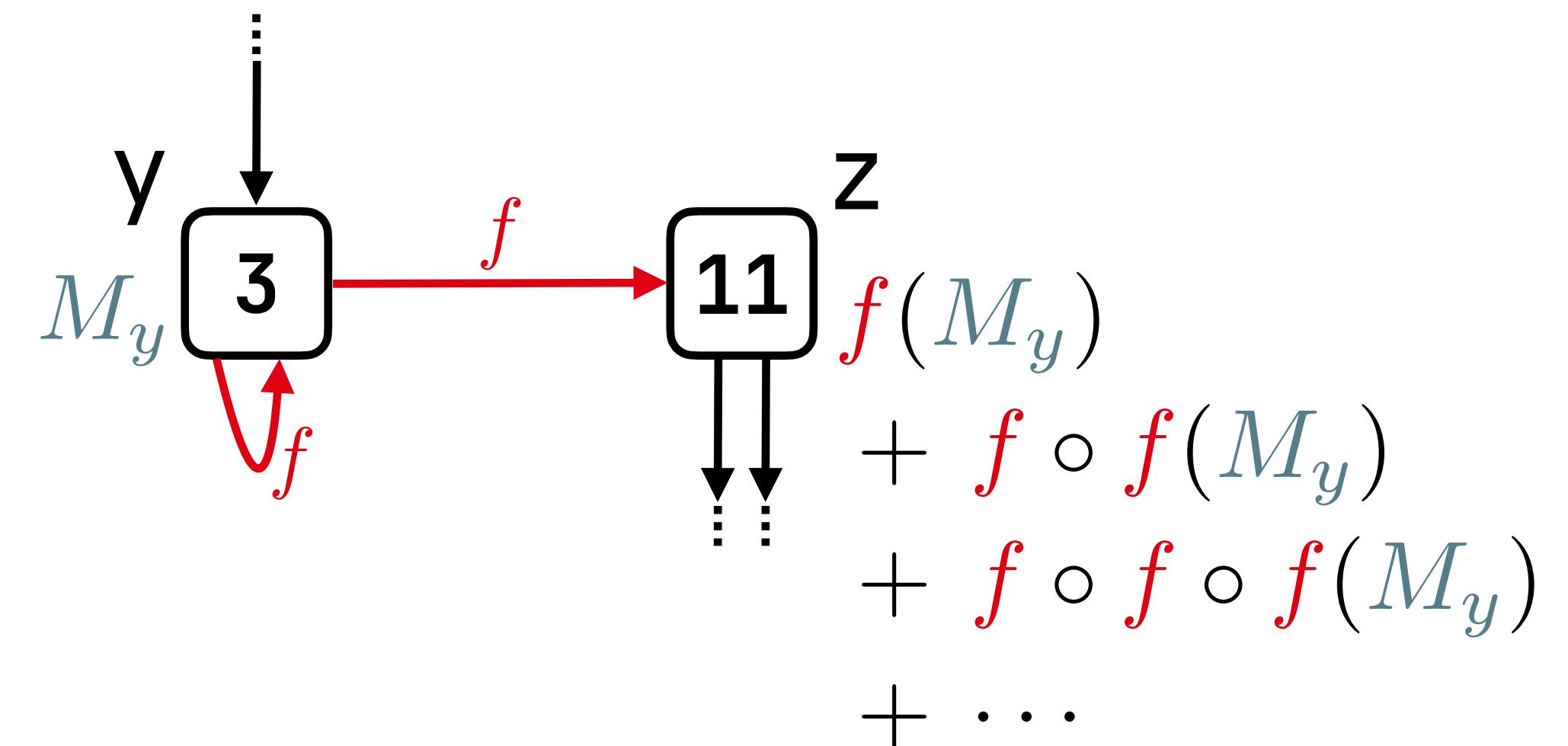
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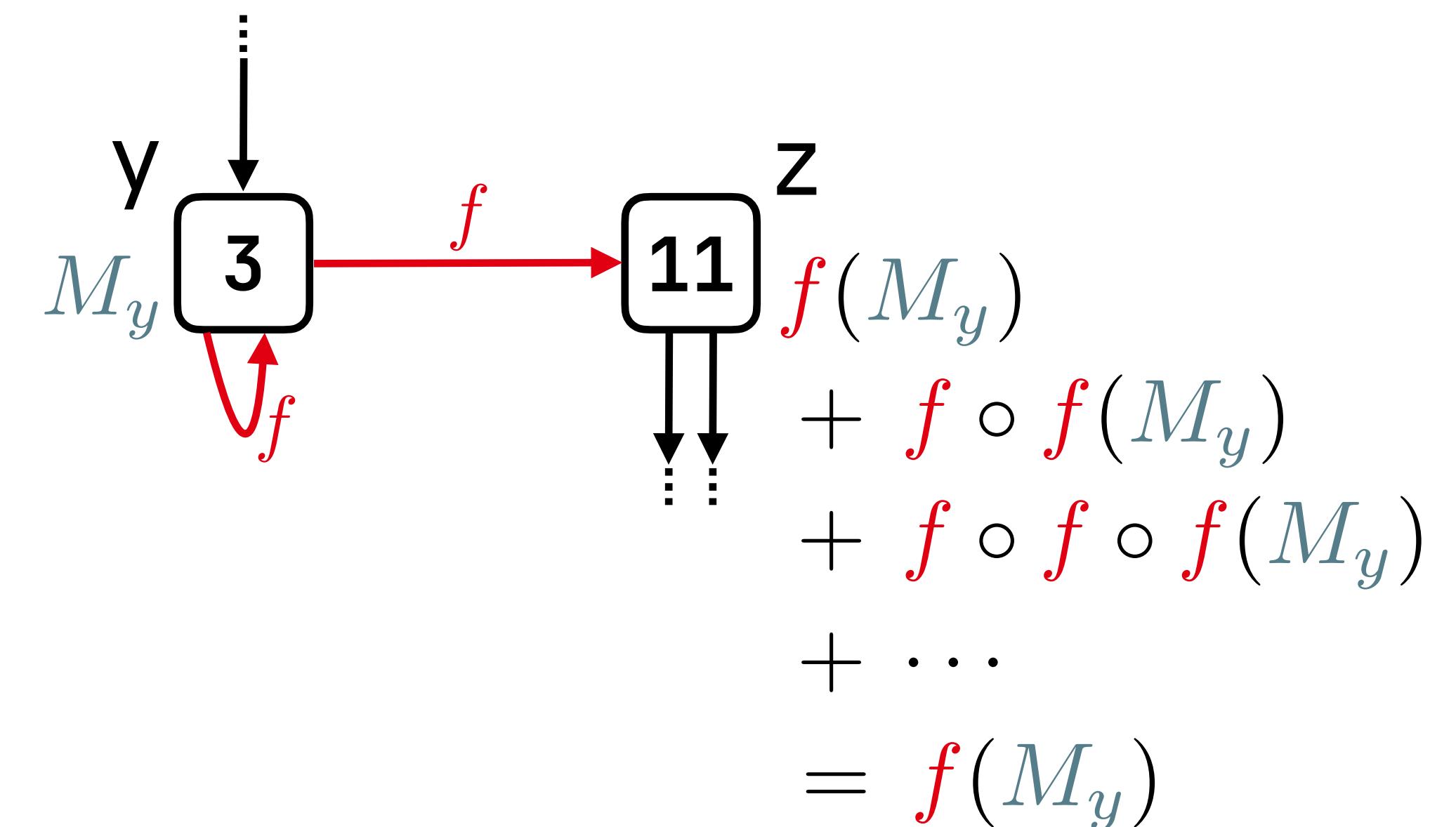
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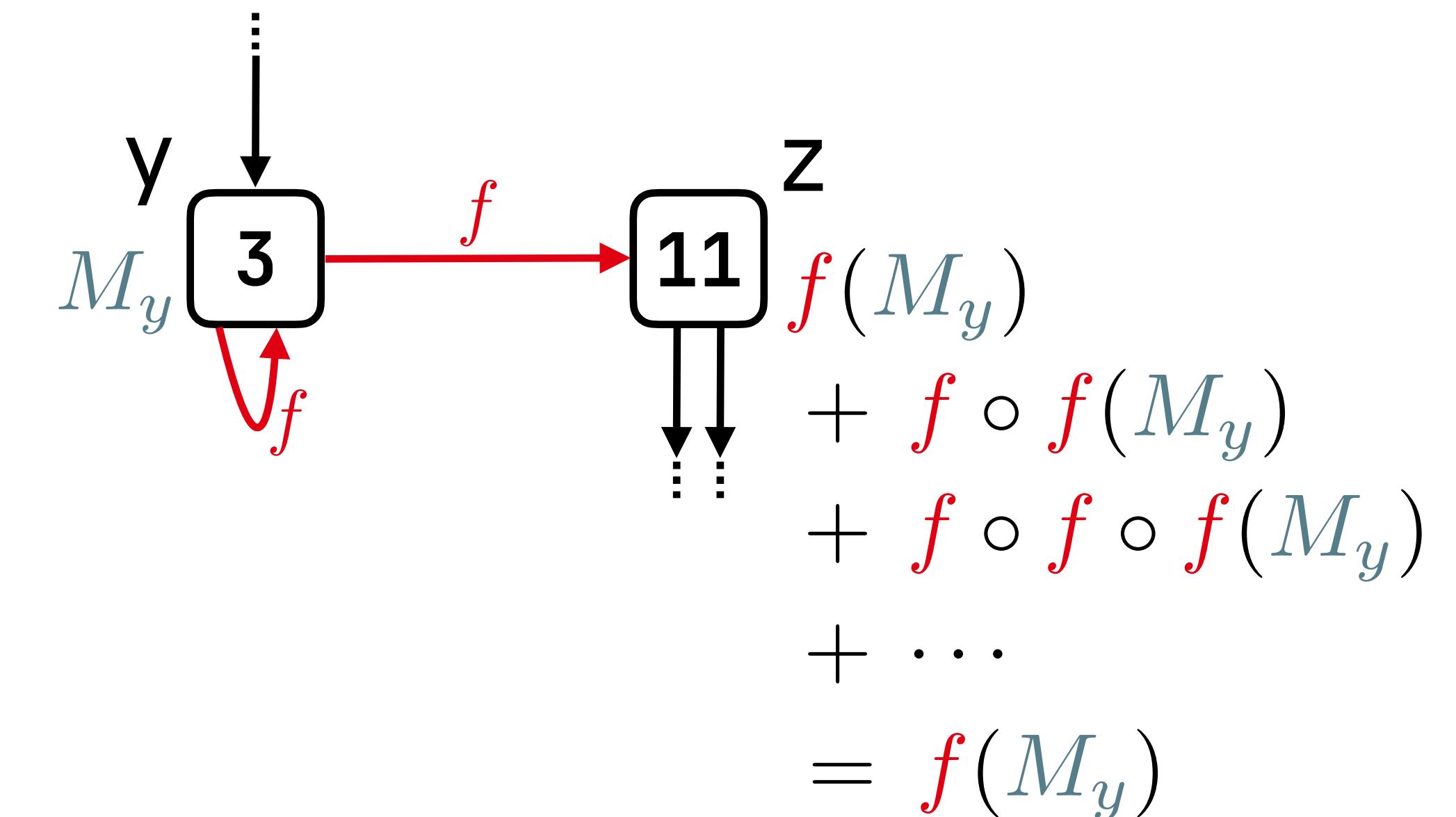
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Finite sum over all simple paths.

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Thanks

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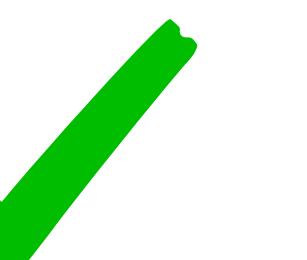
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