

CMPT 125: Introduction to Computing Science and Programming II

Fall 2023

Week 7: Data structures, revisiting recursion

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Fact of the day

The world's first webcam was used to watch a coffee pot to see if it is empty



<https://www.bbc.com/news/technology-20439301>

The webcam takes images 3 times a minute and uploads them

Recap from Last Lecture

- Implementation details
 - Stacks – making them resizable using realloc
 - Queues – using all the available spaces in a static array, making them resizable using malloc
 - Linked Lists – making use of memory allocation to create a data structure that stores items flexibly, with $O(1)$ time complexity for many operations
There is one case where the operation is $O(N)$ given the current implementation:
when we want to remove a node at the end (need to access the second-last node to update its next)

Review from Last Lecture

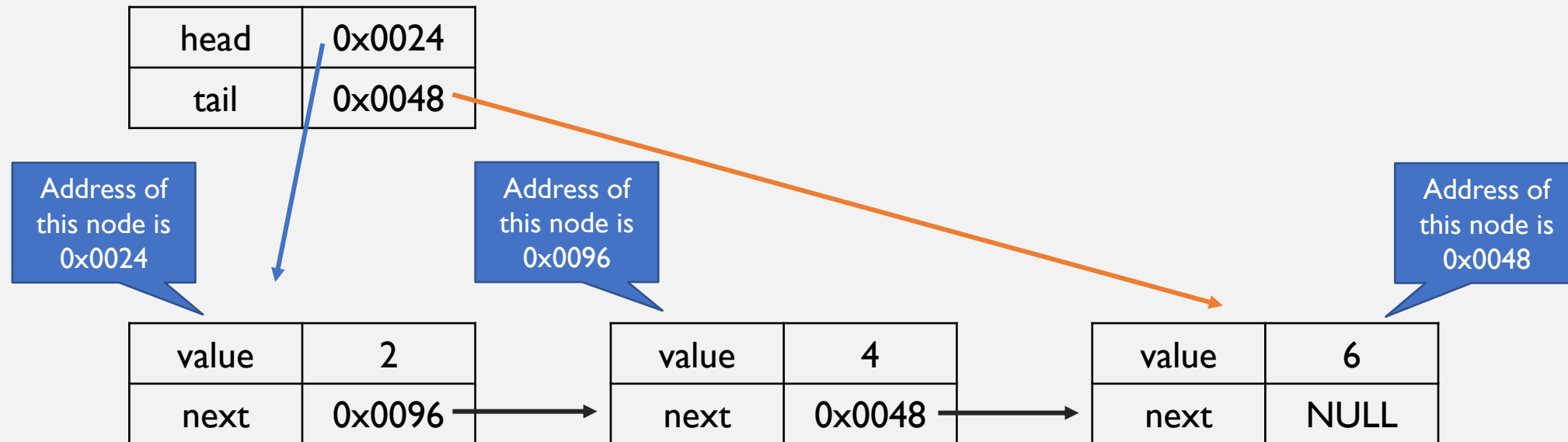
- Implement the queue ADT in C using an int array of size 10 (Size10Queue)
 - **General idea:** careful with the indexing and select a remove strategy that makes the most sense
 - Don't forget to insert a check for null pointers for the create/isEmpty/enqueue/dequeue functions

Today

- Linked Lists (cont'd)
 - Some implementation details (cases need to be considered)
- Live code demo
 - Linked list ADT

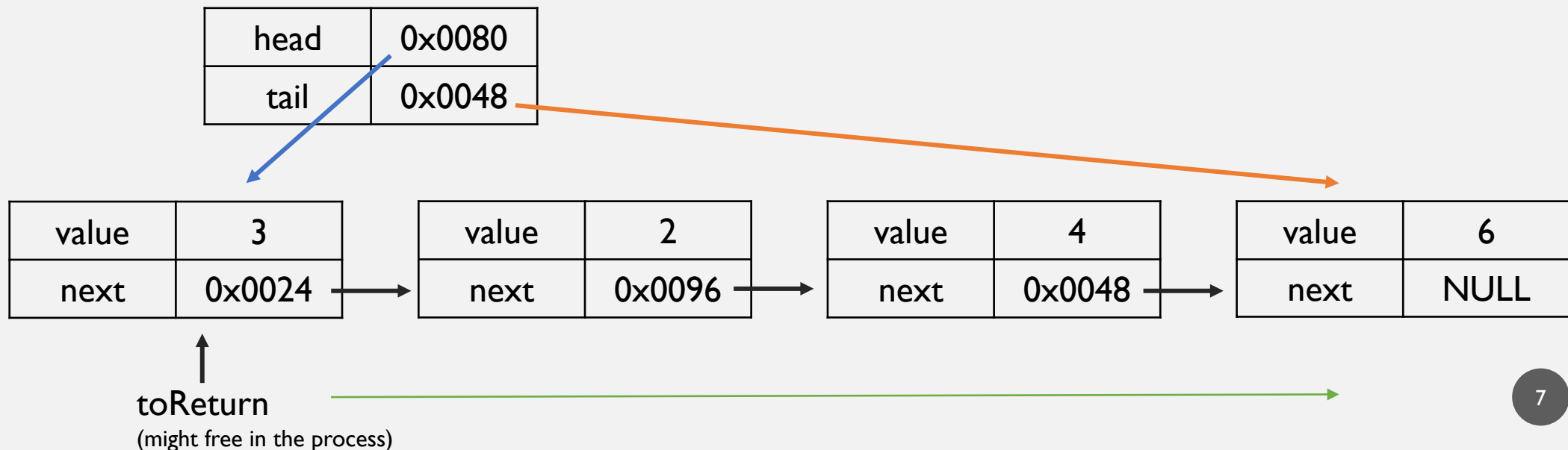
Linked Lists Structure (from Last Lecture)

- To store a linked list in C, we add an extra composite data type that works as a **header** that remembers the first and last node of the list (we call them **head** & **tail**)
 - Some implementations will only have the **head**, which is sufficient. But also having **tail** makes some functions faster



Linked Lists – Basic Operations (6)

- **Search:** Given the key value, do a linear search (*linked lists are not sorted by default*)
 - Create a node pointer called visitor, traverse from head and check each value using a while-loop to advance
 - When found, return the address of that node; if goes to the end of the list (visitor becomes NULL), not found, return NULL



Implementation Details of Linked List

Cases to consider...

2 Cases for InsertFront

- When inserting a node into the front of the linked list, need to consider if the list is empty or not

| | |
|------|------|
| head | NULL |
| tail | NULL |



| | | | | |
|------|---------------|---|-------|------|
| head | 0x0096 | → | value | 4 |
| tail | 0x0096 | → | next | NULL |

Address of
this node is
0x0096

Empty list case

| | | | | |
|------|--------|---|-------|------|
| head | 0x0096 | → | value | 4 |
| tail | 0x0096 | → | next | NULL |



| | | | | |
|------|---------------|---|-------|---------------|
| head | 0x0024 | → | value | 2 |
| tail | 0x0096 | → | next | 0x0096 |

Address of
this node is
0x0096

Address of
this node is
0x0024

Address of
this node is
0x0096

Non-empty list case

2 Cases for InsertEnd

- When inserting a node into the end of the linked list, also need to consider if the list is empty or not

| | |
|------|------|
| head | NULL |
| tail | NULL |



| | | | |
|------|---------------|-------|------|
| head | 0x0096 | value | 4 |
| tail | 0x0096 | next | NULL |

Address of
this node is
0x0096

Empty list case

| | | | |
|------|--------|-------|------|
| head | 0x0096 | value | 4 |
| tail | 0x0096 | next | NULL |



| | | | |
|------|---------------|-------|---------------|
| head | 0x0096 | value | 4 |
| tail | 0x0024 | next | 0x0024 |

Address of
this node is
0x0096

Address of
this node is
0x0096

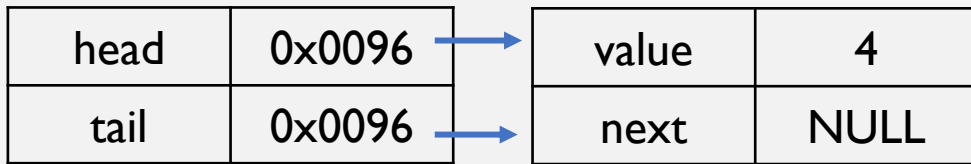
Address of
this node is
0x0024

| | |
|-------|------|
| value | 2 |
| next | NULL |

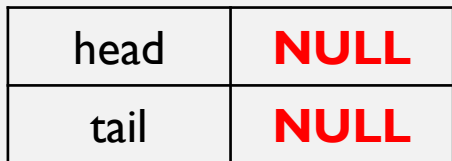
Non-empty list case

2 Cases for RemoveFront

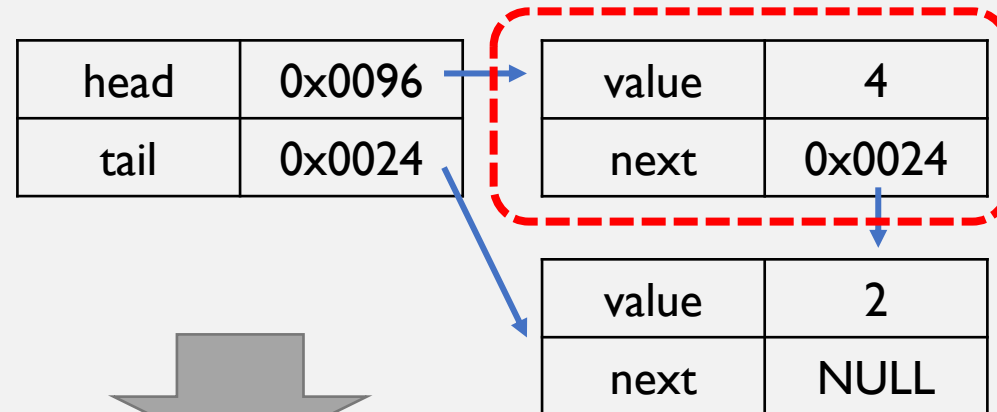
- When removing a node from the front of the linked list, need to consider if the list has 1 item or more



Address of
this node is
0x0096



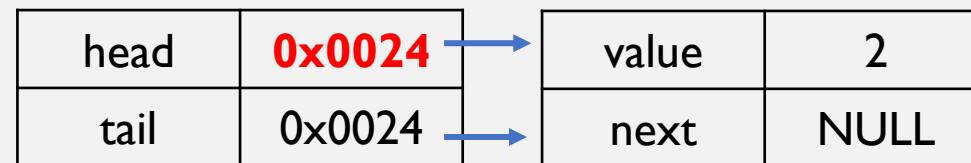
1-item list case



Address of
this node is
0x0096

Address of
this node is
0x0024

head points
to the 2nd
node of the
old list

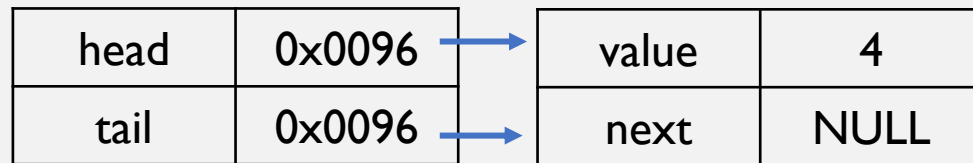


Address of
this node is
0x0024

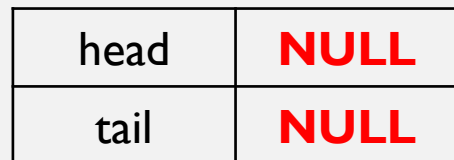
2+-items list case

2 Cases for RemoveEnd

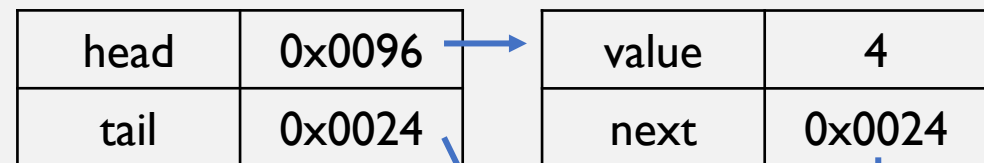
- When removing a node from the end of the linked list, also need to consider if the list has 1 item or more



Address of
this node is
0x0096

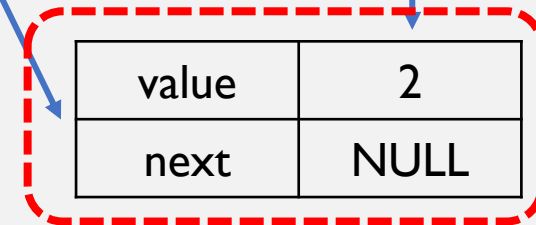


1-item list case

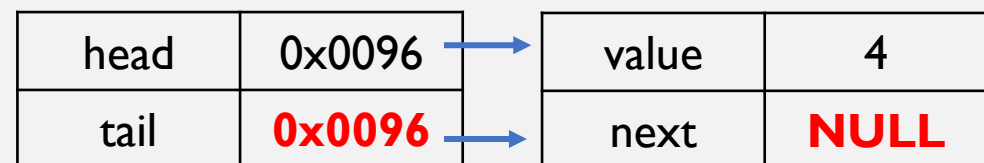


Address of
this node is
0x0096

Address of
this node is
0x0024



tail points to
the 2nd-last
node of the
old list



Address of
this node is
0x0096

2+-items list case

How to Find The 2nd-Last Node?

- **Main idea:** consider the fact that the next node of the 2nd-last node is the last node
 - **Option 1:** check if a node's next is the last node
 - **Option 2:** check if the next pointer of a node's next node is NULL
- This takes $O(n)$ time. Some variations of linked lists allow you to find it faster

```
//myList points to the header of a linked list
if (myList->head == NULL) { //empty list
    //do something for empty list
} else if (myList->head == myList->tail) { //1 node
    //do something for a 1-node list
} else { //2 or more nodes
    //start from the front
    Node_t* secondLast = myList->head;

    while (secondLast->next != myList->tail) {
        //follow the trail
        secondLast = secondLast->next;
    }

    //do something with the secondLast node
}
```

Performance of Linked List Operations

- Most operations with linked list is either constant time (a fixed number of steps), or linear time (number of steps proportional to the size of the list)
 - **create** – $O(1)$ because there is a fixed amount of steps
 - **adding/removing elements at the front or end** – $O(1)$
 - **search** – $O(n)$ because you are doing a linear search
- When compared to arrays:
 - **adding/removing elements anywhere** – $O(n)$ since you might have to resize and copy all the values
 - **search** – $O(n)$ because you are also doing a linear search
- In general traversing though a linked list takes more time because the program needs to trace the address of the nodes, where with arrays it is just pointer arithmetic due to its continuous memory allocation
- Linked lists can grow (and shrink) in size by inserting (and removing) more nodes and no realloc is needed

Typical Linked Lists Implementation in C

- The node is a struct

```
typedef struct Node {  
    int value;  
    struct Node* next;  
} Node_t;
```

- The header is also a struct

```
typedef struct {  
    Node_t* head;  
    Node_t* tail;  
} LList_t;
```

Careful: NOT the
size of a node

```
//start with a list  
LList_t* myList = malloc(sizeof(LList_t));  
myList->head = NULL;  
myList->tail = NULL;  
  
//add a node, typically done in a function  
myList->head = malloc(sizeof(Node_t));  
myList->head->value = 0; //access field of the first node  
myList->head->next = NULL; //access field of the first node  
myList->tail = myList->head;  
  
//don't forget to free the memory created if not used  
free(myList->head);  
free(myList);
```

Commonly Used Code in Linked Lists

- Very often you'll traverse a linked list, for example:
 - to find a node storing a certain value (search)
 - to add a node some where in-between
 - to remove a node some where in-between
 - to print its content
 - ...etc.

```
//start from the front
Node_t* current = myList->head;

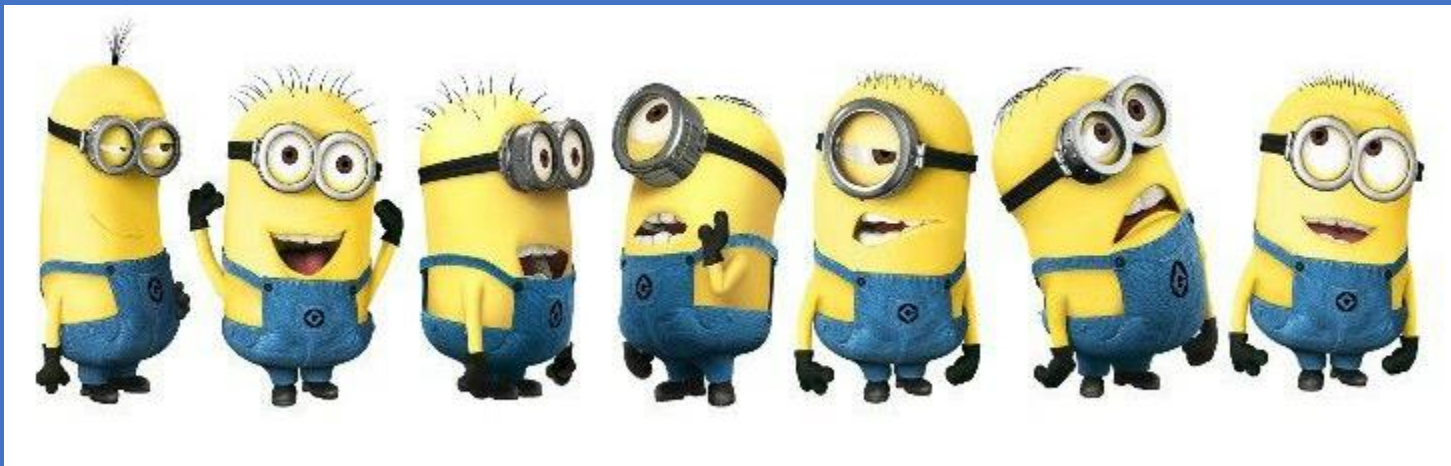
while (current != NULL) { //stop when there is nothing

    //do something with the current value, e.g., print
    printf("%d ", current->value);

    //follow the trail
    current = current->next;
}
```




Break for 10 Minutes



Putting It All Together

Live coding!

Things We Need

- A struct for Node (`Node_t`) + a struct for LinkedList (`LList_t`)
- A function to create an empty list (`LL_create`), returns the address of the list
- A function to free up the list (`LL_free`)
- A function to add an int value to the front of the list as a node (`LL_insertFront`)
- A function to remove the node in the front of the list (`LL_removeFront`), returns the value
- A function to tell if the list is empty (`LL_isEmpty`), returns 1 if empty, 0 otherwise

- Bonus 1: A function to print the content of the list (`LL_print`)
- Bonus 2: A function to get the length of the list (`LL_length`), returns the number of nodes in the list

```
typedef struct Node {
    int value;
    struct Node* next;
} Node_t;
```

```
typedef struct {
    Node_t* head;
    Node_t* tail;
} LList_t;
```

```
LList_t* LL_create() {
    LList_t* theList = malloc(sizeof(LList_t));
    theList->head = NULL;
    theList->tail = NULL;
}
```

It should also do a null test and return theList in the end

```
void LL_free(LList_t* theList) {
    if (theList == NULL) { return; }

    Node_t* current = theList->head;
    while(current != NULL) {
        Node_t* toFree = current;
        current = current->next;
        free(toFree);
    }
    free(theList);
}
```

```
void LL_insertFront(LList_t* theList, int value) {
    if (theList == NULL) { return; }

    Node_t* newNode = malloc(sizeof(Node_t));
    newNode->value = value;
    newNode->next = NULL;

    if (theList->head == NULL) { //empty list
        theList->head = newNode;
        theList->tail = newNode;
    } else { //non-empty list
        newNode->next = theList->head;
        theList->head = newNode;
    }
}
```

```
int LL_removeFront(LList_t* theList) {
    if (theList == NULL || theList->head == NULL) {
        return -99999; //this means error
    }

    Node_t* toRemove = theList->head;
    if (theList->head == theList->tail) { //1 item
        theList->head = NULL;
        theList->tail = NULL;
    } else { //2+ items
        theList->head = theList->head->next;
    }
    int theValue = toRemove->value;
    free(toRemove); //for this implementation
    return theValue;
}
```

```
int LL_isEmpty(LList_t* theList) {
    if (theList == NULL || theList->head == NULL) {
        return 1;
    } else {
        return 0;
    }
}
```

We'll do LL_print and LL_length if there's time otherwise they'll be an exercise 😊

```

void LL_insertEnd(LList_t* theList, int value) {
    if (theList == NULL) { return; }

    Node_t* newNode = malloc(sizeof(Node_t));
    newNode->value = value;
    newNode->next = NULL;

    if (theList->head == NULL) { //empty list
        theList->head = newNode;
        theList->tail = newNode;
    } else { //non-empty list
        theList->tail->next = newNode;
        theList->tail = newNode;
    }
}

```

```

int LL_removeEnd(LList_t* theList) {
    if (theList == NULL || theList->head == NULL) {
        return -99999; //this means error
    }

    Node_t* toRemove = theList->tail;
    if (theList->head == theList->tail) { //1 item
        theList->head = NULL;
        theList->tail = NULL;
    } else { //2+ items
        Node_t* secondLast = theList->head;
        while (secondLast->next != theList->tail) {
            secondLast = secondLast->next;
            printf("on the move\n");
        }
        theList->tail = secondLast;
        secondLast->next = NULL;
    }

    int theValue = toRemove->value;
    free(toRemove); //for this implementation
    return theValue;
}

```

Find the 2nd-
last node

Today's Review

- Linked Lists
 - A data structure that stores items in a specific way to provide flexible sizes with $O(1)$ time complexity for many operations
 - More efficient implementation of ADTs – using `insertFront/insertEnd/removeFront/removeEnd`
 - Need to consider the state of the linked lists for different cases (empty/1-item/2+-items)
- Code live demo
 - Using structs and pointers to implement the list
 - Different functions accessing the list via a pointer to it
 - check for null, check for empty/1-item/2+-items
 - use while-loop to traverse

Homework!

- Implement the rest of the linked list data structure
 - insertEnd, removeEnd (*you'll need the trick to find the 2nd-last node*)
 - move the code to LinkedList.h and LinkedList.c
- Write some testing code to use the linked list
- Think about how to insert/delete a node to/from the inside (not head/tail) of a linked list
 - We'll talk about this next week
- Think about how linked list is related to Stacks & Queues
 - We'll also talk about this next week