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基础

头文件

二进制函数 枚举组合数

```
1 /* builtin clz 前导 0
   * __builtin_ctz 后缀 0
   #define ONES(x) __builtin_popcount(x) // 1 数目
   // 枚举 C(n, k) 所有可能, 复杂度 O(C(n, k))
   int next_combination(int n, int k) {
    int ret, b = k \& -k, t = (k + b);
    ret = (((t ^k) >> 2) / b) | t;
    if (ret >= (1 << n)) return 0;
    return ret;
11
12 }
13
   void run(int n, int k) {
14
     int ik = (1 << k) - 1;
15
16
     do {
17
     } while(ik = next_combination(n, ik));
18
19 }
```

日期时间公式

```
using namespace std;

int zeller(int y, int m, int d) {
   if(m <= 2) y--, m += 12; int c=y/100; y%=100;
   int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
   if(w<0) w+=7; return (w);
}</pre>
```

IO

Cpp 快速读入

```
#include <bits/stdc++.h>
   using namespace std;
   struct FastIO {
     static const int S = 65536;
     char buf[S]:
     int pos, len;
     bool eof;
10
     FILE *in:
11
     FastIO(FILE *_in = stdin) {
12
       in = _in;
13
       pos = len = 0;
14
       eof = false;
15
16
     int nextChar() {
17
       if (pos == len)
18
         pos = 0, len = fread(buf, 1, S, in);
19
       if (pos == len) {eof = true; return -1;}
20
       return buf[pos++];
21
22
     int nextUInt() {
23
       int c = nextChar(), x = 0;
24
       while (c <= 32) c = nextChar();</pre>
       for (;'0' \le c \&\& c \le '10'; c = nextChar()) x = x * 10 + c - '0';
26
       return x;
27
28
     int nextInt() {
29
       int s = 1, c = nextChar(), x = 0;
30
       while (c <= 32) c = nextChar();</pre>
31
       if (c == '-') s = -1, c = nextChar();
       for (; '0' <= c && c <= '9'; c = nextChar()) x = x * 10 + c - '0';
32
33
       return x * s:
34
35
     void nextString(char *s) {
36
       int c = nextChar():
37
       while (c <= 32) c = nextChar();</pre>
38
       for(; c > 32; c = nextChar()) *s++ = c;
39
       *s = 0:
40
41 };
```

Java 模板

```
import java.io.BufferedReader;
import java.io.InputStream;
```

```
3 import java.io.InputStreamReader;
   import java.io.IOException;
   import java.io.PrintWriter;
   import java.util.Arrays;
   import java.util.LinkedList;
   import java.util.Queue;
   import java.util.StringTokenizer;
10
   public class Main {
11
12
     public static void main(String[] args) {
13
       new Main().run();
14
     }
15
     public void run() {
16
17
       InputReader reader = new InputReader(System.in);
18
       PrintWriter writer = new PrintWriter(System.out);
19
       try {
20
       } catch (Exception e) {
21
       } finally {
22
         writer.close();
23
       }
24
25
     class InputReader {
26
27
       InputReader(InputStream in) {
28
         this.reader = new BufferedReader(new InputStreamReader(in));
29
         this.tokenizer = new StringTokenizer("");
       }
30
31
32
       public String nextToken() throws IOException {
33
         while (!tokenizer.hasMoreTokens()) {
34
           tokenizer = new StringTokenizer(reader.readLine());
         }
35
         return tokenizer.nextToken();
36
37
       }
38
39
       public int nextInt() throws IOException {
         return Integer.parseInt(nextToken());
40
       }
41
42
43
       private BufferedReader reader;
44
       private StringTokenizer tokenizer;
45
46
     private void debug(Object...o) {
47
```

```
System.err.println(Arrays.deepToString(o));

49 }
50 }
```

数据结构

BIT

```
* 支持第 k 大的 BIT
    * Kyb
    */
   inline int LB(int x) {return x & (-x):}
   const int MXN = 1e5;
   template <typename T>
   struct BIT {
     T = [MXN+5];
10
     int n:
     void init(int m) {
11
12
       n = m + 5;
13
       for (int i = 0; i \le n; i++)_[i] = 0;
14
15
     T query(int w) {
16
      T ret = 0;
17
       for (w += 3; w > 0; w -= LB(w))ret += [w];
18
       return ret;
19
     }
20
     void update(int w, T d) {
21
       for (w += 3; w < n; w += LB(w))_[w] += d;
22
23
     /*
      * 待验证
24
25
26
     int find_Kth(int k) {
                                  // UESTC_Dagon
27
       int ans=0,cnt=0;
28
       for(int i=22;i>=0;i--){
29
         ans+=1<<i;
         if(ans>=n||cnt+_[ans]>=k)ans-=1<<i;
30
31
         else cnt+=_[ans];
32
33
       return ans - 2;
34
35 };
```

LCA

```
#include <bits/stdc++.h>
const int POW = 18;
const int N = 1e5;
/*
```

```
* p[i][j]: i的第i倍祖先
   * d[i]: i在树中的深度
   * edge[N]: 边集合
   * dfs(u, fa): 求出p,d
    * lca(a, b): 求出(a, b)的最近公共祖先
10
11 int p[N] [POW];
12 | int d[N];
13 | std::vector<int>edge[N];
   void dfs(int u, int fa){
     d[u] = d[fa] + 1;
15
     p[u][0] = fa;
16
     for(int i = 1; i < POW; i++) p[u][i] = p[p[u][i - 1]][i - 1];
17
18
     int sz = edge[u].size();
     for(int i = 0; i < sz; i++){
19
      int v = edge[u][i];
20
       if(v == fa) continue;
21
22
       dfs(v, u);
23
    }
24
25
   int lca(int a, int b) {
     if(d[a] > d[b]) std::swap(a, b);
27
     if(d[a] < d[b]) {</pre>
28
      int del = d[b] - d[a];
29
30
      for(int i = 0; i < POW; i++) if(del & (1 << i)) b = p[b][i];
31
32
     if(a != b) {
      for(int i = POW - 1; i >= 0; i--)
33
         if(p[a][i] != p[b][i])
34
35
           a = p[a][i], b = p[b][i];
36
       a = p[a][0], b = p[b][0];
37
    }
38
     return a;
39 }
```

RMQ

```
13
       for(int i = 0: i < n: i++)
14
         if(i + (1 << (j + 1)) <= n) {
15
           f[i][j + 1] = std::max(f[i][j], f[i + (1 << j)][j]);
16
         }
17
     mm[0] = -1;
     for(int i = 1; i < N; i++) {</pre>
18
       mm[i] = mm[i >> 1] + 1;
19
20
    }
21 }
22
   int cal(int 1, int r) {
23
24
     int k = mm[r - 1 + 1];
25
     return std::max(f[1][k], f[r - (1 << k) + 1][k]);
26
```

二维线段树

```
* 点修改区间查询
   */
   #include <bits/stdc++.h>
   const int N = 5000:
   int c[N][N];
   int lx[N], ly[N];
   int n, Q;
   struct Node
10
11
     struct node
12
13
      int a, b;
14
       int Min, Max;
15
     }f[N];
16
     int a, b;
17
18
     void update(int x)
19
20
       f[x].Min = std::min(f[x << 1].Min, f[x << 1 | 1].Min);
21
       f[x].Max = std::max(f[x << 1].Max, f[x << 1 | 1].Max);
22
     }
23
24
     void build(int x, int a, int b)
25
     {
26
       f[x].a = a; f[x].b = b;
27
       if(a < b)
28
       {
29
         int mid = (a + b) / 2;
30
         build(x << 1, a, mid);
         build(x << 1 | 1, mid + 1, b);
31
32
         update(x);
```

```
} else { f[x].Min = f[x].Max = 0: lv[a] = x: }
34
     }
35
     int queryMin(int x, int a, int b)
36
37
       if (a \le f[x].a \&\& f[x].b \le b) return f[x].Min;
38
39
       else
40
         int mid = (f[x].a + f[x].b) / 2;
41
         if(b <= mid) return queryMin(x << 1, a, b);</pre>
42
43
         else if(a > mid) return queryMin(x << 1 | 1, a, b);</pre>
         else return std::min(queryMin(x << 1, a, b), queryMin(x << 1 | 1, a, b));
44
45
46
     }
47
     int queryMax(int x, int a, int b)
48
49
50
       if(a <= f[x].a && f[x].b <= b) return f[x].Max;</pre>
51
       else
52
         int mid = (f[x].a + f[x].b) / 2;
53
54
         if(b <= mid) return queryMax(x << 1, a, b);</pre>
         else if(a > mid) return queryMax(x << 1 | 1, a, b);</pre>
55
          else return std::max(queryMax(x \leq 1, a, b), queryMax(x \leq 1 | 1, a, b));
56
57
       }
58
    }
59
   }f[N];
60
   void build(int x, int a, int b, int p, int q)
61
62 {
     f[x].a = a; f[x].b = b;
63
     f[x].build(1, p, q);
     if(a < b)
65
66
67
       int mid = (a + b) / 2;
68
       build(x \ll 1, a, mid, p, q);
69
       build(x << 1 | 1, mid + 1, b, p, q);
70
71
     else lx[a] = x;
72 }
73
74
   void update(int x, int y, int c)
75 | {
     x = lx[x]; y = ly[y];
76
77
     for(int X = x; X != 0; X >>= 1)
78
       for(int Y = y; Y != 0; Y >>= 1)
79
       {
80
         if(f[X].f[Y].a == f[X].f[Y].b)
81
82
           if(f[X].a == f[X].b)
```

```
83
             {
 84
               f[X].f[Y].Min = f[X].f[Y].Max = c;
 85
             }
 86
             else
 87
             ł
               f[X].f[Y].Min = std::min(f[X << 1].f[Y].Min, f[X << 1 | 1].f[Y].Min);
 88
               f[X].f[Y].Max = std::max(f[X << 1].f[Y].Max, f[X << 1 | 1].f[Y].Max);
 89
 90
            }
          }
 91
 92
          else { f[X].update(Y); }
 93
 94
 95
 96
    int queryMin(int x, int a, int b, int p, int q)
97
      if (a <= f[x].a && f[x].b <= b) return f[x].queryMin(1, p, q);
 99
      else
100
      {
101
             int mid = (f[x].a + f[x].b) / 2;
             if(b <= mid) return queryMin(x << 1, a, b, p, q);</pre>
102
             else if(a > mid) return queryMin(x << 1 | 1, a, b, p, q);
103
             else return std::min(queryMin(x << 1, a, b, p, q), queryMin(x << 1 | 1, a, b, p, q))
104
               \hookrightarrow q));
        }
105
106
107
    int queryMax(int x, int a, int b, int p, int q)
109
110
        if (a <= f[x].a && f[x].b <= b) return f[x].queryMax(1, p, q);
111
        else
        {
112
             int mid = (f[x].a + f[x].b) / 2:
113
114
             if(b <= mid) return queryMax(x << 1, a, b, p, q);</pre>
115
             else if(a > mid) return queryMax(x << 1 | 1, a, b, p, q);
116
             else return std::max(queryMax(x << 1, a, b, p, q), queryMax(x << 1 | 1, a, b, p,
               \hookrightarrow q));
117
        }
118 }
```

主席树

```
#include <iostream>
#include <cstring>
#include <cstdio>
#include <algorithm>
#include <set>
#include <map>
#include <vector>
#include <queue>
#include <queue>
#include <queue>
```

```
10 typedef std::pair<int, int> PII;
11 #define PB(x) push_back(x)
12 #define SZ size()
13 #define AA first
  #define BB second
  #define MP(x, y) make_pair(x, y)
15
   namespace ST {
17
                 线段树节点数组
18
           e:
                 左右节点指针
19
     * 1, r:
20
         sum:
                 区间和
                 主席树根节点
21
          rt:
                 总节点个数
22
         tot:
23
    struct E {
24
25
      int 1, r;
      int sum:
26
27
    }e[10000000];
28
     int rt[100010];
29
     int tot:
30
31
32
     * 建树 , 初始化调用 , 建立一颗空树
33
     void build(int &rt, int 1, int r) {
34
35
      rt = tot++;
36
      e[rt].sum = 0;
37
      if(1 == r) return;
38
      int mid = (1 + r) >> 1;
      build(e[rt].1, 1, mid);
39
40
      build(e[rt].r, mid + 1, r);
41
42
43
     * 插入新节点,向w位置插入节点,主席书中父亲结点为fa
44
45
46
     void update(int &rt, int 1, int r, int w, int fa) {
      rt = tot++:
47
      e[rt].1 = e[fa].1;
48
      e[rt].r = e[fa].r:
49
50
      e[rt].sum = e[fa].sum + 1;
51
      if(1 == r) return ;
      int mid = (1 + r) >> 1:
52
      if(w <= mid) update(e[rt].1, 1, mid, w, e[fa].1);</pre>
53
      else update(e[rt].r, mid + 1, r, w, e[fa].r);
54
55
56
57
     * 查询, 查询原图中(a,b)点对的信息, 其中c = lca(a,b)
     * 计算时候加入c信息
59
```

```
60
61
     int query(int a, int b, int c, int 1, int r, int k, int nd) {
62
       int mid = (1 + r) >> 1;
       int sum = e[e[a].1].sum + e[e[b].1].sum - e[e[c].1].sum * 2 + (int)(nd <= mid);
63
64
       if(1 == r) return 1;
65
       if(k <= sum) return query(e[a].1, e[b].1, e[c].1, 1, mid, k, nd);</pre>
66
       else return query(e[a].r, e[b].r, e[c].r, mid + 1, r, k - sum, (nd <= mid ? 1e9 :
67
    }
68
   };
```

树链剖分

```
* siz[v]表示以v为根的子树的节点数
   * dep[v]表示v的深度
   * top[v]表示v所在的重链的顶端节点
   * fa[v]表示v的父亲
   * son[v]表示与v在同一重链上的v的儿子节点
   * w[v]表示v与其父亲节点的连边在线段树中的位置
   * 初始需要调用cnt1 = cnt2 = cnt3 = 0; dfs1(ROOT, 0); dfs2(ROOT, 1); bt(1, cnt2);
   * 模板为边带权值,点带权值需要修改query(x,y)
   * update(x, p, c)的p为线段树中的编号,更新x需要调用w[x]
11
   */
  const int N = 1e5:
13 const int M = 2 * N:
14 typedef long long 11;
15 #define MID(x, y) (((x) + (y)) >> 1)
  #include <bits/stdc++.h>
17
18 int fa[N], top[N], w[N], son[N], dep[N], sz[N], r[N];
19 int a[N]. b[N]:
20 ll c[N];
21 int ind[N]:
  int t[M], nt[M];
  int cnt1, cnt2, cnt3;
24 int n. m:
25
26
  struct node
27
    int 1, r;
29
    int a, b;
30
    11 sum:
31 }f[M];
32
  int rt;
34
  void dfs1(int x, int d)
35 {
    dep[x] = d;
    son[x] = 0;
```

```
sz[x] = 1:
39
     for(int k = ind[x]; k != -1; k = nt[k])
40
       if(t[k] != fa[x])
41
       {
42
         fa[t[k]] = x;
         dfs1(t[k], d + 1);
43
         sz[x] += sz[t[k]];
44
45
         if(sz[t[k]] > sz[son[x]]) son[x] = t[k];
       }
46
47
   }
48
   void dfs2(int x, int tt)
49
50
51
     w[x] = ++cnt2;
     top[x] = tt;
     if(son[x]) dfs2(son[x], tt);
53
     for(int k = ind[x]; k != -1; k = nt[k]) if(t[k] != fa[x] && t[k] != son[x])
54
55
       dfs2(t[k], t[k]);
56
   }
57
   void add(int a, int b)
58
59
     t[cnt1] = b;
60
     nt[cnt1] = ind[a];
61
     ind[a] = cnt1++;
62
63
64
   void update(int x)
65
66
     f[x].sum = f[f[x].1].sum + f[f[x].r].sum;
67
68
69
   int bt(int a, int b)
70
71
72
     int x = cnt3++;
     f[x].a = a; f[x].b = b;
73
74
     if(a < b)
75
     int mid = MID(a, b);
76
77
      f[x].l = bt(a, mid);
       f[x].r = bt(mid + 1, b);
78
79
      f[x].sum = 0;
80
81
     else
82
83
      f[x].sum = 0;
84
85
     return x;
86 }
87
```

```
88 // 在线段树上查询,不要直接调用
 89 11 query(int x, int a, int b)
90
      if(a \leq f[x].a && f[x].b \leq b) return f[x].sum:
92
      int mid = MID(f[x].a, f[x].b);
      11 \text{ ans} = 0:
93
      if (a <= mid) ans += query(f[x].1, a, b);
      if(b > mid) ans += query(f[x].r, a, b);
96
      return ans;
 97
 98
    // 单调修改
99
    void update(int x, int p, int cc)
101
102
        if(f[x].a == f[x].b) \{ f[x].sum = cc; return; \}
103
        int mid = MID(f[x].a, f[x].b);
104
        if(p <= mid) update(f[x].1, p, cc);</pre>
105
        else update(f[x].r, p, cc);
106
        update(x);
107
108
    // 树上查询
110 ll query(int x, int y)
111 {
112
        int fx = top[x], fy = top[y];
113
        11 sum = 0;
114
        while(fx != fy)
115
        {
            if(dep[fx] < dep[fy])</pre>
116
117
118
                std::swap(x, y);
119
                std::swap(fx, fy);
120
121
            sum += query(rt, w[fx], w[x]);
122
            x = fa[top[x]];
123
            fx = top[x];
124
125
        if(dep[x] > dep[y]) std::swap(x, y);
126
        if(x == y) return sum;
127
        return sum + query(rt, w[son[x]], w[y]);
128 }
```

树分治

```
/*
* 树分治
3 */
4
5 // 点分治
6 #include <bits/stdc++.h>
```

```
7 | const int N = 100000 + 5:
8 std::vector<int> edges[N];
9 int n;
10 bool vis[N]:
11 int parent[N];
12 int sz[N];
13 int que[N];
   int balance[N];
15
   int bfs(int source,int fa = -1) {
16
     int qf = 0, qe = 0;
17
     que[qe++] = source;
18
     parent[source] = fa;
19
20
     while (qf != qe) {
       int u = que[qf++];
21
       sz[u] = 1;
22
23
       balance[u] = 0;
       for (int v : edges[u]) {
24
         if (!vis[v] && parent[u] != v) {
25
26
           parent[v] = u;
27
           que[qe++] = v;
28
29
       }
30
     for (int i = qe - 1; i > 0; -- i) {
31
32
       int u = que[i];
       sz[parent[u]] += sz[u];
33
       balance[parent[u]] = std::max(balance[parent[u]],sz[u]);
34
35
36
     return qe;
37 }
38
39
   void divide(int root) {
     int tot = bfs(root);
40
     for (int i = tot - 1; i > 0; -- i) {
41
       int u = que[i];
42
       balance[u] = std::max(balance[u],tot - sz[u]);
43
       if (balance[u] < balance[root]) {</pre>
44
45
         root = u;
       }
46
47
     }
48
     bfs(root);
49
50
     // balabalah
51
     vis[root] = true:
52
53
     for (int u : edges[root]) {
54
       if (!vis[u]) {
55
         divide(u);
56
       }
```

```
57 }
58 }
```

Treap

```
* Poj 3481
   #include <bits/stdc++.h>
   using namespace std;
   #define AA first
   #define BB second
   #define MP make_pair
   #define PII pair<int, int>
10
11
   const int N = 1000000+1000, MOD = 1e9+7;
12 const int nil = 0:
   struct node {
13
       int ch[2], key, sz;
14
       PII data:
15
16 } f[N];
17
   int cnt, rt;
18
   void up(int cur) {
       f[cur].sz = f[f[cur].ch[0]].sz + f[f[cur].ch[1]].sz + 1:
20
21 }
22
23
   int newNode(PII data) {
24
       cnt++:
25
       f[cnt].ch[0] = f[cnt].ch[1] = 0;
26
       f[cnt].data = data:
27
       f[cnt].sz = 1:
28
       f[cnt].key = rand();
29
       return cnt:
30
31
   PII split(int p, int n) {
33
       if(n == 0) return MP(nil, p);
34
       int sz = f[f[p].ch[0]].sz;
35
       if(n == sz) {
36
           int x = f[p].ch[0];
37
           f[p].ch[0] = nil;
38
           up(p);
39
           return MP(x, p);
40
       }
41
       else if(n == sz + 1) {
42
           int x = f[p].ch[1];
43
           f[p].ch[1] = nil;
44
           up(p);
45
           return MP(p, x);
```

```
}
46
47
       else if(n < sz) {</pre>
48
           PII res = split(f[p].ch[0], n);
49
           f[p].ch[0] = res.BB;
50
           up(p);
51
           return MP(res.AA, p);
       }
52
53
       else {
           PII res = split(f[p].ch[1], n - sz - 1);
54
55
           f[p].ch[1] = res.AA;
56
           up(p);
57
           return MP(p, res.BB);
58
59
   }
60
   int merge(int p, int q) {
61
       if(p == nil) return q;
62
63
       if(q == nil) return p;
64
       if(f[p].key > f[q].key) {
           f[q].ch[0] = merge(p, f[q].ch[0]);
65
66
           up(q);
67
           return q;
       }
68
69
       else {
70
           f[p].ch[1] = merge(f[p].ch[1], q);
71
           up(p);
72
           return p;
73
       }
74
75
   int getrank(int p, int w) {
76
77
       int ans = 0:
78
       while(p != nil) {
           if(w == f[p].data.BB) return ans + f[f[p].ch[0]].sz;
79
80
           if(w < f[p].data.BB) p = f[p].ch[0];
81
           else {
82
                ans += f[f[p].ch[0]].sz + 1;
83
                p = f[p].ch[1];
           }
84
       }
85
86
       return ans;
87
```

Splay

```
1 // 注意初始化内存池和 null 节点
2 /*
3 * hdu3487 hdu1908
4 */
5 #include <algorithm>
```

```
6 const int MAX NODE = 1e5:
   struct Node{
     int rev,size; Node *ch[2],*p;
     void set(Node*,int); int dir(); void update(); void relax(); void appRev();
10 } nodePool[MAX_NODE],*curNode,*null;
11 Node *newNode(){
     Node *t=curNode++; t->rev=0, t->size=1;
12
13
     t->ch[0]=t->ch[1]=t->p=null; return t;
14 }
15
   struct Splay{
16
     Node *root;
     Splay(){ root=newNode(); root->set(newNode(),0); root->update(); }
17
18
     void rot(Node *t){
19
       Node *p=t->p; int d=t->dir();
       p->relax(); t->relax();
       if(p==root) root=t;
21
22
       p->set(t->ch[!d],d); p->p->set(t,p->dir()); t->set(p,!d);
23
       p->update();
24
     }
25
     void splay(Node *t,Node *f=null){
26
       for(t->relax();t->p!=f;)
27
         if(t->p->p==f) rot(t);
28
         else t->dir()==t->p->dir()?(rot(t->p),rot(t)):(rot(t),rot(t));
29
       t->update();
30
    }
31 };
   void initNull(){ curNode=nodePool;null=curNode++;null->size=0; }
   void Node::set(Node *t,int _d){ ch[_d]=t; t->p=this; }
34 int Node::dir(){ return this==p->ch[1]; }
   void Node::update(){ size=ch[0]->size+ch[1]->size+1;}
   void Node::relax(){ if(rev) ch[0]->appRev(), ch[1]->appRev(), rev=false; }
37 void Node::appRev(){ if(this==null) return; rev^=true; std::swap(ch[0],ch[1]); }
```

KD-Tree

```
1 /*
2 * KDT 2016 Qingdao 11
3 * 估价函数:
4 * 欧几里德距离下界: sqr(max(Max(X-x.Max[0],x.Min[0]-X,0)))+ ...
5 * 欧几里德距离上界: max(sqr(X-x.Max[0]),sqr(x.Min[0]-X))+ ...
6 * 曼哈顿距离下界: max(X-x.Max[0],0)+max(x.Min[0]-X,0)+ ...
7 * 曼哈顿距离上界: max(abs(X-x.Max[0]),abs(x.Min[0]-X))+ ...
8 *
9 */
10
11 #include <bits/stdc++.h>
12 using namespace std;
13 typedef long long LL;
14 #define cmin(x, y) x = min(x, y)
```

```
15 #define cmax(x, v) x = max(x, v)
16
17
   const int N = 1000000;
18 struct point {
     LL _[2], op, p;
19
     LL& operator [] (int x) { return _[x]; }
20
     int operator < (const point &t) const {</pre>
21
22
       return p < t.p;</pre>
23
   } a[N], ans[N];
25
   template <typename T> inline T SQ(T x) { return x * x; }
26
   inline LL dis(point a, point b) {
27
28
     return SQ(1LL * (a[0] - b[0])) + SQ(1LL * (a[1] - b[1]));
29 }
30
   struct node {
31
     int 1, r, fa;
     bool vis, has:
32
33
     point p;
34 } b[N];
   int n, m, cnt;
36
   inline int getnode() {
37
     assert(cnt < N);
38
     b[cnt].l = b[cnt].r = -1;
39
40
     b[cnt].vis = b[cnt].has = false;
41
     b[cnt].fa = -1;
42
     return cnt++;
43 }
44
   void sol(int u, point pt, point& ans, LL &value, int d) {
45
46
     if(u == -1 || !b[u].has) return;
     if(b[u].vis) {
47
48
       LL w1 = dis(b[u].p, pt);
49
       if(w1 < value || (w1 == value && b[u].p.op < ans.op)) {</pre>
50
         ans = b[u].p;
51
         value = w1:
52
       }
53
     int 1 = b[u].1, r = b[u].r;
54
     if(pt[d] > b[u].p[d]) swap(l, r);
56
     sol(1, pt, ans, value, d ^ 1);
     if(1LL * SQ(pt[d] - b[u].p[d]) <= value)</pre>
57
58
       sol(r, pt, ans, value, d ^ 1);
59
60
   void build(int &u, int 1, int r, int d, int fa=-1) {
61
62
63
     int mid = 1:
     if(1 != r) mid += rand() % (r - 1);
```

```
point p = a[mid]:
66
     swap(a[r], a[mid]);
     int j = 1;
     for(int i = 1; i < r; i++)
69
       if(a[i][d] < p[d]) {</pre>
         swap(a[i], a[j++]);
70
71
72
     swap(a[j], a[r]);
     u = getnode();
73
     b[u].fa = fa;
75
     b[u].p = a[j];
76
     if(1 < j) build(b[u].1, 1, j - 1, d ^ 1, u);</pre>
77
78
     if(j < r) build(b[u].r, j + 1, r, d^1, u);
79 }
80
   void add(int u, point p, int d) {
     if(u == -1) return;
83
     if(b[u].p[0] == p[0] \&\& b[u].p[1] == p[1]) {
       b[u].has = b[u].vis = true;
84
85
       return ;
86
     }
     if(p[d] > b[u].p[d]) add(b[u].r, p, d^1);
87
     else add(b[u].1, p, d ^ 1);
     b[u].has |= b[u].vis;
     if("b[u].1) b[u].has |= b[b[u].1].has;
     if("b[u].r) b[u].has |= b[b[u].r].has;
91
92
```

KD-Tree 欧几里得距离

```
Problem: 1941
    User: Ceva
    Language: C++
    Result: Accepted
    Time: 1476 ms
    Memory:63792 kb
  9
  * 求所有点 距离最大和最小值差值的最小值
10
11
12
  #include <iostream>
14 #include <cstdio>
  #include <algorithm>
16 #include <cmath>
17 #include <cstring>
18 #include <string>
19 using namespace std;
```

```
20 #define cmin(x, v) x = min(x, v)
                                                                                                     70
21
   #define cmax(x, y) x = max(x, y)
                                                                                                     71
22
                                                                                                     72
   const int maxn = 1000010:
23
                                                                                                     73
   const int inf = 1e9;
                                                                                                     74
                                                                                                             rt = ++cnt;
25 struct point {
                                                                                                     75
                                                                                                             cur = mk:
       int x[2]:
                                                                                                     76
26
   } p[maxn];
27
                                                                                                     77
   struct node {
                                                                                                             f[rt].p = mid;
28
                                                                                                     78
29
                                                                                                     79
       int 1, r, p;
30
       int mn[2], mx[2];
                                                                                                     80
31 } f[maxn * 2];
                                                                                                     81
                                                                                                     82
                                                                                                             update(rt);
   int n:
33
   int cur;
                                                                                                     83
   int cnt;
                                                                                                     84
35
   int mm, nn;
                                                                                                     85
   int operator < (const point &a, const point &b) { return a.x[cur] < b.x[cur]; }</pre>
                                                                                                     86
                                                                                                             int ans = 0:
36
37
                                                                                                     87
38
   void update(int rt) {
                                                                                                     88
                                                                                                             return ans;
       int id = f[rt].p;
                                                                                                     89
39
       int 1 = f[rt].1, r = f[rt].r;
                                                                                                     90
40
       for(int i = 0; i < 2; i++) {
41
42
           f[rt].mn[i] = f[rt].mx[i] = p[id].x[i];
                                                                                                     92
43
           if(1) cmin(f[rt].mn[i], f[1].mn[i]), cmax(f[rt].mx[i], f[1].mx[i]);
                                                                                                     93
44
           if(r) cmin(f[rt].mn[i], f[r].mn[i]), cmax(f[rt].mx[i], f[r].mx[i]);
                                                                                                     94
45
       }
                                                                                                     95
46
                                                                                                     96
47
                                                                                                     97
                                                                                                             if(a < b) {
   int mn(int rt, point P) {
                                                                                                     98
48
49
       int tot = 0;
                                                                                                     99
                                                                                                            }
50
       for(int i = 0; i < 2; i++) {
                                                                                                    100
           int a = 0:
51
                                                                                                    101
                                                                                                             else {
52
           if(P.x[i] < f[rt].mn[i]) tot += f[rt].mn[i] - P.x[i];
                                                                                                    102
           else if(P.x[i] > f[rt].mx[i]) tot += -f[rt].mx[i] + P.x[i];
53
                                                                                                    103
54
           tot += a:
                                                                                                    104
                                                                                                            }
       }
55
                                                                                                    105 }
56
       return tot:
                                                                                                    106
57
                                                                                                    107
58
                                                                                                    108
   int mx(int rt, point P) {
59
                                                                                                    109
60
       int tot = 0;
                                                                                                    110
61
       for(int i = 0; i < 2; i++) {
                                                                                                    111
           int a = -inf:
62
                                                                                                    112
                                                                                                             if(a > b) {
           cmax(a, abs(P.x[i] - f[rt].mn[i]));
63
                                                                                                    113
           cmax(a, abs(P.x[i] - f[rt].mx[i]));
64
                                                                                                    114
                                                                                                            }
65
           tot += a:
                                                                                                    115
                                                                                                             else {
66
       }
                                                                                                    116
67
                                                                                                    117
       return tot;
68 }
                                                                                                    118
69
                                                                                                    119
                                                                                                            }
```

```
void build(int &rt, int 1, int r, int mk) {
    if(1 > r) return;
    int mid = (1 + r) >> 1:
    nth_element(p + 1, p + mid, p + r + 1);
    //cout << p[rt].x[0] << " " << p[rt].x[1] << " " << l << " " << r << endl;
    f[rt].1 = f[rt].r = 0;
    build(f[rt].1, 1, mid - 1, mk ^ 1);
    build(f[rt].r, mid + 1, r, mk ^ 1);
int getdis(point &x, point &y) {
    for(int i = 0; i < 2; i++) ans += abs(x.x[i] - y.x[i]);
void dfs1(int rt, point &P, int cur, int &m) {
    int 1 = f[rt].1, r = f[rt].r;
    point &t = p[f[rt].p];
    int tot = getdis(t, P);
    if(tot != 0) cmin(m, tot);
    int a = mn(1, P), b = mn(r, P);
        if(1) if(m > a) dfs1(f[rt].1, P, cur ^ 1, m);
        if(r) if(m > b) dfs1(f[rt].r, P, cur ^ 1, m);
        if(r) if(m > b) dfs1(f[rt].r, P, cur ^ 1, m);
        if(1) if(m > a) dfs1(f[rt].1, P, cur ^ 1, m);
void dfs2(int rt, point &P, int cur, int &m) {
    int 1 = f[rt].1, r = f[rt].r;
    point &t = p[f[rt].p];
    cmax(m, getdis(t, P));
    int a = mx(1, P), b = mx(r, P);
        if(1) if(m < a) dfs2(f[rt].1, P, cur ^ 1, m);</pre>
        if(r) if(m < b) dfs2(f[rt].r, P, cur ^ 1, m);</pre>
        if(r) if(m < b) dfs2(f[rt].r, P, cur ^ 1, m);</pre>
        if(1) if(m < a) dfs2(f[rt].1, P, cur ^ 1, m);</pre>
```

```
120 }
121
122 int main() {
123
        scanf("%d", &n);
124
        for(int i = 0; i < n; i++) scanf("%d%d", &p[i].x[0], &p[i].x[1]);
125
        build(root, 0, n - 1, 0);
126
127
        int ans = 1e9;
        cur = 0:
128
129
        cnt = 0;
130
        for(int i = 0; i < n; i++) {
131
            int mx = 0, mn = inf;
            dfs2(root, p[i], 0, mx); dfs1(root, p[i], 0, mn);
132
133
            cmin(ans, mx - mn);
        }
134
        printf("%d\n", ans);
135
136
        return 0;
137 }
```

DLX 精确覆盖

```
1 #include <iostream>
   #include <cstdio>
   #include <cstring>
  const int Maxm = 1000:
   const int Maxn = 1200;
  int ans[Maxn];
8 struct DLX{
     struct Node{
10
      Node *L, *R, *U, *D;
      int col. row:
11
12
    } *head, *row[Maxn], *col[Maxm], node[Maxn * Maxm];
     int colsum[Maxm], cnt;
13
14
15
     /* dancing link
     * 精确覆盖问题
16
17
     * 可以添加迭代加深优化
     * 1) 举深度 h
18
     * 2) 若当前深度 + predeep > h : return false□
19
     * 3) mat 下标 1 开始
20
21
     */
22
     /*
23
        int predeep() {
24
       bool vis[Maxm];
25
        memset(vis. 0. sizeof(vis)):
26
        int ret = 0:
27
        for (Node *p = head->R; p != head; p = p->R)
        if (!vis[p->col]) {
28
29
        ret ++ ;
```

```
30
          vis[p->col] ++ :
31
          for (Node *q = p->D; q != p; q = p->D)
32
          for (Node *r = q->R; r != q; r = r->R)
33
          vis[r->col] = true:
34
35
         return ret:
36
         }
37
38
      void init(int mat[][Maxm], int n, int m) {
39
        cnt = 0:
40
        memset(colsum, 0, sizeof (colsum));
41
        head = &node[cnt ++ ];
        for(int i = 1: i <= n: i ++ )
42
43
          row[i] = &node[cnt ++ ];
44
        for(int j = 1; j <= m; j ++ )
45
           col[j] = &node[cnt ++ ];
46
        head \rightarrow D = row[1], row[1] \rightarrow U = head;
47
        head \rightarrow R = col[1], col[1] \rightarrow L = head;
48
        head -> U = row[n], row[n] -> D = head;
        head->L = col[m], col[m]->R = head;
49
50
        head - row = head - col = 0;
51
        for(int i = 1; i <= n; i ++ ) {
52
           if (i != n) row[i]->D = row[i + 1]:
53
           if (i != 1) row[i]->U = row[i - 1];
54
           row[i] \rightarrow L = row[i] \rightarrow R = row[i];
55
           row[i]->row = i, row[i]->col = 0;
56
57
        for(int i = 1; i <= m; i ++ ) {
58
           if (i != m) col[i]->R = col[i + 1];
           if (i != 1) col[i]->L = col[i - 1];
59
60
           col[i] \rightarrow U = col[i] \rightarrow D = col[i];
           col[i] \rightarrow col = i, col[i] \rightarrow row = 0;
61
62
        for(int i = n; i > 0; i -- )
63
64
           for(int j = m; j > 0; j -- )
65
             if(mat[i][i]) {
                Node *p = &node[cnt ++ ];
66
               p \rightarrow R = row[i] \rightarrow R, row[i] \rightarrow R \rightarrow L = p;
67
                p->L = row[i], row[i]->R = p;
                p->D = col[j]->D, col[j]->D->U = p;
69
70
               p->U = col[j], col[j]->D = p;
71
                p->row = i;
72
                p->col = j;
73
                colsum[j] ++ ;
74
             }
75
76
      void remove(Node *c) {
77
        c\rightarrow L\rightarrow R = c\rightarrow R;
78
        c \rightarrow R \rightarrow L = c \rightarrow L:
79
        for(Node *p = c->D; p != c; p = p->D) {
```

```
for(Node *q = p->R; q != p; q = q->R) {
80
81
             q->U->D = q->D;
82
             q \rightarrow D \rightarrow U = q \rightarrow U;
             colsum[q->col] -- ;
83
84
        }
85
86
87
       void resume(Node *c) {
        for(Node *p = c->U; p != c; p = p->U) {
88
           for(Node *q = p \rightarrow L; q != p; q = q \rightarrow L) {
89
90
             q \rightarrow U \rightarrow D = q;
             q \rightarrow D \rightarrow U = q;
91
92
             colsum[q->col] ++ ;
93
          }
        }
94
         col[c->col]->L->R = col[c->col];
95
         col[c->col]->R->L = col[c->col];
96
97
98
      int dfs(int deep) {
        if(head->R == head) return deep;
99
        Node *p, *q = head->R;
100
             for(p = head \rightarrow R; p != head; p = p \rightarrow R)
101
                 if(colsum[p->col] < colsum[q->col])
102
103
                      q = p;
104
             remove(q);
105
             for(p = q->D; p != q; p = p->D) {
106
                 for(Node* r = p \rightarrow R; r != p; r = r \rightarrow R)
                      if (r->col != 0)
107
                          remove (col[r->col]);
108
                 // ----- 可修改区域 ------
109
                  ans[deep] = p->row;
110
111
                 // -----
112
                 int sta = dfs (deep + 1);
113
                 if(sta) return sta;
114
                 for(Node* r = p->L; r != p; r = r->L)
                      if(r->col != 0)
115
                          resume (col[r->col]):
116
117
             }
             resume(q);
118
119
             return false;
120
        }
121 } dlx;
122 int mat[Maxn][Maxm];
123 int mem[Maxn][3]; // 记录每行代表哪一格填几
124 // col = 324
126 // 数独填充 (x, y) = v
    void addline(int x, int v, int v) {
128
        int i, j;
129
        n++;
```

```
mem[n][0] = x:
130
131
        mem[n][1] = y;
132
        mem[n][2] = v;
133
        for(i = 0; i < Maxm; i++) mat[n][i] = 0;</pre>
134
        mat[n][x * 9 + y + 1] = 1;
135
        mat[n][81 + x * 9 + v] = 1;
136
        mat[n][162 + y * 9 + v] = 1;
137
        mat[n][243 + (3 * (x / 3) + y / 3) * 9 + v] = 1;
138 }
```

DLX 多重覆盖

```
#include <bits/stdc++.h>
   #define cmin(x, y) x = std::min(x, y)
   const int Maxm = 62;
   const int Maxn = 62;
   int K;
   struct DLX{
     struct Node{
       Node *L, *R, *U, *D;
10
       int col, row;
11
    } *head. *row[Maxn]. *col[Maxm]. node[Maxn * Maxm];
12
     int colsum[Maxm], cnt;
     /* dancing link
13
     * 重复覆盖问题
14
15
     * 可以添加迭代加深优化
      * 1) 举深度 h
16
      * 2) 若当前深度 + predeep > h : return false□
17
18
      * 3) mat 下标 1 开始
19
      */
20
21
     int predeep(){
       bool vis[Maxm]:
22
23
       Node * p, *q, *r;
24
       memset(vis, 0, sizeof(vis));
       int ret = 0:
25
26
       for(p = head->R; p != head; p = p->R) {
27
         if(!vis[p->col]) {
28
           ret++:
29
           vis[p->col]++;
30
           for (q = p->D; q != p; q = q->D) {
             for(r = q - R; r != q; r = r - R) {
31
32
               vis[r->col] = true;
33
34
35
36
       }
37
       return ret:
38
```

```
void init(int mat[][Maxm]. int n. int m) {
39
40
        cnt = 0:
41
        int i, j;
42
        Node * p:
43
        memset(colsum, 0, sizeof(colsum));
        head = &node[cnt++]:
44
        for(i = 1; i <= n; i++) row[i] = &node[cnt++];</pre>
45
46
        for(j = 1; j <= m; j++) col[j] = &node[cnt++];</pre>
        head \rightarrow D = row[1], row[1] \rightarrow U = head;
47
48
        head \rightarrow R = col[1], col[1] \rightarrow L = head;
49
        head -> U = row[n], row[n] -> D = head;
50
        head \rightarrow L = col[m], col[m] \rightarrow R = head;
51
        head \rightarrow row = head \rightarrow col = 0:
52
        for(i = 1; i <= n; i++) {
53
          if(i != n) row[i]->D = row[i + 1];
54
          if(i != 1) row[i]->U = row[i - 1];
55
          row[i] -> L = row[i] -> R = row[i];
56
          row[i]->row = i; row[i]->col = 0;
57
        }
        for(i = 1; i <= m; i++) {
58
          if(i != m) col[i] -> R = col[i + 1];
59
60
          if(i != 1) col[i]->L = col[i - 1];
          col[i] -> U = col[i] -> D = col[i];
61
62
          col[i] \rightarrow col = i; col[i] \rightarrow row = 0;
63
        }
64
        for(i = n; i > 0; i--) {
          for(j = m; j > 0; j--) {
65
66
            if(mat[i][j]) {
67
               p = &node[cnt++];
               p->R = row[i]->R, row[i]->R->L = p;
68
69
               p->L = row[i], row[i]->R = p;
70
               p->D = col[j]->D, col[j]->D->U = p;
71
               p->U = col[j], col[j]->D = p;
72
               p->row = i;
73
               p->col = j;
74
               colsum[j]++;
75
76
          }
77
        }
78
79
      void remove(Node *c) {
80
       for(p = c->D; row[p->row] != row[c->row]; p = p->D) {
81
82
          p->R->L = p->L; p->L->R = p->R;
       }
83
84
85
      void resume(Node *c) {
86
        Node * p;
87
        for(p = c->U; row[p->row] != row[c->row]; p = p->U) {
88
          p->L->R = p->R->L = p;
```

```
89
        }
 90
 91
      bool dfs(int deep) {
 92
        if(head->R == head) {
 93
          if(deep <= K) return true;</pre>
 94
 95
          return false:
 96
        if(deep + predeep() > K) return false;
 97
 98
        Node *p, *q = head->R, *r;
 99
        for(p = head->R; p != head; p = p->R) {
          if(colsum[p->col] < colsum[q->col]) q = p;
100
            }
101
102
             for(p = q->D; p != q; p = p->D) {
                 remove(p);
103
                 for(r = p -> R; r != p; r = r -> R) {
104
105
                     if(r->col != 0) remove(r);
106
                 }
                 // ----- 可修改区域 ------
107
                          ans[deep] = p->row;
108
109
110
                 int sta = 0;
111
                 sta = dfs(deep + 1);
112
                 if(sta) return sta;
113
                 for(r = p->L; r != p; r = r->L) {
114
                     if(r\rightarrow col != 0) resume(r):
                 }
115
116
                 resume(p);
117
118
             return false;
119
        }
120 } dlx;
```

图论 **2-SAT**

```
#include <iostream>
#include <cstdio>
#include <iostream>
#include <iostream>
#include <stack>
#include <cstring>
#include <algorithm>
#include <cstring>

const int N = 500;
const int M = 2e6;

/*

* dfn: 标号数组
* low: 回向边数组
* /*
```

```
15 int ind[N]:
16 | int to[M], nt[M];
   int dfn[N], low[N];
18 int color[N]:
   bool vis[N];
19
   int cnt, num, idn;
20
   int ncnt:
21
   std::stack<int>s;
22
23
24
   void add(int i, int j) {
25
     cnt++;
26
     to[cnt] = j;
     nt[cnt] = ind[i];
27
28
     ind[i] = cnt;
29 }
30
31
   void tarjan(int x) {
32
     if(dfn[x] != 0) return;
33
     low[x] = dfn[x] = ++num;
     vis[x] = true:
34
     s.push(x);
35
     for(int k = ind[x]; k != -1; k = nt[k]) {
36
       tarjan(to[k]);
37
       if(vis[to[k]]) low[x] = std::min(low[x], low[to[k]]);
38
39
40
     if(dfn[x] == low[x]) {
41
       ncnt++;
42
       while(true) {
43
         int p = s.top();
44
         s.pop();
45
         color[p] = ncnt;
46
         vis[p] = false;
47
         if(p == x) break;
48
       }
49
    }
50
51
   const int Maxn = 500;
   const int Maxm = 50000;
53
54
   int n. m:
   int a[Maxm], b[Maxm], c[Maxm];
56
57
   bool check() {
58
     memset(dfn, 0, sizeof dfn);
59
     memset(low, 0, sizeof low);
60
     memset(vis, 0, sizeof vis);
     for(int i = 0; i < 2 * n; i++) ind[i] = -1;
61
62
     int tot = n;
63
     num = idn = 0:
     cnt = 0;
```

```
65
    ncnt = 0:
66
    /*
67
     * 这里建图
     * 每个节点分为两个, 选和不选
68
     * (p->q), 表示选p一定要选q
69
     * 检查每个节点洗和不洗是否在同一个集合
70
71
72
    for(int i = 0; i < 2 * n; i++) tarjan(i);
73
    for(int i = 0; i < n; i++) {</pre>
74
      if(color[i] == color[tot + i]) return false;
75
    }
76
    return true;
77
```

KM

```
* 二分图最大权匹配
    * SJTU
    * Hdu 2255
    */
   #include <bits/stdc++.h>
   using namespace std;
   namespace graph {
10
     const int maxn=400; const int oo=0x7ffffffff;
11
     int w[maxn] (maxn], x[maxn], y[maxn], px[maxn], sy[maxn], slack[maxn];
     int par[maxn]; int n; int pa[200][2],pb[200][2],n0,m0,na,nb; char s[200][200];
12
     void adjust(int v){     sy[v]=py[v]; if (px[sy[v]]!=-2) adjust(px[sy[v]]);}
13
14
     bool find(int v){for (int i=0;i<n;i++)</pre>
15
       if (py[i]==-1){
         if (slack[i]>x[v]+y[i]-w[v][i]) slack[i]=x[v]+y[i]-w[v][i], par[i]=v;
16
17
         if (x[v]+y[i]==w[v][i]){
           py[i]=v; if (sy[i]==-1){adjust(i);    return 1;}
18
19
           if (px[sy[i]]!=-1) continue; px[sy[i]]=i;
20
           if (find(sy[i])) return 1;
         }}return 0:}
21
22
     int km(){int i,j,m,flag; for (i=0;i<n;i++) sy[i]=-1,y[i]=0;</pre>
23
       for (i=0;i< n;i++) {x[i]=0; for (j=0;j< n;j++) x[i]=max(x[i],w[i][j]);}
       for (i=0:i<n:i++){</pre>
24
25
         for (j=0;j<n;j++) px[j]=py[j]=-1,slack[j]=oo;</pre>
26
         px[i]=-2; if (find(i)) continue; flag=false;
27
         for (;!flag;){ m=oo;
28
           for (j=0;j<n;j++) if (py[j]==-1) m=min(m,slack[j]);</pre>
29
           for (j=0; j< n; j++){ if (px[j]!=-1) x[j]-=m;
30
             if (py[j]!=-1) y[j]+=m; else slack[j]-=m;}
31
           for (j=0; j< n; j++){ if (py[j]==-1\&\&!slack[j]){
32
             py[j]=par[j];
             if (sy[j]==-1){ adjust(j); flag=true; break;}
33
34
             px[sy[j]]=j; if (find(sy[j])){flag=true;break;}
```

```
35 | }}}}

36 | int ans=0; for (i=0;i<n;i++) ans+=w[sy[i]][i];return ans;}

37 |
38 | using namespace graph;
```

ISAP

```
* 用于边数较多的情况,调用ISAP(intst, inted, intn), n要稍大
   */
   #include <bits/stdc++.h>
   const int Maxn = 1000;
   const int Maxm = Maxn * Maxn:
   struct node {
9
     int u, v, c, next;
10 | Pe[Maxm]:
11 int tot, last[Maxn];
   void adde(int u, int v, int c, int c1) {
12
     e[tot].u = u; e[tot].v = v; e[tot].c = c; e[tot].next = last[u]; last[u] = tot++;
14
     e[tot].u = v; e[tot].v = u; e[tot].c = c1; e[tot].next = last[v]; last[v] = tot++;
15 }
16
   int dist[Maxn], cur[Maxn], gap[Maxn], pre[Maxn];
18 int ISAP(int s. int t. int n) {
19
     int i, j, u, v, det;
    int maxflow = 0;
20
     memset(dist, 0, sizeof(dist[0]) * (n + 3));
21
     memset(gap, 0, sizeof(gap[0]) * (n + 3));
22
23
    for (i = 0; i < n; i ++ )
24
      cur[i] = last[i]:
25
     u = s:
26
     gap[0] = n;
     pre[s] = -1:
27
28
     while (dist[s] <= n) {</pre>
29
       bool flag = false;
      for (j = cur[u]; j != -1; j = e[j].next) {
30
31
         v = e[i].v;
32
         if (e[j].c > 0 && dist[u] == dist[v] + 1) {
          flag = true:
33
34
           pre[v] = u;
35
           cur[u] = j;
36
           u = v:
37
           break;
38
39
       }
40
       if (flag) {
         if (u == t) {
41
42
           int det = INF:
43
           for (i = u; i != s; i = pre[i])
```

```
det = min(det, e[cur[pre[i]]].c);
44
45
            for (i = u; i != s; i = pre[i]) {
46
              e[cur[pre[i]]].c -= det;
47
              e[cur[pre[i]] ^ 1].c += det;
48
49
            maxflow += det;
50
            u = s:
51
       }
52
53
       else {
54
         int mind = n;
55
         for (j = last[u]; j != -1; j = e[j].next ) {
56
           v = e[i].v;
57
           if (e[j].c > 0 && dist[v] < mind) {</pre>
              mind = dist[v];
58
59
              cur[u] = j;
60
           }
61
         }
62
         if (( -- gap[dist[u]]) == 0) break;
         gap[dist[u] = mind + 1] ++ ;
63
         if (u != s) u = pre[u];
64
65
66
67
     return maxflow;
68
```

SAP

```
#include <bits/stdc++.h>
  #define cmin(x, y) x = std::min(x, y)
 3 typedef int ft:
   const ft inf = 0x3f3f3f;
   const int M = 500000+5.N = 20000+5:
   struct SAP{
     int y[M],nxt[M],gap[N],fst[N],c[N],pre[N],q[N],dis[N];
     ft f[M];
10
     int S,T,tot,Tn;
     void init(int s,int t,int tn){
11
12
13
       memset(fst,0,sizeof (int) * tn);
14
       memset(dis, 0, sizeof(int) * tn);
15
       S=s;T=t;Tn=tn;
16
     void add(int u.int v.ft c1.ft c2=0){
18
       tot++;y[tot]=v;f[tot]=c1;nxt[tot]=fst[u];fst[u]=tot;
19
       tot++;y[tot]=u;f[tot]=c2;nxt[tot]=fst[v];fst[v]=tot;
20
21
     ft sap(){
```

```
int u=S.t=1:ft flow=0:
22
23
       for(int i = 0; i < t; i++){
24
         int u=q[i];
25
         for(int j=fst[u];j;j=nxt[j])
           if(dis[y[j]]>dis[u]+1&&f[j^1])
26
27
              q[t++]=y[j],dis[y[j]]=dis[u]+1;
28
       }
29
       for(int i = 0; i < Tn; i++)gap[dis[i]]++;</pre>
30
       while(dis[S]<=Tn){</pre>
31
         while(c[u] \&\& (!f[c[u]] | |dis[y[c[u]]] + 1! = dis[u]))
32
           c[u]=nxt[c[u]];
33
         if(c[u]){
           pre[y[c[u]]]=c[u]^1;
34
35
           u=y[c[u]];
36
           if(u==T){
37
              ft minf=inf;
38
              for(int p=pre[T];p;p=pre[y[p]])
39
                cmin(minf,f[p^1]);
              for(int p=pre[T];p;p=pre[y[p]])
40
41
               f[p^1]-=minf,f[p]+=minf;
42
              flow+=minf;u=S;
43
           }
44
         }else {
45
           if(!(--gap[dis[u]]))break;
46
           int mind=Tn;
47
           c[u]=fst[u];
           for(int j=fst[u];j;j=nxt[j])
48
49
             if(f[j]&&dis[y[j]]<mind)</pre>
50
                mind=dis[y[j]],c[u]=j;
            dis[u]=mind+1;
51
52
            gap[dis[u]]++;
            if(u!=S)u=y[pre[u]];
53
54
         }
55
       }
56
       return flow;
57
    }
58 };
```

Dinic

```
11
       int q[N];
12
       bool vis[N];
13
       int h[N];
14
       int ind[N];
15
       int t[M], c[M], nt[M], opp[M];
16
       int cnt:
17
       int n, m;
18
       int ST, ED;
19
20
       int add(int a, int b, int C) {
21
         t[cnt] = b;
22
         nt[cnt] = ind[a];
23
         ind[a] = cnt:
24
         c[cnt] = C;
25
         return cnt++;
26
       }
27
28
       void add1(int a, int b, int c) {
29
         static int x, y;
30
         x = add(a, b, c);
31
         y = add(b, a, 0);
32
         opp[x] = y; opp[y] = x;
33
34
35
       bool bfs() {
36
         int 1 = 0, r = 0;
37
         q[1] = ST;
38
         memset(h, 0, sizeof h);
39
         h[ST] = 1;
         while(1 <= r) {</pre>
40
41
           int x = q[1++];
           for(int k = ind[x]; k != -1; k = nt[k])
42
43
              if(!h[t[k]] && c[k] > 0) {
44
               h[t[k]] = h[x] + 1;
45
                q[++r] = t[k];
46
47
48
         return h[ED] != 0;
49
       }
50
51
       int dfs(int x, int p) {
52
         if(x == ED) return p;
53
         bool flg = false;
54
         int tot = 0;
55
         for(int k = ind[x]; k != -1; k = nt[k])
           if(h[t[k]] == h[x] + 1) {
56
57
              int d = dfs(t[k], min(p, c[k]));
58
             if(d) {
59
               p -= d;
60
                tot += d;
```

```
c[k] -= d:
61
62
                flg = true;
63
                c[opp[k]] += d;
64
65
           }
66
         return tot;
67
       }
68
69
       int dinic() {
70
         int ans = 0, tmp;
71
         while(bfs()) {
72
           while(true) {
73
              int fw = dfs(ST, INF);
74
             if(!fw) break;
75
              else ans += fw;
76
           }
77
         }
78
         return ans;
79
80 }
```

Dijkstra 费用流

```
* dijkstra 费用流
3
   */
   #include <bits/stdc++.h>
   typedef long long LL;
   typedef std::pair<int, int> PII;
   #define MP(x, y) std::make_pair(x, y)
   #define COST INF 1e9
   #define AA first
10
   #define BB second
11 #define SZ size()
   #define PB(x) push_back(x)
   #define cmin(x, y) x = min(x, y)
14 template <typename T> class MinCostFlow{
15
16
       struct edge{int to;LL cap;T cost;int rev;};
17
18
19
       std::vector<std::vector<edge> >adj;
20
       std::vector<T>pot;
21
22
       std::pair<LL,T>dijkstra(int s,int t,LL FLOW_BOUND){
23
         std::vector<int>used(V.0):
24
         std::vector<T>dist(V,COST_INF);
25
         std::vector<PII>path(V,MP(-1,-1));
         std::priority_queue<std::pair<T,int> >Q;
26
27
         dist[s]=0;
```

```
28
         Q.push(MP(0,s));
29
         while(!Q.empty()){
30
           int x=Q.top().BB;
31
           Q.pop();
32
           if(used[x])continue;
33
           used[x]=1:
34
           for(int i=0;i<adj[x].SZ;i++)if(adj[x][i].cap>0){
35
              edge e=adj[x][i];
36
              int y=e.to;
37
              T d=dist[x]+e.cost+pot[x]-pot[y];
38
              if(d<dist[y]&&!used[y]){</pre>
39
                dist[y]=d;
40
                path[y]=MP(x,i);
41
                Q.push(MP(-d,y));
42
43
           }
44
         }
45
         for(int i=0;i<V;i++)</pre>
46
           pot[i]+=dist[i];
47
         if (dist[t] == COST_INF)
48
           return MP(0,0);
49
         LL f=FLOW_BOUND;
50
         T sum=0:
51
         int x=t;
52
          while(x!=s){
53
           int y=path[x].AA;
54
           int id=path[x].BB;
55
           sum+=adj[y][id].cost;
56
           cmin(f,adj[y][id].cap);
57
           x=y;
58
59
         x=t;
60
         while(x!=s){
           int y=path[x].AA;
61
           int id=path[x].BB;
62
63
           adi[v][id].cap-=f;
           int id2=adj[y][id].rev;
64
65
           adj[x][id2].cap+=f;
66
           x=y;
67
68
         return MP(f,f*sum);
69
       }
70
     public:
       MinCostFlow(int n){//[0,n)
71
72
         adj.resize(V,std::vector<edge>(0));
73
74
         pot.resize(V,0);
75
76
       void add_edge(int s,int t,LL f,T c){
77
          edge e1={t,f,c,(int)adj[t].SZ};
```

```
edge e2={s,OLL,-c,(int)adj[s].SZ};
78
79
         adj[s].PB(e1);
80
         adj[t].PB(e2);
81
82
       std::pair<LL,T>mincostflow(int s,int t,LL FLOW_BOUND=(1LL<<48)){
         std::pair<LL,T>ans=MP(OLL,0);
83
         while(FLOW BOUND>0){
84
85
           std::pair<LL,T>tmp=dijkstra(s,t,FLOW_BOUND);
           if(tmp.AA==0)break;
86
87
           ans.AA+=tmp.AA;
88
           ans.BB+=tmp.BB;
89
           FLOW_BOUND-=tmp.AA;
90
91
         return ans;
92
       }
93 };
```

zkw 费用流

```
* luoshiying 版本
   * 使用前需要给src, des赋值
   * 调用zkw(src, des, n)中的tn为节点数目,要稍大于总数目
   */
   #include <bits/stdc++.h>
   #define MOD 0x3f3f3f3f
   const int Maxn = 3000;
   const int Maxm = 100000:
   struct edge
11
12
     int u, v, c, w, next;
13 | le [Maxm]:
  int last[Maxn];
14
15 int tot:
16 int flow, cost, value;
   int dist[Maxn], visit[Maxn], src, des;
18 std::deque<int> Q:
19
   int n, m;
20
   void adde(int u, int v, int c, int w) {
21
     e[tot].u = u; e[tot].v = v; e[tot].c = c; e[tot].w = w; e[tot].next = last[u]; last[u]
22
     e[tot].u = v: e[tot].v = u: e[tot].c = 0: e[tot].w = -w: e[tot].next = last[v]:
23
       \hookrightarrow last[v] = tot++:
24 }
25
26
   int Aug(int u, int m) {
27
     if(u == des) {
28
       cost += value * m:
29
       flow += m;
```

```
30
       return m:
31
32
     visit[u] = true;
33
     int 1 = m:
34
     int j, v, c, w;
     for(j = last[u]; j != -1; j = e[j].next) {
35
       v = e[j].v; c = e[j].c; w = e[j].w;
36
37
       if(c && !w && !visit[v]) {
38
         int del = Aug(v, 1 < c?1:c);
39
         e[i].c -= del; e[i ^ 1].c += del; l -= del;
40
         if(!1) return m:
       }
41
42
43
     return m - 1;
44
45
46
   bool Modlabel(int src, int des, int n) {
     int i, j, u, v, c, w, del;
48
     memset(dist, 0x3f, sizeof(dist[0])*(n + 3));
49
     dist[src] = 0:
     while(!Q.empty()) Q.pop_back();
50
51
     Q.push_back(src);
52
     while(!Q.empty()) {
53
       u = Q.front(); Q.pop_front();
54
       for(j = last[u]; j != -1; j = e[j].next) {
55
         v = e[j].v; c = e[j].c; w = e[j].w;
56
         if(c && (del = dist[u] + w) < dist[v]) {</pre>
57
           dist[v] = del;
58
           if(Q.empty() || del <= dist[Q.front()]) {</pre>
59
              Q.push_front(v);
60
           }
61
           else {
62
              Q.push_back(v);
63
           }
64
65
66
     for(i = 0; i < n; i++) {</pre>
67
       for(j = last[i]; j != -1; j = e[j].next) {
         e[j].w -= dist[e[j].v] - dist[i];
69
70
       }
71
72
     value += dist[des]:
     return dist[des] < MOD:
74
75
   void zkw(int src, int des, int n) {
     value = cost = flow = 0;
78
     while(Modlabel(src, des, n)){
79
       do {
```

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欧拉回路

```
* 存存在条件为奇数度的点 () 个
   * 有向图所有点入度等于出度
   * 求欧拉回路需要注意访问过的边需要删除
   * 2016 NEERC Moscow G.
   #include <stack>
   const int N = 1000:
   const int M = N * N:
10
   struct Graph {
     int ind[N], vis[N];
11
     int nt[M], t[M], opp[M], chose[M];
12
13
     std::stack<int> stk;
14
15
     void dfs(int x) {
16
      stk.push(x);
       while(stk.size()) {
17
18
        x = stk.top(); stk.pop();
19
        for(int k = ind[x]; ind[x] != -1; ind[x] = nt[k], k = ind[x])
          if(!vis[k]) {
20
            // 无向图加入下面这句
21
            vis[k] = vis[opp[k]] = true;
22
23
            // ========
24
            chose[k] = true;
25
            stk.push(x);
            stk.push(t[k]);
26
27
            break;
28
      }
29
      if(ind[x] == -1) {
30
        // 这里记录 答案
31
32
33
    }
34 };
```

二分图最大匹配 匈牙利算法

```
1/*
* kuangbin模板 ,

* buskin * busk
```

```
#include <bits/stdc++.h>
8 const int MAXN=1000;
9 int uN, vN;
10 int g[MAXN][MAXN];
11 int linker[MAXN];
   bool used[MAXN]:
   bool dfs(int u) {
     int v;
14
15
     for(v = 0; v < vN; v++)
16
       if(g[u][v] && !used[v]) {
17
         used[v] = true;
18
         if(linker[v] == -1 || dfs(linker[v])) {
19
           linker[v] = u;
20
           return true;
21
22
       }
23
     return false;
24
25
   int hungary() {
26
27
     int res = 0;
28
     int u;
     memset(linker, -1, sizeof(linker));
29
     for(u = 0; u < uN; u++)
30
31
32
       memset(used, 0, sizeof(used));
33
       if(dfs(u)) res++;
34
35
     return res;
36
```

最小树形图

```
      1
      /*

      2
      * 最小树形图 , 就是给有向带权图中指定一个特殊的点 root ,

      3
      * 求一棵以 root 为根的有向生成树T ,并且T中所有边的总权值最小 .

      4
      * 最小树形图 ( 根固定 ) ,O(VE) .

      5
      *

      6
      * 求一个有向图的最小生成树 ,

      7
      * 如果根不固定 ,添加一个根节点与所有点连无穷大的边 ,

      8
      * 或者总边权十1 ,如果求出比2 * MOD大或者返回值为一1 ,则不连通 ;

      9
      * 求根 ,则求和虚拟根相连的结点 .

      10
      *

      11
      * 根据pre的信息能构造出这棵树 ,注意结点必须从[0,n) ,

      12
      * 因为要考虑重新标号建图的统一 ,mytype 根据实际情况确定 .

      13
      */

      14
      const int Maxn = 1000;

      15
      const double MOD = 1e9;
```

```
16 struct obi {
17
     int u, v;
18
     double w;
19 | } e[Maxn * Maxn];
20
   int n, m;
21
22
   typedef double mytype;
23
   int visit[Maxn], pre[Maxn], belong[Maxn], ROOT;
24 mytype inv[Maxn];
   mytype dirtree(int n, int m, int root) {
26
     mytype sum = 0;
27
     int i, j, k, u, v;
28
     while (true) {
29
       for (i = 0; i < n; i++) {
30
         inv[i] = MOD;
31
         pre[i] = -1;
32
         belong[i] = -1;
33
         visit[i] = -1;
34
       }
35
       inv[root] = 0;
       for (i = 0; i < m; i++) { // 除原点外, 找每个点的最小入边
36
37
         u = e[i].u; v = e[i].v;
         if (u != v) {
38
39
           if (e[i].w < inv[v]) {</pre>
40
             inv[v] = e[i].w;
             pre[v] = u;
41
42
             if(u == root) ROOT = i; // 记录根所在的边
             // 输出根时利用ROOT - m计算是原图哪个结点
43
           }
44
         }
45
       }
46
47
       for (i = 0; i < n; i++) {
48
49
         if (inv[i] == MOD) return -1;
50
       }
51
       int num = 0;
52
       for (i = 0; i < n; i++) { // 找圈, 收缩圈
53
         if (visit[i] == -1) {
54
55
           for(j = i; j != -1 && visit[j] == -1 && j != root; j = pre[j]) {
             visit[j] = i;
56
57
           }
58
           if (j != -1 && visit[j] == i) {
59
             for (k = pre[j]; k != j; k = pre[k]) {
60
               belong[k] = num;
61
62
             belong[j] = num++;
           }
63
64
65
         sum += inv[i];
```

```
66
67
       if (num == 0) return sum;
68
       for (i = 0; i < n; i++){
69
         if (belong[i] == -1) {
70
           belong[i] = num ++;
71
72
       }
73
       for (i = 0; i < m; i++) { // 重新构图
         e[i].w = e[i].w - inv[e[i].v];
74
75
         e[i].v = belong[e[i].v];
76
         e[i].u = belong[e[i].u];
77
78
       n = num;
79
       root = belong[root];
80
81
```

哈密尔顿回路

```
* 设dp[i][j]表示站在i点,盘面状态为j的最短路
    * 最后答案需要遍历所有i == (1 << n) - 1
   * 注意有的题需要用Floyd预处理最短路
   */
   #include <bits/stdc++.h>
   const int N = 18;
   const int INF = 1e9+7;
   int f[N][N];
   int dp[N][(1 << N)];</pre>
11
   int n, m;
   int main() {
13
     int T; scanf("%d", &T);
14
     while(T--) {
15
       scanf("%d%d", &n, &m);
16
       for(int i = 1; i <= n; i++)
17
         for(int j = 1; j \le n; j++) f[i][j] = INF;
18
       for(int i = 1; i \le n; i++) f[i][i] = 0;
19
       for(int i = 0; i < m; i++) {
20
        int x, y, z; scanf("%d%d%d", &x, &y, &z);
        // 处理重边
21
22
         f[x][y] = std::min(f[x][y], z);
23
        f[y][x] = std::min(f[y][x], z);
       }
24
25
       for(int k = 1; k \le n; k++)
26
        for(int i = 1; i <= n; i++)
27
           for(int j = 1; j <= n; j++)
28
             f[i][j] = std::min(f[i][k] + f[k][j], f[i][j]);
29
30
       for(int i = 1; i <= n; i++)
31
        for(int j = 0; j < 1 << (n + 1); j++) dp[i][j] = INF;
```

```
32
33
       dp[1][2] = 0;
34
35
       for(int j = 0; j < (1 << (n + 1)); j++) {
36
         for(int i = 1; i <= n; i++)
           if(((1 << i) & j) != 0) {
37
             for(int k = 1; k \le n; k++)
38
39
               dp[k][(1 << k) | j] =
                 std::min(dp[i][j] + f[i][k]
40
                      , dp[k][(1 << k) | j]);
41
42
           }
43
       }
44
       int ans = INF;
45
       for(int i = 1; i <= n; i++) {
         ans = std::min(ans, dp[i][(1 << (n + 1)) - 2] + f[i][1]);
46
47
48
       printf("%d\n", ans);
49
50
     return 0;
51 }
```

增广路费用流

```
* 增广路版费用流,复杂度O(n^2m)
   * 调用之前需要清空ind并令cnt = 1
   #include <bits/stdc++.h>
   #define N 3000
   #define M 100000
   #define INF 0x7fffffff
   typedef long long 11;
10
11 int ind[N], pre[N], f[N];
12 bool vis[N];
13 int bg[M], t[M], nt[M], c[M], op[M], v[M];
14 int cnt:
  int S, T;
15
16
17 int add(int a, int b, int C, int V) { // 不要直接调用
    bg[cnt] = a;
18
    t[cnt] = b;
19
     v[cnt] = V:
20
21
     c[cnt] = C;
22
     nt[cnt] = ind[a];
23
     ind[a] = cnt:
24
     return cnt++;
25 }
26
27 | int ADD(int a, int b, int c, int v) {
```

```
int x = add(a, b, c, v):
29
     int y = add(b, a, 0, -v);
30
     op[x] = y; op[y] = x;
31
     return x;
32 }
33
   int h[N], q[N];
34
35
   bool spfa() {
36
     memset(vis, 0, sizeof vis);
     for(int i = S; i <= T; i++) f[i] = INF;</pre>
     for(int i = S; i <= T; i++) pre[i] = -1;</pre>
     pre[S] = -1;
41
     f[S] = 0;
     int 1 = 0, r = 0; q[1] = S; vis[S] = true;
     while(1 \le r)
43
44
45
       int x = q[1 \% N];
46
       l++; vis[x] = false;
47
       for(int k = ind[x]; k != -1; k = nt[k])
48
         if(c[k] > 0 \&\& f[t[k]] > f[x] + v[k])
         {
49
50
           f[t[k]] = f[x] + v[k];
51
           pre[t[k]] = k;
52
           if(!vis[t[k]])
53
           {
54
              r++;
55
              q[r \% N] = t[k];
56
              vis[t[k]] = true;
57
58
         }
59
60
     return pre[T] != -1;
61
62
   int dfs()
64
     int ans = INF;
     int k = pre[T];
     while(k != -1)
67
68
     {
69
       ans = std::min(ans, c[k]);
70
       k = pre[bg[k]];
71
72
     k = pre[T];
     while(k != -1)
73
74
75
       c[k] -= ans;
76
       c[op[k]] += ans;
77
       k = pre[bg[k]];
```

```
}
78
79
     return ans;
80
81
   ll dinic() {// 程序入口
82
      11 \text{ ans} = 0, \text{ tmp};
83
      while(spfa())
84
85
        ans += (11)f[T] * dfs();
86
     }
87
88
     return ans;
89
   }
```

无向图最小割

```
#include <bits/stdc++.h>
   const int maxn = 400;
   const int inf = 1e9;
   int cost[maxn] [maxn], seq[maxn], len[maxn], n, m, pop, ans;
   bool used[maxn];
   void Init(){
     int i,j,a,b,c;
     for(i=0;i<n;i++) for(j=0;j<n;j++) cost[i][j]=0;</pre>
     for(i=0:i<m:i++){</pre>
10
       scanf("%d %d %d",&a,&b,&c);
11
12
       a--; b--;
       cost[a][b]+=c; cost[b][a]+=c;
13
14
15
     pop=n; for(i=0;i<n;i++) seq[i]=i;
16
17
18
   void Work(){
     ans=inf; int i,j,k,l,mm,sum,pk;
19
20
     while(pop > 1){
       for(i=1;i<pop;i++) used[seq[i]]=0; used[seq[0]]=1;</pre>
21
       for(i=1;i<pop;i++) len[seq[i]]=cost[seq[0]][seq[i]];</pre>
22
23
       pk=0; mm=-inf; k=-1;
24
       for(i=1;i<pop;i++) if(len[seq[i]] > mm){ mm=len[seq[i]]; k=i; }
       for(i=1;i<pop;i++){</pre>
25
26
         used[seq[l=k]]=1;
27
         if(i==pop-2) pk=k;
28
         if(i==pop-1) break;
29
         mm=-inf;
30
         for(j=1;j<pop;j++) if(!used[seq[j]])</pre>
31
            if((len[seq[j]]+=cost[seq[1]][seq[j]]) > mm)
32
              mm=len[seq[j]], k=j;
33
       }
34
       sum=0:
35
       for(i=0;i<pop;i++) if(i != k) sum+=cost[seq[k]][seq[i]];</pre>
```

```
36
        ans=std::min(ans.sum):
37
       for(i=0;i<pop;i++)</pre>
38
          cost[seq[k]][seq[i]]=cost[seq[i]][seq[k]]+=cost[seq[pk]][seq[i]];
39
       seq[pk]=seq[--pop];
40
41
     printf("%d\n",ans);
42
43
   void solve() {
44
45
     int K;
     scanf("%d%d%d", &n, &m, &K);
     if(n == 0 \&\& m == 0 \&\& K == 0) exit(0);
     Init():
49
     Work();
50
51
52
   int main() {
53
     while(true) solve();
54
     return 0;
55
```

一般图最大匹配 带花树

```
1 #include <bits/stdc++.h>
^2 /* 求一般图的最大匹配 ,复杂度O(n^3) .
   * g[i][j] 存放关系图 : i,j 否有边 , match[i] 存放i所匹配的点
   * 建图开始初始化q .
   * 最终匹配方案为match .
   * 复杂度O(n^3)
   * 点是从1到n的
   * URAL 1099
   const int MAXN = 500:
   std::deque<int> Q;
   bool g[MAXN] (MAXN], inque [MAXN], inblossom [MAXN], inpath [MAXN];
  int match[MAXN],pre[MAXN],base[MAXN];
14
   // 找公共祖先
15
   int findancestor(int u.int v)
17
18
     memset(inpath,false,sizeof(inpath));
19
     while(1)
20
21
      u=base[u];
22
      inpath[u]=true;
23
      if (match[u] ==-1) break;
24
      u=pre[match[u]];
25
26
     while(1)
```

```
{
27
28
       v=base[v];
29
       if(inpath[v])return v;
       v=pre[match[v]];
30
31
    }
32
   }
33
   // 压缩花
34
   void reset(int u,int anc)
35
36
   ł
     while(u!=anc)
37
38
39
       int v=match[u];
40
       inblossom[base[u]]=1;
       inblossom[base[v]]=1;
41
42
       v=pre[v];
43
       if(base[v]!=anc)pre[v]=match[u];
44
       u=v:
     }
45
46
47
   void contract(int u,int v,int n)
48
49
     int anc=findancestor(u,v);
50
     memset(inblossom,0, sizeof inblossom);
51
52
     reset(u,anc);reset(v,anc);
     if(base[u]!=anc)pre[u]=v;
53
54
     if(base[v]!=anc)pre[v]=u;
     for(int i=1;i<=n;i++)</pre>
55
       if(inblossom[base[i]])
56
       {
57
         base[i]=anc:
58
59
         if(!inque[i])
60
           Q.push_back(i);
61
62
           inque[i]=1;
63
64
       }
65
66
   bool bfs(int S,int n)
68
     for(int i=0;i<=n;i++)pre[i]=-1,inque[i]=0,base[i]=i;</pre>
69
     Q.clear();Q.push_back(S);inque[S]=1;
70
     while(!Q.empty())
71
72
73
       int u=Q.front();Q.pop_front();
74
       for(int v=1;v<=n;v++)</pre>
75
76
         if (g[u][v]&&base[v]!=base[u]&&match[u]!=v)
```

```
77
 78
             if(v==S||(match[v]!=-1&&pre[match[v]]!=-1))contract(u,v,n);
 79
             else if(pre[v]==-1)
            {
 80
81
               pre[v]=u;
               if(match[v]!=-1)Q.push_back(match[v]),inque[match[v]]=1;
 82
 83
               else
 84
               {
 85
                 u=v;
 86
                 while(u!=-1)
 87
                 ł
 88
                   v=pre[u];
                   int w=match[v]:
 89
 90
                   match[u]=v;
 91
                   match[v]=u;
 92
                   u=w;
 93
                 }
 94
                 return true;
 95
 96
 97
 98
 99
      return false;
100
101
102
    int solve(int n)
103
104
105
        memset(match, -1, sizeof match);
106
        int ans=0;
107
        for(int i=1;i<=n;i++)</pre>
108
             if(match[i] == -1 && bfs(i,n))
109
                 ans++;
110
        return ans;
111
```

割点/割边

```
#include <bits/stdc++.h>
using namespace std;

#define MP make_pair

#define AA first

#define BB second

#define SZ size()

#define PB push_back

#define DP begin()

#define ED end()

#define cmin(a, b) a = min(a, b)

typedef std::pair<int, int> PII;
const int N = 50010, M = 1000200, POW = 20, INF = 0x3f3f3f3f;
```

```
13
   class Articulation {
14
   public:
15
     static const int SIZE = N: // 最大结点个数
16
     // keyE 为割边集 , keyV 为割点集
17
     std::vector<PII> keyE;
18
     // cc[p] 表示结点p在哪些双连通分量中
19
     std::vector<int> keyV, cc[SIZE];
20
     int cnt:
21
     // 对于旧版编译器 , 将上面 cc[SIZE] 改成 vector 的形式
22
23
     // 传入结点个数 n 及各结点的出边 e[] , 返回双连通分量的个数 cnt
     int run(int n, const std::vector<int> G[]){
24
       memset(dfn, use = 0, sizeof(dfn[0]) * n);
25
26
       memset(low, cnt = 0, sizeof(low[0]) * n);
27
       keyE.clear();
       fill_n(cc, n, keyV = std::vector<int>());
28
29
       for(int i = 0; i < n; ++i) {</pre>
30
         if(!dfn[i]) dfs(i, 1, G);
31
       }
32
       return cnt:
33
    }
   private:
34
     int dfn[SIZE], low[SIZE], dot[SIZE], use;
35
     void dfs(int x, int dep, const std::vector<int> G[]){
36
37
       int src = 0, out = 1 < dep;
38
       dot[use++] = x;
39
       dfn[x] = low[x] = dep;
40
       for (int i = 0; i < G[x].size(); i++) {</pre>
         int y = G[x][i];
41
42
         if (!dfn[v]){
43
           dfs(y, dep + 1, G);
44
           low[x] = std::min(low[x], low[y]);
           if (low[y] > dfn[x]) keyE.push_back(PII(x, y));
45
46
           if (low[y] >= dfn[x]){
47
             if (++out == 2) keyV.push_back(x);
48
             while (dot[--use] != y) cc[dot[use]].push_back(cnt);
49
             cc[x].push_back(cnt);
50
             cc[y].push_back(cnt++);
51
           }
         } else if (dfn[y] != dfn[x] - 1 || src++){
52
53
           low[x] = std::min(low[x], dfn[y]);
54
55
       }
56
    }
57 };
```

斯坦纳树

```
1 /*
2 * dp[u][i] 表示结点 u 已经和要连通的结点集合 ( 2^k 表示 ) i 连通的最小花费 ,
```

```
* 初始化将 k 个点和 n 个点 dp[u][1<<i] 初始化为最短路 ,
   * , dp[u][0]=0, 加入队列
   * 利用 spfa 求出 dp 数组
   * 状态转移分三部分:
   * 1) dp[u][su] 利用 dp[u][sub] + dp[u][su^sub] 更新 , sub 为 su 子集
   * 2) dp[u][su] 更新相邻的 dp[v][sv]
   * 3) 将 k 中不属于 su 的点与u连接 , 利用 u, k 的最短路
10
11
12
   // hdu 4085
13 // 求斯坦纳森林, 再森林合并为全集
14 #include <bits/stdc++.h>
   #define MOD 0x3f3f3f3f
   typedef std::pair<int, int> PII;
16
17 #define MP std::make_pair
18 #define PB push_back
  #define cmin(x, y) x = std::min(x, y)
20 #define SZ size()
21
  #define AA first
  #define BB second
23 using namespace std;
   const int Maxn = 55;
   const int Maxk = 14;
   const int MSta = 1 << 15;</pre>
27 int dp[Maxn] [MSta];
28 vector<PII> g[Maxn];
   int gg[Maxn][Maxn];
  int mp1[Maxn], mp2[Maxn];
31 int n, m, K;
   vector<int> split[Maxk];
  int gao[MSta];
33
  int has (int sta) {
    int i, a = 0, b = 0, u;
    for (i = 0; i < 2 * K; i++) {
37
      if ((1<<i) & sta) {
38
        u = mp1[i];
39
        if (u <= K) a++:
40
        else b++;
41
      }
42
43
     if (a == b) return true;
44
     else return false;
45 }
   int ans:
   queue<PII> que;
  int inque[Maxn][MSta];
   int doit() {
    int i, j, k, u, v, w, su, sv, sub, vv, tp;
51
   int kk = 2 * K:
   for (i = 1; i <= n; i++) {
```

```
for (j = 0; j < (1 << kk); j++) dp[i][j] = MOD, inque[i][j] = 0;
54
        dp[i][0] = 0; inque[i][0] = 0;
     }
55
      while(!que.empty()) que.pop();
56
57
      memset(mp2, 0, sizeof(mp2));
      for(i = 1, j = 0; i <= K; i++, j++)mp1[j] = i, mp2[i] = 1<<j;
58
      for(i = n - K + 1; i <= n; i++, j++)mp1[j] = i, mp2[i] = 1<<j;
59
60
      for(i = 0; i < kk; i++) {</pre>
        for(u = 1; u <= n; u++) {
61
62
          cmin(dp[u][1<<i], gg[u][mp1[i]]);</pre>
63
          inque[u][1<<i] = 1;
          que.push(MP(u, 1<<i));
64
        }
65
66
      }
      while(!que.empty()) {
67
        u = que.front().AA; su = que.front().BB; que.pop();
68
69
        sub = su:
70
        tp = dp[u][su];
71
        do {
72
          sv = sub ^ su;
          cmin(tp, dp[u][sub] + dp[u][sv]);
73
74
          sub = (sub - 1) \& su;
        }while(sub != su);
75
        if(tp < dp[u][su]) dp[u][su] = tp;</pre>
76
77
        for(j = 0; j < g[u].SZ; j++) {
78
          v = g[u][j].AA; w = g[u][j].BB;
79
          if(su & mp2[v]) continue;
80
          sv = su \mid mp2[v];
81
          if(dp[v][sv] > dp[u][su] + w) {
82
            dp[v][sv] = dp[u][su] + w;
83
            if(!inque[v][sv]) {
84
              inque[v][sv] = 1;
85
               que.push(MP(v, sv));
            }
86
87
          }
88
89
        for(j = 0; j < kk; j++) {
90
          if(su & (1<<j)) continue;
91
          sv = su | (1 << j);
          if(dp[u][sv] > dp[u][su] + gg[u][mp1[j]]) {
92
93
            dp[u][sv] = dp[u][su] + gg[u][mp1[j]];
94
            if(!inque[u][sv]) {
95
              inque[u][sv] = 1;
96
              que.push(MP(u, sv));
97
            }
          }
98
99
100
        inque[u][su] = 0;
101
        }
102
```

```
for(i = 0; i <= kk; i++) split[i].clear();</pre>
103
104
        for(i = 0; i < (1<<kk); i++) {</pre>
105
             split[__builtin_popcount(i)].PB(i);
106
             gao[i] = MOD:
107
             if(has(i)) {
                 for(j = 1; j <= n; j++) {
108
                     cmin(gao[i], dp[j][i]);
109
110
                 }
111
            }
112
        }
113
        for(i = 1; i <= kk; i++) {
114
             for(j = 0; j <split[i].SZ; j++) {</pre>
115
                 su = split[i][j];
116
                 for(sub = su - 1 & su; sub; sub = sub - 1 & su) {
117
                     sv = su ^ sub;
                     cmin(gao[su], gao[sub] + gao[sv]);
118
119
                     sub = (sub - 1) \& su;
120
                }
121
             }
122
123
        return gao[(1<<kk) - 1];
124 }
125
    int main() {
126
        int i,j,k,u,v,w;
127
        int te;
128
        scanf ("%d", &te);
129
        for (int ca = 1; ca <= te; ca ++) {
130
             scanf ("%d%d%d", &n, &m, &K);
131
             for (i = 1; i <= n; i++) g[i].clear();</pre>
132
             for (i = 1; i <= n; i++) {
133
                 for (j = 1; j <= n; j++) gg[i][j] = MOD;</pre>
134
                 gg[i][i] = 0;
135
136
             for (i = 0; i < m; i++) {
137
                 scanf ("%d%d%d", &u, &v, &w);
138
                 cmin (gg[u][v], w);
139
                 cmin (gg[v][u], w);
140
            }
141
             for (i = 1; i <= n; i++) {
142
                 for (j = i + 1; j \le n; j++) {
143
                     if (gg[i][j] < MOD) {</pre>
144
                          g[i].PB (MP (j, gg[i][j]));
145
                          g[j].PB (MP (i, gg[i][j]));
146
                     }
                 }
147
             }
148
             for (k = 1; k <= n; k++) {
149
150
                 for (i = 1; i <= n; i++) {
151
                     for (j = 1; j \le n; j++) {
152
                          cmin (gg[i][j], gg[i][k] + gg[k][j]);
```

```
}
153
154
                }
155
            }
156
            ans = 0:
157
            ans = doit();
158
            if(ans >= MOD) printf("No solution\n");
159
            else printf ("%d\n", ans);
160
        }
161
        return 0;
162 }
```

字符串 Hash

```
* 全局 init()
   * get(u64* H, int L, int w) 获取 L 开头 长度为 w 的 hash
   */
   const int N = 1e5+100;
   typedef unsigned long long ull;
   const ull mod = 1e9 + 7;
8 struct Hash {
    ull P[N];
    ull H[N];
10
11
    void init() {
12
    P[0] = 1:
     for(int i = 1; i < N; i++)
13
14
        P[i] = P[i - 1] * mod;
15
    }
16
17
     template <typename T>
18
       void make(T *a, int n) {
19
        H[n] = 0;
        for(int i = n - 1; i >= 0; i--)
20
           H[i] = H[i + 1] * mod + a[i] + 1;
21
22
      }
23
24
    ull get(int L, int w) {
      return H[L] - H[L + w] * P[w];
25
26
27 | } T;
```

KMP

```
6 void build(T b[], int next[], int m) {
       int k = -1;
       for(int i = 0; i < m; i++) {
          while(k != -1 && b[k + 1] != b[i]) k = next[k];
10
          if(k + 1 != i \&\& b[k + 1] == b[i]) k++;
          next[i] = k:
11
12
       }
13 }
14 // 得到 b 在 a 中的第一个匹配位置
   template <typename T>
   int kmp(T a[], T b[], int next[], int n, int m) {
       int k = -1;
17
18
       for(int i = 0; i < n; i++) {
19
          while(k != -1 && b[k + 1] != a[i]) k = next[k];
20
          if(b[k + 1] == a[i]) k++;
21
          if(k == m - 1) return i - m + 2;
22
      }
23
       return -1;
24 }
```

EXKMP

```
/* Hdu 4300
2 * S 为主串 . T 为子串 .
   * LS 为 S 长度 , LT 为 T 长度 ,
   * B[i] 表示 S[i] 匹配了 B[i] 长度的 T
   #include <bits/stdc++.h>
   const int Maxn = 1e5;
   char S[Maxn], T[Maxn];
   int Next[Maxn], B[Maxn];
10
11 void preExKmp(char T[], int LT, int next[]) {
12
   int i, ind = 0, k = 1;
13
    next[0] = LT;
     while(ind + 1 < LT && T[ind + 1] == T[ind]) ind++:</pre>
15
     next[1] = ind;
16
     for(i = 2; i < LT; i++) {</pre>
      if(i \le k + next[k] - 1 && next[i - k] + i \le k + next[k])
17
18
         next[i] = next[i - k];
19
20
         ind = std::max(0, k + next[k] - i);
21
         while(ind + i < LT && T[ind + i] == T[ind]) ind++;</pre>
22
         next[i] = ind; k = i;
23
24
    }
25 }
26
27 void exKmp(char S[], int LS, char T[], int LT, int next[], int B[]) {
```

```
int i. ind = 0. k = 0:
29
     preExKmp(T, LT, next);
30
     while(ind < LS && ind < LT && T[ind] == S[ind]) ind++;</pre>
     B[0] = ind:
31
32
     for(i = 1; i < LS; i++) {</pre>
       int p = k + B[k] - 1, L = next[i - k];
33
       if((i-1)+L < p)
34
35
        B[i] = L;
36
       else {
37
         ind = std::max(0, p - i + 1);
38
         while(ind + i < LS && ind < LT && S[ind + i] == T[ind]) ind++;</pre>
39
         B[i] = ind:
         k = i;
40
41
       }
42
    }
43
```

SA

```
* 倍增算法 , r 为待匹配数组,
   * n 为总长度 +1 , m 为字符范围 , num 保存字符串
   * 使用时注意 num[] 有效位为 [0, n),
   * 但是需要将 num[n] = 0 ,
   * 另外 , 对于模板的处理将空串也处理了 ,
   * 作为rank最小的串 .
   * 因此有效串为 [0, n] 共 n-1 个 ,
   * 在调用da()函数时, 需要调用 da(num, n + 1, m)
  * 对于 sa[], rank[], height[] 数组都将空串考虑在内, 作为 rank 最小的后缀
11
  * 注意 rank , height 范围从 [0, n]
12
  */
13
14
  const int N = 1e5;
15 namespace SA
16 | €
17
    int len;
    int num[N]:
18
19
    // sa[1~n]value(0~n-1); rank[0..n-1] value(1..n); height[2..n]
20
    int sa[N], rank[N], height[N];
    int wa[N], wb[N], wv[N], wd[N];
21
22
23
    int cmp(int *r, int a, int b, int x) {
      return r[a] == r[b] && r[a + x] == r[b + x]:
24
25
    }
26
27
    // 倍增算法 r 为待匹配数组 n 为总长度 +1 . m 为字符范围
28
    void da(int *r, int n, int m) {
29
     int i, j, k, p, *x = wa, *y = wb, *t;
30
      for(i = 0; i < m; i++) wd[i] = 0;
31
      for(i = 0; i < n; i++) wd[x[i] = r[i]]++;
```

```
32
       for(i = 1: i < m: i++) wd[i] += wd[i - 1]:
33
       for(i = n - 1; i \ge 0; i--) sa[--wd[x[i]]] = i;
34
       for(j = 1, p = 1; p < n; j <<= 1, m = p) {
35
         for(p = 0, i = n - j; i < n; i++) y[p++] = i;
36
         for(i = 0; i < n; i++) if(sa[i] >= j) y[p++] = sa[i] - j;
37
         for(i = 0; i < n; i++) wv[i] = x[y[i]];
38
         for(i = 0; i < m; i++) wd[i] = 0;
39
         for(i = 0; i < n; i++) wd[wv[i]]++;</pre>
40
         for(i = 1; i < m; i++) wd[i] += wd[i - 1];</pre>
41
         for(i = n - 1; i \ge 0; i--) sa[--wd[wv[i]]] = v[i];
         for(t = x, x = y, y = t, p = 1, x[sa[0]] = 0, i = 1; i < n; i++) {
43
           x[sa[i]] = cmp(y, sa[i - 1], sa[i], j) ? p - 1 : p++;
44
         }
45
       }
46
47
       for(i = 0, k = 0; i < n; i++) rank[sa[i]] = i;
48
       for(i = 0; i < n - 1; height[rank[i++]] = k) {</pre>
49
         for(k ? k-- : 0, j = sa[rank[i] - 1]; r[i + k] == r[j + k]; k++);
50
51
52
```

DC3

```
1 // 待排序的字符串放在 r 数组中, 从 r[0] 到 r[n-1], 长度为 n, 且最大值小于 m
2 // 约定除 r[n-1] 外所有的 r[i] 都大于 0, r[n-1]=0
3 // 函数结束后, 结果放在 sa 数组中, 从 sa[0] 到 sa[n-1]
 4 #define maxn 10000
5 #define F(x) ((x)/3+((x)\%3==1?0:tb))
 6 #define G(x) ((x)<tb?(x)*3+1:((x)-tb)*3+2)
 7 int wa[maxn].wb[maxn].wv[maxn].wss[maxn]: // 必须这么大
8 int s[maxn*3],sa[maxn*3];
9 int c0(int *r.int a.int b){return r[a] ==r[b]&&r[a+1] ==r[b+1]&&r[a+2] ==r[b+2]:}
10 int c12(int k,int *r,int a,int b){
   if (k==2) return r[a]<r[b]||r[a]==r[b]&&c12(1,r,a+1,b+1);
     else return r[a]<r[b]||r[a]==r[b]&&wv[a+1]<wv[b+1]:
13 }
14 void sort(int *r,int *a,int *b,int n,int m){
    int i: for(i=0:i<n:i++) wv[i]=r[a[i]]:</pre>
    for(i=0;i<m;i++) wss[i]=0; for(i=0;i<n;i++) wss[wv[i]]++;
17
     for(i=1;i<m;i++) wss[i]+=wss[i-1];</pre>
    for(i=n-1:i>=0:i--) b[--wss[wv[i]]]=a[i]:
18
19 }
20 void dc3(int *r,int *sa,int n,int m){
     int i,j,*rn=r+n,*san=sa+n,ta=0,tb=(n+1)/3,tbc=0,p;
22
    r[n]=r[n+1]=0;
     for(i=0;i<n;i++) if(i%3!=0) wa[tbc++]=i;</pre>
     sort(r+2,wa,wb,tbc,m); sort(r+1,wb,wa,tbc,m); sort(r,wa,wb,tbc,m);
25
     for(p=1,rn[F(wb[0])]=0,i=1;i<tbc;i++)
```

```
rn[F(wb[i])]=c0(r,wb[i-1],wb[i])?p-1:p++;
27
     if(p<tbc) dc3(rn,san,tbc,p);</pre>
28
     else for(i=0;i<tbc;i++) san[rn[i]]=i;</pre>
     for (i=0;i<tbc;i++) if(san[i]<tb) wb[ta++]=san[i]*3;</pre>
29
30
     if(n\%3==1) wb[ta++]=n-1;
     sort(r,wb,wa,ta,m); for(i=0;i<tbc;i++) wv[wb[i]=G(san[i])]=i;</pre>
31
     for(i=0, j=0, p=0; i < ta && j < tbc; p++)</pre>
32
33
       sa[p]=c12(wb[j]%3,r,wa[i],wb[j])?wa[i++]:wb[j++];
     for(;i<ta;p++) sa[p]=wa[i++]; for(;j<tbc;p++) sa[p]=wb[j++];}</pre>
34
     int main(){
35
36
       int n,m=0; scanf("%d",&n);
       for (int i=0;i<n;i++) scanf("%d",&s[i]),s[i]++,m=max(s[i]+1,m);</pre>
37
       printf("d^n,m); s[n++]=0; dc3(s,sa,n,m);
38
39
       for (int i=0;i<n;i++) printf("%d ",sa[i]);printf("\n");</pre>
40
```

Manacher 最长回文串

```
* 求以 i 为中心的最长回文串长度 ,
   * 结果保存在 pk[i] 中 ,下标 [0,n) ,
   * 开头 , 末尾 , 字符之间插入 \# , 例如 ababa 变为 \#a\#b\#a\#b\#a\#
   #include <bits/stdc++.h>
  const int N = 1e5;
  int pk[N];
9 template <typename T>
10 // [0, n)
  void manacher(T *a, int n) {
    int mx = 0:
12
13
    int p;
    for(int i = 0; i < n; i++) {
14
      if(i < mx) pk[i] = std::min(pk[2 * p - i], mx - i);</pre>
15
      else pk[i] = 1;
16
17
      while(i + pk[i] < n && i - pk[i] > -1 && a[i + pk[i]] == a[i - pk[i]]) pk[i]++;
      if(i + pk[i] > mx) { p = i; mx = i + pk[i]; }
18
19
    }
20 }
```

最小表示法

```
/*
    * hdu 3374
    */
#include <string>
#include <algorithm>

std::string find(std::string s) {
    int i,j,k,l,N=s.length(); s+=s;
```

```
for(i=0,j=1;j<N;){
    for(k=0;k<N&&s[i+k]==s[j+k];k++);
    if(k>=N) break;
    if(s[i+k]<s[j+k]) j+=k+1;
    else l=i+k,i=j,j=std::max(l,j)+1;
}
return s.substr(i,N);
}</pre>
```

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SAM

```
* cxlove
   * spoj NSUBSTR
   * 给一个字符串 S,
   * 令 F(x) 表示 S 的所有长度为x的子串中 ,
   * 出现次数的最大值,
   * 求 F(1)..F(Length(S))
   * 结点的 len 值表示那一时刻的后缀长度
   */
   #include <bits/stdc++.h>
11 #define inf 100000005
  #define M 40
13 #define N 510005
14 #define maxn 300005
15 #define eps 1e-10
16 #define zero(a) fabs(a)<eps
17 #define pb(a) push_back(a)
18 #define mp(a,b) make_pair(a,b)
19 #define mem(a,b) memset(a,b,sizeof(a))
20 #define LL unsigned long long
21 #define MOD 1000000007
22 #define lson step<<1
23 #define rson step<<1|1
   #define sqr(a) ((a)*(a))
25 #define Key_value ch[ch[root][1]][0]
   #define test puts("OK");
   #pragma comment(linker, "/STACK:1024000000,1024000000")
28
   struct SAM {
     SAM *pre,*son[26];
29
30
     int len,g;
   }que[N],*root,*tail,*b[N];
  int tot:
33 void add(int c,int 1) {
     SAM *p=tail,*np=&que[tot++];
35
     np->len=1;tail=np;
     while(p&&p->son[c]==NULL) p->son[c]=np,p=p->pre;
     if(p==NULL) np->pre=root;
38
     else {
```

```
SAM *q=p->son[c];
39
40
       if(p->len+1==q->len) np->pre=q;
41
       else {
         SAM *nq=&que[tot++];
42
43
         *nq=*q;
         nq->len=p->len+1;
44
         np->pre=q->pre=nq;
45
46
         while(p&&p->son[c]==q) p->son[c]=nq,p=p->pre;
47
48
    }
49
   char str[N/2];
50
   int dp[N/2];
51
52
   int main() {
     while(scanf("%s",str)!=EOF) {
       int n=strlen(str);
54
55
       tot=0:
56
       root=tail=&que[tot++];
       for(int i=0;i<n;i++) add(str[i]-'a',i+1);</pre>
57
58
       int cnt[N/2];mem(cnt,0);
59
       for(int i=0;i<tot;i++) cnt[que[i].len]++;</pre>
       for(int i=1;i<=n;i++) cnt[i]+=cnt[i-1];</pre>
60
61
       for(int i=0;i<tot;i++) b[--cnt[que[i].len]]=&que[i];</pre>
62
       for(int i=0;i<n;i++) (root=root->son[str[i]-'a'])->g++;
63
       mem(dp,0);
64
       for(int i=tot-1;i>0;i--) {
         dp[b[i]->len]=std::max(dp[b[i]->len],b[i]->g);
65
         if(b[i]->pre) b[i]->pre->g+=b[i]->g;
66
67
       }
       for(int i=n-1;i>0;i--) dp[i]=std::max(dp[i],dp[i+1]);
68
       for(int i=1;i<=n;i++) printf("%d\n",dp[i]);</pre>
69
70
71
     return 0;
72 }
```

Trie

```
1 #include <bits/stdc++.h>
2
3 const int M = 26;
4 const int N = 2000;
5 struct trie {
6    int to[N][M]; // 节点指针
7    int w[256]; // 字母下标
8    int fail[N]; // 失配指针
9    int cnt;
10
11 trie() {
12    // 标注下标
13 }
```

```
14
15
     int newNode() {
       memset(to[cnt], 0, sizeof to[cnt]);
16
17
       return cnt++:
18
19
20
     void init() {
21
       cnt = 0:
22
       memset(to[0], 0, sizeof to[0]);
23
24
25
     void insert(char *a, int mask) {
26
       int p = 0;
27
       while(*a) {
28
         int v = w[*a];
29
         if(!to[p][v]) {
30
           to[p][v] = newNode();
31
32
         p = to[p][v];
33
         a++;
34
35
     }
36
37
     void ac() {
38
       std::queue<int>q;
39
       memset(fail, 0, sizeof fail);
40
       for(int i = 0; i < 4; i++) if(to[0][i]) {</pre>
41
         q.push(to[0][i]);
42
43
       while(q.size()) {
44
         int p = q.front();
45
         q.pop();
46
         for(int i = 0; i < 4; i++) {
47
           int v = to[p][i];
48
           if(v) {
49
              fail[v] = to[fail[p]][i];
              q.push(v);
50
51
           }
52
53
              to[p][i] = to[fail[p]][i];
54
55
56
    }
57
   };
```

数学 Fib

```
#include <bits/stdc++.h>
   using namespace std;
   const int MOD = 1e9+7;
   typedef long long LL;
   LL mul(LL a, LL b) {
     LL tot = 0;
     while(b) {
9
         if(b \& 1) tot = (tot + a):
10
         b >>= 1:
11
         a <<= 1:
12
         if(tot >= MOD) tot -= MOD;
13
         if(a >= MOD) a -= MOD;
       }
14
15
       return tot;
16 };
17
18
   void fib(LL n, LL &x, LL &y) {
     if (n == 1) {
19
20
      x = y = 1;
21
       return:
22
    } else if (n & 1) {
23
      fib(n - 1, y, x);
24
      y += x; y %= MOD;
25
    } else {
26
     LL a. b:
27
      fib(n >> 1, a, b);
28
       y = (mul(a, a) \% MOD + mul(b, b) \% MOD) \% MOD;
29
       x = (mul(a, b) \% MOD + mul(a, (b + MOD - a)) \% MOD) \% MOD;
30
    }
31 }
```

矩阵

```
#include <bits/stdc++.h>
   const int MXN = 4:
   const int mod = 1e4+7;
   template <typename T> class Matrix{
     public:
       int n,m;
8
      T a[MXN][MXN];
9
       Matrix() {memset(a , 0 , sizeof(a));}
10
       Matrix(int _n , int _m) {
11
         n = _n , m = _m , memset(a , 0 , sizeof(a));
12
       }:
13
       T* operator[] (int i) {return a[i];}
```

```
14
       Matrix friend operator * (Matrix A , Matrix B){
15
         Matrix ans(A.n , B.m);
16
         for (int i = 0; i < A.n; ++ i)
17
           for (int j = 0; j < A.m; ++ j)if(A[i][j]!=0)
18
             for (int k = 0; k < B.m; ++ k)
               ans[i][k]=(ans[i][k]+A[i][j]*B[j][k]) % mod;
19
20
         return ans:
21
22 };
```

高斯消元

```
1 // 二进制版本 . hdu5833
  #include <bits/stdc++.h>
   const int N = 1000;
   double a[N][N];
  int gauss(int eq, int var) {
     int line = 0. ret = 0:
     for(int i = 0; i < var; i++) {</pre>
      int target = line;
10
      while(target < eq && a[target][i] == 0) target++;</pre>
11
      if(target == eq) { ret++; continue; }
12
       std::swap(a[target], a[line]);
      for(int j = 0; j < eq; j++) if(j != line && a[j][i]) a[j] ^= a[line];
13
14
15
16
     return ret;
17
18
   // 数值版本
   int Gauss(int equ, int var) {
     int i. i. k:
21
     int max_r; // 当前这列绝对值最大的行 .
23
     int col; // 当前处理的列
24
     // 转换为阶梯阵 .
     col = 0; // 当前处理的列
26
     for(k = 0; k < equ && col < var; <math>k++, col++) {
      // 枚举当前处理的行 .
27
      // 找到该 col 列元素绝对值最大的那行与第 k 行交换 .
28
29
      // ( 为了在除法时减小误差 )
30
      max r=k:
31
       for(i = k + 1; i < equ; i++) {</pre>
32
        if((a[i][col]) > (a[max_r][col])) max_r = i;
33
      }
       if(max_r != k) {
34
35
        // 与第k行交换 .
        for(j=k; j<var+1; j++) swap(a[k][j], a[max_r][j]);</pre>
36
37
      }
```

```
if((a[k][col])==0) {
38
39
         // 说明该 col 列第 k 行以下全是 O 了 , 则处理当前行的下一列 .
40
41
         continue;
42
       for(i=k+1; i<equ; i++) {</pre>
43
         // 枚举要删去的行 .
44
45
         if((a[i][col]) > 0) {
           for(j = var; j >= col; j--) {
46
47
             a[i][j] -= a[k][j] * (a[i][col] / a[k][col]);
48
           }
         }
49
       }
50
51
     // 无穷解 , 返回自由变元个数
     // 会改变解的顺序
53
     if(k < var) {</pre>
54
55
       for(int i=0; i<equ; i++) {</pre>
56
         if(a[i][i]!= 0) continue;
         else {
57
           int flag=0;
58
59
           for(int j=i+1; j<var; j++) {</pre>
             if(a[i][j] != 0) {
60
61
              flag = 1;
               for(int k = 0; k \le i; k++) {
62
63
                 swap(a[k][i], a[i][j]);
64
               }
65
               break;
             }
66
67
68
           if(!flag)
69
             return var - i;
70
         }
71
       }
72
       return var - equ;
73
     }
74
     return 0:
75
```

四边形不等式

```
/*

* Kyb 示意代码 待验证

*/

#define rep(i,a,n) for(int i=(a);i<(int)(n);i++)

#define better true

void dp(int n, int K[][100]) {

rep(r,1,n){

rep(i,1,n-r){
```

Java 开根号

```
* Java 开根号
    * Kyb , 待验证
   public static BigInteger getsqrt(BigInteger n) {
    if (n.compareTo(BigInteger.ZERO) <= 0) return n;</pre>
     BigInteger x, xx, txx;
                                   xx = x = BigInteger.ZERO;
     for (int t = n.bitLength() / 2; t >= 0; t--) {
       txx = xx.add(x.shiftLeft(t + 1)).add(BigInteger.ONE.shiftLeft(t + t));
       if (txx.compareTo(n) <= 0) {</pre>
10
11
         x = x.add(BigInteger.ONE.shiftLeft(t)); xx = txx;
12
13
     }
14
     return x;
15
```

Cpp 大数

```
* C++ 大数运算 , Kyb , 待验证
   */
   #include <bits/stdc++.h>
   class BigNum {
     public:
       static const int MOD = 100000000;
       static const int BIT = 8, SIZE = 105;
       mutable int n, o;
10
       long long u[SIZE];
11
       BigNum() {}
12
       BigNum(const std::string& s) {
13
         memset(this, 0, sizeof(BigNum));
14
         int num = 0, cnt = 1;
15
         for (int i = s.size() - 1: ~i: i--) {
16
           if (s[i] == '-') o ^= 1;
           if (s[i] >= '0' && s[i] <= '9') {
17
             num += (s[i] - '0') * cnt:
18
19
             cnt *= 10;
```

```
if (cnt == MOD) u[n++] = num, num = 0, cnt = 1:
                                                                                                  70
                                                                                                            BigNum s = *this:
20
21
           }
                                                                                                  71
                                                                                                            if (s.u[0] || s.n >= 2) s.o ^= 1;
22
        }
                                                                                                  72
                                                                                                            return s:
23
         if (!n | | cnt >= 10) u[n++] = num:
                                                                                                  73
                                                                                                          }
         if (!u[0] && n == 1) o = 0;
                                                                                                  74
24
                                                                                                          BigNum& operator +=(const BigNum& r) {
25
       }
                                                                                                  75
                                                                                                            if (r.n == 1 && !r.u[0]) return *this:
26
       BigNum(long long x) {
                                                                                                  76
                                                                                                            if (r.o ^ o) return r.o ^= 1. *this -= r. r.o ^= 1. *this:
27
                                                                                                  77
         memset(this, 0, sizeof(BigNum));
                                                                                                            if (r.n > n) n = r.n;
28
         if (x < 0) o = 1, x = -x;
                                                                                                  78
                                                                                                            for (int i = 0; i < r.n; i++) u[i] += r.u[i];</pre>
29
         do u[n++] = x \% MOD; while (x /= MOD);
                                                                                                   79
                                                                                                            for (int i = 0; i < n; i++) if (u[i] >= MOD) u[i + 1]++, u[i] -= MOD;
30
                                                                                                  80
                                                                                                            if (u[n]) n++:
31
       operator std::string() const {
                                                                                                  81
                                                                                                            return *this:
         static char s[SIZE * BIT + 10]:
32
                                                                                                  82
33
         char* c = s + sprintf(s, "sd'", o ? "-" : "", int(u[n - 1]));
                                                                                                  83
                                                                                                          BigNum& operator -=(const BigNum& r) {
34
         for (int i = n - 2; ~i; i--) c += sprintf(c, "%0*d", BIT, int(u[i]));
                                                                                                  84
                                                                                                            if (r.n == 1 && !r.u[0]) return *this;
                                                                                                            if (r.o ^ o) return r.o ^= 1, *this += r, r.o ^= 1, *this;
35
         return s;
36
       }
                                                                                                  86
                                                                                                            if (cmp(*this, r) * (r.o ? -1 : 1) < 0) {
37
       int operator [](int pos) const {
                                                                                                  87
                                                                                                              o^{=1}. n = r.n:
         static int e[BIT] = {1};
                                                                                                              for (int i = 0; i < r.n; i++) u[i] = r.u[i] - u[i];
38
                                                                                                  88
39
         for (static int i = 1; i < BIT; i++) e[i] = e[i - 1] * 10;
                                                                                                   89
40
         return u[pos / BIT] / e[pos % BIT] % 10;
                                                                                                   90
                                                                                                              for (int i = 0; i < r.n; i++) u[i] = u[i] - r.u[i];
                                                                                                  91
41
       }
42
       int length() const {
                                                                                                  92
                                                                                                            for (int i = 0; i < n; i++) if (u[i] < 0) u[i + 1]--, u[i] += MOD;
43
         int ret = (n - 1) * BIT + 1;
                                                                                                  93
                                                                                                            while (!u[n - 1] \&\& n >= 2) --n;
44
         for (int x = u[n - 1] / 10; x; x /= 10) ret++;
                                                                                                  94
                                                                                                            if (!u[0] \&\& n == 1) o = 0;
45
         return ret:
                                                                                                  95
                                                                                                            return *this:
                                                                                                  96
46
                                                                                                          BigNum operator *(const BigNum& r) const {
47
       friend int cmp(const BigNum& 1, const BigNum& r) {
                                                                                                  97
48
         if (l.o != r.o) return (l.o ? -1 : 1):
                                                                                                  98
                                                                                                            BigNum s = 0:
                                                                                                            if (!u[n - 1] || !r.u[r.n - 1]) return s;
         if (l.n != r.n) return (l.o ? -1 : 1) * (l.n - r.n);
                                                                                                  99
49
         for (int i = 1.n - 1; "i; i--) if (1.u[i] - r.u[i])
50
                                                                                                 100
                                                                                                            s.n = r.n + n - 1;
           return (1.0 ? -1 : 1) * (1.u[i] - r.u[i]);
                                                                                                            s.o = r.o \circ o:
51
                                                                                                 101
52
         return 0;
                                                                                                 102
                                                                                                            for (int i = 0; i < n; i++) for (int j = 0; j < r.n; j++)
       }
                                                                                                              s.u[i + j] += u[i] * r.u[j];
53
                                                                                                 103
       // 运算符
                                                                                                            for (int i = 0; i < s.n; i++) if (s.u[i] >= MOD) {
54
                                                                                                 104
       bool operator < (const BigNum& r) const {return cmp(*this, r) < 0;}
                                                                                                 105
                                                                                                              s.u[i + 1] += s.u[i] / MOD;
55
       bool operator > (const BigNum& r) const {return cmp(*this, r) > 0;}
                                                                                                              s.u[i] %= MOD:
56
                                                                                                 106
       bool operator <=(const BigNum& r) const {return cmp(*this, r) <= 0;}
                                                                                                 107
57
                                                                                                              if (i == s.n - 1) s.n++;
58
       bool operator >=(const BigNum& r) const {return cmp(*this, r) >= 0;}
                                                                                                 108
                                                                                                            }
       bool operator ==(const BigNum& r) const {return cmp(*this, r) == 0;}
                                                                                                 109
59
                                                                                                            return s:
60
       bool operator !=(const BigNum& r) const {return cmp(*this, r) != 0;}
                                                                                                 110
       BigNum operator +(const BigNum& r) const {return BigNum(*this) += r;}
                                                                                                 111
                                                                                                          BigNum operator /(const BigNum& r) const {
61
       BigNum operator -(const BigNum& r) const {return BigNum(*this) -= r;}
                                                                                                 112
                                                                                                            BigNum e[35], s = 0, c = 0;
62
       BigNum operator *(int x) const {return BigNum(*this) *= x;}
                                                                                                 113
63
                                                                                                            int m = 0, ro = r.o, lo = o;
       BigNum operator /(int x) const {return BigNum(*this) /= x;}
                                                                                                 114
64
                                                                                                            r.o ^= ro, o ^= lo;
       BigNum& operator *=(const BigNum& r) {return *this = *this * r;}
65
                                                                                                 115
                                                                                                            for (e[m] = r; MOD >> ++m; e[m] = e[m - 1] + e[m - 1]);
       BigNum& operator /=(const BigNum& r) {return *this = *this / r;}
                                                                                                            for (int i = n - 1; ~i; i--) {
66
                                                                                                 116
67
       BigNum& operator %=(const BigNum& r) {return *this = *this % r;}
                                                                                                 117
                                                                                                              int tag = 0;
       BigNum& operator %=(int x) {return *this = *this % x;}
                                                                                                              (s *= MOD) += u[i]:
68
                                                                                                 118
69
       BigNum operator -() const {
                                                                                                 119
                                                                                                              for (int x = m - 1; x; x--) if (s >= e[x]) s -= e[x], tag |= 1 << x;
```

```
(c *= MOD) += tag:
120
121
122
          r.o ^= ro, o ^= lo;
          if (c.u[0] | | c.n >= 2) c.o = r.o ^o:
123
124
          return c;
        }
125
        BigNum operator %(const BigNum& r) const {
126
127
          BigNum e[35], s = 0;
          int m = 0, ro = r.o, lo = o;
128
129
          r.o ^= ro, o ^= lo;
130
          for (e[m] = r; MOD >> ++m; e[m] = e[m - 1] + e[m - 1]);
          for (int i = n - 1; "i; i--) {
131
            (s *= MOD) += u[i]:
132
133
           for (int x = m - 1; x; x--) if (s >= e[x]) s -= e[x];
134
          r.o ^= ro, o ^= lo;
135
          if (s.u[0] || s.n >= 2) s.o = o;
136
137
          return s:
138
        }
        BigNum& operator *=(int x) {
139
          if (!x) return *this = 0;
140
          if (x < 0) o = 1, x = -x;
141
          for (int i = 0; i < n; i++) u[i] *= x;
142
          for (int i = 0; i < n; i++) if (u[i] >= MOD) {
143
            u[i + 1] += u[i] / MOD;
144
145
           u[i] %= MOD;
146
            if (i == n - 1) n++;
147
          }
          if (!u[0] && n == 1) o = 0;
148
          return *this;
149
        }
150
151
        BigNum& operator /=(int x) {
152
          if (x < 0) o = 1, x = -x;
          for (int i = n - 1; i; u[i--] /= x) u[i - 1] += u[i] % x * MOD;
153
154
          for (u[0] /= x; n >= 2; n--) if (u[n - 1]) break;
          if (!u[0] \&\& n == 1) o = 0;
155
156
          return *this:
157
        }
        int operator %(int x) const {
158
159
          long long c = 0;
160
          for (int i = n - 1; "i; i--) c = (c * MOD + u[i]) % x;
          return (1 - o - o) * int(c);
161
        }
162
163 };
```

线性逆元

```
1 // 返回 [0, n] 每个数关于 m 的逆元
2 // x * i = 1 mod m
3 // let m = i * k + r, where k = m / i, r = m % i
```

```
4 // i = (m - r) * k^-1
5 // x = -k * r^{-1}
6 // time complexity : O(n)
   std::vector<LL> linear_inversion(int n, LL m) {
     std::vector<LL> inv(n + 1, 0):
   inv[1] = 1:
10
11
   for (int i = 2; i <= n; ++i) {
12
      inv[i] = (m - m / i) * inv[m % i] % m;
13
    }
14
    return inv;
15 }
```

勒让德定理

欧拉函数

```
2 * 对正整数 n .
3 * 欧拉函数是少于或等于n的数中与n互质的数的数目
4 * P是素数:
   * 若p是x的约数 ,则E(xp) = E(x)p
   * p不是x的约数 ,则E(xp) = E(x)E(p) = E(x)(p-1)
  const int N = 1e5:
  int prime[N];
10 int phi[N];
  bool is_prime[N];
12
13 void get_phi() {
   int i, j, k;
   k = 0:
15
16
   for(i = 2; i < N; i++) {
     if(is_prime[i] == false) {
17
18
       prime[k++] = i;
```

```
phi[i] = i - 1:
19
20
       }
21
       for(j = 0; j < k && i * prime[j] < N; j++) {</pre>
         is_prime[ i * prime[j] ] = true;
22
23
         if( i % prime[j] == 0) {
           phi[ i * prime[j] ] = phi[i] * prime[j];
24
25
           break:
26
         }
         else {
27
           phi[ i * prime[j] ] = phi[i] * ( prime[j] - 1 );
28
29
       }
30
31
32 }
```

行列式求值

```
* 行列式 Mod
   * Kvb 待验证
   */
   #include <algorithm>
   const int MOD = (int)1e9 + 7;
   inline void add(int &a.int b) {
9
     a += b:
     if (a >= MOD) a -= MOD;
10
11 }
12
   int det(int A[16][16],int n) {
     int ret = 1:
13
     for (int i = 0: i < n: ++ i) {
14
15
       if (A[i][i] == 0) {
16
         for (int j = i + 1; j < n; ++ j) {
           if (A[i][i]) {
17
18
             for (int k = i; k < n; ++ k) {
19
               std::swap(A[i][k],A[j][k]);
20
             }
21
             ret = -ret;
22
             break;
23
           }
24
25
         if (A[i][i] == 0) return 0;
26
       }
27
       for (int j = i + 1; j < n; ++ j) {
28
         int a = 1, b = 0, c = 0, d = 1;
29
         int x = A[i][i], y = A[j][i];
30
         while (y) {
31
          int t = x / y;
32
           if (t < 0) t += MOD:
33
           add(a,MOD - c * 111 * t % MOD);
```

```
34
           add(b.MOD - d * 111 * t % MOD):
35
           std::swap(a,c);
36
           std::swap(b,d);
37
           x %= y;
38
           std::swap(x,y);
           ret = -ret:
39
40
41
         for (int k = 0; k < n; ++ k) {
           int q = A[i][k], w = A[j][k];
42
           A[i][k] = (a * 111 * q + b * 111 * w) % MOD;
43
44
           A[j][k] = (c * 111 * q + d * 111 * w) % MOD;
45
46
47
       ret = A[i][i] * 111 * ret % MOD;
48
     if (ret < 0) {
49
50
       ret += MOD:
51
52
     return ret;
53
```

生成树计数

```
* 算法引入:
  * 给定一个无向图 G , 求它生成树的个数 t(G);
  * 算法思想:
  * 1) G 的度数矩阵 D[G] 是一个 n*n 的矩阵 ,
  * 并且满足 : 当 i != j 时 , dij=0 ; 当 i=j 时 , dij 等于 vi 的度数 ;
  * 2) G 的邻接矩阵 A[G] 是一个 n*n 的矩阵 ,
  * 并且满足 : 如果 vi , vj 之间有边直接相连 , 则 aij=1 , 否则为 0 ;
  * 定义图 G 的 Kirchhoff 矩阵 C[G] 为 C[G]=D[G]-A[G]:
  * Matrix-Tree 定理 : G 的所有不同的生成树的个数等于
  * 其 Kirchhoff 矩阵 C[G] 任何一个 n-1 阶主子式的行列式的绝对值 ;
  * 所谓 n-1 阶主子式 , 就是对于 r(1 <= r <= n) , 将 C[G] 的第r行
14
  * 第 r 列同时去掉后得到的新矩阵 , 用 Cr[G] 表示 ;
15
  * Kirchhoff 矩阵的特殊性质:
  * 1) 对于任何一个图 G, 它的 Kirchhoff 矩阵 C 的行列式总是 O,
  * 这是因为 C 每行每列所有元素的和均为 O;
  * 2) 如果 G 是不连通的 , 则它的 Kirchhoff 矩阵 C 的任一个主子式的行列式均为 O ;
  * 3) 如果 G 是一颗树 , 那么它的 Kirchhoff 矩阵 C 的任一个 n-1 阶主子式的行列式均为 1;
20
21
22
  * 算法举例:
23
  * SPOJ HIGH
24
  * 题目地址:
  * http://www.spoj.com/problems/HIGH/
```

```
27
28
   * 题目大意:
   * 一个有 n 座城市的组成国家 , 城市 1 至 n 编号 , 其中一些城市之间可以修建高速公路 ;
29
   * 需要有选择的修建一些高速公路 , 从而组成一个交通网络 ;
30
   * 计算有多少种方案 , 使得任意两座城市之间恰好只有一条路径 ;
31
32
33
   #include <bits/stdc++.h>
34
   const int N=15;
35
   typedef long long 11;
36
37
  int degree[N];
38
39 11 C[N][N];
40 // 生成树计数 : Matrix-Tree 定理
41 | 11 det(11 a[][N], int n) {
     11 \text{ ret} = 1;
42
     for(int i=1; i<n; i++) {</pre>
43
44
      for(int j=i+1; j<n; j++)</pre>
45
        while(a[j][i]) {
          11 t = a[i][i] / a[j][i];
46
          for(int k = i; k < n; k++)
47
            a[i][k] = (a[i][k] - a[j][k] * t);
48
          for(int k = i; k < n; k++)
49
50
            std::swap(a[i][k], a[j][k]);
51
          ret = -ret;
52
53
       if(a[i][i] == 0)
54
        return 0;
       ret = ret * a[i][i];
55
56
57
     if(ret < 0)
58
      ret = -ret:
59
     return ret;
60
61
   int main() {
62
63
     int tcase:
     scanf("%d",&tcase);
64
     while(tcase--) {
65
66
       memset(degree, 0, sizeof(degree) );
67
       memset(C, 0, sizeof(C));
68
       int n ,m;
69
       scanf("%d%d", &n, &m);
70
       int u, v;
71
       while(m--) {
72
        scanf("%d%d", &u, &v);
73
        u--; v--;
74
        C[u][v] = C[v][u] = -1;
75
         degree[u]++; degree[v]++;
76
```

```
for(int i = 0; i < n; ++i)

C[i][i] = degree[i];

printf("%lld\n", det(C, n));

return 0;

}</pre>
```

Simpson 积分

```
* 自适应 Simpson
    * 验题:hdu4978
    * Kyb 待验证
   #include <cmath>
   const double eps = 1e-6;
   double f(double x) { return x; }
   long double simpson(long double a, long double b) {
10
     long double c = a + (b - a) / 2;
11
     return (f(a) + 4 * f(c) + f(b)) * (b - a) / 6:
12
13
   long double asr(long double a, long double b, long double eps, long double A) {
15
     long double c = a + (b - a) / 2;
16
     long double L = simpson(a, c), R = simpson(c, b);
     if (std::fabs(L + R - A) < 15 * eps) return L + R + (L + R - A) / 15;
17
     return asr(a, c, eps / 2, L) + asr(c, b, eps / 2, R);
18
19
   long double asr(long double a, long double b, long double eps) {
     return asr(a, b, eps, simpson(a, b));
22
```

常用结论

vimrc

```
syntax on
set nu
set cindent
set shiftwidth=2
set tabstop=2
set expandtab
```

矩阵乘法

for(i)for(j)if(A[i][j])for(k)C[i][k] += A[i][j]*B[j][k]

四边形体积公式

$$\begin{array}{l} (12V)^2 = a^2d^2(b^2+c^2+e^2+f^2-a^2-d^2) + b^2e^2(c^2+a^2+f^2+d^2-b^2-e^2) + c^2f^2(a^2+b^2+d^2+e^2-c^2-f^2) - a^2b^2c^2-a^2e^2f^2-d^2b^2f^2-d^2e^2c^2 \end{array}$$

卡特兰数

Cat(n) = Comb(2n, n)/(n+1) = Comb(2n, n) - Comb(2n, n+1) = Cat(n-1) * (4n-2)/(n+1)

- n 节点二叉树个数 Cat(n)
- 正 n 边形划分为 n-2 个三角形的种数 Cat(n-2)
- n 个矩阵连乘括号化种数 Cat(n-1)
- n 个元素入栈的出栈顺序种数 Cat(n)
- n 对括号的合法括号序列个数 Cat(n)

牛顿迭代法

$$x_m = \frac{(xf'(x) - f(x))}{f'(x)}$$

康托展开

 $X = a_n n! + a_{n-1}(n-1)! + ... + a_2 \cdot 2! + a_1 \cdot 1!$ a_i 表示 i 开头的后缀中逆序对的个数

高阶等差数列

$$\sum n = \frac{1}{2}n(n+1) \tag{1}$$

$$\sum n^2 = \frac{1}{6}n(n+1)(2n+1) \tag{2}$$

$$\sum n^3 = \left(\sum n\right)^2 \tag{3}$$

$$\sum n^4 = \left(\sum n^2\right) \frac{1}{5} (3n^2 + 3n - 1) \tag{4}$$

$$\sum n^5 = \left(\sum n\right)^2 \frac{1}{3} (2n^2 + 2n - 1) \tag{5}$$

$$\sum n^6 = \left(\sum n^2\right) \frac{1}{7} (3n^4 + 6n^3 - 3n + 1) \tag{6}$$

$$\sum n^7 = \left(\sum n\right)^2 \frac{1}{6} (3n^4 + 6n^3 - n^2 - 4n + 2) \tag{7}$$

$$\sum n^8 = \left(\sum n^2\right) \frac{1}{15} \left(5n^6 + 15n^5 + 5n^4 - 15n^3 - n^2 + 9n - 3\right) \tag{8}$$

割建图

对于图G=(V,E)中的一个点覆盖是一个集合S V使得每一条边至少有一个端点在S中 . 对于二分图 ,

最小路径覆盖 = |P| - 最大匹配数 .

最小点覆盖数 = 最大匹配数、

最大独立顶点集 = 总顶点数 - 最大匹配数,

哈密尔顿判定

对于n >= 3个点的图G, 如果对于任意u, v 都有 deg(u) + deg(v) >= n,则G一定是哈密尔顿图。

弦图

设 next(v) 表示 N(v) 中最前的点 . 令 w* 表示所有满足 $A\in B$ 的 w 中最后的一个点 ,判断 $v\cup N(v)$ 是否为极大团 ,只需判断是否存在一个 $w\in w*$,满足 Next(w)=v 且 $|N(v)|+1\le |N(w)|$ 即可 .

五边形数

$$\prod_{n=1}^{\infty} (1 - x^n) = \sum_{n=0}^{\infty} (-1)^n (1 - x^{2n+1}) x^{n(3n+1)/2}$$
$$f(n) = \frac{n(3n-1)}{2}$$

重心

半径为 r ,圆心角为 θ 的扇形重心与圆心的距离为 $\frac{4r\sin^3(\theta/2)}{3\theta}$ 半径为 r ,圆心角为 θ 的圆弧重心与圆心的距离为 $\frac{4r\sin^3(\theta/2)}{3(\theta-\sin(\theta))}$

第二类 Bernoulli number

$$B_m = 1 - \sum_{k=0}^{m-1} {m \choose k} \frac{B_k}{m-k+1}$$

$$S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m {m+1 \choose k} B_k n^{m+1-k}$$

Stirling 数

第一类:n 个元素的项目分作 k 个环排列的方法数目

$$\begin{aligned} s(n,k) &= (-1)^{n+k} |s(n,k)| \\ |s(n,0)| &= 0 \\ |s(1,1)| &= 1 \\ |s(n,k)| &= |s(n-1,k-1)| + (n-1) * |s(n-1,k)| \end{aligned}$$

第二类:n 个元素的集定义 k 个等价类的方法数

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = S(n-1,k-1) + k * S(n-1,k)$$

数据范围

三角公式

$$\begin{split} & \sin(a \pm b) = \sin a \cos b \pm \cos a \sin b & \cos(a \pm b) = \cos a \cos b \mp \sin a \sin b \\ & \tan(a \pm b) = \frac{\tan(a) \pm \tan(b)}{1 \mp \tan(a) \tan(b)} & \tan(a) \pm \tan(b) = \frac{\sin(a \pm b)}{\cos(a) \cos(b)} \\ & \sin(a) + \sin(b) = 2 \sin(\frac{a + b}{2}) \cos(\frac{a - b}{2}) & \sin(a) - \sin(b) = 2 \cos(\frac{a + b}{2}) \sin(\frac{a - b}{2}) \\ & \cos(a) + \cos(b) = 2 \cos(\frac{a + b}{2}) \cos(\frac{a - b}{2}) & \cos(a) - \cos(b) = -2 \sin(\frac{a + b}{2}) \sin(\frac{a - b}{2}) \\ & \sin(na) = n \cos^{n-1} a \sin a - \binom{n}{3} \cos^{n-3} a \sin^3 a + \binom{n}{5} \cos^{n-5} a \sin^5 a - \dots \\ & \cos(na) = \cos^n a - \binom{n}{2} \cos^{n-2} a \sin^2 a + \binom{n}{4} \cos^{n-4} a \sin^4 a - \dots \end{split}$$

积分表

$$\int \frac{1}{1+x^2} dx = \tan^{-1} x \int \frac{1}{a^2+x^2} dx = \frac{1}{a} \tan^{-1} \frac{x}{a}$$
$$\int \frac{x}{a^2+x^2} dx = \frac{1}{2} \ln|a^2 + x^2| \int \frac{x^2}{a^2+x^2} dx = x - a \tan^{-1} \frac{x}{a}$$

$$\int \sqrt{x^2 \pm a^2} dx = \frac{1}{2} x \sqrt{x^2 \pm a^2} \pm \frac{1}{2} a^2 \ln \left| x + \sqrt{x^2 \pm a^2} \right|$$

$$\int \sqrt{a^2 - x^2} dx = \frac{1}{2} x \sqrt{a^2 - x^2} + \frac{1}{2} a^2 \tan^{-1} \frac{x}{\sqrt{a^2 - x^2}}$$

$$\int \frac{x^2}{\sqrt{x^2 \pm a^2}} dx = \frac{1}{2} x \sqrt{x^2 \pm a^2} \mp \frac{1}{2} a^2 \ln \left| x + \sqrt{x^2 \pm a^2} \right|$$

$$\int \frac{1}{\sqrt{x^2 \pm a^2}} dx = \ln \left| x + \sqrt{x^2 \pm a^2} \right|$$

$$\int \frac{1}{\sqrt{a^2 - x^2}} dx = \sin^{-1} \frac{x}{a} \int \frac{x}{\sqrt{x^2 \pm a^2}} dx = \sqrt{x^2 \pm a^2} \int \frac{x}{\sqrt{a^2 - x^2}} dx = -\sqrt{a^2 - x^2}$$

$$\int \sqrt{ax^2 + bx + c} dx = \frac{b + 2ax}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8a^{3/2}} \ln \left| 2ax + b + 2\sqrt{a(ax^2 + bx + c)} \right|$$

$$\int x^n e^{ax} dx = \frac{x^n e^{ax}}{a} - \frac{n}{a} \int x^{n-1} e^{ax} dx$$

$$\int \sin^2 ax dx = \frac{x}{2} - \frac{1}{4a} \sin 2ax \int \sin^3 ax dx = -\frac{3\cos ax}{4a} + \frac{\cos 3ax}{12a}$$

$$\int \cos^2 ax dx = \frac{x}{2} + \frac{\sin 2ax}{4a} \int \cos^3 ax dx = \frac{3\sin ax}{4a} + \frac{\sin 3ax}{12a}$$

$$\int \tan ax dx = -\frac{1}{a} \ln \cos ax \int \tan^2 ax dx = -x + \frac{1}{a} \tan ax$$

$$\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{x}{a} \sin ax \int x^2 \cos ax dx = \frac{2x \cos ax}{a^2} + \frac{a^2 x^2 - 2}{a^3} \sin ax$$

$$\int x \sin ax dx = -\frac{x \cos ax}{a} + \frac{\sin ax}{a^2} \int x^2 \sin ax dx = \frac{2-a^2 x^2}{a^3} \cos ax + \frac{2x \sin ax}{a^2}$$