

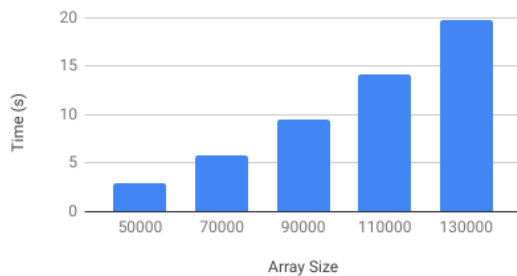
Evan Trout

EECS 268 Lab 7 Report

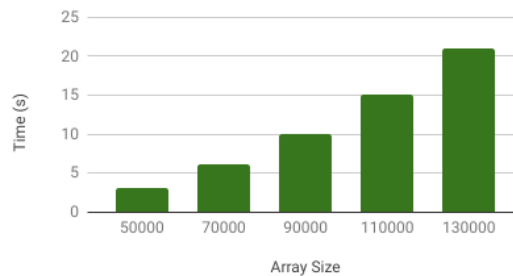
Selection Sort

Size	Ascending (s)	Descending (s)	Random (s)
50000	2.92080	3.14923	3.26441
70000	5.79518	6.11749	6.41747
90000	9.51462	10.0613	10.5619
110000	14.1941	14.9978	15.7432
130000	19.8115	20.9188	22.0235

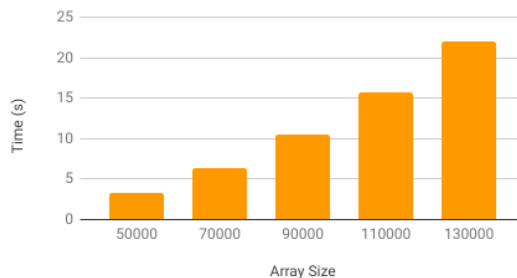
Selection Sort - Ascending



Selection Sort - Descending



Selection Sort - Random



Analysis of the curves for all three cases appear to indicate that selection sort has $O(n^2)$ behavior.

Based on this assumption, the time to sort an an array of size 10,000,000 can be estimated by solving the ratio $t_1/t_2 = O_1/O_2$:

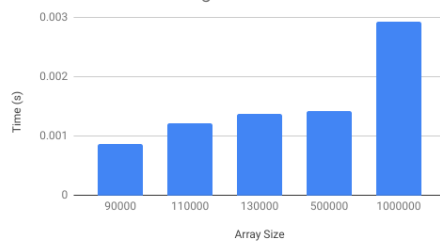
Size	Ascending (s)	Descending (s)	Random (s)
10000000	116832	125969	130576.4

Bubble Sort

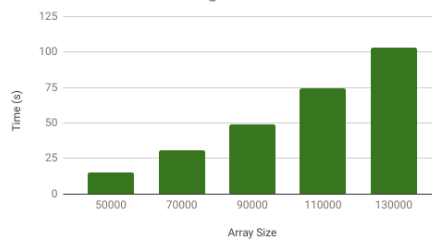
Size	Ascending (s)
90000	0.000874
110000	0.001222
130000	0.001373
500000	0.001418
1000000	0.002922

Size	Descending (s)	Random (s)
50000	15.269	14.0136
70000	30.4555	26.8356
90000	49.2996	44.3323
110000	74.53	66.1054
130000	102.999	93.0703

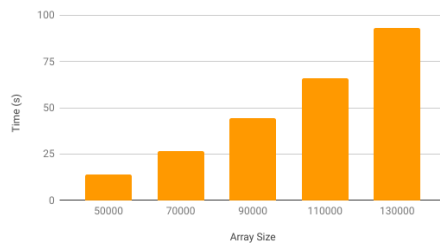
Bubble Sort - Ascending



Bubble Sort - Descending



Bubble Sort - Random



Analysis of the curves appears to indicate that the ascending case has behavior of $O(n)$, while the descending and random cases have behavior of $O(n^2)$.

Based on this assumption, the time to sort an an array of size 10,000,000 can be estimated by solving the ratio $t_1/t_2 = O_1/O_2$:

Size	Ascending (s)	Descending (s)	Random (s)
10000000	0.02836	610760	560544

Insertion Sort

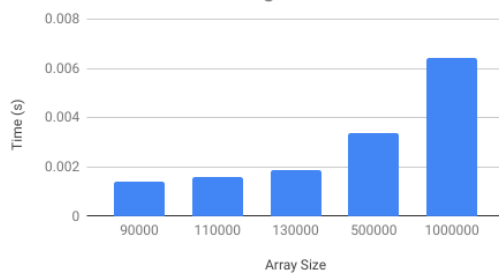
Size Ascending (s)

90000	0.001425
110000	0.001612
130000	0.001866
500000	0.003396
1000000	0.006408

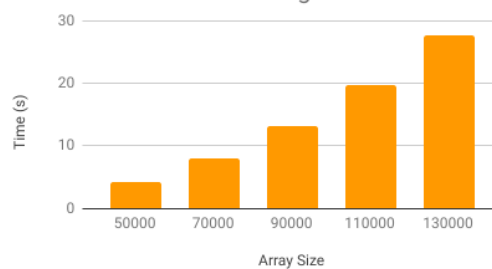
Size Descending (s) Random (s)

50000	4.14375	2.10959
70000	8.00778	4.0564
90000	13.1417	6.66614
110000	19.671	9.88171
130000	27.7176	13.8636

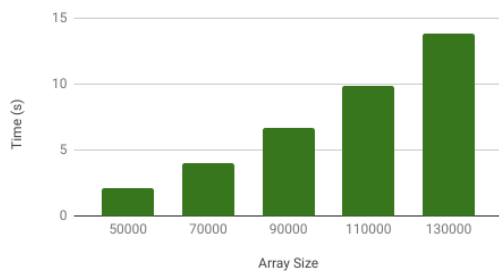
Insertion Sort - Ascending



Insertion Sort - Descending



Insertion Sort - Random



Analysis of the curves appears to indicate that the ascending case has behavior of $O(n)$, while the descending and random cases have behavior of $O(n^2)$.

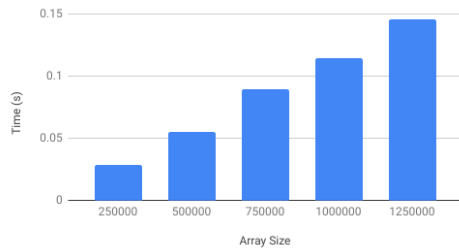
Based on this assumption, the time to sort an array of size 10,000,000 can be estimated by solving the ratio $t_1/t_2 = O_1/O_2$:

Size	Ascending (s)	Descending (s)	Random (s)
10000000	.03732	165750	84383

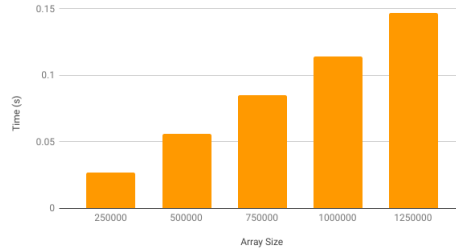
Merge Sort

Size	Ascending (s)	Descending (s)	Random (s)
250000	0.028821	0.026881	0.056024
500000	0.055499	0.05605	0.107812
750000	0.089223	0.084968	0.165064
1000000	0.114636	0.114441	0.223062
1250000	0.146109	0.146994	0.281056

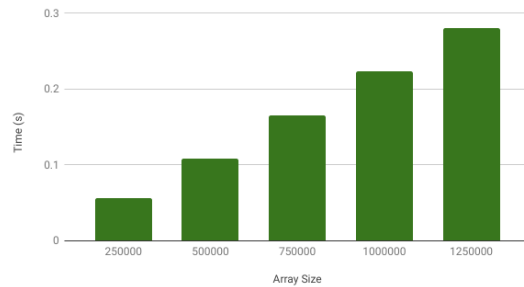
Merge Sort - Ascending



Merge Sort - Descending



Merge Sort - Random



Analysis of the data appears to indicate that all three cases have behavior of $O(n \cdot \log(n))$.

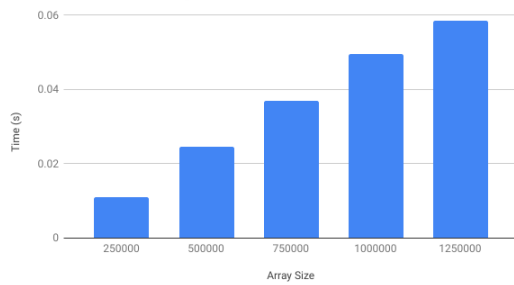
Based on this assumption, the time to sort an array of size 10,000,000 can be estimated by solving the ratio $t_1/t_2 = O_1/O_2$:

Size	Ascending (s)	Descending (s)	Random (s)
10000000	1.33742	1.335145	2.60239

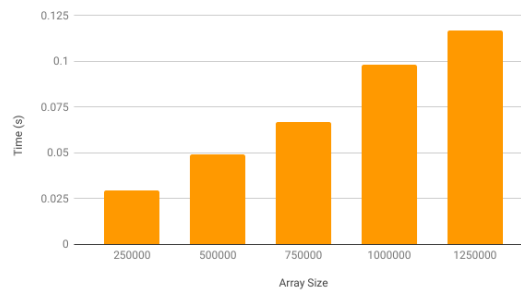
Quick Sort

Size	Ascending (s)	Descending (s)	Random (s)
250000	0.010942	0.029608	0.043318
500000	0.0245	0.049031	0.089779
750000	0.036879	0.066811	0.135556
1000000	0.049661	0.098361	0.184359
1250000	0.058512	0.117093	0.233724

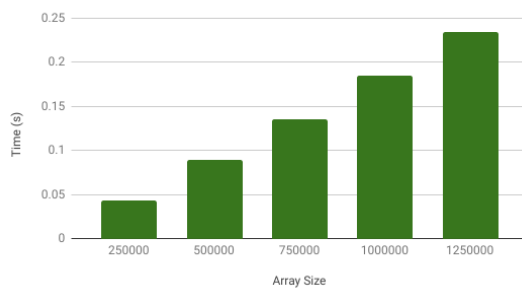
Quick Sort - Ascending



Quick Sort - Descending



Quick Sort - Random



Analysis of the data appears to indicate that all three cases demonstrate $O(n \cdot \log(n))$ behavior.

Based on this assumption, the time to sort an array of size 10,000,000 can be estimated by solving the ratio $t_1/t_2 = O_1/O_2$:

Size	Ascending (s)	Descending (s)	Random (s)
10000000	0.57938	1.1475	2.15086