

Efficient Inference in Probabilistic Computing

M.Eng Proposal

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Abstract

1 Introduction

Probabilistic programming languages are generalizations of programming languages, in which procedures are replaced with random procedures that induce distributions. Probabilistic computing thus allows for easy description and manipulation of probability distributions, letting one describe classical AI models in compact ways.

A core operation of probabilistic programming languages is inference, which is, in general, a difficult computational problem [?]. However, Markov chain Monte Carlo (MCMC) methods converge quite quickly to the posterior distribution for a large class of inference problems. Much effort has been devoted to studying in which cases this inference is efficient, and how to make it so.

2 Language description

2.1 Values

A **value** is a valid binding for a variable, in our language. For example, integers, floats, and booleans are valid values. We call these primitive values. Another type of value is a **procedure** — a function which gets applied to arguments and returns (not necessarily primitive) values.

These procedures should be non-deterministic. To allow for this, we can imagine that there is a special procedure which takes no arguments, and outputs 0 or 1, each with probability $\frac{1}{2}$. Arbitrary random procedures can then be built using this one. We will see later that we actually can and will allow for something much more general than this.

2.2 Expressions and Environments

Expressions are syntactic building blocks, which specified as being one of the following:

- Atomic expressions:
 - A value itself.
 - A variable name.
- Non-atomic expressions
 - An application of an expression to a list of argument expressions.
 - A function, consisting of some variable arguments, and a body expression.
 - An if statement, consisting of a branching expression, and two child expressions.
 - An operator statement, consisting of an operator (e.g. `==`, `<`, `+`, `×`, `AND`, `OR`, `NOT`) and a list of argument expressions.

Expressions are **evaluated**, so that they take on some value. Their evaluation is relative to an **environment**, in which variable names may be bound to values.

Evaluation happens recursively. If an invalid expression is ever given (e.g. the first expression in an application isn't a procedure, or we call the equality operator with three arguments), an error is thrown.

2.3 Directives

There are four primitive operations, called **directives**, which the language supports. They are:

1. `sample(expr)` :
Evaluates with respect to the global environment to sample a value for the expression `expr`.
2. `assume(name, expr)` :
Samples a value for the expression `expr`, and then binds `name`, a string variable name, to the resulting value in the global environment.
3. `observe(expr, val)` :
Observes the expression `expr` to have been evaluated to the value `val`.

4. `infer()` :

Runs a single step of the Markov Chain.

The result of calling `infer()` sufficiently many times, should be that sampling gives the posterior distribution for expressions, conditioned on all the observed values.

3 Inference

Inference is by far the most complicated of the directives, and also the most important. Performing inference is the reason we are interested in probabilistic programming languages in the first place.

3.1 High-level description

We now describe the way in which inference is implemented, at a high level. The idea is to run Metropolis-Hastings on the space of all possible executions of the program.

Our proposal density, which chooses the new state (program execution) to potentially transition to, does the following:

1. Pick, at random, one of the random choices made, at any point in the program's execution history. Make a new choice for this randomness.
2. Rerun the entire program, changing only this choice. If this causes us to evaluate new sections of the code (a different side of an if statement, or a procedure called with different arguments), simply run these new sections.

We then use the update rule of Metropolis-Hastings to decide whether we enter this new state, or whether we keep the old one. If the new execution history doesn't agree with our observations, we assign it a probability of 0, so we never transition to that state.

Unfortunately, as I've described it, this random walk doesn't converge to the posterior. Consider the following example, in which I flip two weighted coins and observe that they came up differently:

```
>>> p, q = 0.6, 0.4
>>> assume('a', bernoulli(p))
>>> assume('b', bernoulli(q))
>>> assume('c', var('a') ^ var('b'))
>>> observe(bernoulli_noise('c', noise_level), True)
```

Notice that if there is no noise in our observations, this algorithm does not necessarily mix in finite time.

Consider the following example:

Some executions are more likely than others, and so
Here pseudocode for a single iteration:

```
stack = globals.db.random_stack()
(xrp, val, prob, args) = globals.db.get(stack)
```

```

old_p = globals.db.prob()
old_to_new_q = - math.log(globals.db.count)

globals.db.save()

globals.db.remove(stack)
new_val = xrp.apply(args)

if val == new_val:
    globals.db.insert(stack, xrp, new_val, args)
    return
globals.db.insert(stack, xrp, new_val, args)

rerun(False)
new_p = globals.db.prob()
new_to_old_q = -math.log(globals.db.count)
old_to_new_q += globals.db.eval_p
new_to_old_q += globals.db.uneval_p
if old_p * old_to_new_q > 0:
    p = random.random()
    if new_p + new_to_old_q - old_p - old_to_new_q < math.log(p):
        globals.db.restore()
globals.db.save()

```

3.2 XRPs

Our language allows for the use of something more general than random procedures, while retaining the ability to perform inference. It allows for what we call an **exchangeable random procedure (XRP)**.

An XRP is like a random procedure, except that different applications of it are not guaranteed to be independent. An XRP is allowed to maintain an internal state, which can be modified. However, an XRP must maintain the exchangeability condition, which says that applications of the XRP gives an exchangeable sequence of random variables.

The role of XRPs will be motivated and discussed in detail later. For now, it may be simpler to just imagine that values are either primitive values or (possibly random) procedures.

The exchangeable sequence

4 Test Cases

In order to ensure the language was working properly, a suite of test cases was developed.

One of the primary techniques used for testing was “following the prior”, using the function `follow_prior(variable, niters, burnin)`. To follow the prior, we draw `niters` samples from the prior, and run each of them `burnin` steps in the random walk. We then view the resulting distribution on `variable`.

Here are some of the more illustrative and interesting test cases.

4.1 Testing a tricky coin

Consider the following scenario: There is a coin that is fair with probability 50%. The rest of the time, its weight is drawn uniformly from the interval $[0, 1]$. The coin is flipped, and we observe it to be heads `nheads` times. We then infer the posterior probability that the coin was fair.

```
>>> noise_level = .001
>>> nheads = 1
>>>
>>> assume('weight', beta(1, 1))
>>> assume('tricky-coin', function([], bernoulli('weight')))
>>> assume('fair-coin', function([], bernoulli(0.5)))
>>> assume('is-fair', bernoulli(0.5))
>>> assume('coin', ifelse('is-fair', 'fair-coin', 'tricky-coin'))
>>>
>>> for i in xrange(nheads):
>>>     observe(bernoulli_noise(apply('coin'), noise_level), True)
>>>
>>> follow_prior('is-fair', 10000, 1000)
```

We should be able to predict the results of running this program for different values of `nheads`. First we compute the probability that the coin comes up heads n times, given that it is tricky. The probability is simply $\beta(1, n+1) = \frac{1}{n+1}$. Thus the probability the coin is fair (not tricky), given it came up heads n times, is, by Baye’s law

$$\frac{\frac{1}{2} \cdot \frac{1}{2^n}}{\frac{1}{2} \cdot \frac{1}{n+1} + \frac{1}{2} \cdot \frac{1}{2^n}} = \frac{n+1}{2^n + n+1}.$$

Here are our results, when running `follow_prior('is-fair', 10000, 1000)`, for the percentage of times the coin was fair:

nheads	Predicted	Actual
0	0.5	0.4982
1	0.5	0.4934
2	0.4286	0.4289
3	0.3333	0.3346
4	0.2381	0.2712
5	0.1579	0.2528

As `nheads` grows, we should expect the mixing time to grow. This is because

4.2 Bayes nets

We first test that inference works, by considering a number of test cases. A classic inference problem which is well understood, is inference in Bayesian networks. Bayesian networks are incredibly easy to describe in our language. Here are some simple examples of inference giving the correct answer in a Bayesian network.

4.2.1 Sprinkler net

Let's start with a simple example, to get familiar with our language. Here's a definition of a bayesian network with just two nodes.

```
>>> assume('cloudy', bernoulli(0.5))
>>> assume('sprinkler', ifelse('cloudy', bernoulli(0.1), bernoulli(0.5)))
```

We then observe that the sprinkler is on. Worlds in which the sprinkler is on are weighted as 100 times more likely than ones in which the sprinkler is off.

```
>>> noise_level = .01
>>> sprinkler_ob = observe(bernoulli_noise('sprinkler', noise_level), True)
```

Now, let's try inferring the weather. We follow the prior 10000 times, going 50 steps each time.

```
>>> follow_prior('cloudy', 10000, 50)}
{False: 0.820100000000000005, True: 0.1799}
```

This is close to the value of $\frac{5}{6} = 0.8\overline{333}$ False, which is what we'd get if there was noise in our observation. However, there is still a number of worlds in which the sprinkler was actually on. In those worlds, the weather was 50/50. Thus we should've expect the answer to be on the order of 0.01 lower.

Now, let's suppose we didn't observe the sprinkler being on, after all.

```
>>> forget(sprinkler_ob)
```

So it should be the case that it's cloudy exactly half the time.

```
>>> follow_prior('cloudy', 10000, 50)}
{False: 0.50429999999999997, True: 0.49569999999999997}
```

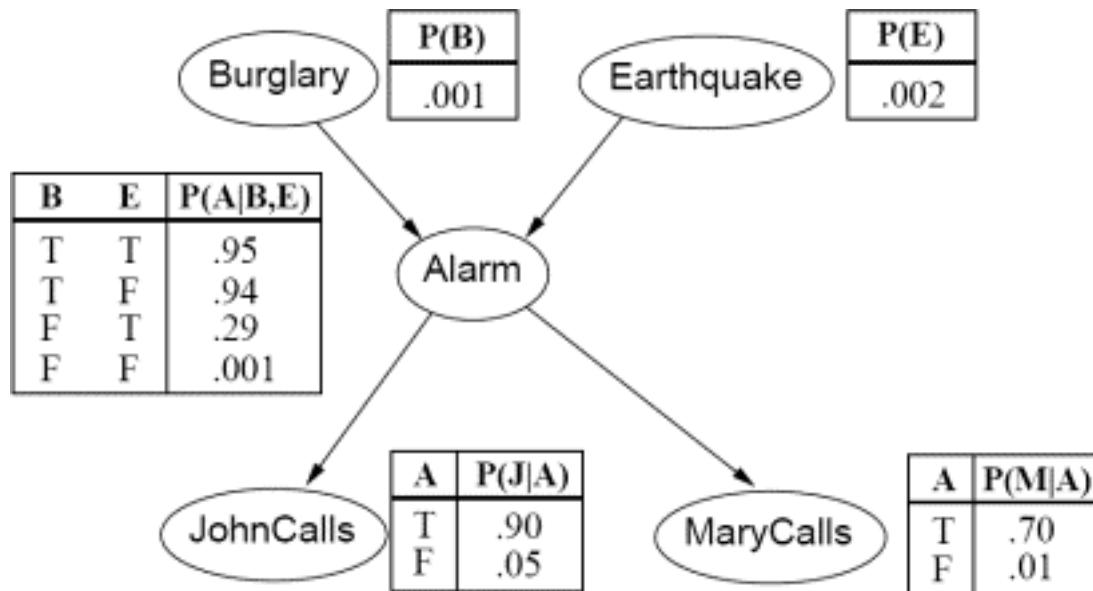
Now, we re-observe the sprinkler to be on. This time, we are much more sure.

```
>>> noise_level = .001
>>> sprinkler_ob = observe(bernoulli_noise('sprinkler', noise_level), True)
>>> follow_prior('cloudy', 10000, 50)}
{False: 0.83999999999999997, True: 0.16}
```

We see that our answer is even closer to the value of $\frac{5}{6} = 0.8\overline{333}$ False, now.

4.2.2 Alarm net

Here's a more complicated Bayesian network, which is given as an example in the classic AI textbook Artificial Intelligence (A Modern Approach).



Let's define this Bayesian network.

```
>>> assume('burglary', bernoulli(0.001))
>>> assume('earthquake', bernoulli(0.002))
>>> assume('alarm', ifelse('burglary', ifelse('earthquake', bernoulli(0.95), bernoulli(0.94)), \
...                               ifelse('earthquake', bernoulli(0.29), bernoulli(0.001))))
>>> assume('johnCalls', ifelse('alarm', bernoulli(0.9), bernoulli(0.05)))
>>> assume('maryCalls', ifelse('alarm', bernoulli(0.7), bernoulli(0.01)))
```

Let's try a couple inference problems.

```
>>> follow_prior('alarm', 1000, 100) # should give 0.002516 True
{False: 0.9973999999999999, True: 0.002599999999999999}
>>>
>>> noise_level = .001
>>> mary_ob = observe(bernoulli_noise('maryCalls', noise_level), True)
>>> follow_prior('johnCalls', 1000, 100) # should give 0.177577 True
{False: 0.9162000000000001, True: 0.08379999999999999}
>>>
>>> forget(mary_ob)
>>> burglary_ob = observe(bernoulli_noise(negation('burglary'), noise_level), True)
>>> follow_prior('johnCalls', 1000, 100) # should give 0.051343 True
{False: 0.9471000000000005, True: 0.05290000000000003}
```

The first and third inferences were on the mark, but the second one was off by a factor of two! The explanation is that Mary calls very rarely calls. Indeed, she only calls about 0.01 of the time, since the alarm almost never sounds. This is rare enough that our observation that she called is

reasonably likely to be wrong. In worlds where she doesn't call, John also tends not to call, thus accounting for the large dip.

To verify that this is the explanation, we can alter the probabilities so that Mary calling is more likely. Here is an example of this being done.

```
>>> reset()
>>>
>>> assume('burglary', bernoulli(0.1))
>>> assume('earthquake', bernoulli(0.2))
>>> assume('alarm', ifelse('burglary', ifelse('earthquake', bernoulli(0.95), bernoulli(0.94)), \
...                               ifelse('earthquake', bernoulli(0.29), bernoulli(0.10))))
>>> assume('johnCalls', ifelse('alarm', bernoulli(0.9), bernoulli(0.5)))
>>> assume('maryCalls', ifelse('alarm', bernoulli(0.7), bernoulli(0.1)))
>>>
>>> follow_prior('alarm', 1000, 100) # should give 0.218400 True
{False: 0.7795999999999996, True: 0.22040000000000001}
>>>
>>> noise_level = .001
>>> mary_ob = observe(bernoulli_noise('maryCalls', noise_level), True)
>>> follow_prior('johnCalls', 1000, 100) # should give 0.764681 True
{False: 0.2432, True: 0.75680000000000003}
>>>
>>> forget(mary_ob)
>>> burglary_ob = observe(bernoulli_noise(negation('burglary'), noise_level), True)
>>> follow_prior('johnCalls', 1000, 100) # should give 0.561333 True
{False: 0.44290000000000002, True: 0.55710000000000004}
```

Success!

4.3 Testing xor

Consider the following program, where we simply draw two booleans `a` and `b`, and observe that their xor is `True`:

```
>>> p, q = 0.6, 0.4
>>> noise_level = .01
>>>
>>> assume('a', bernoulli(p))
>>> assume('b', bernoulli(q))
>>> assume('c', var('a') ^ var('b'))
>>>
>>> follow_prior('a', 10000, 100) # should be 0.60 True
{False: 0.3977, True: 0.6022999999999995}
>>>
>>> xor_ob = observe(bernoulli_noise('c', noise_level), True)
>>> follow_prior('a', 10000, 100) # should be 0.69 True
{False: 0.3342, True: 0.6657999999999995}
```

This second inference is significantly off, and it's not the case that our observation is particularly unlikely. Here, the burn-in was not enough. Notice that in order to mix, the program must walk over states which contradict the observed values.

Suppose we are in the state $\{a:\text{True}, b:\text{True}\}$. Then, the random walk is highly discouraged from entering either of the two adjacent states $\{a:\text{True}, b:\text{False}\}$ and $\{a:\text{False}, b:\text{True}\}$, since worlds in which $a \oplus b$ is **False** are weighted against, by a factor of 100.

Let's estimate the amount of burn-in it should take for this random walk to mix. Suppose we are in a state where $a \oplus b$ is **True**. About half the time, we will attempt to enter one of the two adjacent states (the other half of the time, we will keep the same value for that flip). Entering this state will occur with probability roughly $\frac{1}{100}$. Once we are in such a state, we are given an opportunity to enter either of the two $a \oplus b$ being **True** states, so we have successfully mixed.

Thus we can roughly model this as having a $\frac{1}{200}$ chance of mixing properly, for each iteration. Thus if we have a burn-in of $200 \cdot x$, there is roughly a $1 - \frac{1}{e^x}$ chance of mixing. Here are empirical results of `burnin` against `follow_prior('a', 10000, burnin)` results:

burnin	follow_prior('a', 10000, burnin) percentage True
100	0.66579999999999995
200	0.67869999999999997
300	0.68200000000000005
400	0.68489999999999995
500	0.68759999999999999

These results are slightly better than our prediction, but confirm the overall phenomenon. Because of the noise, we can't expect it to ever converge to exactly 0.69. Instead, we may hope for it to converge to something like $0.01 \cdot 0.60 + 0.99 \cdot 0.69 = 0.6891$.

4.4 Testing a decaying atom

Suppose we are observing a radioactive atom in discrete time intervals, e.g. we look at it each second. Each time we look at the atom, there is some chance p that it decays, called its decay rate.

Suppose there is an atom whose decay rate we don't know. So our prior over decay rates is uniform.

```
>>> assume('decay', beta(1, 1))
>>> assume('geometric', function('x', ifelse(bernoulli('decay'), 'x', apply('geometric', var('x') + 1)))
>>> observe(bernoulli_noise(apply('geometric', 0) == 10, .01), True)
print get_pdf(follow_prior('decay', 100, 100), 0, 1, .1)
```

4.5 Testing mem

Here, we test that the memoization procedure works.

Let's first write the fibonacci function, in the naive way:

```
>>> fibonacci_expr = function('x', ifelse(var('x') <= 1, 1, \
...                                     apply('fibonacci', var('x')-1) + apply('fibonacci', var('x')-2)))
>>> assume('fibonacci', fibonacci_expr)
```

Of course, evaluating fibonacci in this manner is an exponential time operation:

```
>>> t = time(); sample(apply('fibonacci', 20)); time() - t
10946
1.76103687286
```

Mem is an XRP which, when applied to (possibly probabilistic) functions, returns a version of the function which is memoized. That is, function calls are remembered, so that if a function is called with the same arguments, it does not need to recompute. Here is an example of mem being used.

```
>>> assume('bad_mem_fibonacci', mem('fibonacci'))
>>>
>>> t = time(); sample(apply('bad_mem_fibonacci', 20)); time() - t
10946
1.90201187134
>>> t = time(); sample(apply('bad_mem_fibonacci', 20)); time() - t
10946
0.00019097328186
```

Notice that the second call to this mem'd fibonacci is much faster, since mem remembers the value. However, the first call is just as slow as before. Since Fibonacci is recursive, we really want to memoize all the recursive subcalls as well. The canonical introduction to dynamic programming shows how Fibonacci can be computed in linear time this way. We can write the program easily:

```
>>> mem_fibonacci_expr = function('x', ifelse(var('x')<=1, 1, \
...      apply('mem_fibonacci', var('x')-1) + apply('mem_fibonacci', var('x')-2)))
>>> assume('mem_fibonacci', mem(mem_fibonacci_expr))
>>>
>>> t = time(); sample(apply('mem_fibonacci', 20)); time() - t
10946
0.0271570682526
>>> t = time(); sample(apply('mem_fibonacci', 20)); time() - t
10946
0.000293016433716
```

4.6 Testing DPMem

Let's now implement the Dirichlet process.

```
>>> sticks_expr = mem(function('j', beta(1, 'concentration2')))
>>> atoms_expr = mem(function('j', apply('basemeasure2')))
>>> loop_body_expr = function('j', ifelse(bernoulli(apply('sticks', 'j')), apply('atoms', 'j'), \
...      apply(apply('loophelper', ['concentration2', 'basemeasure2']), var('j')+1)))
>>> loop_expr = apply(function(['sticks', 'atoms'], loop_body_expr), [sticks_expr, atoms_expr])
>>> assume('loophelper', function(['concentration2', 'basemeasure2'], loop_expr))
>>> assume('DP', function(['concentration', 'basemeasure'], \
...      apply(apply('loophelper', ['concentration', 'basemeasure']), 1)))
```

We can now use the Dirichlet process to create something we call DPmem. DPmem is a generalization of mem which sometimes returns memoized values, and sometimes resamples new values.

```

"""DEFINITION OF DPMEM"""
restaurants_expr = mem(function('args', apply('DP', ['alpha', function([], apply('proc', 'args'))]))
assume('DPMem', function(['alpha', 'proc'], function('args', apply(restaurants_expr, 'args'))))

"""TESTING DPMEM"""

concentration = 1 # when close to 0, just mem. when close to infinity, sample
assume('DPMemflip', apply('DPMem', [concentration, function(['x'], bernoulli(0.5))]))
print [sample(apply('DPMemflip', 5)) for i in xrange(10)]

print "\n TESTING GAUSSIAN MIXTURE MODEL\n"
assume('concentration', gaussian(1, 0.2)) # use vague-gamma?
assume('expected-mean', gaussian(0, 1))
assume('expected-variance', gaussian(0, 1))
assume('gen-cluster-mean', gaussian(0, 1))
assume('get-datapoint', mem( function(['id'], gaussian('gen-cluster-mean', 1.0))))
assume('outer-noise', gaussian(1, 0.2)) # use vague-gamma?

observe(gaussian(apply('get-datapoint', 0), 'outer-noise'), 1.3)
observe(gaussian(apply('get-datapoint', 0), 'outer-noise'), 1.2)

t = time()
print format(get_pdf(follow_prior('expected-mean', 100, 30), -4, 4, .5), '%0.2f')
print 'time taken', time() - t

#concentration = 1
#uniform_base_measure = uniform_no_args_XRP(2)
#print [sample(apply('DP', [concentration, uniform_base_measure])) for i in xrange(10)]
#expr = beta_bernoulli_1()

```

DPmem can be used, for example, to do mixture modeling. The idea is that we should sometimes

References