

The most powerful unfalsified model

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- Exact identification problems
- Identifiability conditions
- Algorithms
 - from data to kernel representation
 - impulse response identification
 - N4SID-type algorithms
 - MOESP-type algorithms

Exact identification problems ($w_d \mapsto \mathcal{B}$ such that $w_d \in \mathcal{B}$)

An exact identification problem

Problem P1 (Exact identification)

Given two vector time series

$$\begin{aligned} u_d &= (u_d(1), \dots, u_d(T)) \in (\mathbb{R}^m)^T && \text{"inputs"} \\ y_d &= (y_d(1), \dots, y_d(T)) \in (\mathbb{R}^p)^T && \text{"outputs"} \end{aligned}$$

find $n \in \mathbb{N}$ and LTI system \mathcal{B} of order n , with m inputs and p outputs, s.t.

$$w_d := (u_d, y_d) \in \mathcal{B},$$

i.e., w_d is a trajectory of \mathcal{B} .

How can we check that " $w_d \in \mathcal{B}$ "?

Checking that $w_d \in \mathcal{B} = \ker(R(\sigma))$

$$w_d \in \mathcal{B} \iff R(\sigma)w_d = 0$$

$$\iff R_0 w_d(t) + R_1 w_d(t+1) + \dots + R_\ell w_d(t+\ell) = 0 \\ \text{for } t = 1, \dots, T-\ell$$

$$\iff \begin{bmatrix} R_0 & R_1 & \dots & R_\ell & & \\ & R_0 & R_1 & \dots & R_\ell & \\ & & \ddots & \ddots & \ddots & \\ & & & R_0 & R_1 & \dots & R_\ell \end{bmatrix} \begin{bmatrix} w_d(1) \\ w_d(2) \\ \vdots \\ w_d(T) \end{bmatrix} = 0$$

$$\iff [R_0 \ R_1 \ \dots \ R_\ell] \begin{bmatrix} w_d(1) & w_d(2) & \dots & w_d(T-\ell) \\ w_d(2) & w_d(3) & \dots & \\ \vdots & \vdots & \ddots & \\ w_d(\ell+1) & w_d(\ell+2) & \dots & w_d(T) \end{bmatrix} = 0$$

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Checking that $w_d \in \mathcal{B} = \text{image}(M(\sigma))$

$$w_d \in \mathcal{B} \iff \text{there is } v, \text{ such that } w_d = M(\sigma)v$$

$$\iff \text{there is } v, \text{ such that for } t = 1, \dots, T \\ w_d(t) = M_0 v(t) + M_1 v(t+1) + \dots + M_\ell v(t+\ell)$$

$$\iff \text{there is solution } v \text{ of the system}$$

$$\begin{bmatrix} w_d(1) \\ w_d(2) \\ \vdots \\ w_d(T) \end{bmatrix} = \begin{bmatrix} M_0 & M_1 & \dots & M_\ell & & \\ & M_0 & M_1 & \dots & M_\ell & \\ & & \ddots & \ddots & \ddots & \\ & & & M_0 & M_1 & \dots & M_\ell \end{bmatrix} \begin{bmatrix} v(1) \\ v(2) \\ \vdots \\ v(T+\ell) \end{bmatrix}$$

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Checking that $w_d \in \mathcal{B} = \mathcal{B}_{i/s/o}(A, B, C, D)$

Let \mathcal{B} be defined by a minimal input/state/output representation

$$\mathcal{B} := \mathcal{B}_{i/s/o}(A, B, C, D) = \{ (u, y) \mid \sigma x = Ax + Bu, y = Cx + Du \}$$

$$(u_d, y_d) \in \mathcal{B}_{i/s/o}(A, B, C, D) \iff \text{there exists } x_{\text{ini}} \in \mathbb{R}^n, \text{ such that}$$

$$y_d = \underbrace{\begin{bmatrix} C \\ CA \\ CA^2 \\ \vdots \\ CA^{T-1} \end{bmatrix}}_{\theta_T(A, C)} x_{\text{ini}} + \begin{bmatrix} D & & & \\ CB & D & & \\ CAB & CB & D & \\ \vdots & \ddots & \ddots & \ddots \\ CA^{T-1}B & \dots & CAB & CB & D \end{bmatrix} u_d$$

(y_d is the response of \mathcal{B} under input u_d and initial condition x_{ini})

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Comments

- P1 is an **exact fitting problem**, a most basic SYSID problem
- easily generalizable to a **set of N time series**
 $u_{d,1}, \dots, u_{d,N} \in (\mathbb{R}^m)^T$ and $y_{d,1}, \dots, y_{d,N} \in (\mathbb{R}^p)^T$
- the **realization problem**

$$\text{impulse response} \mapsto (A, B, C, D)$$

is a special case of P1 for a set of m time series

- while m is given, **finding n is part of the problem**
any observable system of order $n \geq pT$ is a (trivial) solution
- we are actually interested is a **solution of a minimal order**

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Revised exact identification problem

Problem P1' (Exact identification)

Given two vector time series

$$\begin{aligned} u_d &= (u_d(1), \dots, u_d(T)) \in (\mathbb{R}^m)^T && \text{"inputs"} \\ y_d &= (y_d(1), \dots, y_d(T)) \in (\mathbb{R}^p)^T && \text{"outputs"} \end{aligned}$$

find the smallest $n \in \mathbb{N}$ and LTI system \mathcal{B} of order n , with m inputs and p outputs, such that

$$w_d = (u_d, y_d) \in \mathcal{B}.$$

Set of LTI systems with a bounded complexity

Notation: $\mathcal{L}_{m,\ell}^{w,n}$ is the set of all LTI systems with

- w (external) variables
- at most m inputs
- minimal state dimension at most n and
- lag (= observability index) at most ℓ

For $t \geq n$, the set $\mathcal{B}|_t$ of all t samples long traj. of \mathcal{B} has dimension

$$\dim(\mathcal{B}|_t) \leq tm + n \leq tm + p\ell$$

(where $p(\ell - 1) \leq n \leq p\ell$)

$\implies (m, n)$ and (m, ℓ) specify the complexity of the model class $\mathcal{L}_{m,\ell}^{w,n}$

Another exact identification problem

Problem P2 (Exact identification)

Given a vector time series

$$w_d = (w_d(1), \dots, w_d(T)) \in (\mathbb{R}^w)^T$$

find the smallest $m \in \mathbb{N}$ and $\ell \in \mathbb{N}$ and LTI system $\mathcal{B} \in \mathcal{L}_{m,\ell}^w$, s.t. $w_d \in \mathcal{B}$.

Comments:

- no separation between inputs and outputs
- the complexity is defined by (m, ℓ)

Most powerful unfalsified model

The most powerful unfalsified model in the model class $\mathcal{L}_{m,\ell}^w$ of a time series $w_d \in (\mathbb{R}^w)^T$ is the system $\mathcal{B}_{\text{mpum}}$ that is

1. in the model class, i.e., $\mathcal{B}_{\text{mpum}} \in \mathcal{L}_{m,\ell}^w$,
2. unfalsified, i.e., $w_d \in \mathcal{B}_{\text{mpum}}|_T$, and
3. most powerful among all LTI unfalsified systems, i.e.,

$$\mathcal{B}' \in \mathcal{L}_{m,\ell}^w \text{ and } w_d \in \mathcal{B}'|_T \implies \mathcal{B}_{\text{mpum}}|_T \subseteq \mathcal{B}'|_T.$$

MPUM may not exist, but if it does, then it is unique

Identifiability question

P2 is the problem of computing the MPUM of w_d in \mathcal{L}^w

The following related question is of interest:

Suppose that

$$w_d \in \overline{\mathcal{B}} \in \mathcal{L}^w$$

and **upper bounds** n_{\max} , ℓ_{\max} of the order n and lag ℓ of $\overline{\mathcal{B}}$ are given.

Under what conditions $\mathcal{B}_{\text{mpum}}(w_d)$ is equal to the system \mathcal{B} ?

the answer is given by the following lemma

Identifiability

Fundamental Lemma

Let $\overline{\mathcal{B}} \in \mathcal{L}_m^{w,n}$ be controllable and let $w_d := (u_d, y_d) \in \overline{\mathcal{B}}|_T$.

Then, if u_d is persistently exciting of order $L + n$,

$$\text{image} \left(\begin{bmatrix} w_d(1) & w_d(2) & w_d(3) & \cdots & w_d(T-L+1) \\ w_d(2) & w_d(3) & w_d(4) & \cdots & w_d(T-L+2) \\ w_d(3) & w_d(4) & w_d(5) & \cdots & w_d(T-L+3) \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ w_d(L) & w_d(L+1) & w_d(L+2) & \cdots & w_d(T) \end{bmatrix} \right) = \overline{\mathcal{B}}|_L$$

\Rightarrow under the conditions of the FL, any L samples long response y of \mathcal{B} can be obtained as $y = \mathcal{H}_L(y_d)g$, for certain $g \rightsquigarrow$ **algorithms**

\Rightarrow with $L = \ell_{\max} + 1$, the FL gives **conditions for identifiability**

Persistency of excitation

$u_d = (u_d(1), \dots, u_d(T))$ is **persistently exciting of order L** if

$$\mathcal{H}_L(u_d) := \begin{bmatrix} u_d(1) & u_d(2) & u_d(3) & \cdots & u_d(T-L+1) \\ u_d(2) & u_d(3) & u_d(4) & \cdots & u_d(T-L+2) \\ u_d(3) & u_d(4) & u_d(5) & \cdots & u_d(T-L+3) \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ u_d(L) & u_d(L+1) & u_d(L+2) & \cdots & u_d(T) \end{bmatrix} \quad \text{is full row rank}$$

System theoretic interpretation:

u_d is persistently exciting of order L

\iff

there is no LTI system with # of inputs $< m$ and lag $< L$ for which u_d is a trajectory

Algorithms for exact identification ($w_d \mapsto$ representation of the MPUM)

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$$w_d \mapsto R(\xi)$$

under the assumptions of the FL, $\text{image}(\mathcal{H}_{\ell_{\max}+1}(w_d)) = \mathcal{B}|_{\ell_{\max}+1}$

\Rightarrow a basis for $\text{left ker}(\mathcal{H}_{\ell_{\max}+1}(w_d))$ defines a kernel repr. of \mathcal{B}

let

$$\begin{bmatrix} \tilde{R}_0 & \tilde{R}_1 & \cdots & \tilde{R}_{\ell_{\max}} \end{bmatrix} \mathcal{H}_{\ell_{\max}+1}(w_d) = 0, \quad \text{where } \tilde{R}_i \in \mathbb{R}^{g \times w}$$

and define $\tilde{R}(\xi) = \sum_{i=0}^{\ell_{\max}} \xi^i \tilde{R}_i$

then $\mathcal{B} = \ker(\tilde{R}(\sigma))$ is, in general, a **nonminimal kernel representation**

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1. $w_d \mapsto R(\xi)$
2. $w_d \mapsto$ impulse response H
3. $w_d \mapsto (A, B, C, D)$ (possibly **balanced**)
 - 3.1 $w_d \mapsto R(\xi) \mapsto (A, B, C, D)$ or $w_d \mapsto H \mapsto (A, B, C, D)$
 - 3.2 $w_d \mapsto \mathcal{O}_{\ell_{\max}+1}(A, C) \mapsto (A, B, C, D)$
 - 3.3 $w_d \mapsto (x_d(1), \dots, x_d(\ell_{\max} + m + 1)) \mapsto (A, B, C, D)$

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$$w_d \mapsto R(\xi)$$

\tilde{R} can be made minimal by standard polynomial linear algebra alg.
find a unimodular matrix U , such that

$$U\tilde{R} = \begin{bmatrix} R \\ 0 \end{bmatrix} \quad \text{and } R \text{ is full row rank}$$

then $\ker(R(\sigma)) = 0$ is minimal

Refinements:

- efficient **recursive computation** (exploiting the Hankel structure)
- as a byproduct **find an input/output partition** of the variables
- find a **shortest lag** kernel representation (i.e., R row proper)

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$$w_d \mapsto H$$

Under the conditions of FL, there is G , such that $H = \mathcal{H}_t(y_d)G$
the problem reduces to the one of finding a particular G . Define

$$\begin{bmatrix} \mathcal{H}_{\ell_{\max}+t}(u_d) \\ \mathcal{H}_{\ell_{\max}+t}(y_d) \end{bmatrix} =: \begin{bmatrix} U_p \\ U_f \end{bmatrix} \quad \begin{array}{l} \text{row dim}(U_p) = \text{row dim}(Y_p) = \ell_{\max} \\ \text{row dim}(U_f) = \text{row dim}(Y_f) = t \end{array}$$

Let u_d be p.e. of order $t + \ell_{\max} + n_{\max}$. Then there is G , such that

$$\left. \begin{array}{l} \begin{bmatrix} U_p \\ Y_p \\ U_f \end{bmatrix} G = \begin{bmatrix} 0 \\ 0 \\ \begin{bmatrix} I_m \\ 0 \end{bmatrix} \end{bmatrix} \right\} \begin{array}{l} \text{zero ini. conditions} \\ \leftarrow \text{impulse input} \end{array} \quad (1)$$

$$Y_f G = H$$

$$w_d \mapsto H$$

Block algorithm for computation of $(H(0), \dots, H(t-1))$:

1. Input: u_d, y_d, ℓ_{\max} , and t .
2. Solve the system of eqs (1). Let \bar{G} be the computed solution.
3. Compute $H = Y_f \bar{G}$.
4. Output: the first t samples of the impulse response H .

Refinements:

- solve (1) efficiently by **exploiting the Hankel structure**
- do the computations iteratively for pieces of $H \rightsquigarrow$ **iterative alg.**
- **automatically choose t** , for a sufficient decay of H

$$w_d \mapsto (A, B, C, D)$$

- $w_d \mapsto H(0 : 2\ell_{\max})$ or $R(\xi) \xrightarrow{\text{realization}} (A, B, C, D)$
- $w_d \mapsto \mathcal{O}_{\ell_{\max}+1}(A, C) \xrightarrow{(2)} (A, B, C, D)$
- $w_d \mapsto (x_d(1), \dots, x_d(n_{\max} + m + 1)) \xrightarrow{(3)} (A, B, C, D)$

(2) and (3) are easy:

$$\mathcal{O}_{\ell_{\max}+1}(A, C) \mapsto (A, C) \quad \text{and} \quad (u_d, y_d, A, C) \mapsto (B, C, x_{\text{ini}}) \quad (2)$$

$$\begin{bmatrix} x_d(2) & \cdots & x_d(n_{\max} + m + 1) \\ y_d(1) & \cdots & y_d(n_{\max} + m) \end{bmatrix} = \begin{bmatrix} A & B \\ C & D \end{bmatrix} \begin{bmatrix} x_d(1) & \cdots & x_d(n_{\max} + m) \\ u_d(1) & \cdots & u_d(n_{\max} + m) \end{bmatrix} \quad (3)$$

$$\mathcal{O}_{\ell_{\max}+1}(A, C) \mapsto (A, B, C, D)$$

C is the **first block entry** of $\mathcal{O}_{\ell_{\max}+1}(A, C)$ and A is given by

$$(\sigma^* \mathcal{O}_{\ell_{\max}+1}(A, C))A = (\sigma \mathcal{O}_{\ell_{\max}+1}(A, C)) \quad \text{shift equation}$$

(σ^* removes the last block entry and σ removes the first block entry)

Once C and A are known, the system of equations

$$y_d(t) = CA^t x_d(1) + \sum_{\tau=1}^{t-1} CA^{t-1-\tau} B u_d(\tau) + D \delta(t+1), \quad \text{for } t = 1, \dots, \ell_{\max} + 1$$

is **linear in $D, B, x_d(1)$** (can be solved using Kronecker products)

$$w_d \mapsto \mathcal{O}_{\ell_{\max}+1}(A, C)$$

The columns of $\mathcal{O}_{\ell_{\max}+1}(A, C)$ are n linearly indep. free responses of \mathcal{B}

Under the conditions of FL, **such resp. can be computed from data**

$$\begin{bmatrix} \mathcal{H}_t(u_d) \\ \mathcal{H}_t(y_d) \end{bmatrix} G = \begin{bmatrix} 0 \\ Y_0 \end{bmatrix} \quad \begin{array}{l} \leftarrow \text{zero inputs} \\ \leftarrow \text{free responses} \end{array}$$

in order to obtain lin. indep. free responses, G should be maximal rank

Once we have a maximal rank matrix of free responses Y_0

$$Y_0 = \mathcal{O}_{\ell_{\max}+1}(A, C) \underbrace{\begin{bmatrix} x_{\text{ini},1} & \cdots & x_{\text{ini},j} \end{bmatrix}}_{X_{\text{ini}}} \quad \text{rank revealing factorization}$$

$\rightsquigarrow \mathcal{O}_{\ell_{\max}+1}(A, C)$ and X_{ini} , the factorization fixes the state space basis

Refinements

- Solve (4) efficiently **exploiting the Hankel structure**

- Iteratively compute pieces of Y_0

\rightsquigarrow **iterative algorithm**

- requires smaller persistency of excitation of u_d
- could be more efficient

(Solve a few smaller systems of eqns instead of a single bigger one)

$$w_d \mapsto (x_d(1), \dots, x_d(n_{\max} + m + 1))$$

If the free responses are **sequential**, i.e., if Y_0 is block-Hankel, then X_{ini} is a state sequence of \mathcal{B}

Computation of **sequential free responses** is achieved as follows

$$\begin{array}{l} \left. \begin{array}{l} \begin{bmatrix} U_p \\ Y_p \\ U_f \end{bmatrix} G = \begin{bmatrix} U_p \\ Y_p \\ 0 \end{bmatrix} \\ Y_f G = Y_0 \end{array} \right\} \begin{array}{l} \text{sequential ini. conditions} \\ \leftarrow \text{zero inputs} \end{array} \end{array} \quad (4)$$

Note: now we use the splitting of the data into “past” and “future”

$$Y_0 = \mathcal{O}_{\ell_{\max}+1}(A, C) \begin{bmatrix} x_d(1) & \cdots & x_d(n_{\max} + m + 1) \end{bmatrix} \quad \text{rank revealing factorization}$$

MOESP type algorithms

Orth. projection of the rows of $\mathcal{H}_{n_{\max}}(y_d)$ on $\left(\text{rowspan}(\mathcal{H}_{n_{\max}}(u_d)) \right)^\perp$

$$Y_0 := \mathcal{H}_{n_{\max}}(y_d) \Pi_{u_d}^\perp$$

where

$$\Pi_{u_d}^\perp := \left(I - \mathcal{H}_{n_{\max}}^\top(u_d) (\mathcal{H}_{n_{\max}}(u_d) \mathcal{H}_{n_{\max}}^\top(u_d))^{-1} \mathcal{H}_{n_{\max}}(u_d) \right)$$

Observe that $\Pi_{u_d}^\perp$ is maximal rank and

$$\begin{bmatrix} \mathcal{H}_{n_{\max}}(u_d) \\ \mathcal{H}_{n_{\max}}(y_d) \end{bmatrix} \Pi_{u_d}^\perp = \begin{bmatrix} 0 \\ Y_0 \end{bmatrix}$$

\Rightarrow **the orthogonal projection computes free responses**

Comments

- $T - n_{\max} + 1$ free responses are computed via the orth. proj. while n_{\max} such responses suffice for the purpose of exact identification
- The orth. proj. is a **geometric operation**, whose system theoretic meaning is not revealed
- The **condition for $\text{rank}(Y_0) = n$** , given in the MOESP literature,

$$\text{rank} \left(\begin{bmatrix} X_{\text{ini}} \\ \mathcal{H}_{n_{\max}}(u_d) \end{bmatrix} \right) = n + n_{\max}m$$

is **not verifiable from the data (u_d, y_d)** \implies can not be checked whether the computation gives $\mathcal{O}(A, C)$, cf., p.e. condition of FL

N4SID-type algorithms

Consider the splitting of the data into “past” and “future”

$$\mathcal{H}_{2n_{\max}}(u_d) =: \begin{bmatrix} U_p \\ U_f \end{bmatrix}, \quad \mathcal{H}_{2n_{\max}}(y_d) =: \begin{bmatrix} Y_p \\ Y_f \end{bmatrix}$$

with $\text{rowdim}(U_p) = \text{rowdim}(U_f) = \text{rowdim}(Y_p) = \text{rowdim}(Y_f) = n_{\max}$ and let

$$W_p := \begin{bmatrix} U_p \\ Y_p \end{bmatrix}$$

The key step of the N4SID algorithms is the **oblique projection** of the rows of Y_f along $\text{rowspan}(U_f)$ onto $\text{rowspan}(W_p)$

$$Y_0 := Y_f / U_f W_p := Y_f \underbrace{\begin{bmatrix} W_p^\top & U_f^\top \end{bmatrix} \begin{bmatrix} W_p W_p^\top & W_p U_f^\top \\ U_f W_p^\top & U_f U_f^\top \end{bmatrix}^+ \begin{bmatrix} W_p \\ 0 \end{bmatrix}}_{\Pi_{\text{obl}}}$$

N4SID-type algorithms

Observe that

$$\begin{bmatrix} W_p \\ U_f \\ Y_f \end{bmatrix} \Pi_{\text{obl}} = \begin{bmatrix} W_p \\ 0 \\ Y_0 \end{bmatrix}$$

(in fact Π_{obl} is the least-norm, least-squares solution)

\implies **the oblique projection computes sequential free responses**

Comments

- $T - 2n_{\max} + 1$ sequential free responses are computed via the oblique projection while $n_{\max} + m + 1$ such responses suffice for exact ident.
- The oblique proj. is a **geometric operation**, whose system theoretic meaning is not revealed
- The **conditions for $\text{rank}(Y_0) = n$** , given in the N4SID literature,
 1. u_d persistently exciting of order $2n_{\max}$ and
 2. $\text{rowspan}(X_{\text{ini}}) \cap \text{rowspan}(U_f) = \{0\}$
 are **not verifiable from the data (u_d, y_d)**

References

- **System theoretic interpretation** of the orth. and oblique proj.
- MOESP and N4SID alg. are computationally inefficient; more than what is necessary for exact ident. is computed \rightsquigarrow **cheaper algs**
- The FL gives **conditions for identifiability**, verifiable from the data
- We clarified the **role of the splitting**: the “past” assigns the initial conditions and in the “future” a desired response is computed
 \implies “past” should be chosen at least ℓ samples long; the length of “future” is free as long as the p.e. condition is satisfied

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Software

A MATLAB toolbox for exact SYSID is available from:

`ftp.esat.kuleuven.be/pub/SISTA/markovsky/
abstracts/05-122.html`

In exercise 2 you will use the algorithms

- $w_d \mapsto R(\xi)$ (w2r) and
- $w_d \mapsto (x_d(1), \dots, x_d(n_{\max} + m + 1)) \mapsto (A, B, C, D)$ (uy2x2ss)

in order to find the MPUM for given trajectory of an LTI system.