

Recursive computation of the most powerful unfalsified model

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Recursive LTI modeling

Our aim:

exact recursive modeling of time series by LTI systems

The algorithms can work recursively in:

1. the number of time series,
2. the length of the time series (real-time),
3. the length of the system laws,

or a combination of these.

recursive \approx computationally efficient algorithms

Outline

Introduction

Lag-recursive computation of the MPUM

Simulation example

System identification: $w_d \mapsto \hat{\mathcal{B}} \in \mathcal{M}$

Notation:

- $w_d = (u_d, y_d)$ — given data, in this talk a vector time series
- $\hat{\mathcal{B}}$ — to be found model for w_d , in this talk an LTI system
- \mathcal{M} — model class, in this talk the set of LTI systems \mathcal{L}

System identification

- defines a mapping $w_d \mapsto \mathcal{B}$
- derives effective algorithms that realize the mapping, and
- develops efficient software that implements the algorithms

Exact identification: two points of view

Find the true data generating system

- assume that $w_d \in \tilde{\mathcal{B}} \in \mathcal{L}$
- find back $\tilde{\mathcal{B}}$ from w_d (and an upper bound of the order)
- this is possible provided $\tilde{\mathcal{B}}$ is controllable and an input component of w_d is persistently exciting

Find the least complex LTI system that fits w_d

- no assumption about w_d
- find $\hat{\mathcal{B}} \in \mathcal{L}$ with minimal # of inputs and order, s.t. $w_d \in \hat{\mathcal{B}}$
- $\hat{\mathcal{B}}$ —**most powerful unfalsified model (MPUM)** for w_d in \mathcal{L}

The most powerful unfalsified model (MPUM)

\mathcal{L}^w — LTI model class with w external variables

$\mathcal{B} \in \mathcal{L}^w$ is a linear, shift-invariant, closed subspace of $(\mathbb{R}^w)^{\mathbb{N}}$

\mathcal{B} is the MPUM of w_d in \mathcal{L}^w if

1. \mathcal{B} is in the model class \mathcal{L}^w , i.e., $\mathcal{B} \in \mathcal{L}^w$
2. \mathcal{B} is unfalsified by w_d , i.e., $w_d \subseteq \mathcal{B}$, and
3. any other unfalsified system in \mathcal{L}^w is less powerful, i.e., $\mathcal{B}' \in \mathcal{L}^w$ and $w_d \subseteq \mathcal{B}'$ imply $\mathcal{B} \subseteq \mathcal{B}'$.

LTI model representations

- **Kernel representation** (parameter $R(z) := \sum_{i=0}^1 R_i z^i$)

$$R_0 w(t) + R_1 w(t+1) + \cdots + R_1 w(t+1) = 0$$

- **Impulse response represent** (parameter $h: \mathbb{Z} \rightarrow \mathbb{R}^{p \times m}$)

$$w = \text{col}(u, y), \quad y(t) = \sum_{\tau=-\infty}^t h(\tau) u(t-\tau)$$

- **Input/state/output representation** (parameter (A, B, C, D))

$$w = \text{col}(u, y), \quad \begin{aligned} x(t+1) &= Ax(t) + Bu(t) \\ y(t) &= Cx(t) + Du(t) \end{aligned}$$

$p := \dim(y) = \text{row dim}(R)$ is the # of outputs

$m := \dim(u)$ is the # of inputs, $1 := \text{degree}(R)$ is the lag

Laws of the system

Consider a system $\mathcal{B} \in \mathcal{L}^w$ defined by

$$\mathcal{B} := \{ w \in (\mathbb{R}^w)^{\mathbb{Z}} \mid R_0 w(t) + R_1 w(t+1) + \cdots + R_1 w(t+1) = 0 \}$$

and let $R(z) := \sum_{i=0}^1 R_i z^i \in \mathbb{R}^{p \times w}[z]$.

The rows $r_1(z), \dots, r_p(z)$ of $R(z)$ are **annihilators or laws** of \mathcal{B} .

Define the shift operator σ : $\sigma w(t) := w(t+1)$

$$r_i(\sigma)\mathcal{B} = 0, \quad \text{for } i = 1, \dots, p$$

Hankel matrix of the data

$$\mathcal{H}_{l+1}(w_d) := \begin{bmatrix} w_d(1) & w_d(2) & \cdots & w_d(T-l) \\ w_d(2) & w_d(3) & \cdots & w_d(T-l+1) \\ w_d(3) & w_d(4) & \cdots & w_d(T-l+2) \\ \vdots & \vdots & \cdots & \vdots \\ w_d(l+1) & w_d(l+2) & \cdots & w_d(T) \end{bmatrix}$$

The left kernel of $\mathcal{H}_{l+1}(w_d)$

$$\mathcal{N}_{l+1}(w_d) := \text{left ker}(\mathcal{H}_{l+1}(w_d))$$

contains laws of $\mathcal{B}_{\text{mpum}}(w_d)$.

Moreover, if $\mathcal{B}_{\text{mpum}}(w_d)$ is controllable, all laws of $\mathcal{B}_{\text{mpum}}(w_d)$ are contained in a basis of $\mathcal{N}_{l+1}(w_d)$ for $l \geq 1$ (the lag of $\mathcal{B}_{\text{mpum}}$).

Inefficiency of the algorithm

Suppose that $\mathcal{H}_{l+1}(w_d)$ is rank deficient and let

$$\underbrace{[r_0 \ r_1 \ \cdots \ r_l]}_r \mathcal{H}_{l+1}(w_d) = 0.$$

Then due to the Hankel structure for $L > l$

$$\begin{bmatrix} r_0 & r_1 & \cdots & r_l & & 0 \\ & r_0 & r_1 & \cdots & r_l & \\ & & \cdots & \cdots & & \cdots \\ 0 & & & r_0 & r_1 & \cdots & r_l \end{bmatrix} \mathcal{H}_{L+1}(w_d) = 0.$$

Moreover, extra work has to be done in order to “extract” from $\text{left ker}(\mathcal{H}_{L+1}(w_d))$ a short law r .

In polynomial language, $R(z)$ is not row reduced and an extra work is needed for its row reduction.

Identification algorithm: $w_d \rightarrow R(z)$

Let the rows of R form a basis for $\mathcal{N}_{l+1}(w_d)$, i.e.,

$$R \mathcal{H}_{l+1}(w_d) = 0, \quad \text{rank}(R) = \dim(\mathcal{N}_{l+1}(w_d)).$$

Define a polynomial matrix $R(z)$ constructed from R ,

$$R(z) = \sum_{i=0}^l R_i z^i \in \mathbb{R}^{g \times w}[z],$$

where $R =: [R_0 \ R_1 \ \cdots \ R_l]$, $R_i \in \mathbb{R}^{g \times w}$. Then

$$\ker(R(\sigma)) = \mathcal{B}_{\text{mpum}}(w_d).$$

Lag-recursive algorithm

Assume that the k shortest laws $r^{(1)}, \dots, r^{(k)}$ have been computed

$$r^{(i)}(\sigma) w_d = 0, \quad \text{for } i = 1, \dots, k.$$

Question: How to use this information in order to compute the remaining laws $r^{(k+1)}, \dots, r^{(p)}$ of $\mathcal{B}_{\text{mpum}}(w_d)$ efficiently?

Assuming that the system induced by the matrix

$$R' := \text{col}(r^{(1)}, \dots, r^{(k)})$$

is controllable, this can be done as follows:

1. Compute D , such that $\text{col}(R', D)$ is unimodular.
2. Define $e := D(\sigma) w_d$.
3. Compute E , such that $\mathcal{B}_{\text{mpum}}(e) = \ker(E(\sigma))$.
4. Define $R := (R', ED)$.

Lag-recursive algorithm (cont.)

Lemma

- Assume that $\mathcal{B} := \ker(R(\sigma))$ is controllable.
- Let $R'(\sigma)\mathcal{B} = 0$ and $\mathcal{B}' := \ker(R'(\sigma))$ is controllable.
- Let D be such that $\text{col}(R', D)$ is unimodular.

Then there exist U and F , such that $UR = \text{col}(R', FD)$.

Proposition

- Assume that $\mathcal{B}_{\text{mpum}}(w_d)$ is controllable.
- Let $R'(\sigma)w_d = 0$ and $\ker(R'(\sigma))$ is controllable.
- Let D be such that $\text{col}(R', D)$ is unimodular.
- Define $e := D(\sigma)w_d$ and let $\ker(E(\sigma)) = \mathcal{B}_{\text{mpum}}(e)$.

Then $\ker(\text{col}(R', ED)(\sigma)) = \mathcal{B}_{\text{mpum}}(w_d)$.

Lag-recursive algorithm (cont.)

Computing a D , such that $\text{col}(r, D)$ is unimodular is a **Bézout-type problem**.

Assuming that r_1 and r_2 of r are coprime, we solve the Bézout equation

$$r_1 b - r_2 a = 1,$$

and define

$$D = \begin{bmatrix} a & b & 0 \\ 0 & 0 & I_{w-2} \end{bmatrix},$$

where I_{w-2} is the $(w-2) \times (w-2)$ identity matrix.

Applying the procedure recursively leads to a lag-recursive algorithm for computing a kernel representation of $\mathcal{B}_{\text{mpum}}(w_d)$.

$R = \text{lag_recursive_mpum}(w_d)$

- 1: **if** $w_d = 0$ **then**
- 2: Let $R := []$.
- 3: **else**
- 4: Find the smallest $l \in \{0, 1, 2, \dots\}$, for which $\mathcal{H}_{l+1}(w_d)$ is rank deficient, or define $l := -1$ if there is no such l .
- 5: **if** $l = -1$ **then**
- 6: Let $R := []$.
- 7: **else**
- 8: Compute $r =: [r_0 \ r_1 \ \dots \ r_l]$, $r_i \in \mathbb{R}^{1 \times w}$ in $\text{leftker}(\mathcal{H}_{l+1}(w_d))$ and let $r(z) := \sum_{i=0}^l r_i z^i$.
- 9: Find D , such that $\text{col}(r, D)$ is unimodular.
- 10: Compute $e := D(\sigma)w_d$.
- 11: $r' := \text{lag_recursive_mpum}(e)$.
- 12: **if** $r' = []$, **then** $R := r$ **else** $R := \text{col}(r, r'D)$ **end if**
- 13: **end if**
- 14: **end if**

Simulation example

Consider the 2 outputs 1 input system \mathcal{B} induced by

$$R(z) = \begin{bmatrix} r_1(z) \\ r_2(z) \end{bmatrix} = \begin{bmatrix} -1.73 & 1.37 & -2.03 \\ 0.35 & -0.50 & -0.24 \end{bmatrix} z^0 + \begin{bmatrix} 5.97 & -8.36 & 12.25 \\ 0.16 & 0.65 & 0.19 \end{bmatrix} z^1 + \begin{bmatrix} 2.00 & 7.82 & -11.33 \\ 0.00 & 0.00 & 0.00 \end{bmatrix} z^2.$$

\mathcal{B} has two laws: $r_1(z)$ is of degree 2 and $r_2(z)$ is of degree 1

Simulation example (cont.)

$$\dim(\underbrace{\text{leftker}(\mathcal{H}_2(w_d))}_{\mathcal{N}_2(w_d)}) = 1 \text{ and } \mathcal{N}_2(w_d) = \text{span}(r_2(z))$$

$$\mathcal{N}_2(w_d) = \text{span}\left(\begin{bmatrix} 0.35 & -0.50 & -0.24 & 0.16 & 0.65 & 0.19 \end{bmatrix}^\top\right).$$

$$\dim(\mathcal{N}_3(w_d)) = 3, \text{ and } \mathcal{N}_3(w_d) = \text{span}(r_2(z), zr_2(z), r_1(z))$$

$$\text{span}\left(\begin{bmatrix} 0.35 & -0.5 & -0.24 & 0.16 & 0.65 & 0.19 & 0.00 & 0.00 & 0.00 \\ 0.00 & 0.00 & 0.00 & 0.35 & -0.50 & -0.24 & 0.16 & 0.65 & 0.19 \\ -1.7 & 1.37 & -2.03 & 5.97 & -8.36 & 12.25 & 2.00 & 7.82 & -11.3 \end{bmatrix}^\top\right).$$

$$\text{In general, } \dim(\mathcal{N}_{l+1}(w_d)) = 2l - 1, \text{ for } l \geq 1.$$

Conclusions

- Exact recursive in the lag of the laws modeling.
- Avoid the repeated calculation of one and the same law.
- Main idea: project the time series on the orthogonal complement of the currently found laws.
- This is achieved by solving a Bézout equation.

Simulation example (cont.)

Let w_d be a random trajectory of \mathcal{B} .

Applied on w_d , the proposed algorithm returns the matrix

$$\hat{R}(z) = \begin{bmatrix} 0.36 & -0.52 & -0.24 \\ 0.00 & 0.03 & 0.13 \end{bmatrix} z^0 + \begin{bmatrix} 0.17 & 0.67 & 0.20 \\ -0.03 & -0.12 & -0.74 \end{bmatrix} z^1 + \begin{bmatrix} 0.00 & 0.00 & 0.00 \\ 0.03 & 0.13 & 0.64 \end{bmatrix} z^2$$

which is up to pre-multiplication by unimodular matrix equal to R

$$\mathcal{B} = \ker(\hat{R}(\sigma)).$$

Thank you