

State representations from finite time series

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Exact identification

Let $w_d = (u_d, y_d) \in (\mathbb{R}^w)^T$ be generated by an LTI system \mathcal{B} .

Problem: find \mathcal{B} back from the data $w_d = (w_d(1), \dots, w_d(T))$.

Alternatively, given an arbitrary w_d , find the most powerful unfalsified model $\mathcal{B}_{\text{mpum}}$ for w_d in the LTI model class.

We consider an input/state/output representation of \mathcal{B} :

$$\sigma x = Ax + Bu, \quad y = Cx + Du$$

where σ is the forward shift operator $(\sigma x)(t) = x(t+1)$

m — number of inputs, $p = w - m$ — number of outputs of \mathcal{B}



Outline

Introduction

Past-future intersection

Shift-and-cut map

Simulation example



Approach

we consider N4SID-type algorithms: $w_d \mapsto x_d \mapsto (A, B, C, D)$

x_d — minimal state sequence of \mathcal{B} corresponding to w_d

$(w_d, x_d) \mapsto (A, B, C, D)$ is solving a linear system of equations:

$$\begin{bmatrix} \sigma x_d \\ y_d \end{bmatrix} = \begin{bmatrix} A & B \\ C & D \end{bmatrix} \begin{bmatrix} x_d \\ u_d \end{bmatrix}$$

$w_d \mapsto x_d$ is the heart of the identification problem



Result for infinite time series

Define

$$\begin{bmatrix} \mathcal{H}_p \\ \mathcal{H}_f \end{bmatrix} := \begin{bmatrix} \vdots & \vdots & \vdots \\ \cdots & w_d(-2) & w_d(-1) & w_d(0) & \cdots \\ \cdots & w_d(-1) & w_d(0) & w_d(1) & \cdots \\ \cdots & w_d(0) & w_d(1) & w_d(2) & \cdots \\ \cdots & w_d(1) & w_d(2) & w_d(3) & \cdots \\ \vdots & \vdots & \vdots \end{bmatrix}$$

The minimal state dimension of $\mathcal{B}_{\text{mpum}}(w_d)$ is equal to

$$n = \text{rank}(\mathcal{H}_p) + \text{rank}(\mathcal{H}_f) - \text{rank}(\mathcal{H}).$$

Moreover, a basis for $\text{rowspan}(\mathcal{H}_p) \cap \text{rowspan}(\mathcal{H}_f)$ is a minimal state sequence x_d of $\mathcal{B}_{\text{mpum}}(w_d)$, corresponding to w_d .

Notation

The Hankel matrix with Δ block-rows composed of w_d is

$$\mathcal{H}_\Delta(w_d) := \begin{bmatrix} w_d(1) & w_d(2) & \cdots & w_d(T-\Delta+1) \\ w_d(2) & w_d(3) & \cdots & w_d(T-\Delta+2) \\ w_d(3) & w_d(4) & \cdots & w_d(T-\Delta+3) \\ \vdots & \vdots & \cdots & \vdots \\ w_d(\Delta) & w_d(\Delta+1) & \cdots & w_d(T) \end{bmatrix}$$

Past-future intersection for finite time series

Define the (finite) input and output, “past” and “future” matrices

$$\begin{bmatrix} U_p \\ U_f \end{bmatrix}_{m\Delta} := \mathcal{H}_{2\Delta}(u_d), \quad \begin{bmatrix} Y_p \\ Y_f \end{bmatrix}_{p\Delta} := \mathcal{H}_{2\Delta}(y_d),$$

and the matrix of the “past” and “future” state sequences

$$X_p := [x_d(1) \cdots x_d(T-2\Delta+1)], \quad X_f := [x_d(\Delta+1) \cdots x_d(T-\Delta+1)]$$

In an early subspace ID paper it is shown that

$$\text{rowspan}(X_f) = \text{rowspan}(W_p) \cap \text{rowspan}(W_f)$$

holds under the rank condition $\text{rank}\left(\begin{bmatrix} X_p \\ U_p \end{bmatrix}\right) = n + \Delta m$.

This gives two following **two-stage SVD procedure**.

Input: $w_d = (u_d, y_d)$, and $\Delta \in \mathbb{Z}$, $\Delta \geq 1$ (the lag of \mathcal{B}).

1: Compute the SVD

$$U\Sigma V^\top = \mathcal{H}_{2\Delta}(w_d) =: \begin{bmatrix} W_p \\ W_f \end{bmatrix}_{w\Delta}$$

and define the partitionings

$$\begin{bmatrix} U_{11} & U_{12} \\ U_{21} & U_{22} \end{bmatrix}_{w\Delta}^{w\Delta} := \begin{bmatrix} U_1 & U_2 \end{bmatrix}_{w\Delta}^{w\Delta} := U.$$

2: Compute the matrix $\tilde{X} := U_{12}^\top W_p$.

3: Compute the SVD $\tilde{U}\tilde{\Sigma}\tilde{V}^\top = \tilde{X}$, and define $n := \text{rank}(U_{12}^\top W_p)$

$$\begin{bmatrix} \tilde{U}_1 \\ \tilde{U}_2 \end{bmatrix}_{w\Delta-n}^n := \tilde{U}.$$

Output: \tilde{U}_1 — a minimal state sequence of \mathcal{B} .

Explanation of the algorithm

The first SVD is used to compute a basis for $\text{left ker}(\mathcal{H}_{2\Delta}(w_d))$,

$$\text{colspan}(U_2) = \text{left ker}(\mathcal{H}_{2\Delta}(w_d))$$

From

$$\tilde{X} := U_{12}^\top W_p = -U_{22}^\top W_f,$$

it follows that

$$\text{row span}(\tilde{X}) = \text{row span}(W_p) \cap \text{row span}(W_f),$$

so that the columns of \tilde{X} form a state sequence.

The second SVD is used to compute a minimal state sequence.

There is a nonsingular $S \in \mathbb{R}^{n \times n}$, such that

$$\bar{U}_1 = S X_f, \quad \det(S) \neq 0.$$

Kernel structure of a Hankel matrix

Lemma

$$\begin{aligned} \text{colspan}(N_1) &= \text{left ker}(\mathcal{H}_{\Delta_1}(w_d)) \\ \implies \text{colspan}(N_2) &= \text{left ker}(\mathcal{H}_{\Delta_2}(w_d)) \end{aligned}$$

holds under the conditions that

1. $\Delta_1 > L$,
2. $\mathcal{B}_{\text{mpum}}$ is controllable, and
3. $\mathcal{H}_{L+1+n}(u_d)$ is full rank (persistency of excitation).

Therefore the knowledge of $\text{left ker}(\mathcal{H}_{\Delta}(w_d))$, $\Delta > L$, suffices to construct $\text{left ker}(\mathcal{H}_{2\Delta}(w_d))$.

This gives the following algorithm.

Kernel structure of a Hankel matrix

Let the columns of N_1 be in the left kernel of $\mathcal{H}_{\Delta_1}(w_d)$, i.e.,

$$N_1^\top \mathcal{H}_{\Delta_1}(w_d) = 0$$

and $N_{1,i} \in \mathbb{R}^{w \times \text{col dim}(N_1)}$ be the i th block element of N_1 . Define

$$N_2 := \begin{bmatrix} N_{1,0} & 0 & \cdots & 0 \\ \vdots & N_{1,0} & \ddots & \vdots \\ N_{1,\Delta_1-1} & \vdots & \ddots & 0 \\ 0 & N_{1,\Delta_1-1} & & N_{1,0} \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & N_{1,\Delta_1-1} \end{bmatrix}$$

Then, for $\Delta_2 > \Delta_1$,

$$N_2^\top \mathcal{H}_{\Delta_2}(w_d) = 0.$$

Input: $w_d = (u_d, y_d)$, and $\Delta \in \mathbb{Z}$, $\Delta \geq L$ (the lag of \mathcal{B}).

- 1: Compute a basis N for $\text{left ker}(\mathcal{H}_{\Delta}(w_d))$.
- 2: Compute the matrix

$$\tilde{X} := \begin{bmatrix} N_0^\top & N_1^\top & \cdots & N_{\Delta-2}^\top \\ 0 & N_0^\top & \cdots & N_{\Delta-3}^\top \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & N_0^\top \end{bmatrix} \mathcal{H}_{\Delta-1}(\sigma w_d)$$

- 3: Compute the SVD $\bar{U} \bar{\Sigma} \bar{V}^\top = \tilde{X}$ and define $n := \text{rank}(U_{12}^\top W_p)$

$$\begin{bmatrix} \bar{U}_1 \\ \bar{U}_2 \end{bmatrix}_{\substack{n \\ w\Delta - n}} := \bar{U}.$$

Output: \bar{U}_1 — a minimal state sequence of \mathcal{B} .

Explanation of the algorithm

Step 2 of the new algorithm is the **shift-and-cut map**.

If the kernel computation on step 1 is carried out by the SVD

$$U\Sigma V^T = \mathcal{H}_\Delta(w_d)$$

the new algorithm matches exactly the structure of the past-future intersection algorithm.

Now however we use only matrices with Δ block rows, while before the matrices were with 2Δ block rows.

This leads to improved numerical efficiency.

Simulation example

\mathcal{B} is a random system, $n = 6$, $m = 3$, and $p = 2$

w_d is a random trajectory of \mathcal{B} , $T = 100$

$\hat{\mathcal{B}}_1$ — the system computed by the classical algorithm

$\hat{\mathcal{B}}_2$ — the system computed by the shift-and-cut algorithm

$$\|\mathcal{B} - \hat{\mathcal{B}}_1\|_\infty = 2.4 \times 10^{-15} \quad \text{and} \quad \|\mathcal{B} - \hat{\mathcal{B}}_2\|_\infty = 4.3 \times 10^{-15}$$

$\Rightarrow \mathcal{B}$ is recovered exactly from w_d (up to the numerical errors)

computational requirements of the algorithms:

$$f_1 = 1.9 \times 10^6 \quad \text{and} \quad f_2 = 1.3 \times 10^6$$

Thank you