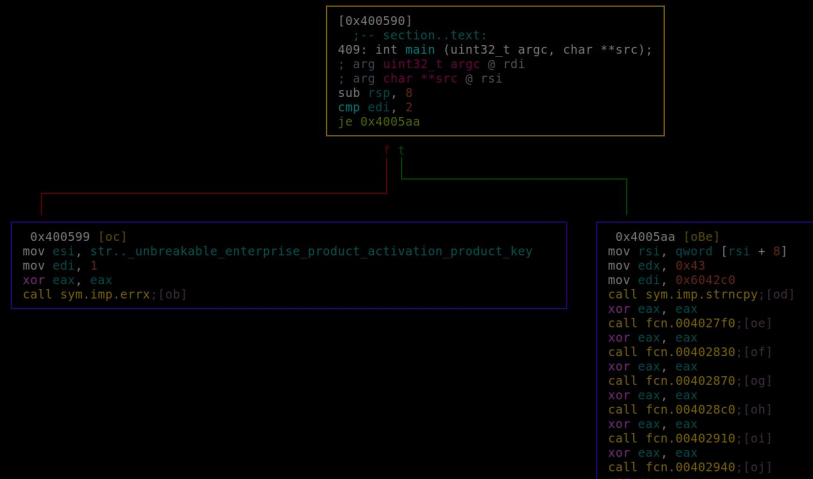




ESILSolve: A Symbolic Execution Engine using ESIL

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\$ whoami

- Mobile Security Researcher at Nowsecure where I work on automating security tests for mobile applications (with Pancake, Edu, Grant, others!)
- Studied Math and Physics at University but sold out and did infosec instead
- Background of doing CTFs and game hacking for fun, as well as some iOS tweak dev
- Have worked as both a Reverse Engineer and boring Software Developer in previous roles
- Twitter: @alkalinesec

Introduction to ESILSolve

```
[0x400590]
;-- section:.text:
400: int main (uint32_t argc, char **src);
; arg uint32_t argc @ rdi
; arg char **src @ rsi
0x00400590 4883ec08      sub rsp, 8
0x00400594 83ff02      cmp edi, 2
0x00400597 7411      je 0x4005aa
```

```
0x400599 [oc]
0x00400599 be58354000    mov esi, str._unbreakable_enterprise_product_activation_product_key
0x0040059e bf01000000    mov edi, 1
0x004005a3 31c0      xor eax, eax
0x004005a5 e86fffff     call sym.imp.errx;[ob]
```

I don't want to manually reverse this.

```
0x4005aa [oBe]
0x004005aa 488b7608    mov rsi, qword [rsi + 8]
0x004005ae ba43000000    mov edx, 0x43
0x004005b3 bfc0426000    mov edi, 0x6042c0
0x004005b8 e853fffff     call sym.imp.strncpy;[od]
0x004005bd 31c0      xor eax, eax
0x004005bf e82c220000    call fcn.004027f0;[oe]
0x004005c4 31c0      xor eax, eax
0x004005c6 e855220000    call fcn.00402830;[of]
0x004005cb 31c0      xor eax, eax
0x004005cd e89e220000    call fcn.00402870;[og]
0x004005d2 31c0      xor eax, eax
0x004005d4 e8e7220000    call fcn.004028c0;[oh]
0x004005d9 31c0      xor eax, eax
0x004005db e830230000    call fcn.00402910;[oi]
0x004005e0 31c0      xor eax, eax
0x004005e2 e859230000    call fcn.00402940;[oj]
0x004005e7 31c0      xor eax, eax
0x004005e9 e892230000    call fcn.00402980;[ok]
0x004005ee 31c0      xor eax, eax
0x004005f0 e8cb230000    call fcn.004029c0;[ol]
0x004005f5 31c0      xor eax, eax
0x004005f7 e884240000    call fcn.00402a00;[om]
0x004005fc 31c0      xor eax, eax
0x004005fe e83d240000    call fcn.00402a40;[on]
0x00400603 31c0      xor eax, eax
0x00400605 e866240000    call fcn.00402a70;[oo]
0x0040060a 31c0      xor eax, eax
0x0040060c e89f240000    call fcn.00402ab0;[op]
0x00400611 31c0      xor eax, eax
0x00400613 e888240000    call fcn.00402af0;[oq]
0x00400618 31c0      xor eax, eax
0x0040061a e801250000    call fcn.00402b20;[or]
0x0040061f 31c0      xor eax, eax
0x00400621 e82a250000    call fcn.00402b50;[os]
0x00400626 31c0      xor eax, eax
0x00400628 e863250000    call fcn.00402b90;[ot]
0x0040062d 31c0      xor eax, eax
0x0040062f e88c250000    call fcn.00402bc0;[ou]
```

Too long !

```
0x4005aa [oBe]
mov rsi, qword [rsi + 8]
mov edx, 0x43
mov edi, 0x6042c0
call sym.imp.strncpy;[od]
xor eax, eax
call fcn.004027f0;[oe]
xor eax, eax
call fcn.00402830;[of]
xor eax, eax
call fcn.00402870;[og]
xor eax, eax
call fcn.004028c0;[oh]
xor eax, eax
call fcn.00402910;[oi]
xor eax, eax
call fcn.00402940;[oj]
xor eax, eax
call fcn.00402980;[ok]
xor eax, eax
call fcn.004029c0;[ol]
xor eax, eax
call fcn.00402a00;[om]
xor eax, eax
call fcn.00402a40;[on]
xor eax, eax
call fcn.00402a70;[oo]
xor eax, eax
call fcn.00402ab0;[op]
xor eax, eax
call fcn.00402af0;[oq]
xor eax, eax
call fcn.00402b20;[or]
xor eax, eax
call fcn.00402b50;[os]
xor eax, eax
call fcn.00402b90;[ot]
xor eax, eax
call fcn.00402bc0;[ou]
```

REASON - 2 - 0 - 2 - 0

Why did you make this? Why should I use it?

- ???
- `ESILSolve` will likely never have the stability or sophistication of established tools. But...
- Tight integration with `r2` makes it usable with all different kinds of IO plugins in situations that might be awkward for more monolithic tools
- Support for architectures not covered by other tools
- It is a great way to learn about symbolic execution if you are familiar with `r2` and `ESIL`, and is a simple but solid framework to test out novel ideas and strategies
- Good for CTFs :)

Symbolic Execution - A very short review

- Symbolic Execution is the emulation of a program with the ability to make some data, memory or register values, symbolic
- Using symbolic values allows exploration of code paths without needing to know concrete values that actually reach them
- Taking branches constrains the symbolic values and an SMT solver (ESILSolve uses Z3) is used to evaluate the resulting expressions to get concrete values
- Some SymEx tools tools you may have heard of and used: angr, KLEE, MAYHEM, more recently SymCC, Sys

Z3 simplified beyond usefulness

```
In [1]: import z3
```

```
In [2]: x = z3.BitVec("x", 32)
```

```
In [3]: y = z3.BitVec("y", 32)
```

```
In [4]: 3*x + y  
Out[4]: 3*x + y
```

```
In [5]: z3.simplify(3*x + x + y)  
Out[5]: 4*x + y
```

```
In [6]: z = z3.If(x/2 == 27, 4*x, x*y)
```

```
In [7]: z  
Out[7]: If(x/2 == 27, 4*x, x*y)
```

```
In [8]: solver = z3.Solver()
```

```
In [9]: solver.add(x == 54)
```

```
In [10]: solver.check()  
Out[10]: sat
```

```
In [11]: model = solver.model()
```

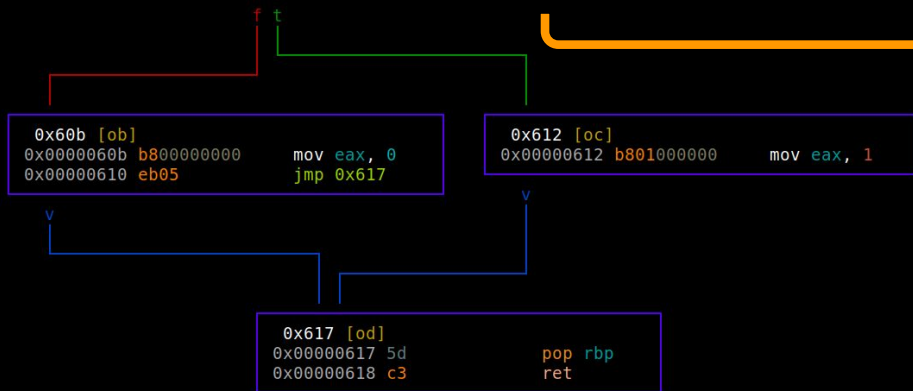
```
In [13]: model.eval(z)  
Out[13]: 216
```

```
In [14]: 54*4  
Out[14]: 216
```

- Z3 can solve expressions involving Ints, Floats, and most importantly Bit Vectors
- Symbolic values are initialized and used in arithmetic and logical expressions
- Solvers (or Optimizers) can create models that satisfy the added constraints

Simple Branch

```
[0x5fa]
31: int main(uint32_t argc, char **argv);
; var char **var_10h @ rbp-0x10
; var uint32_t var_4h @ rbp-0x4
; arg uint32_t argc @ rdi
; arg char **argv @ rsi
0x000005fa 55      push rbp
0x000005fb 4889e5    mov rbp, rsp
0x000005fe 897dfc    mov dword [var_4h], edi
0x00000601 488975f0  mov qword [var_10h], rsi
0x00000605 837dfc02  cmp dword [var_4h], 2
0x00000609 7507      jne 0x612
```



```
expr: rbp,8,rbp,-,[8],8,rbp,-=
000000000000005fa: push rbp
```

```
expr: rsp,rbp,=
000000000000005fb: mov rbp, rsp
```

```
expr: edi,0x4,rbp,-,[4]
000000000000005fe: mov dword [rbp - 4], edi
```

```
expr: rsi,0x10,rbp,-,[8]
00000000000000601: mov qword [rbp - 0x10], rsi
```

```
expr: 2,0x4,rbp,-,[4],==,$z,zf,:=,32,$b,cf,:=,$p,pf,:=,31,$s,sf,:=,2
00000000000000605: cmp dword [rbp - 4], 2
```

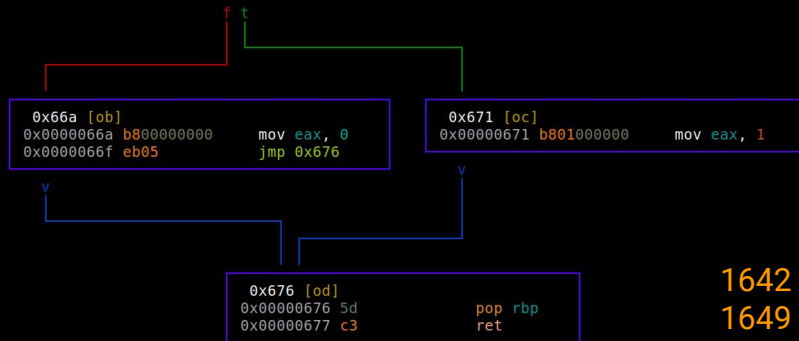
```
expr: zf,!,{,1554,rip,=,}
00000000000000609: jne 0x612
condition val: If(Extract(31, 0, rdi) == 2, 0, 1)
symbolic pc: If(If(Extract(31, 0, rdi) == 2, 0, 1) == 0, 1547, 1554)
```

1547 == 0x60b

1554 == 0x612

Simplish Branch

```
[0x63a]
62: sym.check (int64 t arg1);
; var int64 t var_4h @ rbp-0x4
; arg int64 t arg1 @ rdi
0x0000063a 55      push rbp
0x0000063b 4889e5   mov rbp, rsp
0x0000063e 897dfc   mov dword [var_4h], edi
0x00000641 8b45fc   mov eax, dword [var_4h]
0x00000644 35bebafeca xor eax, 0xcafebabe
0x00000649 89c1     mov ecx, eax
0x0000064b babdbcc8ef mov edx, 0xefc8bcbd
0x00000650 89c8     mov eax, ecx
0x00000652 f7e2     mul edx
0x00000654 89d0     mov eax, edx
0x00000656 c1e80e   shr eax, 0xe
0x00000659 69c054440000 imul eax, eax, 0x4454
0x0000065f 29c1     sub ecx, eax
0x00000661 89c8     mov eax, ecx
0x00000663 3de7110000 cmp eax, 0x11e7
0x00000668 7507     jne 0x671
```



1642 == 0x66a
1649 == 0x671

```
symbolic pc: If(If(Concat(~Extract(31, 30, rdi),
Extract(29, 28, rdi),
~Extract(27, 27, rdi),
Extract(26, 26, rdi),
~Extract(25, 25, rdi),
Extract(24, 24, rdi),
~Extract(23, 17, rdi),
Extract(16, 16, rdi),
~Extract(15, 15, rdi),
Extract(14, 14, rdi),
~Extract(13, 11, rdi),
Extract(10, 10, rdi),
~Extract(9, 9, rdi),
Extract(8, 8, rdi),
~Extract(7, 7, rdi),
Extract(6, 6, rdi),
~Extract(5, 1, rdi),
Extract(0, 0, rdi)) ==
```

```
4583 +
17492*
Concat(0,
Extract(63,
46,
4022910141*
Concat(0,
~Extract(31, 30, rdi),
Extract(29, 28, rdi),
~Extract(27, 27, rdi),
Extract(26, 26, rdi),
~Extract(25, 25, rdi),
Extract(24, 24, rdi),
~Extract(23, 17, rdi),
Extract(16, 16, rdi),
~Extract(15, 15, rdi),
Extract(14, 14, rdi),
~Extract(13, 11, rdi),
Extract(10, 10, rdi),
~Extract(9, 9, rdi),
Extract(8, 8, rdi),
~Extract(7, 7, rdi),
Extract(6, 6, rdi),
~Extract(5, 1, rdi),
Extract(0, 0, rdi))))),
```

```
0,
1) ==
0,
1642,
1649)
```


ESIL - Another even shorter review

- ESIL is radare2's built in intermediate language
- It uses a simple, stack based system to express any instruction as a string of values and operators
- Example: `add x0, x1, x2` in AArch64 becomes `"x2,x1,+,x0,="`
- It's simple, readable, and great
- It is an interesting IR for symbolic execution
- There is a really great in depth talk on ESIL from r2con 2019 by @arnaugamez

```

; CALL XREF from main @ 0x6d8
117: sym.check(int64_t arg1);
; var int64_t var_4h @ rbp-0x4
; arg int64_t arg1 @ rdi
0x0000063a 55          push rbp
0x0000063b 4889e5      mov rbp, rsp
0x0000063e 897dfc      mov dword [var_4h], edi
0x00000641 8b45fc      mov eax, dword [var_4h]
0x00000644 35feca0000 xor eax, 0xcafe
0x00000649 89c1        mov ecx, eax
0x0000064b ba5f5ee477 mov edx, 0x77e45e5f
0x00000650 89c8        mov eax, ecx
0x00000652 f7ea        imul edx
0x00000654 c1fa0d      sar edx, 0xd
0x00000657 89c8        mov eax, ecx
0x00000659 c1f81f      sar eax, 0x1f
0x0000065c 29c2        sub edx, eax
0x0000065e 89d0        mov eax, edx
0x00000660 69c054440000 imul eax, eax, 0x4454
0x00000666 29c1        sub ecx, eax
0x00000668 89c8        mov eax, ecx
0x0000066a 3d53140000 cmp eax, 0x1453
0x0000066f 7537        jne 0x6a8
0x00000671 8b45fc      mov eax, dword [var_4h]
0x00000674 35beba0000 xor eax, 0xbabe
0x00000679 89c1        mov ecx, eax
0x0000067b ba5f5ee477 mov edx, 0x77e45e5f
0x00000680 89c8        mov eax, ecx
0x00000682 f7ea        imul edx
0x00000684 c1fa0d      sar edx, 0xd
0x00000687 89c8        mov eax, ecx
0x00000689 c1f81f      sar eax, 0x1f
0x0000068c 29c2        sub edx, eax
0x0000068e 89d0        mov eax, edx
0x00000690 69c054440000 imul eax, eax, 0x4454
0x00000696 29c1        sub ecx, eax
0x00000698 89c8        mov eax, ecx
0x0000069a 3dbf1f0000 cmp eax, 0x1fbf
0x0000069f 7507        jne 0x6a8
0x000006a1 b800000000 mov eax, 0
0x000006a6 eb05        jmp 0x6ad
; CODE XREFS from sym.check @ 0x66f, 0x69f
0x000006a8 b801000000 mov eax, 1
; CODE XREF from sym.check @ 0x6a6
0x000006ad 5d          pop rbp
0x000006ae c3          ret

; rbp,8,rsp,-,[8],8,rsp,-=
; rsp,rbp,=
; edi,0x4,rbp,-,[4]; arg1
; 0x4,rbp,-,[4],rax,=
; 51966,rax,^,0xffffffff&,rax,=,$z,zf,=,$p,pf,=,31,$s,sf,=,0,cf,=,0,of,=
; eax,rcx,=
; 2011455071,rdx,=
; ecx,rax,=
; 32,32,edx,~,32,eax,~,*,>>,edx,=,edx,eax,*,32,32,eax,~,>>,edx,-,{1,1},{0,0},cf,=,of,=
; 0,cf,=,1,13,-,1,<<,edx,&,{1,cf,=,},13,edx,>>>>,edx,=,$z,zf,=,$p,pf,=,31,$s,sf,=
; ecx,rax,=
; 0,cf,=,1,31,-,1,<<,eax,&,{1,cf,=,},31,eax,>>>>,eax,=,$z,zf,=,$p,pf,=,31,$s,sf,=
; eax,edx,-=,eax,0x80000000,-,!31,$o,^,of,=,31,$s,sf,=,$z,zf,=,$p,pf,=,32,$b,cf,=,3,$b,af,=
; edx,rax,=
; 32,17492,~,32,eax,~,*,DUP,eax,=,eax,-,{1,1},{0,0},cf,=,of,=
; eax,ecx,-=,eax,0x80000000,-,!31,$o,^,of,=,31,$s,sf,=,$z,zf,=,$p,pf,=,32,$b,cf,=,3,$b,af,=
; ecx,rax,=
; 5203,eax,=,$z,zf,=,32,$b,cf,=,$p,pf,=,31,$s,sf,=,5203,0x80000000,-,!31,$o,^,of,=,3,$b,af,=
; zf,!,{1704,rip,=,}
; 0x4,rbp,-,[4],rax,=
; 47806,rax,^,0xffffffff&,rax,=,$z,zf,=,$p,pf,=,31,$s,sf,=,0,cf,=,0,of,=
; eax,rcx,=
; 2011455071,rdx,=
; ecx,rax,=
; 32,32,edx,~,32,eax,~,*,>>,edx,=,edx,eax,*,32,32,eax,~,>>,edx,-,{1,1},{0,0},cf,=,of,=
; 0,cf,=,1,13,-,1,<<,edx,&,{1,cf,=,},13,edx,>>>>,edx,=,$z,zf,=,$p,pf,=,31,$s,sf,=
; ecx,rax,=
; 0,cf,=,1,31,-,1,<<,eax,&,{1,cf,=,},31,eax,>>>>,eax,=,$z,zf,=,$p,pf,=,31,$s,sf,=
; eax,edx,-=,eax,0x80000000,-,!31,$o,^,of,=,31,$s,sf,=,$z,zf,=,$p,pf,=,32,$b,cf,=,3,$b,af,=
; edx,rax,=
; 32,17492,~,32,eax,~,*,DUP,eax,=,eax,-,{1,1},{0,0},cf,=,of,=
; eax,ecx,-=,eax,0x80000000,-,!31,$o,^,of,=,31,$s,sf,=,$z,zf,=,$p,pf,=,32,$b,cf,=,3,$b,af,=
; ecx,rax,=
; 8127,eax,=,$z,zf,=,32,$b,cf,=,$p,pf,=,31,$s,sf,=,8127,0x80000000,-,!31,$o,^,of,=,3,$b,af,=
; zf,!,{1704,rip,=,}
; 0,rax,=
; 0x6ad,rip,=
; 1,rax,=

; rsp,[8],rbp,=,8,rsp,+=
; rsp,[8],rip,=,8,rsp,+=

```

[0x0000063a]>

<code>push rbp</code>	<code>; rbp,8, rsp, -,=[8],8, rsp, -=</code>
<code>mov rbp, rsp</code>	<code>; rsp,rbp,=</code>
<code>mov dword [var_4h], edi</code>	<code>; edi,0x4,rbp, -,=[4] ; arg1</code>
<code>mov eax, dword [var_4h]</code>	<code>; 0x4,rbp, -, [4],rax,=</code>

Challenges of ESIL based SymEx

- ESIL is a lot less verbose / precise than other IR
- “add x0, x1, x2” in VEX ->
- Every value and operation has an explicit size, type, and signedness in other IR
- Turns out they are being overly cautious dorks

```
IRSB {  
    t0:Ity_I64 t1:Ity_I64 t2:Ity_I64  
    t3:Ity_I64 t4:Ity_I64  
  
    00 | ----- IMark(0x400400, 4, 0) -----  
    01 | t0 = GET:I64(x1)  
    02 | t1 = GET:I64(x2)  
    03 | t3 = Add64(t0,t1)  
    04 | PUT(x0) = t3  
    NEXT: PUT(pc) = 0x00000000000400404;  
    Ijk_Boring  
}
```

Challenges of ESIL based SymEx Part II

- <show main parse_expression method>
- Other IR have data only in temporary registers and variables
- ESILSolve needs to have a *symbolic stack*, where for every conditional both IF and ELSE expressions are executed, their stacks are unwound, and a new stack with `If(CONDITION, IF_VALUE, ELSE_VALUE)` values is pushed
- The GOTO operator essentially uses arbitrarily nested conditionals which creates a host of issues, most of which are soluble
- These operations happen to be relatively efficient ones for python

The ESILSolve API

- ESILSolver
 - call_state, init_state, blank_state - create an initial state to begin running
 - register_hook, register_sim - register hooks or simulated methods to manipulate state
 - run, terminate - start symbolically executing, finding find, avoiding avoid

```
from esilsolve import ESILSolver
import r2pipe
import z3

esilsolver = ESILSolver("ais3_crackme")
state = esilsolver.call_state("sym.verify")

addr = 0x1000000
state.registers["rdi"] = addr
flag = z3.BitVec("flag", 24*8)
state.memory[addr] = flag

def check(state):
    state.constrain(state.registers["zf"] == 1)
    if state.is_sat():
        flag_str = state.evaluate_string(flag)
        print("FLAG: %s " % flag_str)
        esilsolver.terminate()

esilsolver.register_hook(0x004005bd, check)
state = esilsolver.run()
```

The ESILSolve API

- ESILState
 - registers - an instance of ESILRegisters with the register values for this state
 - memory - an instance of ESILMemory for this state
 - solver - an instance of a Z3 solver class (Solver, SimpleSolver, Optimizer)
 - constrain - an alias for solver.add
 - constrain_bytes - method to constrain the bytes of a buffer to match a regex range
 - evaluate - model and eval a symbolic expression if it is satisfiable
 - evaluate_buffer, state.evaluate_string - convenience methods for evaluating a BV and casting it to bytes / string
 - clone - make a (COW) copy of the state
 - step - single step the state, executing the current instruction

Demo ESILSolve

Show some cool stuff now

ESILSolve plugin for r2

```
[0x00400729]> aesx?
```

```
Usage: aesx[iscxrelda] # Core plugin for ESILSolve
```

aesxi [debug] [lazy]	Initialize the ESILSolve instance and VM
aesxs[bc] reg addr [name] [length]	Set symbolic value in register or memory
aesxv reg addr value	Set concrete value in register or memory
aesxc sym value	Constrain symbol to be value, min, max, regex
aesxx[ec] expr value	Execute ESIL expression and evaluate/constrain the result
aesxr[ac] target [avoid x,y,z]	Run symbolic execution until target address, avoiding x,y,z
aesxe[j] sym1 [sym2] [...]	Evaluate symbol in current state
aesxb[j] sym1 [sym2] [...]	Evaluate buffer in current state
aesxd[j] [reg1] [reg2] [...]	Dump register values / ASTs
aesxa	Apply the current state, setting registers and memory

```
[0x00400729]> █
```

ESILSolve plugin II

```
undefined8 fcn.004006fd(int64_t arg1)
```

```
{  
    int64_t var_38h;  
    char *var_24h;  
    char *var_18h;  
    char *var_10h;  
  
    stack0xfffffffffffffd8 = "Dufhbm";  
    var_18h = "pG`imos";  
    var_10h = "ewUglpt";  
    var_24h._0_4_ = 0;  
    while( true ) {  
        if (0xb < (int32_t)var_24h) {  
            // target: flag = b'Code_Talkers'  
            return 0;  
        }  
        if (((int32_t)*(char *)((int64_t)((int32_t)var_24h / 3) * 2) +  
            *(int64_t *)((int64_t)&var_24h + (int64_t)((int32_t)var_24h % 3) * 8 + 4)) -  
            (int32_t)*(char *) (arg1 + (int32_t)var_24h) != 1) break;  
        var_24h._0_4_ = (int32_t)var_24h + 1;  
    }  
    return 1;  
}
```

Prove long strings of arithmetical
instructions equivalent to simpler ones

Automatically add comments that show concrete values
that reach any point in code, for every symbol

```
[0x00000610]> s sym.rotateLeft  
[0x000007f2]> pdf  
; CALL XREF from main @ 0x78d  
89: sym.rotateLeft (int64_t arg1, int64_t arg2);  
; var int64_t var_18h @ rbp-0x18  
; var int64_t var_14h @ rbp-0x14  
; var int64_t var_4h @ rbp-0x4  
; arg int64_t arg1 @ rdi  
; arg int64_t arg2 @ rsi  
0x000007f2 55 push rbp  
0x000007f3 4889e5 mov rbp, rsp  
0x000007f6 897dec mov dword [var_14h], edi  
0x000007f9 8975e8 mov dword [var_18h], esi  
0x000007fc 8b4de8 mov ecx, dword [var_18h]  
0x000007ff ba85104208 mov edx, 0x8421085  
0x00000804 89c8 mov eax, ecx  
0x00000806 f7e2 mul edx  
0x00000808 89c8 mov eax, ecx  
0x0000080a 29d0 sub eax, edx  
0x0000080c d1e8 shr eax, 1  
0x0000080e 01d0 add eax, edx  
0x00000810 c1e804 shr eax, 4  
0x00000813 89c2 mov edx, eax  
0x00000815 89d0 mov eax, edx  
0x00000817 c1e005 shl eax, 5  
0x0000081a 29d0 sub eax, edx  
0x0000081c 29c1 sub ecx, eax  
0x0000081e 89c8 mov eax, ecx  
0x00000820 8945e8 mov dword [var_18h], eax  
0x00000823 eb14 jmp 0x839  
; CODE XREF from sym.rotateLeft @ 0x844  
0x00000825 8b45ec mov eax, dword [var_14h]  
0x00000828 c1e81f shr eax, 0x1f  
0x0000082b 8945fc mov dword [var_4h], eax  
0x0000082e 8b45ec mov eax, dword [var_14h]  
0x00000831 01c0 add eax, eax  
0x00000833 0b45fc or eax, dword [var_4h]  
0x00000836 8945ec mov dword [var_14h], eax  
; CODE XREF from sym.rotateLeft @ 0x823  
0x00000839 8b45e8 mov eax, dword [var_18h]  
0x0000083c 8d50ff lea edx, [rax - 1]  
0x0000083f 8955e8 mov dword [var_18h], edx  
0x00000842 85c0 test eax, eax  
0x00000844 75df jne 0x825  
0x00000846 8b45ec mov eax, dword [var_14h]  
0x00000849 5d pop rbp  
0x0000084a c3 ret  
  
[0x000007f2]> aeim  
[0x000007f2]> aei  
[0x000007f2]> aesxi  
[0x000007f2]> aesxs rdi num 8  
[0x000007f2]> aesxv rsi 8  
[0x000007f2]> aesxr 0x00000849  
[0x000007f2]> aesxxc "esi,edi,<<<,eax,-,!" 0  
[0x000007f2]> aesxe eax  
error: state has unsatisfiable constraints
```

Using ESILSolve to improve ESIL

- Symbolic execution of ESIL allows us to convert the string expression to an AST which can be compared to the ASTs generated by other symex tools
- If expressions are not equivalent the SMT solver can give example values that produce different results in the two systems
- ESILCheck is a tool which compares the ASTs generated from ESIL expressions and the ones produced by angr's VEX execution engine
- The claripy backend objects provide conversion to Z3 enabling easy testing
- Other engines, like KLEE, could provide additional checks and coverage, any tool supporting the same SMT solvers can be used.
- A tool called VEX2ESIL also comes with ESILSolve

Future Goals

- Use `ESILSolve` to prove the equality of `ESIL` expressions and other IR expressions, or fix them if necessary, to cover more architectures more faithfully
- Make a modest Sim Procedure library so that common libc functions do not cause state explosion
- Create Sims for syscalls and hopefully have syscalls handled natively in concrete emulation in r2
- Create a fuzzer leveraging `ESILSolve`, concrete `ESIL` emulation in r2, and maybe r2frida to generate amazing coverage guided fuzzing that is snapshot emulated for speed but in-memory fuzzed for accuracy

Thanks To:

Pancake and everyone who has worked on ESIL and radare2 !



Thank You