

# Optimization

732A90

Computational Statistics

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# Plan for today

- Introduction
- Mathematical definition of problem
- 1D optimization, function `optimize()`
- $k$ D optimization, function `optim()`

# Optimization

Nearly everything is optimization

- Chemistry
- Physics
- Economics, **Industry**
- Engineering

## **BUT EVEN**

- Your mobile price plan
- Course scheduling
- Your lunch choice

## **STATISTICS**

- Fit parameters to data
- Propose optimal decision

## Industry

How to produce a cylindrical  $0.5L$  beer can so it requires minimum material?

Given a certain product minimize e.g. material usage, production effort while still meeting consumer requirements.

## Economics/Logistics

- Travelling Salesman Problem
- Windmills
- Flight schedule

# Optimization

- Mathematical optimization
  - Linear optimization
  - Non-linear optimization
  - Combinatorial optimization
  - Integer optimization
  - ...

**Not this lecture**

- Computational approach
    - Algorithms, heuristics
- Convergence ?

**This lecture**

# Mathematical formulation

The goal is to minimize (maximize)

**Objective function:**  $f(\theta)$

(reproduction, chances of survival, quality of life, cost, profit, likelihood, fit to data)

depending on

**Parameters or Unknowns  $\theta$**

(reproduction strategy, resource utilization, consumer choices, height & diameter, production, raw material choice, service times, route, flight routes/times ,parameters)

# Mathematical formulation

$$\min_{\theta \in \Theta} f(\theta) \quad \text{subject to} \quad \begin{array}{ll} c_i(\theta) = 0, & i \in E \\ c_i(\theta) \geq 0, & i \in I \end{array}$$

**QUESTION:** What should we do if we are interested in maximization instead of minimization?

**QUESTION:** What should we do if the constraints are  $c_i(x) \leq 0, i \in I$ ?



# Constraints examples

- Available environment
- Volume: 0.5l of can
- Production: Factories ( $F_1, F_2$ ), retail outlets ( $R_1, R_2, R_3$ ), cost of shipping  $i \rightarrow j$ :  $c_{ij}$ , production  $a_i$  per week, requirement  $b_j$  per week **to optimize:**  $x_{ij}$  amount shipped  $i \rightarrow j$  per week

$$\min_{x \in \mathbb{R}^3} \sum_{ij} c_{ij} x_{ij} \quad \text{minimize shipping costs}$$

$$\sum_{j=1}^3 x_{ij} \leq a_i, i = 1, 2 \quad \text{production capacity}$$

$$\sum_{i=1}^3 x_{ij} \geq b_j, j = 1, 2, 3 \quad \text{demand}$$

$$\forall_{i,j} x_{ij} \geq 0$$

# Optimization approaches

- Constrained optimization
  - Lagrange multipliers, linear programming
  - E.g. LASSO
  - **Not this lecture**
- Unconstrained optimization
  - Steepest descent
  - Newton method
  - Quasi-Newton-Methods
  - Conjugate gradients
  - **This lecture**

# 1D Optimization: analytical approach

- To minimize / maximize

$$f(x), \quad x \in \mathbb{R}$$

$x$  - one dimensional variable

- Compute first derivative:  $\frac{\partial f}{\partial x}$
- Find  $x^*$  for which first derivative is 0
- Compute second derivative:  $\frac{\partial^2 f}{\partial x^2}$
- Compute value of second derivative at  $x = x^*$ 
  - Second derivative positive  
 $\Rightarrow x^*$  local minimum
  - Second derivative negative  
 $\Rightarrow x^*$  local maximum
  - Second derivative 0  
 $\Rightarrow x^*$  saddle point

# 1D Optimization: computational approach

- To minimize / maximize

$$f(x), \quad x \in \mathbb{R}$$

$x$  - one dimensional variable

- Algorithms

*Golden-Section Search* (see next slide)

- local minimum / maximum on interval  $[A, B]$
- Works by narrowing down the search interval with a constant reduction factor

$$1 - \alpha = \frac{\sqrt{5} - 1}{2} \approx 0.62$$

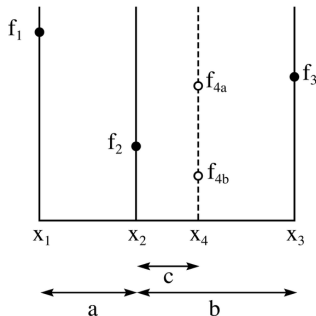
- In **R**:

function `optimize()`

- Brent's method (modified Golden-Section Search)

# Golden-Section Search (minimization)

```
1:  $x_1 = A, x_3 = B,$   
2: while  $x_1 - x_3 > \epsilon$  do  
3:    $a = \alpha(x_3 - x_1)$   
4:    $x_2 = x_1 + a, x_4 = x_3 - a$   
5:   if  $f(x_4) > f(x_2)$  then  
6:      $x_1 = x_1, x_3 = x_4$   
7:   else  
8:      $x_1 = x_2, x_3 = x_3$   
9:   end if  
10: end while
```



Wikipedia, Golden-Section Search

- $f$  has one local minimum  $\rightarrow$  minimum will be found
- $f$  has several local minima  $\rightarrow$  one of them will be found

# 1D Optimization: Example

```
f <- function (x, a) (x - a)^2
xmin <- optimize(f, c(0, 1), tol = 0.0001, a =
  1/3)
xmin
```

```
## "wrong" solution with unlucky interval and
  piecewise constant f():
f <- function(x) ifelse(x > -1, ifelse(x < 4,
  exp(-1/abs(x - 1)), 10), 10)
plot(f, -20, 20)
xmin1<-optimize(f, c(-4, 20))    # doesn't see
  the minimum
xmin1
xmin2<-optimize(f, c(-7, 20))    # ok
xmin2
```

# $k$ D Optimization: analytical approach

- To minimize / maximize

$$f(\vec{x}), \quad \vec{x} = (x_1, \dots, x_n)^\top$$

$\vec{x}$  -  $n$ -dimensional vector

- Compute gradient:

$$\nabla f(\vec{x}) = \left( \frac{\partial f(\vec{x})}{\partial x_1}, \dots, \frac{\partial f(\vec{x})}{\partial x_n} \right)^\top$$

- vector of all partial derivatives

- Find  $\vec{x}^*$  for which gradient is  $\mathbf{0}$ 
  - all partial derivatives are 0

# $k$ D Optimization: analytical approach

- Compute Hessian matrix:

$$\nabla^2 f(\vec{x}) = \left[ \frac{\partial^2 f(\vec{x})}{\partial x_i \partial x_j} \right]_{i,j=1}^n$$

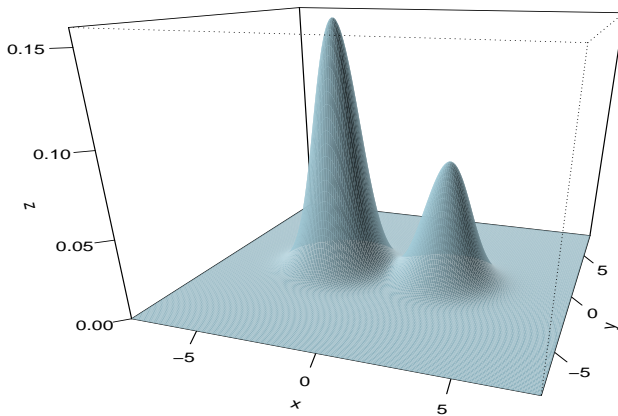
- matrix of all second partial derivatives

- Compute values of Hessian matrix at  $x = x^*$ 
  - Hessian matrix positive definite  
 $\Rightarrow \vec{x}^*$  local minimum
  - Hessian matrix negative definite  
 $\Rightarrow \vec{x}^*$  local maximum
  - Otherwise  
 $\Rightarrow \vec{x}^*$  saddle point



## Statistics

Maximize likelihood



# Maximal (maximum) likelihood

$X_1, \dots, X_n$  i.i.d. - sample is drawn from probability distribution  $P(X|\theta)$ , where  $\theta$  is unknown parameter set

$P(X|\theta)$  - probability function (for discrete random variables)  
- density function (for continuous random variables)

The joint probability / density function for all observations:

$$P(X_1, \dots, X_n|\theta) = \prod_{i=1}^n P(X_i|\theta)$$

Find  $\theta$  that maximizes  $P(X_1, \dots, X_n|\theta)$

**Example:**  $X_1, \dots, X_n \sim N(\mu, \sigma^2) \Rightarrow \theta = (\mu, \sigma^2)$

## General strategy

- 1 Provide a (*good*) starting point  $\vec{x}_0$ ,  
 $\vec{x} = \vec{x}_0$
- 2 Choose a direction  $\vec{p}$  ( $\|\vec{p}\| = 1$ ) and step size  $\alpha$
- 3 Move to  $\vec{x} := \vec{x} + \alpha\vec{p}$
- 4 Repeat step 2 until convergence

# Choice of direction

## Taylor's theorem

$$f(\vec{x} + \alpha\vec{p}) = f(\vec{x}) + \boxed{\alpha\vec{p}^\top \cdot \nabla f(\vec{x})} + o(\alpha^2)$$

$\vec{p}$  s.t.  $\vec{p}^\top \cdot \nabla f(\vec{x}) < 0$  is a *descent* direction.

***Steepest*** descent is

$$\vec{p} = -(\nabla f(\vec{x})) / \|\nabla f(\vec{x})\|$$

*Hessian matrix ignored in steepest descent*

# Choice of step size

- **Expensive way:** find the global minimum in direction  $\vec{p}$
- **Trade-off way:** find a decrease which is *sufficient*

## BACKTRACKING

- 1: Choose (large)  $\alpha_0 > 0$ ,  $\rho \in (0, 1)$ ,  $c \in (0, 1)$ ,
- 2:  $\alpha = \alpha_0$
- 3: **repeat**
- 4:    $\alpha = \rho\alpha$
- 5: **until**  $f(\vec{x} + \alpha\vec{p}) \leq f(\vec{x}) + c\alpha\vec{p}^\top \nabla f(\vec{x})$

# Newton-Raphson Method

- If  $f$  is quadratic

$$f(\vec{p}) = \frac{1}{2} \vec{p}^\top \mathbf{A} \vec{p} + \vec{b}^\top \vec{p} + c,$$

then minimum

$$\vec{p}^* = \mathbf{A}^{-1} \vec{b}.$$

- Taylor expansion of  $f$

$$f(\vec{x} + \alpha \vec{p}) = f(\vec{x}) + \alpha \vec{p}^\top \cdot \nabla f(\vec{x}) + \frac{\alpha^2}{2} \vec{p}^\top \nabla^2 f(\vec{x}) \vec{p} + o(\alpha^3)$$

- $x := x + \alpha \vec{p}$  where

$$\vec{p} = - (\nabla^2 f(\vec{x}))^{-1} \nabla f(\vec{x})$$

# Newton-Raphson Method

”\_”

- $(\nabla^2 f(\vec{x}))^{-1}$  is expensive to compute  
→ quicker approaches  
e.g. Cholesky decomposition
- Hessian should be *positive definite* for  $\vec{p}$  to be a descent direction (if not see book)
- Memory expensive:  
need to store  $O(n^2)$  elements

”+”

- Method converges quickly  
esp. near optimum

# Quasi-Newton Methods

- $k$  iteration number
- Compute an approximation to the Hessian,  $\mathbf{B}$ , that will allow for efficient choice of  $\vec{p}$ .
- **SECANT CONDITION:** (quasi-Newton condition)

$$\mathbf{B}_{k+1} (\vec{x}_{k+1} - \vec{x}_k) = \nabla f(\vec{x}_{k+1}) - \nabla f(\vec{x}_k)$$

## BFGS Algorithm

- 1: Choose  $\mathbf{B}_0 > 0$ ,  $\vec{x}_0$ ,  $k = 0$
- 2: **repeat**
- 3:    $\vec{p}_k$  is solution of  $\mathbf{B}_k \vec{p}_k = -\nabla f(\vec{x}_k)$ , i.e.  
     $\vec{p}_k = -\mathbf{B}_k^{-1} \nabla f(\vec{x}_k)$
- 4:   find suitable  $\alpha_k$
- 5:    $\vec{x}_{k+1} = \vec{x}_k + \alpha_k \vec{p}_k$
- 6:   calculate  $\mathbf{B}_{k+1}$  {next slide}
- 7:    $k = k + 1$
- 8: **until** convergence of  $\vec{x}_k$  at minimum



# Computing $\mathbf{B}_{k+1}$

- We want  $\mathbf{B}_{k+1}$  and  $\mathbf{B}_k$  to be close to each other:

$$\begin{aligned} \min_{\mathbf{B}} \quad & \|\mathbf{B} - \mathbf{B}_k\| \\ \text{s.t.} \quad & \mathbf{B} = \mathbf{B}^\top \end{aligned}$$

- For  $\vec{y}_k = \nabla f(\vec{x}_{k+1}) - \nabla f(\vec{x}_k)$ ,  $\vec{s}_k = \vec{x}_{k+1} - \vec{x}_k$   
We receive

$$\mathbf{B}_{k+1} = \mathbf{B}_k - \frac{\mathbf{B}_k \vec{s}_k \vec{s}_k^\top \mathbf{B}_k}{\vec{s}_k^\top \mathbf{B}_k \vec{s}_k} + \frac{\vec{y}_k \vec{y}_k^\top}{\vec{y}_k^\top \vec{s}_k}$$

- Sherman–Morrison formula for  $\mathbf{B}_{k+1}^{-1}$
- We have to store  $\mathbf{B}_k^{-1}$

- BFGS: Broyden–Fletcher–Goldfarb–Shanno
- More iterations than Newton's method (uses approximation)
- Each iteration quicker, no numeric inversion
- Choice of  $\mathbf{B}_0$ ?

# Conjugate Gradient Method—quadratic case

Minimize

$$f(\vec{x}) = \frac{1}{2} \vec{x}^\top \mathbf{A} \vec{x} - \vec{b}^\top \vec{x}$$

for  $\mathbf{A}$  symmetric positive definite.

Gradient:

$$\nabla f(\vec{x}) = \mathbf{A} \vec{x} - \vec{b} = r(\vec{x})$$

Two vectors  $\vec{p}$  and  $\vec{q}$  are **conjugate** with respect to  $\mathbf{A}$  if

$$\vec{p}^\top \mathbf{A} \vec{q} = 0$$

Method is based on this property

# Conjugate Gradient Method

- $\vec{p}_0 = \vec{r}_0$
- $\vec{p}_{k+1} = -\vec{r}_k + \beta_{k+1}\vec{p}_k,$

where

$$\beta_{k+1} = \frac{\vec{r}_k^\top \mathbf{A} \vec{p}_{k-1}}{\vec{p}_k^\top \mathbf{A} \vec{p}_k}$$

- Convergence in  $\dim(\mathbf{A})$  steps  
(or unless cutoff for  $\vec{r}_k$ )

# Nonlinear CG Method

- If  $f(\cdot)$  general, use  $\nabla f(\cdot)$  instead of  $r(\cdot)$

1: Choose  $\vec{x}_0$ ,  $\vec{p}_0 = -\nabla f(\vec{x}_0)$ ,  $k = 0$

2: **while**  $\nabla f(\vec{x}_k) \neq \vec{0}$  **do**

3:   find suitable  $\alpha_k$

4:    $\vec{x}_{k+1} = \vec{x}_k + \alpha_k \vec{p}_k$  {and now update step}

5:

---

$$\beta_{k+1} = (\nabla f(\vec{x}_{k+1})^\top \nabla f(\vec{x}_{k+1})) / (\nabla f(\vec{x}_k)^\top \nabla f(\vec{x}_k))$$

{Fletcher–Reeves update, other possible}

---

$$D_{k+1} := \nabla f(\vec{x}_{k+1}) - \nabla f(\vec{x}_k)$$

$$\beta_{k+1} = (\nabla f(\vec{x}_{k+1})^\top D_{k+1}) / (\nabla f(\vec{x}_k)^\top \nabla f(\vec{x}_k))$$

{Polak–Ribière update, other possible}

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6:    $\vec{p}_{k+1} = -\nabla f(\vec{x}_{k+1}) + \beta_{k+1} \vec{p}_k$

7:    $k = k + 1$

8: **end while**

# Nonlinear CG Method

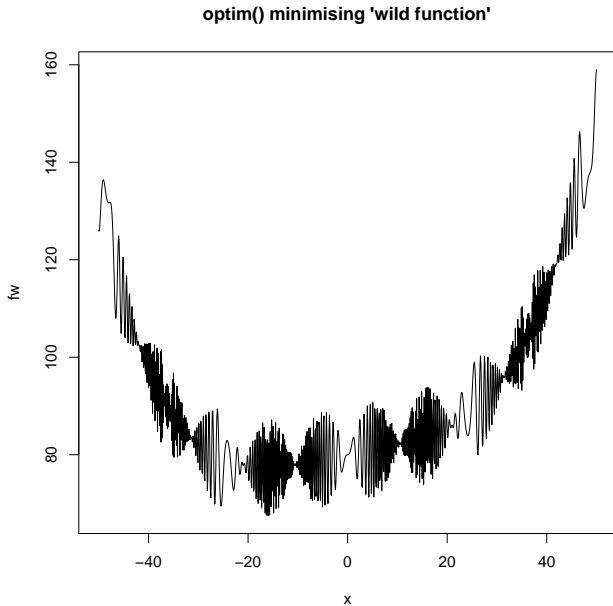
- Local minimum convergence
- Cannot “jump out” of descent path
- Faster than steepest descent
- Slower than Newton and Quasi-Newton but significantly less memory

**Function** `optim()`

Methods:

- Nelder-Mead
- **BFGS**
- **Conjugate Gradient**
- L-BFGS-B
- SANN
- Brent

# *k*D Optimization: Example





## kD Optimization: Example

```
## "wild" function , global minimum at about  
-15.81515  
fw <- function (x)  
      10*sin(0.3*x)*sin(1.3*x^2) + 0.00001*  
      x^4 + 0.2*x+80  
plot(fw, -50, 50, n = 1000, main = "optim() _  
      minimising _ 'wild_function'")  
res <- optim(50, fw, method = "CG", control =  
      list(maxit = 20000, parscale = 20))  
res
```

Thank you for attention!