1 Infinite Words

Theorem 1.1. *Every non-empty ω-regular language contains an ultimately periodic word.*

Theorem 1.2. *For a Kripke structure $\mathcal{K}$ with initial state s and $\varphi \in LTL$, the model checking problem $L(\mathcal{K}, s) \subseteq L(\varphi)$? is PSPACE-complete.*

Proof. **PSPACE** Compute the intersection automaton for $L(\mathcal{K}, s) \cap L(\neg\varphi)$ and test it for emptiness.

PSPACE-hard Encode a poly.-length Turing tape as a Kripke structure and its correct behavior in LTL. □

Theorem 1.3 (Büchi)**.** *The MSO theory of* $(\mathbb{N}, +1, <, 0)$ *is decidable.*

Proof. Corresponds to S1S formula. Can be checked with NBA emptiness test. □

Theorem 1.4. *The FO theory of* $(\mathbb{R}, +, <, 0)$ *is decidable.*

Proof. Encode real numbers x by triples of sets (X_s, X_i, X_f) with the number's sign ($X_s = \emptyset$ or $\{0\}$), the positive decimal digits in binary encoding, and the positive fractional digits in binary encoding. Then an FO sentence can be transformed to an equi-satisfiable MSO sentence over $(\mathbb{N}, +1, <, 0)$. □

Theorem 1.5. *Subset-construction does not suffice to determinize NBAs.*

Theorem 1.6. *For every n, there is $L_n \subseteq \Sigma^\omega$ s.t. there is an NBA that recognizes L_n with $n+2$ states, but every det. Rabin automaton that recognizes L_n has at least $n!$ states.*

Theorem 1.7. *There is a DBA-recog. language which does not have a unique minimal DBA. DBAs minimized with the DFA minimization algorithm can be arbitrarily bad compared to a minimal DBA.*

Theorem 1.8. *Weak DBAs can be minimized uniquely in polynomial time.*

Theorem 1.9. *Given an ABA $\mathcal{A}$, the dual $\tilde{\mathcal{A}}$ is an alternating co-Büchi automaton which accepts $\overline{L(\mathcal{A})}$, with $\tilde{F} = Q \setminus F$ and $\tilde{\delta}$ exchanging true/false and $\wedge/\vee$.*