1 List Of Models

1.1 Infinite Words

- DBA
- NBA
- GBA
- Rabin automaton
- Muller automaton
- Parity automaton
- E automaton
- A automaton
- coBA
- weak BA
- Staiger-Wagner automaton
- ABA
- LTL
- S1S
- ∃S1S
- $S1S_0$

1.2 Finite Trees

- DTA
- NTA
- ↓DTA
- ↓NTA
- DUTA
- NUTA
- deterministic DTD
- DTD
- deterministic EDTD

- single-type EDTD
- EDTD
- Relax NG
- FO
- MSO
- Regular expressions
- DTWA
- TWA

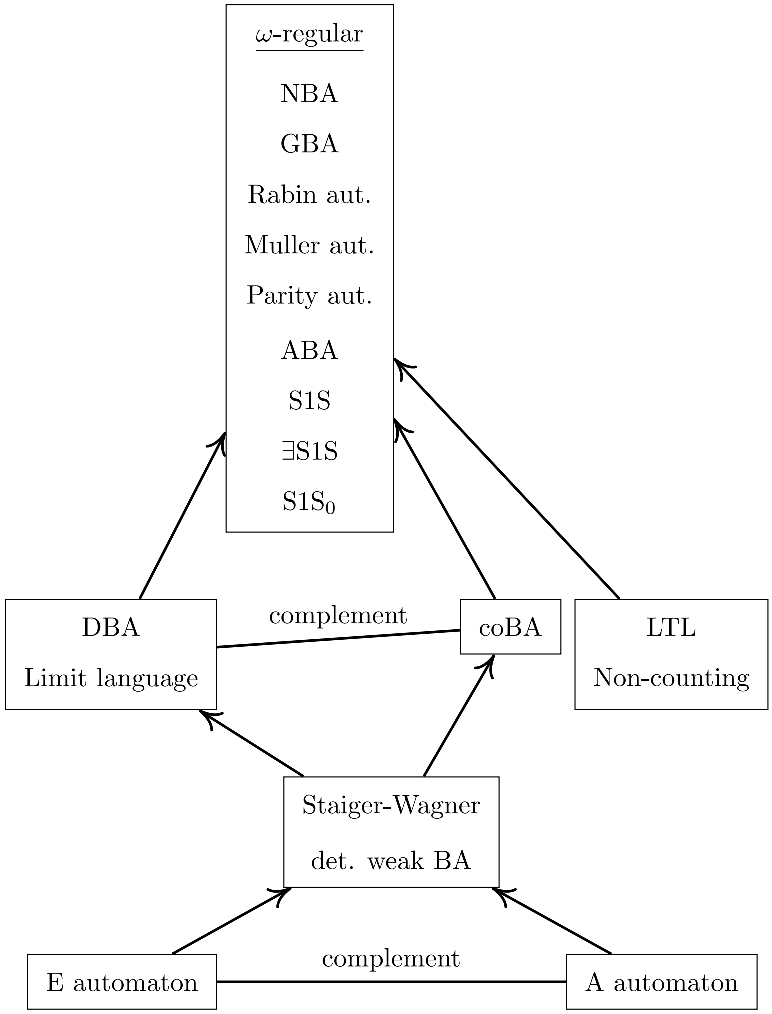
1.3 Infinite Trees

- BTA
- Muller TA
- Parity TA
- DMTA
- S2S (MSO / WMSO)
- $S2S_0$ (MSO / WMSO)

2 List Of Games

- Büchi
- Staiger-Wagner
- weak Parity
- Reachability
- Safety
- Muller
- Parity
- Rabin
- Streett
- Gale-Stewart
- Wadge

3 Infinite Word Models



3.1 Class Inclusions

- E aut. $\subseteq$ Staiger-Wagner
  Proof: SWA with $\mathcal{F} = \{Q' \subseteq Q \mid F \cap Q' \neq \emptyset\}$.

- A. aut. $\subseteq$ Staiger-Wagner
  Proof: SW closed under complement,

- Staiger-Wagner $\subseteq$ DBA / coBA
  Proof: $\mathcal{A}$ SWA $\Rightarrow \mathcal{A}' = (Q \times 2^Q, \Sigma, (q_0, \emptyset), \delta', F')$
  Collect all visited states and accept if that set stays in $\mathcal{F}$.

- DBA $\subseteq$ NBA
  trivial

- coBA $\subseteq$ NBA
  Proof: NBA closed under complement.

- LTL $\subseteq$ NBA
  Proof: **??**

- LTL $\subseteq$ ABA
  Proof: **??**

3.2 Class Exclusions

- E aut. $\not\subseteq$ A aut.
  Example: $(a+b)^*a(a+b)^\omega$
  Proof: **??**

- A aut. $\not\subseteq$ E aut.
  Example: $\{a^\omega\}$
  Proof: **??**

- DBA $\not\subseteq$ coBA
  Example: $(a^*b)^\omega$
  Proof: **??**

- coBA $\not\subseteq$ DBA
  Example: $(a+b)^*a^\omega$
  Proof: **??**

- LTL $\not\subseteq$ NBA
  Example: $((a+b)a)^\omega$
  Proof: **??**

3.3 Class Equalities

3.3.1 NBA

- NBA $\Rightarrow$ ω-regular
  Proof: **??**

- ω-regular $\Rightarrow$ NBA
  Proof: **??**

- NBA $\Rightarrow$ $\exists$S1S
  Proof: **??**

- S1S $\Rightarrow$ S1S$_0$
  Proof: **??**

- $\mathrm{S1S}_0 \Rightarrow$ NBA
  Proof: **??**

- Det. Muller $\Rightarrow$ NBA
  Proof: NBA with $L(\mathcal{A}) = \bigcup_{F \in \mathcal{F}} \left(\bigcap_{q \in F} L(\mathcal{A}_q) \cap \bigcap_{q \notin F} \overline{L(\mathcal{A}_q)} \right)$ where $\mathcal{A}_q$ is $\mathcal{A}$ starting in q.

- NBA $\Rightarrow$ det. Muller
  Proof: **??**

- (det.) Muller $\Rightarrow$ (det.) Parity
  Proof: **??**

- ABA $\Rightarrow$ NBA
  Proof: **??**

3.3.2 LTL

LTL $\Leftrightarrow$ Non-counting
No proof. Remarks in F8.

3.3.3 SW

Staiger-Wagner $\Leftrightarrow$ det. weak BA
Proof: **??**

3.4 Closures

3.4.1 NBA

- Closed under union
  Proof: **??**

- Closed under intersection
  Proof: **??**

- Closed under complement
  Proof: **??**

3.4.2 DBA

- Not closed under complement (inf. many $a \leftrightarrow$ fin. many a)

3.4.3 SW

- Closed under union
  Proof: **??**

- Closed under intersection
  Proof: **??**

- Closed under complement
  Proof: **??**

3.5 Characterizations

- Parity conditions are directly convertible to Rabin chain conditions and vice-versa.
  Proof: Assign priorities in ascending order; $E_k \to 0$, $F_k \setminus E_k \to 1$, $E_{k-1} \setminus F_k \to 2 \ldots$
- U is ω-regular iff U is a Boolean combination of DBA-languages
  Proof: NBAs are closed under Boolean operations.
- U is DBA-recog. iff $U = \lim(L)$ for some regular $L \subseteq \Sigma^*$.
  Proof: **??**
- U is E-recog. iff $U = L \cdot \Sigma^*$ for some regular $L \subseteq \Sigma^*$.
  Proof: **??**
- Landweber's theorem
  Proof: **??**
- DBA $\cap$ coBA $\Rightarrow$ SW
  Proof: **??**

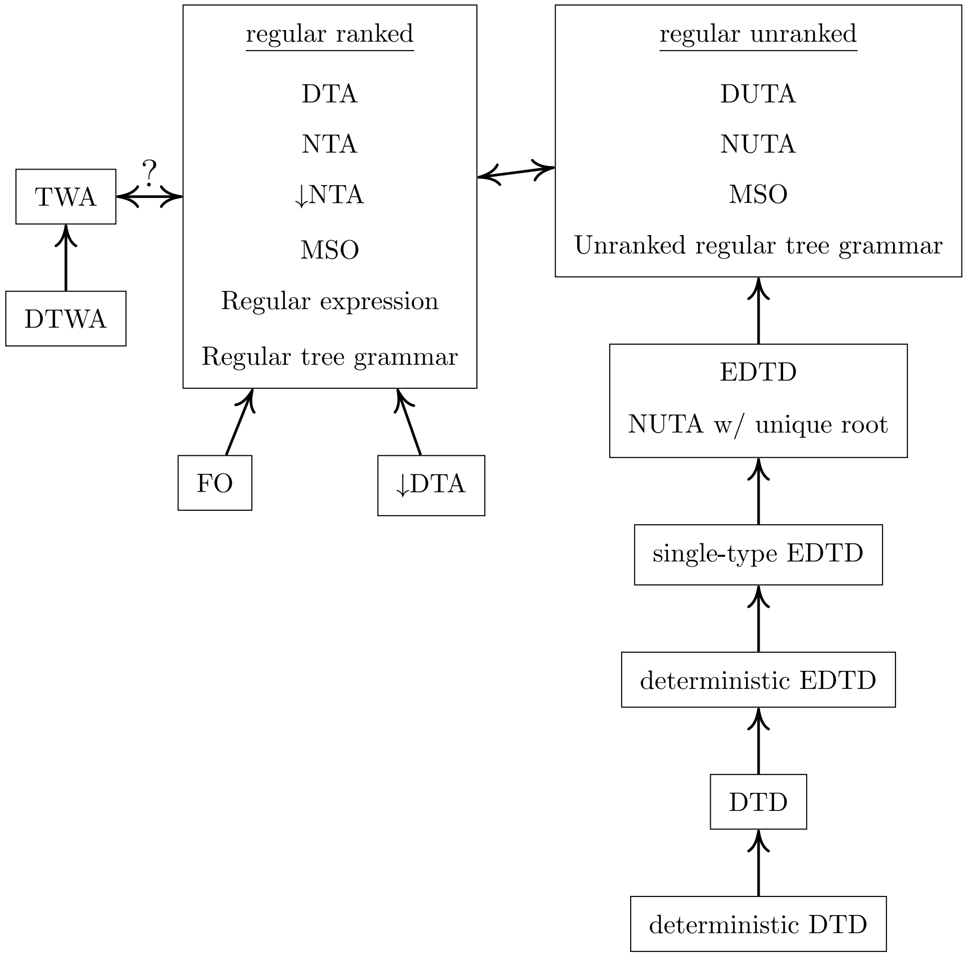
3.6 Duality

- U is A-recog. iff $\Sigma^\omega \setminus U$ is E-recog.
  Proof: **??**
- U is coBA-recog. iff $\Sigma^\omega \setminus U$ is DBA-recog.
  Proof: **??**

3.7 Problems / Complexity

- Emptiness problem for NBAs is decidable in poly. time.
  Proof: **??**
- Emptiness problem for ABAs is PSPACE-complete.
  No proof. Remarks in F23.
- Membership problem for ABAs is decidable in poly. time.

4 Finite Tree Models



4.1 Class Inclusions

- Regular Ranked ⊆ Regular Unranked
  Proof:
- det. DTD ⊆ DTD ⊆ det. EDTD ⊆ single-type EDTD ⊆ EDTD
  trivial
- EDTD ⊆ Regular tree grammar
  Proof: $N = \Sigma'$, $P_{\text{gram}} = P \cup \{a^{(n)} \to a \mid a, n\}$

4.2 Class Exclusions

- ↓DTA $\not\subseteq$ NTA
  Example: $T = \{f(a, b), f(b, a)\}$

Proof: **??**

- DTD $\not\subseteq$ single-type EDTD
Example: $T = \{t \in T_{\{a,b\}} \mid$ there is a path in t on which a occurs exactly twice

- NUTA w/ unique root $\not\subseteq$ NUTA
Example: $T = \{a, b\}$

- FO $\not\subseteq$ MSO
Example: T = positive boolean terms that evaluate to true

4.3 Class Equalities

4.3.1 Regular Ranked

- NTA $\Rightarrow$ DTA
Proof: Subset construction.

- NTA $\Leftrightarrow$ $\downarrow$NTA
Proof: **??**

- $\downarrow$NTA $\Leftrightarrow$ Regular Tree Grammar
Proof: **??**

- MSO $\Leftrightarrow$ NTA
Proof: **??**

4.3.2 Regular Unranked

- NUTA $\Rightarrow$ DUTA
Proof:

4.3.3 EDTD

- NUTA with unique root $\Leftrightarrow$ EDTD (in poly. time)
Proof:

4.4 Closures

4.4.1 Regular Ranked

- Regular (ranked) trees are closed under complement.
Proof: **??**

- Regular (ranked) trees are closed under union.
Proof: **??**

- Regular (ranked) trees are closed under intersection.
Proof: **??**

4.4.2 Regular Unranked

- Regular unranked trees are closed under complement, union, and intersection.
  Proof: via FCNS

4.5 Problems / Complexity

4.5.1 Regular Ranked

- Membership problem for NTAs is decidable.
- Reachable states of NTAs can be computed in linear time in $|\mathcal{A}|$.
- Emptiness of an NTA can be decided in linear time in $|\mathcal{A}|$.
  Algorithm: $T(\mathcal{A}) = \emptyset$ iff Reachable$(\mathcal{A}) \cap F = \emptyset$
- Given a DTA $\mathcal{A}$, $\sim_{T(\mathcal{A})}$ can be computed in time poly$(|Q^m \times \Sigma \times Q|)$ where m is the maximal arity in Σ.
  Algorithm:

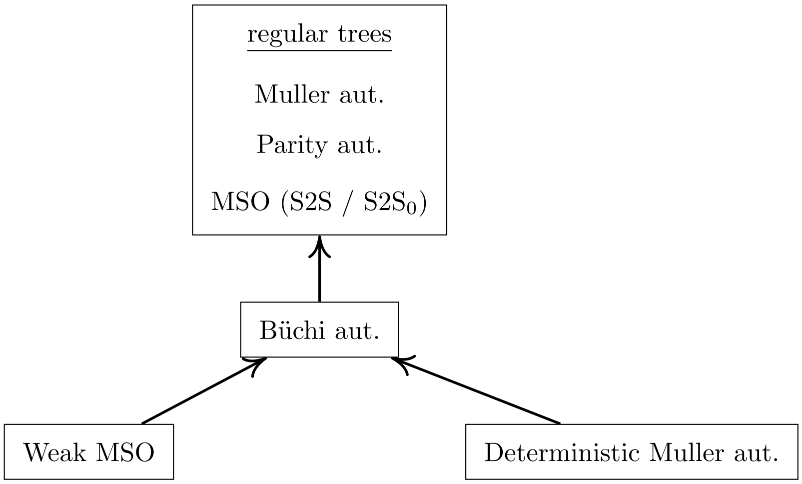
4.5.2 Regular Unranked

- Emptiness / membership / inclusion for NUTAs is decidable in polynomial / polynomial / exponential time.
  Algorithm:
- Inclusion for complete DUTAs is decidable in polynomial time.
  Algorithm:

4.5.3 Grammars

- Emptiness / membership for EDTDs is decidable in polynomial time.
  Proof: EDTD can be converted to NUTA in polynomial time.
- Inclusion for deterministic EDTDs is decidable in polynomial time.
  Proof:

5 Infinite Tree Models



5.1 Class Differences

TODO

5.2 Class Equalities

TODO

5.3 Closures

TODO