Theorem 2. For a Kripke structure K with initial state s and  $\varphi \in LTL$ , the model checking problem  $L(K,s) \subseteq L(\varphi)$ ? is PSPACE-complete.

Proof. PSPACE Compute the intersection automaton for  $L(K,s) \cap L(\neg \varphi)$  and test it for emptiness. PSPACE-hard Encode a poly.-length Turing tape as a Kripke structure and its correct behavior in LTL.

Theorem 3 (Büchi). The MSO theory of  $(\mathbb{N}, +1, <, 0)$  is decidable.

Proof. Corresponds to S1S formula. Can be checked with NBA emptiness test.

Theorem 4. The FO theory of  $(\mathbb{R}, +, <, 0)$  is decidable.

Proof. Encode real numbers x by triples of sets  $(X_s, X_i, X_f)$  with the number's sign  $(X_s = \emptyset)$  or  $\{0\}$ , the positive decimal digits in binary encoding, and the positive fractional digits in binary

encoding. Then an FO sentence can be transformed to an equi-satisfiable MSO sentence over

 $(\mathbb{N}, +1, <, 0).$ 

**Theorem 1.** Every non-empty  $\omega$ -regular language contains an ultimately periodic word.