

# Implementing Tuple Variables in Gecode

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#### Gecode

- Library for constraint programming
- Open source MIT licence
- Implemented in C++



- Declarative style programming
- The programmer states what must hold, without stating how it is achieved
- A constraint programming variable has more in common with mathematical variables than regular variables used in programming



while not done:

propagators prune as much as possible a brancher assigns a value to a variable



Sudoku is a good example of a problem suitable for constraint programming.





# What data types exist in Gecode?

- Boolean
- Integer
- Integer set
- Float



## What about Tuple Variables?

- A tuple is an ordered list of elements.
   The elements can be of any type, and do not need to be of the same type.
- We restrict ourselves to 2dimensional integer tuple variables.



## What about Tuple Variables?

#### Consider the following example:

```
Intvar a,b = {1, ..., 1000}
rel(a ≠ b)
```



## What about Tuple Variables?

- Without tuples variables, it is impossible to prune a combination of values
- Tuples variables move work from searching to propagation



#### Pair Variables

- We restrict ourselves to 2-tuples and call them pairs
- Two variants are created
  - Exact domain representation
  - Approximate representation



#### **Exact Pair**

- Stores a long list of ALL combinations
- Requires more memory
- Some operations are slower
- Can be pruned more efficiently

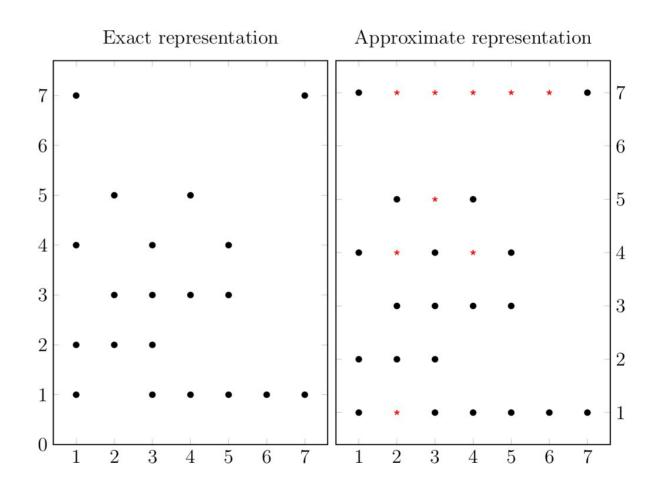


## Approximate Pair

- Does not store all possible combinations
- Requires less memory
- Some operations are faster
- Cannot be pruned as much as the exact representation



## Comparison





#### The cDFA constraint

- A cDFA is a DFA with individual costs for each transition
- A cDFA does not only tell if a string is accepted or not, but also the cost

#### The cDFA constraint



```
cDFA(ps, pc, qs, qc, x, S, C)
- qs: State before the transition
- qc: Cost before the transition
- ps: State after the transition
- qs: Cost after the transition
- x: Next symbol in the string
- S: State function
- C: Cost funtion
  ps=S(qs, x)
  pc=qc+C(qs, x)
```

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#### The cDFA constraint

cDFA(P, Q, x, S, C)

- Q: State and cost before the transition
- P: State and cost after the transition
- x: Next symbol in the string
- S: State function
- C: Cost funtion

$$P=$$



## Comparing performance

```
PairVar P[n+1]
IntVar x[n]
for i=1 to n:
    cDFA(P[n+1], P[n], x[n], S, C)
```

```
IntVar s[n+1], c[n+1], x[n+1]
for i=1 to n:
  cDFA(s[n+1], c[n+1], s[n], c[n],
     x[n], S, C)
```



## Comparing performance

Performance is measured in two ways:

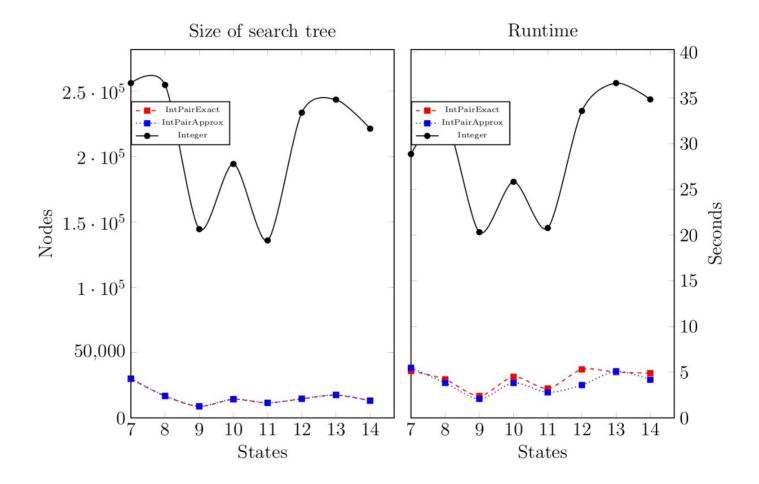
- Size of search tree
- Total runtime

Performance is compared by varying four parameters:

- Number of states
- Size of alphabet
- Cost per transition
- Length of string

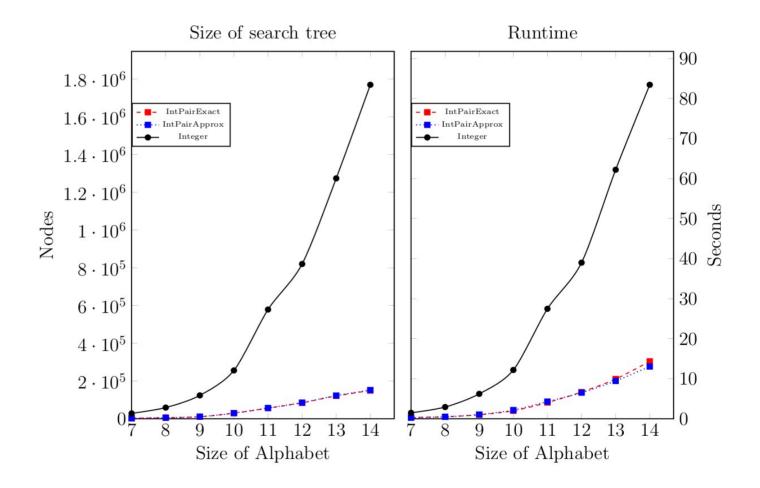
## Varying number of states





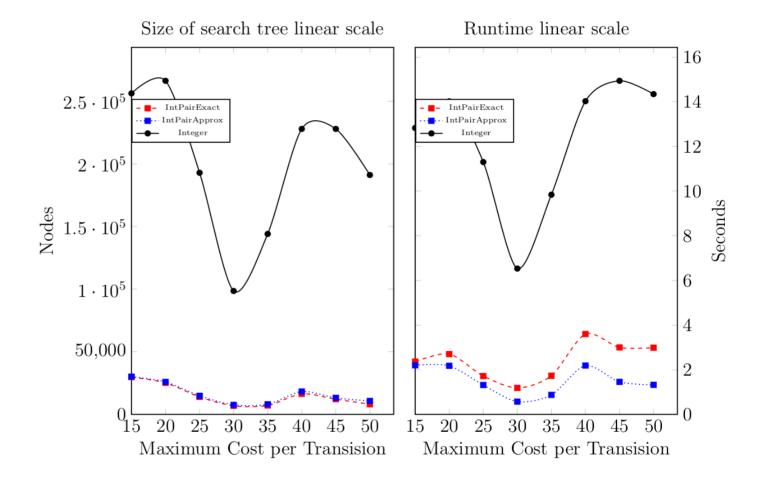
## Varying size of alphabet





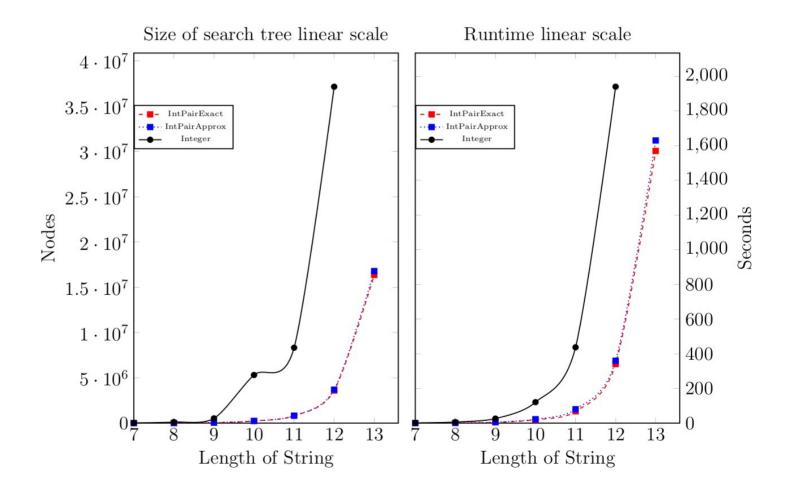
#### Varying cost per transition





## Varying length of string







#### Possible Problems

The propagator for regular integer variables may be optimised further. However, much more work have been spent on optimising this.



#### Conclusions

Tuple variables seem very promising. In these tests, they perform far better than regular integer variables. This is true for both the size of the search tree and the total execution time.