CSC3150 Assignment 4 Report

The bonus is partly implemented.

Part 0: Student Information

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Part 1: Assignment Overview

The assignment is mainly about simulating the file management system in the operating system. The source part simulates only the root directory, while the bonus part considers tree-structured directories. For the source part, **opening files, reading files and listing some file information** in the system are the main functions. The implementation mainly includes **compacting the files and contiguous allocation**. For the bonus part, the key point is still how to maintain a **compact structure of the file system**. Details are shown in the following part.

Part 2: Implementation Details

2.0 Design of System

For the volume of the file system, the first 4KB is the **super block (Volume Control Block VCB)**, in which every byte (8 bits) represents whether the 8 contiguous blocks have been allocated. If allocated, the corresponding bit is 1, otherwise 0. So this area is always interpreted with **binary representation**. The following 32KB is the **file control block (FCB)**. **Each FCB is allocated with 32B, calculated by 32KB/1024(maximum number of files)**. For every FCB, the structure to store the **meta information** of the corresponding file is shown as follows:

File name (20B)	Start block index (2B)	Length of file (2B)	Modified time (2B)	Created time (2B)

While in the bonus part, FCB should store more meta information due to the tree-structured directory, where the **identity of this file** is **a directory (1) or a file (0)**, represented in the **29th byte (though it wastes the other 31 bits)**. The parent directory index of the **root directory is 0xffff**. The depth of the file or directory is the distance between itself and the root directory + 1. For example, the **root directory** has **depth 1**.

File name (20B) Start block index (2B) Length of file (2B) Modified time (2B) Created time (2B)

2.1 fs_open

fs_open function returns a **pointer to the physical address** of the wanted file. Naturally, first consider whether this file is in the FCB. So try to search for its name in FCB. If the file already exists in FCB, then return its starting address with the help of metadata stored in FCB.

```
if ((!found) && (op == G_WRITE)){
  // printf("[fs_open] : G_WRITE new file: %s.\n", s);
/* Handle the G_WRITE new file. */
  /* traverse the VCB to record the first available block */
  int first avai block = 0xfffff;
  for (int i = 0; i < fs->SUPERBLOCK_SIZE; i++){
    if (first_avai_block != 0xffff) break;
    if (fs->volume[i] != 0xff){
      for (int j = 7; j >= 0; j--){
        if (((fs-)volume[i] & (1 << j)) >> j) == 0) {
          first_avai_block = i*8 + (7-j);
          // write to VCB to allocate memory.
          fs \rightarrow volume[i] += (1 \leftrightarrow j);
          // printf("[fs_open] : find available block %u\n", first_avai_block);
          // printf("[fs_open] : update VCB %d to %x\n", i, fs->volume[i]);
          break;
  if (first_avai_block == 0xffff) {
    printf("[ERROR fs_open] : no available block.\n");
    return 0XFFFFFFF;
```

```
/* traverse the FCB to find first available FCB*/
int first_avai_FCB = 0xfff;
for (int i = 0; i < fs->FCB_ENTRIES; i++){
    if (first_avai_FCB == 0xfff){
        for (int j = 0; j < fs->MAX_FILENAME_SIZE; j++){ // browse the file name of this FCB to judge whether available.
        if (fs->volume[fs->SUPERBLOCK_SIZE + i*(fs->FCB_SIZE) + j] != 0) break;
        if (j == fs->MAX_FILENAME_SIZE - 1) {
            first_avai_FCB = i;
            // printf("[fs_open] : find available FCB %x\n", first_avai_FCB);
            break;
        }
    }
    else break;
}

if (first_avai_FCB == 0xfff) {
    printf("[ERROR fs_open] : no available FCB.\n");
    return 0XFFFFFFFF;
}
```

If not found, then consider the **G_WRITE** mode. Now the program tries to find the first available block and available FCB to store the newly created file. To find an available place, just search the FCB and VCB from head to tail. During the process of searching for VCB, also update the first available VCB to allocate a new block for the newly created file.

After finding the available block and FCB, store the metadata of the file in FCB. Since it's a vacant file, there is no need to update the file content.

By now, <code>fs_open</code> in the source part ends. However, in the <code>bonus</code> part, not only considering whether the file name matches, whether the filename is in the same directory and whether the file number limit is reached also matter. Only when the matched name file is in the current directory is the file considered found. What's more, the meta data of this newly created file is more than the source part. But the extras in the bonus part mentioned above are trivial. <code>More importantly</code>, after opening this newly created file, the <code>parent directory data should be updated</code>. Since this is a new file, the <code>file content</code> of the directory is added with the name of the new file. And the metadata 'length' and 'modified time' should also be updated. <code>But take care, every time the length changes, consider whether the block size of the file changes. If the block size changes, compaction is needed both in the bonus part and source part. And the flow of compaction is the same.</code>

```
fs->volume[fs->SUPERBLOCK SIZE + first avai FCB*(fs->FCB SIZE) + 20] = (first avai block >> 8);
 fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 21] = (first_avai_block & 0xff);
fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 24] = (gtime >> 8);
fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 25] = (gtime & 0xff);
 fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 26] = (gtime >> 8);
fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 27] = (gtime & 0xff);
fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 28] = current_depth + 1;
 fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 38]
fs->volume[fs->SUPERBLOCK_SIZE + first avai FCB*(fs->FCB_SIZE) + 31] = (current dir index & 0xff):
fp = fs->FILE_BASE_ADDRESS + ((fs->STORAGE_OUCK_SIZE)*(((fs->volume[fs->SUPERBLOCK_SIZE + first_avai_FCB*(fs->FCB_SIZE) + 20]) << 8) + fs->volume[fs->SUPERBLOCK_SIZE + first_
int parent_dir_length = (fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 22] << 8) + fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 23] int parent_dir_length == 0) ? 1 : 1 + (parent_dir_length - 1) / fs->STORAGE_BLOCK_SIZE;
int parent_dir_start_block = (fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 20] << 8) + fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE)
int new_parent_dir_length = parent_dir_length + name_length;
int new_parent_dir_block_size = (new_parent_dir_length == 0) ? 1 : 1 + (new_parent_dir_length - 1) / fs->STORAGE_BLOCK_SIZE;
int parent_parent_index = (fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 30] << 8) + fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 30] = (8) + fs->volume[fs->FCB_SIZE + current_dir_index*(fs->FCB_SIZE) + 30] = (8) + fs->volume[fs->FCB_SIZE + current_dir_index*(fs->FCB_SIZE) + 30] = (8) + fs->volume[fs->FCB_SIZE + current
char new dir contents[1024];
 char parent_dir_name[20];
char parent dir invariant[4];
if (parent_dir_block_size == new_parent_dir_block_size) 
     fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 22] = (new_parent_dir_length >> 8);
     fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 23] = (new_parent_dir_length & 0xff);
     fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 24] = (gtime >> 8);
fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + 25] = (gtime & 0xff);
      for (int i = 0; i < name_length; i++){</pre>
           fs->rolume[fs->FILE_BASE_ADDRESS + parent_dir_start_block*(fs->STORAGE_BLOCK_SIZE) + parent_dir_length + i] = s[i];
      for (int i = 0; i < parent dir_length; i++){
          new_dir_contents[i] = fs->volume[fs->FILE_BASE_ADDRESS + parent_dir_start_block*(fs->STORAGE_BLOCK_SIZE) + i];
      for (int i = 0; i < name_length; i++){
           new_dir_contents[parent_dir_length + i] = s[i];
       .
for (int i = 0; i < 20; i++){
          parent_dir_name[i] = fs->volume[fs->SUPERBLOCK_SIZE + current_dir_index*(fs->FCB_SIZE) + i];
```

The process of handling the length change is as follows:

- 1. investigating whether the block size changes ->
- 2. if no change or this is the last file in the FCB, thank god, no need to compact ->
- 3. else remove this changed size file and do the compaction (compact the file contents, compact the VCB and update the FCB information)->
- 4. add this changed size file to the end of the block and FCB

The key point of the implementation is **how to compact**. First, compact the file contents block by block from the starting block, which is often the starting block of the removed file. Rewrite the block contents. It is intuitional that the **new contents are just from the block after an offset (equal to the removed block size).**

```
// compact the VCB (from the backward direction)
// find the last allocated block
int last_block = (1 << 15);
for (int i = 0; i < (fs->MAX_FILE_SIZE / fs->STORAGE_BLOCK_SIZE); i++){
    int VCB_index = i / 8;
    int VCB_offset = i % 8;
    if ((fs->volume[VCB_index] & (1 << (7-VCB_offset))) == 0) {
        last_block = i - 1;
        // printf("write compact: 最后一个block index为 %d\n", last_block);
        break;
    }
}

for (int i = 0; i < parent_dir_block_size; i++){
    int VCB_index = (last_block - i) / 8;
    int VCB_offset = (last_block - i) % 8;
    fs->volume[VCB_index] &= (0xff - (1 << (7 - VCB_offset)));
    // printf("压VCB: VCB[0], VCB[1]: %x %x\n", fs->volume[0], fs->volume[1]);
}
```

Secondly, compact the VCB. The VCB is just all 1 in the previous state, so find the last VCB, which in the byte representation is not 0xff. Then, from the last VCB to compact backward, the **last few** bits (the number of bits equal to the removed block size) should be updated from 1 to 0.

```
// Consect to ACID
for (Int 1 = 0); for 5-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
for (Int 1 = 0); for for 5-0); for for 5-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-000_BITTERS; i+>)
if (correct dir_index + 1 + 1 > 6-0-0
```

```
// update first half parents
for (int i = 0; i < current_dir_index; i++){
    int this_parent_index = (fs->volume[fs->SUPERBLOCK_SIZE + i*(fs->FCB_SIZE) + 30] << 8) + fs->volume[fs->SUPERBLOCK_SIZE + i*(fs->FCB_SIZE) + 31];
    if (this_parent_index == 0xfffff){
        fs->volume[fs->SUPERBLOCK_SIZE + (current_dir_index + i)*(fs->FCB_SIZE) + 30] = (0xff);
        fs->volume[fs->SUPERBLOCK_SIZE + (current_dir_index + i)*(fs->FCB_SIZE) + 31] = (0xff);
    }
    else if (this_parent_index == current_dir_index){
        fs->volume[fs->SUPERBLOCK_SIZE + i*(fs->FCB_SIZE) + 30] = (new_avai_FCB >> 8);
        fs->volume[fs->SUPERBLOCK_SIZE + i*(fs->FCB_SIZE) + 31] = (new_avai_FCB & 0xff);
    }
    else if (this_parent_index > current_dir_index){
        fs->volume[fs->SUPERBLOCK_SIZE + i*(fs->FCB_SIZE) + 30] = ((this_parent_index-1) >> 8);
        fs->volume[fs->SUPERBLOCK_SIZE + i*(fs->FCB_SIZE) + 31] = ((this_parent_index-1) & 0xff);
    }
}
```

Finally, update FCB. The FCB contains the information mentioned in the design part. However, one more thing needs to be consisted is that, in bonus part, the parent directory index is stored in every FCB. But think about the process of compaction, some FCB are moved forward 1 unit! And when the removed file is a directory, its index is newly allocated! So when updating the FCB, the parent directory index should be handled carefully. And a trivial case is that the root directory's parent index should remain the same.

```
int new_VCB_index;
 int new VCB offset;
 for (int i = 0; i < new_parent_dir_block_size; i++){
   new_VCB_index = (i + new_avai_block) / 8;
   new_VCB_offset = (i + new_avai_block) % 8;
   fs->volume[new_VCB_index] |= (1 << (7 - new_VCB_offset));</pre>
   // printf("写VCB: VCB[0], VCB[1]: %x %x\n", fs->volume[0], fs->volume[1]);
 for (int i = 0; i < 20; i++) {
   fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + i] = parent_dir_name[i];
   if (parent_dir_name[i] == '\0') {
     break;
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 20] = (new_avai_block >> 8);
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 21] = (new_avai_block & 0xff);
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 22] = (new_parent_dir_length >> 8);
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 23] = (new_parent_dir_length & 0xff);
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 24] = (gtime >> 8);
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 25] = (gtime & 0xff);
 for (int i = 0; i < 4; i++){
   fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 26 + i] = parent_dir_invariant[i];
   // update parent in FCB
 if ((parent_parent_index > current_dir_index) && (parent_parent_index != 0xffff)) parent_parent_index--;
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 30] = (parent_parent_index >> 8);
 fs->volume[fs->SUPERBLOCK_SIZE + new_avai_FCB*(fs->FCB_SIZE) + 31] = (parent_parent_index & 0xff);
 for (int i = 0; i < new_parent_dir_length; i++){
   fs->volume[fs->FILE_BASE_ADDRESS + new_avai_block*(fs->STORAGE_BLOCK_SIZE) + i] = new_dir_contents[i];
 current_dir_index = new_avai_FCB;
return fp;
```

After compaction, add the removed file to the end of other files. This is generally the whole process and the important things of compaction. The following statement will not repeat it.

2.2 fs write

Generally, the writing process has a similar flow to the opening process (fetech the file in FCB -> check if there is any block size change -> handle it). But since the opening process has created the file already, the writing process should **definitely verify the file to write, otherwise raise an error.**

2.3.1 LS_D and LS_S

The main idea is to search in FCB iteratively. In every iteration, **find the most fit file** to print out (how to define the 'most fit' depends on **file's modified time, size, created time, LS_S and LS_D**. But it's trivial here.) To filter and find the candidate to print out, **maintaining the most fit file information from the last iteration** is needed.

```
int prev max;
int now_max;
int now_max_index;
printf("[fs_gsys] : ===sort by modified time===\n");
for (int i = 0; i < fs->FCB ENTRIES; i++){
 int vacant = 1;
  for (int n = 0; n < fs->MAX_FILENAME_SIZE; n++){
    if (fs\rightarrow volume[fs\rightarrow SUPERBLOCK\_SIZE + i*(fs\rightarrow FCB\_SIZE) + n] != 0){
      vacant = 0;
      break:
  if (vacant) return;
  int complete = 1;
  if (i == 0) prev_max = 2147483647;
  for (int j = 0; j < fs->FCB_ENTRIES; j++){
    int vacant = 1;
    for (int n = 0; n < fs->MAX_FILENAME_SIZE; n++){
      if (fs\rightarrow volume[fs\rightarrow SUPERBLOCK\_SIZE + j*(fs\rightarrow FCB\_SIZE) + n] != 0){}
        vacant = 0:
        break;
    if (vacant) break;
    int time = (fs->volume[fs->SUPERBLOCK SIZE + j*(fs->FCB SIZE) + 24] << 8) + fs->volume[fs->SUPERBLOCK SIZE + j*(fs->FCB SIZE) + 25];
    if (time < prev_max){</pre>
      now_max = time;
      now_max_index = j;
      complete = 0;
  if (complete) return;
```

```
for (int j = 0; j < fs->FCB_ENTRIES; j++){
  int vacant = 1;
  for (int n = 0; n < fs->MAX_FILENAME_SIZE; n++){
    if (fs\rightarrow volume[fs\rightarrow SUPERBLOCK\_SIZE + j*(fs\rightarrow FCB\_SIZE) + n] != 0){}
      vacant = 0;
  if (vacant) break;
  int time = (fs->volume[fs->SUPERBLOCK_SIZE + j*(fs->FCB_SIZE) + 24] << 8) + fs->volume[fs->SUPERBLOCK_SIZE + j*(fs->FCB_SIZE) + 25];
  if ((time < prev_max) && (time > now_max)){
   now_max = time;
   now_max_index = j;
char name[20];
for (int j = 0; j < fs->MAX_FILENAME_SIZE; j++){
 name[j] = fs->volume[fs->SUPERBLOCK_SIZE + now_max_index*(fs->FCB_SIZE) + j];
  if (name[j] == '\0') break;
printf("[fs_gsys] : %s\n", name);
prev_max = now_max;
```

For example, in LS_D, the most fit file is the **most recently modifed file**, which means **modified time is largest**. So look at the modified time in FCB. In the first iteration, set **an upper bound of 0xffff** to ensure **every FCB is qualified** to print out. **Record the largest modified time as the second iteration upperbound**. In the second iteration, only the files with modified time less than the upperbound are qualified to print out. Iteration by iteration, all files can be printed out.

In the bonus part, the procedure is the same, but there is one more restriction: only the files in the current directory can be printed out. Sorry that when implementing this part of bonus, it's near 12pm on November 30th, so I have no time to fix it.

2.3.2 RM

The removal part is simple if the implementation of compaction is successful. The only difference is that there is no need to add the file back after compaction. Other things are the same.

2.3.3 RM RF

I have no time to implement this part of bonus. I'm so sorry about it. My mind goes like this. Since this bonus requires the depth of directory to be no more than 3, there is no need to use recursive methods to solve it. When the removed directory depth is 1, it's not possible to remove the root directory, so raise an error. When the depth is 2, first check all its children by searching the whole FCB. All children files should be removed. If there are any children directories, the children of them should only be files, so remove them directly using RM. Then remove the children directories, lastly remove itself. When the depth is 3, remove its children files first then itself.

2.3.4 CD and CD_P

CD and CD_P are both simple since the program maintains a global variable to record the current directory. Only updating that variable is fine.

2.3.5 MKDIR

MKDIR is the directory version of fs_open . Just need one more step to check whether it is valid. So no more words here.

Part 3: How to Run My Program

1. Make sure you are in the directory: Assignment_4_120090575\source or Assignment_4_120090575\bonus, type sbatch ./slurm.sh in the terminal.

2. You will get a result.out file to see the fs_gsys output and a snapshot.bin file as the dump data. snapshot.bin should be the same as data.bin in test case 4.

Part 4: Environment

CUDA version:

nvcc: NVIDIA (R) Cuda compiler driver

Copyright (c) 2005-2022 NVIDIA Corporation

Built on Wed_Jun__8_16:49:14_PDT_2022

Cuda compilation tools, release 11.7, V11.7.99

Build cuda_11.7.r11.7/compiler.31442593_0

- GPU: Quadro RTX 4000
- VS Code version: version 1.70
- OS version:
- CentOS Linux release 7.5.1804 (Core)
- kernel version:

Linux version 3.10.0-862.el7.x86_64 (builder@kbuilder.dev.centos.org) (gcc version 4.8.520150623 (Red Hat 4.8.5-28) (GCC)) #1 SMP Fri Apr 20 16:44:24 UTC 2018

Part 5: What I've Learnt

- 1. The file system management is well learnt and familiar.
- 2. Some allocation, like contiguous allocation and linked allocation are more familiar, expecially contiguous allocation.
- 3. The debug skill is enhanced.

Part 6: Screenshot

Test case 3:

```
[fs gsys] : ===sort by modified time===
41
42
     [fs gsys] : t.txt
43
     [fs gsys] : b.txt
44
     [fs gsys] : ===sort by file size===
45
     [fs gsys] : t.txt 32
     [fs gsys] : b.txt 32
     [fs_gsys] : ===sort by file size===
47
     [fs_gsys] : t.txt 32
     [fs gsys] : b.txt 12
     [fs_gsys] : ===sort by modified time===
50
51
     [fs gsys] : b.txt
52
     [fs gsys] : t.txt
     [fs gsys] : ===sort by file size===
     [fs gsys] : b.txt 12
54
     [fs gsys] : ===sort by file size===
56
     [fs_gsys] : *ABCDEFGHIJKLMNOPQR 33
57
     [fs gsys]: )ABCDEFGHIJKLMNOPQR 32
     [fs_gsys] : (ABCDEFGHIJKLMNOPQR 31
58
     [fs gsys] : 'ABCDEFGHIJKLMNOPQR 30
     [fs gsys]: &ABCDEFGHIJKLMNOPQR 29
61
     [fs_gsys] : %ABCDEFGHIJKLMNOPQR 28
62
     [fs_gsys] : $ABCDEFGHIJKLMNOPQR 27
     [fs_gsys] : #ABCDEFGHIJKLMNOPQR 26
     [fs_gsys] : "ABCDEFGHIJKLMNOPQR 25
64
     [fs_gsys] : !ABCDEFGHIJKLMNOPQR 24
65
```

```
[fs gsys] : !ABCDEFGHIJKLMNOPQR 24
65
     [fs_gsys] : b.txt 12
67
     [fs gsys] : ===sort by modified time===
    [fs gsys] : *ABCDEFGHIJKLMNOPOR
68
     [fs_gsys] : )ABCDEFGHIJKLMNOPQR
     [fs_gsys] : (ABCDEFGHIJKLMNOPQR
70
     [fs_gsys] : 'ABCDEFGHIJKLMNOPQR
71
72
     [fs gsys] : &ABCDEFGHIJKLMNOPQR
    [fs gsys] : b.txt
74
     [fs gsys] : ===sort by file size===
    [fs gsys] : ~ABCDEFGHIJKLM 1024
75
    [fs_gsys] : }ABCDEFGHIJKLM 1023
76
    [fs gsys] : ABCDEFGHIJKLM 1022
77
78
     [fs_gsys] : {ABCDEFGHIJKLM 1021
79
     [fs gsys] : zABCDEFGHIJKLM 1020
    [fs gsys] : yABCDEFGHIJKLM 1019
     [fs gsys] : xABCDEFGHIJKLM 1018
81
     [fs_gsys] : wABCDEFGHIJKLM 1017
82
     [fs_gsys] : vABCDEFGHIJKLM 1016
83
     [fs_gsys] : uABCDEFGHIJKLM 1015
84
    [fs_gsys] : tABCDEFGHIJKLM 1014
85
86
     [fs gsys] : sABCDEFGHIJKLM 1013
87
    [fs_gsys] : rABCDEFGHIJKLM 1012
     [fs gsys] : qABCDEFGHIJKLM 1011
88
     [fs_gsys] : pABCDEFGHIJKLM 1010
89
     [fs gsys] : oABCDEFGHIJKLM 1009
```

```
Assignment_4_120090575 > source > 1 result.out
      [fs gsys] : ?A 37
1062
      [fs_gsys] : >A 36
1063
1064
      [fs gsys] : =A 35
1065
      [fs_gsys] : <A 34
      [fs_gsys] : *ABCDEFGHIJKLMNOPQR 33
1066
      [fs gsys] : ;A 33
1067
      [fs_gsys] : )ABCDEFGHIJKLMNOPQR 32
1068
      [fs_gsys] : :A 32
1069
      [fs_gsys] : (ABCDEFGHIJKLMNOPQR 31
1070
      [fs_gsys] : 9A 31
1071
      [fs_gsys] : 'ABCDEFGHIJKLMNOPQR 30
1072
1073
      [fs gsys] : 8A 30
1074
      [fs_gsys] : &ABCDEFGHIJKLMNOPQR 29
1075
      [fs_gsys] : 7A 29
      [fs_gsys] : 6A 28
1076
1077
      [fs_gsys] : 5A 27
      [fs_gsys] : 4A 26
1078
1079
      [fs_gsys] : 3A 25
1080
      [fs_gsys] : 2A 24
1081
      [fs_gsys] : b.txt 12
1082
      [fs_gsys] : ===sort by file size===
      [fs_gsys] : EA 1024
1083
      [fs_gsys] : ~ABCDEFGHIJKLM 1024
1084
      [fs_gsys] : aa 1024
1085
      [fs_gsys] : bb 1024
1086
1087
      [fs_gsys] : cc 1024
```

```
Assignment 4 120090575 > source > 1 result.out
       [fs gsys] : ===sort by file size===
1082
1083
      [fs gsys] : EA 1024
1084
      [fs gsvs] : ~ABCDEFGHIJKLM 1024
1085
       [fs_gsys] : aa 1024
1086
       [fs gsys] : bb 1024
1087
      [fs gsys] : cc 1024
1088
      [fs_gsys] : dd 1024
1089
      [fs gsys] : ee 1024
1090
      [fs_gsys] : ff 1024
      [fs gsys] : gg 1024
1091
      [fs_gsys] : hh 1024
1092
      [fs gsys] : ii 1024
1093
1094
      [fs_gsys] : jj 1024
1095
      [fs_gsys] : kk 1024
1096
      [fs gsys] : 11 1024
1097
      [fs_gsys] : mm 1024
      [fs gsys] : nn 1024
1098
1099
       [fs_gsys] : oo 1024
1100
      [fs gsys] : pp 1024
1101
      [fs_gsys] : qq 1024
      [fs gsys] : }ABCDEFGHIJKLM 1023
1102
      [fs_gsys] : ABCDEFGHIJKLM 1022
1103
1104
       [fs_gsys] : {ABCDEFGHIJKLM 1021
       [fs gsys] : zABCDEFGHIJKLM 1020
1105
       [fs_gsys] : yABCDEFGHIJKLM 1019
1106
       [fs_gsys] : xABCDEFGHIJKLM 1018
1107
```

```
Assignment 4 120090575 > source > 1 result.out
      [fs_gsys] : CA 41
2083
2084
      [fs gsys] : BA 40
2085
      [fs gsvs] : AA 39
2086
      [fs_gsys] : @A 38
      [fs_gsys] : ?A 37
2087
2088
      [fs_gsys] : >A 36
      [fs_gsys] : =A 35
2089
2090
      [fs gsys] : <A 34
2091
      [fs_gsys] : *ABCDEFGHIJKLMNOPQR 33
2092
      [fs gsys] : ;A 33
      [fs gsys] : )ABCDEFGHIJKLMNOPQR 32
2093
      [fs gsys] : :A 32
2094
      [fs_gsys] : (ABCDEFGHIJKLMNOPQR 31
2095
2096
      [fs gsys] : 9A 31
2097
      [fs_gsys] : 'ABCDEFGHIJKLMNOPQR 30
2098
      [fs gsys] : 8A 30
      [fs_gsys] : &ABCDEFGHIJKLMNOPQR 29
2099
2100
      [fs_gsys] : 7A 29
2101
      [fs_gsys] : 6A 28
      [fs_gsys] : 5A 27
2102
2103
      [fs_gsys] : 4A 26
2104
      [fs_gsys] : 3A 25
      [fs_gsys] : 2A 24
2105
2106 [fs_gsys] : b.txt 12
```

Test case 4:

```
25
     triggering gc
     [fs gsys] : ===sort by modified time===
     [fs gsys] : 1024-block-1023
27
     [fs gsys] : 1024-block-1022
28
     [fs_gsys] : 1024-block-1021
29
30
     [fs gsys] : 1024-block-1020
     [fs gsys] : 1024-block-1019
31
     [fs gsys] : 1024-block-1018
32
     [fs gsys] : 1024-block-1017
     [fs_gsys] : 1024-block-1016
     [fs gsys] : 1024-block-1015
35
     [fs_gsys] : 1024-block-1014
36
37
     [fs gsys] : 1024-block-1013
     [fs gsys] : 1024-block-1012
     [fs_gsys] : 1024-block-1011
39
     [fs gsys] : 1024-block-1010
     [fs_gsys] : 1024-block-1009
41
     [fs_gsys] : 1024-block-1008
42
     [fs_gsys] : 1024-block-1007
43
44
     [fs gsys] : 1024-block-1006
     [fs_gsys] : 1024-block-1005
45
     [fs gsys] : 1024-block-1004
47
     [fs_gsys] : 1024-block-1003
     [fs_gsys] : 1024-block-1002
     [fs gsys] : 1024-block-1001
49
     [fs_gsys] : 1024-block-1000
50
```

```
Assignment_4_120090575 > source > 📘 result.out
1027
       [fs gsys] : 1024-block-0023
       [fs_gsys] : 1024-block-0022
1028
       [fs gsys] : 1024-block-0021
1029
1030
       [fs gsys] : 1024-block-0020
       [fs_gsys] : 1024-block-0019
1031
1032
       [fs_gsys] : 1024-block-0018
       [fs gsys] : 1024-block-0017
       [fs_gsys] : 1024-block-0016
1034
1035
       [fs gsys] : 1024-block-0015
1036
       [fs_gsys] : 1024-block-0014
       [fs gsys] : 1024-block-0013
1037
       [fs_gsys] : 1024-block-0012
1038
       [fs_gsys] : 1024-block-0011
1039
       [fs_gsys] : 1024-block-0010
1040
1041
       [fs_gsys] : 1024-block-0009
1042
       [fs_gsys] : 1024-block-0008
       [fs_gsys] : 1024-block-0007
1043
       [fs_gsys] : 1024-block-0006
1044
1045
       [fs_gsys] : 1024-block-0005
       [fs gsys] : 1024-block-0004
1046
       [fs_gsys] : 1024-block-0003
1047
       [fs_gsys] : 1024-block-0002
1048
       [fs gsys] : 1024-block-0001
1049
       [fs_gsys] : 1024-block-0000
1050
1051
```