

# Pre Quiz

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**1**

(a), and (b) if "is friends" a reflexive relation.

**2**

Cannot conclude

**3**

$$(\exists f \in F, f(CS4269)) \wedge (TY \in F)$$

a)  $TY(CS4269)$  cannot conclude

b)  $\forall f \in F, f(CS4269)$  cannot conclude

**4**

$$\forall n \in N, \exists r \in R, r = \sqrt{n}$$

$5.5 \notin N$ , cannot conclude  $\exists r \in R, r = \sqrt{5.5}$

$5 \in N \implies \exists r \in R, r = \sqrt{5}$

**5**

A set is a collection of unique elements. like  $\mathbb{N}$ . The cardinality of a set describes its size. like  $|\mathbb{N}| = \aleph_0$

$$\exists \text{ injection } f : A \mapsto B, |A| \leq |B|$$

$$\exists \text{ surjection } f : A \mapsto B, |A| \geq |B|$$

**6**

A finite set is a set that contains, well finite amount of element. For example  $\emptyset$

A countable set is either finite or has the same cardinality as  $\mathbb{N}$ . for example  $\mathbb{Z}^+$

An uncountable set is a set that has the same cardinality as  $\mathbb{R}$ . for example  $\mathcal{P}(\mathbb{N})$

**7**

A function is a mapping from a set to another set.  $f : D_f \mapsto R_f$ , where  $R_f \subseteq CoDomain_f$ . Codomain is the set where the elements a function can possibly maps to, while range is the set of elements which the function actually maps to.

## 8

The power set of  $S$  is a set of all subsets of  $S$ , which has cardinality  $2^{|S|}$  For example  $f : \emptyset \mapsto \{\emptyset\}$

## 9

Base case:

$$\sum_{k=0}^0 k^2 = \frac{0(0+1)(2 \times 0 + 1)}{6} = 0$$

Inductive case: Suppose

$$\sum_{k=0}^n k^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\begin{aligned} \sum_{k=0}^{n+1} k^2 &= \frac{n(n+1)(2n+1)}{6} + (n+1)^2 \\ &= \frac{1}{6}(2n^3 + 3n^2 + n + 6n^2 + 12n + 6) \\ &= \frac{1}{6}(2n^3 + 9n^2 + 13n + 6) \\ &= \frac{1}{6}((n+1)(2n^2 + 7n + 6)) \\ &= \frac{1}{6}(n+1)(n+2)(2n+3) \end{aligned}$$

## 10

Base case:  $h = 1, |V| = 2^1 - 1 = 1$  Inductive steps: strong induction:

Suppose a CBST of height  $h$  has  $2^h - 1$  nodes, and a CBST with height  $h-1$  has  $2^{h-1} - 1$  nodes, The number of leaves  $= 2^h - 1 - (2^{h-1} - 1) = 2^{h-1}$ , to add a new level each leaf grows 2 new leaf, total  $2 \cdot 2^{h-1} = 2^h$  Thus a CBST of height  $h+1$  has  $2^h - 1 + 2^h = 2^{h+1} - 1$

## 11

A string  $w$  over an alphabet  $\Sigma$  is a finitely long sequence of characters from  $\Sigma$ . A language is all strings formed from an alphabet based on certain rules.  $\Sigma^*$  is a regular language that consists of all finitely long sequence from  $\Sigma$ .

## 12

A regular language: the empty language  $\emptyset$ . proof: trivial, a DFA that accepts nothing.

A context free language: valid parentheses, rules:

$$\begin{aligned} S &\rightarrow \epsilon \\ S &\rightarrow SS \\ S &\rightarrow (S) \end{aligned}$$

## 13

A Turing machine is a state machine defined as such:  $\{Q, \Gamma, \Sigma, \delta, b, s, f\}$  each represents: the set of states, alphabet symbols, tape symbols, transition rules  $(Q - F \times \Gamma) \mapsto (Q \times \Gamma \times \{L, R\})$ , blank symbol, starting and final states. Languages that can be either accepted or rejected by a Turing machine are called recursive. Languages that may be accepted but not guaranteed to terminate on a Turing machine are recursively enumerable.

## 14

Decidable means always terminate of a Turing machine, either accepted or rejected in finite amount of time. For example  $\emptyset$ , as regular languages are decidable.

## 15

Post's correspondence problem is undecidable.

## 16

One Turing machine terminates on a string in time polynomial wrt input size. Meaning the language is in complexity class P

## 17

NP complexity class consists of languages that can be accepted or rejected by a non-deterministic Turing machine in polynomial time wrt input size.