极客时间算法训练营 第六课 树与图

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《算法竞赛进阶指南》作者

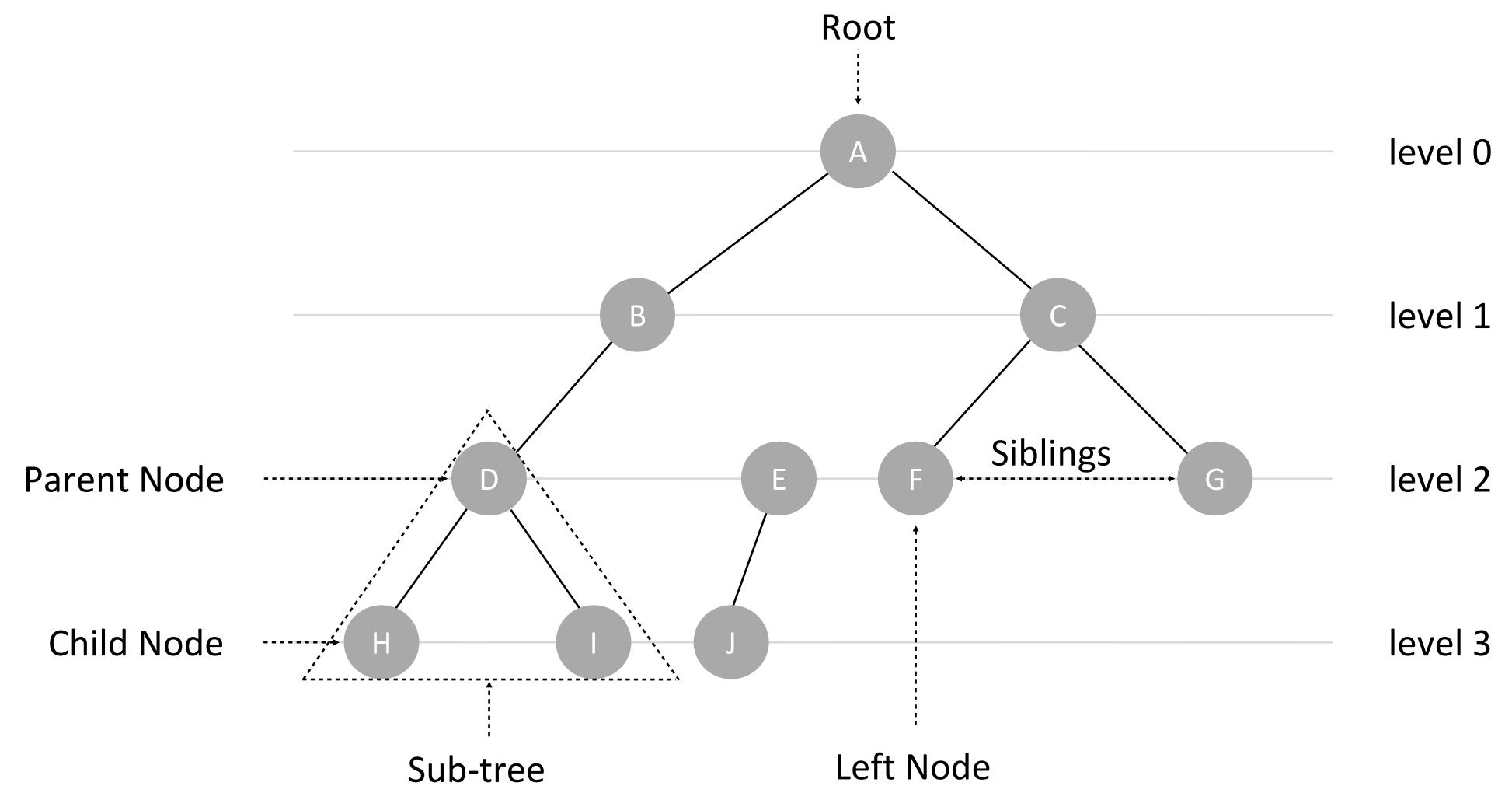


日灵

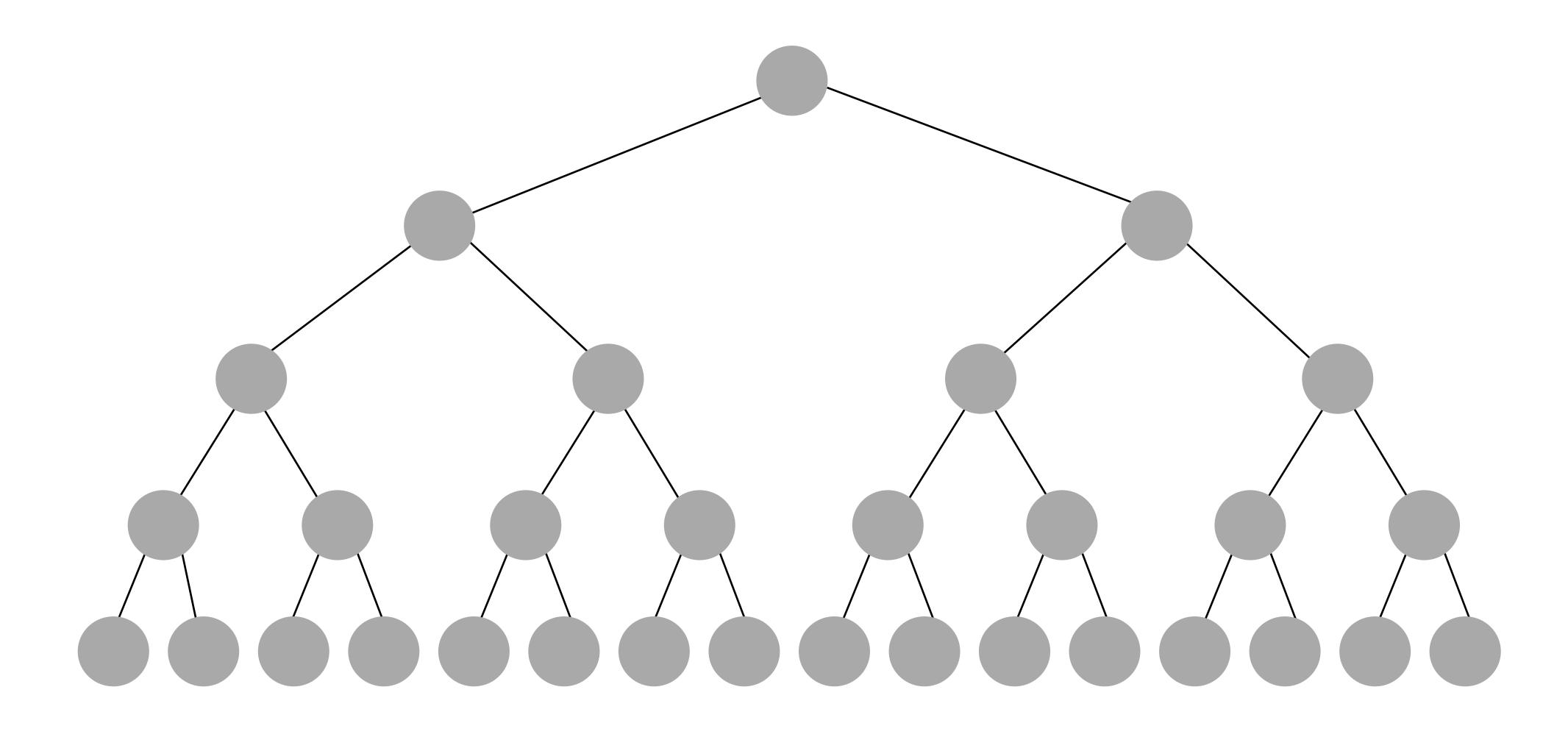
- 1. 树、二叉树、树的遍历、树的序列化
- 2. 树的直径、最近公共祖先
- 3. 树的变形(基环树)
- 4. 图、图的遍历、拓扑排序

树、二叉树、树的遍历

树 Tree



二叉树 Binary Tree



代码 - 定义树的结点

```
C++
struct TreeNode {
  int val;
  TreeNode *left;
  TreeNode *right;
  TreeNode(int x)
    : val(x), left(nullptr), right(nullptr) {}
Python
class TreeNode:
   def __init__(self, val):
     self.val = val
     self.left, self.right = None, None
```

```
public class TreeNode {
  public int val;
  public TreeNode left, right;
  public TreeNode(int val) {
     this.val = val;
     this.left = null;
     this.right = null;
  }
}
```

二叉树的遍历

前序 Pre-order: 根 - 左子树 - 右子树

A B D H I E J C F G

中序 In-order: 左子树 - 根 - 右子树

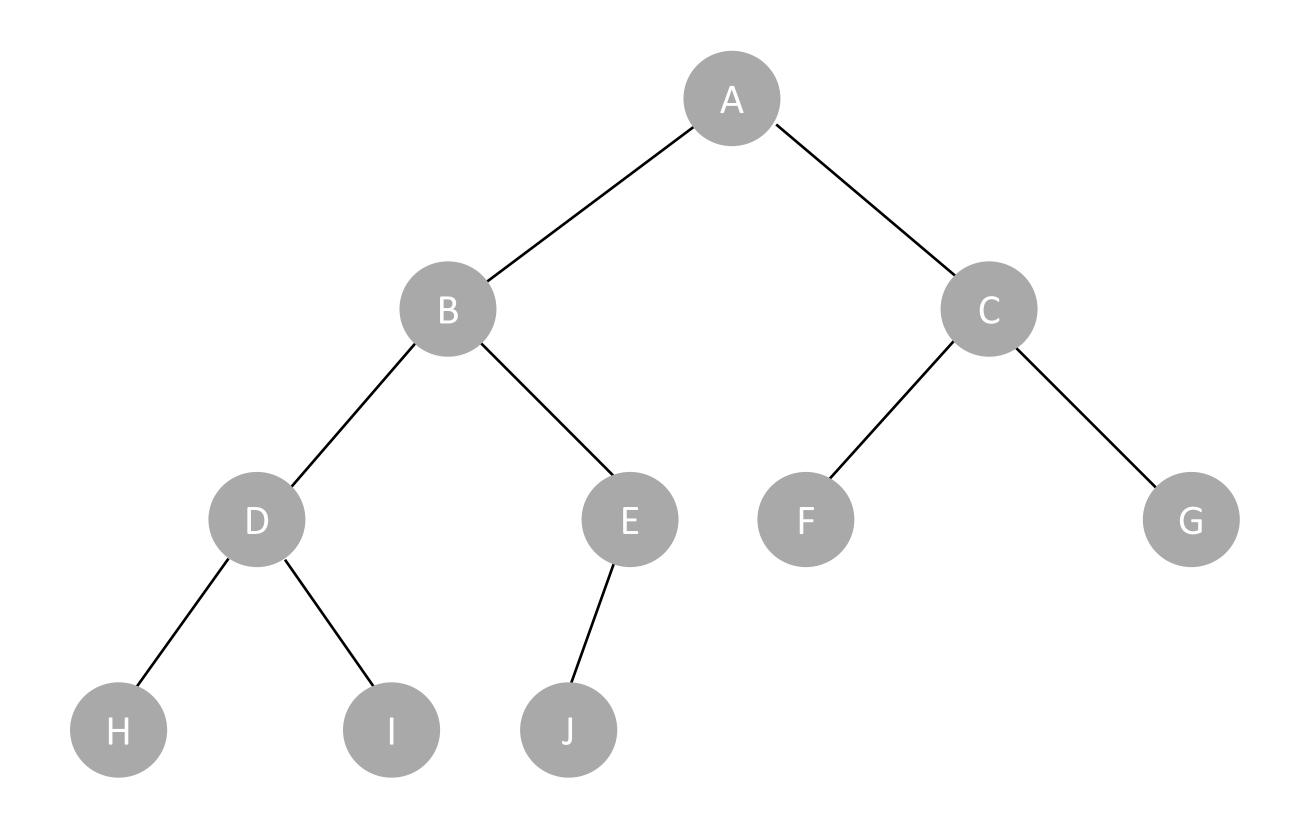
<u>HDIBJE</u>AFCG

后序 Post-order: 左子树 - 右子树 - 根

<u>HIDJE</u>BFGCA

层次序

A B C D E F G H I J

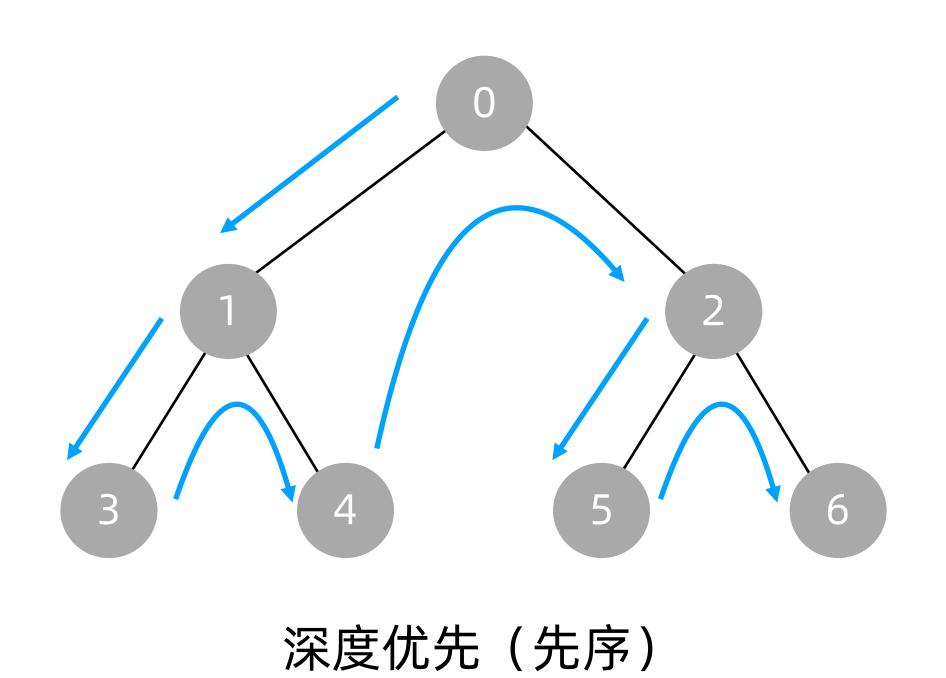


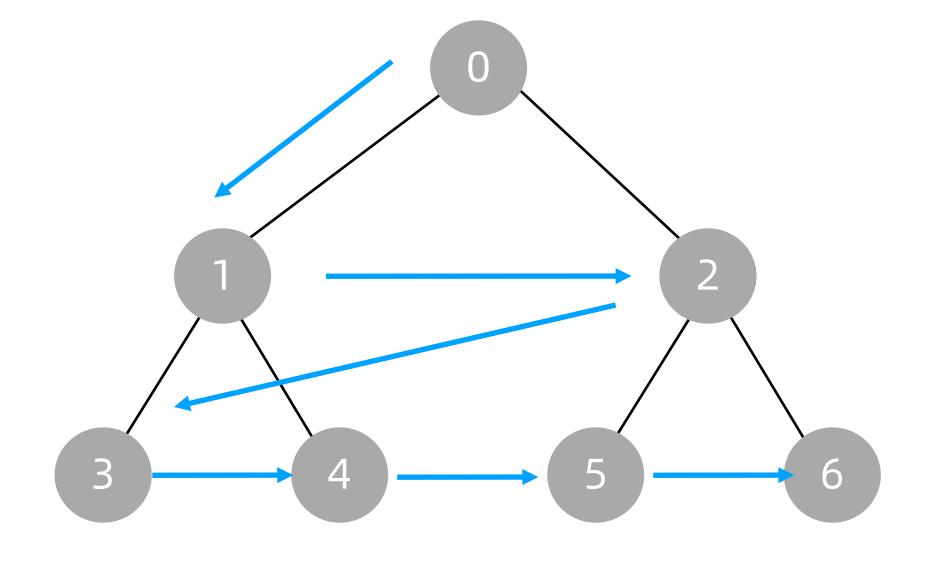
树的遍历

先序、中序、后序一般用递归来求 树的先序遍历又称树的深度优先遍历

层次序一般借助<mark>队列</mark>来求 树的层序遍历又称树的广度优先遍历

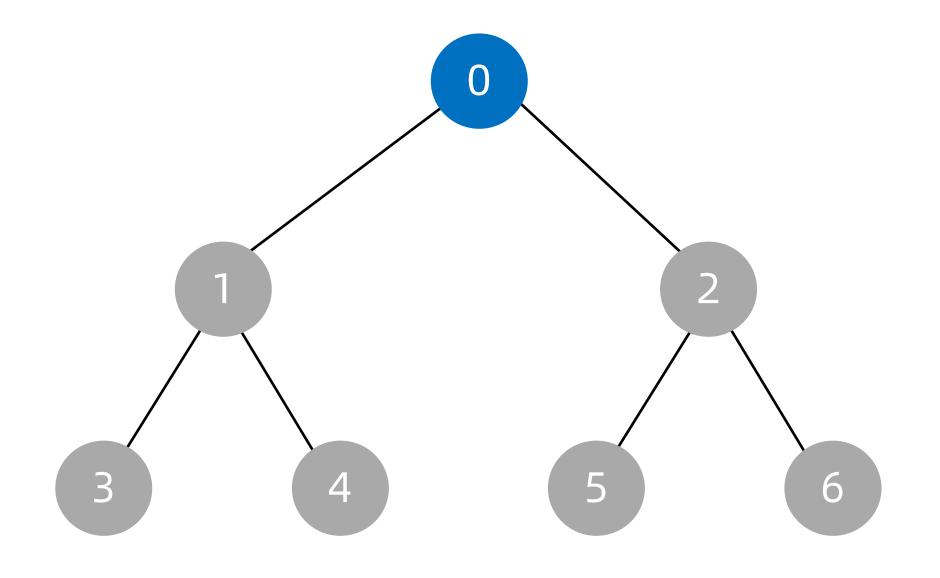
树的遍历

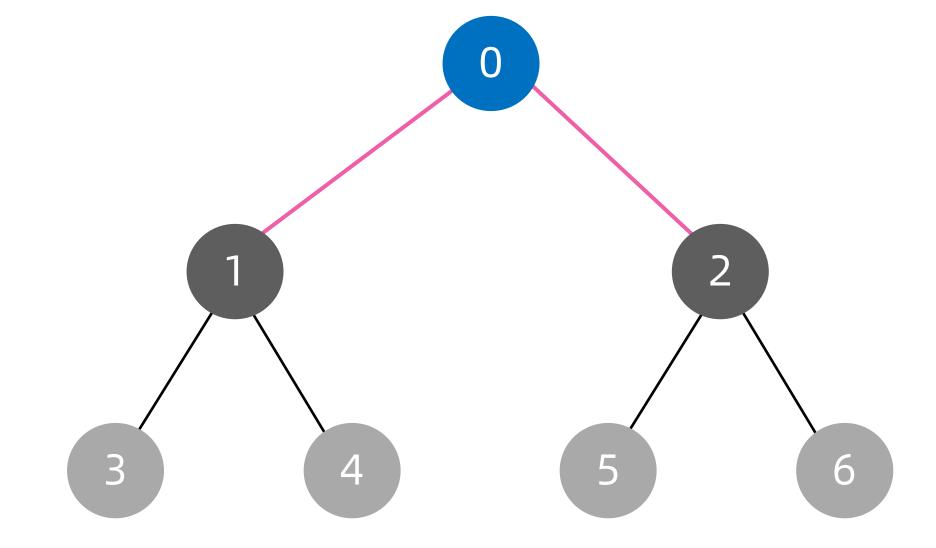




广度优先(层序)

广度优先遍历

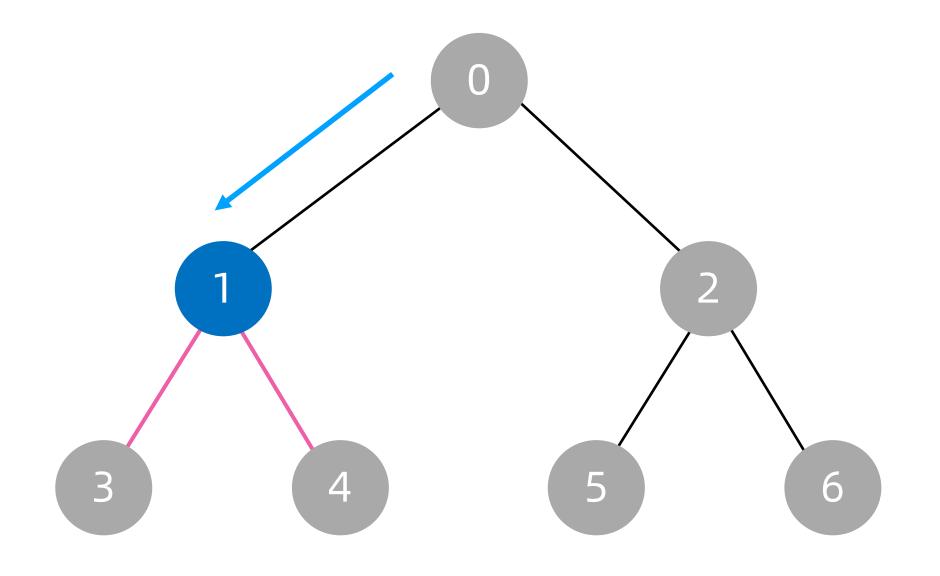


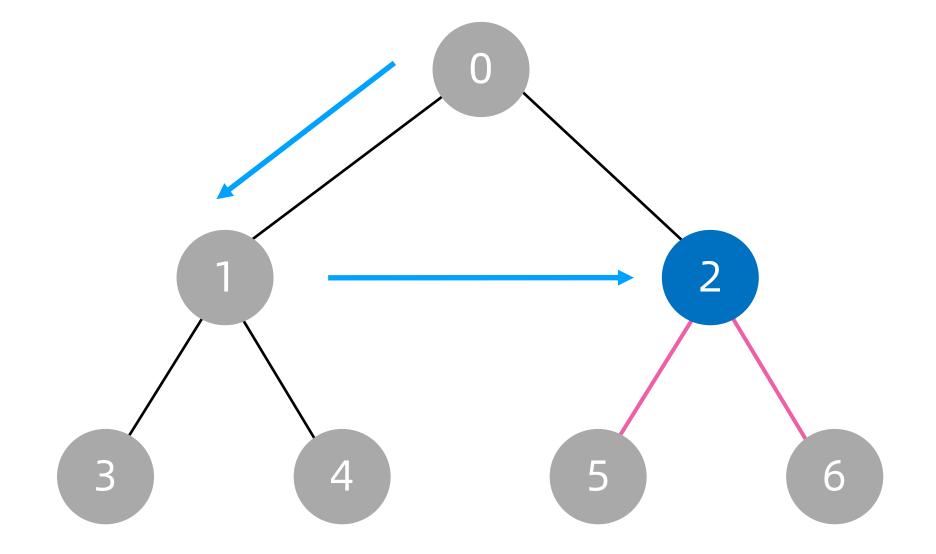


队列 O

队列 1 2

广度优先遍历



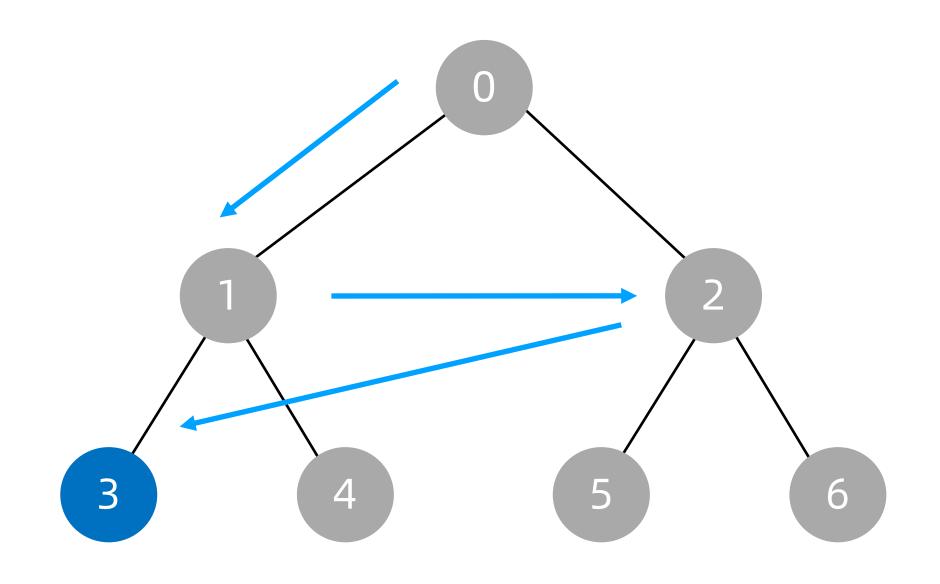


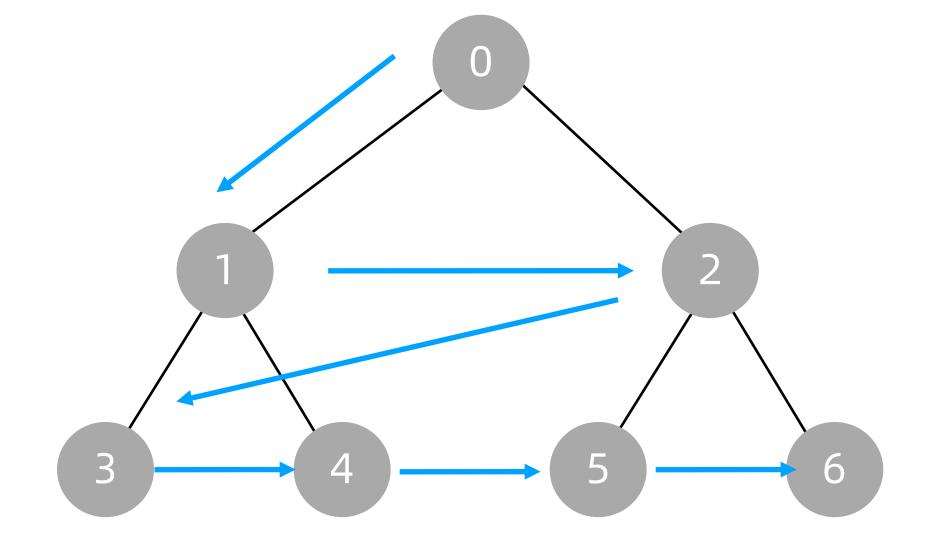
队列 234

队列

3 4 5 6

广度优先遍历





队列 4 5 6

队列

实战

二叉树的中序遍历

https://leetcode-cn.com/problems/binary-tree-inorder-traversal/

N叉树的前序遍历

https://leetcode-cn.com/problems/n-ary-tree-preorder-traversal/description/

N叉树的层序遍历

https://leetcode-cn.com/problems/n-ary-tree-level-order-traversal/

二叉树的序列化与反序列化

https://leetcode-cn.com/problems/serialize-and-deserialize-binary-tree/

实战

从前序与中序遍历序列构造二叉树

https://leetcode-cn.com/problems/construct-binary-tree-from-preorder-and-inorder-traversal/

从中序与后序遍历序列构造二叉树 (Homework)

https://leetcode-cn.com/problems/construct-binary-tree-from-inorder-and-postorder-traversal/

注意: 从前序与后序遍历序列,并不能唯一确定一棵二叉树

树的直径、最近公共祖先、树的变形

树的直径(选做)

树的直径

https://leetcode-cn.com/problems/tree-diameter/

两次深度优先遍历

- 第一次从任意一个点出发,找到距离它最远的点 p
- 第二次从 p 出发,找到距离它最远的点 q
- 连接 p,q 的路径即为树的直径

最近公共祖先 (LCA)

二叉树的最近公共祖先

https://leetcode-cn.com/problems/lowest-common-ancestor-of-a-binary-tree/

先求出父结点,然后用向上标记法

- p向上一直到 root 标红色
- q向上,第一次遇到红色时,就找到了LCA

基环树

向一棵树添加一条边,就形成了一个环 此时整个结构被称为基环树(pseudotree / unicyclic graph)

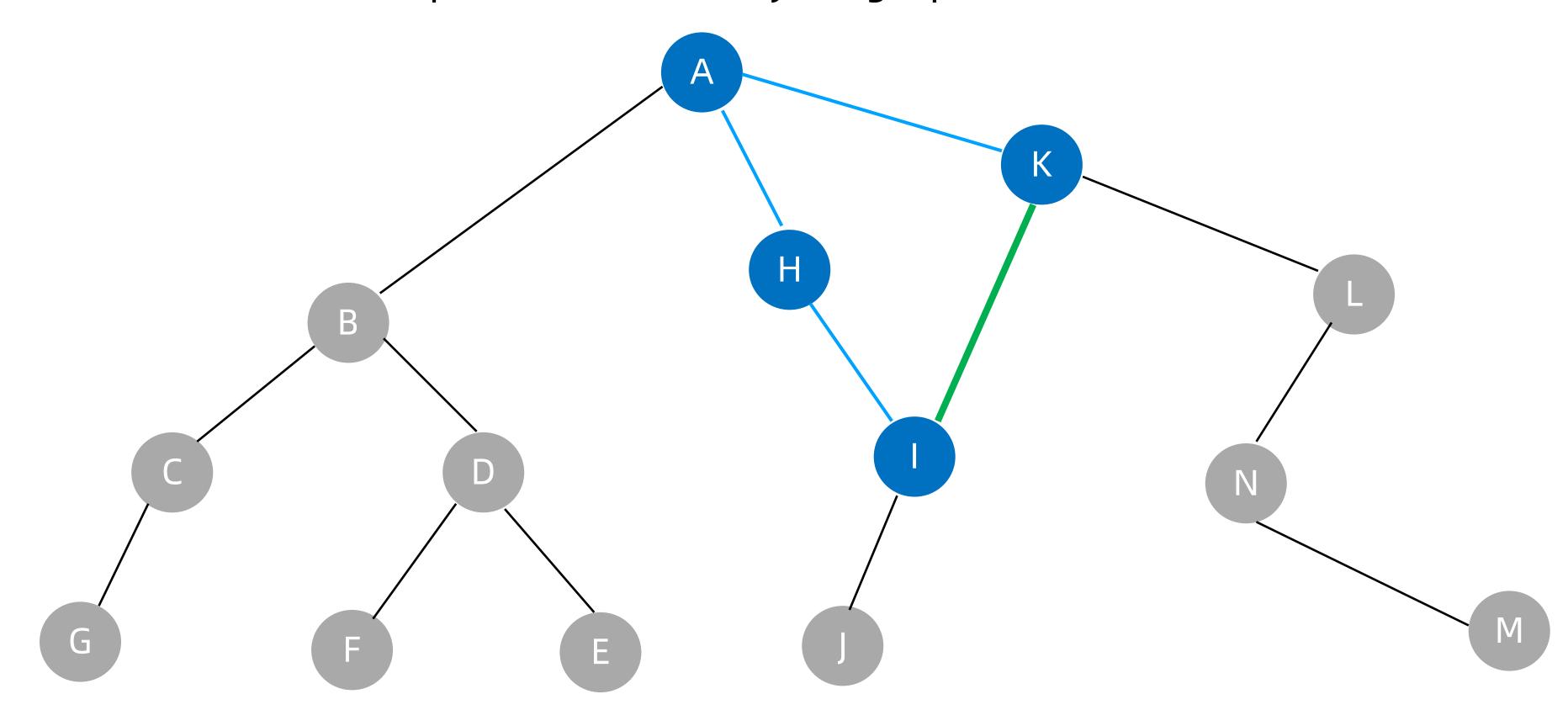
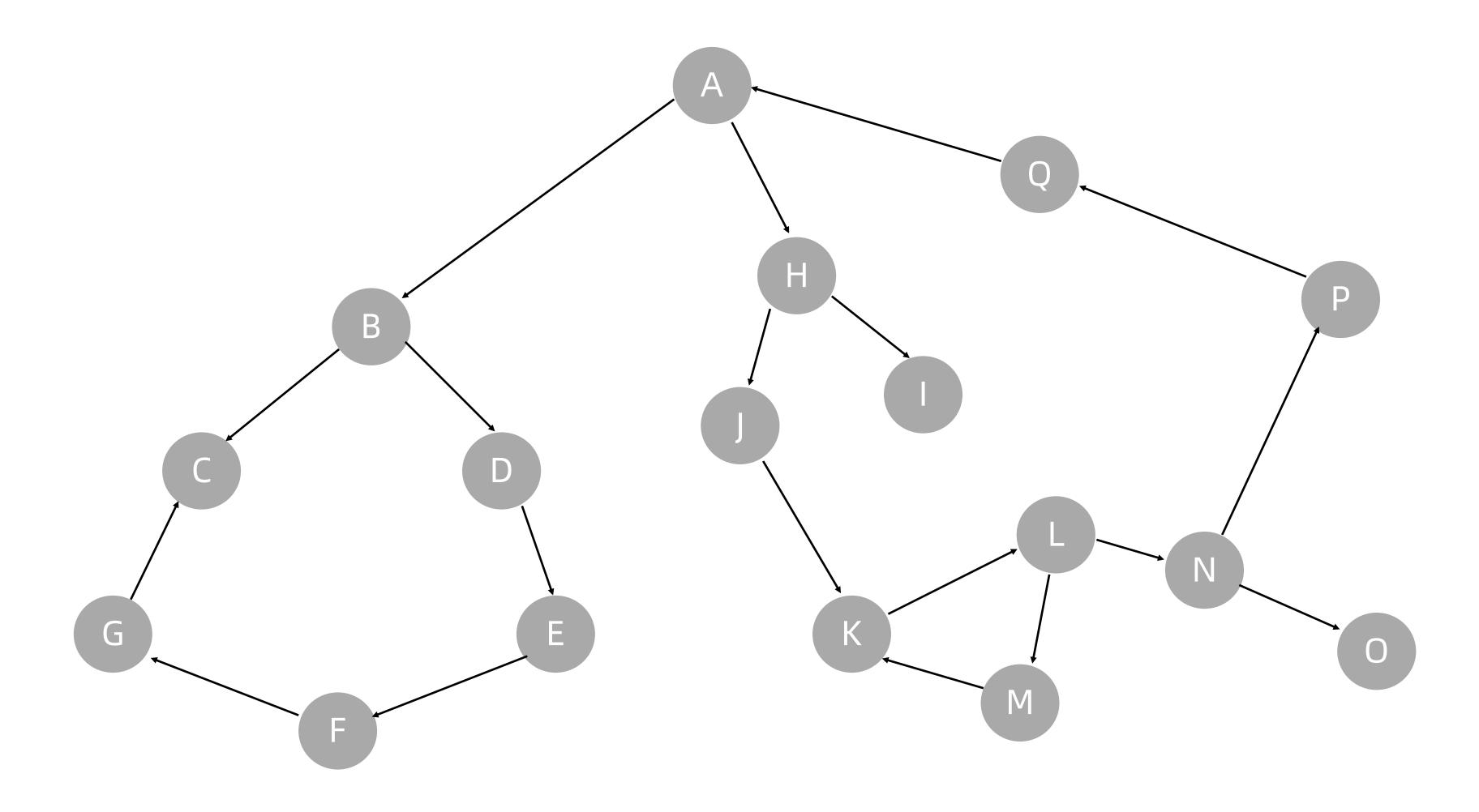


图 Graph



图的遍历

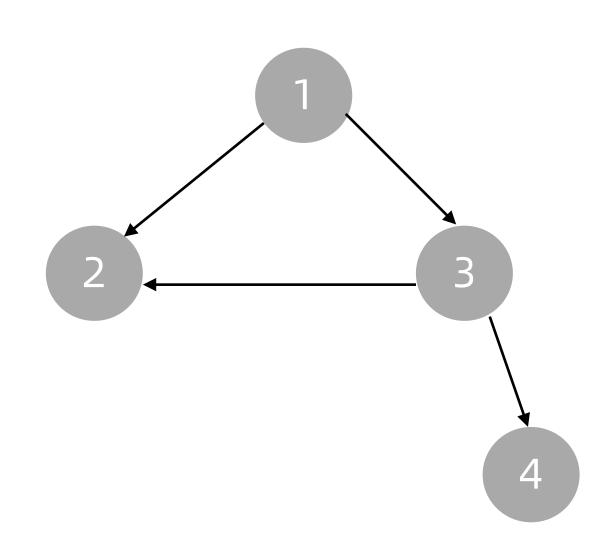
链表、树、图的关系

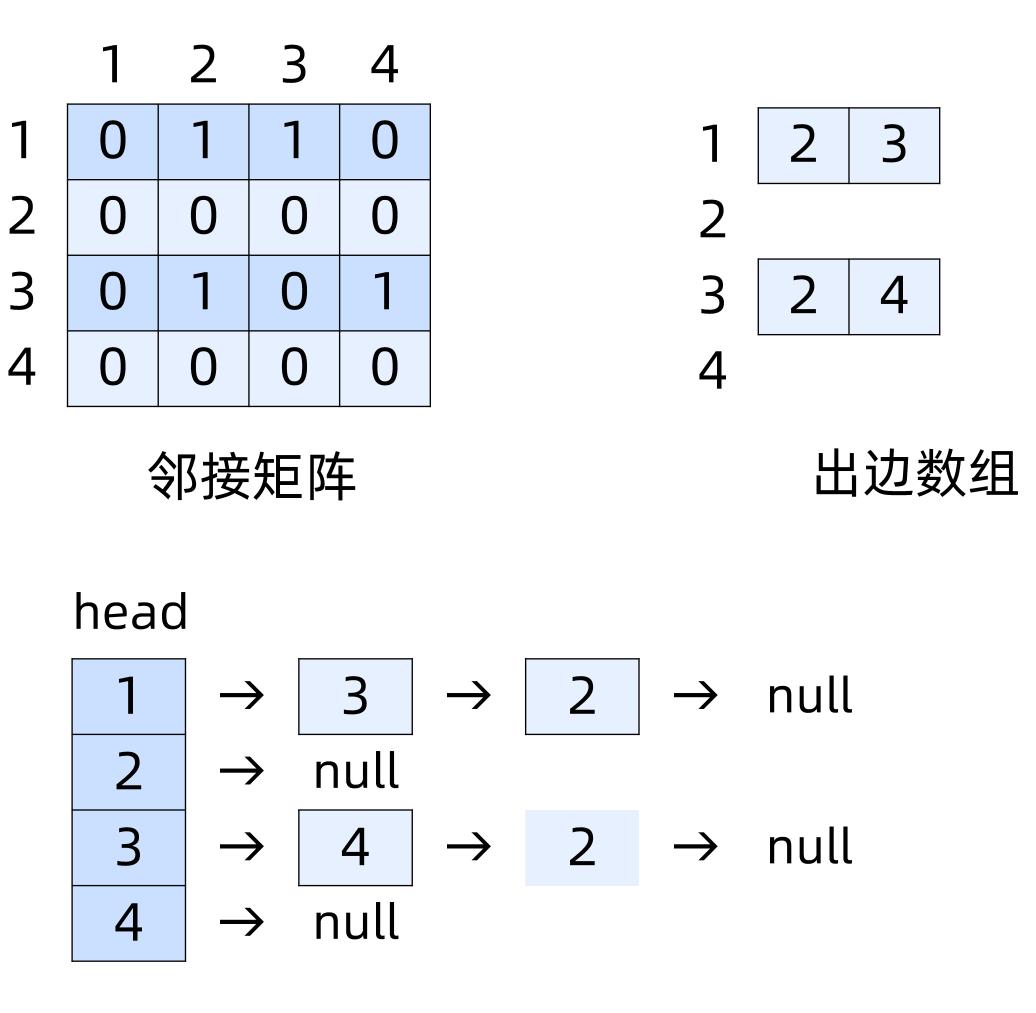
链表是特殊化的树

树是特殊化的图

- N个点 N 1 条边的连通无向图——树
- N 个点 N 条边的连通无向图——基环树

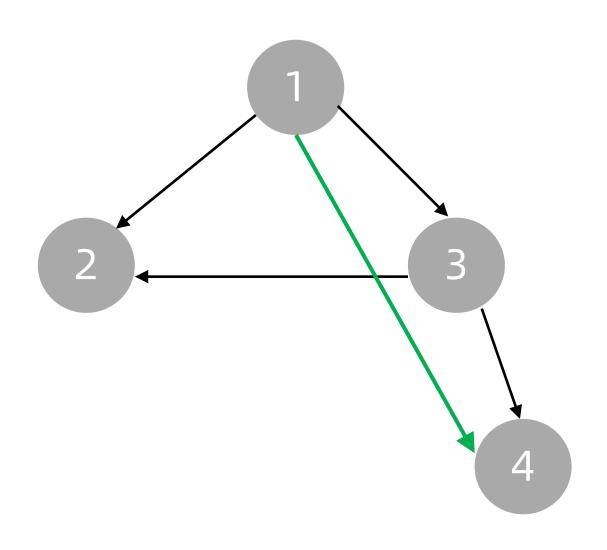
图的存储

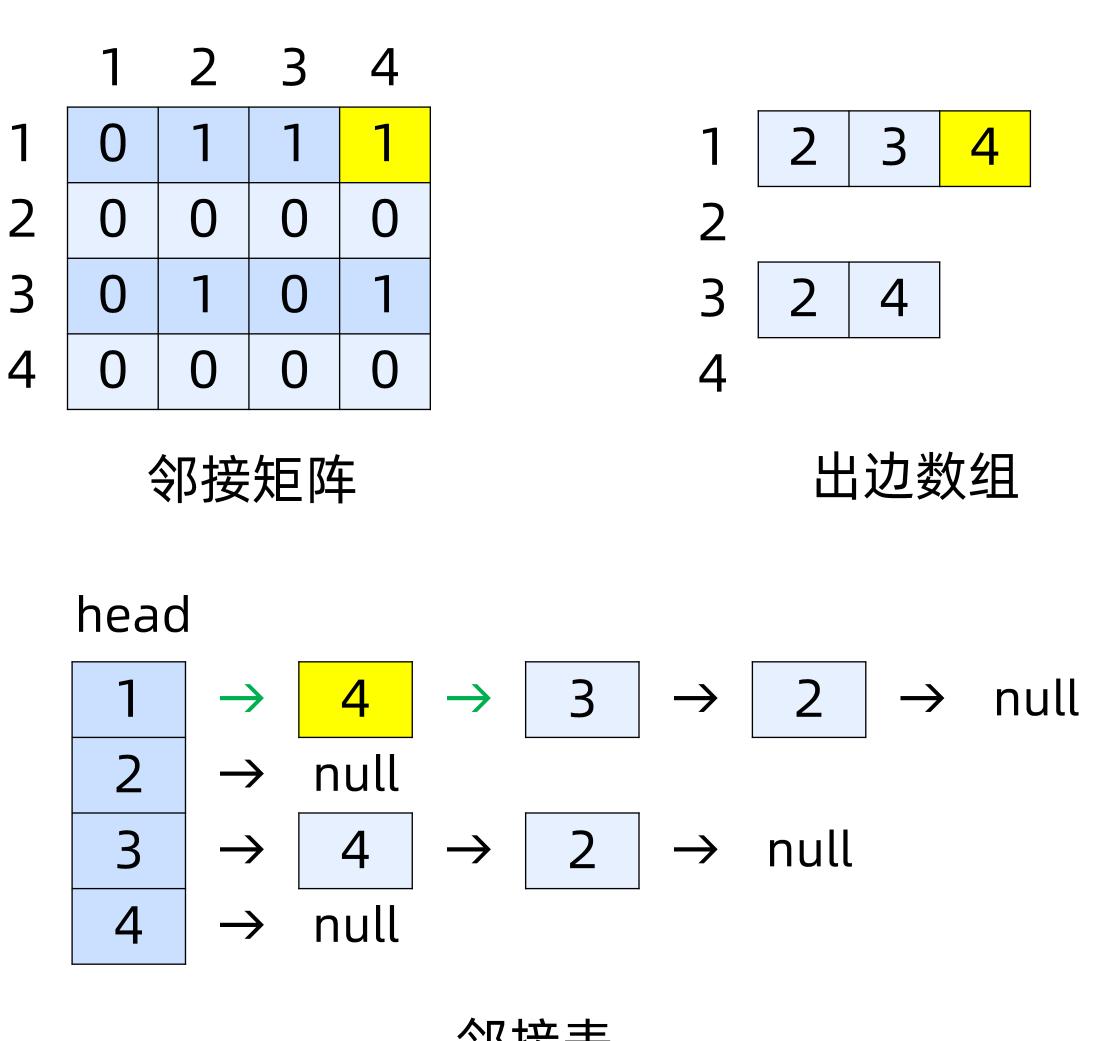




邻接表

图的存储





邻接表

图的存储

定义

```
邻接矩阵 O(n²): int graph[MAX_N][MAX_N];
出边数组 O(n+m): vector<int> graph[MAX_N];
邻接表 O(n+m): struct Node { int to; Node* next; };
Node* head[MAX_N];
```

新增边 (x,y)

- 邻接矩阵: graph[x][y] = 1;
- 出边数组: graph[x].push_back(y);
- 邻接表: Node* node = new Node();
 node->to = y;
 node->next = head[x];
 head[x] = node;

图的遍历

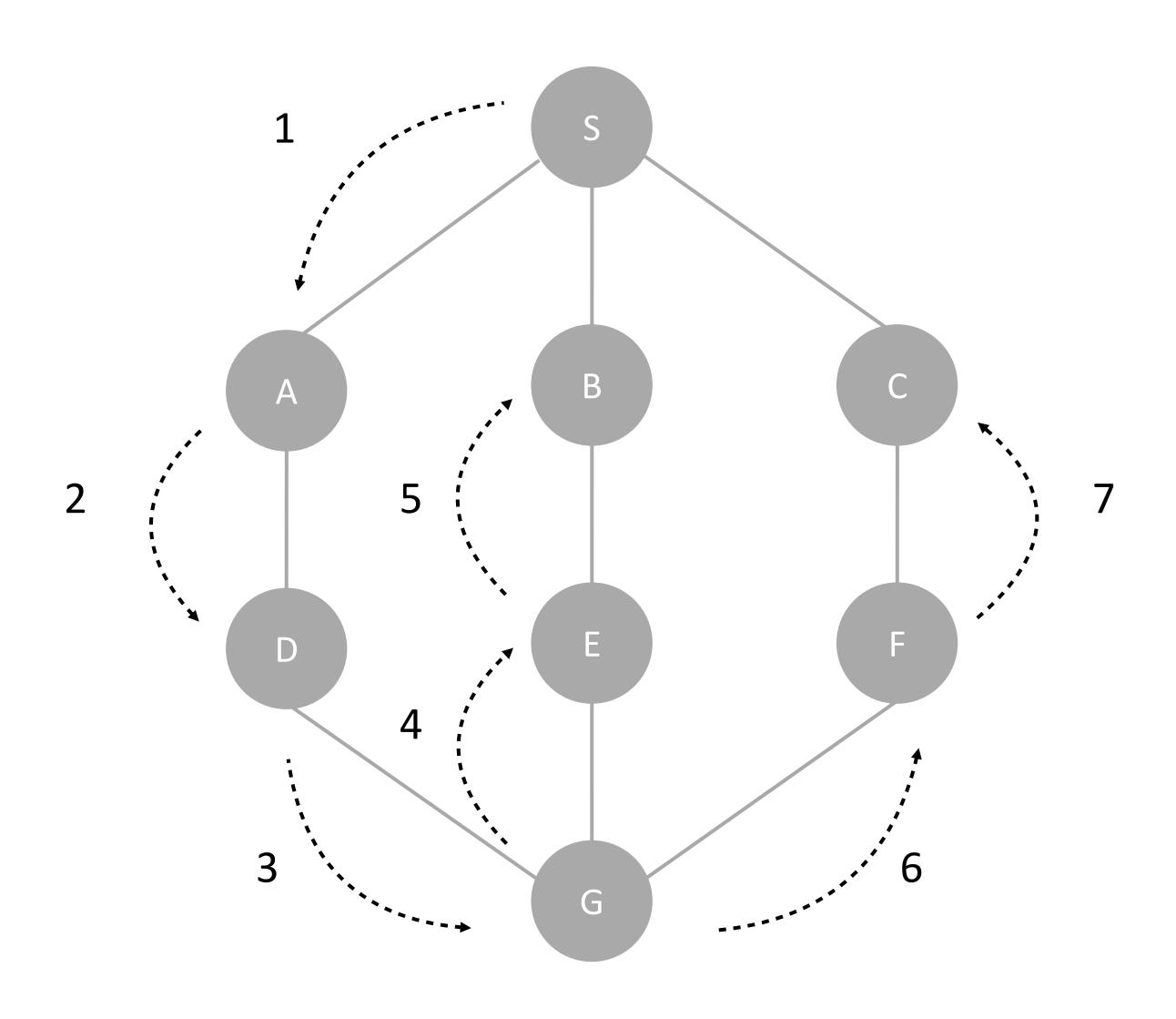
深度优先遍历

• 划分连通块

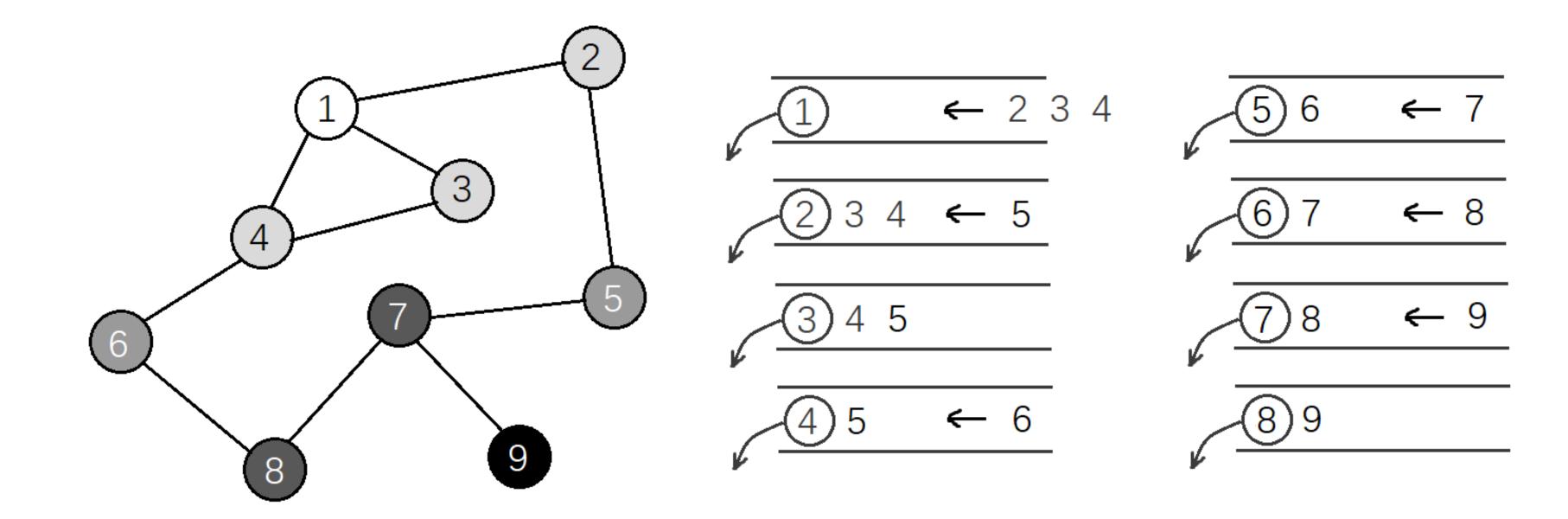
广度优先遍历

• 拓扑排序

图的深度优先遍历



图的广度优先遍历



实战

课程表

https://leetcode-cn.com/problems/course-schedule/

两种做法:

- 深度优先遍历 找环
- 广度优先遍历 拓扑排序

课程表 II (Homework)

https://leetcode-cn.com/problems/course-schedule-ii/

实战

冗余连接

https://leetcode-cn.com/problems/redundant-connection/description/

数据实际上是一棵基环树 深度优先遍历找环,环上删除一条边

冗余连接Ⅱ(选做)

https://leetcode-cn.com/problems/redundant-connection-ii/

₩ 极客时间 训练营