

Week 1

Software testing: process of executing program/system with intent of finding errors

Fault: incorrect portions of code (can be missing as well as incorrect)

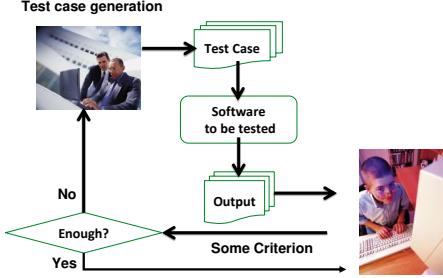
Failure: observable correct behaviour of program

Error: cause of fault, something bad programmer did (conceptual, typo, etc)

Bug: informal term for fault

Test case: set of test inputs, execution conditions, expected results developed for particular objective, such as to exercise particular program path or verify compliance with specific requirement

A Typical Software Testing Process



Testing: find inputs that cause failure of software, failure unknown, performed by testers

Debugging: process of finding & fixing fault given failure, failure is known, performed by devs

Black-Box/Functional Testing: identify functions & design test cases to check whether functions are correctly performed by software (formal & informal specs)

Equivalence partitioning: divide into partitions, select 1 test case from each partition, partitions must be disjointed (no input belongs to more than 1 partition) & all partitions must cover entire input domain

Equivalence partitioning examples: isEven then even & odd, password min 8 & max 12 characters then less than, valid, more than

Boundary-Value analysis: partition input domain, identify boundaries, select test data (for range $[R_1, R_2]$ less than R_1 , equal to R_1 , between, equal to R_2 , greater than R_2 , for unordered sets select in & not in, for equality select equal & not equal, for sets, lists select empty & not empty)

White box/structural testing: generate test cases based on program structure, abstract program to control flow graph (node is sequence of statements, edge is transfers of control)

Week 2

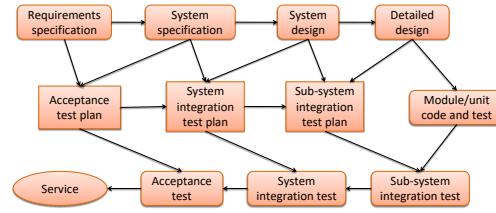
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Test oracle: expected output of software for given input, part of test case

Test driver: software framework that can load collection of test cases or test suite

Test suite: collection of test cases

The V-model of development



Testing types:

Unit/Module: test single module in isolated environment, use drivers & stubs for isolation

Integration: test parts of system by combining modules, integrated collection of modules tested as group or partial system

System: test system as a whole after integration phase

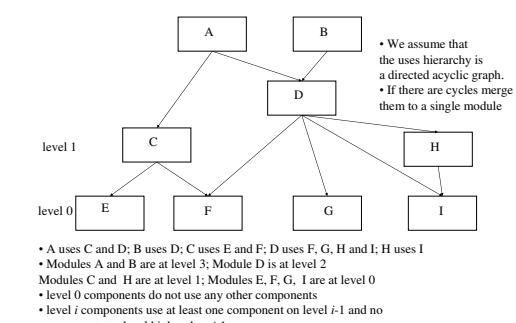
Acceptance: test system as a whole to find out if it satisfies requirements specifications

Driver: program that calls interface procedures of module being tested & reports results, simulates module that calls module currently being tested, also provides access to global variables for module under test

Stub: program that has same interface as module being used by module being tested but simpler, simulates module called by module being tested

Mock objects: create object that mimics behaviour needed for testing

Module Structure



Integration types:

Bottom-Up: only terminal modules tested in isolation, requires drivers but not stubs (since lower levels are tested already)

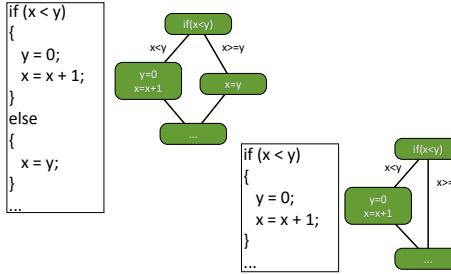
Top-down: modules tested in isolation are modules at highest level, requires stubs but not drivers

Sandwich: begin both bottom-up & top-down, meet at predetermined point in middle

Big bang: every module unit tested, then integrate all at once, no driver or stub needed but may be hard to isolate bugs

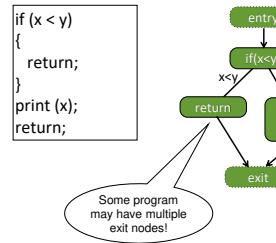
System/Acceptance testing: can construct test case based on requirements specifications, main purpose is to assure that system meets requirements, alpha testing performed within development organisation,

CFG : The if statement



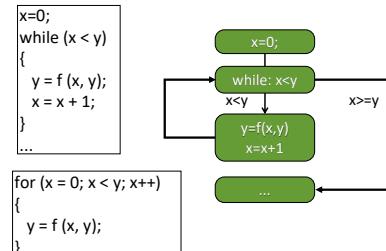
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CFG : The dummy nodes



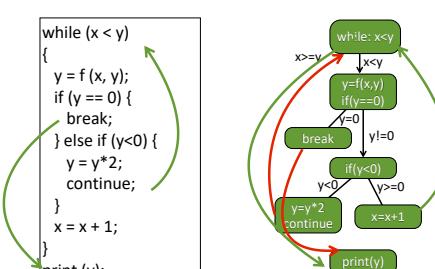
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CFG : while and for loops



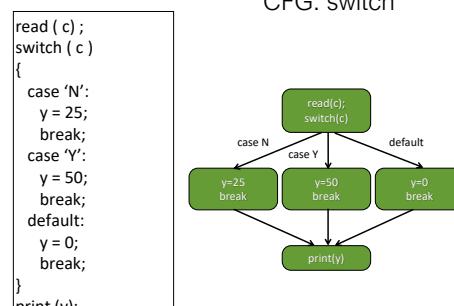
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CFG: break and continue



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CFG: switch



Coverage types: statement, branch, path (infinite if loop exists), strictly subsumes all beforehand

beta testing performed by select group of friendly customers

Week 3

Basis Path Testing: between branch & path coverage, fulfills branch testing & tests all independent paths that could be used to construct any arbitrary path through computer program

Independent path: includes some vertices/edges not covered in other path

Cyclomatic complexity: $e - n + 2p$, e is edges, n is nodes, p is number of connected components, or $1 + d$, d is loops or decision points, upper bound on number of test cases to guarantee coverage of all statements

Decision Coverage: executing true/false of decision

Condition Coverage: executing true/false of each condition

Condition/Decision Coverage: DC & CC, better than either

Multiple Condition Coverage: whether every possible combination of boolean sub-expressions occurs, test cases are truth table, 2^n test cases for n conditions

Modified C/DC: for each basic condition C , 2 test cases, values of all evaluated conditions except C are the same, compound decision as a while evaluates to true for 1 & false for the other, subsumed by MCC & subsumes CC, DC, C/DC, stronger than statement & branch

MC/DC coverage: each entry & exit point invoked, each decision takes every possible outcome, each condition in a decision takes every possible outcome, each condition in decision is shown to independently affect outcome of decision, independence of condition is shown by proving that only one condition changes at a time

MC/DC: linear complexity

- N+1 test cases for N basic conditions

$((a \mid\mid b) \And c) \mid\mid d) \And e$

Test Case	a	b	c	d	e	outcome
(1)	true	--	true	--	true	true
(2)	false	true	true	--	true	true
(3)	true	--	false	true	true	true
(4)	true	--	true	--	false	false
(5)	true	--	false	false	--	false
(6)	false	false	--	false	--	false

Underlined values independently affect the output of the decision

Week 4

Dataflow Coverage: considers how data gets accessed & modified in system & how it can get corrupted

Common access-related bugs: using undefined/uninitialized variable, deallocating/reinitialising variable before constructed/initialised/used, deleting collection object leaving members unaccessible

Variable definition: defined whenever value modified (LHS of assignment, input statement, call-by-reference)

Variable use: used whenever value read (RHS of assignment, call-by-value, branch statement predicate)

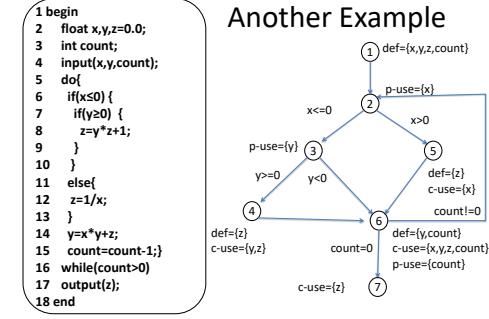
p-use: use in predicate of branch statement

c-use: any other use

Use & redefine in single statement: both sides of assignment, call-by-reference

du-pair: with respect to variable v is a pair (d, u) such that d is a node defining v , u is a node/edge using v (if p-use u is outgoing edge of predicate), there is a def-clear path with respect to v from d to u

Definition clear: with respect to variable v if no variable re-definition of v on path

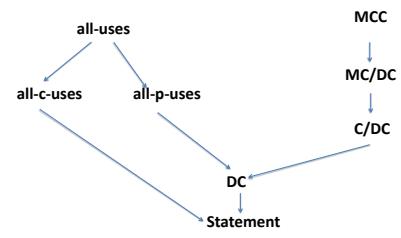


Another Example

All-C-Uses for above: $dcu(x, 1) + dcu(y, 1) + dcu(y, 6) + dcu(z, 1) + dcu(z, 4) + dcu(z, 5) + dcu(count, 1) + dcu(count, 6) = 2 + 2 + 2 + 3 + 3 + 3 + 1 + 1 = 17$

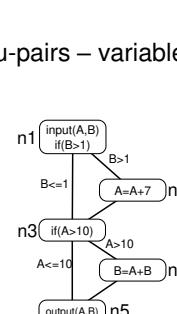
All-P-Uses for above: $dpu(x, 1) + dpu(y, 1) + dpu(y, 6) + dpu(z, 1) + dpu(z, 4) + dpu(z, 5) + dpu(count, 1) + dpu(count, 6) = 2 + 2 + 2 + 0 + 0 + 0 + 2 + 2 = 10$ (note this includes using the initial count definition even though it will always be redefined (-1) before the comparison)

Relationships among some of the coverage criteria



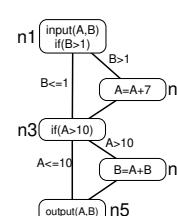
Identifying du-pairs – variable A

du-pair	path(s)
(1,2)	<1,2>
(1,4)	<1,3,4>
(1,5)	<1,3,4> <1,3,5>
(1,<3,4>)	<1,3,4>
(1,<3,5>)	<1,3,5>
(2,4)	<2,3,4>
(2,5)	<2,3,4> <2,3,5>
(2,<3,4>)	<2,3,4>
(2,<3,5>)	<2,3,5>



Identifying du-pairs – variable B

du-pair	path(s)
(1,4)	<1,2,3,4> <1,3,4>
(1,5)	<1,2,3,5> <1,3,5>
(1,<1,2>)	<1,2>
(1,<1,3>)	<1,3>
(4,5)	<4,5>



Dataflow test coverage criteria:

All-Defs: for every variable v , at least one def-clear path from every definition of v to at least one c-use or p-use of v must be covered

All-P/C-Uses: for every variable v , at least one def-clear path from every definition of v to every p/c-use of v must be covered

All-Uses: all du-pairs covered

Notations:

$d_1(x)$: definition of variable x in node i

$u_i(x)$: use of variable x in node i

$dcu(d_i(x)) = dcu(x, i)$: set of c-uses with respect to $d_1(x)$

$dpu(d_i(x)) = dpu(x, i)$: set of p-uses with respect to $d_1(x)$

Table 1. Types of Structural Coverage						
Coverage Criteria	Statement Coverage	Decision Coverage	Condition Coverage	Condition/Decision Coverage	MC/DC	Multiple Condition Coverage
Every point of entry and exit in the program has been invoked at least once		*	*	*	*	*
Every statement in the program has been invoked at least once	*					
Every decision in the program has taken all possible outcomes at least once		*	*	*	*	*
Every condition in a decision in the program has taken all possible outcomes at least once			*	*	*	*
Every condition in a decision has been shown to independently affect that decision				*		
Every combination of condition outcomes within a decision has been invoked at least once						*

Week 5

Program mutation: create artificial bugs by injecting changes to statements of programs, simulate subtle bugs in real programs

Mutation testing: software testing technique based on program mutation, can be used to evaluate test effectiveness & enhance test suite, can be stronger than control/data-flow coverage, extremely costly since need to run whole test suite against each mutant

Mutation testing steps: applies artificial changes based on mutation operators to generate mutants (each mutant with only one artificial bug), run test suite against each mutant (if any test fails mutant killed, else survives), compute mutation score

Symbolic execution/evaluation: analyse program to determine what inputs cause each part of program to execute, execute programs with symbols (track symbolic state rather than concrete input, when execute one path actually simulate many test inputs (since considering all inputs that can exercise same path))

Problems with symbolic execution:

Path explosion: 2^n paths for n branches, infinite paths for unbounded loops, calculate constraints for all paths is infeasible for real software

Constraint too complex: especially for large programs, also it is NP-complete

Input sub-domain: set of inputs satisfying path condition

Searching input to execute path: equivalent to solving associated path condition

Example

```

y = read();
y= s is a symbolic variable for input
p = 1;
while(y < 10){
    y = y + 1;
    if y >2
        p = p + 1;
    else
        p = p + 2;
}
print(p);
    
```

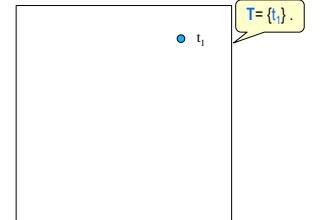
Week 6

Random testing: random number generator (monkeys) to generate test cases, also called fuzz testing, monkey testing, select tests from entire input domain (set of all possible inputs) randomly & independently, no guide towards failure-causing inputs

Adaptive Random Testing: achieve even spread of test cases

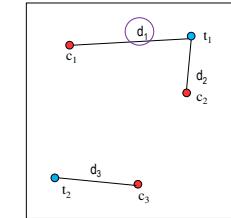
Fixed Size Candidate Set ART

Step 1. Randomly select the first input, namely t_1 , and store it in a list (called T)



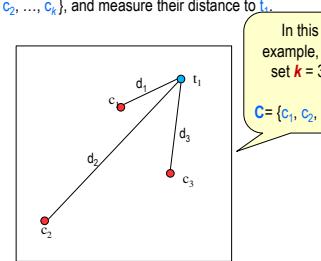
Fixed Size Candidate Set ART (cont.)

Step 6. Select the candidate with the longest distance to its nearest neighbour.



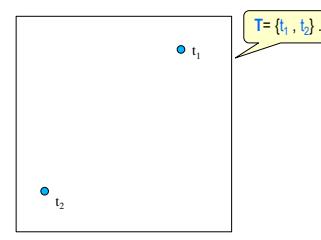
Fixed Size Candidate Set ART (cont.)

Step 2. Construct k random inputs to form a candidate set $C = \{c_1, c_2, \dots, c_k\}$, and measure their distance to t_1 .



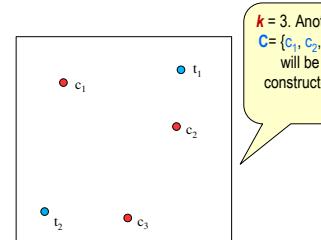
Fixed Size Candidate Set ART (cont.)

Step 3. Select the candidate which is the farthest away from t_1 to be the next test case. We name it t_2 and store it in T .



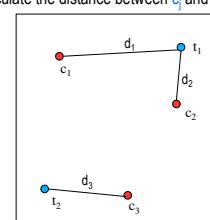
Fixed Size Candidate Set ART (cont.)

Step 4. Re-construct another candidate set C with k random inputs.



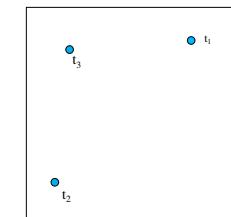
Fixed Size Candidate Set ART (cont.)

Step 5. For each candidate c_i in C , do the following
1. find which test case in T is the nearest neighbour of c_i
2. calculate the distance between c_i and its nearest neighbour.



Fixed Size Candidate Set ART (cont.)

Step 7. Store the selected candidate in T



Distance in ART:

can use Euclidean distance, if $p = (p_1, p_2, \dots, p_n)$ & $q = (q_1, q_2, \dots, q_n)$ are 2 points in n -dimensional space, $d(p, q) = d(q, p) = \sqrt{(q_1 - p_1)^2 + (q_2 - p_2)^2 + \dots + (q_n - p_n)^2}$

$$\sqrt{\sum_{i=1}^n (q_i - p_i)^2}$$

Algorithm 2:

```

initial_test.data := randomly generate a test data from the input domain;
selected.set := { initial.test.data };
counter := 1;
total_number_of.candidates := 10;
use initial.test.data to test the program;
if (program output is incorrect) then
    reveal_failure := true;
else
    reveal_failure := false;
end_if
while (not reveal_failure) do
    candidate.set := {};
    test.data := Select.The_Best_Test_Data(selected.set, candidate.set,
                                           total_number_of.candidates);
    use test.data to test the program;
    if (program output is incorrect) then
        reveal.failure := true;
    else
        selected.set := selected.set + { test.data };
        counter := counter + 1;
    end_if
end_while
output counter;
    
```

T. Y. Chen, H. Leung, and I. K. Mak. Adaptive random testing. In Proceedings of the 8th Asian Computing Science Conference, pages 320–329, 2004

ART

Algorithm /2

Random white-box testing: generate random method invocations & random parameters

Fuzz testing: random testing technique that involves providing invalid, unexpected or random data as inputs to program, commonly used to discover coding errors & unknown vulnerabilities in software, OS or networks by inputting massive amounts of random data (fuzz) to system in attempt to make it crash, cost-effective alternative to more systematic testing techniques

Fuzz Testing Example /1

- Standard HTTP GET request
 - GET /index.html HTTP/1.1

Anomalous requests:

```

GET //////////////index.html HTTP/1.1
GET %n%n%n%n%n%n.html HTTP/1.1
GET /AAAAAAAAAAAAAA.html HTTP/1.1
GET /index.html HTTTTTTTTTTTTTT/1.1
GET /index.html HTTP/1.1.1.1.1.1.1.1
    
```

Ways to generate inputs:

Mutation based/dumb fuzzing: little/no

do
Stress testing tool report: number of requests, transactions, KBps, round trip time (from user making request to receiving result), number of concurrent connection, performance degradation, types of visitors to site & number, CPU & memory use of app server

Top 2 Web App Security Risks:

Injection: SQL, OS, LDAP, occur when untrusted data sent to interpreter as part of command/query

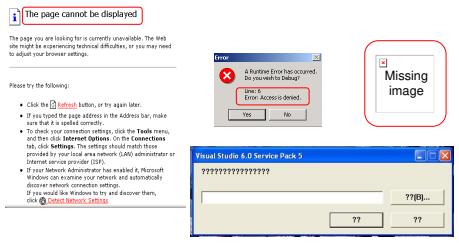
Cross Site Scripting: occur whenever app takes untrusted data & send to web browser without proper validation & escaping

Injection Protection: validate input (careful with special characters, whitelist, validate length, type, syntax), avoid use of interpreter (use stored procedures), otherwise use safe APIs (strongly typed parameterised queries, such as PreparedStatement), use Object Relational Manager

XSS Protection: appropriate encoding of all output data (HTML/XML depending on output mechanism, encode all characters other than very limited subset, specify character encoding)

Usability Testing Steps: identify website purpose, identify intended users, define tests & conduct usability testing, analyze acquired info

Usability Examples



Compatibility Testing: ensures product functionality & reliability on supported browsers & platforms that exist on customer computer

Week 8

Test management: manage test plans & cases, track requirements & defects, execute tests, measure progress

Test Report Contents:

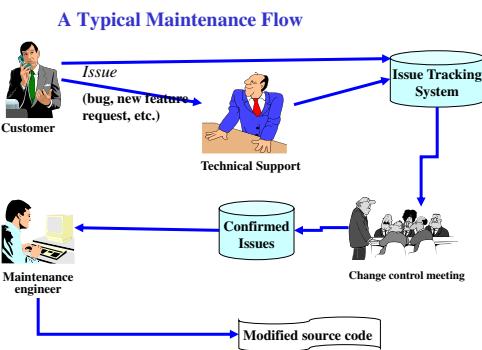
Test objective: identifying objectives of testing, should be planned so all requirements individually tested, state exit criteria

Test environment: description of test tools, required packages, hardware & software environment

Modules (features) under test & test cases: software modules to be tested & corresponding test cases, expected & actual test results

Testing schedule: overall test schedule & resource allocation

Test Summary: summary & analysis of test results, test coverage report



Issue Tracking System: software tool designed to help devs keep track of reported software issues (including bugs & new feature requests), extremely valuable in software development, used extensively by companies developing software products

Issue Tracking System Uses: track issues/bugs, communicate with teammates, submit & review patches, manage quality assurance

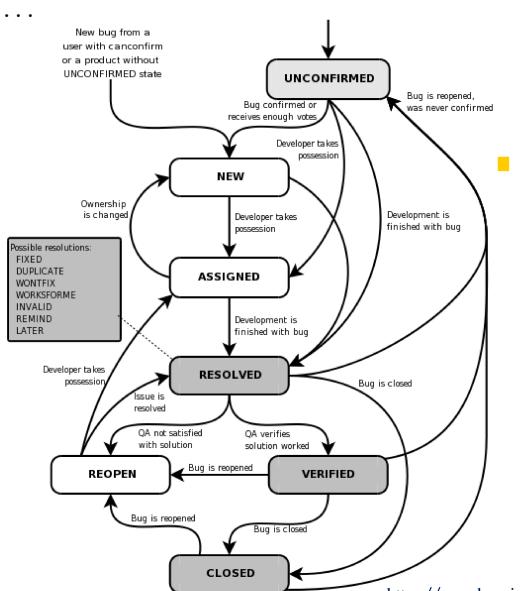
Bug Data:

Description: summary, product/module it belongs to, detailed descriptions

Role: reporter/submitter, assignee, QA contact

Status: current status, resolution, priority, severity

Time: open date, changed date, closed date, ...



Issue Tracking System Design: database, business logic, client

Useful bug reports: ones that get bugs fixed, reproducible, specific

Regression Testing: run old test cases on new version, establishing policy for regular regression testing is key for achieving successful, reliable & predictable software development projects

Regression Test Selection: speed up regression testing by only rerunning tests affected by code changes

Regression Test Prioritisation: rank all test cases (to discover bugs sooner or achieve higher coverage sooner), run test cases according to ranked sequence, stop when all resources used up

Reliability: probability that system/capability of system functions without failure for specified time in specified environment, ex: reliability of 0.92 for 8 hours means when executed for that long would operate without failure for 92 out of 100 periods

Single failure specification: what is probability of failure of a system/component

Multiple failure specification: if system/component fails at time t_1, t_2, \dots, t_{i-1} what is probability of failure at time t_i

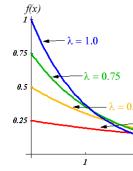
Reliability formulas: $R(t) = P(T > t), F(t) = P(T \leq t), t \geq 0$

Probability Density Function (PDF)



Probability density function (PDF):

depicting changes of the probability of failure up to a given time t



Exponential PDF:

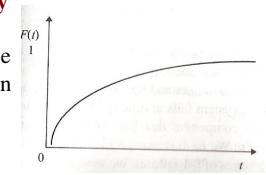
$$f(t) = \begin{cases} \lambda e^{-\lambda t} & t \geq 0 \\ 0 & t < 0 \end{cases}$$

Cumulative Density Function (CDF)



Cumulative density function (CDF):

depicting cumulative failures up to a given time t



For exponential distribution, CDF is:

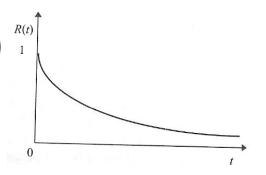
$$F(t) = \begin{cases} 1 - e^{-\lambda t} & t \geq 0 \\ 0 & t < 0 \end{cases}$$

Reliability Function (R)



Reliability function (R): depicting a component functioning without failure until time t , also called (survival function)

$$R(t) = 1 - F(t)$$



For exponential distribution, R is:

$$R(t) = e^{-\lambda t}$$



Reliability: Example

- Assume the failures of a computer system follows exponential distribution, and the probability of failure is 5% within 100 hours. What's the probability of the system working without failures in 1000 hours? (Hint: $\ln(0.95) = -0.05$; $\ln(0.61) = -0.5$).

$$P(t \leq 100) = F(100) = 0.05$$

$$1 - e^{-100\lambda} = 0.05$$

$$\lambda = 0.0005$$

$$R(1000) = e^{-1000 \times 0.0005} = 0.6065$$

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Reliability Distributions: exponential, Poisson, lognormal, normal, Pareto, Weibull

Software Reliability Models

- Goel-Okumoto (G-O) model

$$m(t) = a(1 - \exp[-bt]), a > 0, b > 0,$$

- Gompertz growth curve model

$$m(t) = ab^t, a > 0, 0 < b < 1, 0 < k < 1,$$

- Logistic growth curve model

$$m(t) = \frac{a}{1 + k \exp[-bt]}, a > 0, b > 0, k > 0,$$

- Yamada delayed S-shaped model

$$m(t) = a(1 - (1 + bt) \exp[-bt]), a > 0, b > 0,$$

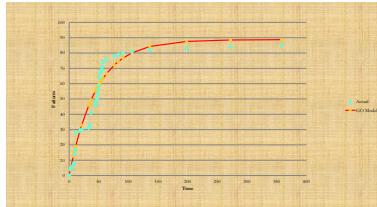
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Goel-Okumoto (G-O) Model /1

- Proposed by Amrit Goel of Syracuse University and Kazu Okumoto in 1979.
- Based on NHPP(Non-Homogeneous Poisson Process)
- An example: the number of failures (y) of a software system at time t:

i	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
t(i)	0	2	3	7	8	9	10	11	18	21	33	35	37	44	45	47	48
y(i)	0	4	5	7	8	14	17	28	29	30	31	33	41	46	48	50	53
i	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33
t(i)	49	50	51	52	53	55	56	57	63	76	83	91	106	136	198	272	358
y(i)	56	59	64	67	68	69	71	74	76	78	79	80	81	82	83	84	85

Goel-Okumoto (G-O) Model /2



$$a = 88.733, b = 0.022$$

G-O model: #Failures = $88.733 \cdot (1 - \exp(-0.022t))$

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Mean time to failure: mean of probability density, expected value of T, average lifetime of system, $E(T) = \int_0^\infty t f(t) dt = \int_0^\infty R(t) dt$, for exponential is $\frac{1}{\lambda}$

Mean time between failures: MTTF + MTTR (mean time to repair)

Software reliability tools tasks: collecting failure & test time info, calculating estimates of model parameters using this info, testing to fit model against collected info, selecting model to make predictions of remaining faults, time to test, apply model

Week 9

Software reviews: quality improvement processes for written material, by detecting defects early & preventing leakage downstream higher cost of later detection & rework eliminated

Software products that can be reviewed: requirements specifications, design descriptions, source code (code review), release notes

Code review types: ad-hoc review, pass-round, walkthrough, group review, formal inspection

Formal Inspection: planning/overview, preparation (product docs, rules/checklist), inspection, rework

Code review steps: perform examination of software products, detect defects (bugs), violation of coding standards, code smells, other problems, look for code patterns that indicate problems based on prior xp, static analysis tools can also help

Bug patterns: infinite recursion, null pointer bugs, SQL injection, divide by 0, buffer overflow, memory leak, deadlock, infinite loop, XSS

Code smells: indications of poor coding & design choices that can cause problems during later phase of development, hint something gone wrong somewhere

A List of Code Bad Smells

- Duplicated Code
- Long Method
- Large Class (Too many responsibilities)
- Long Parameter List (Object is missing)
- Feature Envy (Method needing too much information from another object)
- Lazy Class (Do not do too much)
- Middle Man (Class with too much delegating methods)
- Temporary Field (Attributes only used partially under certain circumstances)
- Message Chains (Coupled classes, internal representation dependencies)
- Data Classes (Only accessors)
- ...

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Checking Coding Conventions /1

- Your code should be readable!
- Formatting conventions ensure consistency and therefore, familiarity for readers
- Indentation: Indent when starting out new blocks of code

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Checking Coding Conventions /2

- Sun's Coding Conventions for Java
→ <http://www.oracle.com/technetwork/java/codeconvtoc-136057.html>
- JavaScript Coding Style
→ <https://google-styleguide.googlecode.com/svn/trunk/javascriptguide.xml>
- Code Conventions for the JavaServerPages Technology
→ http://java.sun.com/developer/technicalArticles/javalaserverpages/code_convention/

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Checking Coding Conventions /3

- Checking identifier names:**
 - Choosing meaningful names
 - Class names start upper-cased
 - e.g., BankAccount, Vehicle, VehicleApplet
 - Variable and Method names start lower-cased
 - e.g., aliceAccount, hondaCivic, nCount, etc.
 - Constants are ALL_CAPS and words in name are separated by an underscore
 - e.g., PI, MAX_WIDTH, DEFAULT_WIDTH

Checking Coding Conventions /3

- Block-style comments

/* This is a multi-line comment.

* Use when you need to write a long comment about a fragment

*/

- One-line comments

/* C-style comments */

* use to put descriptive notes before a code fragment

// C++ style comment

* use at the end of the line, to describe variables or short pieces of code

- javadoc comments

- like block-style, but starts with /** instead

- use immediately before classes, methods, and fields

Checking comments

Code Review benefits: can find 60–100% of defects, can assess/improve quality of work product, software development process & review process itself, reduce total project cost but have non-trivial cost (15%), early defect removal is 10–100 times cheaper, reviews distribute domain knowledge, dev skills, corporate culture

Common problems in code review: insufficient preparation, moderator domination, incorrect review rate, ego involvement & personality conflict, issue resolution & meeting digression, recording difficulties & clerical overhead

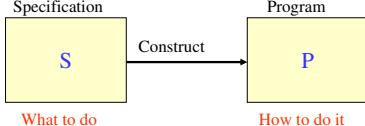
Static Analysis: analyse program without executing, doesn't depend on test cases, generally doesn't know what the software is supposed to do, looks for bug patterns, no replacement for testing, many defects can't be found with static analysis

Patterns to be checked: bad practice, correctness, performance, dodgy code, vulnerability to malicious code

Pattern examples: equals method should not assume type of object argument, collection should not contain themselves (`!s.contains(s)`), should not use `String.toString()`

Week 10

Problems in software development



What to do How to do it

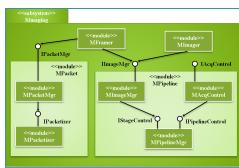
- How to ensure that S is not ambiguous so that it can be correctly understood by all the people involved?
- How can S be effectively used for inspecting and testing P?
- How can software tools effectively support the analysis of S, transformation from S to P, and verification of P against S?

Example of informal specifications

Natural language:

"A software system for an Automated Teller Machine (ATM) needs to provide services on various accounts. The services include ... The operations on a current or savings account include deposit, withdrawal, show balance, and print out transaction records."

Graphical:



The major problems with informal specifications:

- Informal specifications are likely to be **ambiguous**, which is likely to cause **misinterpretations**.
- Informal specifications are **difficult** to be used for **inspection and testing of programs** because of the **big gap** between the functional descriptions in the specifications and the program structures.
- Informal specifications are **difficult to be analyzed** for their **consistency and validity**.

Formal Methods

Formal methods include

- Formal specification
- Formal verification

These are all based on mathematical representation and analysis of software

- Set theory, (temporal/first-order/higher-order) logics, automata, etc.
- Unambiguous syntax & sound semantics

Formal Specifications

- Formal specification is concerned with producing an unambiguous set of product specifications so that customer requirements, as well as environmental constraints and design intentions are correctly reflected, thus reducing the chances of accidental fault injections.
- Formal specifications typically focus on the functional aspect or the correctness of expected program behaviour.
- With formal specifications, the desirable properties for software specifications can be more easily and sometimes formally analyzed and assured.

Formal Specifications

- Many formal specifications are descriptive.
- Descriptive specifications** focus on the properties or conditions associated with software products and their components.
 - Model (or state-based) specifications focus on states and models (e.g Z and VDM languages).
 - Algebraic specifications focus on functional computation carried out by a program or program-segment and related properties. (e.g., Larch, Common Algebraic Specification Language)
- Foundations: set theory and logic

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Formal Methods – Benefits / limitations

- + Requirements and specifications are unambiguous
- + Errors due to misunderstandings are reduced
- + Eases implementation
- + Correctness proofs can be carried out
- + Validation of requirements specifications is possible
- Formal specifications are difficult to read
- Can't model all aspects of real world
- Correctness proofs are resource intensive
- Development cost increases

Sets

- Are collections of elements
- Are represented by standard notation
- Are manipulated by standard elementary operations
- Are entities which can interact with each other

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Formal Verification

- Formal verification techniques attempt to show the absence of faults.
- Software testing shows the presence of defects, not their absence.
- The basic idea is to verify the correctness, or absence of faults, of some program code or design, against its formal specifications.
- The presence of a formal specification is a prerequisite for formal verifications.

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Examples of Sets

Notation: {}

- Set of colours: {green,blue,yellow}
- Set of sports: {tennis,football,equestrian}
- Set of courses: {SENG3130,SENG6140}
- Empty set: {} or Ø

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An Example of Formal Specification in Z

<i>BirthdayBook</i>	Gives the name of the schema
<i>known : P NAME</i>	Variable declarations go above the dividing line
<i>birthday : NAME → DATE</i>	
<hr/>	
<i>known = dom birthday</i>	Conditions go beneath the line
<hr/>	
<i>AddBirthday</i>	Name of the Schema
<i>known : P NAME</i>	Explicit inclusion of all declarations
<i>known' : P NAME</i>	
<i>birthday : NAME → DATE</i>	
<i>birthday' : NAME → DATE</i>	
<i>name? : NAME</i>	
<i>date? : DATE</i>	
<hr/>	
<i>known = dom birthday</i>	Invariants included by hand
<i>known' = dom birthday'</i>	
<i>name? ⊏ known</i>	Predcondition
<i>birthday' = birthday ∪ {name? ↦ date?}</i>	Postcondition

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Set Construction

{set : range | condition • operation}

- In notational form (aka comprehensive specification):
(Signature | Predicate • Term)

{x : X | P(x) • E(x)}

Alternate even numbers:

*{x : N | x mod 2 = 0 • 2 * x} = {0,4,8,12,16,20,...}*

Tens:

*{x : N | x · 10 * x} = {0,10,20,30,40,...}*

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Sets:

- S : P X* S is declared as a set of X's.
- x ∈ S* x is a member of S.
- x ∉ S* x is not a member of S.
- S ⊆ T* S is a subset of T: Every member of S is also in T.
- S ∪ T* The union of S and T: It contains every member of S or T or both.
- S ∩ T* The intersection of S and T: It contains every member of both S and T.
- S \ T* The difference of S and T: It contains every member of S except those also in T.
- Ø* Empty set: It contains no members.
- {x}* Singleton set: It contains just x.
- N* The set of natural numbers 0, 1, 2,
- S : F X* S is declared as a finite set of X's.
- max (S)* The maximum of the nonempty set of numbers S.

Functions:

- f:X→Y* f is declared as a partial injection from X to Y.
- dom f* The domain of f: the set of values x for which f(x) is defined.
- ran f* The range of f: the set of values taken by f(x) as x varies over the domain of f.
- f ⊕ {x → y}* A function that agrees with f except that x is mapped to y.
- {x} ⊑ f* A function like f, except that x is removed from its domain.

Logics:

- P ∧ Q* P and Q: It is true if both P and Q are true.
- P ⇒ Q* P implies Q: It is true if either Q is true or P is false.
- ØS' = ØS* No components of schema S change in an operation.

Set Operators

- Examples:

- Newcastle $\in \{\text{NSW cities}\}$
- Melbourne $\notin \{\text{NSW cities}\}$
- $\{\text{Newcastle, Sydney}\} \subseteq \{\text{NSW cities}\}$
- $\#\{\text{Newcastle, Sydney}\} = 2$
- $\{\text{Tom, Jim, James}\} \cup \{\text{James, Kathy}\} = \{\text{Tom, Jim, James, Kathy}\}$
- $\{\text{Tom, Jim, James}\} \cap \{\text{James, Kathy}\} = \{\text{James}\}$
- $\{\text{Tom, Jim, James}\} \setminus \{\text{James, Kathy}\} = \{\text{Tom, Jim}\}$

An Example

<i>BirthdayBook</i>	Gives the name of the schema
<i>known : P NAME</i>	Variable declarations go above the dividing line
<i>birthday : NAME → DATE</i>	
<i>known = dom birthday</i>	Conditions go beneath the line

■ One possible state of the system has three people in the set *known*, with their birthdays recorded by the function *birthday*:

```
known = {John; Mike; Susan}
birthday = {John 25-Mar,
            Mike 20-Dec,
            Susan 20-Dec}
```

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The Z Language

- A well-known language formal specification
- Based on set theory
- Equally used to model (specify) state as well as operations on states
- First developed in 1977–1990 at the University of Oxford with industrial partners (IBM, Inmos)

<i>AddBirthday</i>	Name of the Schema
<i>known : P NAME</i>	Explicit inclusion of all declarations
<i>known' : P NAME</i>	
<i>birthday : NAME → DATE</i>	
<i>birthday' : NAME → DATE</i>	
<i>name? : NAME</i>	
<i>date? : DATE</i>	
<i>known = dom birthday</i>	Invariants included by hand
<i>known' = dom birthday'</i>	
<i>name? ⊏ known</i>	Precondition
<i>birthday' = birthday ∪ {name? ↦ date?}</i>	Postcondition

- *name* is the before state
- *name'* is the after state
- *name?* is an input variable
- *name!* is an output variable

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An Example

Standard Example from the Z reference manual:

The Birthday Book

- Stores a list of names with associated birthdays
- Operations for:
 - Add new people
 - Query for a given persons birthday
 - Query for all birthdays on a given date
 - ...

For the birthday book, we have to deal with two basic types:

- Names of people
- Birthday dates (day and month)

To introduce the two types, we just specify them: *[NAME, DATE]*.

The simplified version:

<i>AddBirthday</i>	
Δ <i>BirthdayBook</i>	
<i>name? : NAME</i>	
<i>date? : DATE</i>	
<i>name? ⊏ known</i>	
<i>birthday' = birthday ∪ {name? ↦ date?}</i>	

Δ is implicitly used to:

- Include SchemaName and SchemaName'
- Signal to the reader that this modifies the state

A modifying operation describes the relationship of two states, one before and one after the operation

Pre- and Postconditions

Precondition:

- *known* = dom *birthday*

- *name?* $\not\in$ *known*

Postcondition:

- *birthday'* = *birthday* $\cup \{\text{name?} \mapsto \text{date?}\}$

Formal verification can be performed based on Precondition and Postcondition