

Lectures 2&3: Introduction to SQL

Lecture 2: SQL Part I

Today's Lecture

1. SQL introduction & schema definitions
 - ACTIVITY: Table creation
2. Basic single-table queries
 - ACTIVITY: Single-table queries!
3. Multi-table queries
 - ACTIVITY: Multi-table queries!

1. SQL Introduction & Definitions

What you will learn about in this section

1. What is SQL?
2. Basic schema definitions
3. Keys & constraints intro
4. ACTIVITY: CREATE TABLE statements

SQL Motivation

- Dark times 5 years ago.
 - Are databases dead?
- Now, as before: everyone sells SQL
 - Pig, Hive, Impala
- “Not-Yet-SQL?”



Basic SQL

SQL Introduction

- SQL is a standard language for querying and manipulating data
- SQL is a **very high-level** programming language
 - This works because it is optimized well!
- Many standards out there:
 - ANSI SQL, SQL92 (a.k.a. SQL2), SQL99 (a.k.a. SQL3),
 - Vendors support various subsets

SQL stands for
Structured Query Language

NB: Probably the world's most successful **parallel** programming language (multicore?)

SQL is a...

- Data Definition Language (DDL)
 - Define relational *schemata*
 - Create/alter/delete tables and their attributes
- Data Manipulation Language (DML)
 - Insert/delete/modify tuples in tables
 - Query one or more tables – discussed next!

Tables in SQL

Product

PName	Price	Manufacturer
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

A relation or table is a multiset of tuples having the attributes specified by the schema

Let's break this definition down

Tables in SQL

Product

PName	Price	Manufacturer
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

A multiset is an unordered list (or: a set with multiple duplicate instances allowed)

List: [1, 1, 2, 3]

Set: {1, 2, 3}

Multiset: {1, 1, 2, 3}

i.e. no *next()*, etc. methods!

Tables in SQL

Product

PName	Price	Manufacturer
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

An attribute (or column) is a typed data entry present in each tuple in the relation

NB: Attributes must have an atomic type in standard SQL, i.e. not a list, set, etc.

Tables in SQL

Product

PName	Price	Manufacturer
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

Also referred to sometimes as a record

A tuple or row is a single entry in the table having the attributes specified by the schema

Tables in SQL

Product

PName	Price	Manufacturer
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks
SingleTouch	\$149.99	Canon
MultiTouch	\$203.99	Hitachi

The number of tuples is the cardinality of the relation

The number of attributes is the arity of the relation

Data Types in SQL

- Atomic types:
 - Characters: CHAR(20), VARCHAR(50)
 - Numbers: INT, BIGINT, SMALLINT, FLOAT
 - Others: MONEY, DATETIME, ...
- Every attribute must have an atomic type
 - Hence tables are flat

Why?

Table Schemas

- The **schema** of a table is the table name, its attributes, and their types:

```
Product(Pname: string, Price: float, Category:  
string, Manufacturer: string)
```

- A **key** is an attribute whose values are unique; we underline a key

```
Product(Pname: string, Price: float, Category:  
string, Manufacturer: string)
```

Key constraints

A key is a **minimal subset of attributes** that acts as a unique identifier for tuples in a relation

- A key is an implicit constraint on which tuples can be in the relation
 - i.e. if two tuples agree on the values of the key, then they must be the same tuple!

`Students(sid:string, name:string, gpa: float)`

1. Which would you select as a key?
2. Is a key always guaranteed to exist?
3. Can we have more than one key?

NULL and NOT NULL

- To say “don’t know the value” we use **NULL**
 - NULL has (sometimes painful) semantics, more detail later

Students(sid:string, name:string, gpa: float)

sid	name	gpa
123	Bob	3.9
143	Jim	NULL

Say, Jim just enrolled in his first class.

In SQL, we may constrain a column to be NOT NULL, e.g., “name” in this table

General Constraints

- We can actually specify arbitrary assertions
 - E.g. “*There cannot be 25 people in the DB class*”
- In practice, we don’t specify many such constraints. Why?
 - Performance!

Whenever we do something ugly (or avoid doing something convenient) it’s for the sake of performance

Summary of Schema Information

- Schema and Constraints are how databases understand the semantics (meaning) of data
- They are also useful for optimization
- SQL supports general constraints:
 - Keys and foreign keys are most important
 - We'll give you a chance to write the others

ACTIVITY: [Activity-2-1.ipynb](#)

2. Single-table queries

What you will learn about in this section

1. The SFW query
2. Other useful operators: LIKE, DISTINCT, ORDER BY
3. ACTIVITY: Single-table queries

SQL Query

- Basic form (there are many many more bells and whistles)

```
SELECT <attributes>
FROM   <one or more relations>
WHERE  <conditions>
```

Call this a SFW query.

Simple SQL Query: Selection

Selection is the operation of filtering a relation's tuples on some condition

PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

```
SELECT *
FROM Product
WHERE Category = 'Gadgets'
```



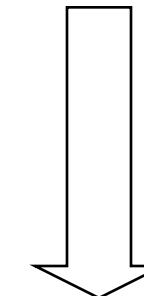
PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks

Simple SQL Query: Projection

Projection is the operation of producing an output table with tuples that have a subset of their prior attributes

PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

```
SELECT Pname, Price, Manufacturer  
FROM Product  
WHERE Category = 'Gadgets'
```



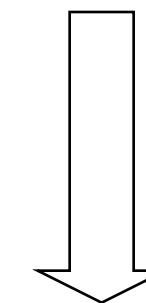
PName	Price	Manufacturer
Gizmo	\$19.99	GizmoWorks
Powergizmo	\$29.99	GizmoWorks

Notation

Input schema

Product(PName, Price, Category, Manufacturer)

```
SELECT Pname, Price, Manufacturer  
FROM Product  
WHERE Category = 'Gadgets'
```



Output schema

Answer(PName, Price, Manufacturer)

A Few Details

- SQL **commands** are case insensitive:
 - Same: SELECT, Select, select
 - Same: Product, product
- **Values** are **not**:
 - Different: ‘Seattle’, ‘seattle’
- Use single quotes for constants:
 - ‘abc’ - yes
 - “abc” - no

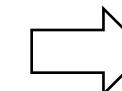
LIKE: Simple String Pattern Matching

```
SELECT *
FROM   Products
WHERE  PName  LIKE  '%gizmo%'
```

- $s \text{ } \text{LIKE} \text{ } p$: pattern matching on strings
- p may contain two special symbols:
 - $\%$ = any sequence of characters
 - $_$ = any single character

DISTINCT: Eliminating Duplicates

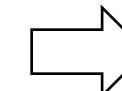
```
SELECT DISTINCT Category  
FROM Product
```



Category
Gadgets
Photography
Household

Versus

```
SELECT Category  
FROM Product
```



Category
Gadgets
Gadgets
Photography
Household

ORDER BY: Sorting the Results

```
SELECT      PName, Price, Manufacturer  
FROM        Product  
WHERE       Category='gizmo' AND Price > 50  
ORDER BY    Price, PName
```

Ties are broken by the second attribute on the ORDER BY list, etc.

Ordering is ascending, unless you specify the DESC keyword.

ACTIVITY: [Activity-2-2.ipynb](#)

3. Multi-table queries

What you will learn about in this section

1. Foreign key constraints
2. Joins: basics
3. Joins: SQL semantics
4. ACTIVITY: Multi-table queries

Foreign Key constraints

- Suppose we have the following schema:

`Students(sid: string, name: string, gpa: float)`

`Enrolled(student_id: string, cid: string, grade: string)`

- And we want to impose the following constraint:

- ‘Only bona fide students may enroll in courses’ i.e. a student must appear in the Students table to enroll in a class

Students

sid	name	gpa
101	Bob	3.2
123	Mary	3.8

Enrolled

student_id	cid	grade
123	564	A
123	537	A+

student_id alone is not a key- what is?

We say that student_id is a foreign key that refers to Students

Declaring Foreign Keys

```
Students(sid: string, name: string, gpa: float)
Enrolled(student_id: string, cid: string, grade: string)

CREATE TABLE Enrolled(
    student_id CHAR(20),
    cid          CHAR(20),
    grade        CHAR(10),
    PRIMARY KEY (student_id, cid),
    FOREIGN KEY (student_id) REFERENCES Students
)
```

Foreign Keys and update operations

`Students(sid: string, name: string, gpa: float)`

`Enrolled(student_id: string, cid: string, grade: string)`

- What if we insert a tuple into Enrolled, but no corresponding student?
 - INSERT is rejected (foreign keys are constraints)!
- What if we delete a student?
 - 1. Disallow the delete
 - 2. Remove all of the courses for that student
 - 3. SQL allows a third via *NULL* (*not yet covered*)

DBA chooses (syntax in the book)

Keys and Foreign Keys

Company

CName	StockPrice	Country
GizmoWorks	25	USA
Canon	65	Japan
Hitachi	15	Japan

What is a foreign key vs. a key here?

Product

PName	Price	Category	Manufacturer
Gizmo	\$19.99	Gadgets	GizmoWorks
Powergizmo	\$29.99	Gadgets	GizmoWorks
SingleTouch	\$149.99	Photography	Canon
MultiTouch	\$203.99	Household	Hitachi

Joins

Product(PName, Price, Category, Manufacturer)
Company(CName, StockPrice, Country)

Ex: Find all products under \$200 manufactured in Japan;
return their names and prices.

*Note: we will often omit
attribute types in schema
definitions for brevity, but
assume attributes are
always types*

```
SELECT PName, Price
FROM Product, Company
WHERE Manufacturer = CName
      AND Country='Japan'
      AND Price <= 200
```

Joins

```
Product(PName, Price, Category, Manufacturer)  
Company(CName, StockPrice, Country)
```

Ex: Find all products under \$200 manufactured in Japan;
return their names and prices.

```
SELECT PName, Price  
FROM Product, Company  
WHERE Manufacturer = CName  
      AND Country='Japan'  
      AND Price <= 200
```

A join between tables returns
all unique combinations of
their tuples which meet
some specified join condition

Joins

```
Product(PName, Price, Category, Manufacturer)  
Company(CName, StockPrice, Country)
```

Several equivalent ways to write a basic join in SQL:

```
SELECT PName, Price  
FROM Product, Company  
WHERE Manufacturer = CName  
      AND Country='Japan'  
      AND Price <= 200
```

```
SELECT PName, Price  
FROM Product  
JOIN Company ON Manufacturer = Cname  
              AND Country='Japan'  
WHERE Price <= 200
```

A few more later on...

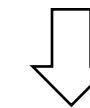
Joins

Product

PName	Price	Category	Manuf
Gizmo	\$19	Gadgets	GWorks
Powergizmo	\$29	Gadgets	GWorks
SingleTouch	\$149	Photography	Canon
MultiTouch	\$203	Household	Hitachi

Company

Cname	Stock	Country
GWorks	25	USA
Canon	65	Japan
Hitachi	15	Japan



```

SELECT PName, Price
FROM Product, Company
WHERE Manufacturer = CName
AND Country='Japan'
AND Price <= 200
    
```

PName	Price
SingleTouch	\$149.99

Tuple Variable Ambiguity in Multi-Table

```
Person(name, address, worksfor)  
Company(name, address)
```

```
SELECT DISTINCT name, address  
FROM Person, Company  
WHERE worksfor = name
```

Which “address” does this refer to?

Which “name”s??

Tuple Variable Ambiguity in Multi-Table

```
Person(name, address, worksfor)  
Company(name, address)
```

```
SELECT DISTINCT Person.name, Person.address  
FROM Person, Company  
WHERE Person.worksfor = Company.name
```

```
SELECT DISTINCT p.name, p.address  
FROM Person p, Company c  
WHERE p.worksfor = c.name
```

Both equivalent ways to resolve variable ambiguity

Meaning (Semantics) of SQL Queries

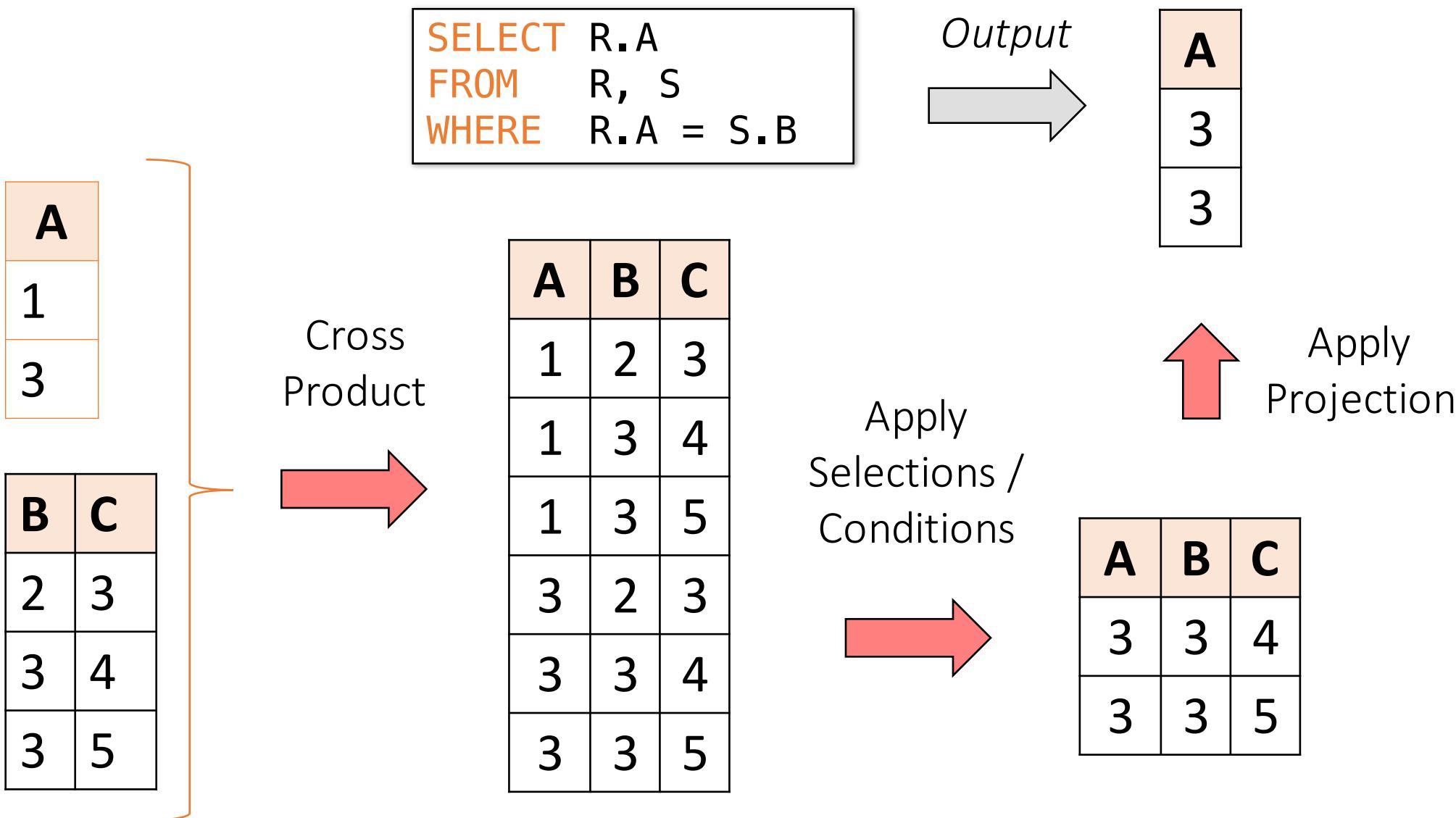
```
SELECT x1.a1, x1.a2, ..., xn.ak  
FROM R1 AS x1, R2 AS x2, ..., Rn AS xn  
WHERE Conditions(x1,..., xn)
```

Almost never the *fastest* way
to compute it!

```
Answer = {}  
for x1 in R1 do  
  for x2 in R2 do  
    ....  
    for xn in Rn do  
      if Conditions(x1,..., xn)  
        then Answer = Answer  $\cup$  {(x1.a1, x1.a2, ..., xn.ak)}  
return Answer
```

Note: this is a *multiset* union

An example of SQL semantics



Note the *semantics* of a join

```
SELECT R.A
FROM R, S
WHERE R.A = S.B
```

1. Take cross product:

$$X = R \times S$$

Recall: Cross product ($A \times B$) is the set of all unique tuples in A, B

Ex: $\{a,b,c\} \times \{1,2\}$
 $= \{(a,1), (a,2), (b,1), (b,2), (c,1), (c,2)\}$

2. Apply selections / conditions:

$$Y = \{(r, s) \in X \mid r.A == r.B\}$$

= Filtering!

3. Apply projections to get final output:

$$Z = (y.A,) \text{ for } y \in Y$$

= Returning only *some* attributes

Remembering this order is critical to understanding the output of certain queries (see later on...)

Note: we say “semantics” not “execution order”

- The preceding slides show *what a join means*
- Not actually how the DBMS executes it under the covers

A Subtlety about Joins

```
Product(PName, Price, Category, Manufacturer)  
Company(CName, StockPrice, Country)
```

Find all countries that manufacture some product
in the ‘Gadgets’ category.

```
SELECT Country  
FROM Product, Company  
WHERE Manufacturer=CName AND Category='Gadgets'
```

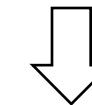
A subtlety about Joins

Product

PName	Price	Category	Manuf
Gizmo	\$19	Gadgets	GWorks
Powergizmo	\$29	Gadgets	GWorks
SingleTouch	\$149	Photography	Canon
MultiTouch	\$203	Household	Hitachi

Company

Cname	Stock	Country
GWorks	25	USA
Canon	65	Japan
Hitachi	15	Japan



```
SELECT Country
FROM Product, Company
WHERE Manufacturer=Cname
AND Category='Gadgets'
```

Country
?
?

What is the problem ?
What's the solution ?

ACTIVITY: [Lecture-2-3.ipynb](#)

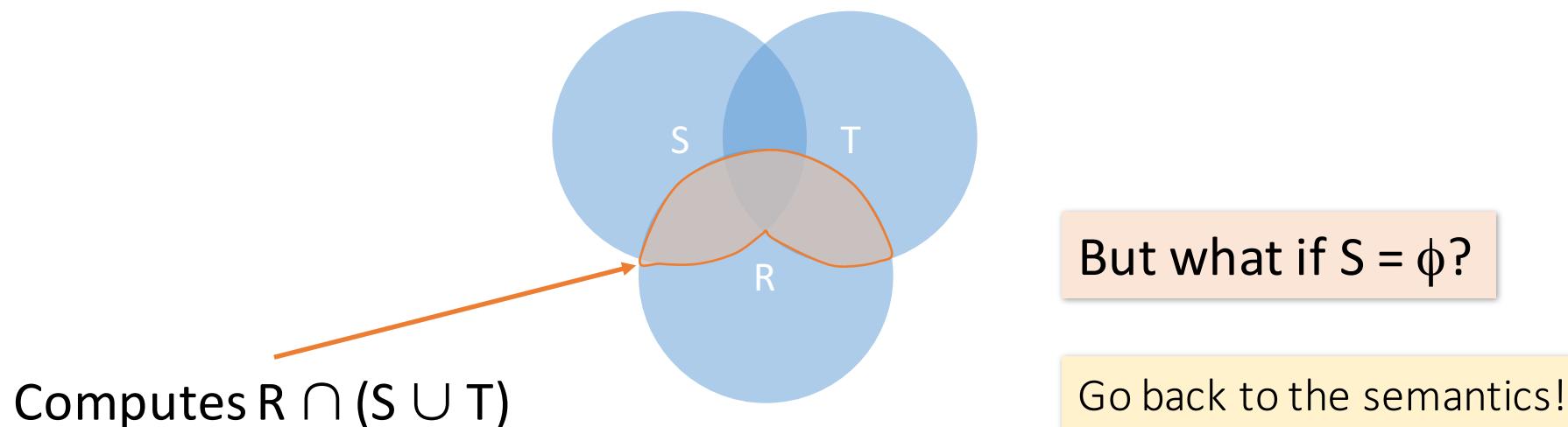
An Unintuitive Query

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```

What does it compute?

An Unintuitive Query

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```



An Unintuitive Query

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```

- Recall the semantics!
 1. Take cross-product
 2. Apply selections / conditions
 3. Apply projection
- If $S = \{\}$, then the cross product of $R, S, T = \{\}$, and the query result = $\{\}$!

Must consider semantics here.

Are there more explicit way to do set operations like this?

Lecture 3: SQL Part II

Today's Lecture

1. Set operators & nested queries
 - ACTIVITY: Set operator subtleties
2. Aggregation & GROUP BY
 - ACTIVITY: Fancy SQL Part I
3. Advanced SQL-izing
 - ACTIVITY: Fancy SQL Part II

1. Set Operators & Nested Queries

What you will learn about in this section

1. Multiset operators in SQL
2. Nested queries
3. ACTIVITY: Set operator subtleties

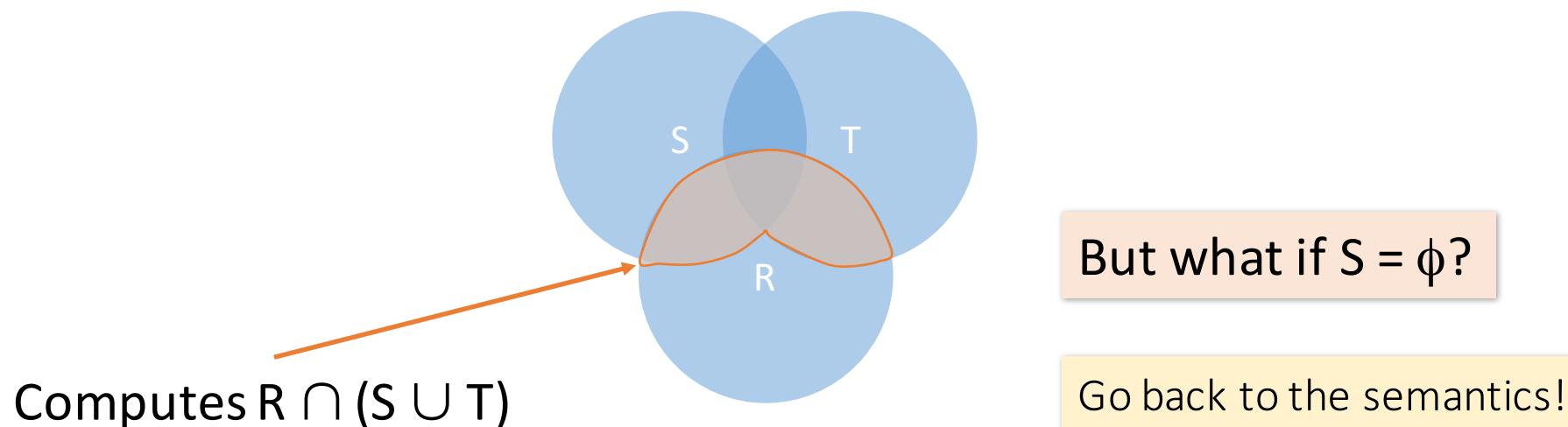
An Unintuitive Query

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```

What does it compute?

An Unintuitive Query

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```



An Unintuitive Query

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```

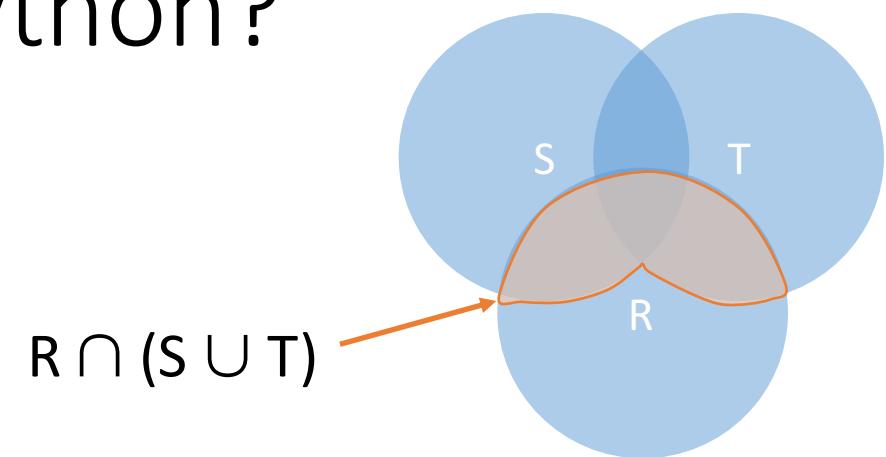
- Recall the semantics!
 1. Take cross-product
 2. Apply selections / conditions
 3. Apply projection
- If $S = \{\}$, then the cross product of $R, S, T = \{\}$, and the query result = $\{\}$!

Must consider semantics here.

Are there more explicit way to do set operations like this?

What does this look like in Python?

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```



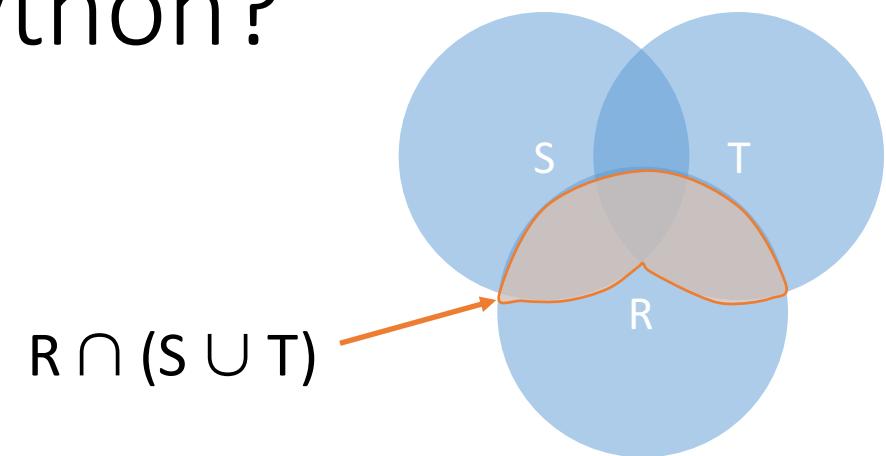
- Semantics:
 1. Take cross-product
 2. Apply selections / conditions
 3. Apply projection

Joins / cross-products are just nested for loops (in simplest implementation)!

If-then statements!

What does this look like in Python?

```
SELECT DISTINCT R.A  
FROM   R, S, T  
WHERE  R.A=S.A OR R.A=T.A
```



```
output = []  
  
for r in R:  
    for s in S:  
        for t in T:  
            if r['A'] == s['A'] or r['A'] == t['A']:  
                output.append(r['A'])  
return list(output)
```

Can you see now what happens if $S = []$?

See bonus activity on website!

Multiset Operations

Recall Multisets

Multiset X

Tuple
(1, a)
(1, a)
(1, b)
(2, c)
(2, c)
(2, c)
(1, d)
(1, d)



Equivalent
Representations
of a Multiset

$\lambda(X)$ = “Count of tuple in X”
(Items not listed have implicit count 0)

Multiset X

Tuple	$\lambda(X)$
(1, a)	2
(1, b)	1
(2, c)	3
(1, d)	2

Note: In a set all counts are {0,1}.

Generalizing Set Operations to Multiset Operations

Multiset X

Tuple	$\lambda(X)$
(1, a)	2
(1, b)	0
(2, c)	3
(1, d)	0

Multiset Y

Tuple	$\lambda(Y)$
(1, a)	5
(1, b)	1
(2, c)	2
(1, d)	2

 \cap $=$

Multiset Z

Tuple	$\lambda(Z)$
(1, a)	2
(1, b)	0
(2, c)	2
(1, d)	0

$$\lambda(Z) = \min(\lambda(X), \lambda(Y))$$

For sets, this is
intersection

Generalizing Set Operations to Multiset Operations

Multiset X

Tuple	$\lambda(X)$
(1, a)	2
(1, b)	0
(2, c)	3
(1, d)	0

 \cup

Multiset Y

Tuple	$\lambda(Y)$
(1, a)	5
(1, b)	1
(2, c)	2
(1, d)	2

 $=$

Multiset Z

Tuple	$\lambda(Z)$
(1, a)	7
(1, b)	1
(2, c)	5
(1, d)	2

$$\lambda(Z) = \lambda(X) + \lambda(Y)$$

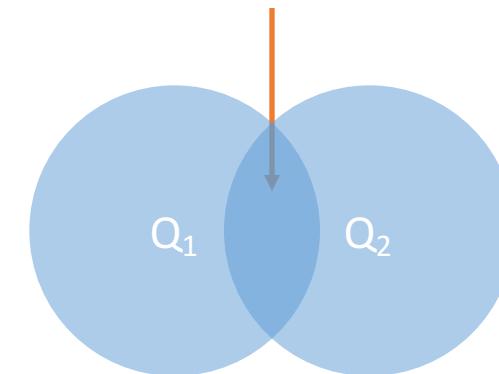
For sets,
this is **union**

Multiset Operations in SQL

Explicit Set Operators: INTERSECT

```
SELECT R.A  
FROM   R, S  
WHERE  R.A=S.A  
INTERSECT  
SELECT R.A  
FROM   R, T  
WHERE  R.A=T.A
```

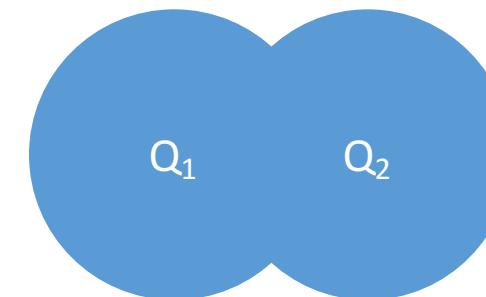
$$\{r.A \mid r.A = s.A\} \cap \{r.A \mid r.A = t.A\}$$



UNION

```
SELECT R.A  
FROM R, S  
WHERE R.A=S.A  
UNION  
SELECT R.A  
FROM R, T  
WHERE R.A=T.A
```

$$\{r.A \mid r.A = s.A\} \cup \{r.A \mid r.A = t.A\}$$



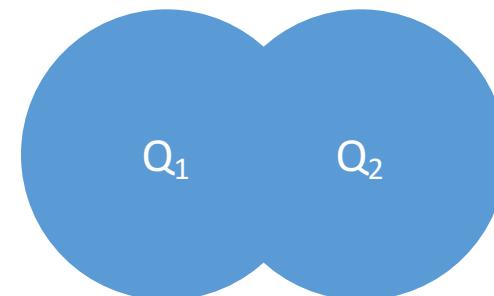
Why aren't there duplicates?

What if we want duplicates?

UNION ALL

```
SELECT R.A  
FROM R, S  
WHERE R.A=S.A  
UNION ALL  
SELECT R.A  
FROM R, T  
WHERE R.A=T.A
```

$$\{r.A \mid r.A = s.A\} \cup \{r.A \mid r.A = t.A\}$$

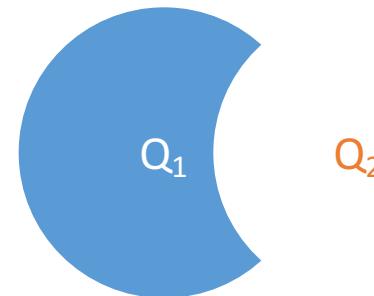


*ALL indicates
Multiset
operations*

EXCEPT

```
SELECT R.A  
FROM   R, S  
WHERE  R.A=S.A  
EXCEPT  
SELECT R.A  
FROM   R, T  
WHERE  R.A=T.A
```

$$\{r.A \mid r.A = s.A\} \setminus \{r.A \mid r.A = t.A\}$$



What is the multiset version?

INTERSECT: Still some subtle problems...

```
Company(name, hq_city)  
Product(pname, maker, factory_loc)
```

```
SELECT hq_city  
FROM Company, Product  
WHERE maker = name  
      AND factory_loc = 'US'  
INTERSECT  
SELECT hq_city  
FROM Company, Product  
WHERE maker = name  
      AND factory_loc = 'China'
```

“Headquarters of companies which make gizmos in US AND China”

What if two companies have HQ in US: BUT one has factory in China (but not US) and vice versa? **What goes wrong?**

INTERSECT: Remember the semantics!

```
Company(name, hq_city) AS C
Product(pname, maker,
factory_loc) AS P
```

```
SELECT hq_city
FROM Company, Product
WHERE maker = name
AND factory_loc='US'
```

INTERSECT

```
SELECT hq_city
FROM Company, Product
WHERE maker = name
AND factory_loc='China'
```

Example: $C \times P =$

C.name	C.hq_city	P.pname	P.maker	P.factory_loc
X Co.	U.S.	X	X Co.	U.S.
X Co.	U.S.	Y	Y Inc.	China
Y Inc.	U.S.	Y	Y Inc.	China
Y Inc.	U.S.	X	X Co.	U.S.

INTERSECT: Remember the semantics!

```
Company(name, hq_city) AS C
Product(pname, maker,
factory_loc) AS P
```

```
SELECT hq_city
FROM Company, Product
WHERE maker = name
AND factory_loc='US'
INTERSECT
SELECT hq_city
FROM Company, Product
WHERE maker = name
AND factory_loc='China'
```

Example: $C \times P =$

C.name	C.hq_city	P.pname	P.maker	P.factory_loc
X Co.	U.S.	X	X Co.	U.S.
X Co.	U.S.	Y	Y Inc.	China
Y Inc.	U.S.	Y	Y Inc.	China
Y Inc.	U.S.	X	X Co.	U.S.

Hq_city = U.S.

Hq_city = U.S.

INTERSECT: Remember the semantics!

```
Company(name, hq_city) AS C
Product(pname, maker,
factory_loc) AS P
```

```
SELECT hq_city
FROM Company, Product
WHERE maker = name
AND factory_loc='US'
INTERSECT
SELECT hq_city
FROM Company, Product
WHERE maker = name
AND factory_loc='China'
```

Example: $C \times P =$

C.name	C.hq_city	P.pname	P.maker	P.factory_loc
X Co.	U.S.	X	X Co.	U.S.
X Co.	U.S.	Y	Y Inc.	China
Y Inc.	U.S.	Y	Y Inc.	China
Y Inc.	U.S.	X	X Co.	U.S.

Hq_city = U.S.

\cap

Hq_city = U.S.



Hq_city = U.S.

We did the INTERSECT
on the wrong attributes!

One Solution: Nested Queries

```
Company(name, hq_city)  
Product(pname, maker, factory_loc)
```

```
SELECT DISTINCT hq_city  
FROM Company, Product  
WHERE maker = name  
AND name IN (  
    SELECT maker  
    FROM Product  
    WHERE factory_loc = 'US')  
AND name IN (  
    SELECT maker  
    FROM Product  
    WHERE factory_loc = 'China')
```

“Headquarters of companies which make gizmos in US AND China”

Note: If we hadn't used DISTINCT here, how many copies of each hq_city would have been returned?

High-level note on nested queries

- We can do nested queries because SQL is ***compositional***:
 - Everything (inputs / outputs) is represented as multisets- the output of one query can thus be used as the input to another (nesting)!
- This is extremely powerful!

Nested queries: Sub-queries Returning Relations

Another example:

```
Company(name, city)
Product(name, maker)
Purchase(id, product, buyer)
```

```
SELECT c.city
FROM Company c
WHERE c.name IN (
    SELECT pr.maker
        FROM Purchase p, Product pr
    WHERE p.product = pr.name
    AND p.buyer = 'Joe Blow')
```

“Cities where one can find companies that manufacture products bought by Joe Blow”

Nested Queries

Is this query equivalent?

```
SELECT c.city  
FROM Company c,  
      Product pr,  
      Purchase p  
WHERE c.name = pr.maker  
  AND pr.name = p.product  
  AND p.buyer = 'Joe Blow'
```

Beware of duplicates!

Nested Queries

```
SELECT DISTINCT c.city
FROM Company c,
      Product pr,
      Purchase p
WHERE c.name = pr.maker
  AND pr.name = p.product
  AND p.buyer = 'Joe Blow'
```

```
SELECT DISTINCT c.city
FROM Company c
WHERE c.name IN (
    SELECT pr.maker
    FROM Purchase p, Product pr
    WHERE p.product = pr.name
      AND p.buyer = 'Joe Blow')
```

Now they are equivalent

Subqueries Returning Relations

You can also use operations of the form:

- s > ALL R
- s < ANY R
- EXISTS R

ANY and ALL not supported by
SQLite.

Ex: **Product(name, price, category, maker)**

```
SELECT name
FROM Product
WHERE price > ALL(
    SELECT price
    FROM Product
    WHERE maker = 'Gizmo-Works')
```

Find products that
are more expensive
than all those
produced by
“Gizmo-Works”

Subqueries Returning Relations

You can also use operations of the form:

- $s > \text{ALL } R$
- $s < \text{ANY } R$
- EXISTS R

Ex: **Product(name, price, category, maker)**

```
SELECT p1.name
FROM Product p1
WHERE p1.maker = 'Gizmo-Works'
AND EXISTS(
    SELECT p2.name
    FROM Product p2
    WHERE p2.maker <> 'Gizmo-Works'
    AND p1.name = p2.name)
```

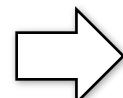
<> means \neq

Find ‘copycat’ products, i.e. products made by competitors with the same names as products made by “Gizmo-Works”

Nested queries as alternatives to INTERSECT and EXCEPT

INTERSECT and EXCEPT not in some DBMSs!

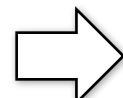
```
(SELECT R.A, R.B
FROM R)
INTERSECT
(SELECT S.A, S.B
FROM S)
```



```
SELECT R.A, R.B
FROM R
WHERE EXISTS(
    SELECT *
    FROM S
    WHERE R.A=S.A AND R.B=S.B)
```

If R, S have no duplicates, then can write without sub-queries (HOW?)

```
(SELECT R.A, R.B
FROM R)
EXCEPT
(SELECT S.A, S.B
FROM S)
```



```
SELECT R.A, R.B
FROM R
WHERE NOT EXISTS(
    SELECT *
    FROM S
    WHERE R.A=S.A AND R.B=S.B)
```

A question for Database Fans & Friends

- Can we express the previous nested queries as single SFW queries?
- Hint: show that all SFW queries are **monotone** (roughly: more tuples, more answers).
 - A query with **ALL** is not monotone

Correlated Queries

Movie(title, year, director, length)

```
SELECT DISTINCT title
FROM Movie AS m
WHERE year <> ANY(
    SELECT year
    FROM Movie
    WHERE title = m.title)
```

Find movies whose title appears more than once.

Note the scoping of the variables!

Note also: this can still be expressed as single SFW query...

Complex Correlated Query

Product(name, price, category, maker, year)

```
SELECT DISTINCT x.name, x.maker
FROM Product AS x
WHERE x.price > ALL(
    SELECT y.price
    FROM Product AS y
    WHERE x.maker = y.maker
    AND y.year < 1972)
```

Find products (and their manufacturers) that are more expensive than all products made by the same manufacturer before 1972

Can be very powerful (also much harder to optimize)

Basic SQL Summary

- SQL provides a high-level declarative language for manipulating data (DML)
- The workhorse is the SFW block
- Set operators are powerful but have some subtleties
- Powerful, nested queries also allowed.

[Activity-3-1.ipynb](#)

2. Aggregation & GROUP BY

What you will learn about in this section

1. Aggregation operators
2. GROUP BY
3. GROUP BY: with HAVING, semantics
4. ACTIVITY: Fancy SQL Pt. I

Aggregation

```
SELECT AVG(price)  
FROM Product  
WHERE maker = "Toyota"
```

```
SELECT COUNT(*)  
FROM Product  
WHERE year > 1995
```

- SQL supports several **aggregation** operations:
 - SUM, COUNT, MIN, MAX, AVG

Except COUNT, all aggregations apply to a single attribute

Aggregation: COUNT

- COUNT applies to duplicates, unless otherwise stated

```
SELECT COUNT(category)
FROM Product
WHERE year > 1995
```

Note: Same as COUNT().
Why?*

We probably want:

```
SELECT COUNT(DISTINCT category)
FROM Product
WHERE year > 1995
```

More Examples

```
Purchase(product, date, price, quantity)
```

```
SELECT SUM(price * quantity)  
FROM Purchase
```

What do these mean?

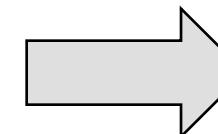
```
SELECT SUM(price * quantity)  
FROM Purchase  
WHERE product = 'bagel'
```

Simple Aggregations

Purchase

Product	Date	Price	Quantity
bagel	10/21	1	20
banana	10/3	0.5	10
banana	10/10	1	10
bagel	10/25	1.50	20

```
SELECT SUM(price * quantity)
FROM Purchase
WHERE product = 'bagel'
```



50 (= 1*20 + 1.50*20)

Grouping and Aggregation

Purchase(product, date, price, quantity)

```
SELECT      product,  
            SUM(price * quantity) AS TotalSales  
FROM        Purchase  
WHERE       date > '10/1/2005'  
GROUP BY    product
```

Find total sales
after 10/1/2005
per product.

Let's see what this means...

Grouping and Aggregation

Semantics of the query:

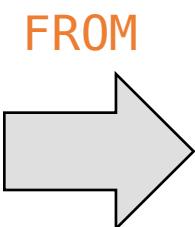
1. Compute the **FROM** and **WHERE** clauses

2. Group by the attributes in the **GROUP BY**

3. Compute the **SELECT** clause: grouped attributes and aggregates

1. Compute the **FROM** and **WHERE** clauses

```
SELECT product, SUM(price*quantity) AS TotalSales  
FROM Purchase  
WHERE date > '10/1/2005'  
GROUP BY product
```



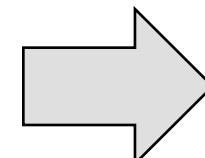
Product	Date	Price	Quantity
Bagel	10/21	1	20
Bagel	10/25	1.50	20
Banana	10/3	0.5	10
Banana	10/10	1	10

2. Group by the attributes in the GROUP BY

```
SELECT product, SUM(price*quantity) AS TotalSales  
FROM Purchase  
WHERE date > '10/1/2005'  
GROUP BY product
```

Product	Date	Price	Quantity
Bagel	10/21	1	20
Bagel	10/25	1.50	20
Banana	10/3	0.5	10
Banana	10/10	1	10

GROUP BY

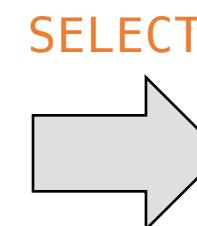


Product	Date	Price	Quantity
Bagel	10/21	1	20
	10/25	1.50	20
Banana	10/3	0.5	10
	10/10	1	10

3. Compute the **SELECT** clause: grouped attributes and aggregates

```
SELECT product, SUM(price*quantity) AS TotalSales  
FROM Purchase  
WHERE date > '10/1/2005'  
GROUP BY product
```

Product	Date	Price	Quantity
Bagel	10/21	1	20
	10/25	1.50	20
Banana	10/3	0.5	10
	10/10	1	10



Product	TotalSales
Bagel	50
Banana	15

GROUP BY v.s. Nested Quereis

```
SELECT      product, Sum(price*quantity) AS TotalSales  
FROM        Purchase  
WHERE       date > '10/1/2005'  
GROUP BY    product
```

```
SELECT DISTINCT x.product,  
              (SELECT Sum(y.price*y.quantity)  
               FROM Purchase y  
              WHERE x.product = y.product  
                    AND y.date > '10/1/2005') AS TotalSales  
FROM        Purchase x  
WHERE       x.date > '10/1/2005'
```

HAVING Clause

```
SELECT      product, SUM(price*quantity)
FROM        Purchase
WHERE       date > '10/1/2005'
GROUP BY    product
HAVING     SUM(quantity) > 100
```

Same query as before, except that we consider only products that have more than 100 buyers

HAVING clauses contains conditions on **aggregates**

Whereas WHERE clauses condition on *individual tuples*...

General form of Grouping and Aggregation

SELECT	S
FROM	R_1, \dots, R_n
WHERE	C_1
GROUP BY	a_1, \dots, a_k
HAVING	C_2

Why?

- S = Can ONLY contain attributes a_1, \dots, a_k and/or aggregates over other attributes
- C_1 = is any condition on the attributes in R_1, \dots, R_n
- C_2 = is any condition on the aggregate expressions

General form of Grouping and Aggregation

SELECT	S
FROM	R_1, \dots, R_n
WHERE	C_1
GROUP BY	a_1, \dots, a_k
HAVING	C_2

Evaluation steps:

1. Evaluate **FROM-WHERE**: apply condition C_1 on the attributes in R_1, \dots, R_n
2. **GROUP BY** the attributes a_1, \dots, a_k
3. **Apply condition C_2 to each group (may have aggregates)**
4. Compute aggregates in S and return the result

Group-by v.s. Nested Query

```
Author(login, name)  
Wrote(login, url)
```

- Find authors who wrote ≥ 10 documents:
- Attempt 1: with nested queries

```
SELECT DISTINCT Author.name  
FROM Author  
WHERE COUNT(  
    SELECT Wrote.url  
    FROM Wrote  
    WHERE Author.login = Wrote.login) > 10
```

This is
SQL by
a novice

Group-by v.s. Nested Query

- Find all authors who wrote at least 10 documents:
- Attempt 2: SQL style (with GROUP BY)

```
SELECT      Author.name  
FROM        Author, Wrote  
WHERE       Author.login = Wrote.login  
GROUP BY    Author.name  
HAVING     COUNT(Wrote.url) > 10
```

This is
SQL by
an expert

No need for **DISTINCT**: automatically from **GROUP BY**

Group-by vs. Nested Query

Which way is more efficient?

- Attempt #1- *With nested*: How many times do we do a SFW query over all of the Wrote relations?
- Attempt #2- *With group-by*: How about when written this way?

With GROUP BY can be much more efficient!

[Activity-3-2.ipynb](#)

3. Advanced SQL-izing

What you will learn about in this section

1. Quantifiers
2. NULLs
3. Outer Joins
4. ACTIVITY: Fancy SQL Pt. II

Quantifiers

Product(name, price, company)
Company(name, city)

```
SELECT DISTINCT Company.cname  
FROM Company, Product  
WHERE Company.name = Product.company  
AND Product.price < 100
```

An existential quantifier is a logical quantifier (roughly) of the form “there exists”

Existential: easy ! 😊

Find all companies that make some products with price < 100

Quantifiers

Product(name, price, company)
Company(name, city)

Find all companies
with products all
having price < 100

```
SELECT DISTINCT Company.cname
FROM Company
WHERE Company.name NOT IN(
    SELECT Product.company
    FROM Product
    WHERE Product.price >= 100)
```



Equivalent
Find all companies
that make only
products with price
< 100

A universal quantifier is of
the form “for all”

Universal: hard ! 😞

NULLS in SQL

- Whenever we don't have a value, we can put a NULL
- Can mean many things:
 - Value does not exists
 - Value exists but is unknown
 - Value not applicable
 - Etc.
- The schema specifies for each attribute if can be null (*nullable* attribute) or not
- How does SQL cope with tables that have NULLs?

Null Values

- *For numerical operations*, $\text{NULL} \rightarrow \text{NULL}$:
 - If $x = \text{NULL}$ then $4*(3-x)/7$ is still NULL
- *For boolean operations*, in SQL there are three values:

FALSE = 0

UNKNOWN = 0.5

TRUE = 1

- If $x = \text{NULL}$ then $x = \text{"Joe"}$ is UNKNOWN

Null Values

- $C1 \text{ AND } C2 = \min(C1, C2)$
- $C1 \text{ OR } C2 = \max(C1, C2)$
- $\text{NOT } C1 = 1 - C1$

```
SELECT *
FROM Person
WHERE (age < 25)
    AND (height > 6 AND weight > 190)
```

Won't return e.g.
(age=20
height=NULL
weight=200)!

Rule in SQL: include only tuples that yield TRUE / 1.0

Null Values

Unexpected behavior:

```
SELECT *
FROM Person
WHERE age < 25 OR age >= 25
```

Some Persons are not included !

Null Values

Can test for NULL explicitly:

- x IS NULL
- x IS NOT NULL

```
SELECT *
FROM Person
WHERE age < 25 OR age >= 25
OR age IS NULL
```

Now it includes all Persons!

RECAP: Inner Joins

By default, joins in SQL are “**inner joins**”:

```
Product(name, category)  
Purchase(prodName, store)
```

```
SELECT Product.name, Purchase.store  
FROM Product  
JOIN Purchase ON Product.name = Purchase.prodName
```

```
SELECT Product.name, Purchase.store  
FROM Product, Purchase  
WHERE Product.name = Purchase.prodName
```

Both equivalent:
Both INNER JOINS!



Inner Joins + NULLS = Lost data?

By default, joins in SQL are “**inner joins**”:

```
Product(name, category)  
Purchase(prodName, store)
```

```
SELECT Product.name, Purchase.store  
FROM Product  
JOIN Purchase ON Product.name = Purchase.prodName
```

```
SELECT Product.name, Purchase.store  
FROM Product, Purchase  
WHERE Product.name = Purchase.prodName
```

However: Products that never sold (with no Purchase tuple) will be lost!

Outer Joins

- An **outer join** returns tuples from the joined relations that don't have a corresponding tuple in the other relations
 - I.e. If we join relations A and B on $a.X = b.X$, and there is an entry in A with $X=5$, but none in B with $X=5$...
 - A LEFT OUTER JOIN will return a tuple (a, NULL) !
- Left outer joins in SQL:

```
SELECT Product.name, Purchase.store  
FROM Product  
LEFT OUTER JOIN Purchase ON  
Product.name = Purchase.prodName
```

Now we'll get products even if they didn't sell

INNER JOIN:

Product

name	category
Gizmo	gadget
Camera	Photo
OneClick	Photo

Purchase

prodName	store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

```
SELECT Product.name, Purchase.store  
FROM Product  
INNER JOIN Purchase  
ON Product.name = Purchase.prodName
```

Note: another equivalent way to write an
INNER JOIN!



name	store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

LEFT OUTER JOIN:

Product

name	category
Gizmo	gadget
Camera	Photo
OneClick	Photo

Purchase

prodName	store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz

```
SELECT Product.name, Purchase.store  
FROM Product  
LEFT OUTER JOIN Purchase  
ON Product.name = Purchase.prodName
```



name	store
Gizmo	Wiz
Camera	Ritz
Camera	Wiz
OneClick	NULL

Other Outer Joins

- Left outer join:
 - Include the left tuple even if there's no match
- Right outer join:
 - Include the right tuple even if there's no match
- Full outer join:
 - Include the both left and right tuples even if there's no match

[Activity-3-3.ipynb](#)

Summary

SQL is a rich programming language
that handles the way data is processed
declaratively