PS1 Recap Slides

Stanford CS145 Fall 2015

PS1: What you learned

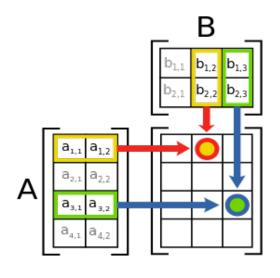
• This was a tough problem set- congratulations on doing so well!

- You used a declarative programming language (SQL) to
 - do linear algebra
 - answer *questions* about data
 - do *graph* operations
 - Cool stuff! However the point is not these <u>specific</u> applications...
- Less tricky versions of these same types of queries will be fair game for exams

Linear Algebra, Declaratively

• Matrix multiplication & other operations = just **joins**!

• The shift from procedural to declarative programming



$$C_{ij} = \sum_{k=1}^{m} A_{ik} B_{kj}$$

```
C = [[0]*p for i in range(n)]
for i in range(n):
   for j in range(p):
     for k in range(m):
        C[i][j] += A[i][k] * B[k][j]
```



```
SELECT A.i, B.j, SUM(A.x * B.x)
FROM A, B
WHERE A.j = B.i
GROUP BY A.i, B.j;
```

Declare a desired output set

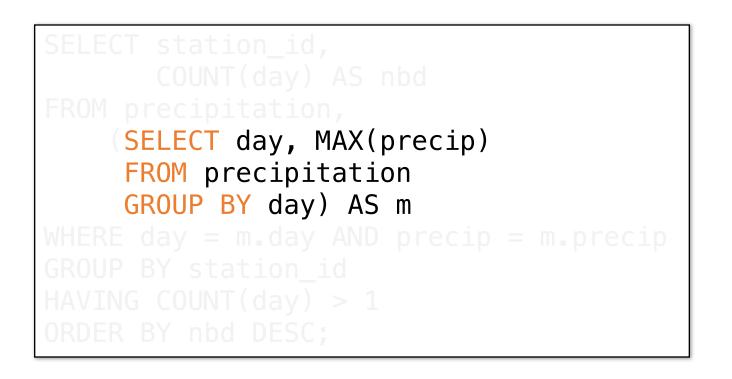
Proceed through a series of instructions

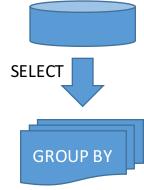
GROUP BY / HAVING + Aggregators + Nested queries

Think about order*!

*of the semantics, not the actual execution

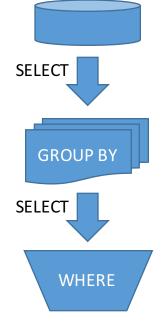
GROUP BY / HAVING + Aggregators + Nested queries





Get the max precipitation by day

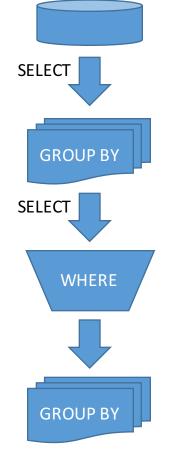
GROUP BY / HAVING + Aggregators + Nested queries



Get the max precipitation by day

Get the station, day pairs where / when this happened

GROUP BY / HAVING + Aggregators + Nested queries

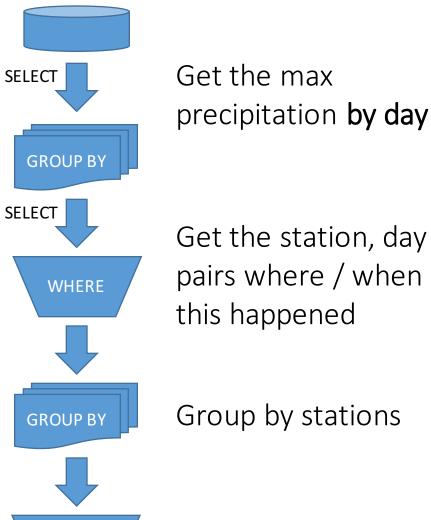


Get the max precipitation by day

Get the station, day pairs where / when this happened

Group by stations

GROUP BY / HAVING + Aggregators + Nested queries



HAVING

Having > 1 such day

Complex correlated queries

```
SELECT x1.p AS median
FROM x AS x1
WHERE

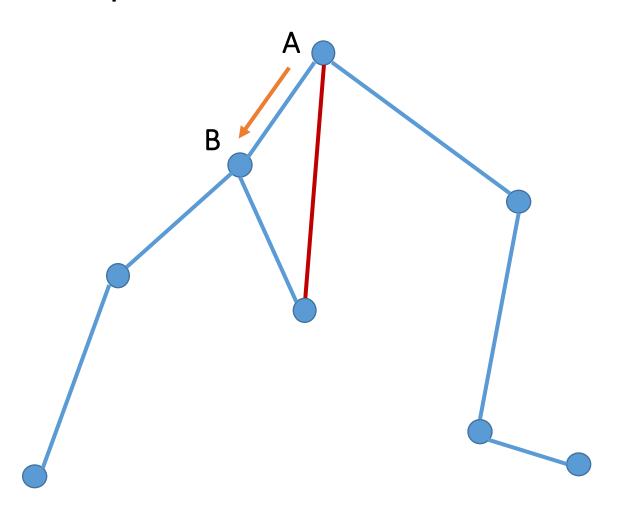
    (SELECT COUNT(*)
    FROM X AS x2
    WHERE x2.p > x1.p)
=
    (SELECT COUNT(*)
    FROM X AS x2
    WHERE x2.p < x1.p);</pre>
```

This was a tricky problem- but good practice in thinking about things declaratively

Nesting + EXISTS / ANY / ALL

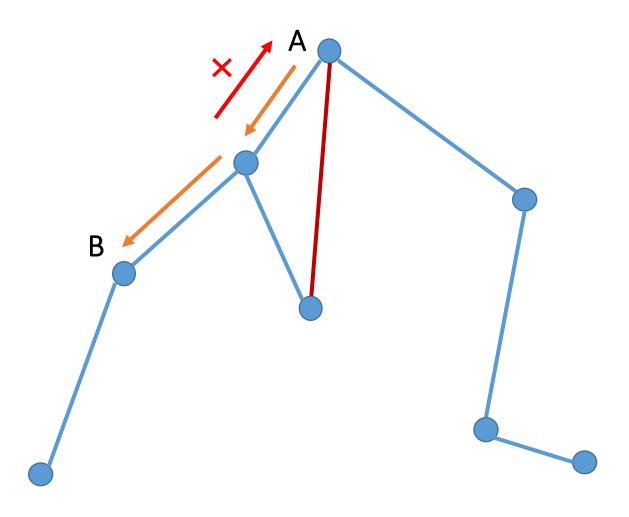
```
SELECT sid, p3.precip
FROM
      SELECT sid, precip
      FROM precipitation AS p1
      WHERE precip > 0 AND NOT EXISTS (
            SELECT p2.precip
            FROM precipitation AS p2
            WHERE p2.sid = p1.sid
              AND p2_precip > 0
              AND p2_precip < p1_precip)) AS p3
WHERE NOT EXISTS
      SELECT p4 precip
      FROM precipitation AS p4
      WHERE p4_precip -400 > p3_precip);
```

More complex, but again just think about order!



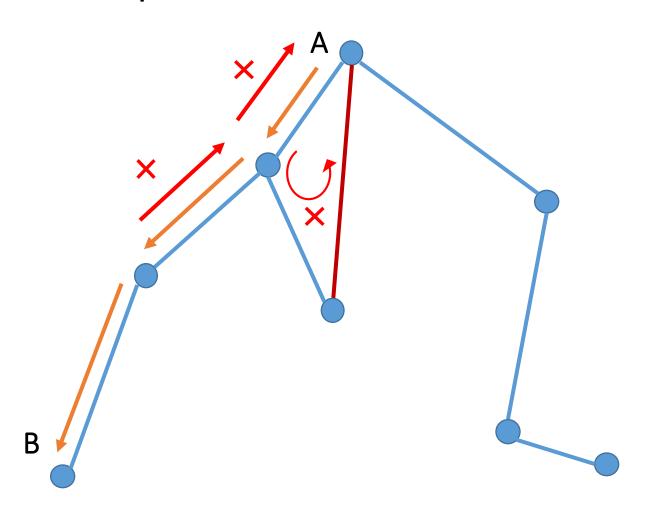
For fixed-length paths

SELECT A, B, d FROM edges



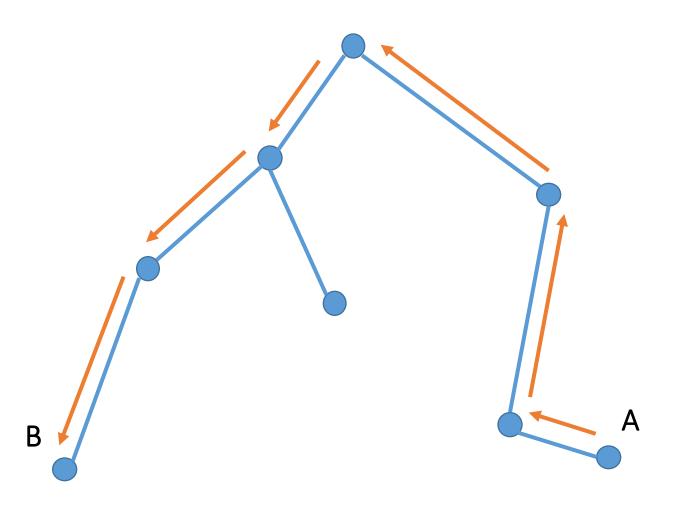
For fixed-length paths

```
SELECT A, B, d
FROM edges
UNION
SELECT e1.A, e2.B,
        e1.d + e2.d AS d
FROM edges e1, edges e2
WHERE e1.B = e2.A
AND e2.B <> e1.A
```



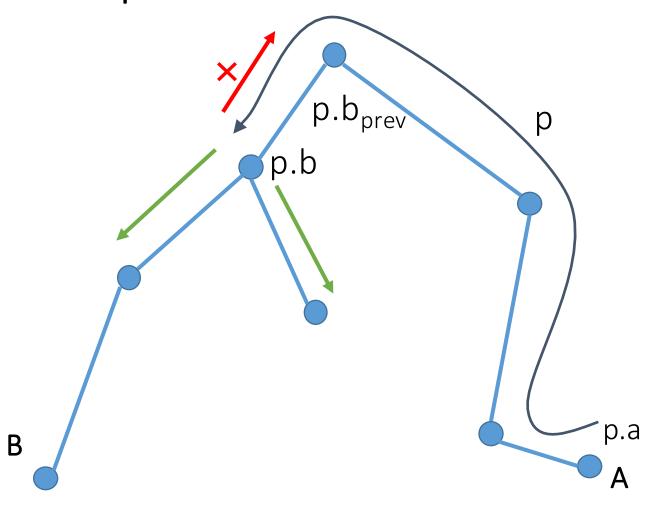
For fixed-length paths

```
SELECT A, B, d
FROM edges
UNION
SELECT e1.A, e2.B,
       e1.d + e2.d AS d
FROM edges e1, edges e2
WHERE e1.B = e2.A
  AND e2.B <> e1.A
UNTON
SELECT e1.A, e3.B,
       e1.d + e2.d + e3.d AS d
FROM edges e1, edges e2, edges e3
WHERE e1.B = e2.A
  AND e2.B = e3.A
  AND e2.B <> e1.A
  AND e3.B \ll e2.A
  AMD e3.B <> e1.A
```



For variable-length paths on trees

```
WITH RECURSIVE
paths(a, b, b_prev, d) AS (
      SELECT A, B, A
      FROM edges
      UNION
      SELECT p.a, e.B, e.A,
             p.d + e.d
      FROM paths p, edges e
      WHERE p.b = e.A
        AND s.B <> p.b_prev)
SELECT a, b, MAX(d)
FROM paths;
```



For variable-length paths on trees

```
WITH RECURSIVE
paths(a, b, b_prev, d) AS (
      SELECT A, B, A
      FROM edges
      UNION
      SELECT p.a, e.B, e.A,
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      FROM paths p, edges e
      WHERE p.b = e.A
        AND e_B <> p_b_prev)
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```