

HW #1

Mohit Joshi
msj696

3. msj696

4. a) a process sets "turn" to itself

Suppose that process 0 sets $wantCS[0] = true$ and then sets $turn = 0$ instead of 1

Then process 1 sets $wantCS[1] = true$ in another thread

Then process 0 evaluates $(wantCS[1] \&\& turn == 1)$ to false and enters its critical section

Then process 1 sets $turn = 1$ instead of 0, evaluates $(wantCS[0] \&\& turn == 0)$ to false and enters its critical section

Both processes are now in their critical sections at the same time, breaking mutual exclusion //

b) a process sets "turn" before setting "wantCS"

Suppose that process 0 sets $turn = 1$ instead of setting $wantCS$ first

Then process 1 sets $turn = 0$, sets $wantCS[1] = true$, evaluates $(wantCS[0] \&\& turn == 0)$ to false, and enters its critical section

Then process 0 sets $wantCS[0] = true$, evaluates $(wantCS[1] \&\& turn == 1)$ to false, and enters its critical section

Both processes are now in their critical sections at the same time, breaking mutual exclusion //

5. Change `int turn = 1` to `boolean turn[2] = {false, true}`

```
public void requestCS(int i) {  
    int j = 1 - i;  
    wantCS[i] = true;  
    boolean temp = turn[j];  
    turn[i] = temp;  
    while (wantCS[j] && temp == turn[j]);  
}
```

6. Consider the scenario where some processes p_0 and p_1 start choosing numbers concurrently while all other numbers are still 0

Then p_1 reaches the `number[i]++` line when p_0 finishes its second iteration of the previous line

In this scenario, their numbers will be the same when they do the check to enter their critical sections

Before p_0 increments its number, p_1 does the check and enters its critical section, then p_0 increments, checks, and enters its own critical section

Both processes are in their critical sections at the same time, breaking mutual exclusion. //