

## 第 3 次课堂作业

股票买卖:

解法:

```
def findbest(input):
    # best_couple = ""
    if len(input) <= 1:
        return 0, ""
    mid = len(input)//2
    left = input[:mid]
    right = input[mid:]

    # print(left, right)
    best_couple = ""

    best_left = findbest(left)[0]
    best_right = findbest(right)[0]
    best_cross = max(right) - min(left)
    if max(best_cross, best_right, best_left) == best_cross:
        best_couple = (str(max(right)) + '-' + str(min(left)))
        # print(best_couple)

    # if left == input[:len(input)//2]:
    #     print(max(right), '-', min(left))
    # print([max(best_cross, best_right, best_left), max(right), min(left)])
    # print([max(best_cross, best_right, best_left), max(right), min(left)][-1])
    # print(max(right), min(left))
    return max(best_cross, best_right, best_left), best_couple

A = [2, 5, 8, 18, 12, 9, 16]
print(findbest(A))
```

时间复杂度分析:

$\log n$  次递归

每次时间复杂度为  $n$

$O(n \log n)$

测试结果:

```
(16, '18-2')
[Finished in 0.1s]
```

Counting sort:

解法:

```
def counting_sort(a):  
    n=max(a)  
    b=[0]*n  
    for i in a:  
        b[i-1]=b[i-1]+1;  
  
    for i in range(n):  
        if b[i] != 0:  
            print(i+1)  
  
counting_sort([1,21,32,21,27])
```

时间复杂度分析:

两个并列的 for 循环

一个 for 循环遍历所有元素

一个 for 循环遍历整数范围 k 次

$O(n+k)$  (其中 k 是整数的范围)

测试结果:

```
1  
21  
27  
32  
[Finished in 0.1s]
```

Dijkstra 算法：  
解法：

```
import heapq
G = {1:{1:0, 2:1, 3:12},
     2:{2:0, 3:9, 4:3},
     3:{3:0, 5:5},
     4:{3:4, 4:0, 5:13, 6:15},
     5:{5:0, 6:4},
     6:{6:0}}

def dijkstra(G,start):    ###dijkstra算法
    INF = 999
    dis = dict((key,INF) for key in G)    # 初始化每个点的距离
    dis[start] = 0
    ###堆优化

    dis_un = {}    #存放未访问的点的距离
    pq = []    #存放排序后的值
    for node,d in dis.items():    #构造最小堆
        entry = [d,node]
        dis_un[node] = entry
        heapq.heappush(pq,entry)

    print("dis_un:",dis_un)
    print(pq)

    while len(pq)>0:
        v_dis,v = heapq.heappop(pq)    #未访问点中距离最小的点和对应的距离
        for node in G[v]:    #与v直接相连的点
            new_dis = dis[v] + G[v][node]
            if new_dis < dis[node]:    #如果与v直接相连的node通过v到src的距离小于dis中对应的node的值,则用小的值替换
                dis[node] = new_dis    #更新所有点的距离
                dis_un[node][0] = new_dis    #更新未访问的点到start的距离 改变dis_un中的entry从而改变pq
                print("_____")
                print("dis_un:",dis_un)
                print(pq)

    return dis
```

时间复杂度分析：  
遍历每个点：n  
用堆获取最小值：1  
堆跟新值：logn  
O（nlogn）

测试结果：

```
dis_un: {1: [0, 1], 2: [999, 2], 3: [999, 3], 4: [999, 4], 5: [999, 5], 6: [999, 6]}
[[0, 1], [999, 2], [999, 3], [999, 4], [999, 5], [999, 6]]

dis_un: {1: [0, 1], 2: [1, 2], 3: [999, 3], 4: [999, 4], 5: [999, 5], 6: [999, 6]}
[[1, 2], [999, 4], [999, 3], [999, 6], [999, 5]]

dis_un: {1: [0, 1], 2: [1, 2], 3: [12, 3], 4: [999, 4], 5: [999, 5], 6: [999, 6]}
[[1, 2], [999, 4], [12, 3], [999, 6], [999, 5]]

dis_un: {1: [0, 1], 2: [1, 2], 3: [10, 3], 4: [999, 4], 5: [999, 5], 6: [999, 6]}
[[10, 3], [999, 4], [999, 5], [999, 6]]

dis_un: {1: [0, 1], 2: [1, 2], 3: [10, 3], 4: [4, 4], 5: [999, 5], 6: [999, 6]}
[[10, 3], [4, 4], [999, 5], [999, 6]]

dis_un: {1: [0, 1], 2: [1, 2], 3: [10, 3], 4: [4, 4], 5: [15, 5], 6: [999, 6]}
[[4, 4], [999, 6], [15, 5]]

dis_un: {1: [0, 1], 2: [1, 2], 3: [8, 3], 4: [4, 4], 5: [15, 5], 6: [999, 6]}
[[15, 5], [999, 6]]

dis_un: {1: [0, 1], 2: [1, 2], 3: [8, 3], 4: [4, 4], 5: [15, 5], 6: [19, 6]}
[[15, 5], [19, 6]]
结果 {1: 0, 2: 1, 3: 8, 4: 4, 5: 15, 6: 19}
```