此章收录算法模板

高精度加法

高精度减法

```
vector<int> sub(vector<int> &A, vector<int> &B) // C = A - B, 满足A >= B, A >=
{
    vector<int> C;
    for (int i = 0, t = 0; i < A.size(); i ++ )
    {
        t = A[i] - t;
        if (i < B.size()) t -= B[i];
        C.push_back((t + 10) % 10);
        if (t < 0) t = 1;
        else t = 0;
    }
    while (C.size() > 1 && C.back() == 0) C.pop_back();
    return C;
}
```

高精度乘低精度

```
vector<int> mul(vector<int> &A, int b) // C = A * b, A >= 0, b >= 0
{
   vector<int> C;

   int t = 0;
   for (int i = 0; i < A.size() || t; i ++ )
   {
      if (i < A.size()) t += A[i] * b;
        C.push_back(t % 10);
      t /= 10;
   }</pre>
```

```
while (C.size() > 1 && C.back() == 0) C.pop_back();
return C;
```

高精度除以低精度

```
vector<int> div(vector<int> &A, int b, int &r) // A / b = C ... r, A >= 0, b >
{
    vector<int> C;
    r = 0;
    for (int i = A.size() - 1; i >= 0; i -- )
    {
        r = r * 10 + A[i];
        C.push_back(r / b);
        r %= b;
    }
    reverse(C.begin(), C.end());
    while (C.size() > 1 && C.back() == 0) C.pop_back();
    return C;
}
```

lowbit运算

```
int lowbit(int x) // 返回末尾的1
{
    return x & -x;
}
```

马拉车算法

```
void init() // a[]为原串, b[]为插入'#'后的新串
{
   int k = 0;
   b[k ++ ] = '
```

```
, b[k ++ ] = '#';
    for (int i = 0; i < n; i ++ ) b[k ++ ] = a[i], b[k ++ ] = '#';
    b[k ++ ] = '^';
    n = k;
}

void manacher() // 马拉车算法, b[]为插入'#'后的新串
{
    int mr = 0, mid;
    for (int i = 1; i < n; i ++ )
    {
        if (i < mr) p[i] = min(p[mid * 2 - i], mr - i);
        else p[i] = 1;</pre>
```

归并排序

```
void merge_sort(int q[], int l, int r) // 归并排序
{
    if (l >= r) return;
    int mid = l + r >> 1;
    merge_sort(q, l, mid);
    merge_sort(q, mid + 1, r);

    int k = 0, i = l, j = mid + 1;
    while (i <= mid && j <= r)
        if (q[i] <= q[j]) tmp[k ++ ] = q[i ++ ];
        else tmp[k ++ ] = q[j ++ ];

    while (i <= mid) tmp[k ++ ] = q[i ++ ];
    while (j <= r) tmp[k ++ ] = q[j ++ ];

    for (i = l, j = 0; i <= r; i ++, j ++ ) q[i] = tmp[j];
}</pre>
```

DLX重复覆盖

```
int l[N], r[N], u[N], d[N], col[N], row[N], s[N], idx;
int ans[N], top; // 记录选择了哪些行
bool st[M]; // N为节点数, M为列数
void init() // 初始化十字链表
    for (int i = 0; i <= m; i ++ )
       l[i] = i - 1, r[i] = i + 1;
       u[i] = d[i] = i;
       s[i] = 0, col[i] = i;
   1[0] = m, r[m] = 0;
   idx = m + 1;
}
void add(int& hh, int& tt, int x, int y) // 在十字链表中插入节点
   row[idx] = x, col[idx] = y, s[y] ++ ;
   u[idx] = y, d[idx] = d[y], u[d[y]] = idx, d[y] = idx;
   r[hh] = l[tt] = idx, r[idx] = tt, l[idx] = hh;
   tt = idx ++;
}
int h() // IDA*的启发函数
{
```

```
int res = 0;
    memset(st, 0, sizeof st);
    for (int i = r[0]; i; i = r[i])
        if (st[col[i]]) continue;
        res ++ ;
        st[col[i]] = true;
        for (int j = d[i]; j != i; j = d[j])
            for (int k = r[j]; k != j; k = r[k])
               st[col[k]] = true;
    }
    return res;
void remove(int p)
    for (int i = d[p]; i != p; i = d[i])
        r[l[i]] = r[i];
        l[r[i]] = l[i];
    }
}
void resume(int p)
    for (int i = u[p]; i != p; i = u[i])
        r[l[i]] = i;
        l[r[i]] = i;
    }
}
bool dfs(int k)
    if (k + h() > top) return false;
    if (!r[0])
        top = k;
        return true;
    int p = r[0];
    for (int i = r[0]; i; i = r[i])
        if (s[i] < s[p])
           p = i;
    for (int i = d[p]; i != p; i = d[i])
    {
        ans[k] = row[i];
        remove(i);
        for (int j = r[i]; j != i; j = r[j]) remove(j);
        if (dfs(k + 1)) return true;
        for (int j = l[i]; j != i; j = l[j]) resume(j);
        resume(i);
   return false;
```

DLX精确覆盖

```
int l[N], r[N], u[N], d[N], col[N], row[N], s[N], idx;
```

```
int ans[N], top; // 记录选择了哪些行
void init() // 初始化十字链表
    for (int i = 0; i <= m; i ++ )
       l[i] = i - 1, r[i] = i + 1;
       u[i] = d[i] = i;
    1[0] = m, r[m] = 0;
   idx = m + 1;
}
void add(int& hh, int& tt, int x, int y) // 在十字链表中添加节点
    row[idx] = x, col[idx] = y, s[y] ++;
   u[idx] = y, d[idx] = d[y], u[d[y]] = idx, d[y] = idx;
   r[hh] = l[tt] = idx, r[idx] = tt, l[idx] = hh;
    tt = idx ++;
}
void remove(int p)
    r[l[p]] = r[p], l[r[p]] = l[p];
    for (int i = d[p]; i != p; i = d[i])
        for (int j = r[i]; j != i; j = r[j])
            s[col[j]] -- ;
            d[u[j]] = d[j], u[d[j]] = u[j];
        }
void resume(int p)
    for (int i = d[p]; i != p; i = d[i])
        for (int j = r[i]; j != i; j = r[j])
            s[col[j]] ++ ;
           d[u[j]] = j, u[d[j]] = j;
    r[l[p]] = p, l[r[p]] = p;
bool dfs()
    if (!r[0]) return true;
    int p = r[0];
    for (int i = r[0]; i; i = r[i])
       if (s[i] < s[p])
           p = i;
    if (!s[p]) return false;
    remove(p);
    for (int i = d[p]; i != p; i = d[i])
    {
        ans[ ++ top] = row[i];
        for (int j = r[i]; j != i; j = r[j]) remove(col[j]);
        if (dfs()) return true;
        for (int j = r[i]; j != i; j = r[j]) resume(col[j]);
        top -- ;
    }
```

```
resume(p);
return false;
```

并查集+路径压缩

```
int find(int x) // 并查集
{
    if (p[x] != x) p[x] = find(p[x]);
    return p[x];
}
```

字符串哈希

```
ULL get(int l, int r) // 计算子串 str[l ~ r] 的哈希值 {
    return h[r] - h[l - 1] * p[r - l + 1];
```

Trie插入

```
int son[N][26], cnt[N], idx;
void insert(char *str) // 插入字符串
   int p = 0;
   for (int i = 0; str[i]; i ++ )
       int u = str[i] - 'a';
       if (!son[p][u]) son[p][u] = ++ idx;
       p = son[p][u];
   cnt[p] ++ ;
}
int query(char *str) // 查询字符串出现次数
   int p = 0;
   for (int i = 0; str[i]; i ++ )
       int u = str[i] - 'a';
       if (!son[p][u]) return 0;
       p = son[p][u];
   return cnt[p];
```

邻接链表 (无权)

```
void add(int a,int b)
{
    e[idx] = b, next[idx] = h[a], h[a] = idx++;
}
```

邻接链表 (带权)

```
void add(int a,int b,int c)
{
    e[idx] = b, next[idx] = h[a], w[idx] = c, h[a] = idx++;
}
```

dijkstra算法

```
int dijkstra() // 求1号点到n号点的最短路距离,如果从1号点无法走到n号点则返回-1
   memset(dist, 0x3f, sizeof dist);
   dist[1] = 0;
   priority queue<PII, vector<PII>, greater<PII>> heap;
   heap.push(\{0, 1\});
   while (heap.size())
       auto t = heap.top();
       heap.pop();
       int ver = t.second, distance = t.first;
       if (st[ver]) continue;
       st[ver] = true;
       for (int i = h[ver]; i != -1; i = ne[i])
           int j = e[i];
           if (dist[j] > dist[ver] + w[i])
               dist[j] = dist[ver] + w[i];
               heap.push({dist[j], j});
           }
       }
    }
   if (dist[n] == 0x3f3f3f3f) return -1;
   return dist[n];
void dijkstra() // 求1号点到n号点的最短路距离
   memset(dist, 0x3f, sizeof dist);
   dist[1] = 0;
   priority queue<PII, vector<PII>, greater<PII>> heap;
   heap.push(\{0, 1\});
   while (heap.size())
       auto t = heap.top();
       heap.pop();
       int ver = t.second, distance = t.first;
       if (st[ver]) continue;
       st[ver] = true;
```

```
for (int i = h[ver]; i != -1; i = ne[i])
{
    int j = e[i];
    if (dist[j] > dist[ver] + w[i])
    {
        dist[j] = dist[ver] + w[i];
        heap.push({dist[j], j});
    }
}
}
```

匈牙利算法(NTR算法)

```
bool find(int x)
{
    for (int i = h[x]; i != -1; i = ne[i])
    {
        int j = e[i];
        if (!st[j])
        {
            st[j] = true;
            if (match[j] == 0 || find(match[j]))
            {
                match[j] = x;
                return true;
            }
        }
    }
    return false;
}
```

spfa算法 (最短路)

```
int spfa() // 求1号点到n号点的最短路距离,如果从1号点无法走到n号点则返回-1
{
   int hh = 0, tt = 0;
   memset(dist, 0x3f, sizeof dist);
   dist[1] = 0;
   q[tt ++] = 1;
   st[1] = true;
   while (hh != tt)
       int t = q[hh ++];
       if (hh == N) hh = 0;
       st[t] = false;
       for (int i = h[t]; i != -1; i = ne[i])
           int j = e[i];
           if (dist[j] > dist[t] + w[i])
              dist[j] = dist[t] + w[i];
                          // 如果队列中已存在;,则不需要将;重复插入
              if (!st[j])
               {
                  q[tt ++ ] = j;
```

```
if (tt == N) tt = 0;
    st[j] = true;
}

}

if (dist[n] == 0x3f3f3f3f) return -1;
return dist[n];
```

spfa算法 (判断负环)

```
bool spfa() // 如果存在负环,则返回true, 否则返回false。
   // 不需要初始化dist数组
   // 原理: 如果某条最短路径上有n个点(除了自己),那么加上自己之后一共有n+1个点,
   // 由抽屉原理一定有两个点相同,所以存在环。
   int hh = 0, tt = 0;
   for (int i = 1; i \le n; i ++ ) q[tt ++ ] = i, st[i] = true;
   while (hh != tt)
       int t = q[hh ++];
       if (hh == N) hh = 0;
       st[t] = false;
       for (int i = h[t]; \sim i; i = ne[i])
           int j = e[i];
           if (dist[j] > dist[t] + w[i])
               dist[j] = dist[t] + w[i];
               cnt[j] = cnt[t] + 1;
               if (cnt[j] >= n) return true;
               if (!st[j])
                  st[j] = true;
                  q[tt ++ ] = j;
                  if (tt == N) tt = 0;
               }
          }
      }
   return false;
```

括扑排序

```
void topsort()
{
   int hh = 0, tt = -1;

   // d[i] 存储点i的入度
   for (int i = 1; i <= n; i ++ )</pre>
```

欧拉函数

```
int phi(int x) // 欧拉函数
{
   int res = x;
   for (int i = 2; i <= x / i; i ++ )
      if (x % i == 0)
      {
        res = res / i * (i - 1);
        while (x % i == 0) x /= i;
      }
   if (x > 1) res = res / x * (x - 1);
   return res;
}
```

线性筛+欧拉函数

```
void get_eulers(int n)  // 线性筛法求1~n的欧拉函数
{
    euler[1] = 1;
    for (int i = 2; i <= n; i ++ )
    {
        if (!st[i])
        {
            primes[cnt ++ ] = i;
            euler[i] = i - 1;
        }
        for (int j = 0; primes[j] <= n / i; j ++ )
        {
            int t = primes[j] * i;
            st[t] = true;
            if (i % primes[j] == 0)
            {
                euler[t] = euler[i] * primes[j];
                break;
            }
            euler[t] = euler[i] * (primes[j] - 1);
        }
}</pre>
```

欧几里得算法

```
int gcd(int a, int b) // 欧几里得算法 {
    return b ? gcd(b, a % b) : a;
}
```

扩展欧几里得算法

```
int exgcd(int a, int b, int &x, int &y) // 扩展欧几里得算法, 求x, y, 使得ax + by =
{
    if (!b)
    {
        x = 1; y = 0;
        return a;
    }
    int d = exgcd(b, a % b, y, x);
    y -= (a / b) * x;
    return d;
}
```

判定质数

```
bool is_prime(int x) // 判定质数 {
    if (x < 2) return false;
    for (int i = 2; i <= x / i; i ++ )
        if (x % i == 0)
        return false;
    return true;
}
```

线性质数筛

```
void get_primes(int n) // 线性筛质数
{
    for (int i = 2; i <= n; i ++ )
    {
        if (!st[i]) primes[cnt ++ ] = i;
        for (int j = 0; primes[j] <= n / i; j ++ )
        {
            st[primes[j] * i] = true;
            if (i % primes[j] == 0) break;
        }
    }
}</pre>
```

快速幂算法

```
int quick_power(int a, int k, int p) // 求a^k mod p
{
   int res = 1 % p;
   while (k)
   {
```

```
if (k & 1) res = (LL)res * a % p;
    a = (LL)a * a % p;
    k >>= 1;
}
return res;
}
```

高斯消元 (浮点型)

```
int gauss() // 高斯消元, 答案存于a[i][n]中, 0 <= i < n
   int c, r;
   for (c = 0, r = 0; c < n; c ++)
       int t = r;
       for (int i = r; i < n; i ++ ) // 找绝对值最大的行
           if (fabs(a[i][c]) > fabs(a[t][c]))
       if (fabs(a[t][c]) < eps) continue;</pre>
       for (int i = c; i <= n; i ++ ) swap(a[t][i], a[r][i]); // 将绝对值最大的作
       for (int i = n; i >= c; i -- ) a[r][i] /= a[r][c]; // 将当前行的首位变成1
       for (int i = r + 1; i < n; i ++ ) // 用当前行将下面所有的列消成0
           if (fabs(a[i][c]) > eps)
               for (int j = n; j \ge c; j -- )
                   a[i][j] = a[r][j] * a[i][c];
      r ++ ;
   }
   if (r < n)
       for (int i = r; i < n; i ++)
           if (fabs(a[i][n]) > eps)
               return 2; // 无解
       return 1; // 有无穷多组解
   }
   for (int i = n - 1; i >= 0; i -- )
       for (int j = i + 1; j < n; j ++ )
           a[i][n] -= a[i][j] * a[j][n];
   return 0; // 有唯一解
```

高斯消元 (布尔型)

```
int gauss() // 高斯消元,答案存于a[i][n]中,0 <= i < n
{
   int c, r;
   for (c = 0, r = 0; c < n; c ++ )
   {
     int t = r;
     for (int i = r; i < n; i ++ ) // 找非零行
        if (a[i][c])
        t = i;</pre>
```

```
if (!a[t][c]) continue;
    for (int i = c; i <= n; i ++ ) swap(a[r][i], a[t][i]); // 将非零行换到最頂
    for (int i = r + 1; i < n; i ++ ) // 用当前行将下面所有的列消成0
        if (a[i][c])
           for (int j = n; j >= c; j -- )
               a[i][j] ^= a[r][j];
   r ++ ;
}
if (r < n)
    for (int i = r; i < n; i ++)
       if (a[i][n])
           return 2; // 无解
    return 1; // 有多组解
}
for (int i = n - 1; i >= 0; i -- )
    for (int j = i + 1; j < n; j ++ )
        a[i][n] ^= a[i][j] * a[j][n];
return 0; // 有唯一解
```

Lucas定理

```
int qmi(int a, int k, int p) // 快速幂模板
   int res = 1 \% p;
   while (k)
       if (k & 1) res = (LL)res * a % p;
       a = (LL)a * a % p;
       k >>= 1;
   }
   return res;
int C(int a, int b, int p) // 通过定理求组合数C(a, b)
   if (a < b) return 0;
   LL x = 1, y = 1; // x是分子, y是分母
   for (int i = a, j = 1; j \le b; i ---, j ++-)
       x = (LL)x * i % p;
       y = (LL) y * j % p;
   return x * (LL)qmi(y, p - 2, p) % p;
int lucas(LL a, LL b, int p)
   if (a < p && b < p) return C(a, b, p);
   return (LL)C(a % p, b % p, p) * lucas(a / p, b / p, p) % p;
```

LCIS (最长上升子序列)

```
#include <iostream>
#include <algorithm>
const int N = 3010;
int dp[N][N];
int n;
int number A[N], number B[N];
int main()
    scanf("%d", &n);
    for (int i = 1; i \le n; i++) scanf("%d", &number A[i]);
    for (int i = 1; i \le n; i++) scanf("%d", &number B[i]);
    for(int i = 1;i<=n;i++)
        int maxlen = 1;
        for (int j = 1; j \le n; j++)
            dp[i][j] = dp[i - 1][j];
            if (number A[i] == number B[j]) dp[i][j] = std::max(dp[i][j], maxle:
            if (number A[i] > number B[j]) maxlen = std::max(maxlen, dp[i-1][j]
    }
    int res = 0;
    for (int i = 1; i \le n; i++) res = std::max(res, dp[n][i]);
    std::cout << res;</pre>
```

LIS+最少LIS组合

```
#include <iostream>
#include <algorithm>
const int N = 1010;
int dp[N];
int h[N],q[N];
int main()
    int n = 0;
    while(std::cin>>h[n]) n++;
    int len = 0, res = 0;
    for (int i = 0; i < n; i++)
        //LIS
        dp[i] = 1;
        for(int j = 0; j < i; j++)
            if(h[j] >= h[i]) dp[i] = std::max(dp[i], dp[j] + 1);
        res = std::max(res,dp[i]);
        //最少LIS
```

```
int k = 0;
while(k<len&&q[k]<h[i]) k++;
if( k== len) q[len++] = h[i];
else q[k] = h[i];
//如此建立出的数组一定是一个单调递增的数组
}
printf("%d\n", res);
printf("%d\n", len);
```

最快斐波那契数列

```
#include <iostream>
const int MOD = 1000000007;
//模拟矩阵X乘
void matrix mul(int a[][2],int b[][2],int M[][2])
    int temp[][2] = { \{0,0\},\{0,0\} };
    for (int i = 0; i < 2; i++)
        for(int j = 0; j < 2; j++)
            for (int k = 0; k<2; k++)
                 long long x = temp[i][j] + static cast < long long > (a[i][k]) * b[i]
                 temp[i][j] = x \% MOD;
        }
    for (int i = 0; i < 2; i++)
        for (int j = 0; j < 2; j++)
            M[i][j] = temp[i][j];
}
int f(long long n)
    if (n < 2) return 1;
    int M O[2] = \{ 1,1 \};
    int A[][2] = { \{1,1\},\{1,0\} };
    //将答案初始化为单位矩阵I
    int res[][2] = { \{1,0\},\{0,1\} };
    long long k = n - 1;
    while(k)
        if (k & 1) matrix mul(res, A, res);
        matrix mul(A, A, A);
        k >>= 1;
    int M n[2] = \{ 0, 0 \};
    for (int i = 0; i < 2; i++)
        for (int j = 0; j < 2; j++)
            long long x = M_n[i] + static_cast < long long > (M_0[j]) * res[j][i];
            M n[i] = x % MOD;
    }
```

```
return M_n[0];
}

int main()
{
    //从下标为1开始计算
    long long n;
    scanf_s("%lld", &n);
    printf("%d", f(n-1));
}
```

求约数

```
#include <iostream>
#include <algorithm>
#include <vector>
using namespace std;
vector<int> get divisors(int x)
{
    vector<int> res;
    for (int i = 1; i \le x / i; i ++ )
        if (x % i == 0)
        {
            res.push back(i);
            if (i != x / i) res.push back(x / i);
    sort(res.begin(), res.end());
    return res;
int main()
    int n;
    cin >> n;
    while (n -- )
        int x;
        cin >> x;
        auto res = get divisors(x);
        for (auto x : res) cout << x << ' ';
        cout << endl;</pre>
   return 0;
```

KMP

next数组的一些性质供大家复习:

- 1. next数组记录的就是最长相等前后缀的长度,如果next[len-1]!=0/-1,则说明字符串有相同的前后缀。
- 2. 最长相等前后缀的长度为next[len-1] + 1,数组长度为len。

- 3. 如果len % (len (next[len 1] + 1)) == 0,则说明(数组长度-最长相等前后缀的长度)正好可以被数组的长度整除,该字符串中有重复的子字符串。
- 4. 数组长度减去最长相等前后缀的长度相当于第一个重复子字符串的长度,也就是一个重复周期的长度,如果这个周期可以被整除,则说明整个数组就是这个周期的循环。

```
int next[N];
//一般型
void getNext(){
    int j = 0;
    next[0] = 0;
    for (int i = 1; i < n; i++)
        while (j>0 \& p[j]!=p[i]) j = next[j-1];
        if(p[j] == p[i]) j++;
        next[i] = j;
   }
}
//右移型
void getNext(){
    int j = -1;
    next[0] = j;
    for (int i = 1; i < n; i++)
    {
        while (j \ge 0 \& p[j+1]! = p[i]) j = next[j+1];
        if(p[j+1]==p[i]) j++;
        next[i] = j;
    }
}
```