

# RECURSION, DICTIONARIES

(download slides and .py files and follow along!)

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6.0001 LECTURE 6

# QUIZ PREP

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- a paper and an online component
- open book/notes
- not open Internet, not open computer
- start printing out whatever you may want to bring

# LAST TIME

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- tuples - immutable
- lists - mutable
- aliasing, cloning
- mutability side effects

# TODAY

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- recursion – divide/decrease and conquer
- dictionaries – another mutable object type

# RECURSION

wikipedia.

Version. Recursion is the process of repeating items in a self-similar way.

# WHAT IS RECURSION?

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- Algorithmically: a way to design solutions to problems by **divide-and-conquer** or **decrease-and-conquer**
  - reduce a problem to simpler versions of the same problem
- Semantically: a programming technique where a **function calls itself**.
  - in programming, goal is to NOT have infinite recursion
    - must have **1 or more base cases** that are easy to solve
    - must solve the same problem on **some other input** with the goal of simplifying the larger problem input

# ITERATIVE ALGORITHMS SO FAR

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- looping constructs (`while` and `for` loops) lead to **iterative** algorithms
- can capture computation in a set of **state variables** that update on each iteration through loop

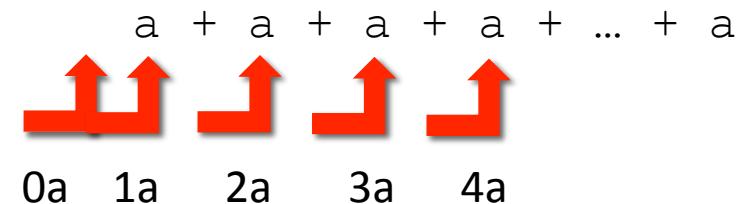
# MULTIPLICATION – ITERATIVE SOLUTION

- “multiply  $a * b$ ” is equivalent to “add  $a$  to itself  $b$  times”

- capture **state** by

- an **iteration** number ( $i$ ) starts at  $b$   
 $i \leftarrow i-1$  and stop when  $0$

- a current **value of computation** (result)  
 $\text{result} \leftarrow \text{result} + a$



```
def mult_iter(a, b):
    result = 0
    while b > 0:
        result += a
        b -= 1
    return result
```

iteration  
current value of computation,  
a running sum  
current value of iteration variable

# MULTIPLICATION – RECURSIVE SOLUTION

## ■ recursive step

- think how to reduce problem to a **simpler/ smaller version** of same problem

## ■ base case

- keep reducing problem until reach a simple case that can be **solved directly**
- when  $b = 1$ ,  $a^*b = a$

$$\begin{aligned} a^*b &= \underbrace{a + a + a + a + \dots + a}_{b \text{ times}} \\ &= a + \underbrace{a + a + a + a + \dots + a}_{b-1 \text{ times}} \\ &= a + \boxed{a * (b-1)} \end{aligned}$$

recursive reduction

```
def mult(a, b):
```

```
    if b == 1:  
        return a
```

```
    else:
```

```
        return a + mult(a, b-1)
```

# FACTORIAL

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$$n! = n * (n-1) * (n-2) * (n-3) * \dots * 1$$

Ex)  $4! = 4 \times 3 \times 2 \times 1$

$3! = 3 \times 2 \times 1$

- for what  $n$  do we know the factorial?

$n = 1 \rightarrow$  if  $n == 1:$   
                          return 1

*base case*

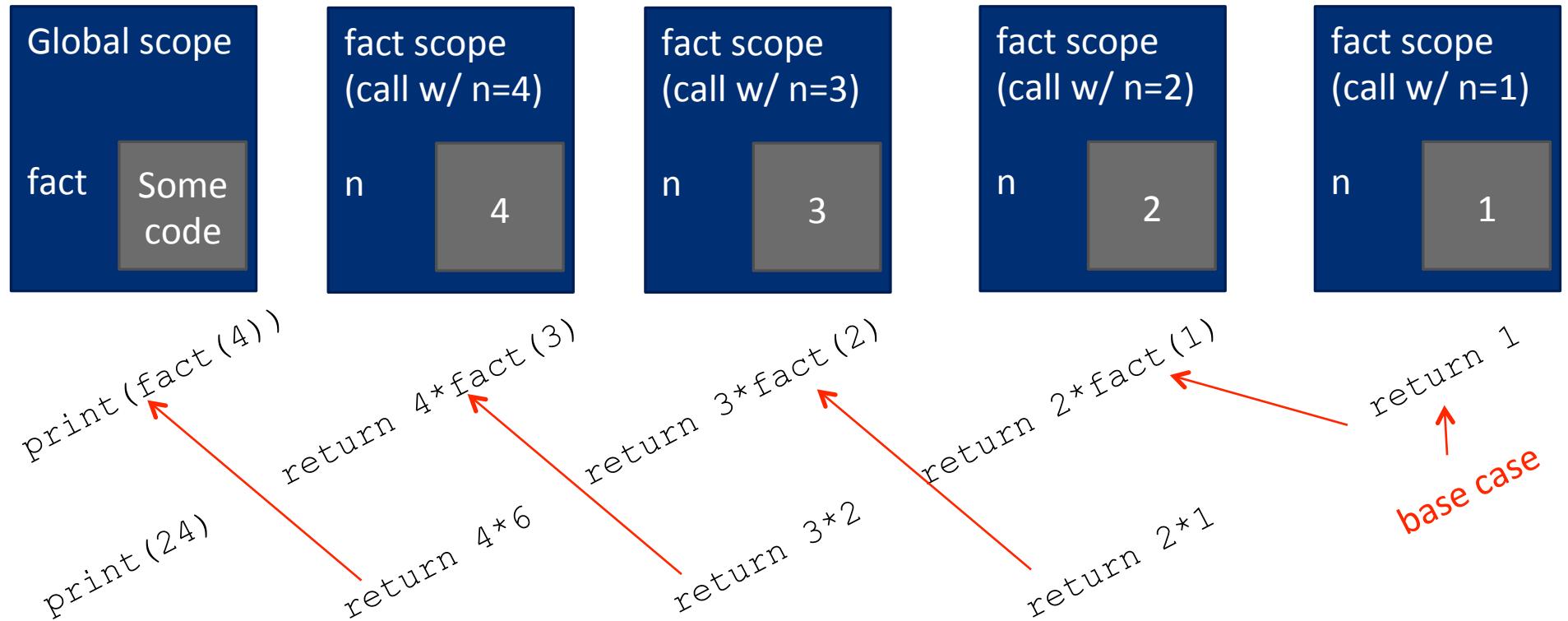
- how to reduce problem? Rewrite in terms of something simpler to reach base case

$n * (n-1)! \rightarrow$  else:  
                          return  $n * \text{factorial}(n-1)$

*recursive step*

# RECURSIVE FUNCTION SCOPE EXAMPLE

```
def fact(n):  
    if n == 1:  
        return 1  
    else:  
        return n*fact(n-1)  
  
print(fact(4))
```



# SOME OBSERVATIONS

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- each recursive call to a function creates its **own scope/environment**
- **bindings of variables** in a scope are not changed by recursive call
- flow of control passes back to **previous scope** once function call returns value

using the same variable  
names but they are different  
objects in separate scopes

# ITERATION vs. RECURSION

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```
def factorial_iter(n):      def factorial(n):  
    prod = 1                  if n == 1:  
    for i in range(1,n+1):     return 1  
        prod *= i              else:  
    return prod                  return n*factorial(n-1)
```

- recursion may be simpler, more intuitive
- recursion may be efficient from programmer POV
- recursion may not be efficient from computer POV

# INDUCTIVE REASONING

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- How do we know that our recursive code will work?
- `mult_iter` terminates because  $b$  is initially positive, and decreases by 1 each time around loop; thus must eventually become less than 1
- `mult` called with  $b = 1$  has no recursive call and stops
- `mult` called with  $b > 1$  makes a recursive call with a smaller version of  $b$ ; must eventually reach call with  $b = 1$

```
def mult_iter(a, b):  
    result = 0  
    while b > 0:  
        result += a  
        b -= 1  
    return result
```

```
def mult(a, b):  
    if b == 1:  
        return a  
    else:  
        return a + mult(a, b-1)
```

# MATHEMATICAL INDUCTION

---

- To prove a statement indexed on integers is true for all values of  $n$ :
  - Prove it is true when  $n$  is smallest value (e.g.  $n = 0$  or  $n = 1$ )
  - Then prove that if it is true for an arbitrary value of  $n$ , one can show that it must be true for  $n+1$

# EXAMPLE OF INDUCTION

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- $0 + 1 + 2 + 3 + \dots + n = (n(n+1))/2$
- Proof:
  - If  $n = 0$ , then LHS is 0 and RHS is  $0*1/2 = 0$ , so true
  - Assume true for some  $k$ , then need to show that
$$0 + 1 + 2 + \dots + k + (k+1) = ((k+1)(k+2))/2$$
    - LHS is  $k(k+1)/2 + (k+1)$  by assumption that property holds for problem of size  $k$
    - This becomes, by algebra,  $((k+1)(k+2))/2$
    - Hence expression holds for all  $n \geq 0$

# RELEVANCE TO CODE?

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- Same logic applies

```
def mult(a, b):  
    if b == 1:  
        return a  
  
    else:  
        return a + mult(a, b-1)
```

- Base case, we can show that `mult` must return correct answer
- For recursive case, we can assume that `mult` correctly returns an answer for problems of size smaller than  $b$ , then by the addition step, it must also return a correct answer for problem of size  $b$
- Thus by induction, code correctly returns answer

# TOWERS OF HANOI

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- The story:
  - 3 tall spikes
  - Stack of 64 different sized discs – start on one spike
  - Need to move stack to second spike (at which point universe ends)
  - Can only move one disc at a time, and a larger disc can never cover up a small disc

# TOWERS OF HANOI

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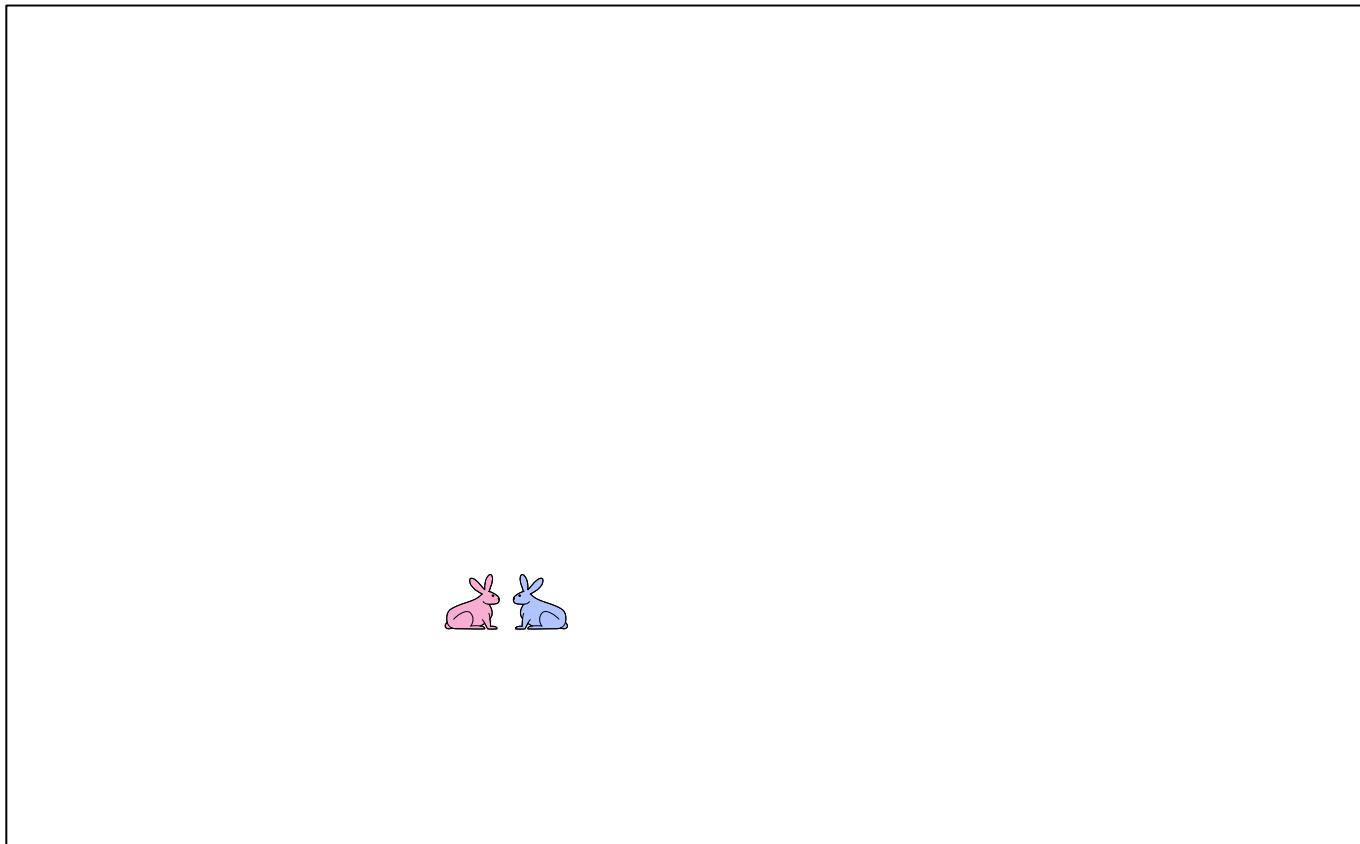
- Having seen a set of examples of different sized stacks, how would you write a program to print out the right set of moves?
- **Think recursively!**
  - Solve a smaller problem
  - Solve a basic problem
  - Solve a smaller problem

```
def printMove(fr, to):  
    print('move from ' + str(fr) + ' to ' + str(to))  
  
def Towers(n, fr, to, spare):  
    if n == 1:  
        printMove(fr, to)  
    else:  
        Towers(n-1, fr, spare, to)  
        Towers(1, fr, to, spare)  
        Towers(n-1, spare, to, fr)
```

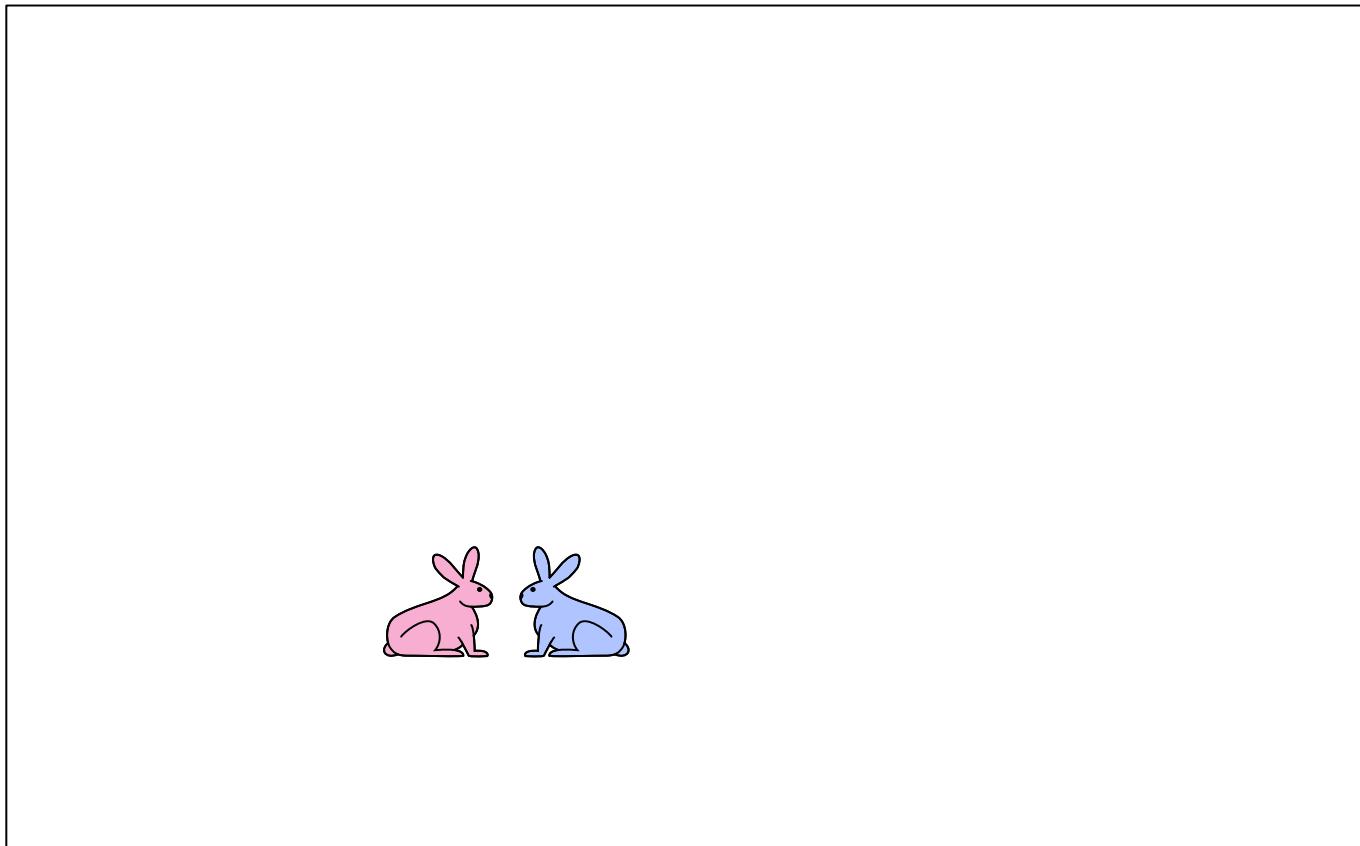
# RECURSION WITH MULTIPLE BASE CASES

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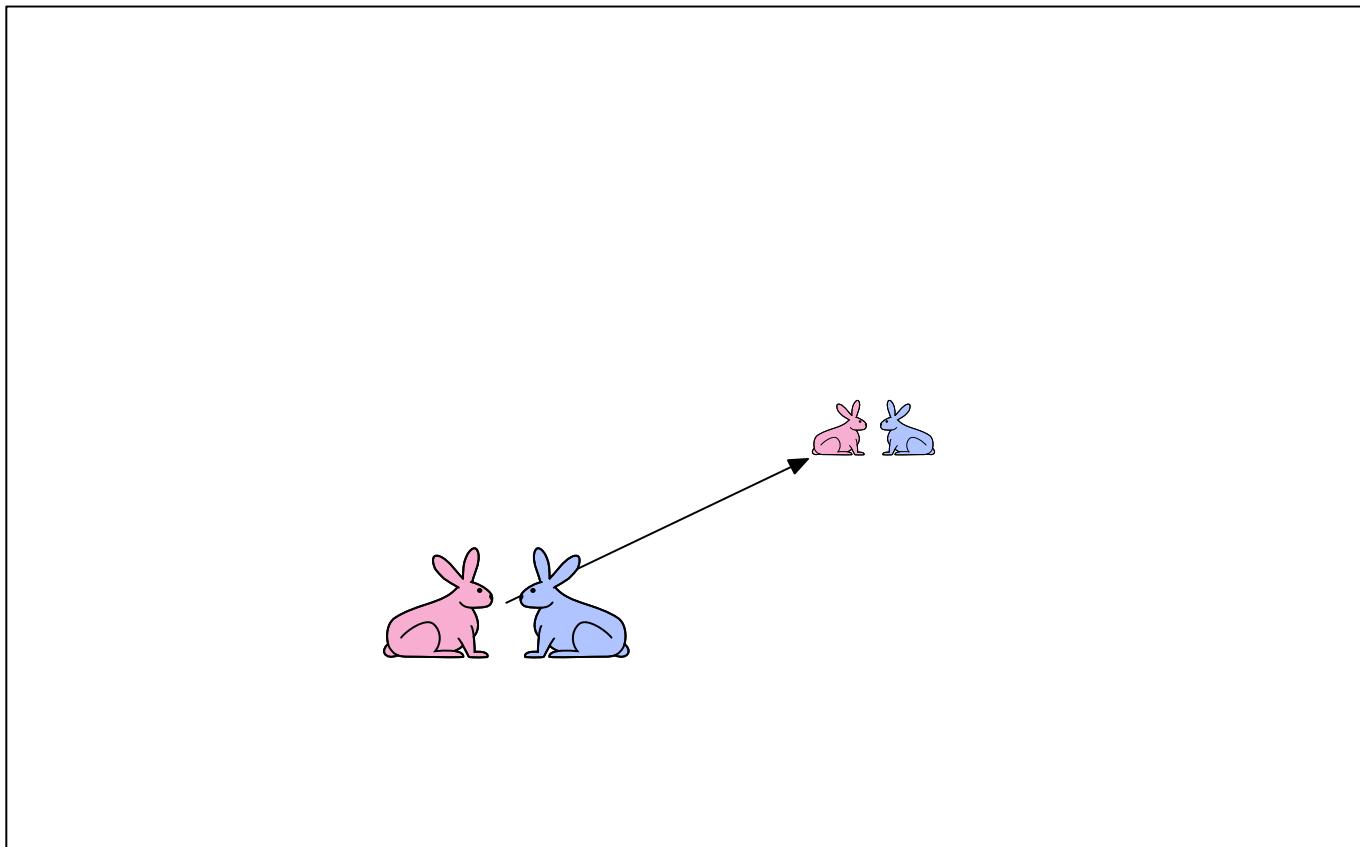
- Fibonacci numbers
  - Leonardo of Pisa (aka Fibonacci) modeled the following challenge
    - Newborn pair of rabbits (one female, one male) are put in a pen
    - Rabbits mate at age of one month
    - Rabbits have a one month gestation period
    - Assume rabbits never die, that female always produces one new pair (one male, one female) every month from its second month on.
    - How many female rabbits are there at the end of one year?



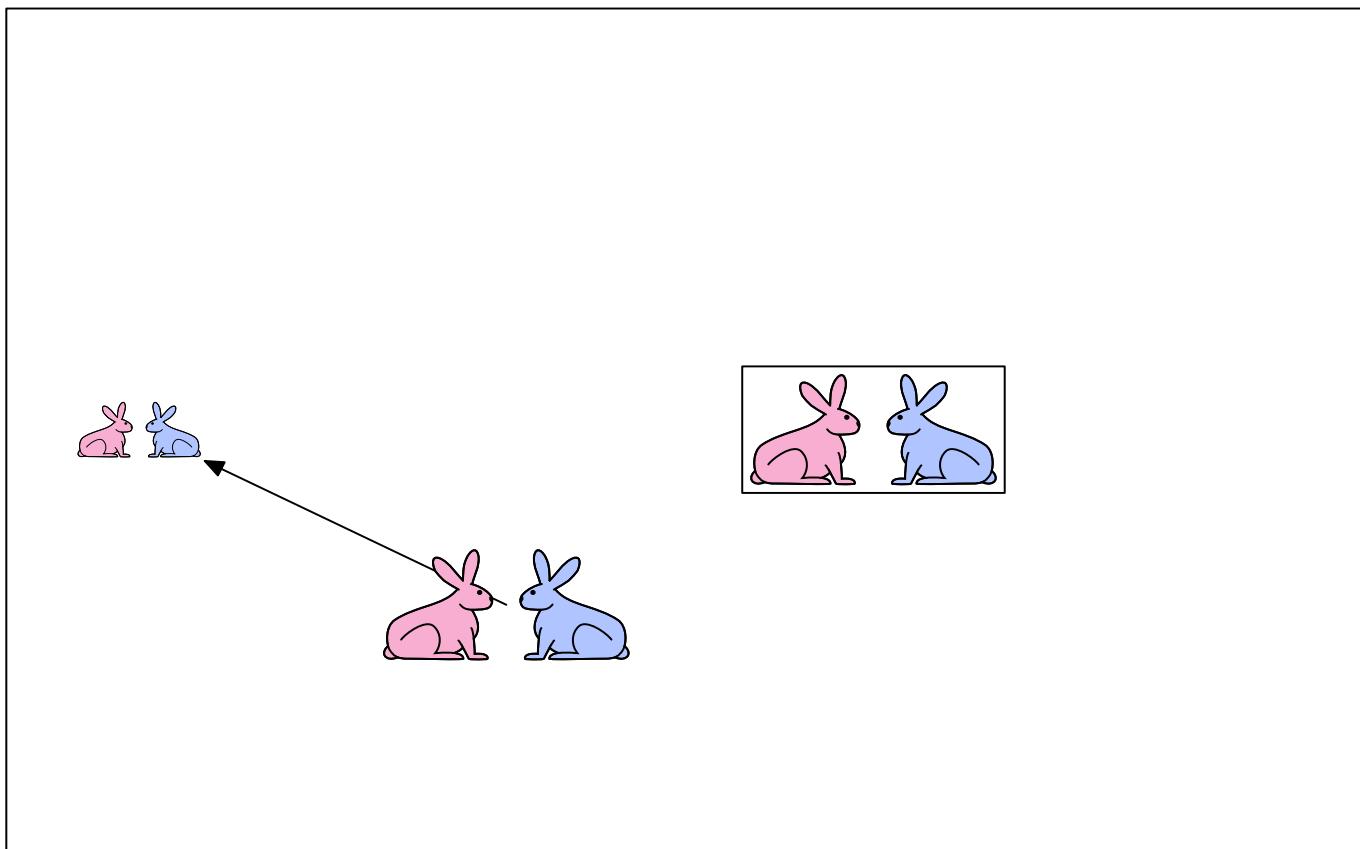
Demo courtesy of Prof. Denny Freeman and Adam Hartz



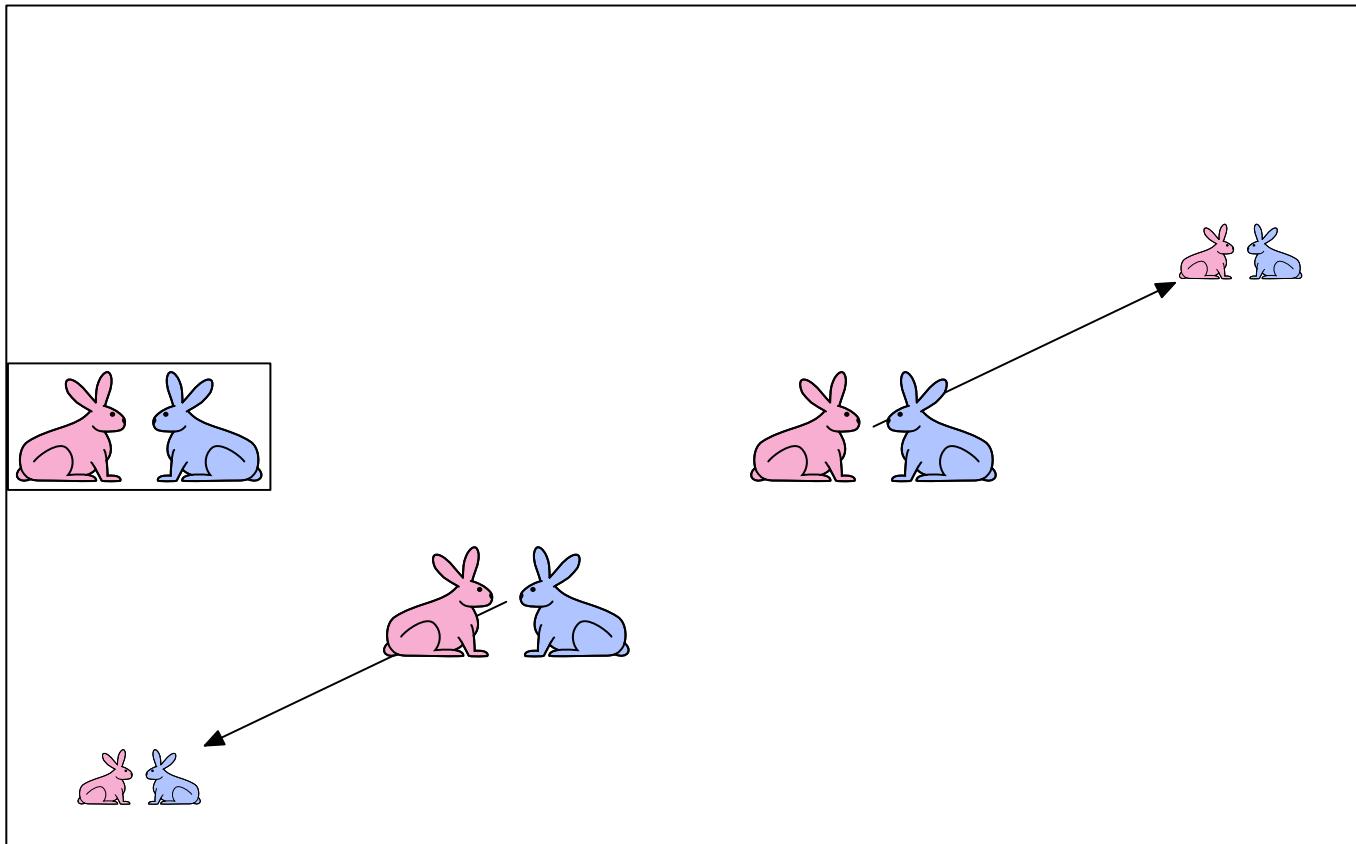
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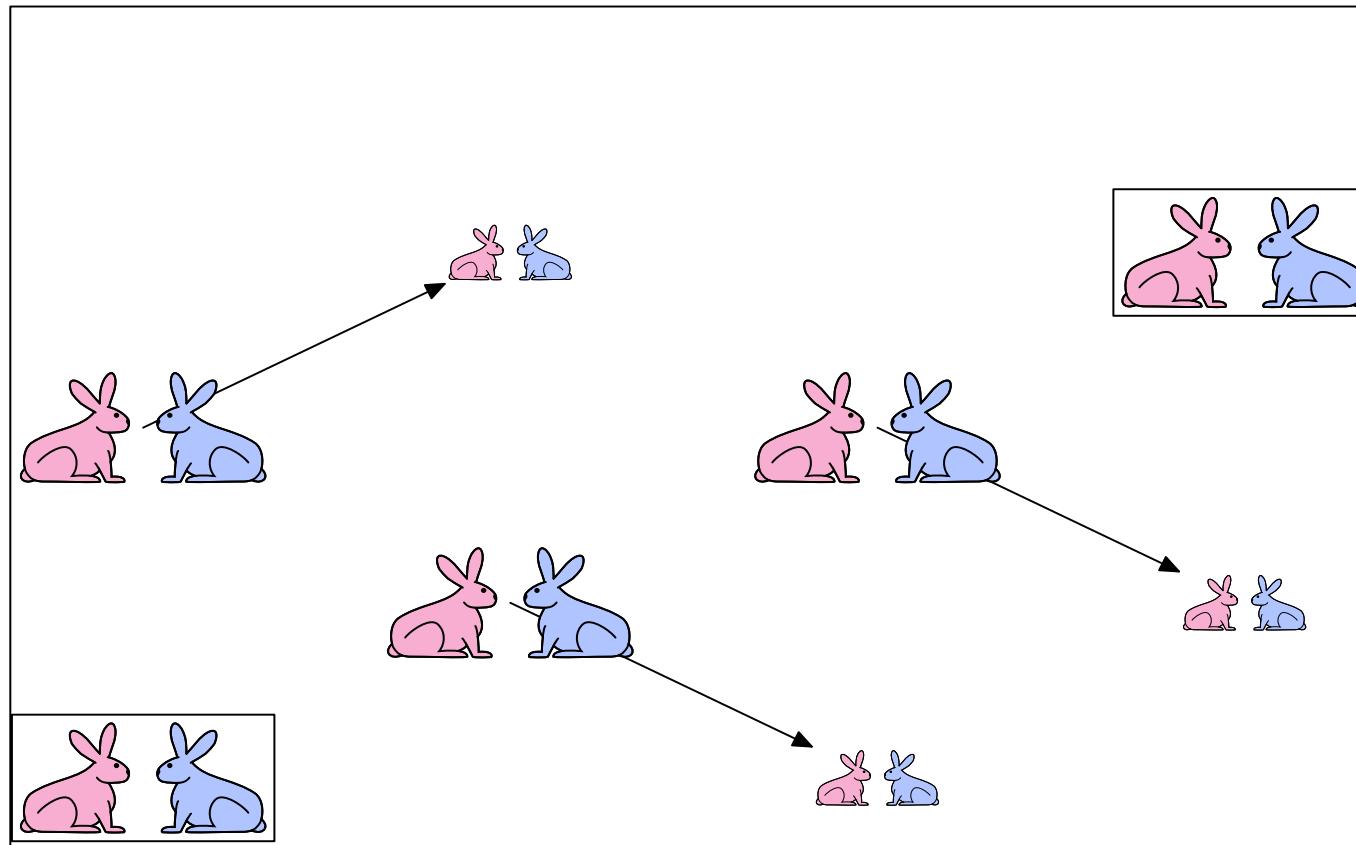
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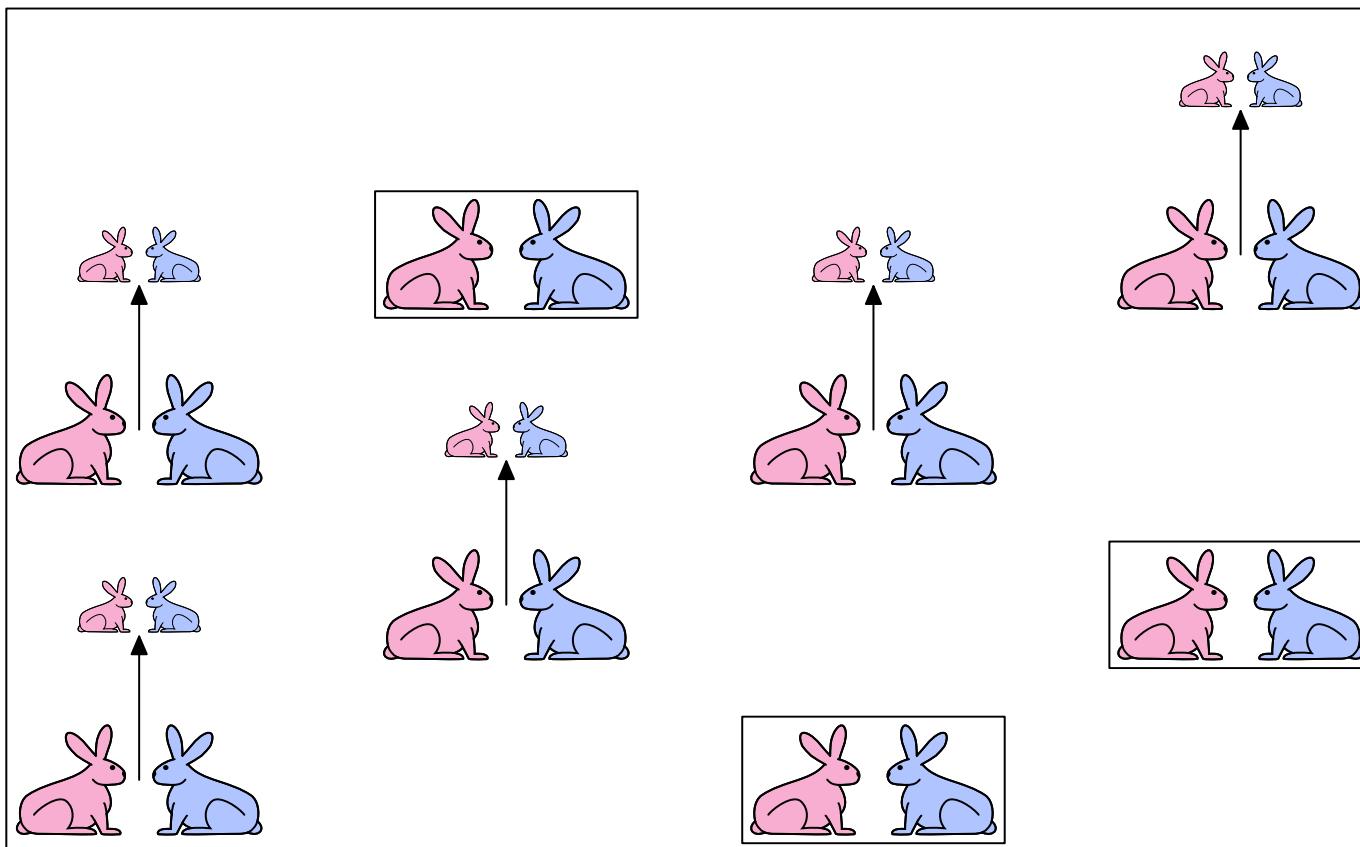
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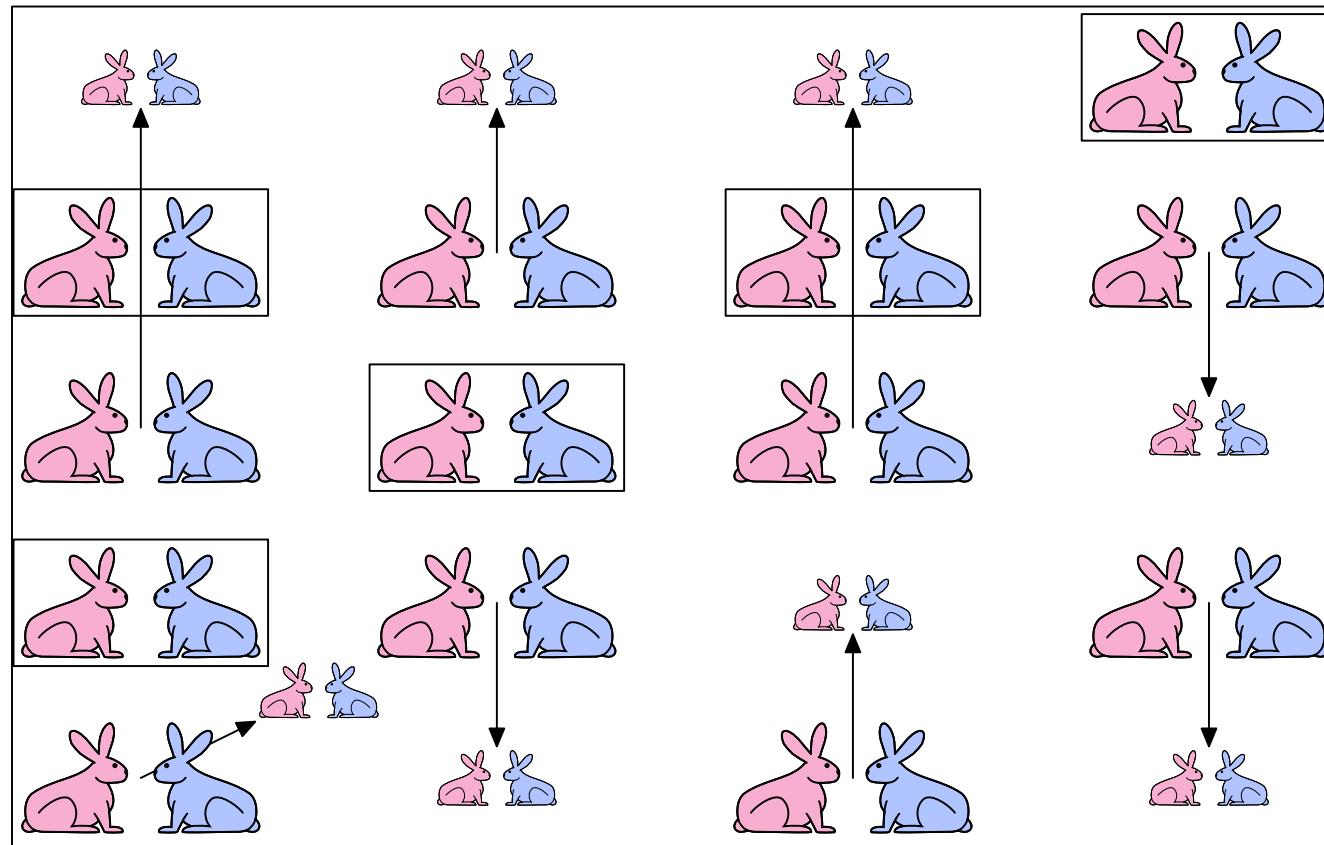
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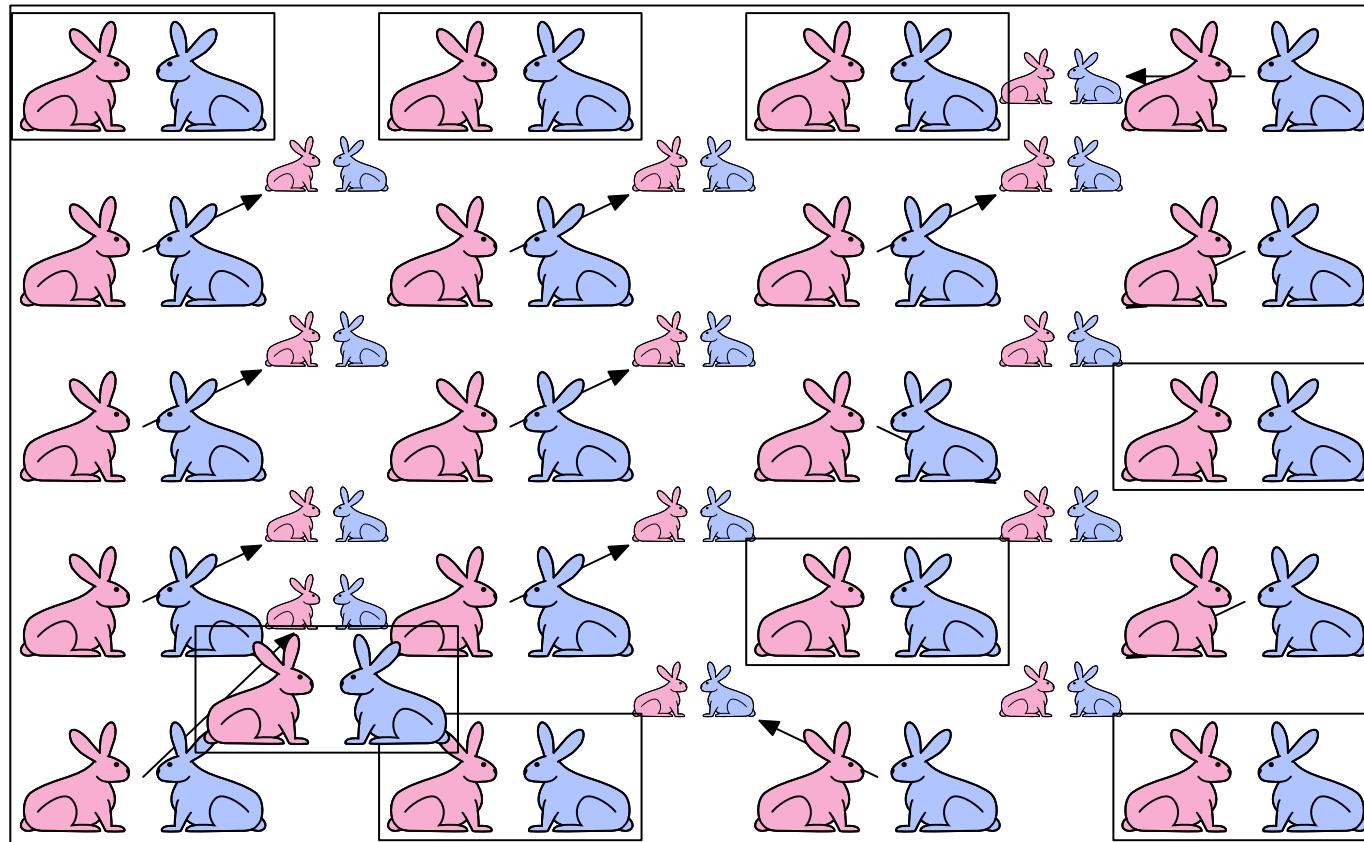


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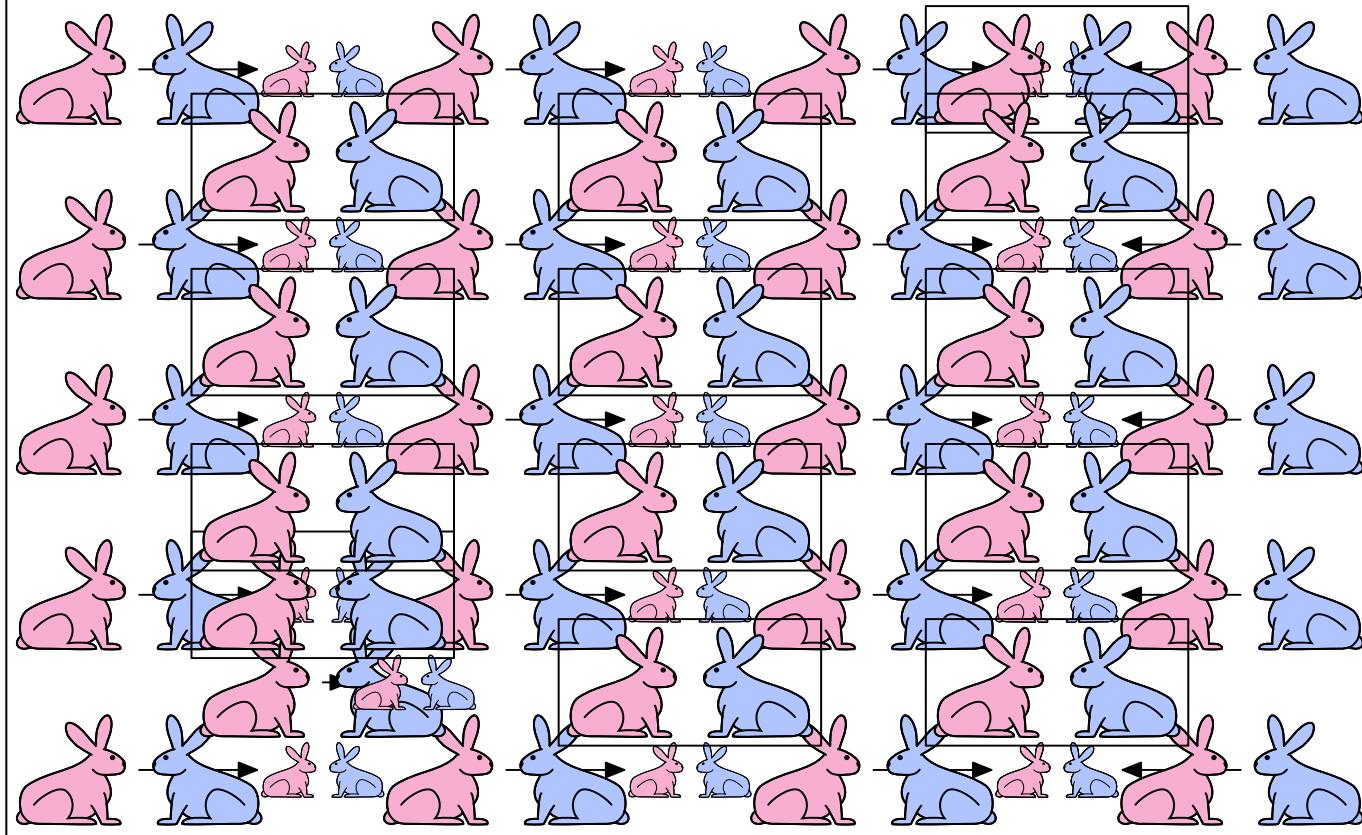


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# FIBONACCI

After one month (call it 0) – 1 female

After second month – still 1 female (now pregnant)

After third month – two females, one pregnant, one not

In general, females( $n$ ) = females( $n-1$ ) + females( $n-2$ )

- Every female alive at month  $n-2$  will produce one female in month  $n$ ;
- These can be added those alive in month  $n-1$  to get total alive in month  $n$

Month	Females
0	1

# FIBONACCI

---

- Base cases:
  - Females(0) = 1
  - Females(1) = 1
- Recursive case
  - $\text{Females}(n) = \text{Females}(n-1) + \text{Females}(n-2)$

# FIBONACCI

---

```
def fib(x):  
    """assumes x an int >= 0  
        returns Fibonacci of x"""  
  
    if x == 0 or x == 1:  
  
        return 1  
  
    else:  
  
        return fib(x-1) + fib(x-2)
```

# RECURSION ON NON-NUMERICS

---

- how to check if a string of characters is a palindrome, i.e., reads the same forwards and backwards
  - “Able was I, ere I saw Elba” – attributed to Napoleon
  - “Are we not drawn onward, we few, drawn onward to new era?” – attributed to Anne Michaels



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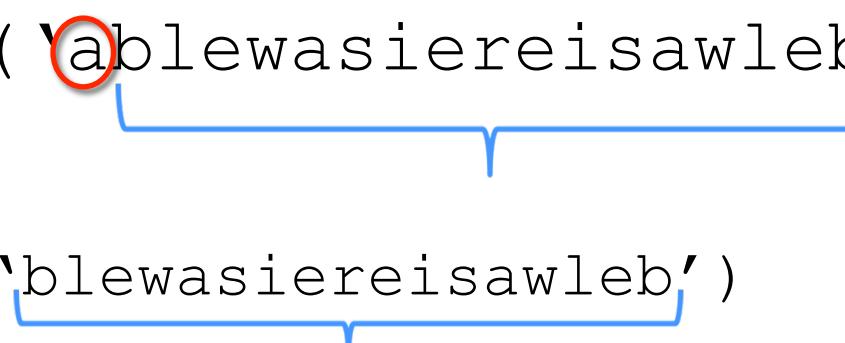
# SOLVING RECURSIVELY?

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- First, convert the string to just characters, by stripping out punctuation, and converting upper case to lower case
- Then
  - Base case: a string of length 0 or 1 is a palindrome
  - Recursive case:
    - If first character matches last character, then is a palindrome if middle section is a palindrome

# EXAMPLE

---

- ‘Able was I, ere I saw Elba’ → ‘ablewasiereisawleba’
  - isPalindrome ( ‘ablewasiereisawleba’ )  
is same as
    - ‘a’ == ‘a’ and  
isPalindrome ( ‘blewasiereisawleb’ )
- 

```
def isPalindrome(s):

    def toChars(s):
        s = s.lower()
        ans = ''
        for c in s:
            if c in 'abcdefghijklmnopqrstuvwxyz':
                ans = ans + c
        return ans

    def isPal(s):
        if len(s) <= 1:
            return True
        else:
            return s[0] == s[-1] and isPal(s[1:-1])

    return isPal(toChars(s))
```

# DIVIDE AND CONQUER

---

- an example of a “divide and conquer” algorithm
- solve a hard problem by breaking it into a set of sub-problems such that:
  - sub-problems are easier to solve than the original
  - solutions of the sub-problems can be combined to solve the original

# DICTIONARIES

---

# HOW TO STORE STUDENT INFO

---

- so far, can store using separate lists for every info

```
names = ['Ana', 'John', 'Denise', 'Katy']  
grade = ['B', 'A+', 'A', 'A']  
course = [2.00, 6.0001, 20.002, 9.01]
```

- a **separate list** for each item
- each list must have the **same length**
- info stored across lists at **same index**, each index refers to info for a different person

# HOW TO UPDATE/RETRIEVE STUDENT INFO

---

```
def get_grade(student, name_list, grade_list, course_list):  
    i = name_list.index(student)  
    grade = grade_list[i]  
    course = course_list[i]  
    return (course, grade)
```

- **messy** if have a lot of different info to keep track of
- must maintain **many lists** and pass them as arguments
- must **always index** using integers
- must remember to change multiple lists

# A BETTER AND CLEANER WAY – A DICTIONARY

---

- nice to **index item of interest directly** (not always int)
- nice to use **one data structure**, no separate lists

**A list**

0	Elem 1
1	Elem 2
2	Elem 3
3	Elem 4
...	...

index  
element

**A dictionary**

Key 1	Val 1
Key 2	Val 2
Key 3	Val 3
Key 4	Val 4
...	...

custom  
index by  
label  
element

# A PYTHON DICTIONARY

- store pairs of data
  - key
  - value

'Ana'	'B'
'Denise'	'A'
'John'	'A+'
'Katy'	'A'

my\_dict = { }  *empty dictionary*

grades = {'Ana': 'B', 'John': 'A+', 'Denise': 'A', 'Katy': 'A'}

*custom index by label*

*element*

↑      ↑      ↑      ↑      ↑      ↑      ↑      ↑

key1 val1      key2 val2      key3 val3      key4 val4

# DICTIONARY LOOKUP

---

- similar to indexing into a list
- **looks up the key**
- **returns the value** associated with the key
- if key isn't found, get an error

'Ana'	'B'
'Denise'	'A'
'John'	'A+'
'Katy'	'A'

```
grades = {'Ana':'B', 'John':'A+', 'Denise':'A', 'Katy':'A'}  
grades['John']      → evaluates to 'A+'  
grades['Sylvan']    → gives a KeyError
```

# DICTIONARY OPERATIONS

'Ana'	'B'
'Denise'	'A'
'John'	'A+'
'Katy'	'A'
'Sylvan'	'A'

```
grades = {'Ana': 'B', 'John': 'A+', 'Denise': 'A', 'Katy': 'A'}
```

- **add** an entry

```
grades['Sylvan'] = 'A'
```

- **test** if key in dictionary

'John' in grades  
'Daniel' in grades

→ returns True  
→ returns False

- **delete** entry

```
del(grades['Ana'])
```

# DICTIONARY OPERATIONS

'Ana'	'B'
'Denise'	'A'
'John'	'A+'
'Katy'	'A'

```
grades = {'Ana': 'B', 'John': 'A+', 'Denise': 'A', 'Katy': 'A'}
```

- get an **iterable that acts like a tuple of all keys**

grades.keys() → returns ['Denise', 'Katy', 'John', 'Ana']

*no guaranteed  
order*

- get an **iterable that acts like a tuple of all values**

grades.values() → returns ['A', 'A', 'A+', 'B']

*no guaranteed  
order*

# DICTIONARY KEYS and VALUES

---

- values
  - any type (**immutable and mutable**)
  - can be **duplicates**
  - dictionary values can be lists, even other dictionaries!
- keys
  - must be **unique**
  - **immutable** type (`int`, `float`, `string`, `tuple`, `bool`)
    - actually need an object that is **hashable**, but think of as immutable as all immutable types are hashable
  - careful with `float` type as a key
- **no order** to keys or values!

```
d = {4:{1:0}, (1,3):"twelve", 'const':[3.14,2.7,8.44]}
```

# list

vs

# dict

- **ordered** sequence of **elements**
- look up elements by an integer index
- indices have an **order**
- index is an **integer**

- **matches** “**keys**” to “**values**”
- look up one item by another item
- **no order** is guaranteed
- key can be any **immutable** type

# EXAMPLE: 3 FUNCTIONS TO ANALYZE SONG LYRICS

---

- 1) create a **frequency dictionary** mapping str:int
- 2) find **word that occurs the most** and how many times
  - use a list, in case there is more than one word
  - return a tuple (list, int) for (words\_list, highest\_freq)
- 3) find the **words that occur at least X times**
  - let user choose “at least X times”, so allow as parameter
  - return a list of tuples, each tuple is a (list, int) containing the list of words ordered by their frequency
  - IDEA: From song dictionary, find most frequent word. Delete most common word. Repeat. It works because you are mutating the song dictionary.

# CREATING A DICTIONARY

---

```
def lyrics_to_frequencies(lyrics):  
    myDict = {}  
    for word in lyrics:  
        if word in myDict:  
            myDict[word] += 1  
        else:  
            myDict[word] = 1  
    return myDict
```

can iterate over list  
in dictionary  
update value  
associated with key

# USING THE DICTIONARY

---

```
def most_common_words(freqs):
    values = freqs.values()
    best = max(values)
    words = []
    for k in freqs:
        if freqs[k] == best:
            words.append(k)
    return (words, best)
```

can iterate over keys  
in dictionary

this is an iterable, so can  
apply built-in function

# LEVERAGING DICTIONARY PROPERTIES

---

```
def words_often(freqs, minTimes):
    result = []
    done = False
    while not done:
        temp = most_common_words(freqs)
        if temp[1] >= minTimes:
            result.append(temp)
            for w in temp[0]:
                del(freqs[w])
        else:
            done = True
    return result

print(words_often(beatles, 5))
```

can directly mutate  
dictionary; makes it  
easier to iterate

# FIBONACCI RECURSIVE CODE

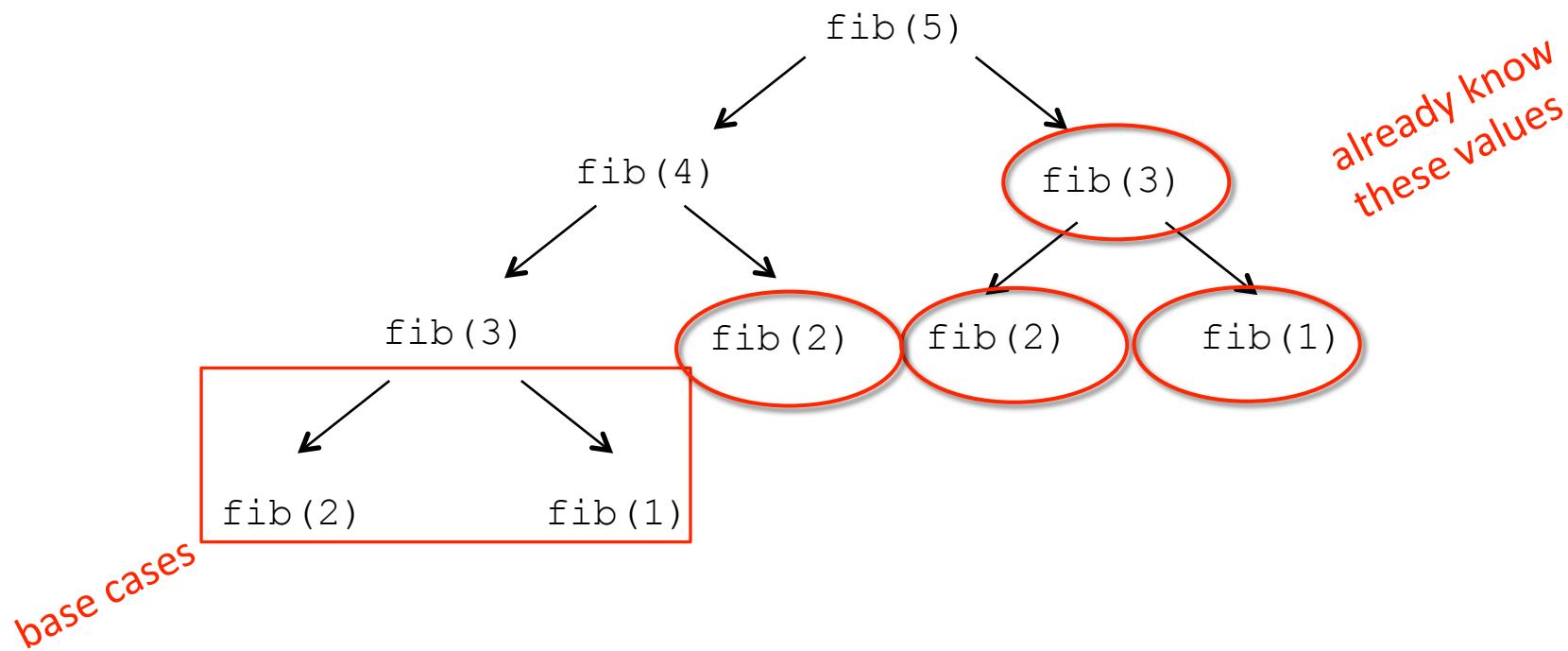
---

```
def fib(n):  
    if n == 1:  
        return 1  
    elif n == 2:  
        return 2  
    else:  
        return fib(n-1) + fib(n-2)
```

- two base cases
- calls itself twice
- this code is inefficient

# INEFFICIENT FIBONACCI

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$$



- **recalculating** the same values many times!
- could keep **track** of already calculated values

# FIBONACCI WITH A DICTIONARY

---

```
def fib_efficient(n, d):
    if n in d:
        return d[n]
    else:
        ans = fib_efficient(n-1, d) + fib_efficient(n-2, d)
        d[n] = ans
        return ans

d = {1:1, 2:2}
print(fib_efficient(6, d))
```

Method sometimes  
called "memoization"

Initialize dictionary  
with base cases

- do a **lookup first** in case already calculated the value
- **modify dictionary** as progress through function calls

# EFFICIENCY GAINS

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- Calling `fib(34)` results in 11,405,773 recursive calls to the procedure
- Calling `fib_efficient(34)` results in 65 recursive calls to the procedure
- Using dictionaries to capture intermediate results can be very efficient
- But note that this only works for procedures without side effects (i.e., the procedure will always produce the same result for a specific argument independent of any other computations between calls)

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