Omni Compiler Runtime Library for Coarray

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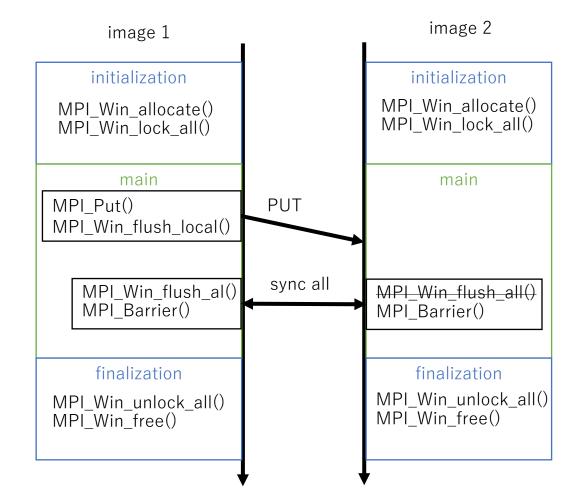
Omni compiler runtime library for coarray

- Supported one-sided communication library
 - MPI 3
 - GASNet
 - FJRDMA

Overview of implementation over MPI-3

• For the left program, the omni runtime library calls MPI functions as shown in the right figure.

```
int array[N]:[*];
void main()
{
   int buf[N] = ...;
   if (image_num == 1){
      array[:]:[2] = buf[:];
   }
   xmp_sync_all();
}
```



Initialization of coarray

xmp_init_all() calls _XMP_mpi_onesided_initialize() which setups MPI window

- Parameter heap_size is coarray heap size. We can specify it by environment variable XMP_ONESIDED_HEAP_SIZE
- MPI_Win_allocate() creates a MPI window and allocates memory for coarray. All coarrays will belong to this window.
- MPI_Win_lock_all() begins an access epoch for the program. All coarray communication will be issued in this epoch.

Finalization of coarray

xmp_finalize_all() calls _XMP_mpi_onesided_finalize() which frees MPI window

```
void _XMP_mpi_onesided_finalize(void)
{
    MPI_Win_unlock_all(_xmp_onesided_win);
    _XMP_mpi_destroy_shift_queue(false); // finalize coarray heap
    MPI_Win_free(&_xmp_onesided_win);
}
```

- MPI_Win_unlock_all() ends the access epoch began at _XMP_mpi_onesided_initialize().
- MPI_Win_free() frees the window and heap for coarray.

Allocation of coarray

_XMP_coarray_malloc()

- To create a coarray descriptor, _XMP_coarray_set_info() set coarray information specified by _XMP_coarray_malloc_info_X() and _XMP_coarray_malloc_image_info_X()
- _XMP_mpi_coarray_malloc() allocate the coarray from the coarray heap
- _push_coarray_queue() register the coarray descriptor.

PUT (continuous)

_XMP_coarray_continuous_put (special pattern, high-speed)

- If target_rank is equal to my rank, it copies memory in local.
- laddr is local memory address and raddr is remote memory offset from the window base address.
- MPI_Win_flush_local() waits end of the PUT locally. (doesn't wait remote completion)

GET (continuous)

_XMP_coarray_continuous_get (special pattern, high-speed)

- It is almost the same as _XMP_coarray_continuous_put()
- MPI_Win_flush_local() waits both local and remote completion of the GET

PUT (normal)

_XMP_coarray_put (general pattern)

```
void _XMP_coarray_put(void *dst_desc, void *src_addr, void *src_desc)
  int target_rank = calc_target_rank(_image_num, ...);
  if(target_rank == _XMP_wold_rank){
    _XMP_local_put(...); return; //local memory copy
  char *laddr = calc_local_addr(src_addr, _src_info);
  char *raddr = calc_remote_addr(dst_desc, target_rank, _dst_info);
 MPI_Datatype src_types[_XMP_N_MAX_DIM], dst_types[_XMP_N_MAX_DIM];
  create_datatypes(src_types, _src_dims, _src_info, element_size);
  create_datatypes(dst_types, _dst_dims, _dst_info, element_size);
 MPI Put(
                    laddr, 1, src_types[0], target_rank,
          (MPI_Aint)raddr, 1, dst_types[0], win);
 MPI_Win_Flush_local(target_rank, _xmp_onesided_win);
  free_datatypes(src_types, _src_dims);
  free_datatypes(dst_types, _dst_dims);
```

- create_datatypes() makes N-dimensional array datatype with MPI_Type_create_hvector()
 - src_types[0] is N-dim datatype, src_types[1] is (N-1)-dim datatype, ...
 - Base datatype is MPI_BYTE

GET (normal)

_XMP_coarray_get (general pattern)

```
void _XMP_coarray_get(void *src_desc, void *dst_addr, void *dst_desc)
  int target_rank = calc_target_rank(_image_num, ...);
  if(target == _XMP_wold_rank){
    _XMP_local_put(...); return; //local memory copy
  char *laddr = calc_local_addr(dst_addr, _dst_info);
  char *raddr = calc_remote_addr(src_desc, target_rank, _src_info);
 MPI_Datatype src_types[_XMP_N_MAX_DIM], dst_types[_XMP_N_MAX_DIM];
  create_datatypes(src_types, _src_dims, _src_info, element_size);
  create_datatypes(dst_types, _dst_dims, _dst_info, element_size);
 MPI_Get(
                    laddr, 1, src_types[0], target_rank,
          (MPI_Aint)raddr, 1, dst_types[0], win);
 MPI_Win_Flush_local(target_rank, _xmp_onesided_win);
  free_datatypes(src_types, _src_dims);
  free_datatypes(dst_types, _dst_dims);
```

It is almost the same with _XMP_coarray_put()

Synchronization (1/2)

xmp_sync_all

```
void xmp_sync_all(int *status)
{
   if(!_is_win_flushed){
      MPI_Win_flush_all(_xmp_mpi_onesided_win);
      _is_win_flushed = true;
   }
   MPI_Barrier(MPI_COMM_WORLD);

   //this function returns no status now
}
```

- _is_win_flushed is false if some PUT/GETs are issued after last MPI_Win_flush_all()
- MPI_Win_flush_all() waits all end of preceding PUTs issued by this rank
- MPI_Barrier() synchronizes all images

Synchronization (2/2)

xmp_sync_images

```
void xmp_sync_images(int num, int *image_set, int *status)
{
  int rank_set[num];
  for(int i = 0; i < num; i++) rank_set[i] = image_set[i] - 1;

  if(!_is_win_flushed){
    MPI_Win_flush_all(_xmp_mpi_onesided_win);
    _is_win_flushed = true;
  }

  //send signals
  for(int i = 0; i < num; i++){
    MPI_Send(NULL, 0, MPI_BYTE, rank_set[i], TAG_SYNCREQ, COMM_WORLD);
  }

  //continue to the next slide</pre>
```

- MPI_Win_flush_all() waits end of all preceding PUTs
- MPI_Send() sends signals to ranks in rank_set

Synchronization (2/2, cont'd)

xmp_sync_images

```
// cont'd
// wait signals
                                              _signal_table[] has # of
for(int rank : rank_set){
  if( _signal_table[rank] > 0 ){
                                              received signals
    _signal_table[rank] -= 1;
    remove(rank_set, rank);
while( size(rank_set) > 0 ){
 MPI_Recv(NULL, 0, MPI_BYTE, MPI_ANY_SOURCE, TAG_SYNCREQ, MPI_COMM_WORLD, &st);
 int rank = st.MPI_SOURCE;
  if( contains(rank_set, rank) )
    remove(rank_set, rank);
  else
    _signal_table[rank] += 1;
//this function returns no status now
```

- Firstly, matching with already received signals in _signal_table
- MPI_Recv() receives a signal from any ranks. If rank_set contains the rank, it is removed from rank_set, else it is added to _signal_table. This is repeated while rank_set is not empty.

Question

 MPI 3.0 spec. p.459, Example 11.9: the following example is unsafe even in the unified model, because the load of X can not be guaranteed to occur after the MPI_BARRIER. ... Compiler and hardware specific notations could ensure the load occurs after the data is updated, or <u>explicit one-sided</u> <u>synchronization calls</u> can be used to ensure the proper result.

```
Process A:

MPI_Win_lock_all

MPI_Put(X) /* update to window */

MPI_Win_flush(B)

MPI_Barrier

MPI_Barrier

MPI_Barrier

MPI_Win_unlock_all
```

 We have to insert one-sided synchronization call to ensure result. Is the following code correct?

```
Process A:

MPI_Win_lock_all

MPI_Put(X) /* update to window */

MPI_Win_flush(B)

MPI_Barrier

MPI_Win_sync
load X

MPI_Win_unlock_all
```

 and, should we change synchronization functions like the following?

```
void xmp_sync_all(int *status)
{
   if(!_is_win_flushed){
     MPI_Win_flush_all(_xmp_mpi_onesided_win);
     _is_win_flushed = true;
   }
   MPI_Barrier(MPI_COMM_WORLD);
}
```



```
void xmp_sync_all(int *status)
{
   if(!_is_win_flushed){
      MPI_Win_flush_all(_xmp_mpi_onesided_win);
      _is_win_flushed = true;
   }
   MPI_Win_sync(_xmp_mpi_onesided_win);
   MPI_Barrier(MPI_COMM_WORLD);
}
```

