INFORMED SEARCH

Heuristic search



INFORMED SEARCH STRATEGIES

- Uses domain-specific hints about the location of goals.
- Greedy best-first search
- A* search
- Search contours
- Satisficing search
- Memory-bound search
- Bidirectional heuristic search

INFORMED (HEURISTIC) SEARCH

- Best-first search with evaluation function f(n).
- Most best-first algorithms include as a component of f a heuristic function h(n)
- h(n) = estimated cost of the cheapest path from the state at node n to a goal state.

Straight line between Arad to Bucharest!

GREEDY BEST-FIRST SEARCH

- Tries to expand the node that is closest to the goal.
- Let h(n) be straight line distance heuristic $h_{SLD}(n)$.

GREEDY BEST-FIRST SEARCH

Figure 3.22

Arad	366	Mehadia	241
Bucharest	0	Neamt	234
Craiova	160	Oradea	380
Drobeta	242	Pitesti	100
Eforie	161	Rimnicu Vilcea	193
Fagaras	176	Sibiu	253
Giurgiu	77	Timisoara	329
Hirsova	151	Urziceni	80
Iasi	226	Vaslui	199
Lugoj	244	Zerind	374

Values of h_{SLD} —straight-line distances to Bucharest.

BEST-FIRST SEARCH

CHOOSE A NODE, N, WITH MINIMUM VALUE OF SOME EVALUATION FUNCTION, F(N)=H(N).

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function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure
  node \leftarrow Node(State=problem.INITIAL)
  frontier \leftarrow a priority queue ordered by f, with node as an element
  reached ← a lookup table, with one entry with key problem. INITIAL and value node
  while not IS-EMPTY(frontier) do
     node \leftarrow Pop(frontier)
     if problem.IS-GOAL(node.STATE) then return node
     for each child in Expand(problem, node) do
       s \leftarrow child.STATE
       if s is not in reached or child.PATH-COST < reached[s].PATH-COST then
          reached[s] \leftarrow child
          add child to frontier
  return failure
function EXPAND(problem, node) yields nodes
  s \leftarrow node.STATE
  for each action in problem.ACTIONS(s) do
     s' \leftarrow problem.RESULT(s, action)
     cost \leftarrow node.PATH-COST + problem.ACTION-COST(s, action, s')
```

yield NODE(STATE=s', PARENT=node, ACTION=action, PATH-COST=cost)

GREEDY BEST-FIRST SEARCH

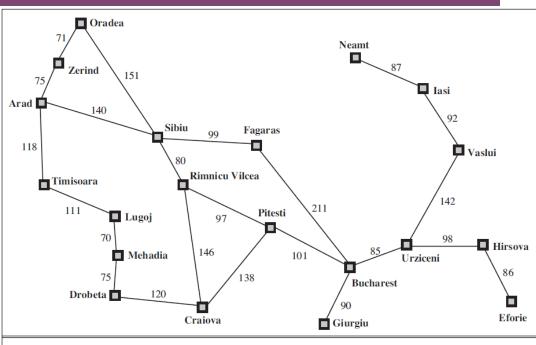


Figure 3.2 A simplified road map of part of Romania.

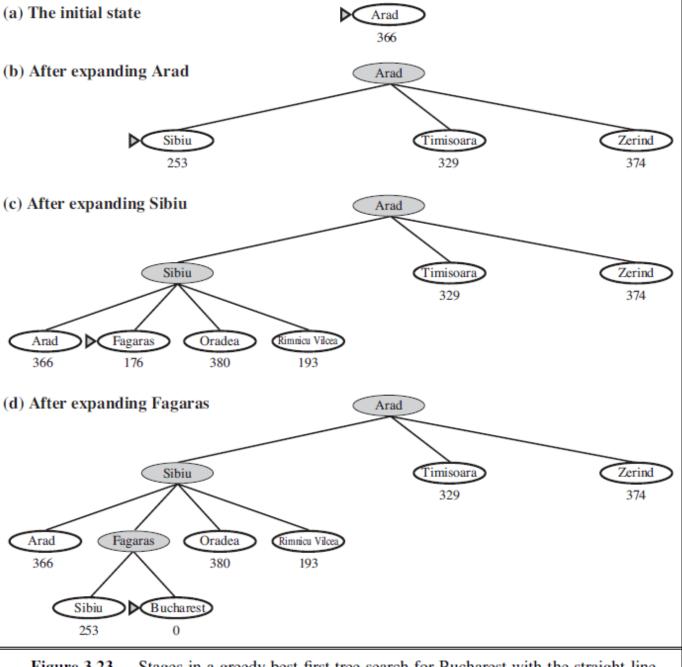


Figure 3.23 Stages in a greedy best-first tree search for Bucharest with the straight-line distance heuristic h_{SLD} . Nodes are labeled with their h-values.

GREEDY BEST-FIRST SEARCH

- Challenges:
- Not optimal
- Tree search is not optimal and incomplete.
- Worst-case time and space complexity for Tree version is O(b^m). m is maximum depth.
- Graph search version is complete in finite space.