Cryptography

Classification of Cryptography Cryptoanalysis Symmetric Key Cryptography

M S Vilkhu

Symmetric Key Cryptography

- Symmetric encryption, also referred to as conventional encryption or single-key encryption, was the only type of encryption in use prior to the development of publickey encryption in the 1970s.
- It remains by far the most widely used of the two types of encryption.
- Approach
- Study some of the of symmetric ciphers.
- look at a general model for the symmetric encryption process;
- Help to understand the context within which the algorithms are used.
- Next, we examine a variety of algorithms in use before the computer era.
- Finally, we look briefly at a different approach known as steganography.
- Study two most widely used symmetric cipher: DES and AES

Basic Terms

- Cryptography The art or science encompassing the principles and methods of transforming an intelligible message into one that is unintelligible, and then retransforming that message back to its original form
- Cryptology Both cryptography and cryptanalysis
- Cryptanalysis The study of principles and methods of transforming an unintelligible message back into an intelligible message without knowledge of the key. Also called code breaking

Basic Terms

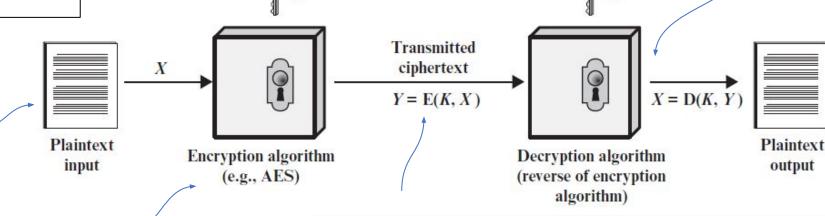
- Plaintext The original intelligible message
- Cipher text The transformed message
- Cipher An algorithm for transforming an intelligible message into one that is unintelligible by transposition and/or substitution methods
- Key Some critical information used by the cipher, known only to the sender& receiver
- Encipher (encode) The process of converting plaintext to cipher text using a cipher and a key
- Decipher (decode) the process of converting cipher text back into plaintext using a cipher and a key
- Code An algorithm for transforming an intelligible message into an unintelligible one using a code-book

Symmetric Cypher Model

Secret key: The secret key is also input to the **encryption algorithm**.

The key is a value **independent** of the plaintext and of the algorithm. Different key different output

Plaintext: This is the original intelligible message or data input to algo.



Secret key shared by

sender and recipient

Encryption algorithm:

The encryption algorithm performs transformation

Ciphertext: This is the scrambled message (output).

depends on the plaintext and the secret key.

For a given message, two different keys will produce two different ciphertexts. The ciphertext is an apparently random stream of data and, as it stands, is unintelligible.

Secret key shared by

sender and recipient

Decryption algorithm: This is

essentially the encryption

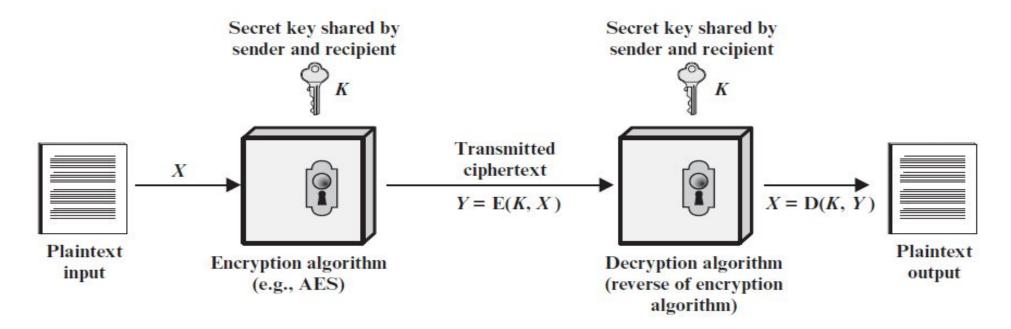
algorithm run in reverse. It

takes the ciphertext and the

secret key and produces the

original plaintext.

Symmetric Cypher Model



5 components

- Plaintext
- Encryption algorithm
- Secret key
- Ciphertext
- Decryption algorithm

Symmetric Cypher Model

- Plaintext: This is the original intelligible message or data input to algo.
- Encryption algorithm: The encryption algorithm performs transformation
- Secret key: The secret key is also input to the encryption algorithm.
 - The key is a value **independent** of the plaintext and of the algorithm.
 - The algorithm will produce a **different output depending** on **the specific key** being used at the time. The **exact substitutions and transformations** performed by the algorithm **depend on the key**.
- Ciphertext: This is the scrambled message produced as output. It depends on the plaintext and the secret key. For a given message, two different keys will produce two different ciphertexts. The ciphertext is an apparently random stream of data and, as it stands, is unintelligible.
- **Decryption algorithm:** This is essentially the encryption algorithm run in reverse. It takes the ciphertext and the secret key and produces the original plaintext.

symmetric encryption: Two requirements

- Two requirements for secure use of symmetric encryption:
- A strong encryption algorithm

we would like the algorithm to be such that an opponent who knows the algorithm and has access to one or more ciphertexts would be unable to decipher the ciphertext or figure out the key.

A secret key known only to sender / receiver

Obtain copy of key through secure means.

Keep the **key secure**

Algo – known to both user and attacker

Key – to be kept secret

With the use of symmetric encryption, the principal security problem is maintaining the secrecy of the key.

symmetric encryption

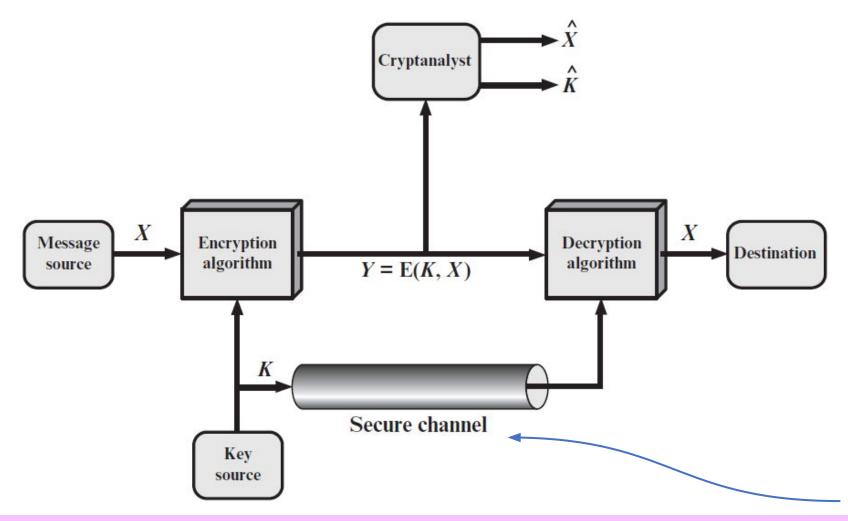
- Algorithms for symmetric encryption is not required to be kept secret
- It means that chip can be made with data encryption algorithms,
- requirement is to keep the **key secret**

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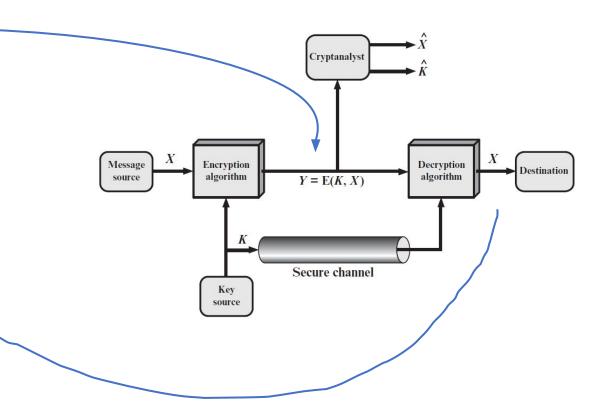
Symmetric Key Model



Key generated at source should be shared with the destination through secure channel

Model

- Y = E (X, K)
- Y is crypto text produced by encryption algo E with message X and Key K
- X = D(Y, K)
- X message, D decrypt algorithm,
 Y is crypto text, K key



Cryptography Characterized by dimensions

Cryptographic systems are characterized along three independent dimensions:

- 1. The type of operations used for transforming plaintext to ciphertext.
- 2. The number of keys used.
- 3. The way in which the plaintext is processed.

Cryptography Characterized by dimensions

Cryptographic systems are **characterized** along **three independent dimensions**:

1. The type of operations used for transforming plaintext to ciphertext.

All encryption algorithms are based on two general principles:

substitution, in which each element in the plaintext (bit, letter, group of bits or letters) is mapped into another element, and

transposition, in which elements in the plaintext are rearranged.

The fundamental requirement - **no information be lost** (i.e., that all **operations are reversible**).

Most systems, referred to as product systems, involve multiple stages of substitutions and transpositions.

Cryptography Characterized by dimensions

Cryptographic systems are characterized along three independent dimensions:

2. The number of keys used.

If both sender and receiver use the **same key**, the system is referred to as **symmetric**, single-key, secret-key, or conventional encryption.

If the sender and receiver use different keys, the system is referred to as asymmetric, two-key, or public-key encryption.

3. The way in which the plaintext is processed.

A **block cipher** processes the input **one block of elements** at a time, producing an output block for each input block.

A **stream cipher** processes the input elements **continuously**, producing output **one element at a time**, as it goes along.

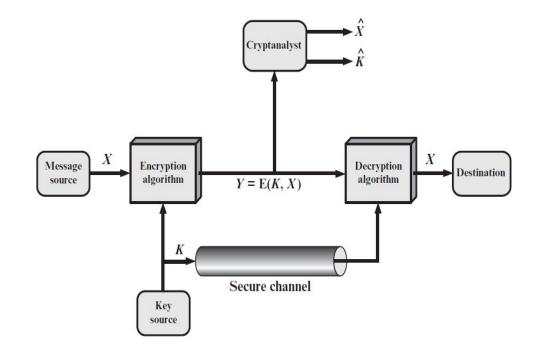
Break the code

• With help of symmetric algorithm and the cyphertext – is it possible to decrypt the message (plaintext)?

- Algo is know
- Cypher text can be captured or read while the message is transmitted
- Is it possible to determine plain text?

Break the code

- An opponent, observing Y but not having access to K or X,
- may attempt to recover X or K or both X and K.
- assumed opponent knows the encryption (E) and decryption (D) algorithms.
- Two scenarios
- If Interest & focus is X particular message then efforts to estimate plain text X
- If opponent is interested in being able to read future messages - an attempt is made to recover K by generating an estimate K



- What is Cryptanalysis?
- What is Brute-Force Attack
- Difference
- **objective** of attacking an encryption system is to **recover the key** in use rather than simply to recover the plaintext of a single ciphertext.
- two general approaches to attacking a conventional encryption scheme:
- Cryptanalysis:
- Brute-force attack:

- **Cryptanalysis**: Cryptanalytic attacks
 - rely on the nature of the algorithm
 - some knowledge of the general characteristics of the plaintext or
 - even some sample plaintext-ciphertext pairs.
- This type of attack exploits the characteristics of the algorithm to attempt to deduce a
 - specific plaintext or
 - to **deduce the key** used.

- **Cryptanalysis**: Cryptanalytic attacks
 - rely on the nature of the algorithm
 - some knowledge of the general characteristics of the plaintext or
 - even some sample plaintext-ciphertext pairs.
- This type of attack exploits the characteristics of the algorithm to attempt to deduce a
 specific plaintext or to deduce the key being used.
- Brute-force attack: The attacker tries every possible key on a piece of ciphertext until an intelligible translation into plaintext is obtained.
- On average, half of all possible keys must be tried to achieve success.

If either type of **attack succeeds** in deducing the key, the effect is **catastrophic**: All future and past messages encrypted with that key are compromised.

Aspect	Cryptanalysis	Brute Force Attack
Definition	Analyzing cryptographic systems to find weaknesses or vulnerabilities.	Systematically trying every possible key or password.
Method	Uses analytical techniques, such as pattern recognition, statistical analysis, and algebraic methods.	Exhaustively tests all possible keys or passwords.
Goal	To find a more efficient way to decrypt data or discover the key.	To find the correct key or password through exhaustive search.
Efficiency	Can be more efficient if vulnerabilities are found ; may reduce the amount of work needed.	Can be very inefficient , especially with strong encryption, as it involves testing all possibilities.
Knowledge Required	Requires understanding of the cryptographic algorithm and its potential weaknesses .	No need for knowledge about the algorithm, just a systematic approach to trying all possible keys.
Examples of Techniques	Differential cryptanalysis, linear cryptanalysis, known plaintext attacks.	Trying every possible key for a given encryption algorithm, e.g., a 128-bit key requires 2 ¹²⁸ attempts.
Computational Demand	Can vary widely depending on the weakness exploited; sometimes less computational power is needed compared to brute force.	High computational demand, especially for strong encryption with large key sizes.
Practicality	Effective if weaknesses in the encryption are found; can be feasible with sophisticated techniques.	Often impractical for strong encryption due to the vast number of combinations.

Cryptography

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- 1. Pattern Recognition:
- 2. Statistical Analysis:
- 3. Algebraic Methods:

1. Pattern Recognition:

- Definition:
- The process of identifying regularities, structures, or repeating elements within a cryptographic message or encrypted data.
- In cryptanalysis, this often involves looking for repeated sequences, common phrases, or structural patterns, visual pattern that can reveal information about the encryption method or the plaintext.
- to infer plaintext or encryption rules.
- Application:
- For example, in a substitution cipher, the frequency of letter or word patterns in the ciphertext can be analyzed to guess the substitution mappings.

It's more qualitative and relies on recognizing specific features or sequences.

2. Statistical Analysis:

Definition: The application of statistical methods to **analyze the frequency and distribution of symbols or patterns** in ciphertext.

This can help identify anomalies or deviations from expected distributions that can be exploited to break the cipher.

Application: In classical cryptography, **frequency analysis** involves studying the **frequency of letters** or groups of letters in ciphertext.

For instance, in **English text, the letter 'E' is the most common**, so an analyst might look for the most frequent letters in the ciphertext and hypothesize that they correspond to 'E' or other common letters.

Focuses on quantitative analysis of the frequency and distribution of symbols.

It involves applying statistical methods to detect deviations from expected patterns and deduce information about the encryption.

3. Algebraic Methods:

Definition: Techniques that **use algebraic structures and equations** to analyze and solve cryptographic problems.

This involves **applying mathematical** operations and theories to **deduce the encryption keys** or to **understand the transformation functions** used in encryption.

Application:

In the case of linear cryptanalysis, algebraic techniques are used to create linear approximations of the encryption function.

For **block ciphers** like **DES**, algebraic methods might involve **solving systems of linear equations to determine the key.**

Aspect	Pattern Analysis	Statistical Analysis
Focus	Repetitions and patterns in the ciphertext	Frequency distribution of letters or letter
		groups

For instance, the word

MI LI TA RY

will always produce the same ciphertext letter in **the first and third** positions regardless of the keywords used.

Patterns like these can be catalogued and matched against single-letter repeats in the ciphertext.

Analyzes the **frequency** of letters, digraphs, or other elements in the ciphertext based on known statistical properties of the language in the plaintext (e.g., in English, the letter "E" is the most frequent, followed by "T" and "A").

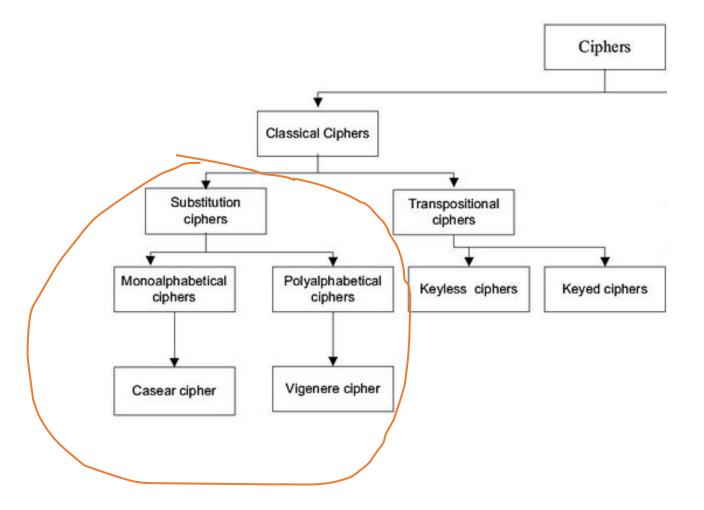
if "X" appears most frequently in the ciphertext, it might correspond to "E" in English since "E" is the most common letter in English.

- examine a sampling of what might be called classical encryption techniques.
- two basic building blocks of all encryption techniques
 - Substitution Techniques
 - Transposition techniques
 - There are systems which combine both

study of these techniques enables us

- to illustrate the **basic approaches** to symmetric encryption used today and
- the types of **cryptanalytic attacks** that must be anticipated.

Classification



- Substitution Techniques
- the letters of plaintext are replaced by other letters or by numbers or symbols
- If the **plaintext** is viewed as a **sequence of bits**, then substitution involves replacing plaintext bit patterns with ciphertext bit patterns.
- replacing each element of the plaintext with another element according to a fixed system or rule.
- The idea to disguise the original text by substituting each letter or symbol with another.

• Substitution Techniques

Monoalphabetic Ciphers

- Caesar Cipher
- Atbash Cipher
- Playfair Cipher

Polyalphabetic ciphers

• The Vigenère cipher

- Substitution Techniques
- Caesar Cipher
- The Caesar cipher is a simple substitution cipher where each letter in the plaintext is shifted a fixed number of places down or up the alphabet.
- use of a substitution cipher was by Julius Caesar. The Caesar cipher involves replacing each letter of the alphabet with the letter standing three places further down the alphabet

```
plain: a b c d e f g h i j k l m n o p q r s t u v w x y z cipher: D E F G H I J K L M N O P Q R S T U V W X Y Z A B C
```

• Substitution Techniques

Caesar Cipher

```
Plaintext: HELLO Shift: 3  
Ciphertext: KHOOR  
In this example, each letter in "HELLO" is shifted by 3 positions in the alphabet H \rightarrow K  
E \rightarrow H  
L \rightarrow O  
L \rightarrow O  
plain: a b c d e g h i j k I m n o p q r s t u v w x y z cipher: D E F G H J K L M N O P QR S T U V W X Y Z A B C
```

- Substitution Techniques
- Caesar Cipher

Mathematically

$$C = E(3, p) = (p + 3) \mod 26$$
 where p-plaintext, c ciphertext

A shift may be of any amount, so that the general Caesar algorithm is $C = E(k, p) = (p + k) \mod 26$

k takes the value from 1 to 25

Decrypt algorthym

$$p = D(k, C) = (C - k) \mod 26$$

- Substitution Techniques
- Caesar Cipher

If it is known that a given ciphertext is a **Caesar cipher**, then a **brute-force cryptanalysis** is easily performed:

simply try all the 25 possible keys. shows the results of applying this strategy to the example ciphertext.

.

Three important characteristics of this problem enabled us to use a bruteforce cryptanalysis:

- 1. The encryption and decryption algorithms are known.
- 2. There are only 25 keys to try.
- 3. The language of the **plaintext is known** and easily **recognizable**.

- Substitution Techniques
- Caesar Cipher

Very important

3. The language of the **plaintext is known** and easily **recognizable**.

If the file is compressed and then encrypted, then recognizing the plaintext is difficult

- Substitution Techniques
- Caesar Cipher

Very important

3. The language of the plaintext is known and easily recognizable.

If the file is compressed and then encrypted, then recognizing the plaintext is difficult

File compression is the process of **reducing the size** of a file by encoding its data more efficiently, typically by **eliminating redundancies, re-encoding repeated patterns, or using mathematical algorithms.** The goal of compression is to make files smaller, so they use less storage space and can be transmitted more quickly over networks

- Substitution Techniques
- Atbash Cipher
- The Atbash cipher is a substitution cipher where **each letter is mapped to its reverse** in the alphabet.

```
Plaintext: HELLO
```

Ciphertext: SVOOL

In this example, each letter in "HELLO" is replaced with its reverse:

 $H \rightarrow S$

 $E \rightarrow V$

 $L \rightarrow 0$

 $L \rightarrow 0$

 $\mathsf{O} \to \mathsf{L}$

- Substitution Techniques
- Monoalphabetic Ciphers
- With only 25 possible keys, the Caesar cipher is far from secure.
- increase in the key space can be achieved by allowing an arbitrary substitution.

approach is referred to as a **monoalphabetic substitution** cipher, because a **single cipher** alphabet (mapping from plain alphabet to cipher alphabet) is used per message

line of attack.

If the cryptanalyst knows the nature of the plaintext (e.g., noncompressed English text), then the analyst can exploit the **regularities of the language**.

- Substitution Techniques
- Monoalphabetic Ciphers
- Variation to Caesar Cipher
- assignment for the Caesar cipher:

```
plain: abcd efghij kl mnopqrstuvwxyz
```

cipher: DEFGHIJKLMNOPQRSTUVWXYZABC

- instead, the "cipher" line can be any permutation of the 26 alphabetic characters,
- then there are 26! or greater than 4 x 10²⁶ possible keys

Permutation.

A **permutation** of a finite set of elements S is an **ordered sequence** of all the elements of S, with each element **appearing exactly once**. For example, if $S = \{a, b, c\}$, six permutations of S For n, there are n! combinations

- Regularities of the language.
- Regularities of a language refer to the consistent patterns and structures found within a language that can be exploited in various fields, including cryptography, linguistics, and natural language processing.
- Understanding these regularities helps in areas such as **cryptanalysis**, text analysis, and language modeling.

Regularities of the language.

- 1. Frequency of Letters and Words certain letters appear more frequently than others
- **2. Common Letter Combinations** -Certain pairs (bigrams) or triplets (trigrams) of letters appear frequently in text. For instance, 'TH', 'HE', and 'ING
- **3. Morphology** -Languages have regular patterns for creating new words or altering word forms ed for past tense,
- **4. Grammar and Syntax** -regular syntactic rules governing sentence structure
- **5.** Patterns in Punctuation
- **6. Textual Cohesion and Coherence** structured, different parts of a text are linked together

- Substitution Techniques
- Monoalphabetic Ciphers
- Cypher text to be solved –
- UZQSOVUOHXMOPVGPOZPEVSGZWSZOPFPESXUDBMETSXAIZVUEPHZHMDZSHZOWSFP APPDTSVPQUZWYMXUZUHSXEPYEPOPDZSZUFPOMB**ZWP**FUPZHMDJUDTMOHMQ
- Relative frequency of letters.

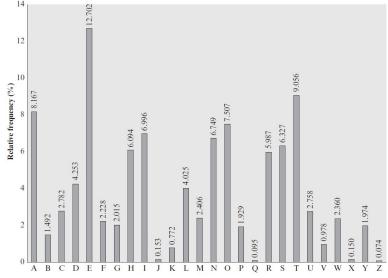
P 13.33	H 5.83	F 3.33	B 1.67	C 0.00
Z 11.67	D 5.00	W 3.33	G 1.67	K 0.00
S 8.33	E 5.00	Q 2.50	Y 1.67	L 0.00
U 8.33	V 4.17	T 2.50	I 0.83	N 0.00
O 7.50	X 4.17	A 1.67	J 0.83	R 0.00
M 6.67				

- Substitution Techniques
- Monoalphabetic Ciphers
- P and Z are the equivalents of plain letters e and t, but it is not certain which is which
- The letters S, U, O, M, and H are all of relatively high frequency and probably correspond to plain letters from the set {a, h, i, n, o, r, s}.

• The letters with the lowest frequencies (namely, A, B, G, Y, I, J) are likely included in the

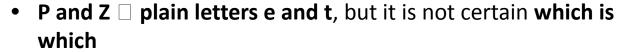
set {b, j, k, q, v, x, z}.

P 13.33	H 5.83	F 3.33	B 1.67	C 0.00
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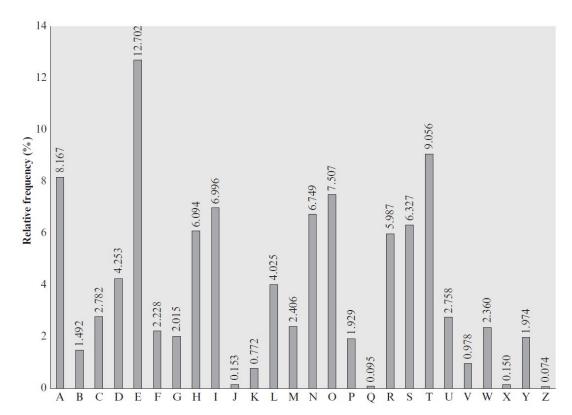


- Substitution Techniques
- Monoalphabetic Ciphers

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S 8.33	E 5.00	Q 2.50	Y 1.67	L 0.00
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- S, U, O, M, and H are all of relatively high frequency □ set {a, h, i, n, o, r, s}.
- lowest frequencies (namely, A, B, G, Y, I, J) are likely included in the □ set {b, j, k, q, v, x, z}.



- Substitution Techniques
- Monoalphabetic Ciphers
- powerful tool is to look at the for Continued analysis of frequencies plus trial and error should easily yield a
- Solution sequency of two-letter combinations, known as digrams
- In our **ciphertext**, the most common **digram is ZW**, which appears **three times**. So we make the correspondence of **Z with t and W with h**.
- Then look at frequency of trigram ZWP
- Continued analysis of frequencies plus trial and error should easily yield a solution

- Substitution Techniques
- Monoalphabetic Ciphers

Ciphertext

UZQSOVUOHXMOPVGPOZPEVSGZWSZOPFPESXUDBMETSXAIZ VUEPHZHMDZSHZOWSFPAPPDTSVPQUZWYMXUZUHSX EPYEPOPDZSZUFPOMBZWPFUPZHMDJUDTMOHMQ



UZQSOVUOHXMOPVGPOZPEVSGZWSZOPFPESXUDBMETSXAIZ

ta e e te a that e e a VUEPHZHMDZSHZOWSFPAPPDTSVPQUZWYMXUZUHSX

et tathaeee ae th t a EPYEPOPDZSZUFPOMBZWPFUPZHMDJUDTMOHMQ

e e e tat e the t

Plain text

it was disclosed yesterday that several informal but direct contacts have been made with political representatives of the viet cong in moscow

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- Substitution Techniques
- Monoalphabetic Ciphers
- For Cipher text

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Plain text is

it was disclosed yesterday that several informal but direct contacts have been made with political representatives of the viet cong in moscow

- Substitution Techniques
- Monoalphabetic Ciphers
- Monoalphabetic ciphers are easy to break because they reflect the frequency data of the original alphabet.
- A countermeasure is to provide multiple substitutes, known as homophones, for a single letter.
- For example, the letter **e** could be assigned a **number of different cipher symbols**, such as 16, 74, 35, and 21, with each homophone assigned to a letter in **rotation or randomly**.
- If the
- number of symbols assigned to each letter is proportional to the relative frequency of that letter,

then **single-letter frequency** information is completely **obliterated**.

- Substitution Techniques
- Monoalphabetic Ciphers
- The great mathematician Carl Friedrich Gauss believed that he had devised an unbreakable cipher using homophones.
 - However, even with homophones, each element of plaintext affects only one element of ciphertext, and
 - multiple-letter patterns (e.g., digram frequencies) still survive in the ciphertext,
- making cryptanalysis relatively straightforward.
- Two principal methods are used in substitution ciphers to lessen the extent to which the structure of the plaintext survives in the ciphertext:
 - One approach is to encrypt multiple letters of plaintext, and
 - the other is to use multiple cipher alphabets.

- Substitution Techniques
- Monoalphabetic Ciphers
- Two principal methods are used in substitution ciphers to lessen the extent to which the structure of the plaintext survives in the ciphertext:
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 - the other is to use multiple cipher alphabets.

Playfair Cipher

- Substitution Techniques
- Playfair Cipher
- The best-known multiple-letter encryption cipher is the Playfair,
- which treats digrams in the plaintext as single units and translates these units into ciphertext digrams
- invented by Charles Wheatstone in 1854 and popularized by Lord Playfair.
- The Playfair cipher is more complex than a monoalphabetic substitution cipher and offers better security.
- based on 5 x 5 matrix of letters
- constructed using a keyword.
- Here is an example, step by step method

- Substitution Techniques
- Playfair Cipher

Step 1.

Create a **5x5 Key Square**: The first step is to generate a 5x5 grid of letters using a keyword (**minus the duplicate**). (left to right and top to bottom)

The **keyword** is written into the grid first, without repeating letters. (Left to right)

After the **keyword, the rest of the alphabet is filled in**, omitting the letter "J" (which is combined with "I" to fit in the 25-letter grid)

Keyword - MONARCHY

М	0	N	Α	R
С	Н	Υ	В	D
E	F	G	I	K
L	Р	Q	S	Т
U	V	W	Х	Z

- Substitution Techniques
- Playfair Cipher

Step 2.

Prepare the Plaintext: The plaintext is **divided into pairs** of letters (**digraphs**).

If a pair contains **two identical letters**, **Of** if there's an **odd number of letters**, an extra filler letter like "X" is added.

"I" and "J" are treated as the same letter.

Example:

Plaintext: BALLOON

Split into digraphs: BA LL OO N

After handling double letters and odd length: BA LX LO ON

М	0	N	Α	R
С	Н	Υ	В	D
E	F	G	I	K
L	Р	Q	S	Т
U	V	W	Х	Z

- Substitution Techniques
- Playfair Cipher

Step 3.

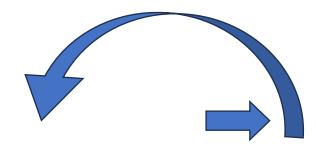
Encryption Rules:

Repeating plaintext letters that are in the same pair are separated with a filler letter, such as x,

balloon □ treated as **BA LX LO ON**

Rules

Same row. Two plaintext letters that fall in the same row of the matrix are each replaced by the letter to the right, with the first element of the row circularly following the last. For example, $ar \square$ is encrypted as RM



M	0	N	A	R
С	Н	Υ	В	D
Е	F	G	I	K
L	Р	Q	S	Т
U	V	W	Х	Z

- Substitution Techniques
- Playfair Cipher

Same column: Two plaintext letters that fall in the same column are each replaced by the letter beneath, with the top element of the column circularly following the last. For example, $\underline{\mathbf{mu}} \square$ is encrypted as $\underline{\mathbf{CM}}$.

Different row and column:

letters form a rectangle, replace them with the letters on the same row, but at the opposite corners of the rectangle.

Otherwise, each plaintext letter in a pair is **replaced by the letter that lies in its own row** and the column occupied by the other plaintext letter.

Thus,

hs BP and

ea \square **IM** (or JM, as the encipherer wishes).

If the

M	0	Ζ	Α	R
C	Ι	Υ	В	D
E	F	G	I	K
L	Р	Q	S	Т
U	V	W	Х	Z

M	0	N	Α	R
С	I	Υ	B	D
Ε	F	G	Ι	K
L	P	Q	S	Т
U	V	W	Х	Z

<mark>∠</mark>	0	N	A	R
C	I	Υ	В	D
E	F	G	_	K
L	Р	Q	S	Т
U	V	W	Χ	58 Z

- Substitution Techniques
- Playfair Cipher

Step 3.

Encryption Rules:

Repeating plaintext letters that are in the same pair are separated with a filler letter, such as x,

balloon ☐ treated as BA LX LO ON

M	0	N	A	R
C	H	Υ	B	D
E	F	G	I	K
L	P	Q	S	Т
U	V	W	Х	Z

<mark>✓</mark>	0	Z	A	R
C	I	Υ	В	D
E	F	G		К
L	Р	Q	S	Т
U	V	W	X 51	Z

- Substitution Techniques
- Playfair Cipher

Step 4.

Encrypt the Digraphs

Digraph 1: BA

B is in row 2, column 4.

A is in row 1, column 4.

Since they are in the **same column**, move down:

B -> I, A -> B.

Encrypted digraph: IB

М	0	N	A	R
С	H	Υ	B	D
E	F	G	_	K
L	Р	Q	S	Т
U	V	W	Х	Z

- Substitution Techniques
- Playfair Cipher

Step 4.

Encrypt the Digraphs

Digraph 2:

L is in row 4, column 1.

X is in row 5, column 4.

They form a rectangle, so swap the columns: L -> S, X -> U. Encrypted diagram SU

Digraph 3:

LO is in row 4, column 1.

O is in row 1, column 2.

They form a **rectangle**, so swap the columns: L -> P, O -> M. Encrypted digraph: PM

Digraph 4: ON

O is in row 1, column 2.

N is in row 1, column 3.

Since they are in the **same row**, move right one position: **O -> N, N -> A. Encrypted digraph: NA**

М	0	N	Α	R
С	Н	Υ	В	D
Ε	F	G	-	K
L	Р	Q	S	Т
U	V	W	Х	Z

М	0	Z	Α	R
С	Ι	Υ	В	D
Ε	F	G	I	K
L	Р	Q	S	Т
U	V	W	Х	Z

М	O	N	A	R
С	Ι	Υ	В	D
E	F	G	I	K
L	Р	Q	S	Т
U	V	W	Х	Z

- Substitution Techniques
- Playfair Cipher
- Final Ciphertext:
- After encrypting all the digraphs, the final ciphertext for
- BALLOON becomes: IBSU PMNA

- Substitution Techniques
- Playfair Cipher
- Decryption:
- Decryption follows the same **steps but in reverse**:
- If the digraph letters are in the same row, move left one position.
- If they are in the same column, move up one position.
- If they form a rectangle, reverse the corner swapping process.

The Playfair cipher provides **greater security** than monoalphabetic ciphers because it operates on digraphs, making it harder to perform frequency analysis. However, with modern techniques, it can still be broken relatively easily compared to more advanced ciphers.

- Substitution Techniques
- Playfair Cipher
- identification of individual digrams is more difficult.
- the **relative frequencies** of individual letters exhibit a **much greater range** than that of digrams, **making frequency analysis much more difficult**.
- Playfair cipher was for a long time considered unbreakable.
- It was used as the standard **field system** by the **British Army** in **World War I** and enjoyed considerable use by the U.S. Army and other Allied forces during **World War II.**
- Strength of the Playfair Cipher
- The Playfair cipher is a great advance over simple monoalphabetic ciphers. For one thing, whereas there are only 26 letters, there are 26 * 26 = 676 digrams
- easy to break, -- much of the structure of the plaintext language is intact.

Thank You