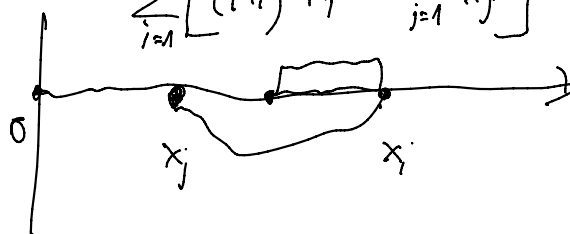


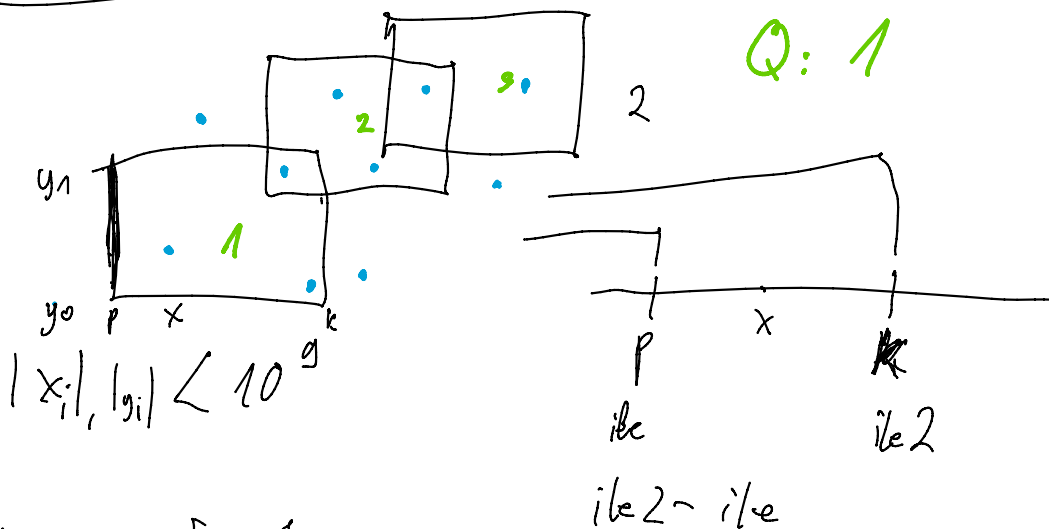
$$\sum_{i=0}^{n-1} \sum_{j=0}^{i-1} (x_i - x_j) \quad O(n^2)$$

$$\sum_{i=1}^n \left[(i-1) x_i - \sum_{j=1}^{i-1} x_j \right] \quad O(n^2)$$


$sum = 0$

$i-1 \text{ razy}$

for $i=1$ $i \leq N$
 res += $(i-1) \cdot x[i] - sum$; $O(n)$
 $sum += x[i]$;



U: $sum[p.y]++$

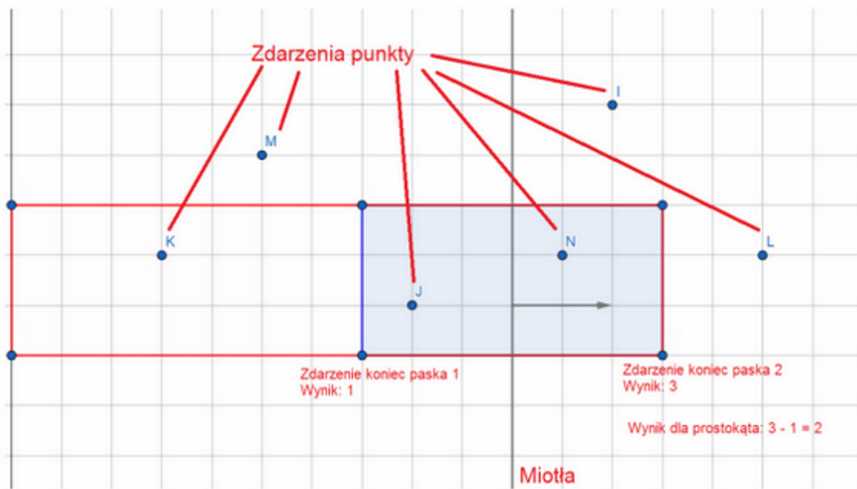
Q: for $i=y_0$ $i \leq y_1$
 res += $sum[i]$;

$V(x): tab[x]++$

$Q(a, b): sum(tab[a], tab[b])$

$O(n \log n)$

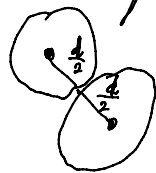
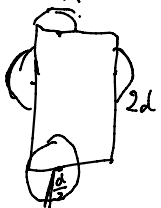




- $O(n^2)$
- usunąć z kolejki i seta pkt takie że $x < x_m - d$
 - spr. punkty o y miesz (y_m-d, y_m+d) i zaktualizuj d jeśli trzeba
 - wrzuć $p(x_m, y_m)$ na kolejke i set



$2d \times 3d$

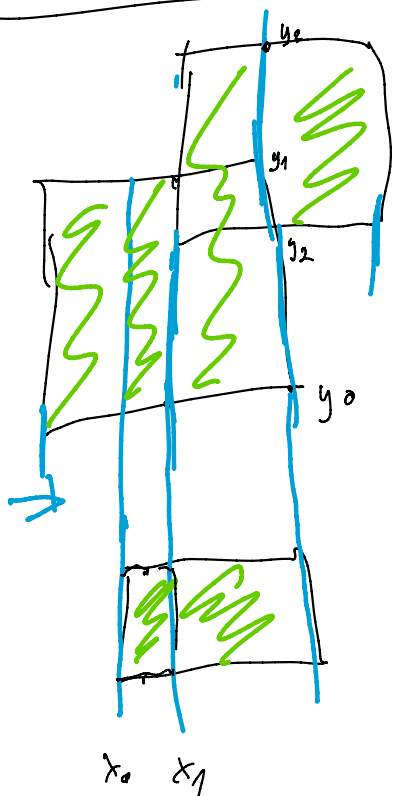


$$\frac{\pi d^2}{4} \cdot k \leq 6d^2$$

$$k \leq \frac{24d^2}{\pi}$$

$$k \leq \frac{24}{\pi} < \frac{24}{3} = 8$$

$$O(n \log n)$$



$$(b-a) * \text{sum } y$$

val - ilość prost na danym przedziale bezog
sum - ile do wyniku

if (val[x] > 0)
sum[x] = en - beg + 1;

else
sum[x] = sum[2x] + sum[2x+1];

nie trzeła (azy



$$x_0 = x_1$$

$$(x_1 - x_0) \cdot \text{sum}[1]$$

$$(x_1 - x_0) \cdot d_1 + (x_1 - x_0) \cdot d_2 = (x_1 - x_0) (d_1 + d_2)$$