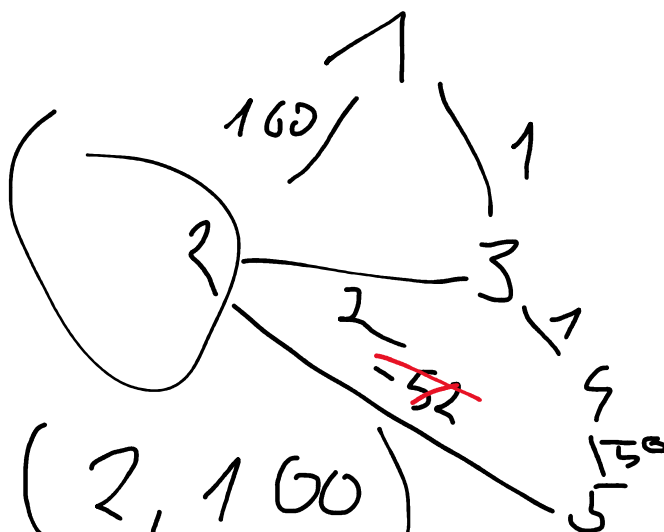




$$\text{nei}[a][b] = 1$$

$$\text{cost}[a][b] = 3$$

$$g[a].PB(\{b, 3\})$$



$$(2, 5), \quad \underline{(2, 160)}$$

↓

$$\underline{O(E \log V)}$$

$$p-g \left\langle \underset{\text{cost}}{p \langle i, i \rangle}, \underset{\text{node}}{vec \langle p \langle i, i \rangle \rangle}, \text{gather} \langle p \langle i, i \rangle \rangle \right\rangle$$

$$\text{set} \left\langle p \langle i, i \rangle \right\rangle$$

```

1 function Dijkstra(Graph, source):
2     dist[source] ← 0                                // Initialization
3
4     create vertex priority queue Q
5
6     for each vertex v in Graph.Vertices:
7         if v ≠ source
8             dist[v] ← INFINITY                    // Unknown distance from source to v
9             prev[v] ← UNDEFINED                    // Predecessor of v
10
11         Q.add_with_priority(v, dist[v])
12
13
14     while Q is not empty:                            // The main loop
15         u ← Q.extract_min()                        // Remove and return best vertex
16         for each neighbor v of u:                  // Go through all v neighbors of u
17             alt ← dist[u] + Graph.Edges(u, v)
18             if alt < dist[v]:
19                 dist[v] ← alt
20                 prev[v] ← u
21             Q.decrease_priority(v, alt)
22
23     return dist, prev

```

B F
I

1

$V-1$

$1 \sim n$

```

// funkcja BF zwraca true jeśli znalazł najkrótsze ścieżki
// lub false jeśli istnieje ujemny cykl
bool BF () {
    for (int i = 1; i <= n; i++)
        D[i] = INF;
    D[x] = 0;
    for (int I = 1; I < n; I++) {
        for (int v = 1; v <= n; v++)
            for (int i = 0; i < sasiedzi[v].size(); i++)
                if (D[v] + waga[v][i] < D[sasiedzi[v][i]])
                    D[sasiedzi[v][i]]
                        = D[v] + waga[v][i];
    }
    //sprawdanie, czy istnieje ujemny cykl
    for (int v = 1; v <= n; v++)
        for (int i = 0; i < sasiedzi[v].size(); i++)
            if (D[v] + waga[v][i] < D[sasiedzi[v][i]])
                return false;
    return true;
}

```

$d(x, x) = 0$

$$\begin{cases} d(x, x) = 0 \\ d(x, y) = d(y, x) \end{cases}$$

$$d(a, b) \leq d(a, c) + d(c, b)$$

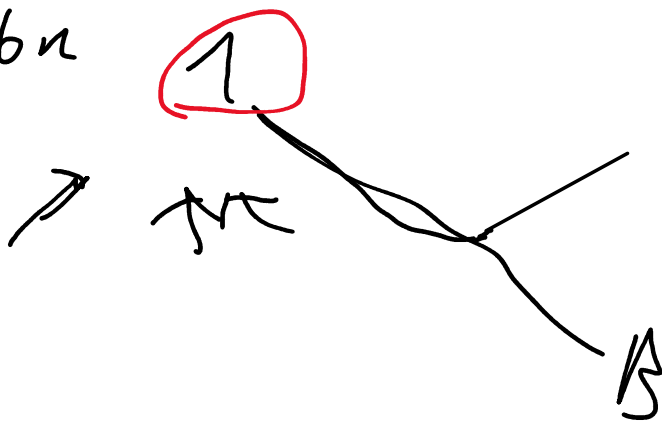
$$dp[a][b]$$

```
for(int u = 1; u <= n; u++)
    for(int x = 1; x <= n; x++)
        for(int y = 1; y <= n; y++)
            D[x][y] = min(D[x][y], D[x][u] + D[u][y]);
```

$$\min(d[x][x]) \quad O(n^3)$$

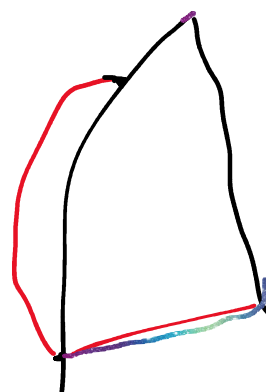
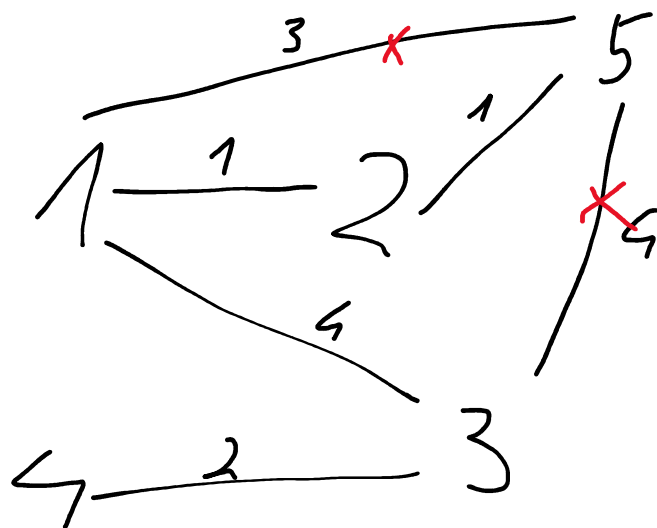
$$-C \quad \begin{array}{c} \uparrow \\ \text{---} \\ \downarrow \end{array} \quad O(|V||E| \cdot \log n)$$

Pubn



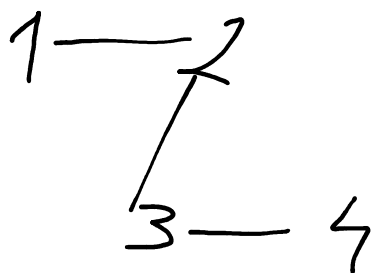
bellman ford - spr yjeung, uypisz odd

Naj cykl - $F \cup d[a][v]$



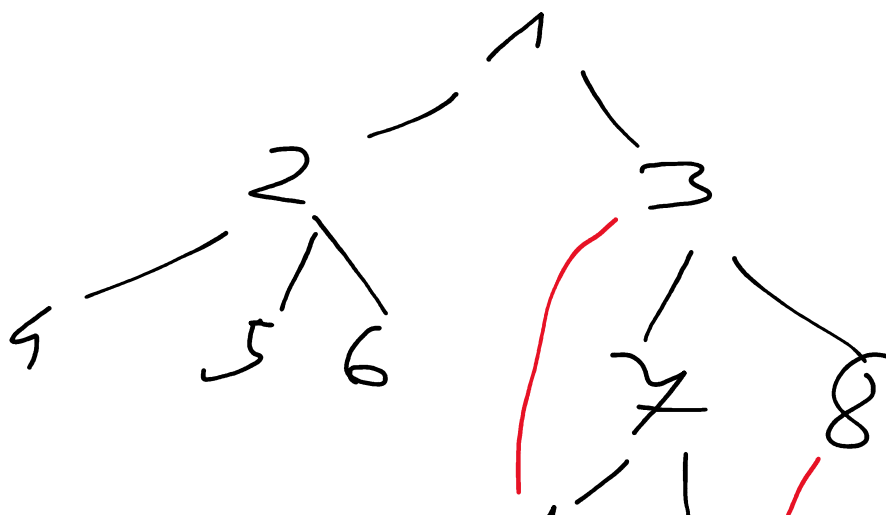
$O(E \log E)$

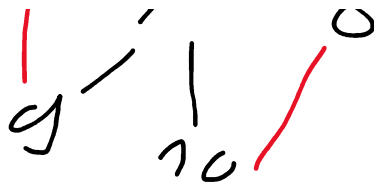
$F \cup$



MST - od $Kruskala$

Nawetnica





$$O(E \log E + V \log V + E \log V)$$

$Cy \leq 1$

$$A = \frac{\sum a_i}{n} = \frac{a_1 + a_2 + \dots + a_n}{n}$$

$$nA = \sum a_i$$

$$\sum a_i - nA = 0$$

$$a_1 + \dots + a_n - nA = 0$$

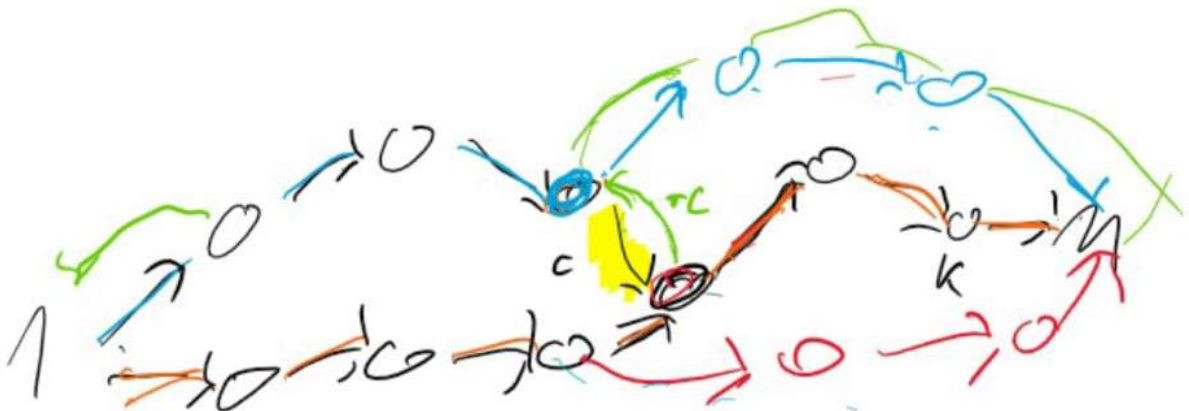
$$a_1 - A + a_2 - A + \dots + a_n - A = 0$$

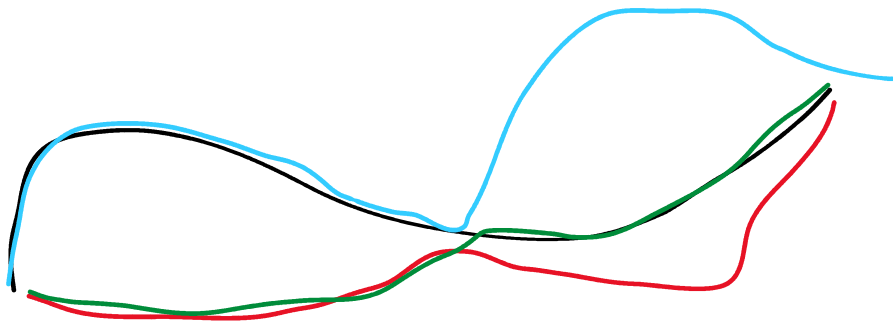
$$\sum a_i - A = 0$$

while [abs (R-L) > eps]

while (ct-- > 0)

ct = 50





.