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HW5b
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Part A:

15.3.2

10,000/999 * (999 + 10,000) \appro 10 * 11,000 = 110,000 disk I/O's

15.3.3

a).
$$10,000/(M-1) * (M-1 + 10,000) \le 100,000 \Leftrightarrow M >= 1112$$

16.2.6

- a). $\phi_{b+c\to x, c+d\to y}(R(a,b,c) \text{ join } S(b,c,d,e)) = \phi_{b+c\to x, c+d\to y}(\phi_{b,c}(R(a,b,c)) \text{ join } \phi_{b,c,d}(S(b,c,d,e))$ because b,c,d are the input attributes and join attributes.
- b). $\pi_{a, b, a+d \to z}(R(a,b,c) \text{ join } S(b,c,d,e)) = \pi_{a, b, a+d \to z}(\pi_{a,b,c}(R(a,b,c)) \text{ join } \pi_{b,c,d}(S(b,c,d,e))) = \pi_{a, b, a+d \to z}(R(a,b,c) \text{ join } \pi_{b,c,d}(S(b,c,d,e)))$ Because a,c,c,d are the input attributes and join attributes.

Part B:

In linear cost model, let c = average seek time, so, cost of reading n blocks = n*c In affine cost model, let c = average seek time for the fist block and c' = average rotational latency for each additional block, so, cost of reading n (contiguous) blocks = c + (n-1)*c'

- 1). B(R) * c = 1000c
- 2). c + (B(R) 1)c' = c + 999c'
- 3). This is an ill-posed question because if W.c is the primary key, then W.c should not have duplicates. Anyway, I will still do it. The general approach is the same for subquestion 3) through 7).
- a). 10 blocks
- b). 1,000 records
- 4).
- a). 1 block
- b). 1 record
- 5).
- a). 1 block
- b). 100 records

- 6).
- a). 50 records
- b). max(500, 1000) = 1000 records
- 7).
- a). 10,000/(3*10) = 333 records
- b). 10,000/3 = 3333 records
- 8). 0 because S and U don't have any shared attributes
- 9). 100*10,000/10 = 100,000 records
- 10). 1,000,000*10,000/10,000 = 1,000,000 records
- 11). 1,000,000/3 * 10,000 / 200 = 17,000,000records
- 12). I don't know how to solve this question.