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HW11

Part 1:

Review question 1:

a).  $\text{OpenLate}(r) \leftarrow \text{Restaurant}(r, d, o, c) \text{ AND } d = \text{'Friday'} \text{ AND } o > 9:00 \text{ AM}$   
 $\text{OpenLate}(r) \leftarrow \text{Restaurant}(r, d, o, c) \text{ AND } d = \text{'Saturday'} \text{ AND } o > 9:00 \text{ AM}$

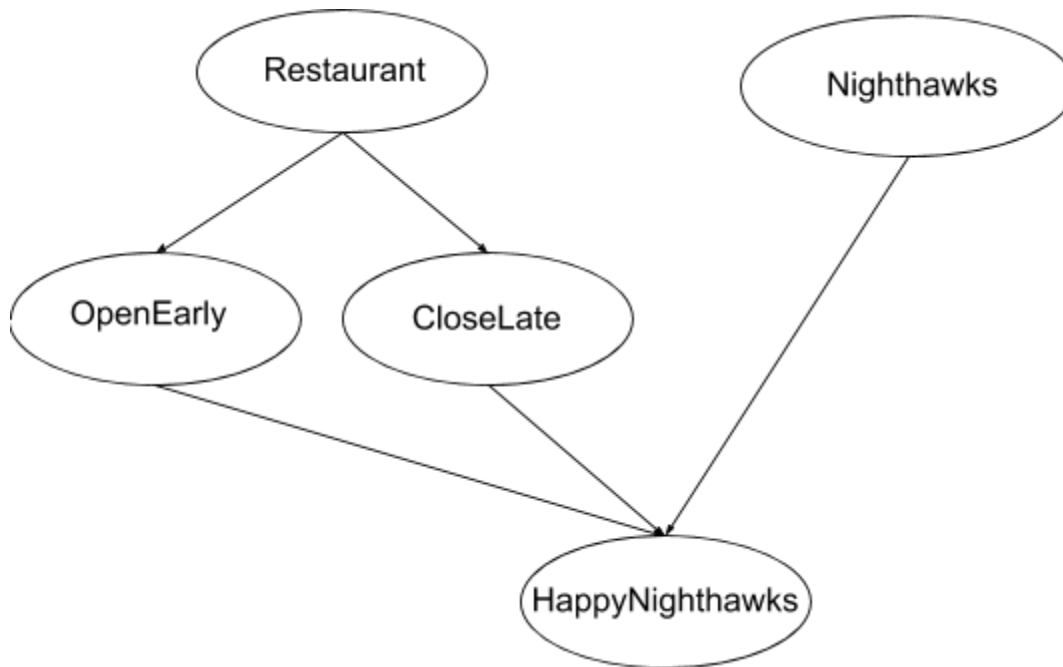
b).  $\text{CloseLate}(r) \leftarrow \text{Restaurant}(r, d, o, c) \text{ AND } d = \text{'Friday'} \text{ AND } c > 9:00 \text{ PM}$   
 $\text{CloseLate}(r) \leftarrow \text{Restaurant}(r, d, o, c) \text{ AND } d = \text{'Saturday'} \text{ AND } c > 9:00 \text{ PM}$

$\text{OpenEarly}(r) \leftarrow \text{Restaurant}(r, d, o, c) \text{ AND } d = \text{'Saturday'} \text{ AND } o < 8:00 \text{ AM}$   
 $\text{OpenEarly}(r) \leftarrow \text{Restaurant}(r, d, o, c) \text{ AND } d = \text{'Sunday'} \text{ AND } o < 8:00 \text{ AM}$

$\text{HappyNighthawks}(p) \leftarrow \text{Nighthawks}(p, r) \text{ AND } \text{CloseLate}(r) \text{ AND } \text{OpenEarly}(r)$

Review question 2:

a).



b). No, they will be no happy nighthawks because the OpenEarly relation is empty (i.e. no restaurant in Restaurant opens before 8:00 AM on Saturdays and Sundays).

c). Yes, now Tom is a happy nighthawk!

	OpenEarly	CloseLate	HappyNighthawks
Before	{}	{"UNOs 360"}	{}
After	{"Magnolia Cafe"}	{"UNOs 360", "Magnolia Cafe"}	{Tom}

### Review Question 3:

To simplify the computation notation-wise, I will use integer numbers to represent UT employees as follows:

Employee Name	Corresponding Integer
Fenves	0
Glodbart	1
Wood	2
Fussell	3
Beckner	4
Tewfik	5
Miranker	6
Mok	7
Alcook	8
Ghosh	9

Then  $\text{Arc}(x,y)$  can be represented as  $\{(0,1), (0,2), (1,3), (1,4), (2,5), (3,6), (3,7), (4,8), (5,9)\}$ . The following sequence shows what relation  $\text{Reaches}(x,y)$  contains at each iteration. Tuples in boldface are new ones generated at that iteration.

$t = 0$ :  $\text{Reaches}(x,y) = \{\}$

$t = 1$ :  $\text{Reaches}(x,y) = \{(\mathbf{0,1}), (\mathbf{0,2}), (\mathbf{1,3}), (\mathbf{1,4}), (\mathbf{2,5}), (\mathbf{3,6}), (\mathbf{3,7}), (\mathbf{4,8}), (\mathbf{5,9}), (\mathbf{0,3}), (\mathbf{0,4}), (\mathbf{0,5}), (\mathbf{1,6}), (\mathbf{1,7}), (\mathbf{1,8}), (\mathbf{2,9})\}$

$t = 2$ :  $\text{Reaches}(x,y) = \{(0,1), (0,2), (1,3), (1,4), (2,5), (3,6), (3,7), (4,8), (5,9), (0,3), (0,4), (0,5), (1,6), (1,7), (1,8), (2,9), (\mathbf{0,6}), (\mathbf{0,7}), (\mathbf{0,8}), (\mathbf{0,9})\}$

$t = 3$ :  $\text{Reaches}(x,y) = \{(0,1), (0,2), (1,3), (1,4), (2,5), (3,6), (3,7), (4,8), (5,9), (0,3), (0,4), (0,5), (1,6), (1,7), (1,8), (2,9), (0,6), (0,7), (0,8), (0,9)\}$  (no new tuples generated, terminate!)

Part 2:

17.4.1

a). <START T>; <T, A, 5, 15>; <T, B, 10, 25>; <COMMIT T>

b). <START T>; <T, B, 10, 15>; <T, A, 5, 20>; <COMMIT T>

c). <START T>; <T, A, 5, 11>; <T, B, 10, 12>; <COMMIT T>

17.4.3

b).

- 1) U is identified as a committed transaction; T is identified as an incomplete transaction.
- 2) Redo U in an earliest-first order:
  - write 21 to B on disk
  - write 41 to D on disk
- 3) Undo T in a latest-first order:
  - write 30 to C on disk
  - write 10 to A on disk

d).

- 1) U, T are both identified as committed transactions.
- 2) Redo U in an earliest-first order:
  - write 21 to B on disk
  - write 41 to D on disk
- 3) Redo T in an earliest-first order:
  - write 11 to A on disk
  - write 31 to C on disk
  - write 51 to E on disk

17.4.4

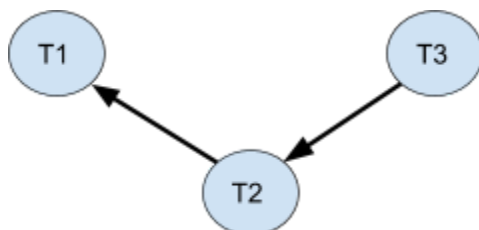
b). Value A, B, C, D *might* appear on disk

18.1.1

r(A); r(B); w(B); w(C); w(D); w(E)

18.2.4

a).



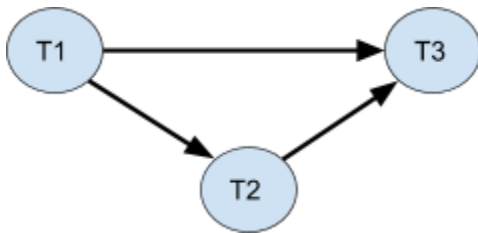
i.

ii. Yes, T3; T2; T1

iii. No, because

T1; T2; T3	No, r <sub>2</sub> (A) will read w <sub>1</sub> (A) while r <sub>2</sub> (A) originally reads initial A.
T1; T3; T2	No, r <sub>2</sub> (C) will read w <sub>1</sub> (C) while r <sub>2</sub> (C) originally reads initial C
T2; T1; T3	No, r <sub>3</sub> (B) will read w <sub>2</sub> (B) while r <sub>3</sub> (B) originally reads initial B
T2; T3; T1	No, r <sub>3</sub> (B) will read w <sub>2</sub> (B) while r <sub>3</sub> (B) originally reads initial B
T3; T1; T2	No, r <sub>2</sub> (C) will read w <sub>1</sub> (C) while r <sub>2</sub> (C) originally reads initial C
T3; T2; T1	Conflict-equivalent to S

b).



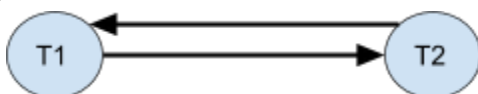
i.

ii. Yes, T1; T2; T3

lii. No, because

T1; T2; T3	Conflict-equivalent to S
T1; T3; T2	No, r <sub>3</sub> (C) will read initial C while r <sub>3</sub> (C) originally reads w <sub>2</sub> (C)
T2; T1; T3	No, r <sub>2</sub> (B) will read initial B while r <sub>2</sub> (B) originally reads w <sub>1</sub> (B)
T2; T3; T1	No, same as above
T3; T1; T2	No, r <sub>3</sub> (C) will read initial C while r <sub>3</sub> (C) originally reads w <sub>2</sub> (C)
T3; T2; T1	No, same as above

d).



i.

ii. No, there is a cycle in the precedence graph.

iii. No, because

T1; T2	No, r_1(B) will read w_1(B) while r_1(B) originally reads w_2(B)
T2; T1	No, same as above

### 18.3.3

b). No request will get delayed

d). w\_2(B) will be delayed. It (T2) will be allowed to resume after r\_1(B) is finished and thereafter lock on B is released.