

HW7

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Part A

15.3.2

$10,000/999 * (999 + 10,000) \approx 10 * 11,000 = 110,000$ disk I/O's

15.3.3

a). $10,000/(M-1) * (M-1 + 10,000) \leq 100,000 \Leftrightarrow M \geq 1112$

15.4.2

a). set union: $3 * (B(R) + B(S)) = 3 * 20,000 = 60,000$ disk I/O's

b). simple sort-join: $5 * (B(R) + B(S)) = 5 * 20,000 = 100,000$ disk I/O's

c). Merge-join: $3 * (B(R) + B(S)) = 3 * 20,000 = 60,000$ disk I/O's

15.4.3

We can modify phase 1 such that some blocks are not written back to secondary storage, which was written out and read back to main-memory in the original two-pass algorithm.

Concretely, we can

- 1). Set a smaller chunk size to read into main-memory in phase 1;
- 2). Perform the same sort process as before;
- 3). For the last chunk, do not write it back to secondary storage. Just leave it in main memory;
- 4). In phase 2, do the same thing as before but acknowledging the fact that what was originally the last sublist is now already in main-memory.

Part B

16.4.1

a). $T(W \text{ join } X \text{ join } Y \text{ join } Z) = 100 * 200 * 300 * 400 / (60 * 100 * 50) = 8,000$

b). $T(\text{select}_{\{c=20\}}(Y) \text{ join } Z) = T(\text{select}_{\{c=20\}}(Y) * T(Z) / \max\{V(\text{select}_{\{c=20\}}(Y), d), V(Z, d)\}) = T(Y)/V(Y, c) * T(Z) / \max\{V(Y, d), V(Z, d)\} = 300/50 * 400 / 50 = 48$

c). $T(X \text{ join}_{\{X.c < Y.c\}} Y) = T(\text{select}_{\{X.c < Y.c\}}(X \text{ CartesianProduct } Y)) = T(X \text{ CartesianProduct } Y) / 3 = T(X) * T(Y) / 3 = 200 * 300 / 3 = 20,000$

Part C

1.

a). Restricting to only left-deep trees:

Set	{W}	{X}	{Y}	{Z}		
Cost	0	0	0	0		
Best Plan	W	X	Y	Z		
Set	{W, X}	{W, Y}	{W, Z}	{X, Y}	{X, Z}	{Y, Z}
Cost	0	0	0	0	0	0
Best Plan	W join X	W join Y	W join Z	X join Y	X join Z	Y join Z
Set	{W, X, Y}	{W, X, Z}	{W, Y, Z}	{X, Y, Z}		
Cost	1000/3	1000/3	2400	600		
Best Plan	(W join X) join Y	(W join X) join Z	(Y join Z) join W	(X join Y) join Z		
Set	{W, X, Y, Z}					
Cost	1000 * 4/3					
Best Plan	((W join X) join Y) join Z					

b). Consider all trees:

(red cells are the only different cells in this answer compared to the above table)

Set	{W}	{X}	{Y}	{Z}		
Cost	0	0	0	0		
Best Plan	W	X	Y	Z		
Set	{W, X}	{W, Y}	{W, Z}	{X, Y}	{X, Z}	{Y, Z}
Cost	0	0	0	0	0	0
Best Plan	W join X	W join Y	W join Z	X join Y	X join Z	Y join Z
Set	{W, X, Y}	{W, X, Z}	{W, Y, Z}	{X, Y, Z}		
Cost	1000/3	1000/3	2400	600		
Best Plan	Y join (W join X)	(W join X) join Z	W join (Y join Z)	Z join (X join Y)		
Set	{W, X, Y, Z}					
Cost	1000 * 4/3					
Best Plan	Z join (Y join (W join X))					

2. In this problem, we restrict our attention to left-deep trees.

Query graph looks like (there is no query here but we will draw a link between two nodes if those two relations have shared attributes)

W-----X-----Y-----Z

Copied from the table in Part C::1.a, red strikethrough cells are those eliminated by the query graph because the nodes are not adjacent, effectively reducing the size of the table from 15 entries to 10 entries.

Set	{W}	{X}	{Y}	{Z}		
Cost	0	0	0	0		
Best Plan	W	X	Y	Z		
Set	{W, X}	{W, Y}	{W, Z}	{X, Y}	{X, Z}	{Y, Z}
Cost	0	0	0	0	0	0
Best Plan	W join X	W join Y	W join Z	X join Y	X join Z	Y join Z
Set	{W, X, Y}	{W, X, Z}	{W, Y, Z}	{X, Y, Z}		
Cost	1000/3	1000/3	2400	600		
Best Plan	(W join X) join Y	(W join X) join Z	(Y join Z) join W	(X join Y) join Z		
Set	{W, X, Y, Z}					
Cost	1000 * 4/3					
Best Plan	((W join X) join Y) join Z					