

# **DTS203TC**

# **Design and Analysis of Algorithms**

## **Lecture 4: Lab introduction and Quicksort**

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# Why does sorting matter?

My answer: sorting can help you get what you want much **faster** and **cheaper**

# Videos

- [https://www.bilibili.com/video/BV1Ws411f7aJ/?spm\\_id\\_from=333.337.search-card.all.click](https://www.bilibili.com/video/BV1Ws411f7aJ/?spm_id_from=333.337.search-card.all.click)
- [https://www.bilibili.com/video/BV1ex411e7eb/?spm\\_id\\_from=333.337.search-card.all.click](https://www.bilibili.com/video/BV1ex411e7eb/?spm_id_from=333.337.search-card.all.click)

# Learning outcomes

- Jupyter Notebook
- Insertion Sort
- Merge Sort
- Quicksort

# Jupyter notebook

- We assume you can do Python programming
- Jupyter notebook is very convenient for demonstrating and managing your coursework and projects
- Installation and practice

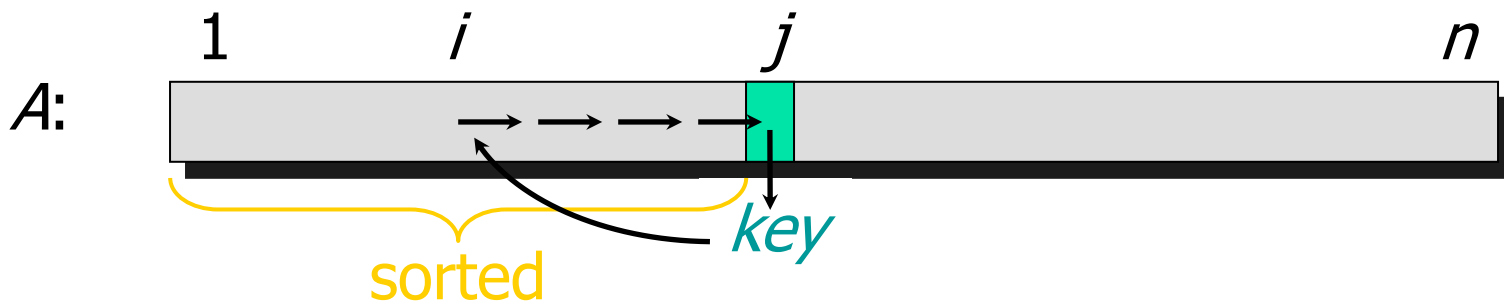
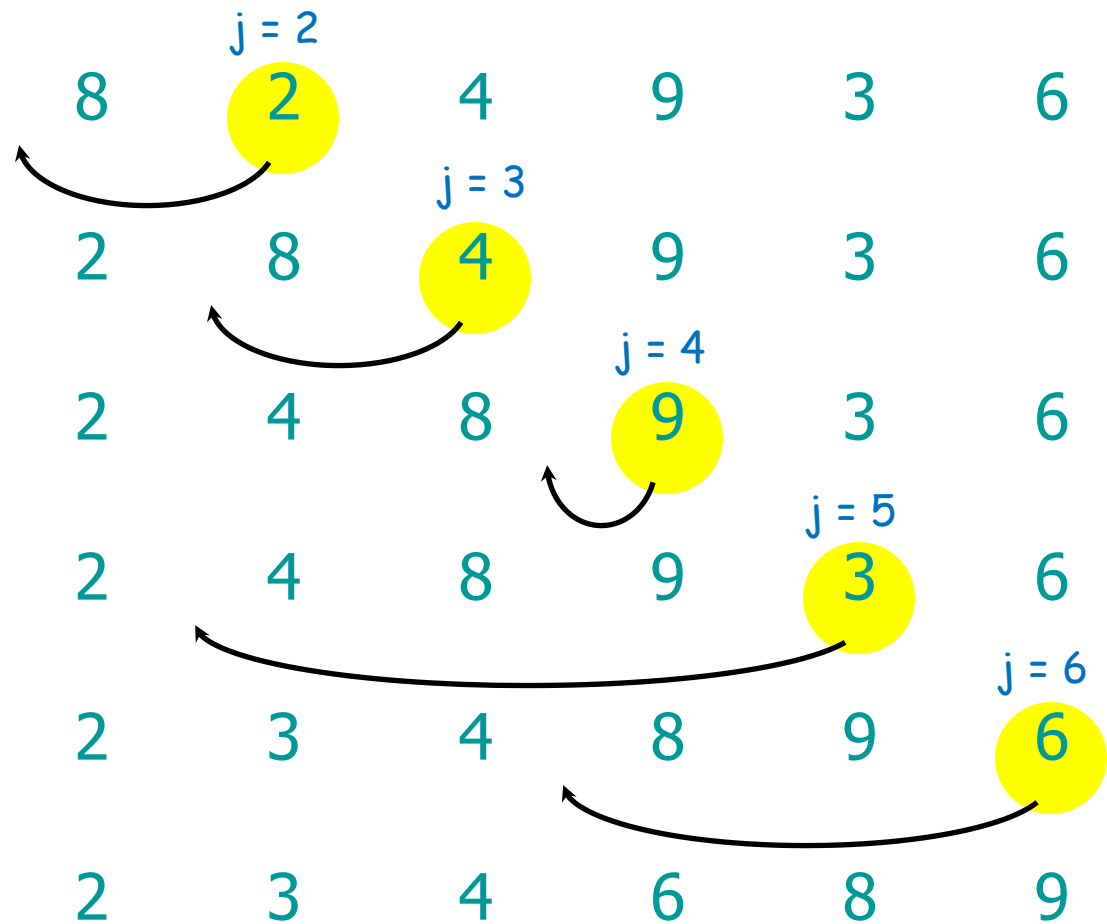
# Insertion Sort

# Insertion sort

## ■ Pseudocode

INSERTION-SORT( $A$ )

```
1  for  $j = 2$  to  $A.length$ 
2     $key = A[j]$ 
3    // Insert  $A[j]$  into the sorted sequence  $A[1..j-1]$ .
4     $i = j - 1$ 
5    while  $i > 0$  and  $A[i] > key$ 
6         $A[i + 1] = A[i]$ 
7         $i = i - 1$ 
8     $A[i + 1] = key$ 
```



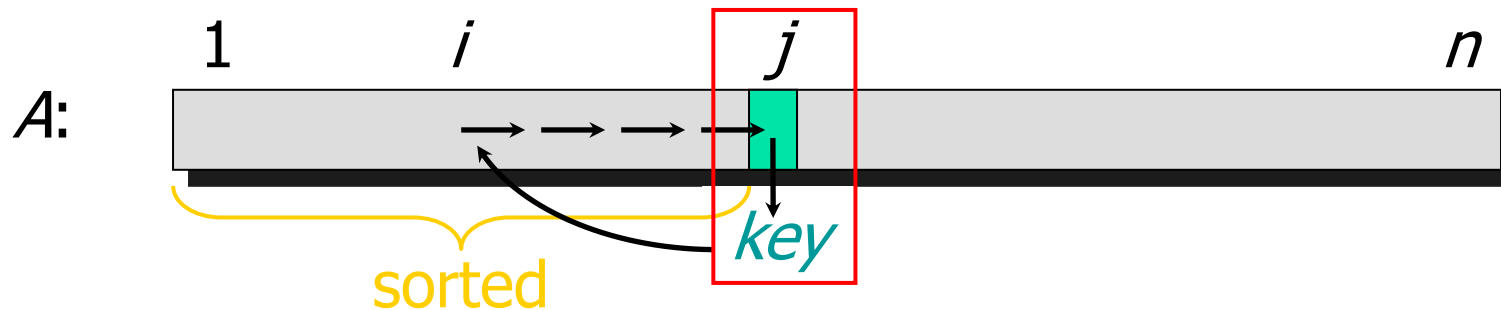
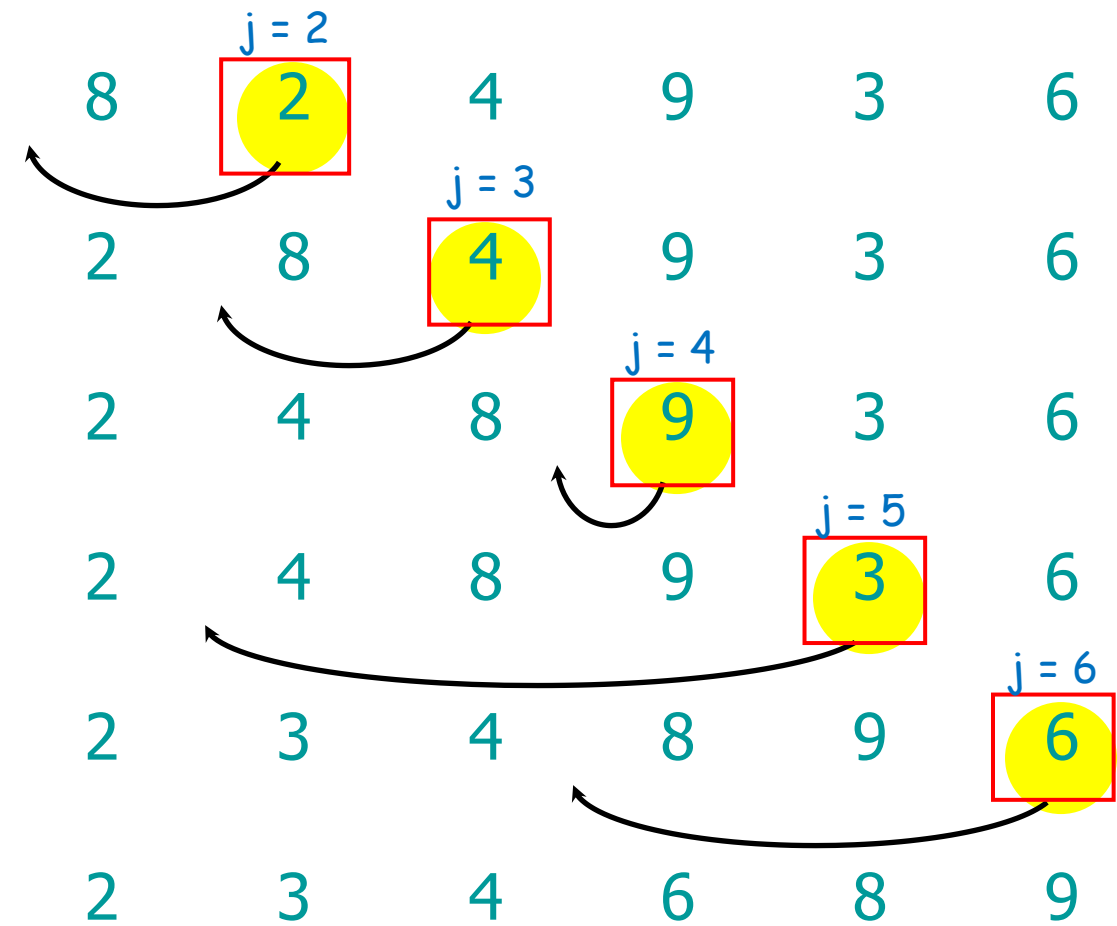


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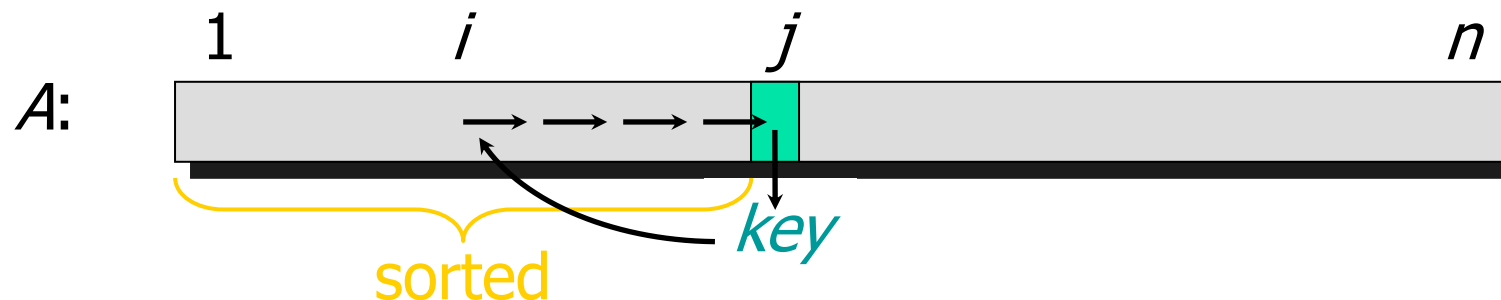
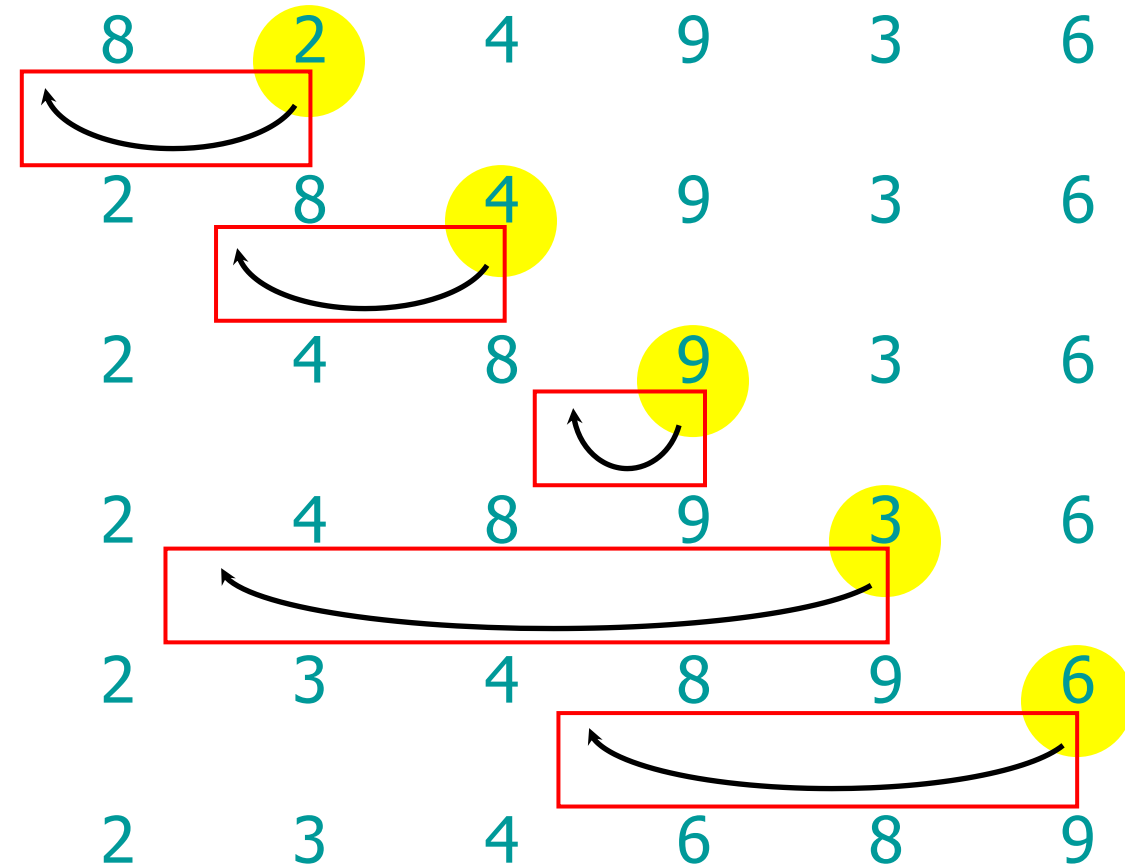


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8       $A[i+1] = key$ 
```

```
def insertion_sort(InputList):
    for j in range(1, len(InputList)):
        key = InputList[j]
        i = j-1
        # Compare the current element with next one
        while (InputList[i] > key) and (i >= 0):
            InputList[i+1] = InputList[i]
            i=i-1
        InputList[i+1] = key

arr = [19,2,31,45,30,11,121,27]

insertion_sort(arr)

print(arr)

[2, 11, 19, 27, 30, 31, 45, 121]
```

# Merge Sort

# Two steps

## 1. Divide

- dividing a sequence of  $n$  numbers into **two** smaller sequences

## 2. Conquer

- merging two sorted sequences of **total length  $n$**

# First step

## Divide

- dividing a sequence of  $n$  numbers into **two** smaller sequences is straightforward

51, 13, 10, 64, 34, 5, 32, 21

we want to sort these 8 numbers,  
divide them into two halves

51, 13, 10, 64, 34, 5, 32, 21

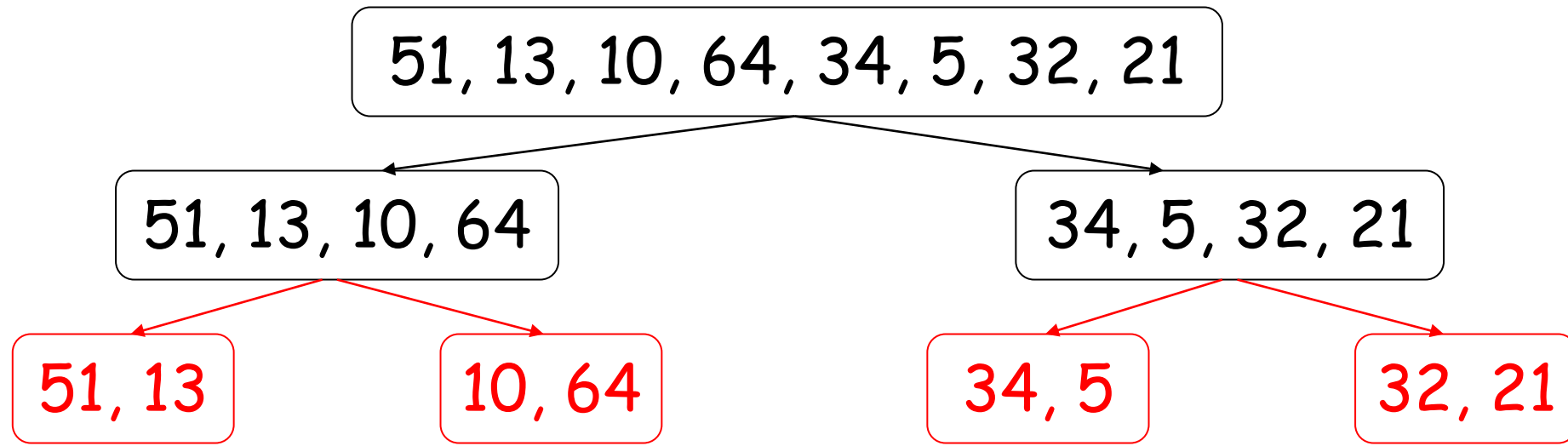
51, 13, 10, 64

34, 5, 32, 21

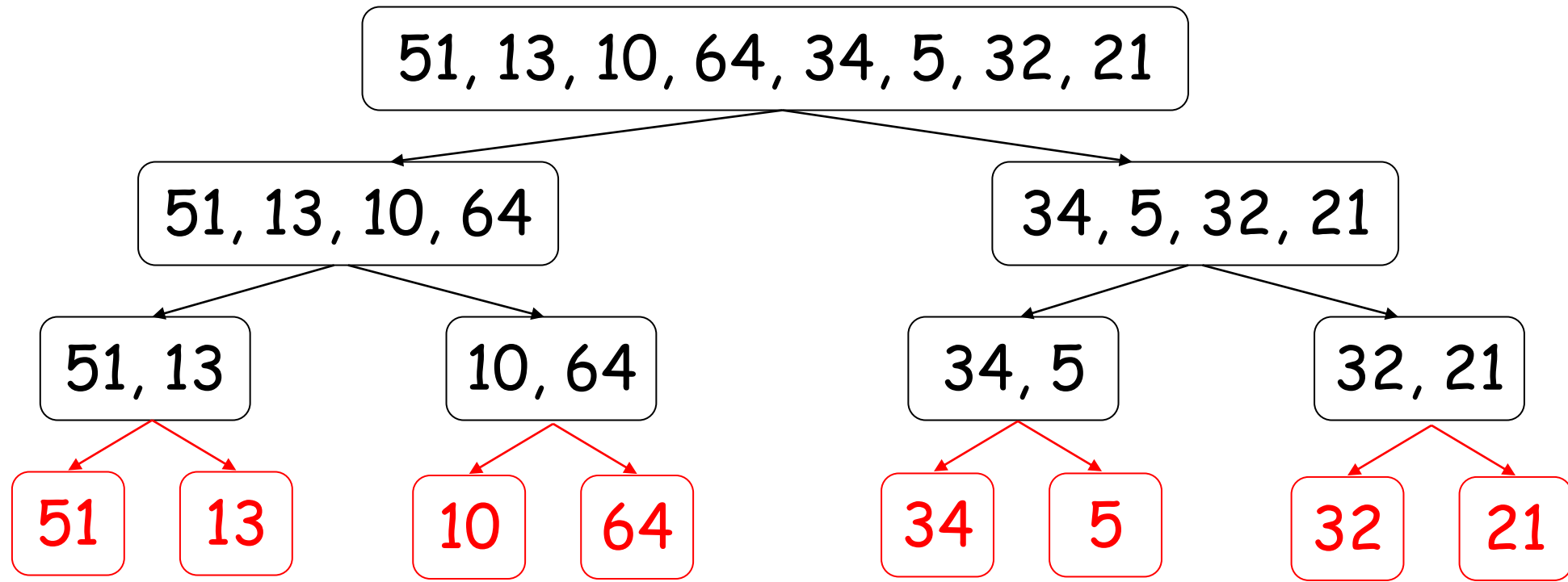
divide these 4  
numbers into  
halves

similarly for  
these 4





further halve each sequence ...  
until we get sequence with only **1** number

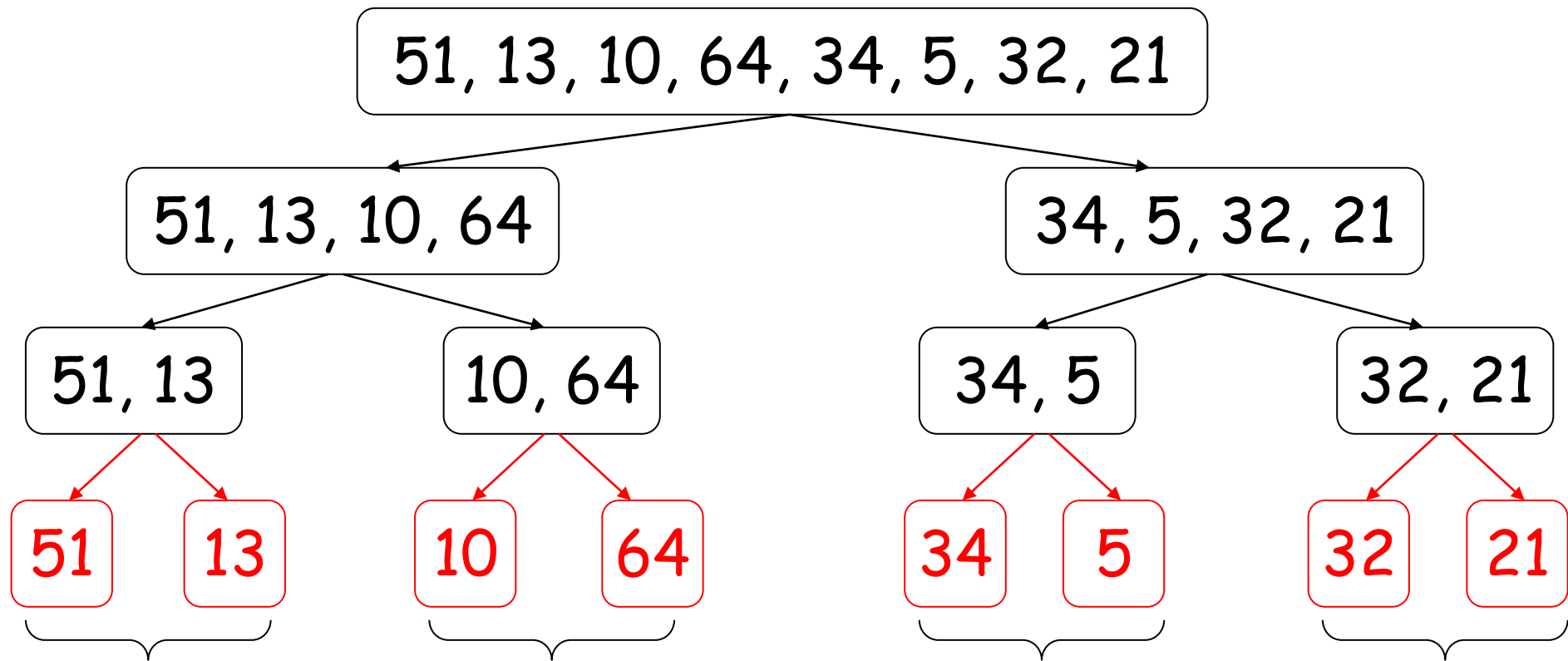


further divide each shorter sequence ...  
until we get sequence with only **1** number

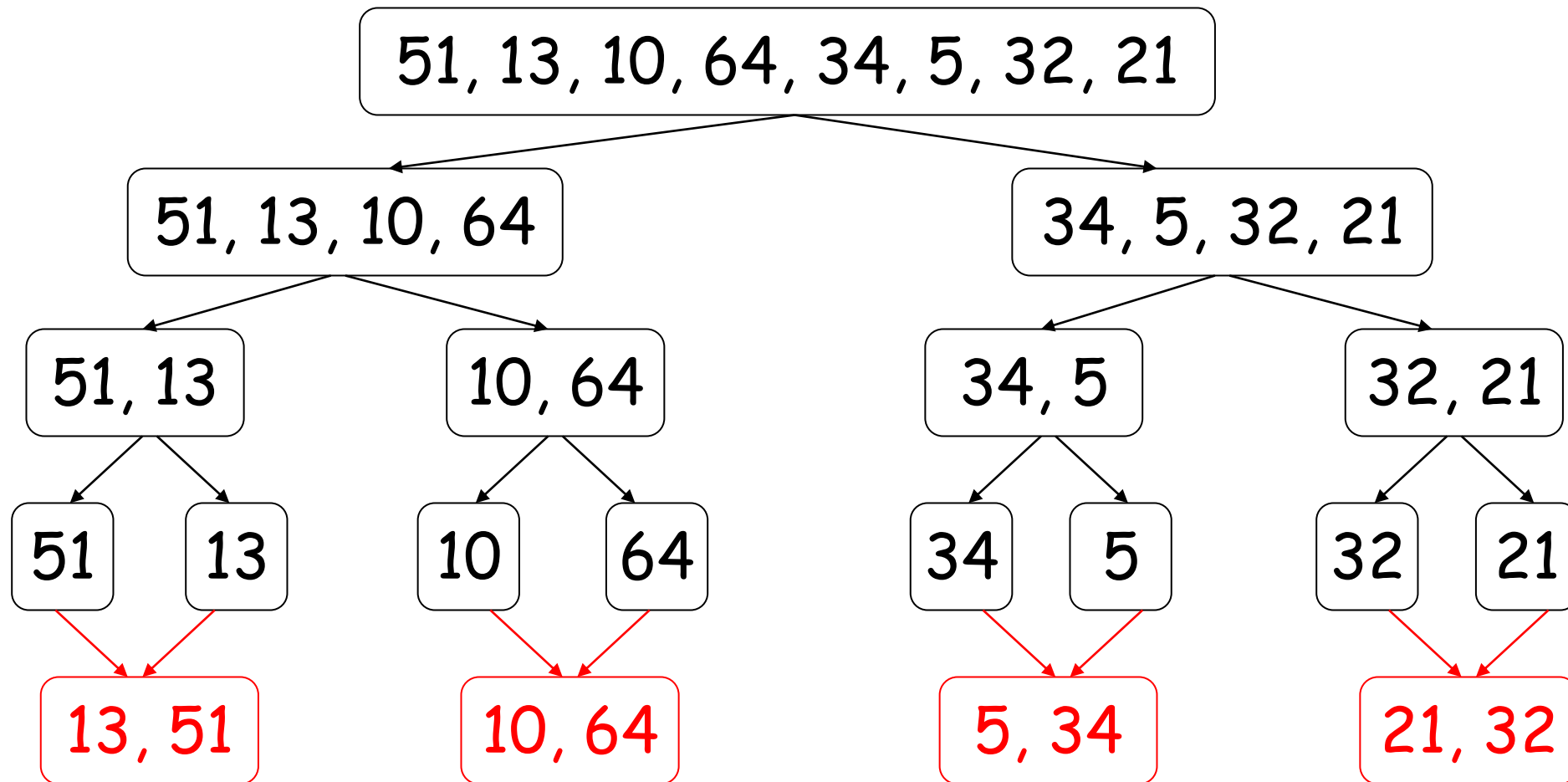
# Second step

## Conquer

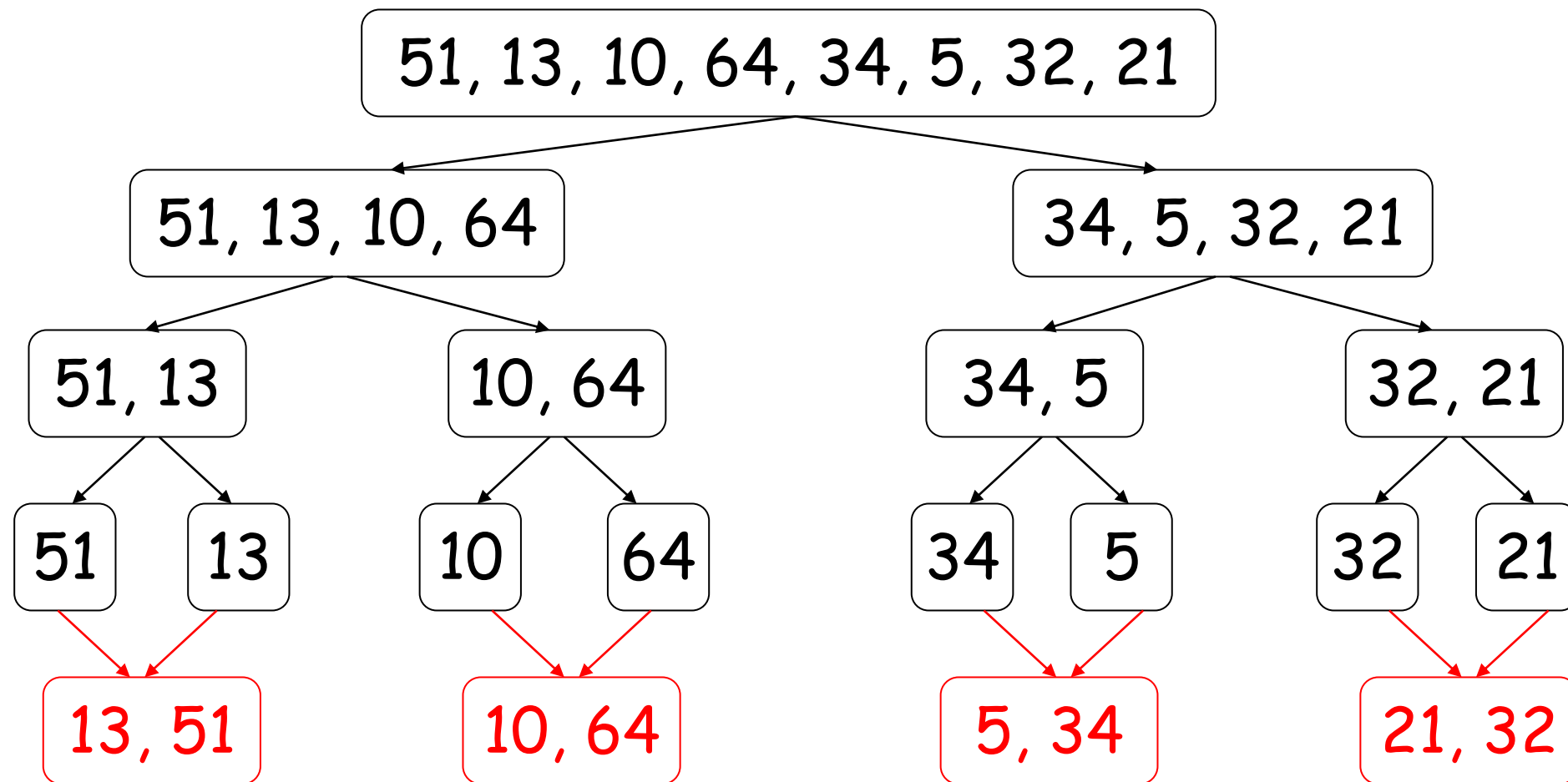
- merge two sorted sequences of **total length**  $n$  can also be done easily, at most  $n-1$  comparisons



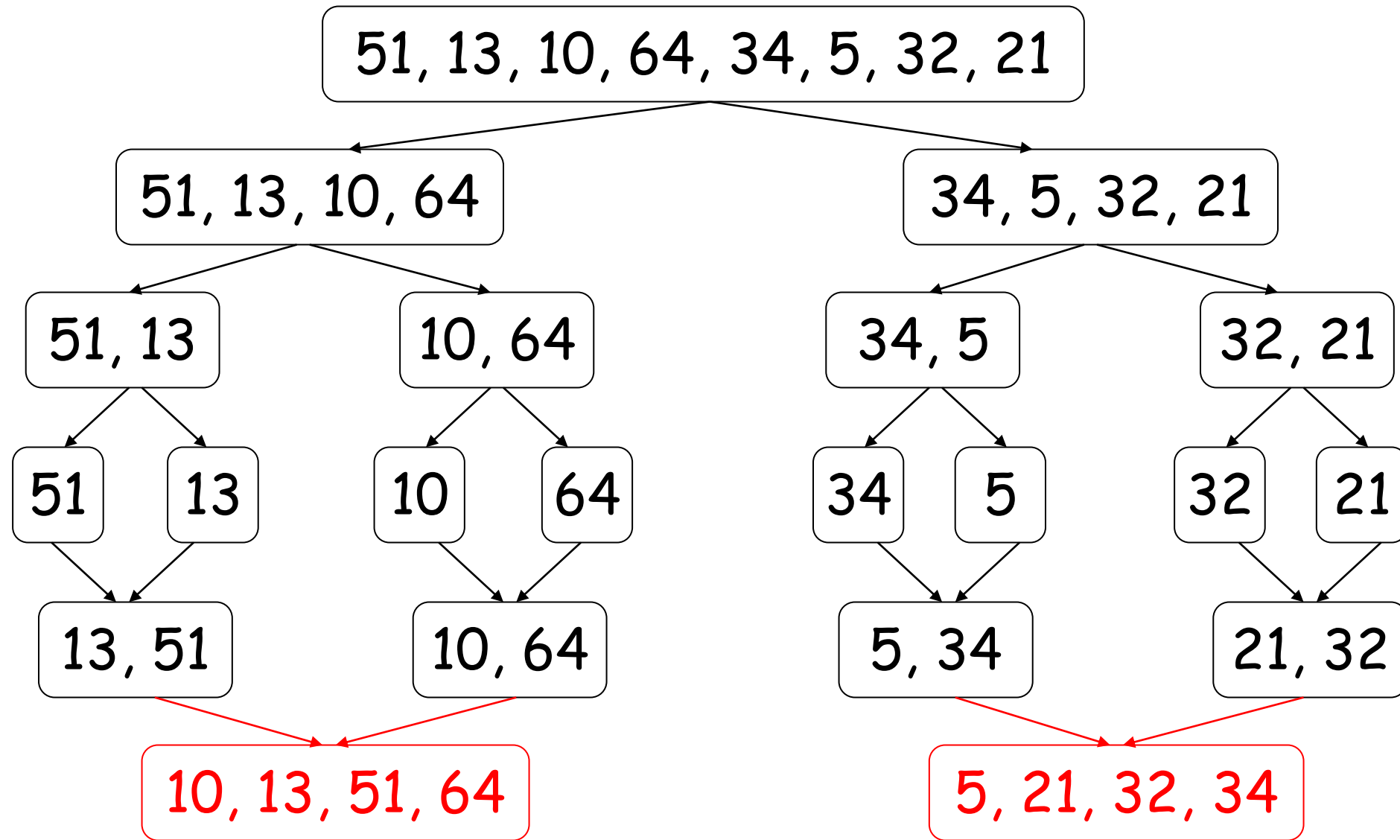
**merge** pairs of single  
number into a sequence  
of 2 sorted numbers



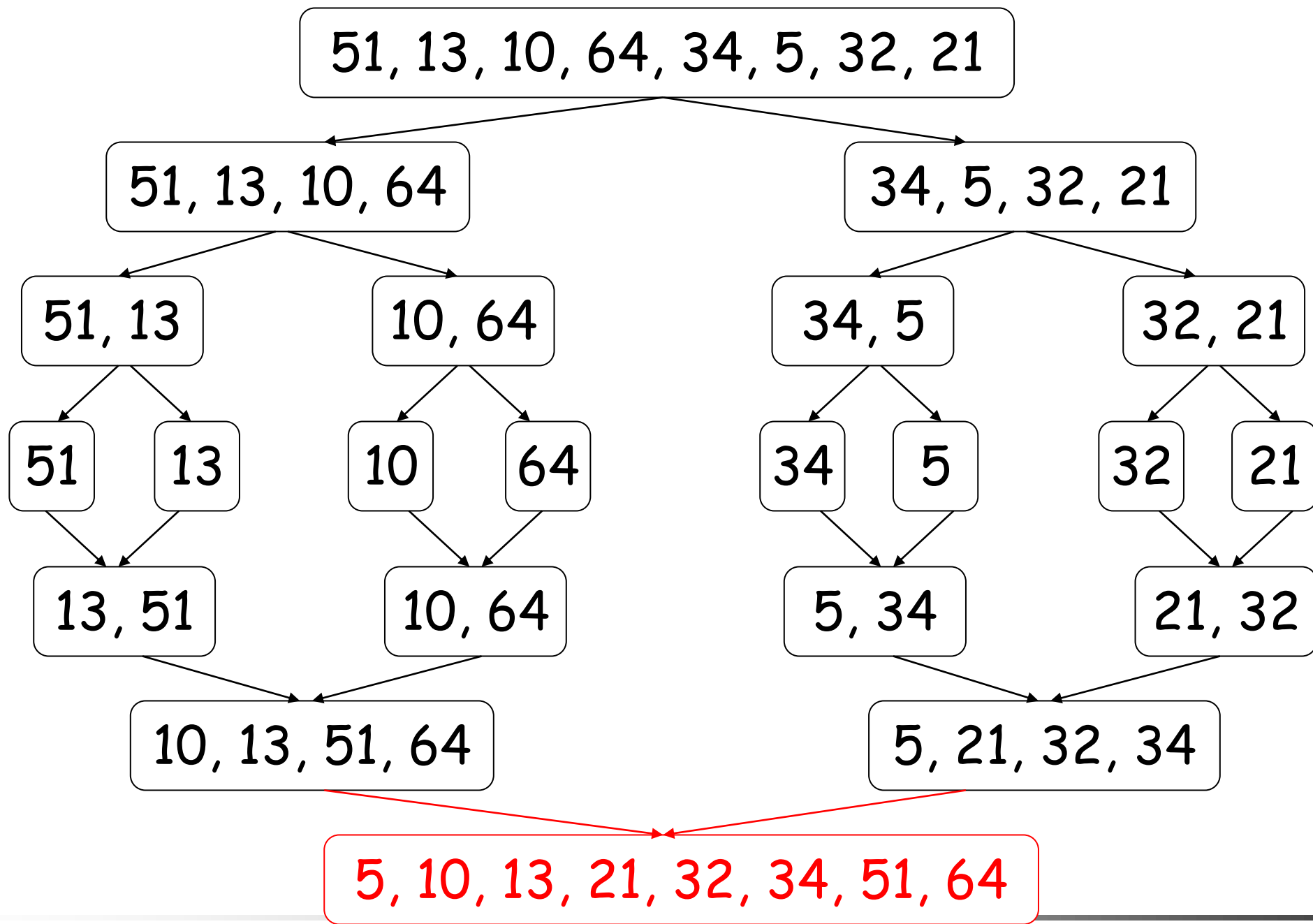
merge pairs of single  
number into a sequence  
of 2 sorted numbers



then **merge** again into sequences of  
4 sorted numbers



one more merge give the **final** sorted sequence

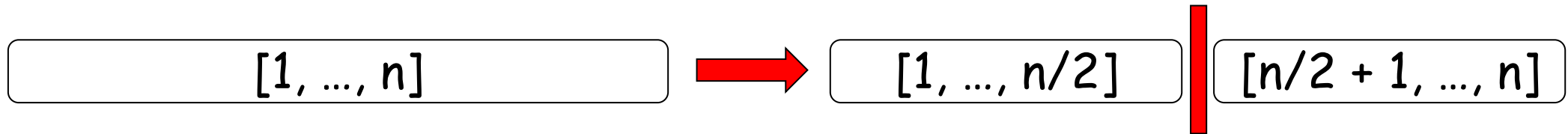




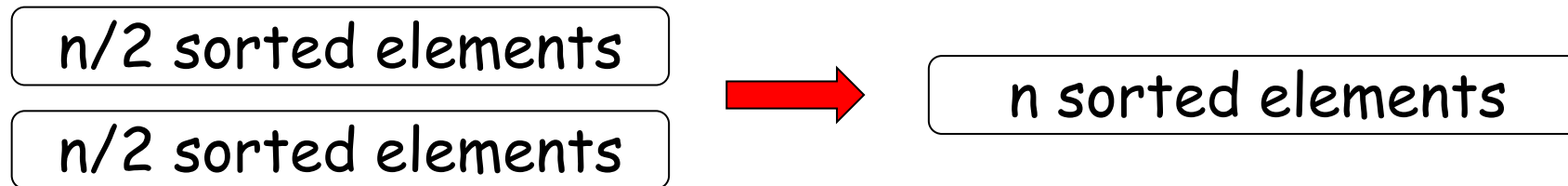
# Summary

## ■ Divide

- divide a sequence of  $n$  numbers into **two** smaller sequences



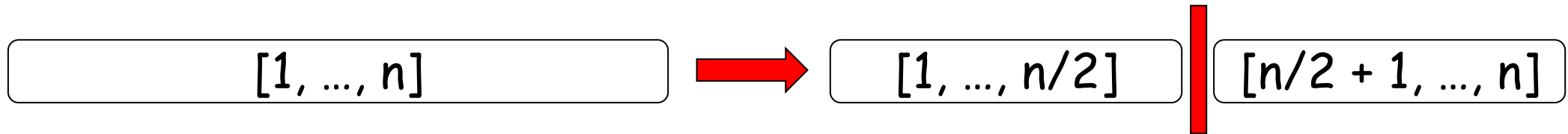
## ■ Conquer



# Summary

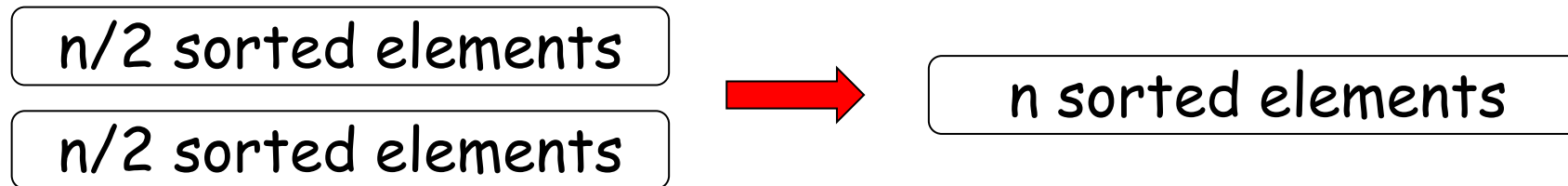
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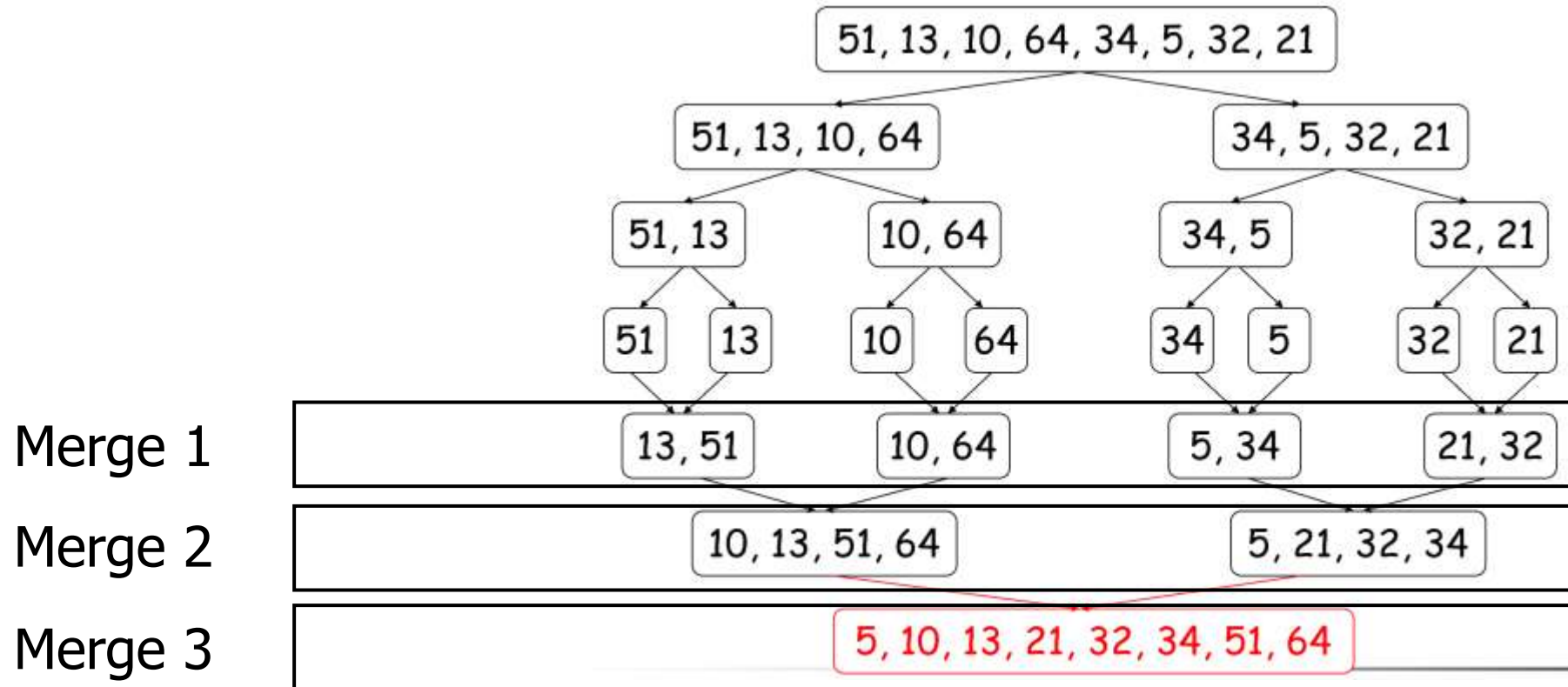
## ■ Conquer

- **merge** two sorted sequences of total length  $n$



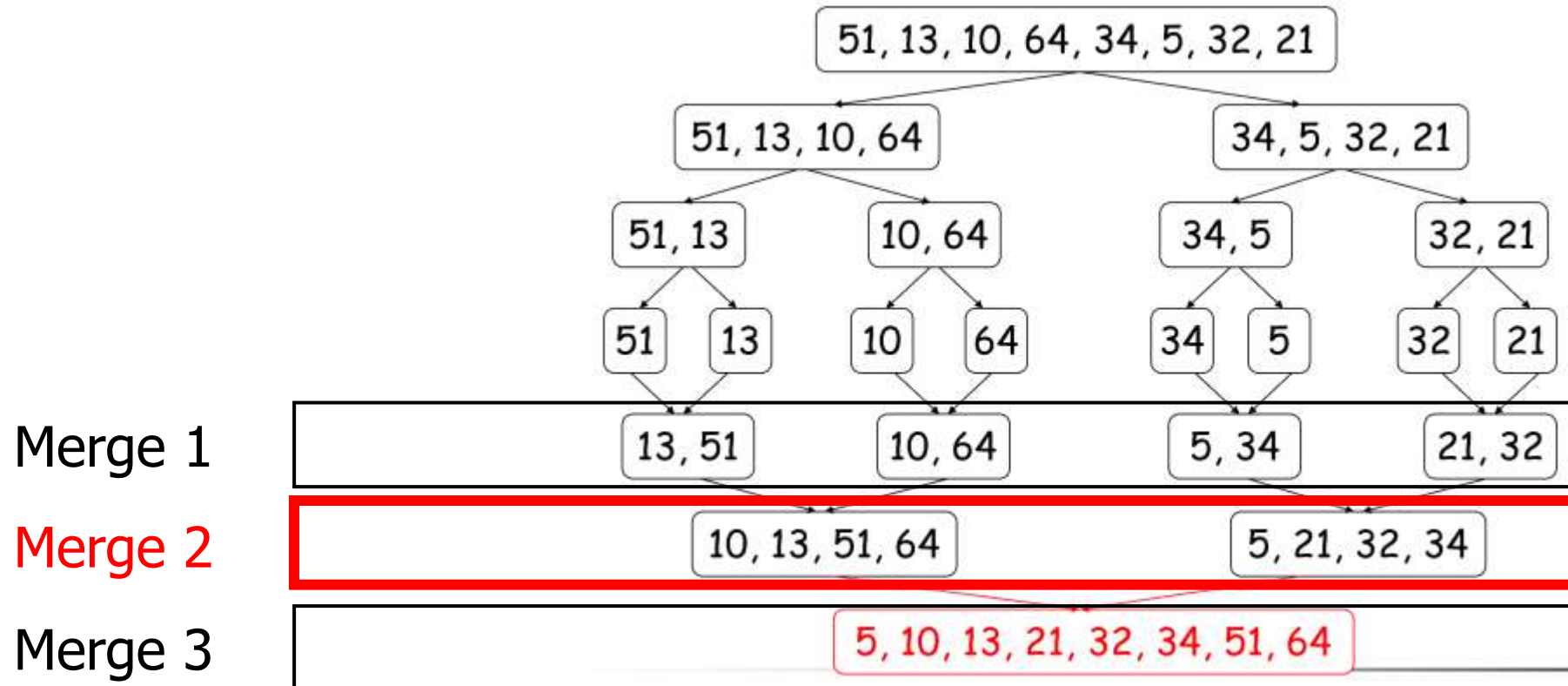
# Merge algorithm (intuition)

## ■ Example



# Merge algorithm (intuition)

## ■ Example



10, 13, 51, 64



5, 21, 32, 34



Result:

To merge two sorted sequences,  
we keep two **pointers**, one to each sequence

Compare the two numbers pointed,  
copy the **smaller** one to the result  
and **advance** the corresponding pointer

10, 13, 51, 64



5, 21, 32, 34



Result:

5,

Then compare again the two numbers  
pointed to by the pointer;  
copy the smaller one to the result  
and advance that pointer

10, 13, 51, 64

5, 21, 32, 34



Result:

5, 10,

Repeat the same process ...

10, 13, 51, 64

5, 21, 32, 34

Result:

5, 10, 13

Again ...



10, 13, 51, 64

5, 21, 32, 34

Result:

5, 10, 13, 21

and again ...

10, 13, 51, 64

5, 21, 32, 34

Result:

5, 10, 13, 21, 32

...

10, 13, 51, 64

5, 21, 32, 34

Result: 5, 10, 13, 21, 32, 34

When we reach the end of one sequence,  
simply copy the **remaining** numbers in the other  
sequence to the result

10, 13, 51, 64

5, 21, 32, 34



Result: 5, 10, 13, 21, 32, 34, 51, 64

Then we obtain the final sorted sequence

# Merge algorithm (implementation)

- Pseudo-code
- Python

Algorithm Mergesort( $A[1..n]$ )  
  if  $n > 1$  then begin  
    copy  $A[1..\lfloor n/2 \rfloor]$  to  $B[1..\lfloor n/2 \rfloor]$   
    copy  $A[\lfloor n/2 \rfloor + 1..n]$  to  $C[1..\lceil n/2 \rceil]$   
    **Mergesort**( $B[1..\lfloor n/2 \rfloor]$ )  
    **Mergesort**( $C[1..\lceil n/2 \rceil]$ )  
    **Merge**( $B, C, A$ )  
  end

```
def merge_sort(unsorted_list):  
    if len(unsorted_list) <= 1:  
        return unsorted_list  
  
    # Find the middle point and divide it  
    middle = len(unsorted_list) // 2  
  
    #divide the array into two halves  
    left_list = unsorted_list[:middle]  
    right_list = unsorted_list[middle:]  
  
    #call merge_sort for the first half  
    left_list = merge_sort(left_list)  
    #call merge_sort for the second half  
    right_list = merge_sort(right_list)  
  
    return list(merge(left_list, right_list))
```

Algorithm Mergesort( $A[1..n]$ )  
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```

Questions:

What are A, B, C? How about the Python code?

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```

Questions:  
How many functions are there?



Algorithm Merge( $B[1..p]$ ,  $C[1..q]$ ,  $A[1..p+q]$ )

set  $i=1$ ,  $j=1$ ,  $k=1$

while  $i \leq p$  and  $j \leq q$  do

begin

if  $B[i] \leq C[j]$  then

set  $A[k] = B[i]$  and  $i = i+1$

else set  $A[k] = C[j]$  and  $j = j+1$

$k = k+1$

end

if  $i = p+1$  then copy  $C[j..q]$  to  $A[k..(p+q)]$

else copy  $B[i..p]$  to  $A[k..(p+q)]$

```
def merge(left_half, right_half):
    res = left_half + right_half
    i=j=k=0
    while(i<=len(left_half)-1) and (j<=len(right_half)-1):

        if left_half[i] <= right_half[j]:
            res[k] = left_half[i]
            i = i + 1
        else:
            res[k] = right_half[j]
            j = j + 1
        k=k+1
    if(i==len(left_half)):
        res[k:len(res)] = right_half[j:len(right_half)]
    if(j==len(right_half)):
        res[k:len(res)] = left_half[i:len(right_half)]
    return res

arr = [64, 34, 25, 12, 22, 11, 90]
print(merge_sort(arr))

[11, 12, 22, 25, 34, 64, 90]
```

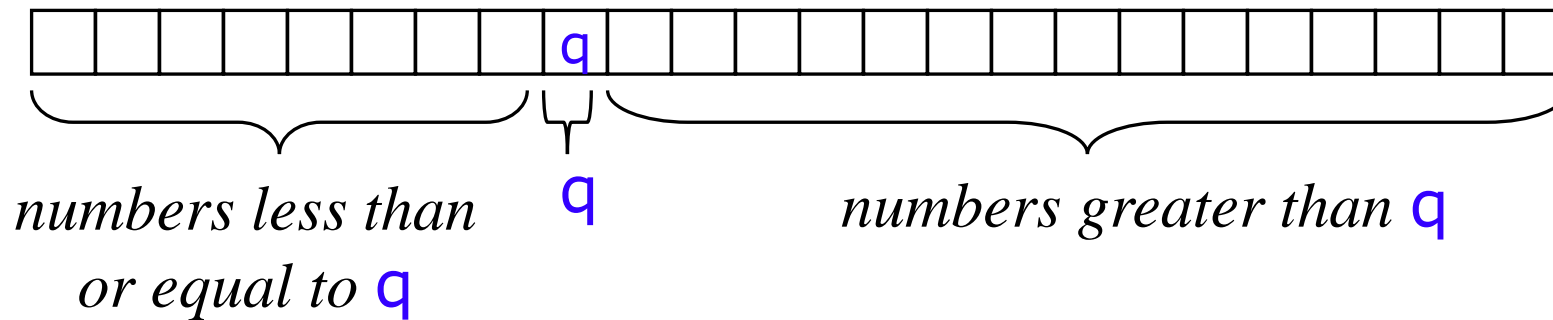
# Quicksort

# Quicksort

- Divide-and-conquer algorithm
- Sorts "in place" (like insertion sort, but not like merge sort).

# Quicksort

- Pick some number  $q$  from the array
- Move all numbers less than (or equal to)  $q$  to the beginning of the array
- Move all number greater than  $q$  to the end of the array
- Quicksort the numbers less than or equal to  $q$
- Quicksort the numbers greater than or equal to  $q$



# Quicksort pseudo code

Quicksort( $A, p, r$ )

if  $p < r$

    // partition an array such that:

    // all  $A[p, \dots, q-1]$  are less than or equal to  $A[q]$

    // all  $A[q+1, \dots, r]$  are larger than  $A[q]$

$q = \text{Partition}(A, p, r)$

    Quicksort( $A, p, q-1$ )

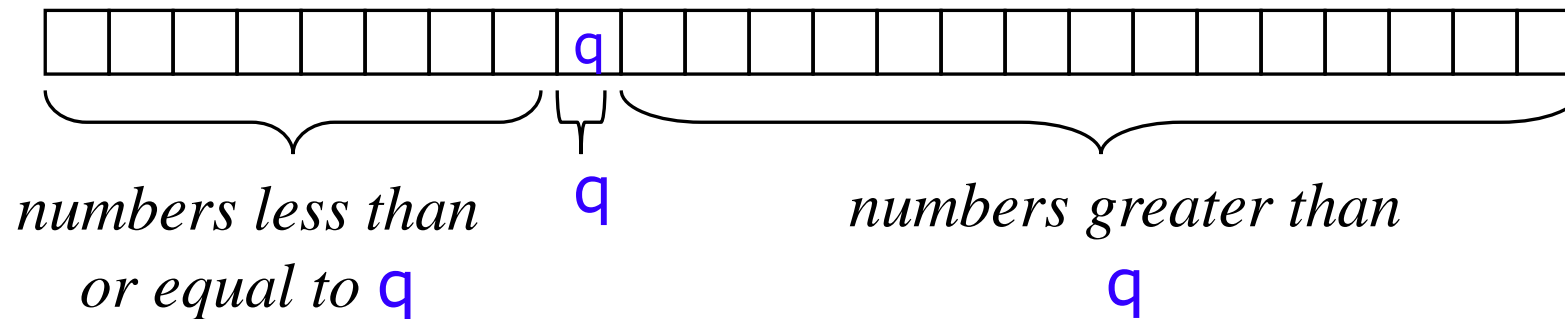
    Quicksort( $A, q+1, r$ )

---

Initial call: Quicksort( $A, 1, n$ )

# Partitioning

- A key step in the Quicksort algorithm is **partitioning** the array
  - We choose some (any) number  $q$  in the array to use as a **pivot**
  - We **partition** the array into three parts:



# Partitioning

Partition( $A, p, r$ )

$x = A[p]$

$i = p$

for  $j = p+1$  to  $r$

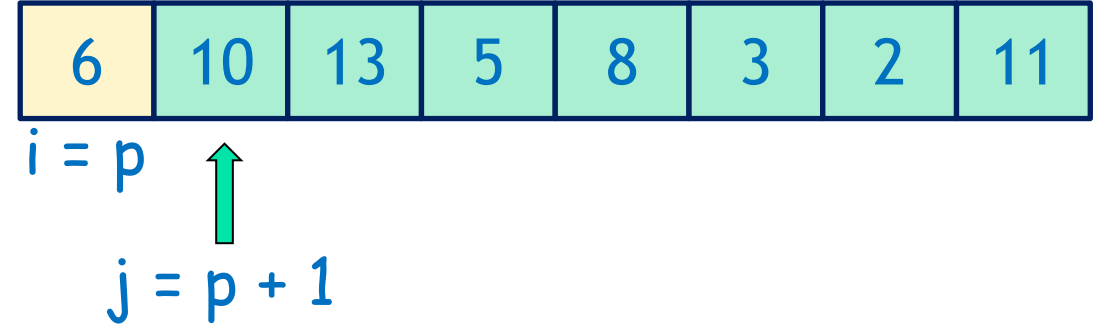
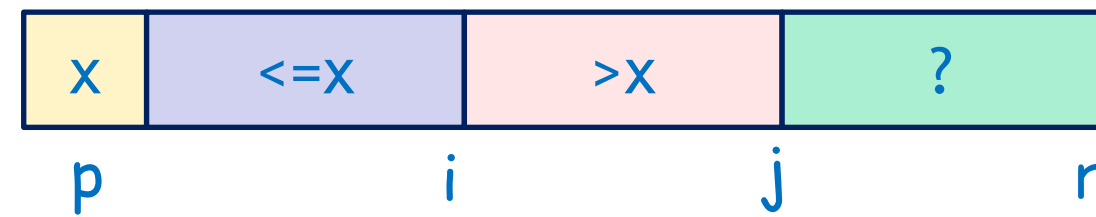
    if  $A[j] \leq x$

$i = i + 1$

        exchange  $A[i]$  with  $A[j]$

exchange  $A[p]$  and  $A[i]$

Return  $i$



# Partitioning

Partition( $A, p, r$ )

$x = A[p]$

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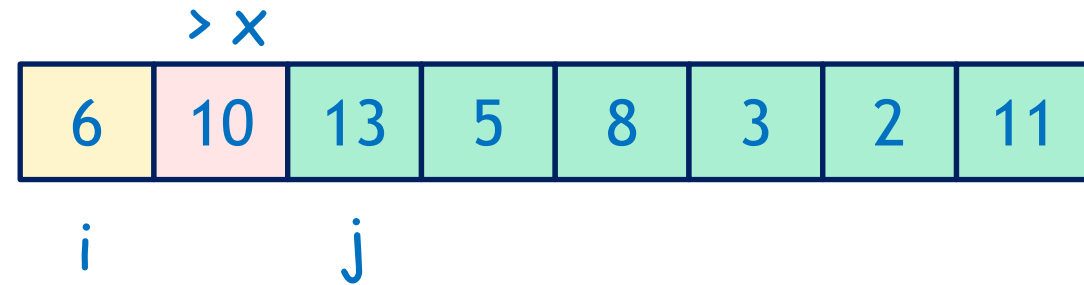
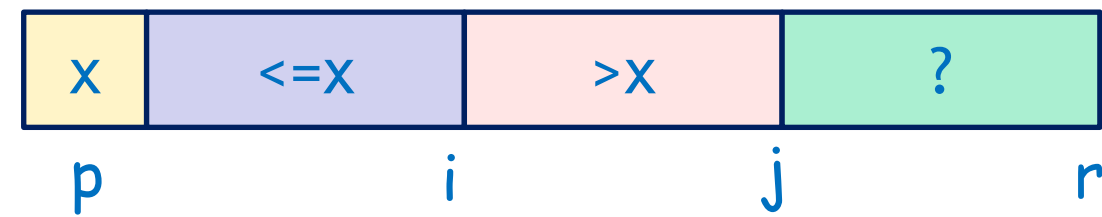
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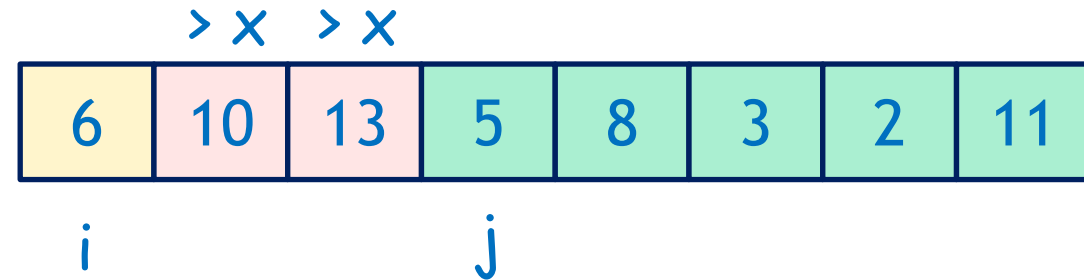
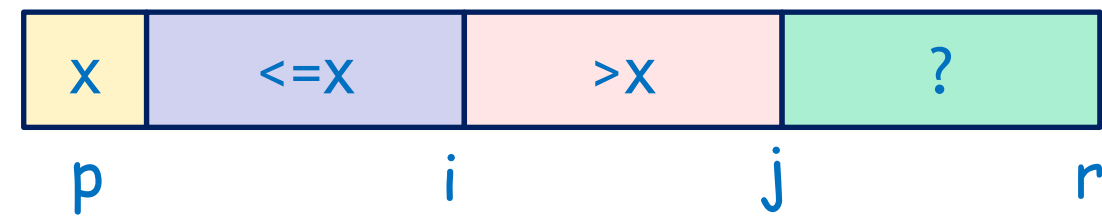
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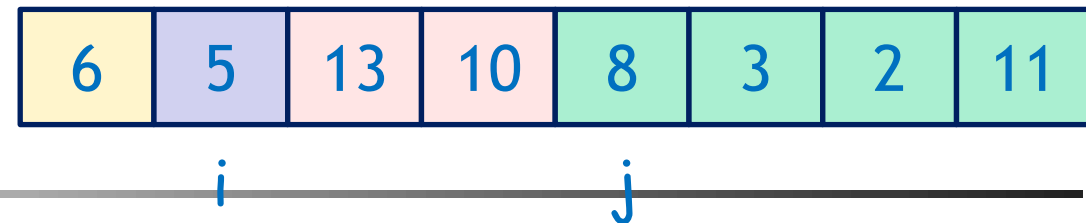
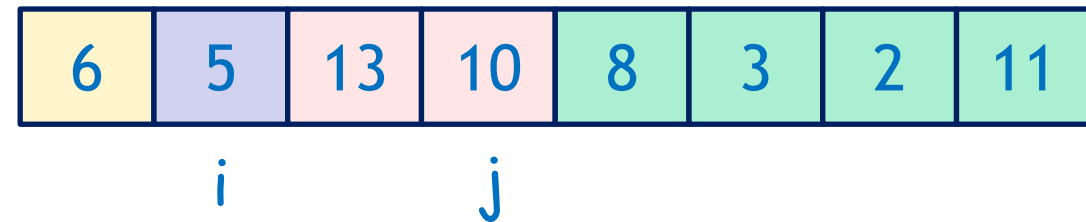
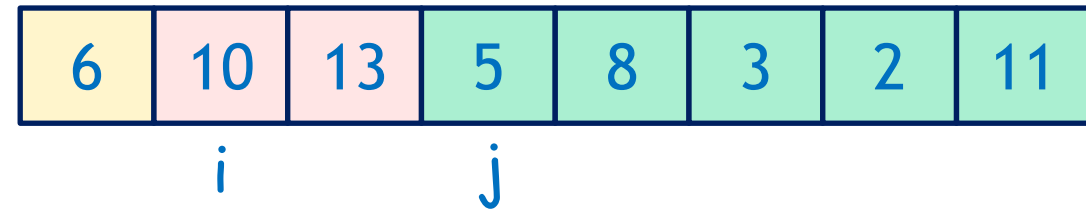
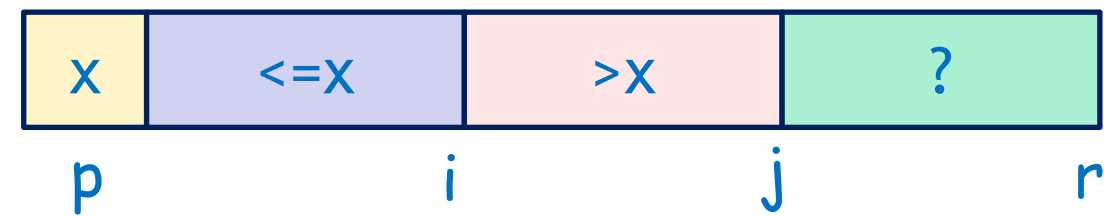
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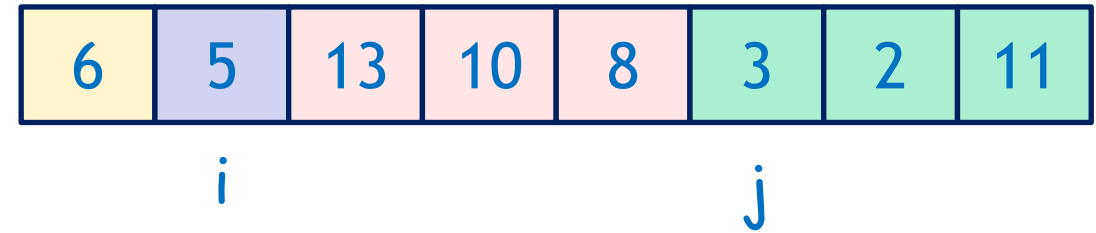
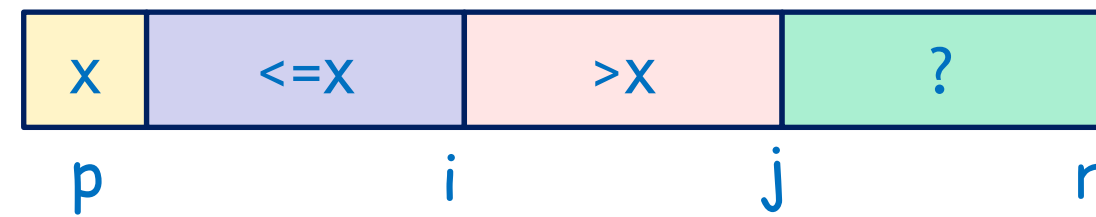
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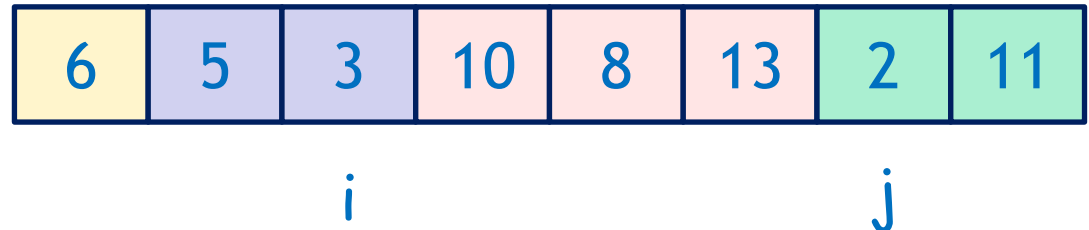
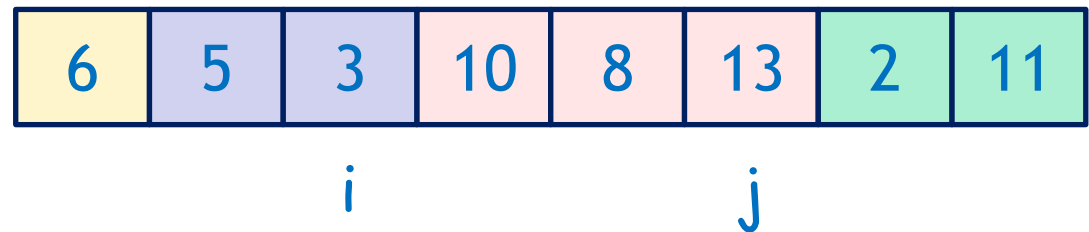
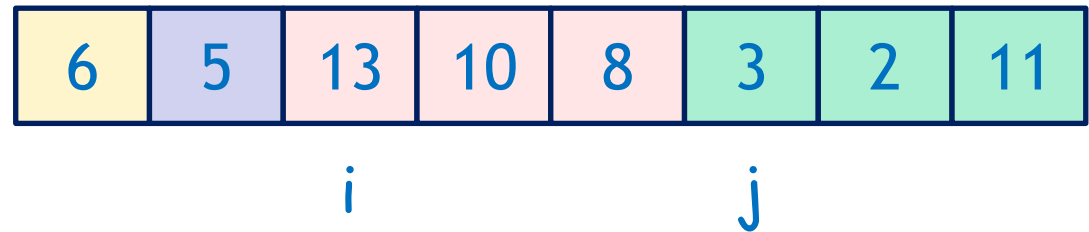
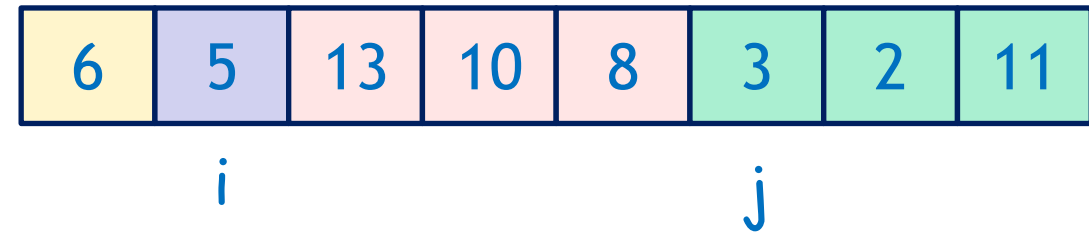
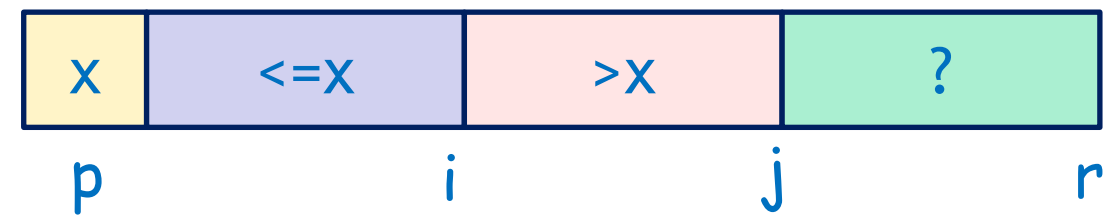
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for  $j = p+1$  to  $r$

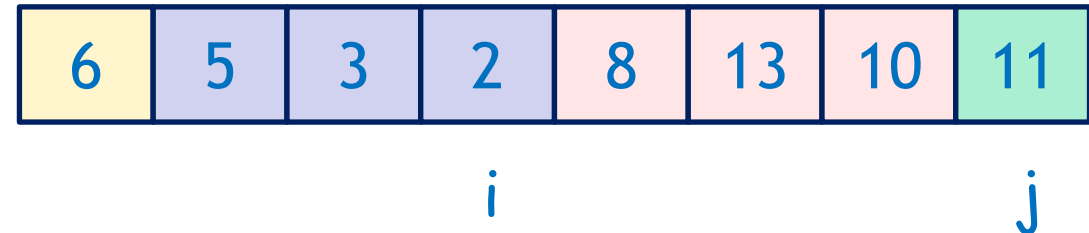
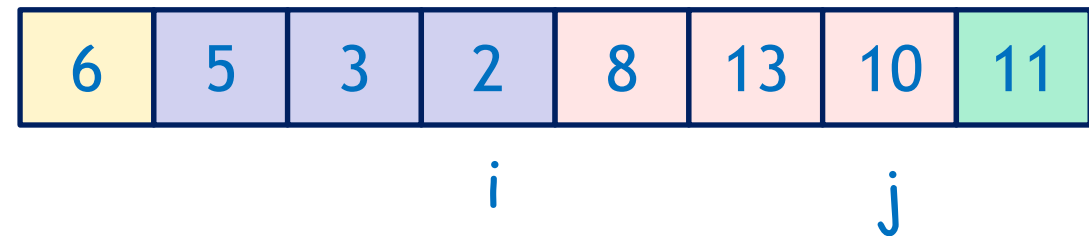
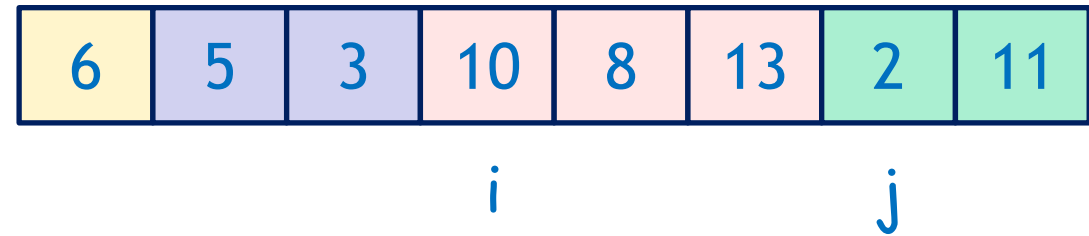
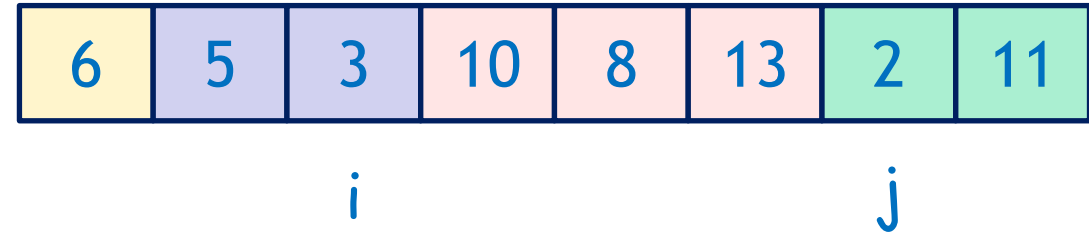
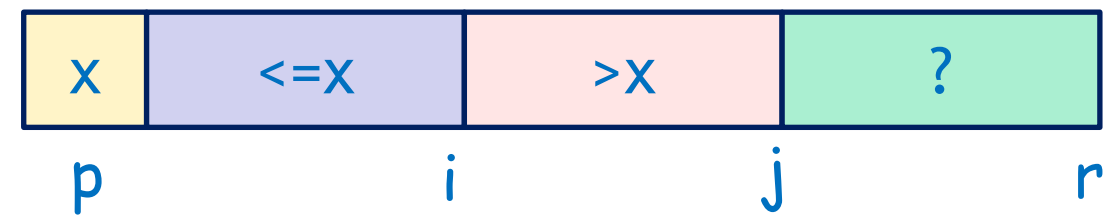
    if  $A[j] \leq x$

$i = i + 1$

        exchange  $A[i]$  with  $A[j]$

exchange  $A[p]$  and  $A[i]$

Return  $i$



# Partitioning

Partition( $A, p, r$ )

$x = A[p]$

$i = p$

for  $j = p+1$  to  $r$

    if  $A[j] \leq x$

$i = i + 1$

        exchange  $A[i]$  with  $A[j]$

exchange  $A[p]$  and  $A[i]$

Return  $i$

