

Sheet 4: Spectral radius and randomness

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Let G be a finite undirected d -regular graph on n vertices. Let A denote the adjacency matrix of G , and let b_0, b_1, \dots, b_{n-1} be an orthonormal eigenbasis for A with real eigenvalues $\lambda_0 \geq \lambda_1 \geq \dots \geq \lambda_{n-1}$.

Definition 1 Let the spectral radius of G be

$$\rho_0(G) = \max \{|\lambda_i| \mid 1 \leq i \leq n-1\}.$$

That is, we exclude the trivial eigenvalue d .

Lemma 2 If G is connected, then $b_0 = (1/\sqrt{n}, \dots, 1/\sqrt{n})$.

Solution 3 If G is connected, then

$$b_0 = \frac{1}{\sqrt{n}} \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} \quad (1)$$

Since G is a finite, undirected, d -regular graph on n vertices, thus, by results from the previous sheet, d regular and connectedness implies $\lambda_0(G) = d$, where d is also maximal degree.

The associated eigenvector for eigenvalue d is the all 1's vector,

$$\begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} \quad (2)$$

since every row of the adjacency matrix of a regular graph sum to d . Thus, we just need to normalize the all one's vector to get

$$b_0 = \frac{1}{\sqrt{n}} \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} \quad (3)$$

Let S, T be subsets of $V(G)$. Assume now that the edges of G are choosed randomly.

Exercise 4 What value do we expect for $e(S, T)$ for a random G ?

solution:

Originally, $e(S, T) = \chi_S^T A \chi_T$, but here, the edges between the 2 vertex set, S, T are chosen randomly.

since the graph G is d -regular on n vertices, there must be $\frac{d \cdot n}{2}$ edges in total, and we have $\binom{n}{2}$ possible “slots” to put those edges. And we are doing this randomly. Thus, for each “slot”, the possibility that there is an edge there, is

$$\mathbb{P}(\exists \text{ an edge}) = \frac{\frac{d \cdot n}{2}}{\frac{n(n-1)}{2}} = \frac{d}{n-1} \quad (4)$$

Thus, we should have a matrix

$$P = \begin{bmatrix} 0 & & & \frac{d}{n-1} \\ & 0 & & \\ & & \ddots & \\ \frac{d}{n-1} & & & 0 \end{bmatrix} \quad (5)$$

with zero on the diagonal and $\frac{d}{n-1}$ everywhere except the diagonal. $e(S, T) = \chi_S^T P \chi_T$ would take into account that an edge is in place with possibility $\frac{d}{n-1}$.

end solution

It turns out that when the spectral radius is small, the graph mimicks the random behaviour.

Exercise 5 For $S \subseteq V(G)$ decompose the characteristic function X_S as

$$X_S = \sum_{i=0}^{n-1} \alpha_i b_i.$$

What is α_0 ? What is $\sum_{i=0}^{n-1} \alpha_i^2$?

Solution 6 $S \subseteq V(G)$ decompose the characteristic function χ_s as $\chi_s = \sum_{i=0}^{n-1} \alpha_i b_i$ where

$$\chi_s(v) = \begin{cases} 1 & \text{if } v \in S \\ 0 & \text{if } v \notin S \end{cases} \quad (6)$$

We could think of χ_s as a column vector, where we have the n vertices, and we put 1 and 0 according to whether v_i is in S or not.

Since χ_s is a vector, it could be expressed as a linear combination of eigenvectors, since those eigenvectors form an orthonormal eigenbasis:

$$b_0, b_1, \dots, b_{n-1}. \quad (7)$$

Thus, we have

$$\chi_s = \alpha_0 b_0 + \alpha_1 b_1 + \dots + \alpha_{n-1} b_{n-1} \quad (8)$$

α_0 is the coefficient corresponding to the eigenvector

$$b_0 = \frac{1}{\sqrt{n}} \begin{bmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{bmatrix} \quad (9)$$

Also,

$$\langle \chi_s, b_0 \rangle = \left\langle \sum_i \alpha_i b_i, b_0 \right\rangle \quad (10)$$

$$= \sum_i \alpha_i \langle b_i, b_0 \rangle \quad (11)$$

$$= \sum_i \alpha_i \delta_{i,0} \quad (12)$$

$$= \alpha_0 \quad (13)$$

$$\sum_{i=0}^{n-1} \alpha_i^2 = \alpha_0^2 + \alpha_1^2 + \dots + \alpha_{n-1}^2 \quad (14)$$

Theorem 7 For all $S, T \subseteq V(G)$ we have

$$\left| e(S, T) - \frac{d|S||T|}{n} \right| \leq \rho_0(G) \sqrt{|S||T|}.$$

Proof. As suggested in exercises that builds up to this theorem, write

$$\chi_S = \sum_{i=0}^{n-1} \alpha_i b_i \quad (15)$$

$$\chi_T = \sum_{i=0}^{n-1} \beta_i b_i \quad (16)$$

in eigenbasis. The α_i and β_i are coefficients, and

$$\alpha_i = \langle \chi_S, b_i \rangle \quad (17)$$

$$\beta_i = \langle \chi_T, b_i \rangle \quad (18)$$

also, in particular,

$$\alpha_0 = \langle \chi_S, b_0 \rangle \quad (19)$$

$$= \left\langle \chi_S, \frac{1}{\sqrt{n}} \mathbf{1} \right\rangle \text{ where } \mathbf{1} \text{ is } n \times 1 \text{ vector of all } 1\text{'s} \quad (20)$$

$$= \frac{1}{\sqrt{n}} \langle \chi_S, \mathbf{1} \rangle \quad (21)$$

$$= \frac{1}{\sqrt{n}} |S| \quad (22)$$

similarly, $\beta_0 = \frac{1}{\sqrt{n}}|T|$.
Thus

$$e(S, T) = \chi_S^T A \chi_T \quad (23)$$

$$= \sum_{i=0}^{n-1} \alpha_i b_i^T A \beta_j b_j \quad (24)$$

$$= \sum_{i=0}^{n-1} \alpha_i \beta_i \lambda_i \text{ since } b_i \text{ form orthonormal basis} \quad (25)$$

$$= \alpha_0 \beta_0 \lambda_0 + \sum_{i=1}^{n-1} \alpha_i \beta_i \lambda_i \quad (26)$$

$$= \frac{|S|}{\sqrt{n}} \frac{|T|}{\sqrt{n}} d + \sum_{i=1}^{n-1} \alpha_i \beta_i \lambda_i. \quad (27)$$

Thus, we have

$$|e(S, T) - \frac{d|S||T|}{n}| = \left| \sum_{i=1}^{n-1} \alpha_i \beta_i \lambda_i \right|. \quad (28)$$

Now, LHS of equation we want to show

$$= \left| \sum_{i=1}^{n-1} \alpha_i \beta_i \lambda_i \right| \quad (29)$$

$$\leq \left| \sum_{i=1}^{n-1} \alpha_i \beta_i \right| \rho_0(G) \quad (30)$$

$$\leq \left| \sum_{i=0}^{n-1} \alpha_i \beta_i \right| \rho_0(G) \quad (31)$$

$$= |\langle \chi_S, \chi_T \rangle| \rho_0(G) \quad (32)$$

$$\leq \|\chi_S\| \|\chi_T\| \rho_0(G) \quad (33)$$

$$= \sqrt{|S|} \sqrt{|T|} \rho_0(G) \quad (34)$$

$$= \rho_0(G) \sqrt{|S||T|} \quad (35)$$

$$= RHS \quad (36)$$

observe

$$|\langle \chi_S, \chi_T \rangle| \quad (37)$$

$$= \left| \left\langle \sum_i \alpha_i b_i, \sum_j \beta_j b_j \right\rangle \right| \quad (38)$$

$$= \left| \sum_{i,j} \alpha_i \beta_j \langle b_i, b_j \rangle \right| \quad (39)$$

$$= \left| \sum_{i,j} \alpha_i \beta_j \delta_{i,j} \right| \quad (40)$$

$$= \left| \sum_i \alpha_i \beta_i \right| \quad (41)$$

and

$$\|\chi_S\| = \sqrt{\langle \chi_S, \chi_S \rangle} \quad (42)$$

$$= \sqrt{|S|} \quad (43)$$

$$\|\chi_T\| = \sqrt{\langle \chi_T, \chi_T \rangle} \quad (44)$$

$$= \sqrt{|T|} \quad (45)$$

$$(46)$$

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A subset $S \subseteq V(G)$ is *independent*, if $E(S, S)$ is empty. Using the above theorem, one can easily get an upper bound for the maximal size of an independent subset of G .

Corollary 8 *An independent subset of G has size at most*

$$\frac{\rho_0(G)}{d}n.$$

Proof. From the previous theorem, we have $\forall S, T \subseteq V(G)$,

$$|e(S, T) - \frac{d|S||T|}{n}| \leq \rho_0(G)\sqrt{|S||T|} \quad (47)$$

Let $S \subseteq V(G)$ be independent, we have $e(S, S) = 0$. So the above formula reduce to

$$|e(S, S) - \frac{d|S||S|}{n}| \leq \rho_0(G)\sqrt{|S||S|} \quad (48)$$

$$\implies \left| \frac{d}{n}||S|^2 \right| \leq \rho_0(G)|S| \quad (49)$$

$$\implies \left| \frac{d}{n}||S| \right| \leq \rho_0(G) \quad (50)$$

$$|S| \leq \rho_0(G)\frac{n}{d}. \quad (51)$$

Thus, an independent subset S has size at most $\frac{\rho_0(G)}{d}n$.

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Of course, this bound is only meaningful, if $\rho_0(G)$ can go below, say, $d/2$. It is also not clear how small $\rho_0(G)$ can get for large d -regular graphs – we will get a good estimate on this later, using graph limits.