



Lec 4: Sweep Operator

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Agenda

- Gauss Jordan wrap up
- The Sweep Operator

Matrix Decomposition

For linear system, $Ax = b$, one never actually computes $A^{-1}b$.

When $A = \mathbf{X}^\top \mathbf{X}$ is symmetric, some special computational methods are available:

- Sweep operator
- QR decomposition
- Cholesky decomposition

Gauss Jordan

$$\text{GJ}[1 : n][A|b] = [I|A^{-1}b] = A^{-1}[A|b],$$

$$\text{GJ}[1 : n][A|I] = [I|A^{-1}] = A^{-1}[A|I].$$

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Matrix version:

$$\begin{aligned} & \text{GJ}[1 : m] \left[\begin{array}{cc|cc} A_{11} & A_{12} & I_1 & 0 \\ A_{21} & A_{22} & 0 & I_2 \end{array} \right] \\ &= \left[\begin{array}{cc|cc} I_1 & A_{11}^{-1}A_{12} & A_{11}^{-1} & 0 \\ 0 & A_{22} - A_{21}A_{11}^{-1}A_{12} & -A_{21}A_{11}^{-1} & I_2 \end{array} \right] \end{aligned}$$

Sweep operator

The sweep operator arises from a close analysis of elimination applied to the equation $\mathbf{X}^\top \mathbf{X} \beta = \mathbf{X}^\top \mathbf{Y}$. For a symmetric matrix A , $B = \text{Sweep}(A, k)$ is obtained as follows:

- 1 Divide k^{th} row and k^{th} column by a_{kk} :

$$b_{ik} = a_{ik}/a_{kk}; \quad b_{kj} = a_{kj}/a_{kk}$$

- 2 Subtract $a_{ik}a_{kj}/a_{kk}$ from the other entries:

$$b_{ij} = a_{ij} - a_{ik}a_{kj}/a_{kk}$$

- 3 Invert k^{th} diagonal element:

$$b_{kk} = -1/a_{kk}$$

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Question: What is the time complexity? A single sweep operation visits every element of the matrix and requires roughly one divide per element. This leads to roughly $\mathcal{O}(n^2)$ divides.

R code

```
mySweep <- function(A, m)
{
  n <- dim(A)[1]
  for (k in 1:m)
  {
    for (i in 1:n)
      for (j in 1:n)
        if (i!=k & j!=k)
          A[i,j] <- A[i,j] - A[i,k]*A[k,j]/A[k,k]
    for (i in 1:n)
      if (i!=k)
        A[i,k] <- A[i,k]/A[k,k]
    for (j in 1:n)
      if (j!=k)
        A[k,j] <- A[k,j]/A[k,k]
    A[k,k] <- - 1/A[k,k]
  }
  return(A)
}
A = matrix(c(1,2,3,7,11,13,17,21,23), 3,3)
solve(A)
```

```
##      [,1] [,2] [,3]
## [1,]  1.00 -3.0  2.00
## [2,] -0.85  1.4 -0.65
## [3,]  0.35 -0.4  0.15
```

```
mySweep(A,3)
```

Proposition 1

Proposition

Suppose $V_{p \times m} = U_{p \times m} A_{m \times m}$, and $B = \text{Sweep}(A, k)$. Then $\hat{V} = \hat{U}B$, where:

- $\hat{U} = U$ except that its k^{th} column is v_k ;
- $\hat{V} = V$ except that its k^{th} column is $-u_k$.

Proof:

Proposition 2

Proposition

If A is a symmetric invertible matrix and we sweep on each diagonal element of A the result is $B = \text{Sweep}(A, 1 : n) = -A^{-1}$

Proof:

Proposition 3

Proposition

If $A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}$ and we sweep on each of the diagonal entries of A_{11} , we get $B = \begin{bmatrix} -A_{11}^{-1} & A_{11}^{-1}A_{12} \\ A_{21}A_{11}^{-1} & A_{22} - A_{21}A_{11}^{-1}A_{12} \end{bmatrix}$

Proof: Start with an simple example $A = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}$

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Proof: Take advantage of proposition 2.

Proposition 4

We construct a matrix $\mathbf{Z} = [\mathbf{X}\mathbf{Y}]$, and let

$$\mathbf{A} = \mathbf{Z}^\top \mathbf{Z} = \begin{bmatrix} \mathbf{X}^\top \mathbf{X} & \mathbf{X}^\top \mathbf{Y} \\ \mathbf{Y}^\top \mathbf{X} & \mathbf{Y}^\top \mathbf{Y} \end{bmatrix}$$

be the cross-product matrix. Then

$$\begin{aligned} \text{SWP}[1 : p] \mathbf{A} &= \begin{bmatrix} -(\mathbf{X}^\top \mathbf{X})^{-1} & (\mathbf{X}^\top \mathbf{X})^{-1} \mathbf{X}^\top \mathbf{Y} \\ \mathbf{Y}^\top \mathbf{X} (\mathbf{X}^\top \mathbf{X})^{-1} & \mathbf{Y}^\top \mathbf{Y} - \mathbf{Y}^\top \mathbf{X} (\mathbf{X}^\top \mathbf{X})^{-1} \mathbf{X}^\top \mathbf{Y} \end{bmatrix} \\ &= \begin{bmatrix} -\frac{\text{Var}(\hat{\beta})}{\hat{\beta}^\top} & \hat{\beta} \\ \hat{\beta}^\top & \text{RSS} \end{bmatrix} \end{aligned}$$

where $\text{RSS} = \|\mathbf{Y} - \mathbf{X}\hat{\beta}\|_{\ell_2}^2$ is the residual sum of squares.

Other Application of sweep operator

Perform sweep operator on multivariate normal distribution

$$\begin{bmatrix} \Omega & x - \mu \\ x^t - \mu^t & 0 \end{bmatrix}$$

Why Sweep Operator?

- ① Efficient for computing the central statistics used in multiple regression
 - Start from a square correlation or covariance matrix.
 - Compute multiple correlation, residual variance, regression slopes, and standard errors of slopes, plus some other
- ② An efficient way to compute a whole series of regressions, as in stepwise regression.
- ③ (Drawback) If it is applied many times to the same matrix, rounding error can accumulate.