

Tools for OpenMP Programming

PPCES 2018

Tim Cramer



OpenMP Tools

ThreadSanitizer

- → Overview
- Intel Inspector XE
 - → Overview
 - → Live Demo
- Intel VTune Amplifier XE
 - → Overview
 - → Live Demo





Race Condition

- Data Race: the typical OpenMP programming error, when:
 - → two or more threads access the same memory location, and
 - → at least one of these accesses is a write, and
 - → the accesses are not protected by locks or critical regions, and
 - → the accesses are not synchronized, e.g. by a barrier.
- Non-deterministic occurrence: e.g. the sequence of the execution of parallel loop iterations is non-deterministic and may change from run to run
- In many cases private clauses, barriers or critical regions are missing
- Data races are hard to find using a traditional debugger
 - → Use the Intel Inspector XE or the ThreadSanitizer





OpenMP Tools

ThreadSanitizer

- → Overview
- Intel Inspector XE
 - → Overview
 - → Live Demo
- Intel VTune Amplifier XE
 - → Overview
 - → Live Demo





OpenMP Correctness checking: ThreadSanitizer – Overview

- Correctness checking for threaded applications
- Integrated in clang and gcc compiler
- Low runtime overhead: 2x 15x
- Used to find data races in browsers like Chrome and Firefox





OpenMP Correctness checking: ThreadSanitizer – Usage

module load archer gcc/6

- Compile the program with clang compiler:

 clang -fsanitize=thread -fopenmp -g myprog.c -o myprog

 clang++ -fsanitize=thread -fopenmp -g myprog.cpp

 -o myprog

 gfortran -fsanitize=thread -fopenmp -g myprog.f -c

 clang -fsanitize=thread -fopenmp -lgfortran myprog.o

 -o myprog
 - Execute: OMP_NUM_THREADS=4 ./myprog
 - Understand and correct the detected threading errors
 - Edit the source code
 - Repeat until no errors reported





OpenMP Correctness checking: ThreadSanitizer – Result Summary

```
1 #include <stdio.h>
   int main(int argc, char **argv)
 4
       int a = 0;
       #pragma omp parallel
 6
          if (a < 100) \{ \leftarrow
             #pragma omp critical
 9
             a++; ←
10
12 }
```

```
WARNING: ThreadSanitizer: data race
```

```
Read of size 4 at 0x7fffffffdcdc by thread T2:

#0 .omp_outlined. race.c:7

(race+0x0000004a6dce)

#1 __kmp_invoke_microtask <null>
(libomp_tsan.so)
```

 Previous write of size 4 at 0x7fffffffdcdc by main thread:

```
(race+0x0000004a6e2c)
#1 __kmp_invoke_microtask <null>
(libomp_tsan.so)
```

#0.omp outlined.race.c:9





OpenMP Tools

- ThreadSanitizer
 - → Overview
- Intel Inspector XE
 - → Overview
 - → Live Demo
- Intel VTune Amplifier XE
 - → Overview
 - → Live Demo





Intel Inspector XE

Detection of

- → Memory Errors
- → Deadlocks
- → Data Races

Support for

- → Linux (32bit and 64bit) and Windows (32bit and 64bit)
- → WIN32-Threads, Posix-Threads, Intel Threading Building Blocks and OpenMP

Features

- → Binary instrumentation gives full functionality
- → Independent stand-alone GUI for Windows and Linux





PI Example Code

```
double f(double x)
                                                           3.5
  return (4.0 / (1.0 + x*x));
                                                           2.5
double CalcPi (int n)
                                                           1.5
  const double fH = 1.0 / (double) n;
  double fSum = 0.0;
  double fX;
                                                           0.5
  int i;
                                                           -0.5
#pragma omp parallel for private(fX,i) reduction(+:fSum)
  for (i = 0; i < n; i++)
     fX = fH * ((double)i + 0.5);
fSum += f(fX);
  return fH * fSum;
```





0.5

-10.5

PI Example Code

```
double f(double x)
  return (4.0 / (1.0 + x*x));
                                                  What if we
double CalcPi (int n)
                                                  would have
  const double fH = 1.0 / (double) n;
                                                   forgotten
  double fSum = 0.0;
  double fX;
                                                      this?
  int i;
#pragma omp parallel for private(fX,i) reduction(+:fSum)
  for (i = 0; i < n; i++)
    fX = fH * ((double)i + 0.5);
    fSum += \dot{f}(fX);
  return fH * fSum;
```





Live Demo

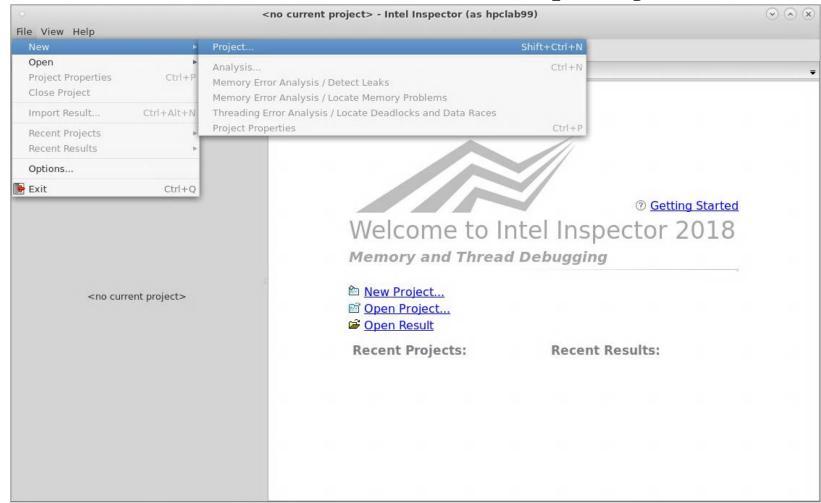
Intel Inspector XE





Inspector XE – Create Project

\$ module load intelixe ; inspxe-gui



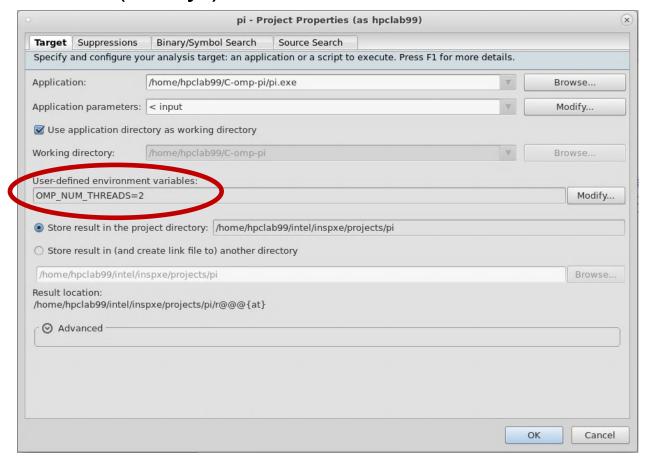


Inspector XE – Create Project

ensure that multiple threads are used

- choose a small dataset (really!), execution time can increase

10X - 1000X



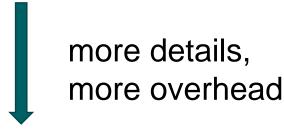


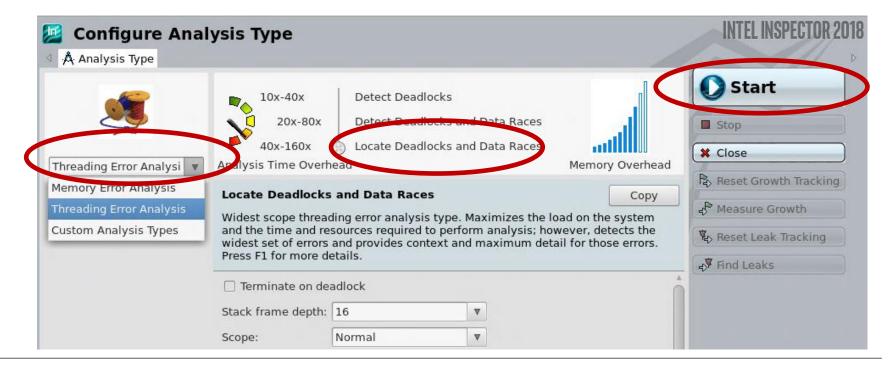


Inspector XE – Configure Analysis

Threading Error Analysis Modes

- Detect Deadlocks
- 2. Detect Deadlocks and Data Races
- Locate Deadlocks and Data Races



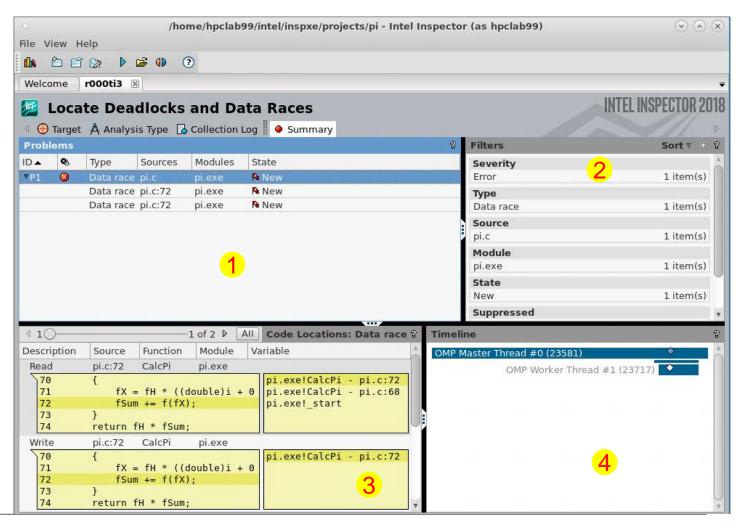






Inspector XE – Results

- detected problems
- 2 filters
- 3 code location
- 4 Timeline

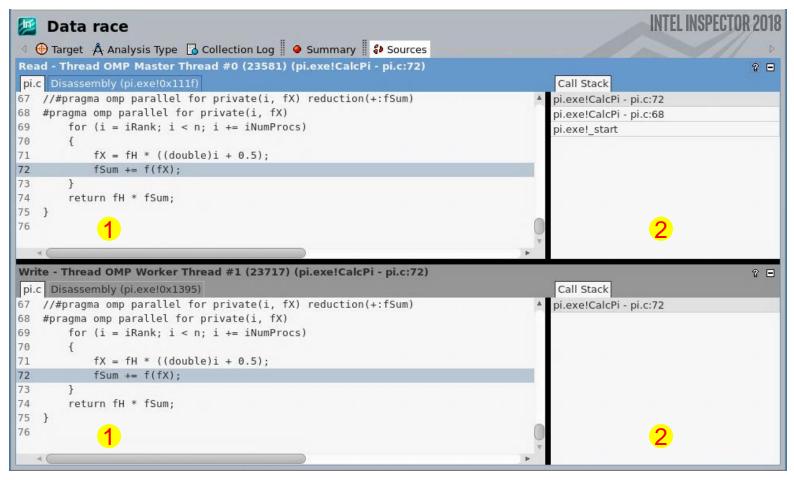






Inspector XE – Results

- Source Code producing the issue double click opens an editor
- Corresponding Call Stack







Inspector XE – Results

Source Code producing the issue – double click opens an editor

Corresponding Call Stack The missing reduction is **Data race** detected. √ ⊕ Target A Analysis Type Collection Log Summary Sources Read - Thread OMP Master Thread #0 (23581) (pi.exe!CalcPi - pi.c:72) 8 = pi.c Disassembly (pi.exe!0x111f) Call Stack //#pragma omp parallel for private(i, fX) reduction(+:fSum) pi.exe!CalcPi - pi.c:72 #pragma omp parallel for private(i, fX) pi.exe!CalcPi - pi.c:68 69 for (i = iRank; i < n; i += iNumProcs) pi.exe! start 70 71 fX = fH * ((double)i + 0.5);72 fSum += f(fX);73 74 return fH * fSum; 75 } 76 Write - Thread OMP Worker Thread #1 (23717) (pi.exe!CalcPi - pi.c:72) pi.c Disassembly (pi.exe!0x1395) Call Stack //#pragma omp parallel for private(i, fX) reduction(+:fSum) pi.exe!CalcPi - pi.c:72 #pragma omp parallel for private(i, fX) 69 for (i = iRank; i < n; i += iNumProcs) 70 71 fX = fH * ((double)i + 0.5);72 fSum += f(fX);73 return fH * fSum; 75 } 76





Command Line Tool – inspxe-cl

Threading Error Analysis Modes

- 1. Detect Deadlocks (ti1)
- 2. Detect Deadlocks and Data Races (ti2)
- 3. Locate Deadlocks and Data Races (ti3)

\$ inspxe-cl -collect ti3 -- pi.exe < input

Data collection without GUI allows to use batch jobs.

\$ inspxe-cl -report problems ...

Viewing results in text mode is helpful, when remote connections are slow.





OpenMP Tools

- ThreadSanitizer
 - → Overview
- Intel Inspector XE
 - → Overview
 - → Live Demo
- Intel VTune Amplifier XE
 - → Overview
 - → Live Demo





Intel VTune Amplifier XE

Performance Analyses for

- → Serial Applications
- → Shared Memory Parallel Applications
- Sampling-based measurements
- Features:
 - → Hot Spot Analysis
 - → Concurrency Analysis
 - → Wait
 - → Hardware Performance Counter Support

Prerequisites:

- → ssh -Y login-t
- → module load intelytune

Application

- → Use representative data set
- → Build with debug information (-g) AND optimization (-O3)
- → Set environment variables (e.g., OMP_NUM_THREADS, OMP_PROC_BIND, OMP_PLACES, etc.)





Stream

- Standard Benchmark to measure memory performance.
- Version is parallelized with OpenMP.

Measures Memory bandwidth for:

```
y=x (copy)
y=s*x (scale)
y=x+z (add)
y=x+s*z (triad)
```

#pragma omp parallel for
 for (j=0; j<N; j++)
 b[j] = scalar*c[j];</pre>

for double vectors x,y,z and scalar double value s

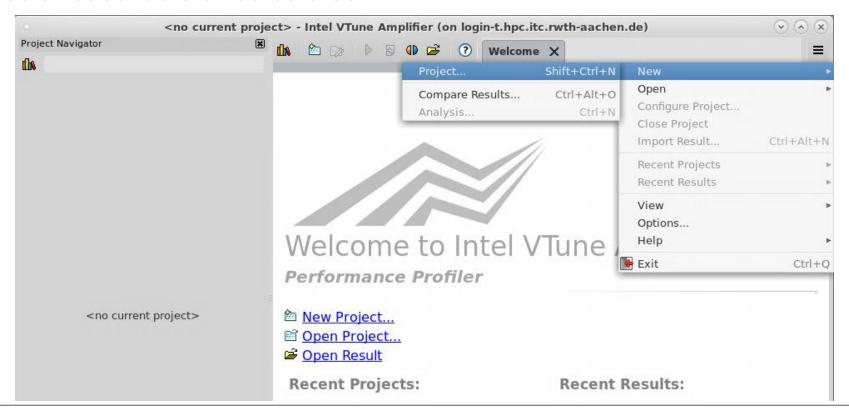
Function	Rate (MB/s)	Avg time	Min time	Max
time				
Copy:	121755.3877	0.0119	0.0118	0.0122
Scale:	122280.4397	0.0118	0.0118	0.0119
Add:	124754.8422	0.0173	0.0173	0.0176
Triad:	124797.8048	0.0174	0.0173	0.0178





Intel VTune Amplifier XE – Create Project

- \$ module load intelvtune; amplxe-gui
- Create a Project in the same way as with the inspector.
- Executable should be build with optimization.
- Use a reasonable sized data set.

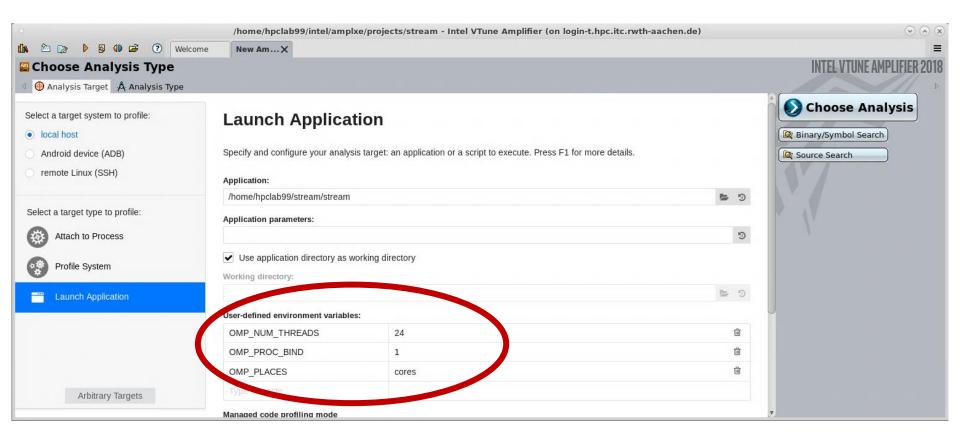






Intel VTune Amplifier XE – Create Project

Set environment variables







Amplifier XE – Algorithm Analysis

Different Analysis Types

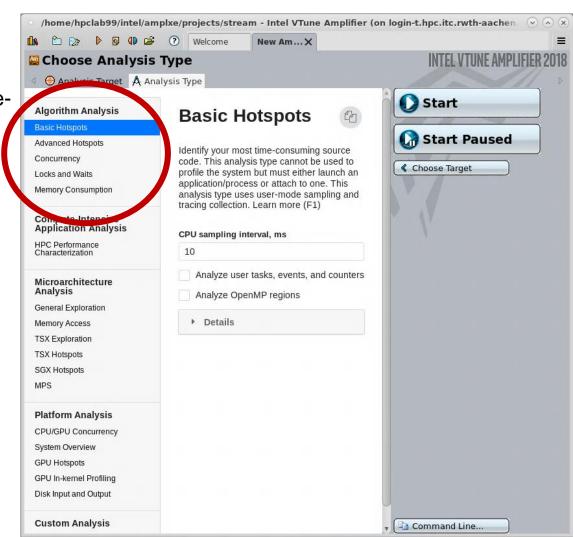
Basic Hotspots: Identify most timeconsuming source code

Advanced Hotspots: Uses additional OS kernel support and analyses the CPI (Cycles Per Instruction) metric

Concurrency: Identify synchronization overhead and potential candidates for parallelism

Locks and Waits: Identify waiting times and discover how this affects your application performance

Memory consumption





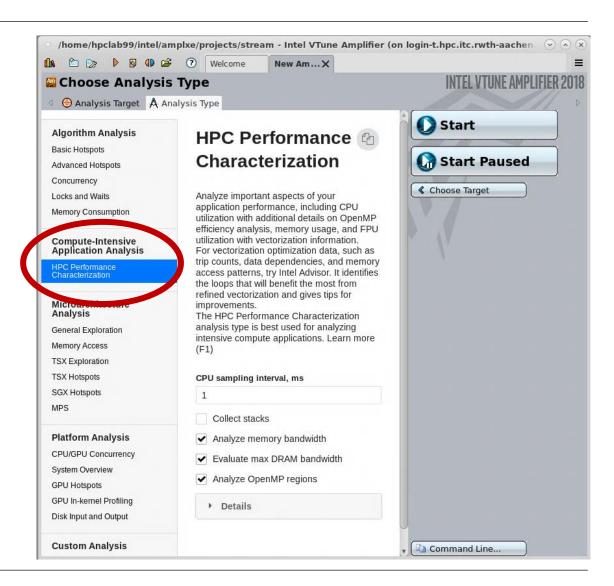


Amplifier XE – HPC Performance Characterization

Different Analysis Types

HPC Performance Characterization: Analyze

- CPU utilization
- OpenMP efficiency
- memory usage
- FPU utilization
- vectorization





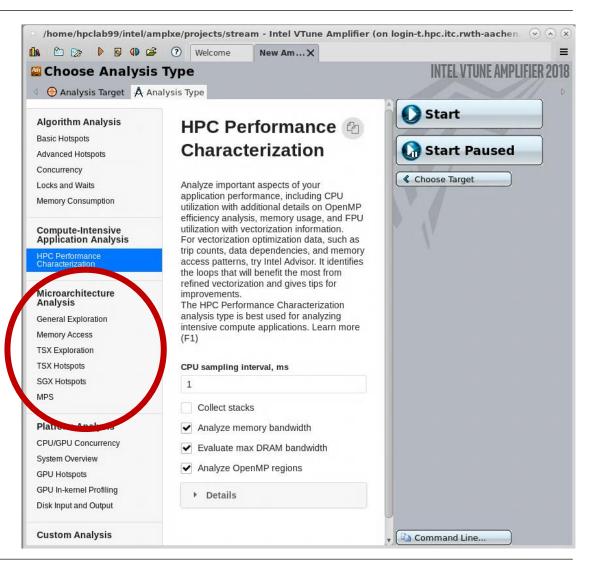


Amplifier XE – Microarchitecture Analysis

Different Analysis Types

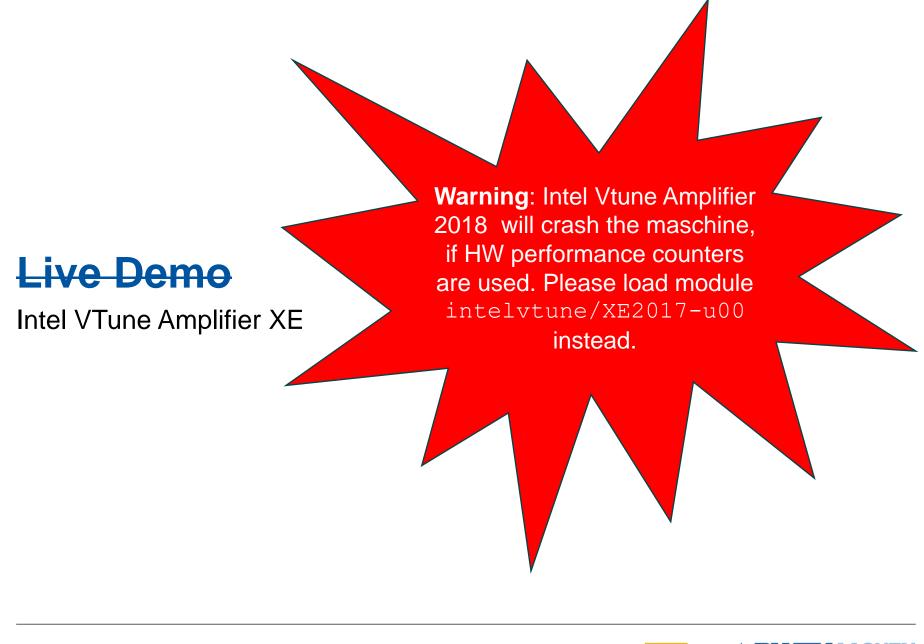
Microarchitecture Analysis:

- Based on hardware sampling
- Identify memory access related issues (e.g., NUMA effects)
- Analyze Intel Transactional Synchronization Extensions (TSX)
- MPI Performance Snapshot (MPS)



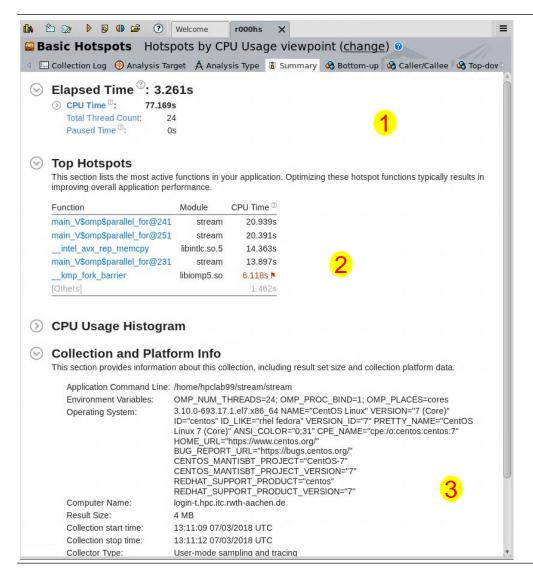












Summary:

- 1 General Timing Information
- 2 Top Hotspots
- 3 Platform Information

Tool marks locations which might have a performance issue with a red flag.

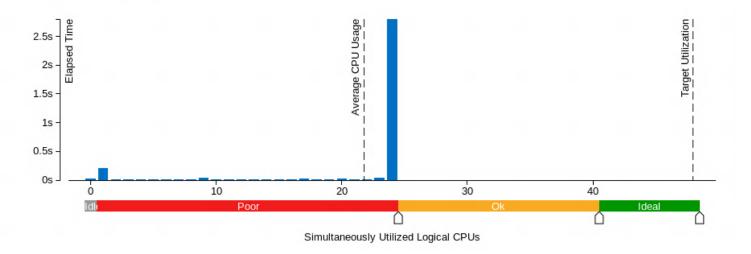




CPU Usage Histogram

- Show CPU Utilization
- Tool counts logical cores
 - → Hyperthreads are not used in most HPC application
 - → Use slider to change

This histogram displays a percentage of the wall time the specific number of CPUs were running simultaneously. Spin and Overhead time adds to the Idle CPU usage value.

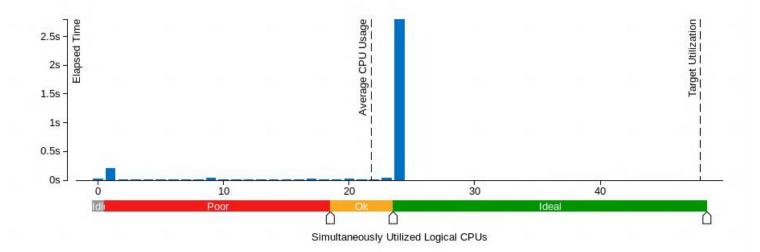




CPU Usage Histogram

- Show CPU Utilization
- Tool counts logical cores
 - → Hyperthreads are not used in most HPC application
 - → Use slider to change
 - CPU Usage Histogram

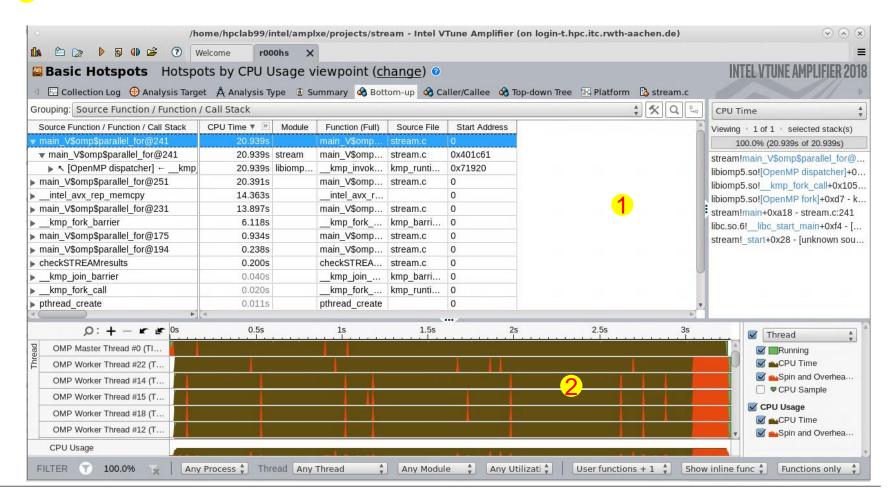
This histogram displays a percentage of the wall time the specific number of CPUs were running simultaneously. Spin and Overhead time adds to the Idle CPU usage value.







- 1 Function Summary
- 2 Timeline View



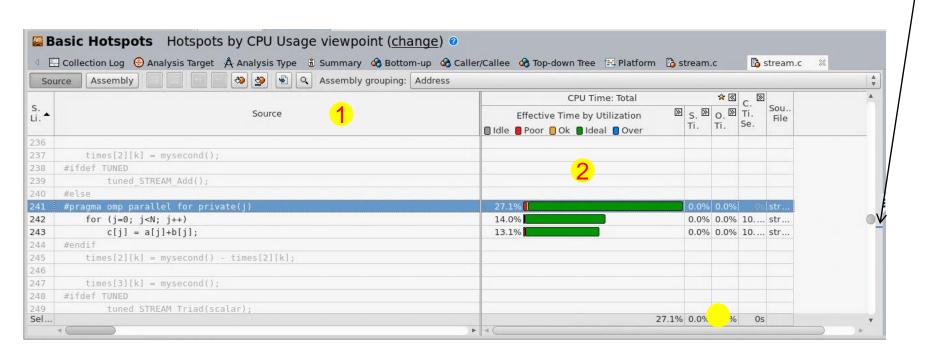




Double clicking on a function opens source code view.

- Source Code View (only if compiled with -g)
- 2 CPU Utilization (would be all "poor", if slider in histogram is unmodified)

Hotspots







Hardware Performance Counters

Hardware Counters of our Intel Nehalem Processor:

SB_DRAIN.ANY, STORE_BLOCKS,AT_RET, DTLB_LOAD_MISSES.LARGE_W, MEM_INS MEM_UNCORE_RETIRED.L3_D, MEM_UNC FP_COMP_OPS_EXE.MMX, FP_COMP_OPS SIMD_INT_128.PACKED_MPY, SIMD_INT_1 LOAD_DISPATCH.RS_DELAYED, LOAD_DIS L2_RQSTS.LD_HIT, L2_RQSTS.LD_MISS, L2 L2_RQSTS.PREFETCH_MISS, L2_RQSTS.P L2_DATA_RQSTS.PREFETCH.L_, L2_DATA_ L2_WRITE.RFO.M_STATE, L2_WRITE.RFO.

L1I.HITS:

Counts all instruction fetches that hit the L1 instruction cache.

S.STLB_HIT, DTLB_LOAD_MISSES.PDE_MIS,
OPS_ISSUED.STALLED_CYCLE, UOPS_ISSUED.FUSED,
RETIRED.UNCA, FP_COMP_OPS_EXE.X87,
E_SING, FP_COMP_OPS_EXE.SSE_DOU,
IT_128.SHUFFLE_MOV, LOAD_DISPATCH.RS,
DED, INST_QUEUE_WRITE_CYCLES, LSD_OVERFLOW,
ITS.IFETCHES, L2_RQSTS.PREFETCH_HIT,
M_, L2_DATA_RQSTS.DEMAND.ME,
ITATE, L2_WRITE.RFO.S_STATE,
WRITE.LOCK.HIT. L2_WRITE_LOCK.MESI.

L1D WB L2.I STATE, L1D WB L2.S STATE, L1D WB L2.E STATE, L1D WB L2.M STATE, L1D WB L2.M STATE, L1D WB L2.MESI, CPU CLK UNHALTED. THREAD, CPU CLK UNH L1D CACHE LD.E STATE, L1D CACHE LD.M STATE, L1D CACHE LD.MESI, L1D CACHE ST.S STATE, L1D CACHE ST.E STATE, L1D CACHE ST.M STATE, L1D CACHE LOCK.HIT, L1D CACHE LOCK.S STATE, L1D CACHE LOCK.E STATE, L1D CACHE LOCK.M STATE, L1D ALL REF.ANY, L1D ALL REF.CACHEABLE, DTLB MISSES.ANY, DTLB MISSES.WALK COMPLET, DTLB MISSES.STLB HIT, DTLB MISSES.PDE MISS. DTLB_MISSES.LARGE_WALK_C, LOAD_HIT_PRE, L1D_PREFETCH.REQUESTS, L1D_PREFETCH.MISS, L1D_PREFETCH.TRIGGERS, L1D.M_EPL, L1D.M_EVICT, L1D.M_SNOOP_EVICT, L1D_CACHE_PREFETCH_LOCK L1D_CACHE_LOCK_FB_HIT, CACHE_LOCK_CYCLES.L1D_L2, CACHE_LOCK_CYCLES.L1D, IO_TRANSACTIONS, L1I.CYCLES_STALLED, LARGE_ITLB.HIT, ITLB_MISSES.ANY, ITLB_MISSES.WALK_COMPLET, ILD_STALL.LCP, ILD_STALL.MRU, ILD_STALL.IQ_FULL, ILD_STALL.REGEN, ILD_STALL.ANY, BR_INST_EXEC.COND, BR_INST_EXEC.DIRECT_NEAR, BR_INST_EXEC.INDIRECT_NEA, BR_INST_EXEC.NEAR_CALLS, BR_INST_EXEC.TAKEN, BR_MISP_EXEC.COND, BR_MISP_EXEC.DIRECT, BR_MISP_EXEC.INDIRECT_NO, BR_MISP_EXEC.NON_CALLS, BR_MISP_EXEC.RETURN_NEA, BR MISP EXEC.DIRECT NEAR, BR MISP EXEC.INDIRECT NEA, BR MISP EXEC.NEAR CALLS, BR MISP EXEC.TAKEN, RESOURCE STALLS.ANY, RESOURCE STALLS.LOAD, RESOURCE STALLS.RS FULL, RESOURCE STALLS.STORE, RESOURCE_STALLS.ROB_FULL, RESOURCE_STALLS.FPCW, RESOURCE_STALLS.MXCSR, RESOURCE_STALLS.OTHER, MACRO_INSTS.FUSIONS_DECO, BACLEAR_FORCE_IQ, ITLB_FLUSH, OFFCORE_REQUESTS.L1D_WR, UOPS EXECUTED.PORT0, UOPS EXECUTED.PORT1, UOPS EXECUTED.PORT2 COR, UOPS EXECUTED.PORT3 COR, UOPS EXECUTED.PORT4 COR, UOPS EXECUTED.PORT5, UOPS EXECUTED.PORT5015. UOPS_EXECUTED.PORT234, OFFCORE_REQUESTS_SQ_FUL, OFF_CORE_RESPONSE_0, SNOOP_RESPONSE.HIT, SNOOP_RESPONSE.HITE, SNOOP_RESPONSE.HITM, OFF_CORE_RESPONSE_1, INST_RETIRED.ANY_P, INST RETIRED.X87, INST RETIRED.MMX, UOPS RETIRED.ANY, UOPS RETIRED.RETIRE SLOTS, UOPS RETIRED.MACRO FUSE, MACHINE CLEARS.CYCLES, MACHINE CLEARS.MEM ORDE, MACHINE CLEARS.SMC, BR INST RETIRED.ALL BRAN, BR INST RETIRED.CONDITION, BR INST RETIRED.NEAR CAL, BR MISP RETIRED.ALL BRAN, BR MISP RETIRED.NEAR CAL, SSEX UOPS RETIRED.PACKED, SSEX UOPS RETIRED.SCALAR, SSEX_UOPS_RETIRED.PACKED, SSEX_UOPS_RETIRED.SCALAR, SSEX_UOPS_RETIRED.VECTOR, ITLB_MISS_RETIRED, MEM_LOAD_RETIRED.L1D_HIT, MEM_LOAD_RETIRED.L2_HIT, MEM_LOAD_RETIRED.L3_UNS, MEM_LOAD_RETIRED.OTHER_, MEM_LOAD_RETIRED.L3_MISS, MEM_LOAD_RETIRED.HIT_LFB, MEM_LOAD_RETIRED.DTLB_MI, FP_MMX_TRANS.TO_FP, FP_MMX_TRANS.TO_FMX, FP_MMX_TRANS.ANY, MACRO_INSTS.DECODED, UOPS DECODED.MS, UOPS DECODED.ESP FOLDING, UOPS DECODED.ESP SYNC, RAT STALLS.FLAGS, RAT STALLS.REGISTERS, RAT STALLS.ROB READ POR, RAT STALLS.SCOREBOARD, RAT STALLS.ANY, SEG RENAME STALLS, ES REG RENAMES, UOP UNFUSION, BR INST DECODED, BPU MISSED CALL RET, BACLEAR.BAD TARGET, BPU CLEARS.EARLY, BPU CLEARS.LATE, L2 TRANSACTIONS.LOAD, L2 TRANSACTIONS.RFO, L2_TRANSACTIONS.IFETCH, L2_TRANSACTIONS.PREFETCH, L2_TRANSACTIONS.L1D_WB, L2_TRANSACTIONS.FILL, L2_TRANSACTIONS.WB, L2_TRANSACTIONS.ANY, L2_LINES_IN.S_STATE, L2_LINES_IN.E_STATE, L2_LINES_IN.ANY, LZ_LINES_OUT.DEMAND_CLEA, LZ_LINES_OUT.DEMAND_DIRT, LZ_LINES_OUT.PREFETCH_CLE, LZ_LINES_OUT.PREFETCH_DIR, LZ_LINES_OUT.ANY, SQ_MISC.SPLIT_LOCK, SQ_FULL_STALL_CYCLES, FP_ASSIST.ALL, FP_ASSIST.OUTPUT, FP_ASSIST.INPUT, SIMD_INT_64.PACKED_MPY, SIMD_INT_64.PACKED_ARITH, SIMD_INT_64.PACK, SIMD_INT_64.PACKED_LOGICA, CPUID, SIMD_INT_64.PACKED_ARITH, SIMD_INT_64.SHUFFLE_MOVE, UNC_GQ_CYCLES_FULL.READ_, UNC_GQ_CYCLES_FULL.WRITE, UNC_GQ_CYCLES_FULL.PEER_, UNC_GQ_CYCLES_NOT_EMPTY, UNC_GQ_CYCLES_NOT UNC_GQ_ALLOC.READ_TRACK, UNC_GQ_ALLOC.RT_L3_MISS, UNC_GQ_ALLOC.RT_TO_L3_RE, UNC_GQ_ALLOC.RT_TO_RTID_, UNC_GQ_ALLOC.WT_TO_RTID_, UNC_GQ_ALLOC.WT_TO_R

UNC GQ DATA.FROM QPI, UNC GQ DAT UNC_SNP_RESP_TO_LOCAL_H, UNC_SNP UNC SNP RESP TO REMOTE, UNC SNP UNC_L3_HITS.PROBE, UNC_L3_HITS.ANY, UNC_L3_LINES_IN.F_STATE, UNC_L3_LINE UNC_L3_LINES_OUT.ANY, UNC_QHL_REQ UNC_QHL_CYCLES_FULL.IOH, UNC_QHL_ UNC_QHL_OCCUPANCY.REMOT, UNC_QH UNC QHL TO QMC BYPASS, UNC QMC UNC QMC ISOC FULL.READ.C, UNC QMC UNC QMC BUSY.READ.CH1, UNC QMC B UNC QMC OCCUPANCY.CH2. UNC QMC UNC_QMC_NORMAL_READS.C, UNC_QMC UNC_QMC_CRITICAL_PRIORIT, UNC_QMC UNC QMC WRITES.PARTIAL.C, UNC QMC UNC QMC PRIORITY UPDATE, UNC QMC

BR_MISP_EXEC.COND: Counts the number of mispredicted conditional near branch instructions executed, but not necessarily retired.

UNC GQ DATA.TO L3, UNC GQ DATA.TO CORES, TO_LOCAL_H, UNC_SNP_RESP_TO_REMOTE, REMOTE, UNC_L3_HITS.READ, UNC_L3_HITS.WRITE, STATE, UNC_L3_LINES_IN.S_STATE, NC_L3_LINES_OUT.F_STATE, LOCAL_, UNC_QHL_REQUESTS.LOCAL_ S_NOT_EMPT, UNC_QHL_OCCUPANCY.IOH IFLICT_CYCLES., UNC_QHL_CONFLICT_CYCLES., ULL.WRI, UNC QMC NORMAL FULL.WRI, C FULL.WRITE.C, UNC QMC BUSY.READ.CH0, UNC QMC OCCUPANCY.CH1, DRMAL READS.C. UNC QMC NORMAL READS.C. PRIORITY_RE, UNC_QMC_CRITICAL_PRIORIT, ULL.CH2, UNC_QMC_WRITES.FULL.ANY, QMC CANCEL.ANY, UNC QMC PRIORITY UPDATE, LLED SINGL, UNC QPI TX STALLED SINGL,

UNC_QPL_TX_STALLED_SINGL, UNC_QPL_TX_STALLED_SINGL, UNC_QPL_TX_STALLED_SINGL, UNC_QPL_TX_STALLED_SINGL, UNC_QPL_TX_STALLED_MULTI, UNC_QPL_TX_STALLED





Hardware Performance Counters

Derived Metrics

- Clock cycles per Instructions (CPI)
 - → CPI indicates if the application is utilizing the CPU or not
 - → Take care: Doing "something" does not always mean doing "something useful".
- Floating Point Operations per second (FLOPS)
 - → How many arithmetic operations are done per second?
 - → Floating Point operations are normally really computing and for some algorithms the number of floating point operations needed can be determined.





Amplifier XE – Hardware Counter

OPI rate (Clock cycles per instruction):

- In theory modern processors can finish 4 instructions in 1 cycle, so a CPI rate of 0.25 is possible
- A value between 0.25 and 1 is often considered as good for HPC applications
- Determine with "Advanced Hotspot" analysis

Elapsed Time 3: 3.327s () CPU Time 1: 76.810s Instructions Retired: 45,579,600,000 1 CPI Rate 2: 4.213 CPU Frequency Ratio :: 1.136 Total Thread Count: 24 Paused Time 10:

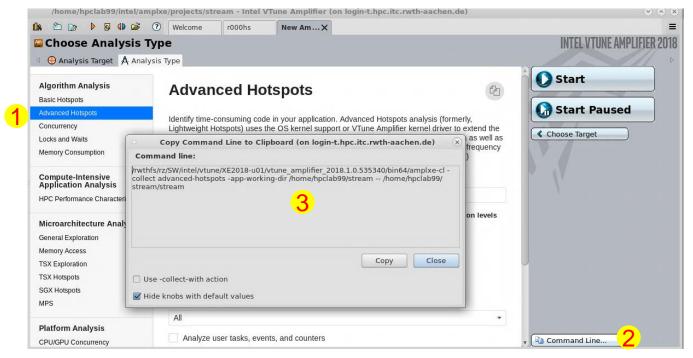


0s



VTune Amplifier in Command Line

- Real-world analysis should be execute in the batch system
 - ↑→ Choose analysis type
 - 2→ Push "Command Line..."
 - 3→ Copy & paste command line and use in batch script







Summary

Correctness:

- Data Races are very hard to find, since they do not show up every program run.
- Intel Inspector XE or ThreadSanitizer help a lot in finding these errors.
- Use really small datasets, since the runtime increases significantly.
- If possible use non optimized code.

Performance:

- Start with simple performance measurements like hotspots analyses and then focus on these hot spots.
- In OpenMP applications analyze the waiting time of threads. Is the waiting time balanced?
- Hardware counters might help for a better understanding of an application, but they might be hard to interpret.



